

Recursion

- In C/C++, a function may call itself either directly or indirectly.
- When a function calls itself recursively, each invocation gets a fresh set of all explicit parameters and automatic (local) variables, independent of the previous set.

Example: Compute the greatest common divisor (GCD) of two numbers.

Let $M > N > 0$, and $M = QN + R$ such that $0 \leq R < N$

- if $R = 0$, then $M = QN$, and $\text{gcd}(M, N) = N$
- if $R > 0$, then $N > R$, and $\text{gcd}(M, N) = \text{gcd}(N, R)$

Non-recursive algorithm to compute the GCD of 2 numbers.

```
unsigned gcd(unsigned M, unsigned N)
/* given M >= 0 and N > 0, compute gcd of M & N */
{
    unsigned R;
    while (R = M % N)    // R = M % N and R != 0
    {
        M = N;
        N = R;
    }
    return N;
}
```

Recursive algorithm to compute the GCD of 2 numbers

```
unsigned gcd(unsigned M, unsigned N)
/* given M >= 0 and N > 0, compute gcd of M & N */
{
    unsigned R = M % N;
    if (R == 0)           // base case
        return N;
    else
        return gcd(N, R); // recursion
}
```

Two fundamental rules of recursion:

1. Base cases

You must have some base cases, which can be solved without recursion.

2. Making progress (recursion)

For the cases that are to be solved recursively, the recursive call must always be to a case that makes progress toward a base case.

Example: Fibonacci number

$$\begin{aligned}f(0) &= 0; \\f(1) &= 1; \\f(n) &= f(n-2) + f(n-1) \text{ for } n \geq 2;\end{aligned}$$

```
unsigned fib(unsigned n)
{  // precondition: given n >= 0

    if (n == 0)      // base case #1
        return 0;

    if (n == 1)      // base case #2
        return 1;

    return fib(n-2) + fib(n-1);  // recursion

} // this program works but is unacceptably slow
```

Usually, looping (iteration) is more efficient than recursion.

Computing the Fibonacci number using iteration is much more efficient.

Recursion is preferred when the problem is recursively defined, or when the data structure that the algorithm operates on is recursively defined.

Non-Recursive binary search

```
// Precondition: A[] is sorted in ascending order
int binSearch(int A[], int N, int key)
{
    int low = 0;
    int high = N-1;

    while (low <= high)
    {
        int mid = (low + high) / 2;

        if (A[mid] < key)
            low = mid+1;
        else if (A[mid] > key)
            high = mid-1;
        else
            return mid; // key == A[mid]
    }

    return -1;
}
```

Recursive binary search

```
// Precondition: A[] is sorted in ascending order
// A[low..high] represents the search range

int binSearch_R(int A[], int low, int high, int key)
{
    if (low > high) // search range is empty
        return -1; // key is not found

    int mid = (low + high) / 2; // probe the mid-point of
    if (A[mid] == key) // the search range
        return mid;

    if (A[mid] < key) // recursion with reduced search range
        return binSearch_R(A, mid+1, high, key);
    else
        return binSearch_R(A, low, mid-1, key);
}
```

Recursive function to print the permutations of an array of N symbols

//Consider the case for N=3

```
char x[3] = { 'a', 'b', 'c' };
```

Permutations of the 3 symbols are:

```
a b c  
a c b  
b a c  
b c a  
c b a  
c a b
```

Pseudo code:

```
void permute(char x[], int s, int e)
// Enumerate the permutations of the symbols in x[s..e],
// x[s] to x[e] inclusive.

// Symbols in x[0..s-1] have already been fixed.
{
    if (s == e) // base case
        // only 1 permutation is possible
        print the contents of x[];

    else // recursion
    {
        // x[s] can have e-s+1 choices
        for each choice of x[s]
            find the permutations of x[s+1..e] by recursion;
    }
}
```

```

//Detailed codes

void swap(char& a, char& b) // swap a with b
{
    char t = a;
    a = b;
    b = t;
}

void permute(char x[], int s, int e)
// Enumerate the permutations of the symbols in x[s..e]
{
    if (s == e) // base case
        cout << x << endl;
    else //recursion
    {
        for (int k = s; k <= e; k++)
        {
            //1. there are e-s+1 choices for the first symbol
            swap(x[s], x[k]);

            //2. once we fix the first symbol,
            //   find the permutations of the remaining e-s
            //   symbols by recursion
            permute(x, s+1, e);

            //3. restore the original order of the array
            swap(x[s], x[k]);
        }
    }
}

void testPermute()
{
    char x[] = "abcd";

    permute(x, 0, 3); //print permutations of 4 symbols
}

```

Example: Towers of Hanoi

- There are 3 towers, labeled L, M and R
- Initially, tower L contains a stack of N disks of graded sizes, stacked in order of size with the smallest disk at the top.
- A move consists of taking the top disk from one tower and placing it on another tower, subject to the constraint that a disk may never be placed above a disk smaller than itself.
- The problem is to move the N disks from tower L to tower R in the smallest number of moves.

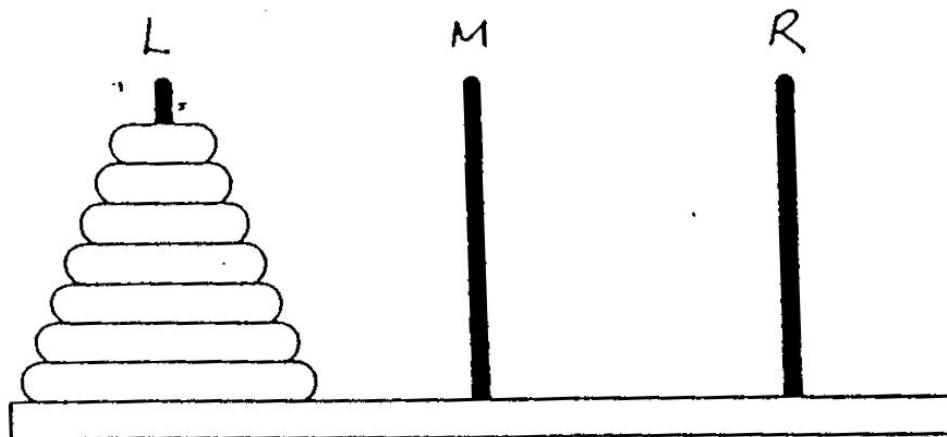


Figure 3.8. The Towers of Hanoi

Sequence of moves:

N = 1	N = 2	N = 3
 L → R	  	      

The problem can be solved in 3 “steps” using recursion:

1. Move the top N-1 disks from L to M using R as a temporary working space.
2. Move the bottom disk from L to R.
3. Move the remaining N-1 disks from M to R using L as a temporary working space.

Let `move(N, F, T, U)` denotes the operation of moving `N` disks from tower `F` to tower `T` using tower `U`.

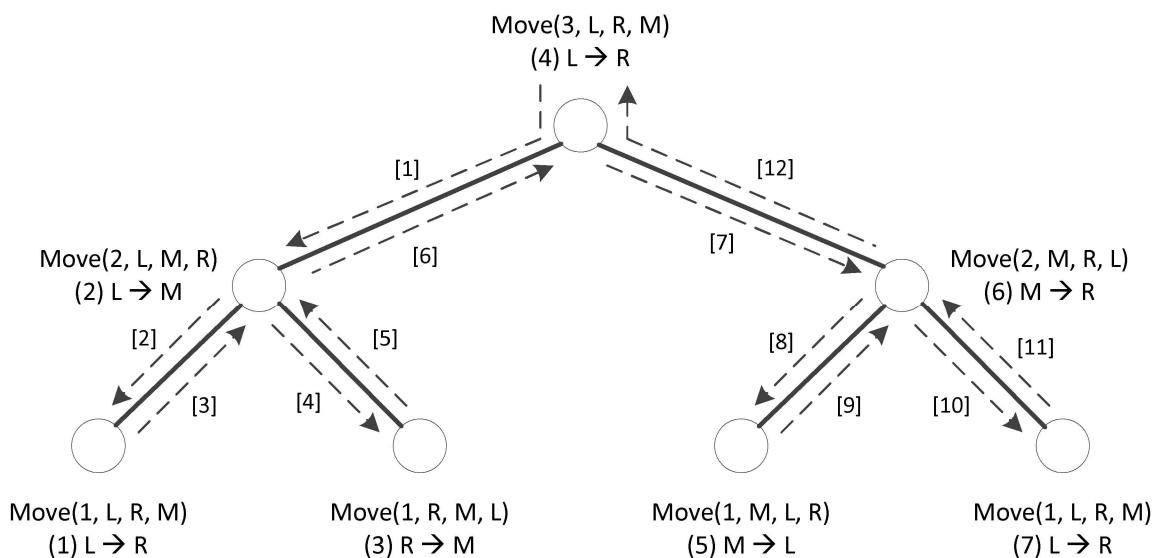
Recursive function to print out the sequence of moves:

```
void move(int N, char F, char T, char U)
{
  if (N == 1)
    cout << F << " -> " << T << endl;
  else if (N > 1)
  {
    move(N-1, F, U, T);
    cout << F << " -> " << T << endl;
    move(N-1, U, T, F);
  }
}
```

```
//version 2, without the else statement
```

```
void move_2(int N, char F, char T, char U)
{
    if (N > 0)
    {
        move_2(N-1, F, U, T);
        cout << F << " -> " << T << endl;
        move_2(N-1, U, T, F);
    }
}
```

Recursion tree for moving 3 disks



depth of the recursion = number of disks to be moved = N

$$\begin{aligned} \text{number of moves} &= \text{number of nodes in the recursion tree} \\ &= 1 + 2 + 4 + \dots + 2^{N-1} = 2^N - 1 \end{aligned}$$