## Homework 1

# Description

The goal of this assignment is to master the following skills:

- be able to write basic algorithms using recursion
- understand and be able to use asymptotic notation  $(O, \Theta, \Omega)$ 
  - o given two functions, recognize all their possible relationships in terms of O,  $\Theta$  and  $\Omega$
  - simplify functional expression with O-notation

### Submission

### Problems to be submitted

#### **Problem 1**

Complete the following implementation in Java.

```
/* Given a string and a character, find the number of times the character
 * appears in the string. Matches are case-sensitive.
 * Give a recursive implementation.
 */
public static int countChar(String str, char c) {
    return 0;
}
```

#### Problem 2

Complete the following implementation in Java.

```
/* Find the maximum value in a list of Integers, using recursion.
    *
    * Hint: to keep track of which parts of the list still need to be visited,
    * which works a lot like cleanHotel(int lo, int hi)
```

```
* use a recursive helper function: recursiveMaxHelper(List<Integer> li, int lo, int

*

*/
public static int recursiveMax(List<Integer> li) {
    return 0;
}
```

#### **Problem 3**

Complete the following implementation in Java.

#### **Problem 4**

Show that  $3^{\log_4 n} = n^{\log_4 3}$  using log rules. Give an identity that generalizes this equality. (This problem is an exercise is using log/exponent rules. The identity will also be useful for simplifying expressions.)

#### **Problem 5**

a) If f(n) is O(g(n)), is it necessarily true that  $2^{f(n)}$  is  $O(2^{g(n)})$ ? Justify your answer.

b) Explain why the statement "The running time of algorithm A is at least  $O(n^2)$ " is meaningless.

#### Problem 6

Order by  $\Theta$ -notation, and indicate those have the same asymptotic growth rate (i.e., that are  $\Theta$  of each other):  $e^n$ , n,  $2^n$ ,  $n \log n$ ,  $\log n$ ,  $n \log n$ ,  $n \log$ 

**Note:** intuitively, "f is  $\Theta(g)$ " roughly means "f=g", the same way that "f is O(g)" roughly means " $f \leq g$ ".

Formally, "f is  $\Theta(g)$ " means "f is O(g)" and "g is O(f)".

#### **Problem 7**

Give a big-O estimate for each of the following functions. For the function g in your estimate that f(n) is O(g(n)) use as simple of a function g as possible, and of the smallest order.

a) 
$$5n^3 - 7n^2 + 88$$

**b)** 
$$(n \log n + n^2)(n^3 + 2)$$

c) 
$$\log(n^3 + 1) + (\log n)^2$$

d) 
$$(n \log n + 1)^2 + (\log n + 1)(n^2 + 1)$$

**e)** 
$$n^{2n} + n^{n^2}$$

f) 
$$(n! + 2^n)(n^3 + \log(n^2 + 1))$$

g) 
$$n \cdot (5/4)^{\log_4 n + 1}$$

**Hint**: you can simplify **g**) to  $O(n^?)$  where ? is something for you to determine.

### **Problem 8**

Give a big-O estimate for each of the following functions. For the function g in your estimate that f(n) is O(g(n)) use as simple of a function g as possible, and of the smallest order.

a) 
$$\sum_{i=10}^{n} i$$

b) 
$$\sum_{i=1}^{\log_2 n} ni$$

c) 
$$\sum_{i=1}^n 4^n$$

d) 
$$\sum_{i=1}^n 4^i$$
 e)  $\sum_{i=1}^n 4^i$ 

f) 
$$\sum_{i=1}^{\log_2 n} n(1/3)^i$$

# Additional problems (not to be submitted)

### need more practice with recursion?

- finish the practice problems in
- o at codingbat.com, solve Recursion-1 problems

### challenging asymptotic growth rates

The following functions commonly appear in the analysis of more advanced algorithms:

- $\circ \log \log n$ ,
- $\circ$   $n/(\log\log n)$ ,
- $\circ \log^* n$  (you can also find an explanation in the textbook)

Can you order these with respect to the functions in Problem 6?