

Assignment - 2

Total Marks: 120

Deadline: 14 May 2025, 11:55 pm

You have to upload the scanned PDF of your handwritten assignment.

1. Proof that the following languages are non-regular using the pumping lemma. [30 marks]

- a. $L = \{w \in \{0, 1\}^*, w = 0^i 1^j 0^k \text{ where } i + k = j \text{ and } i, j, k > 0\}$
- b. $L = \{w \in \{0, 1\}^*, w = 0^i 1^j 0^k \text{ where } i + k > j \text{ and } i, j, k > 0\}$
- c. $L = \{w \in \{0, 1\}^*, w = 0^i 1^j 0^k \text{ where } i + k < j \text{ and } i, j, k > 0\}$
- d. $L = \{w \in \{0, 1\}^*, ww^R\}$ (w^R denotes the reverse string of w)
- e. $L = \{w \in \{0, 1\}^*, ww'\}$ (w' denotes a string where every bit is flipped from w)
- f. $L = \{w \in \{0, 1\}^*, w1^n \text{ where } n > 0 \text{ and } |w| = n\}$

2. Design the Context Free Grammar (CFG) for the following languages. [30 marks]

- a. $L = \{w \in \{0, 1\}^*, \text{ the length of } w \text{ is odd and the mid is } 0\}$
- b. $L = \{w \in \{0, 1\}^*, w \text{ contains even number of } 0\text{s}\}$
- c. $L = \{w \in \{0, 1\}^*, w \text{ starts and ends with the same symbol}\}$
- d. $L = \{w_1, w_2 \in \{0, 1\}^*, w_1 \# w_2 \text{ where } w_2 \text{ starts with "01" and } |w_1| = |w_2|\}$
- e. $L = \{w \in \{0, 1\}^*, w \text{ contains exactly three } 1\text{s}\}$
- f. $L = \{w \in \{0, 1\}^*, w = 0^i 1^j 0^k \text{ where } i + k < j \text{ and } i, j, k > 0\}$

3. [10 marks]

Take a look at the grammar below and solve the following problems.

$$\begin{aligned} S &\rightarrow AX \\ A &\rightarrow 0A \mid 1B \mid \epsilon \\ B &\rightarrow 0B \mid 1A \\ X &\rightarrow 00X \mid 01X \mid 10X \mid 11X \mid 0 \mid 1 \end{aligned}$$

- (a) Give a leftmost derivation for the string 01101010. (3 points)
- (b) Sketch the parse tree corresponding to the derivation you gave in (a). (2 points)
- (c) Demonstrate that the given grammar is ambiguous by showing two more parse trees (apart from the one you already found in (b)) for the same string. (4 points)
- (d) Find a string w of length six such that w has exactly one parse tree in the grammar above. (1 point)

4. Construct Push-Down Automata for the following languages. [30 marks]

- a. $L = \{w \in \{0, 1\}^*, w = 0^i 1^j 0^k \text{ where } i + k = j \text{ and } i, j, k > 0\}$
- b. $L = \{w \in \{0, 1\}^*, w = 0^i 1^j 0^k \text{ where } i + k > j \text{ and } i, j, k > 0\}$
- c. $L = \{w \in \{0, 1\}^*, w = 0^i 1^j 0^k \text{ where } i + k < j \text{ and } i, j, k > 0\}$
- d. $L = \{w \in \{0, 1\}^*, ww^R\}$ (w^R denotes the reverse string of w)
- e. $L = \{w\#x, \text{ where } x \in \{0, 1\}^* \text{ and } w^R \text{ is a substring of } x\}$ (w^R denotes the reverse string of w)
- f. $L = \{w \in \{0, 1\}^*, w1^n \text{ where } n > 0 \text{ and } |w| = n\}$

5. Design the Turing Machines for the following languages. [20 marks]

- $L = \{w \in \{a, b\}^*, w \text{ is a palindrome}\}$
- Design a Turing Machine (TM) that, given an input string w over the alphabet $\{a, b\}$, writes the same string again after a delimiter symbol $\#$. The TM should halt with the tape containing $w\#w$, where the first occurrence of w is the input and the second is its copy.