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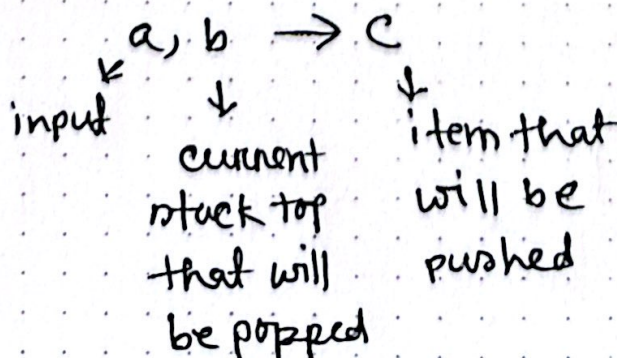
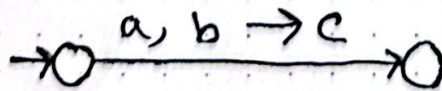
Pushdown Automata

- Computational model that recognizes CFL (context-free language)

- PDA - NFA paired with stack

It is similar to NFA

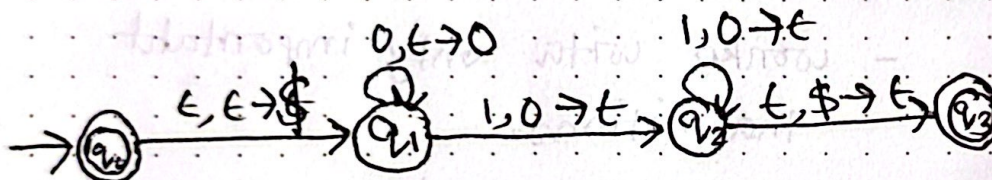
- works with only important transitions
- can have epsilon transition
- CFG has memory of at most two quantity.
- γ - stack input [different from alphabet]



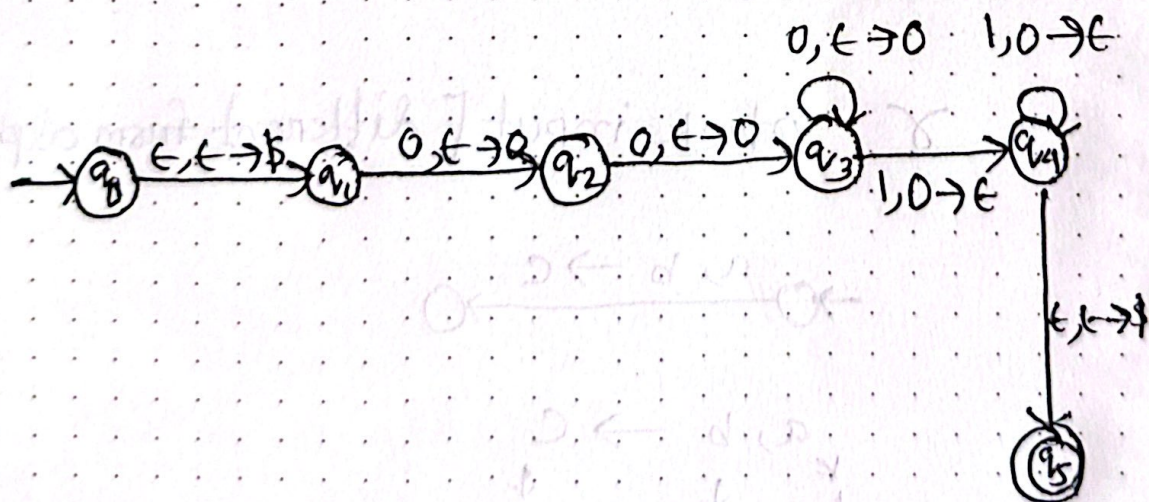
- moving from one state to another means $\epsilon, \epsilon \rightarrow \epsilon$

PDA Examples

$$L = \{0^n 1^n; n \geq 0\}$$

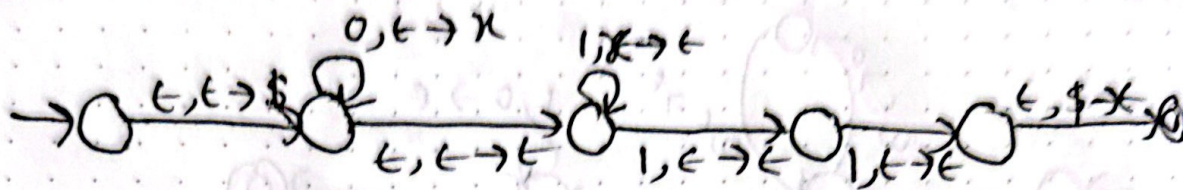


$$L = \{0^n 1^n; n \geq 2\}$$

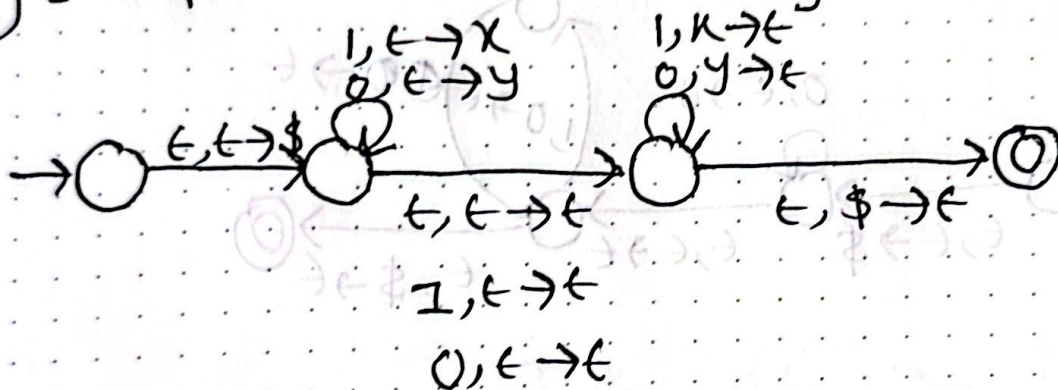


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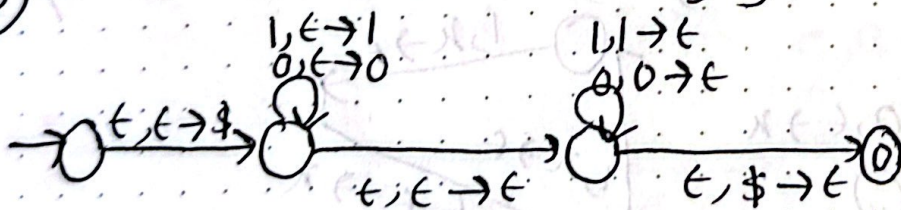
③ $L = \{ w = 0^n 1^{n+2} ; n \geq 0 \}$



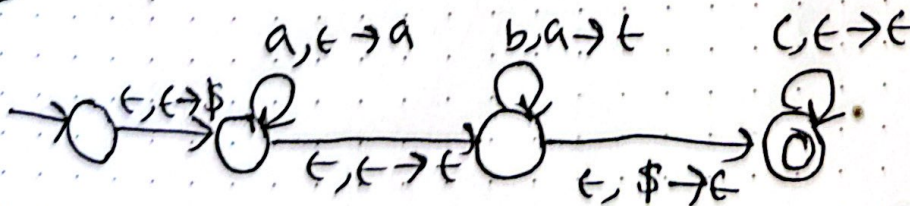
④ $L = \{ w \text{ is a palindrome} \}$



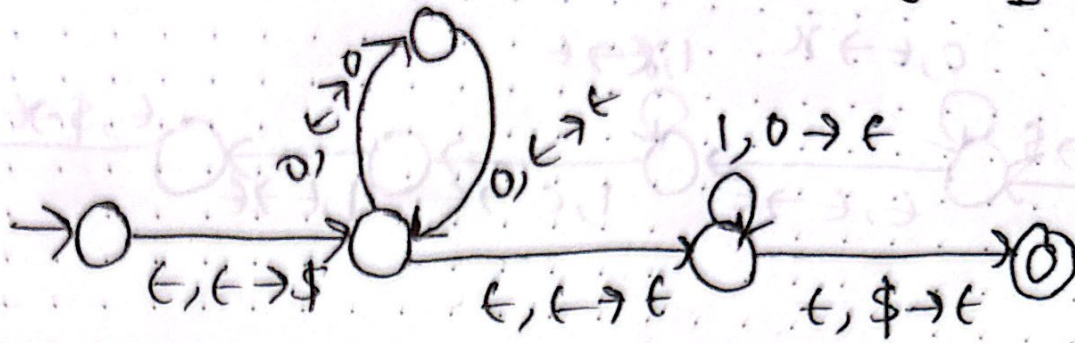
⑤ $L = \{ w_1 w_2 ; w_1 \in \{0,1\}^* \}$



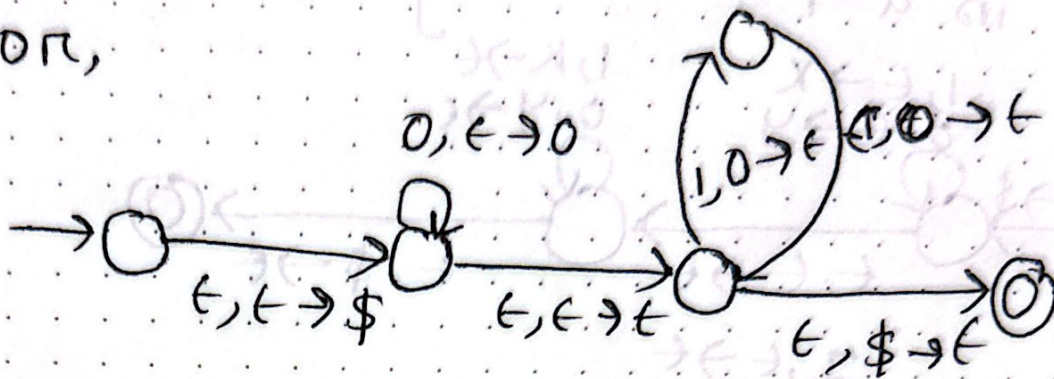
⑥ $L = \{ w = a^i b^j c^k \text{ where } i = j \}$



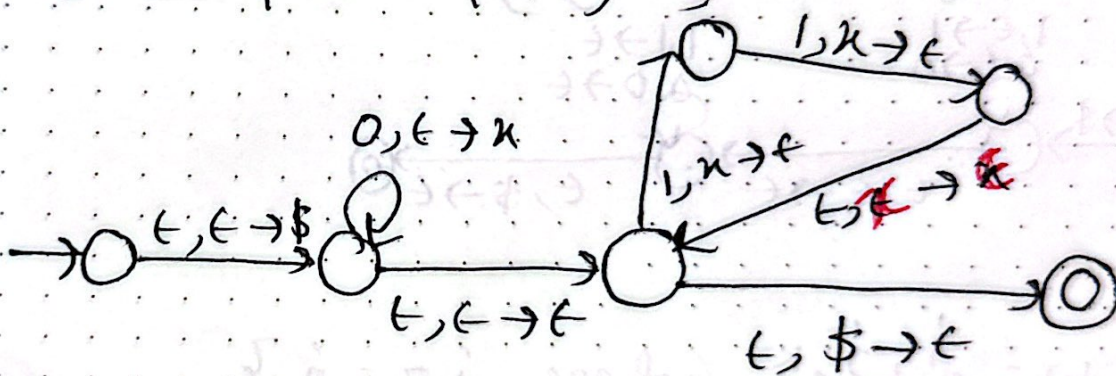
⑦ $L = \{ w \in \{0,1\}^* : w = 0^{2n} 1^n ; n > 0 \}$



or,



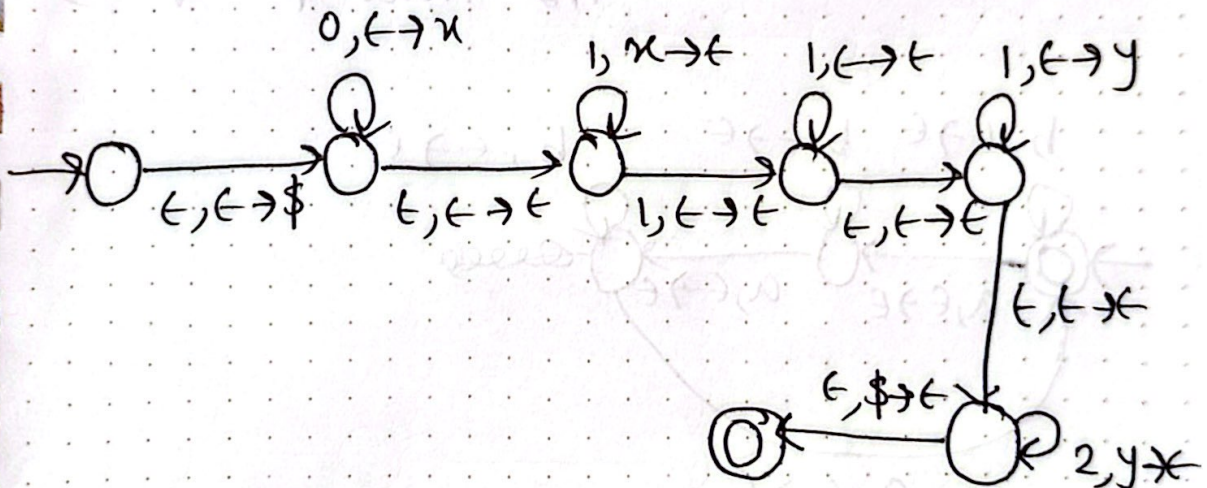
⑧ $L = \{ w \in \{0,1\}^* : w = 0^{3n} 1^{2n}, n > 0 \}$



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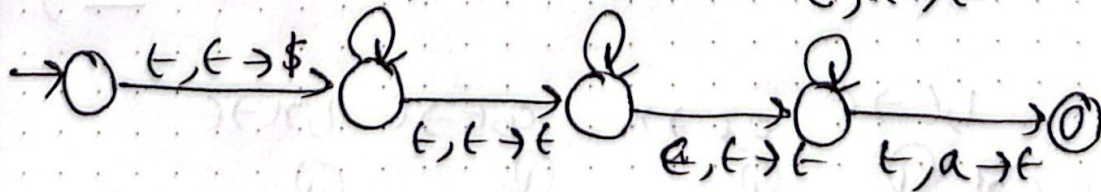
⑨ $L = \{ w = 0^i 1^j 2^k \mid j > i + k ; i, k > 0 \}$

$0^i 1^i \underline{1}^p 1^k 2^k$

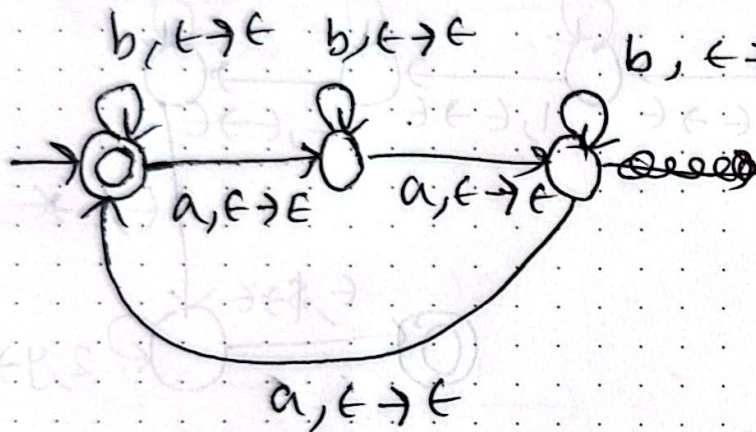


⑩ $L = \{ w = a^i b^j c^k \mid i + j > k \}$

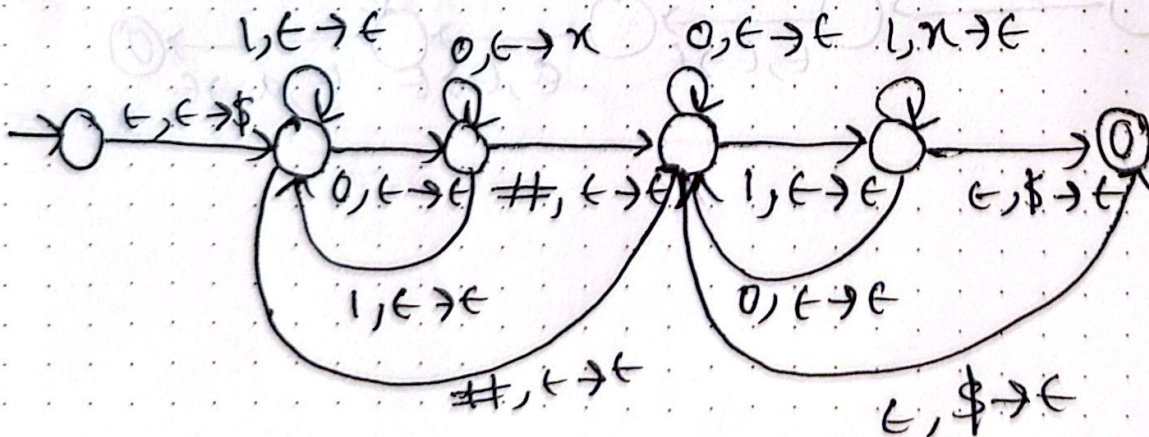
$a, \epsilon \rightarrow a \quad b, \epsilon \rightarrow a \quad c, a \rightarrow \epsilon$



- (41) $L = \{ \text{the count of 'a' in } w \text{ is multiple of 3} \}$



- (12) $\{ w_1 \# w_2 ; \text{the number of '00' in } w_1 \text{ is the same as the number of '11' in } w_2 \}$



$L = \{w \in \{0,1\}^* : \text{no of 0's \& 1's are equal in } w\}$

Hint - push all 0's & 1's in stack. [Pop 1's for 0's & 0's for 1's]

