



Inspiring Excellence

Network Layer

Routing Algorithm
Distance Vector Routing

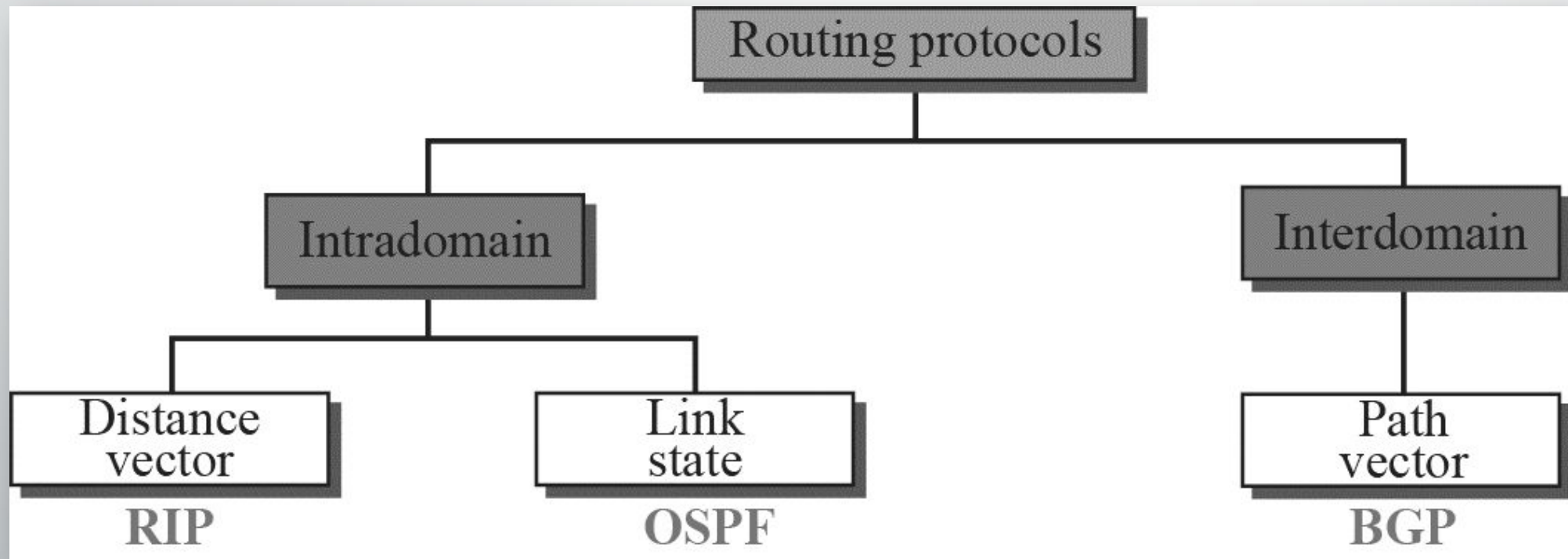
Lecture 11 | CSE421 – Computer Networks

Department of Computer Science and Engineering
School of Data & Science

Objectives

- understand principles behind network layer services:
 - routing algorithms
 - distance vector
 - link state

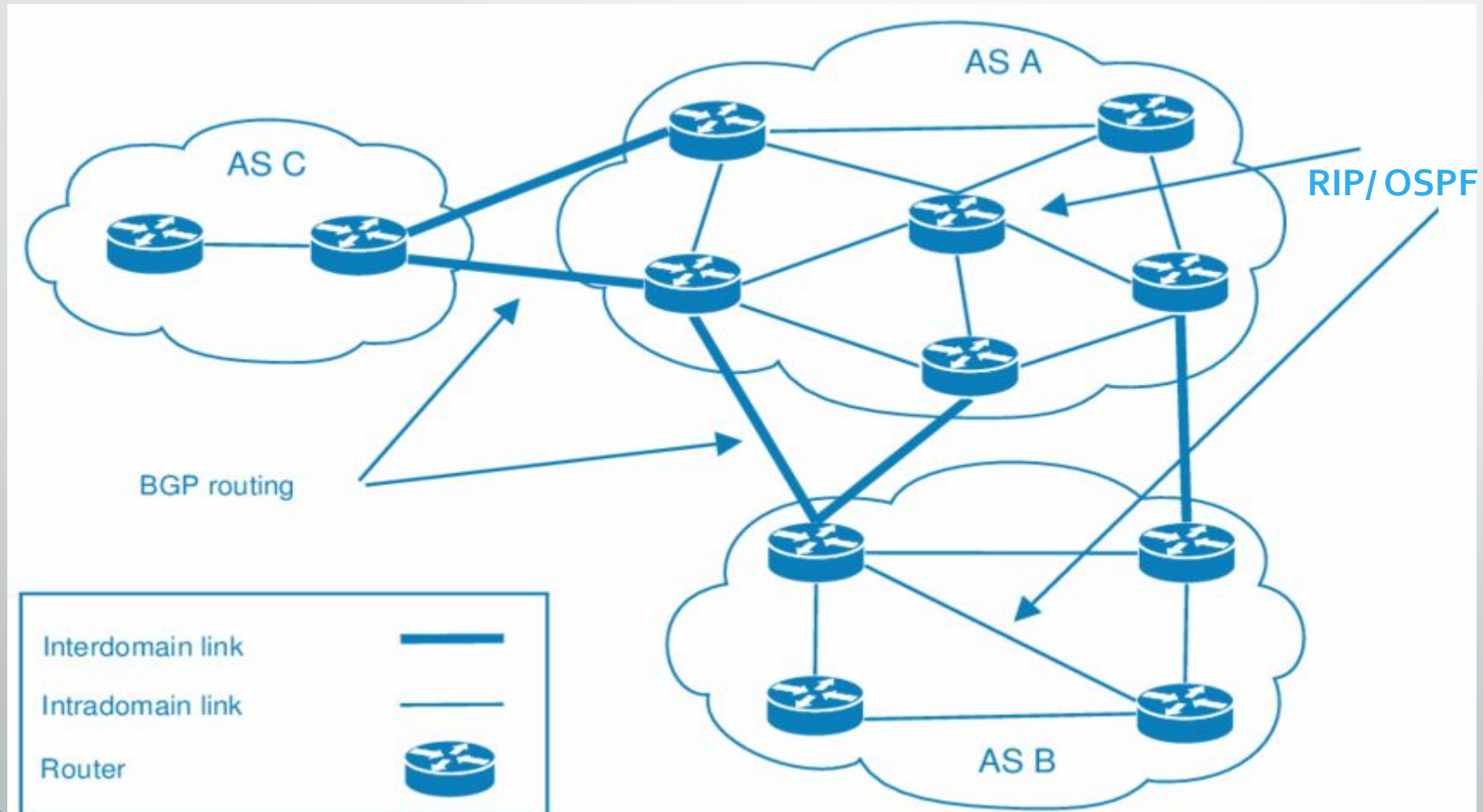
Popular Routing Protocols



Autonomous Systems

- Internet is divided into autonomous systems.
- An autonomous system (AS) is a group of networks and routers under the authority of a single administration.
- Routing *inside* an autonomous system is called **intra-domain routing**. Routing *between* autonomous systems is called **inter-domain routing**.

Autonomous Systems



Routing Algorithms

- Given a set of routers and links connecting the routers.
- Routing algorithm finds a “good” path from the source to destination router.
- Good path = Least cost path

Routing Algorithm classification

Global and Decentralized

Global:

- all routers have complete topology and link cost info
- “link state” algorithms

Decentralized:

- router knows physically-connected neighbors, link costs to neighbors
- iterative process of computation, exchange of info with neighbors
- “distance vector” algorithms



Distance Vector Algorithm

Distance Vector Algorithm

Based on Bellman-Ford Algorithm

- ☐ computes **shortest paths** from a single source node to all of the other nodes in a **weighted topology**.
- ☐ Shortest paths are based on the cost calculated from source node to destination node.
- ☐ Cost for now is arbitrary. It can be calculated using hop counts, link delay, link bandwidth etc.
- ☐ Distributed route computation using only neighbor's info

Distance vector algorithm

Basic idea:

- Each node periodically sends its own distance vector estimate to neighbors
- When a node x receives new DV estimate from neighbor;
- It updates its own DV using B-F equation

Operation of Distance Vector

- Periodic Updates:

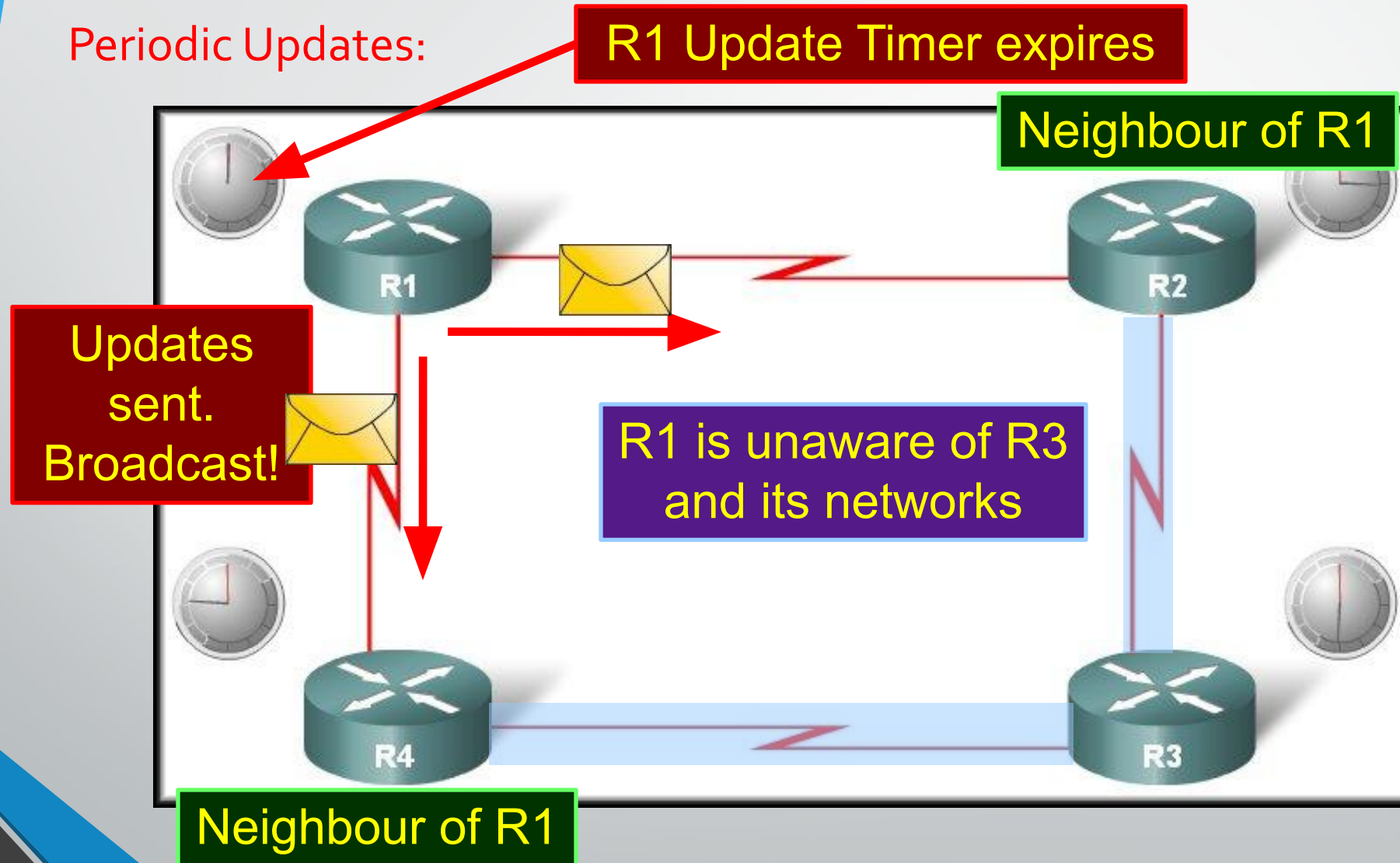
- Periodically broadcast the entire routing table to each of its neighbors (RIP – every 30 seconds).

- Inefficient

- Router is only aware of the:

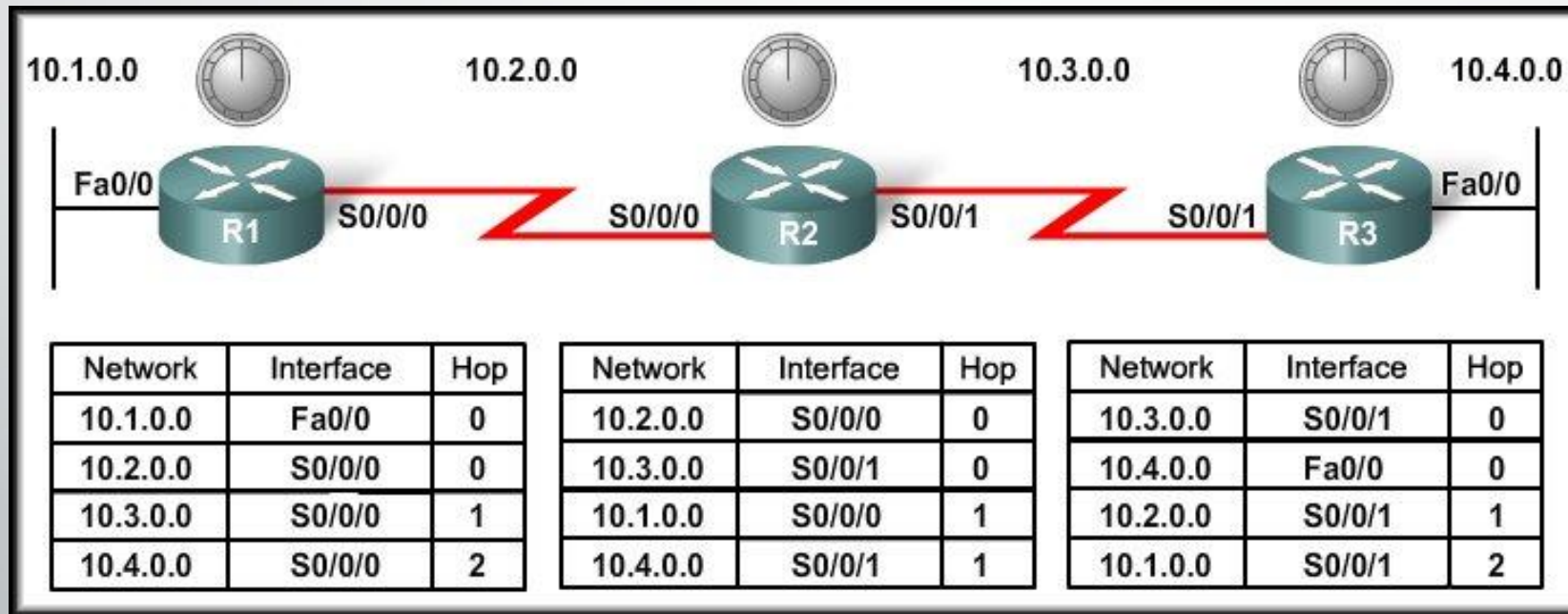
- Network addresses of its own interfaces.
 - Network addresses the neighbors running the same routing protocol.

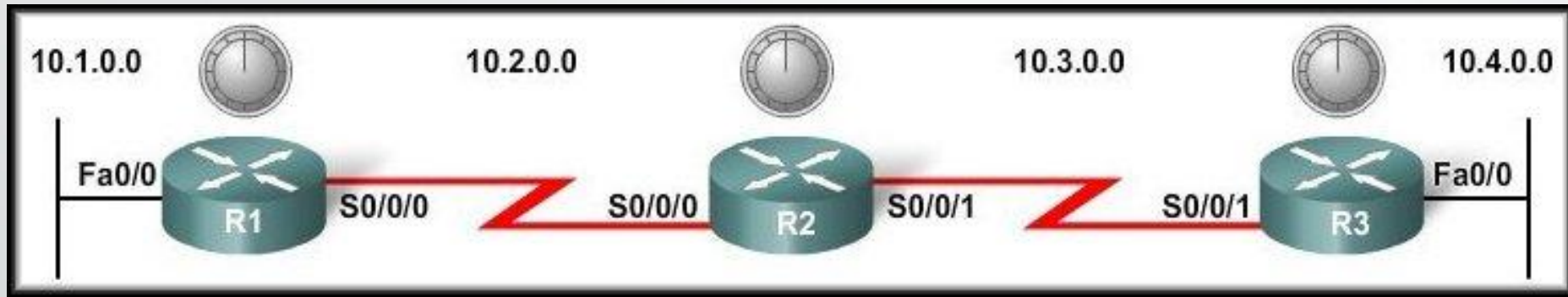
Operation of Distance Vector



Distance Vector Routing Protocols

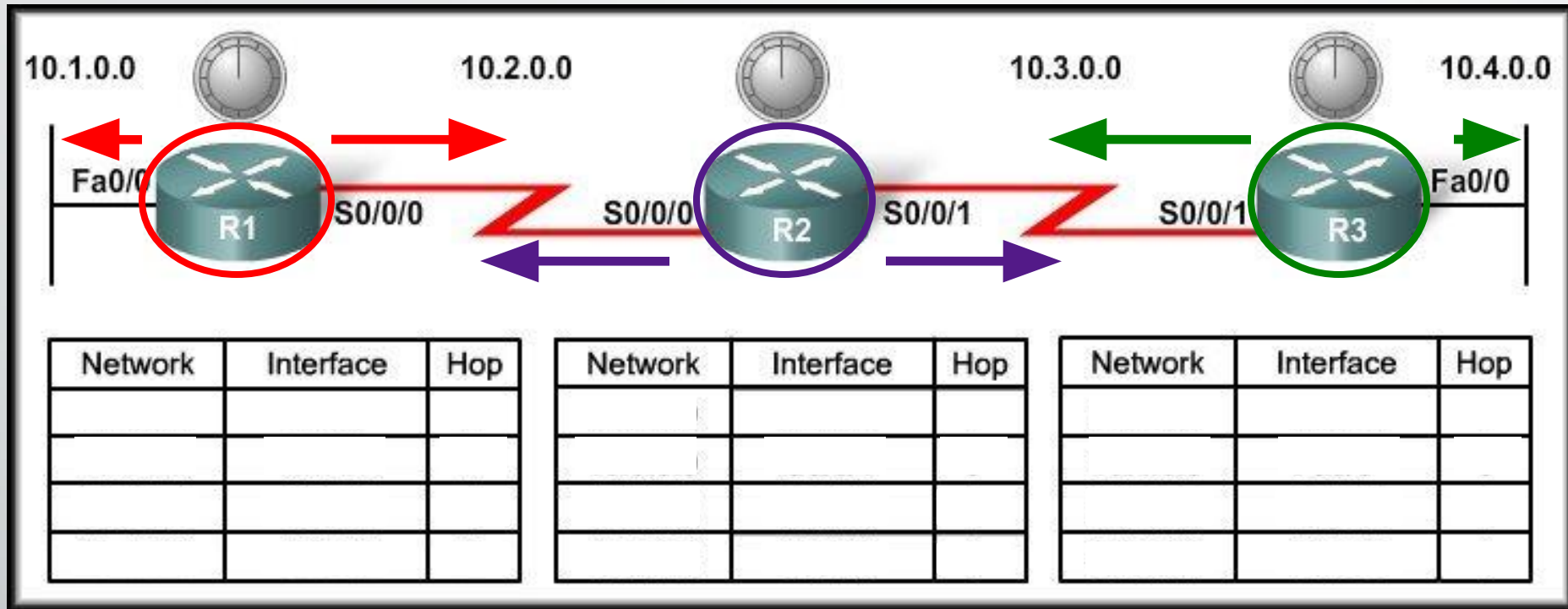
Network Discovery





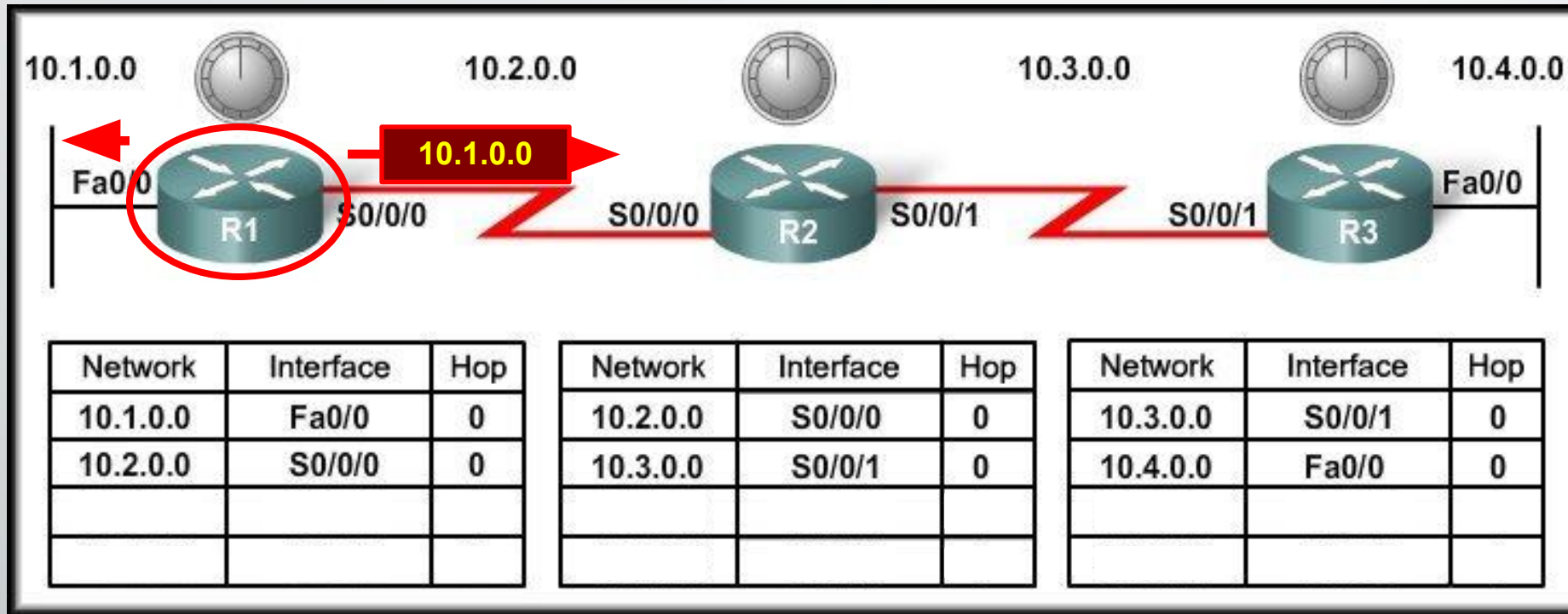
- Network Discovery:
 - Is part of the process of the routing protocol algorithm that enables routers to **learn about remote networks for the first time.**

Cold Start



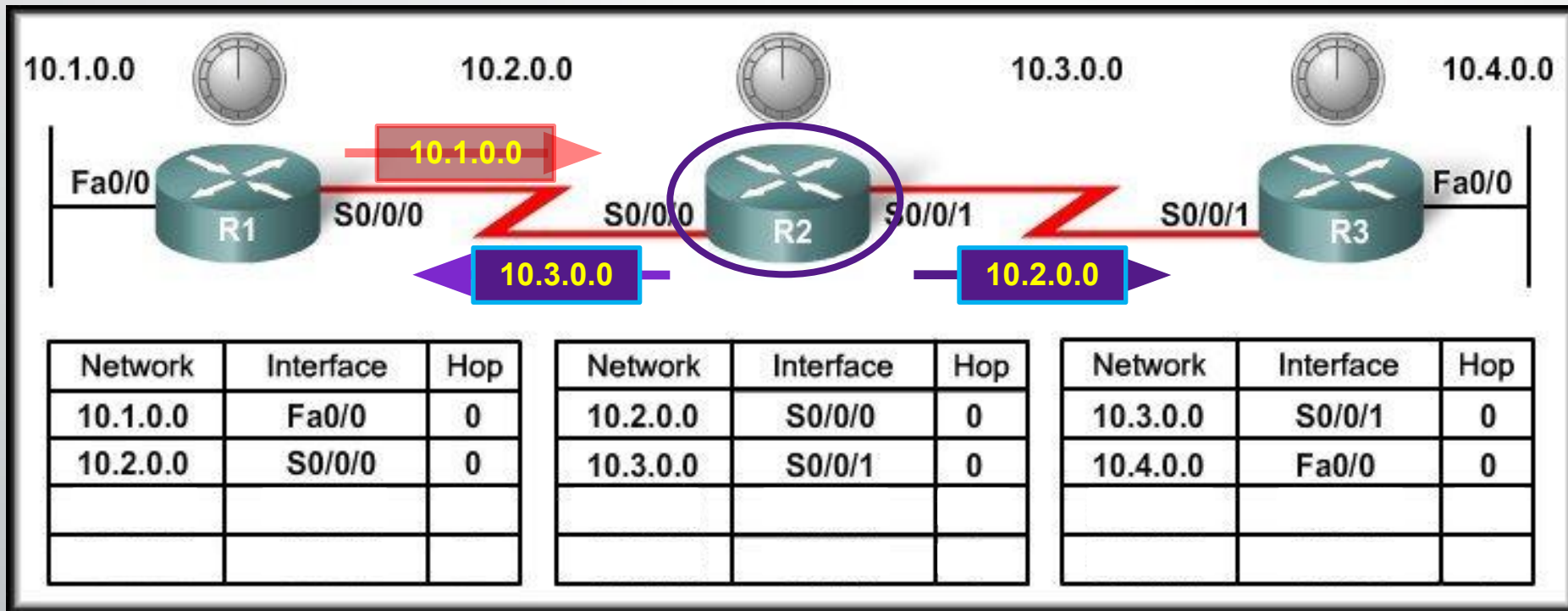
- When a router powers up:
 - Knows nothing about the network topology.
 - Knows only the information saved in NVRAM.
 - Sends updates about its known networks out all ports.

Initial Exchange of Routing Information

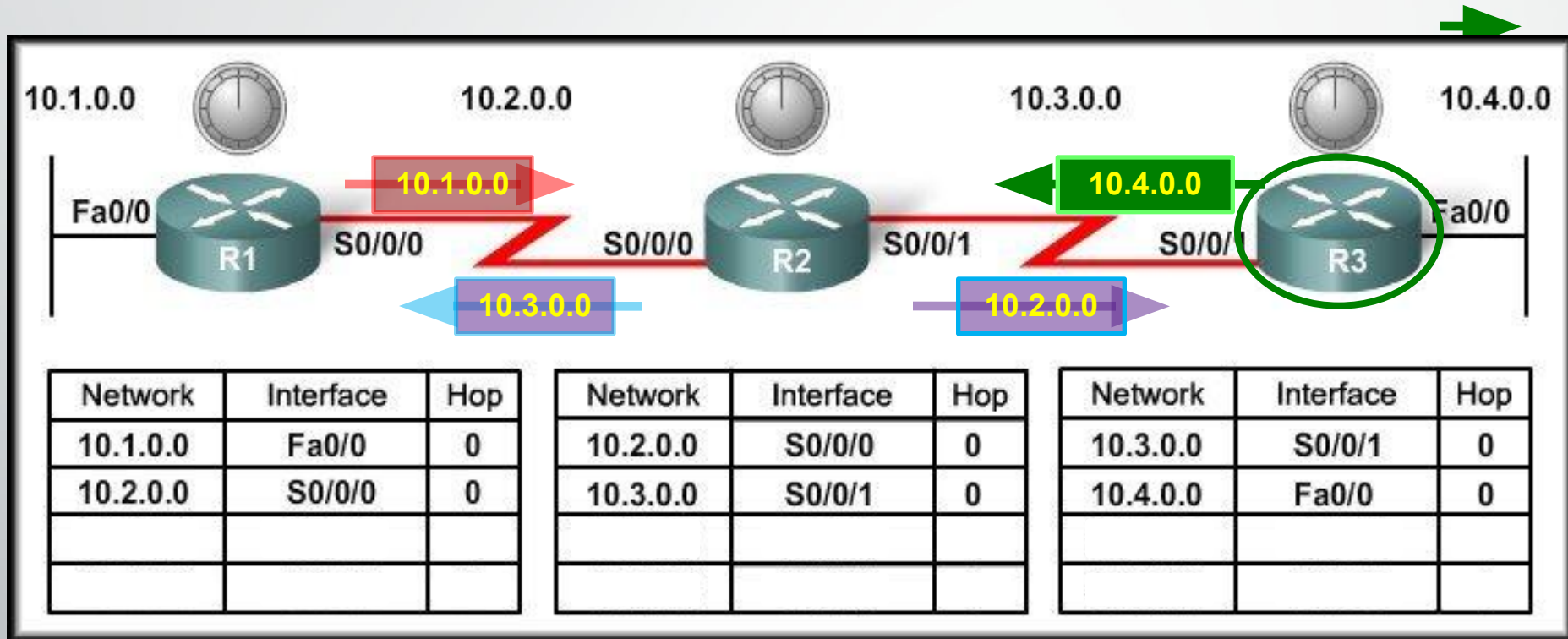


- Sends an update about network **10.1.0.0** out the **Serial o/o/o** interface with a metric of 1.
- Sends an update about network **10.2.0.0** out the **Fao/o** interface with a metric of 1.

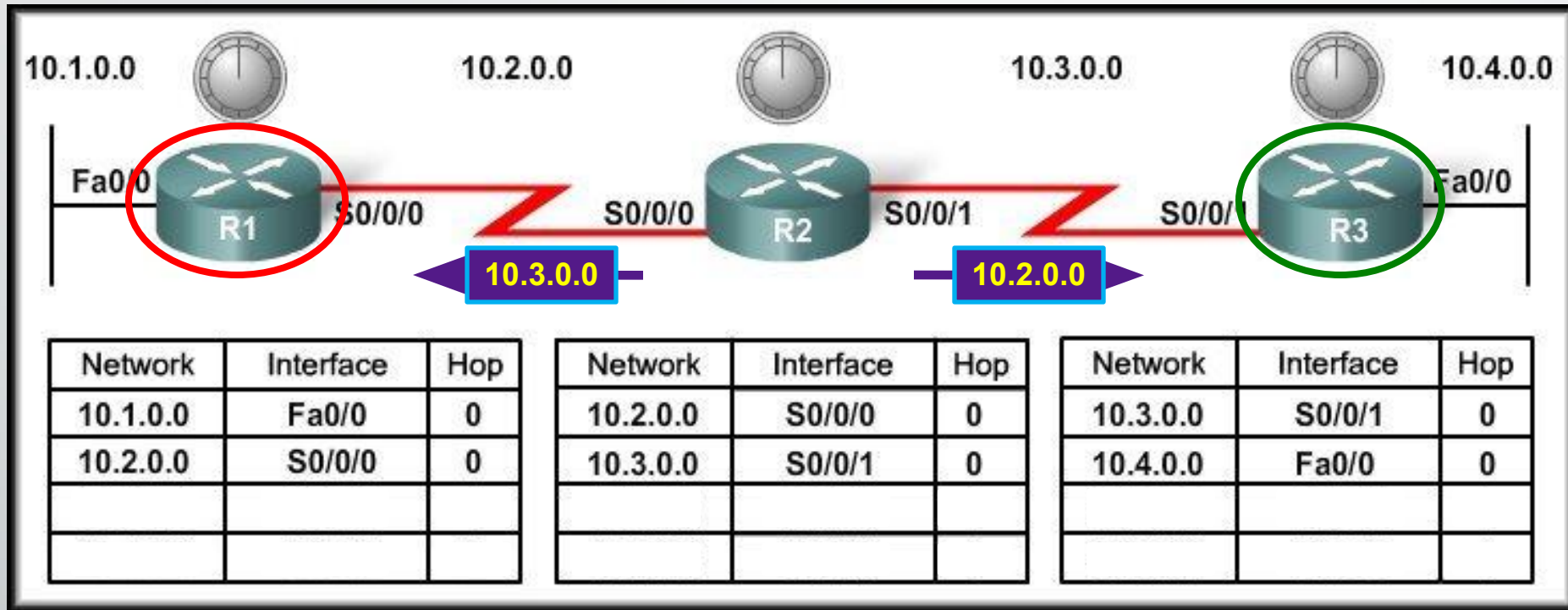
Initial Exchange of Routing Information



Initial Exchange of Routing Information

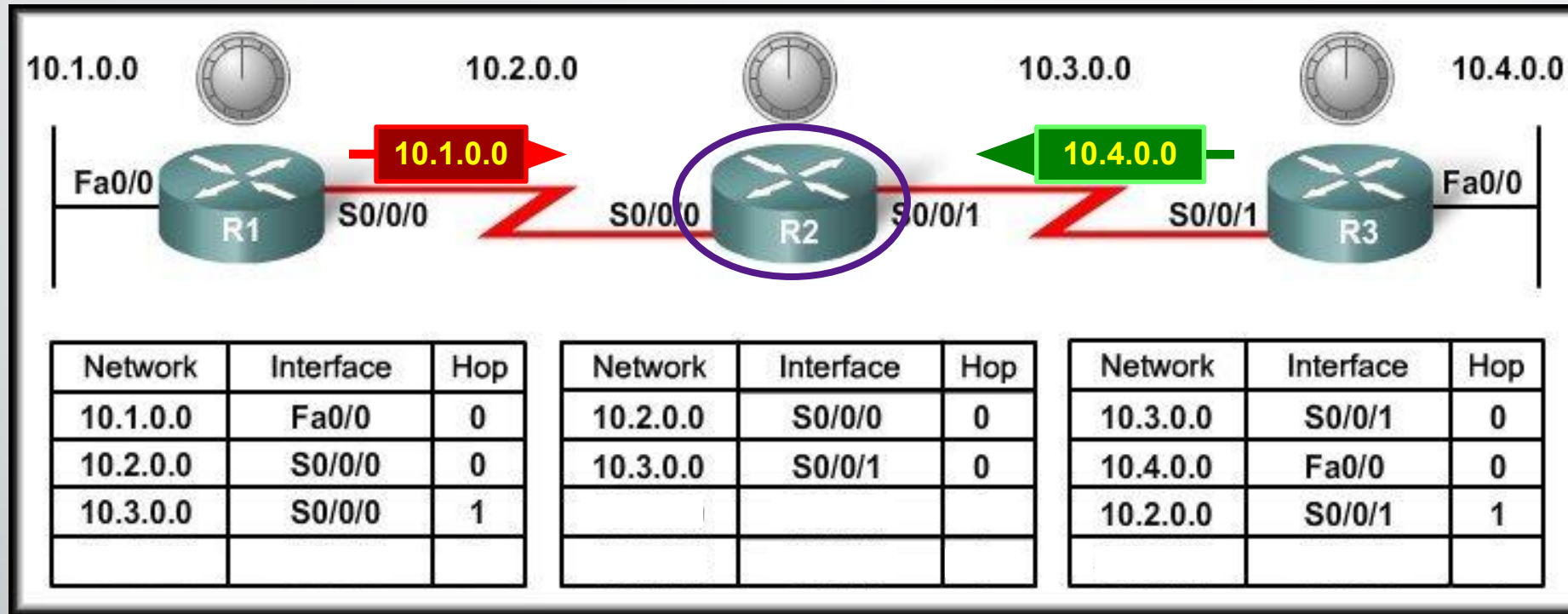


Initial Exchange of Routing Information



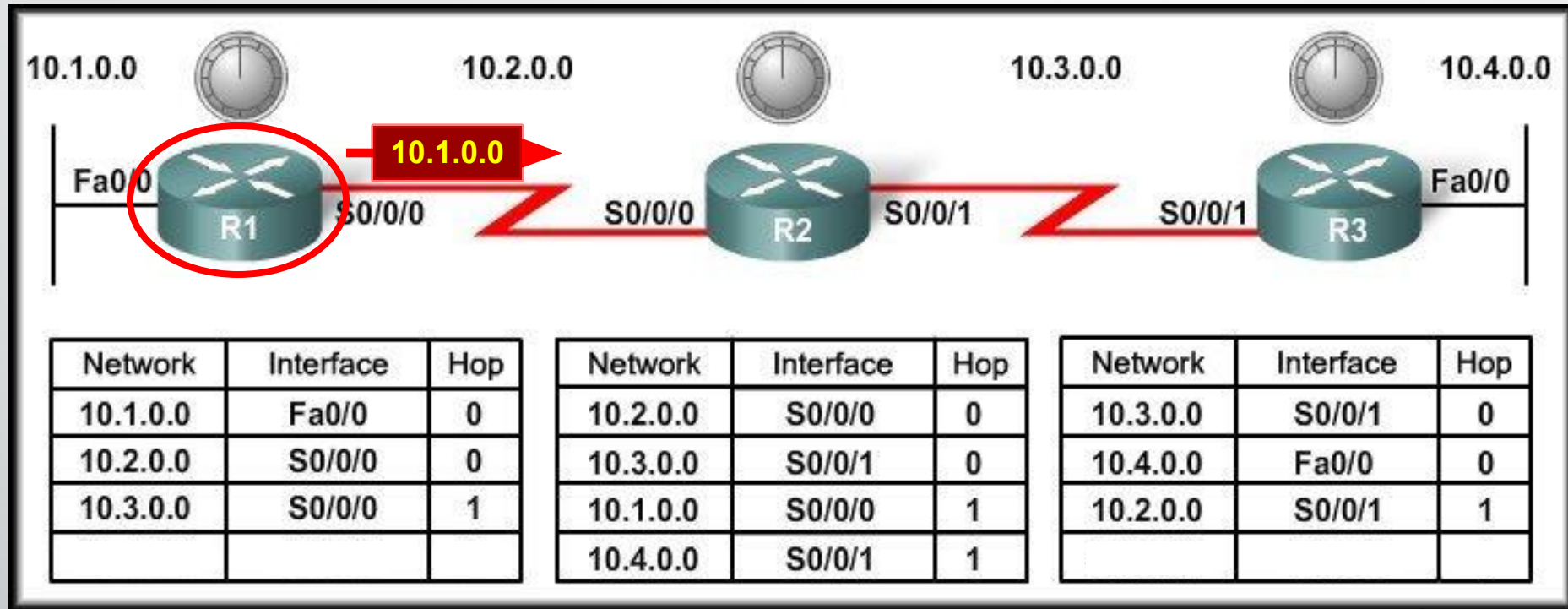
- **R1 Receives** the update from R2 about network 10.3.0.0 and adds it to its routing table.
- **R3 Receives** the update from R2 about network 10.2.0.0 and adds it to its routing table.

Initial Exchange of Routing Information



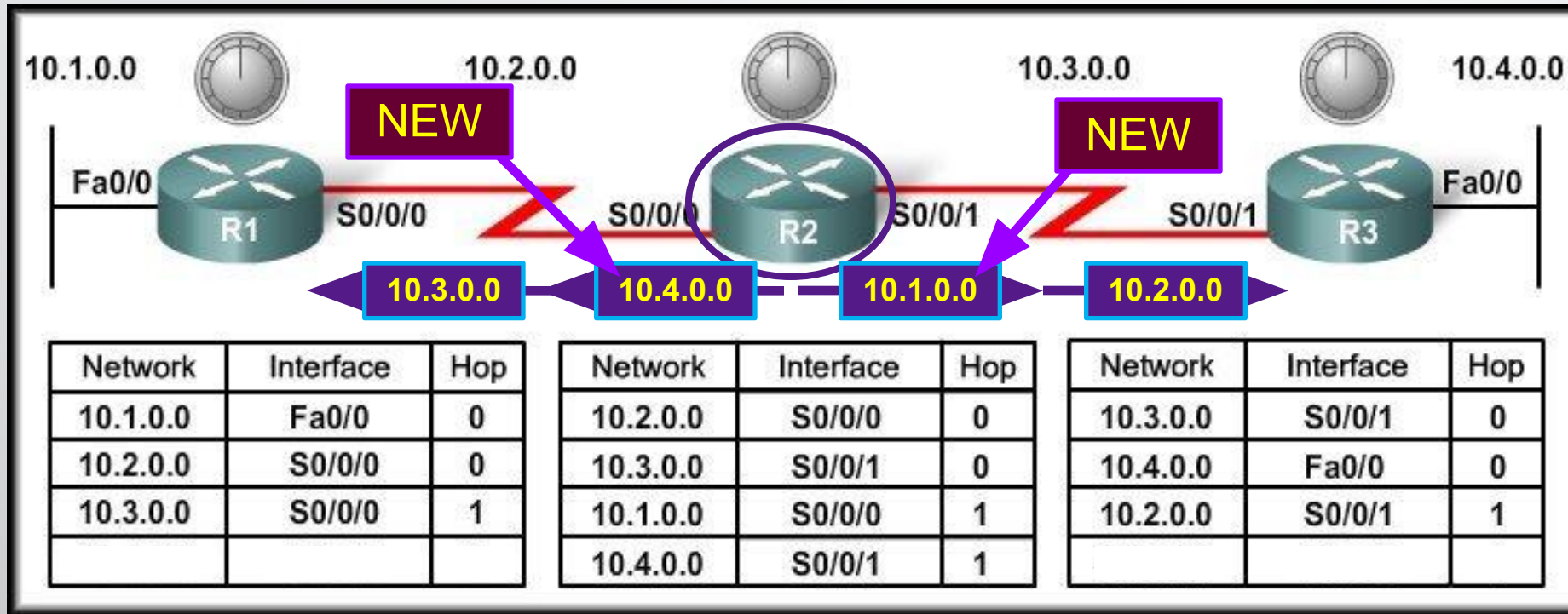
- **R2 Receives** the update from R1 about network 10.1.0.0 and adds it to its routing table.
- **R2 Receives** the update from R3 about network 10.4.0.0 and adds it to its routing table.

Next Exchange of Routing Information



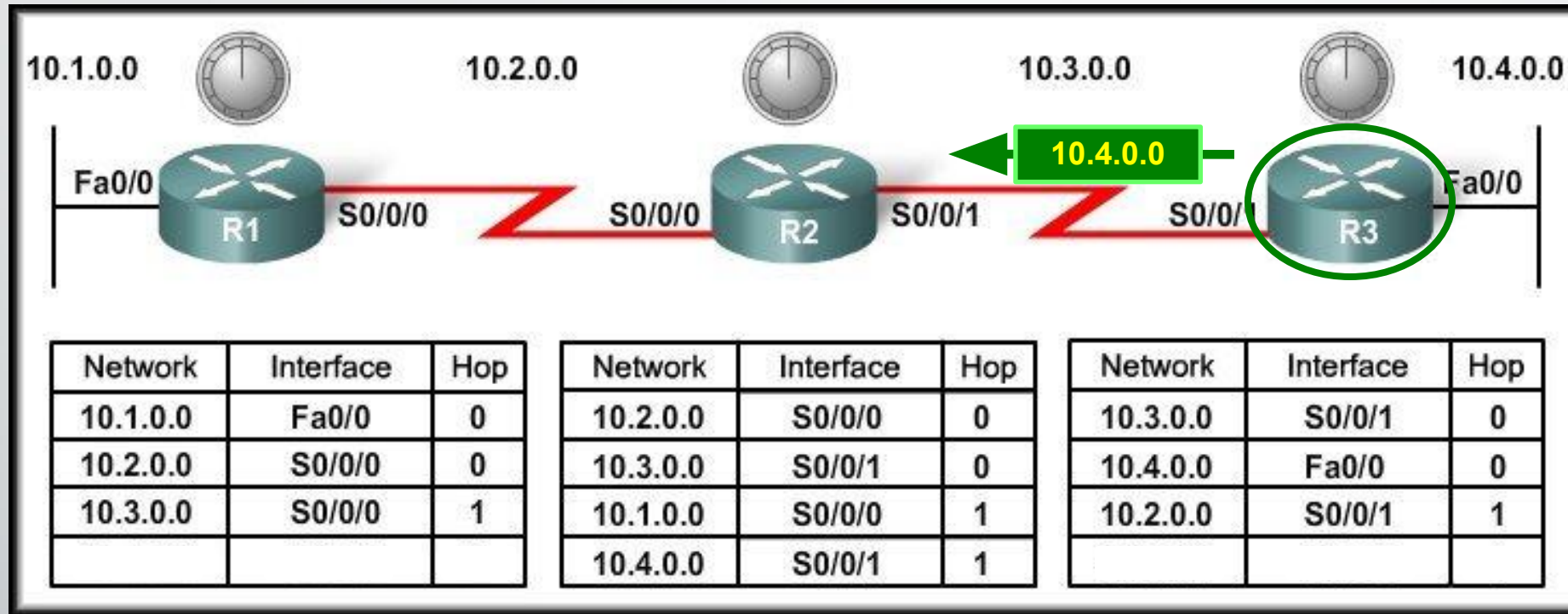
- Sends an update about network **10.1.0.0** out the S0/0/0 interface with a metric of 1 - **AGAIN!**
- **When R2 receives the update**, there is **no change** in information so the update is ignored.

Next Exchange of Routing Information



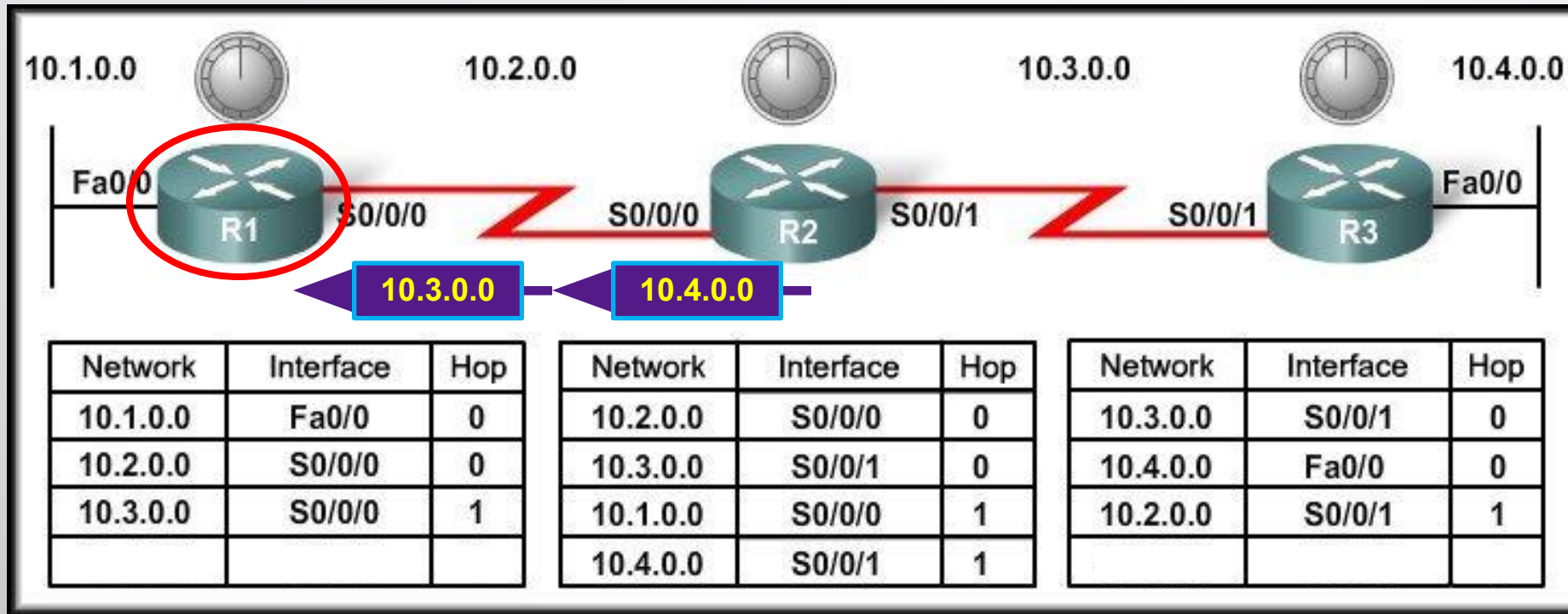
- Sends an update about networks **10.3.0.0** with a metric of 1 and **10.4.0.0** with a metric of 2 out the **Serial o/o/o** interface.
- Similarly sends updates about networks **10.1.0.0** with a metric of 2 and **10.2.0.0** with a metric of 1 out the **Serial o/o/1** interface.

Next Exchange of Routing Information



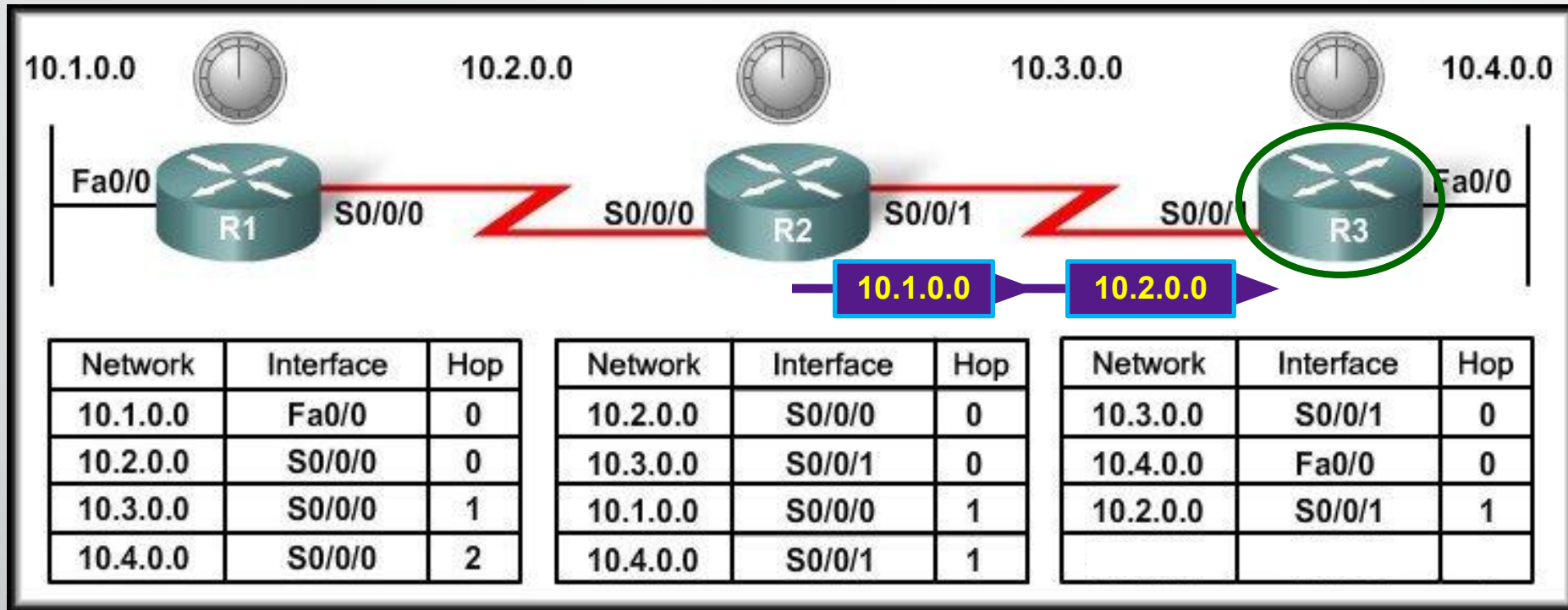
- r Sends an update about network 10.4.0.0 out the So/o/o interface with a metric of 1 - **AGAIN!**
- r When R2 receives the update, there is no change in information so the update is ignored.

Next Exchange of Routing Information



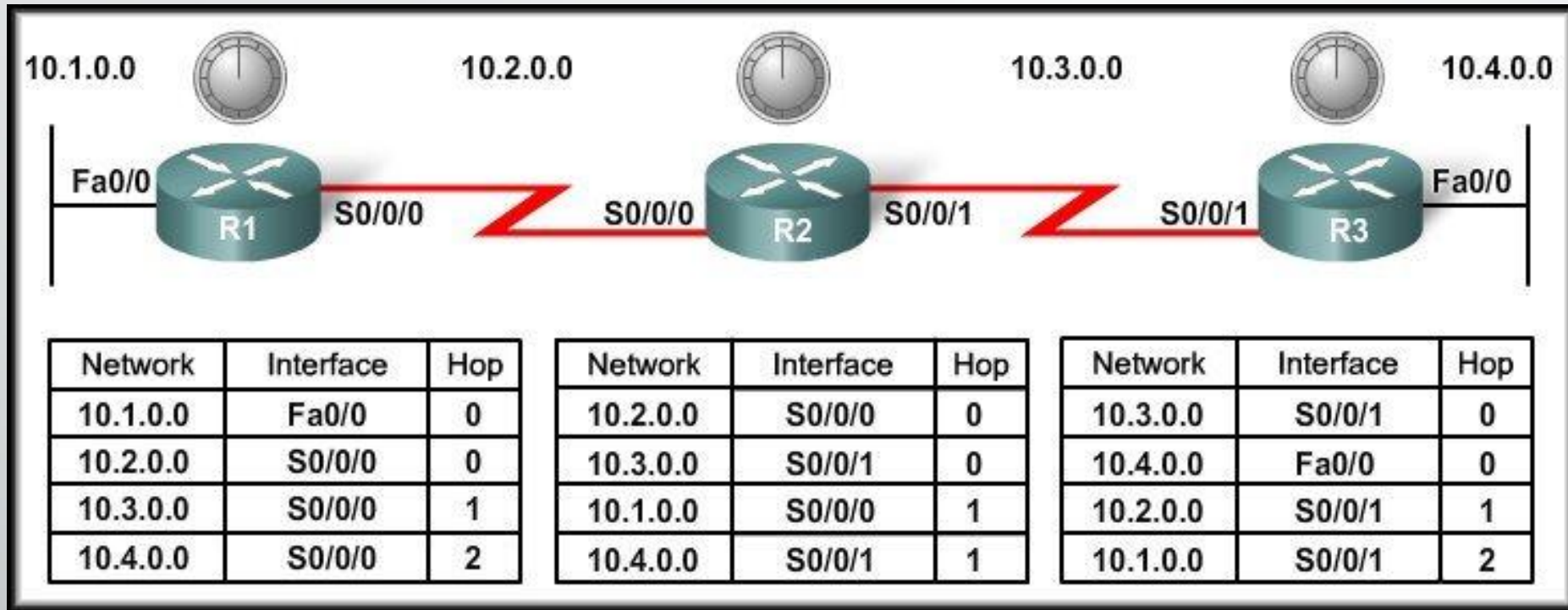
- R1 receives an update from R2 about network 10.3.0.0 and there is no change – update ignored.
- R1 receives an update from R2 about network 10.4.0.0 (new) and adds it to its routing table.

Next Exchange of Routing Information



- R3 receives an update from R2 about network 10.2.0.0 and there is no change – update ignored.
- R3 receives an update from R2 about network 10.1.0.0 (new) and adds it to its routing table.

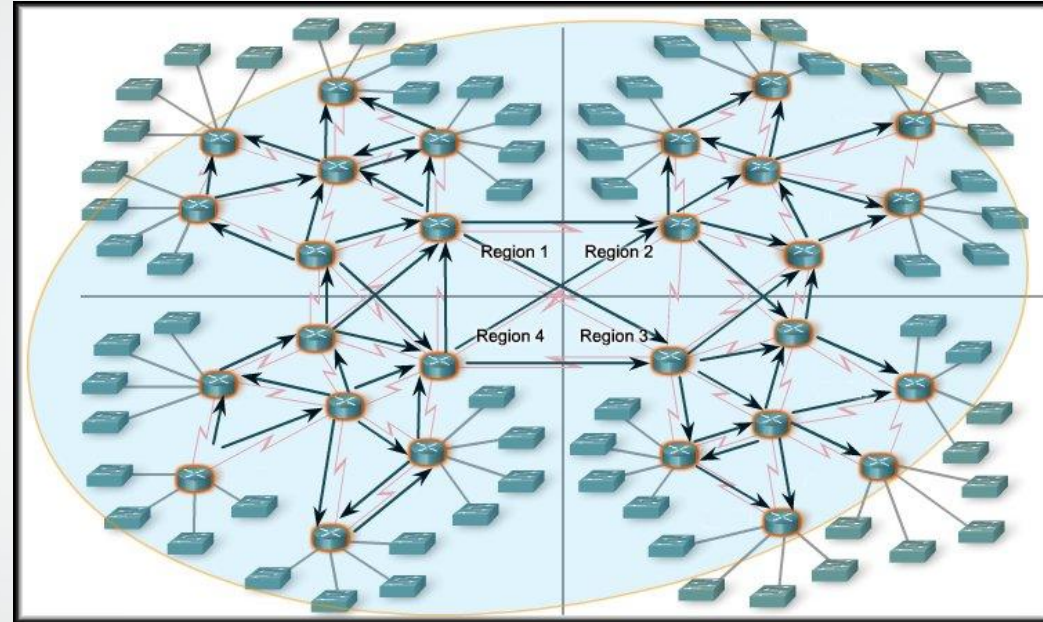
Next Exchange of Routing Information



- The network has **CONVERGED!**
 - All routers now know about all of the networks attached to all of their neighbouring routers.

Convergence

- The amount of time it takes for a network to converge is **directly proportional to the size of that network.**
- Routing protocols are compared based on how fast they can propagate this information - their **speed to convergence.**
- A network is not completely operable until it has converged.
 - Network administrators prefer routing protocols with shorter convergence times.





THE END