



Inspiring Excellence

Network Layer: IPv4 Functions

Lecture 7 | CSE421 – Computer Networks

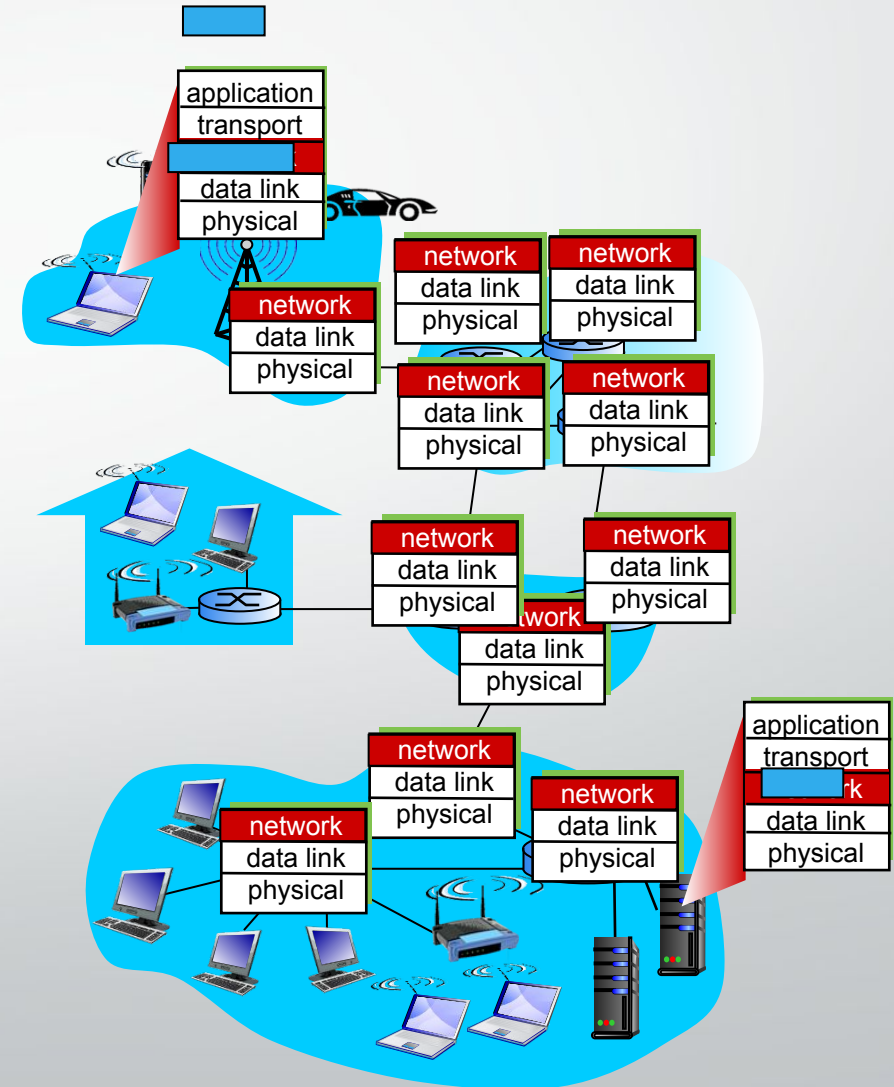
Department of Computer Science and Engineering
School of Data & Science

Objectives

- Short overview of the Network Layer
- Packet Switching: Virtual Circuits & Datagram Networks
- IP Fragmentation & Reassembly
- ICMP
 - Ping
 - Traceroute

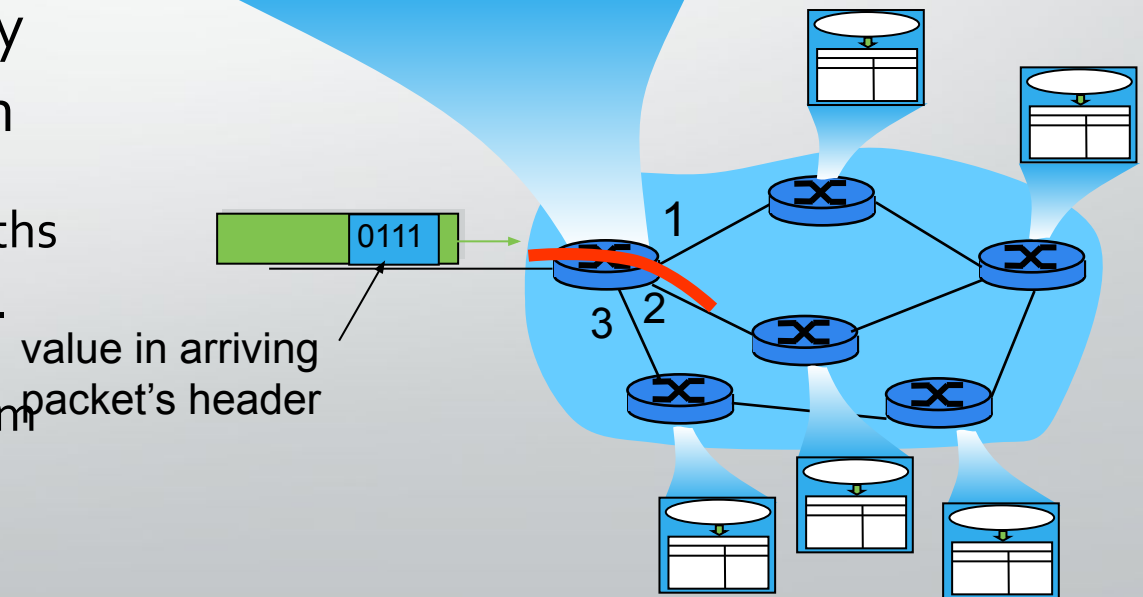
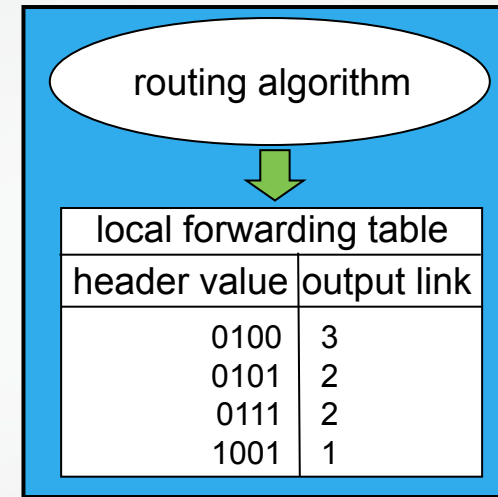
The Network Layer

- Transport segment from sending to receiving host
- On sending side encapsulates segments into packets
- Network layer protocols in every host, router
- Router examines header fields in all IP packets passing through it
- On receiving side, delivers segments to transport layer



Functions of Network Layer

- **Forwarding:** move packets from router's input to appropriate router output
 - Analogy: process of getting through a single interchange
- **Routing:** determine route taken by packets from source to destination
 - The algorithms that calculate the paths are referred to as routing algorithms.
 - Analogy: process of planning trip from source to destination
 - Has various routing algorithms



Packet Switching:

Virtual Circuits

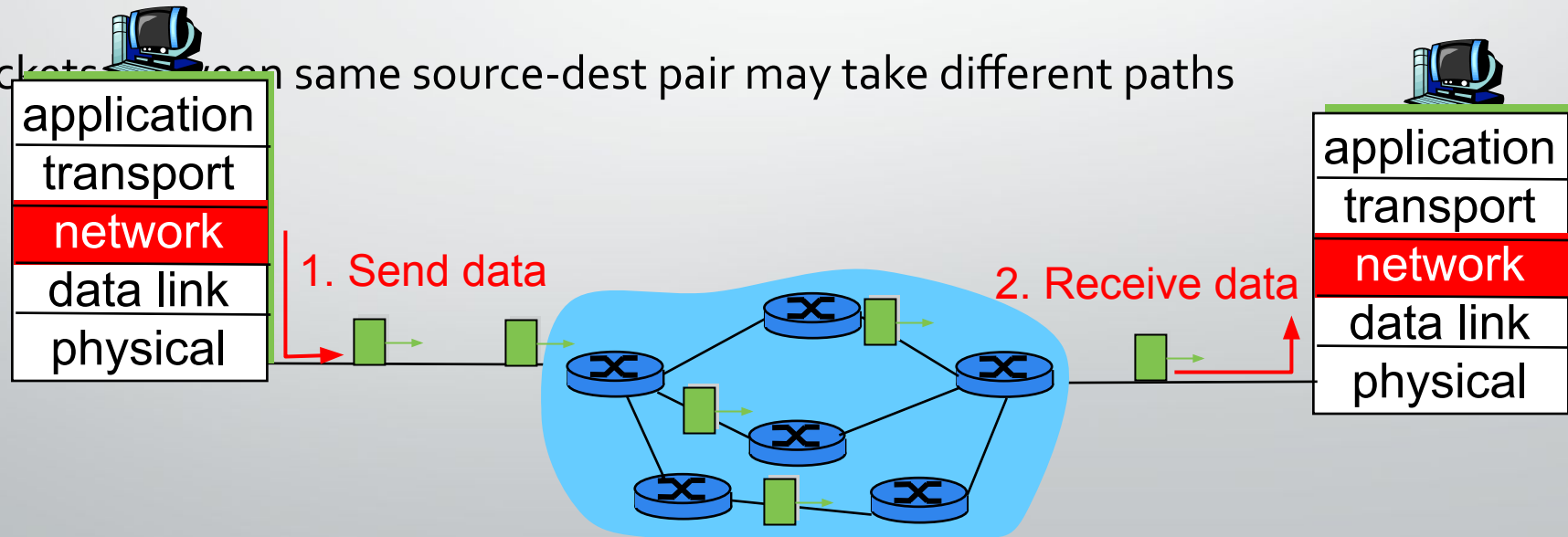
Datagram Network

Connection and Connection-less service

- Datagram network => network-layer connectionless service
- VC network => network-layer connection service
 - analogous to the transport-layer services, but:
 - **service:** host-to-host
 - **no choice:** network provides one or the other
 - **implementation:** in network core

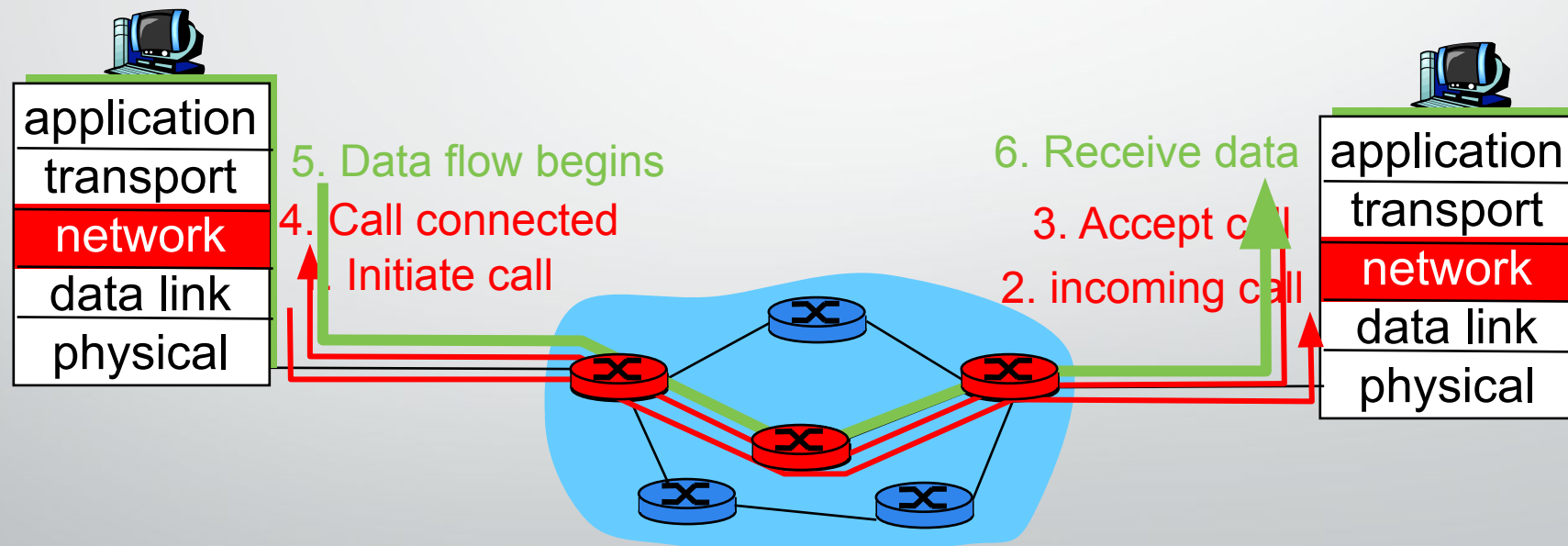
Datagram networks

- No call setup at network layer
- Routers: no state about end-to-end connections
 - no network-level concept of “connection”
- Packets forwarded using destination host address
- packets from same source-dest pair may take different paths



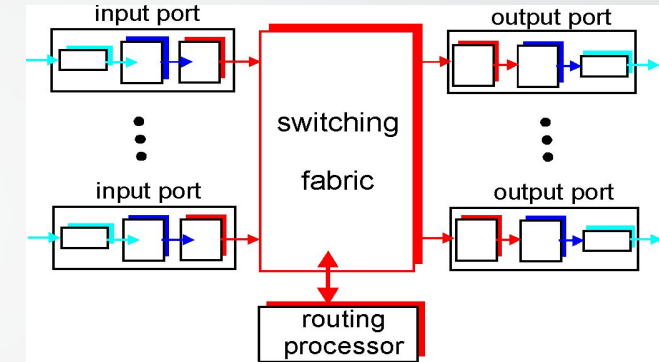
Virtual Circuits: Signaling Protocols

- Used to setup, maintain teardown VC
- Used in ATM, frame-relay, X.25
- Not used in today's Internet

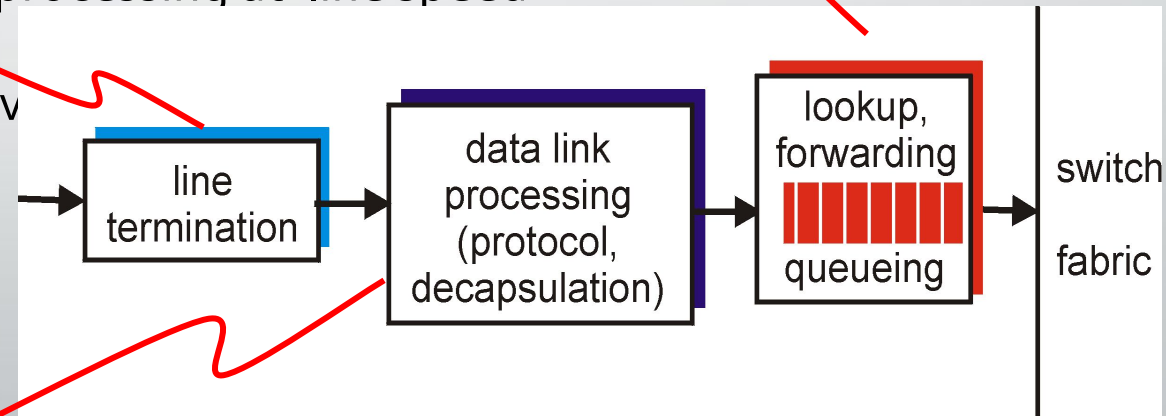


Functions of a Router

- Run routing algorithms/protocol (RIP, OSPF, BGP)
- **Forwarding** datagrams from incoming to outgoing link
- **Decentralized switching:**
 - Given datagram dest., lookup output port using forwarding table in input port memory
 - **Goal:** complete input port processing at 'line speed'
 - **Physical layer:** bit-level reception
 - **Queueing:** datagrams arrive



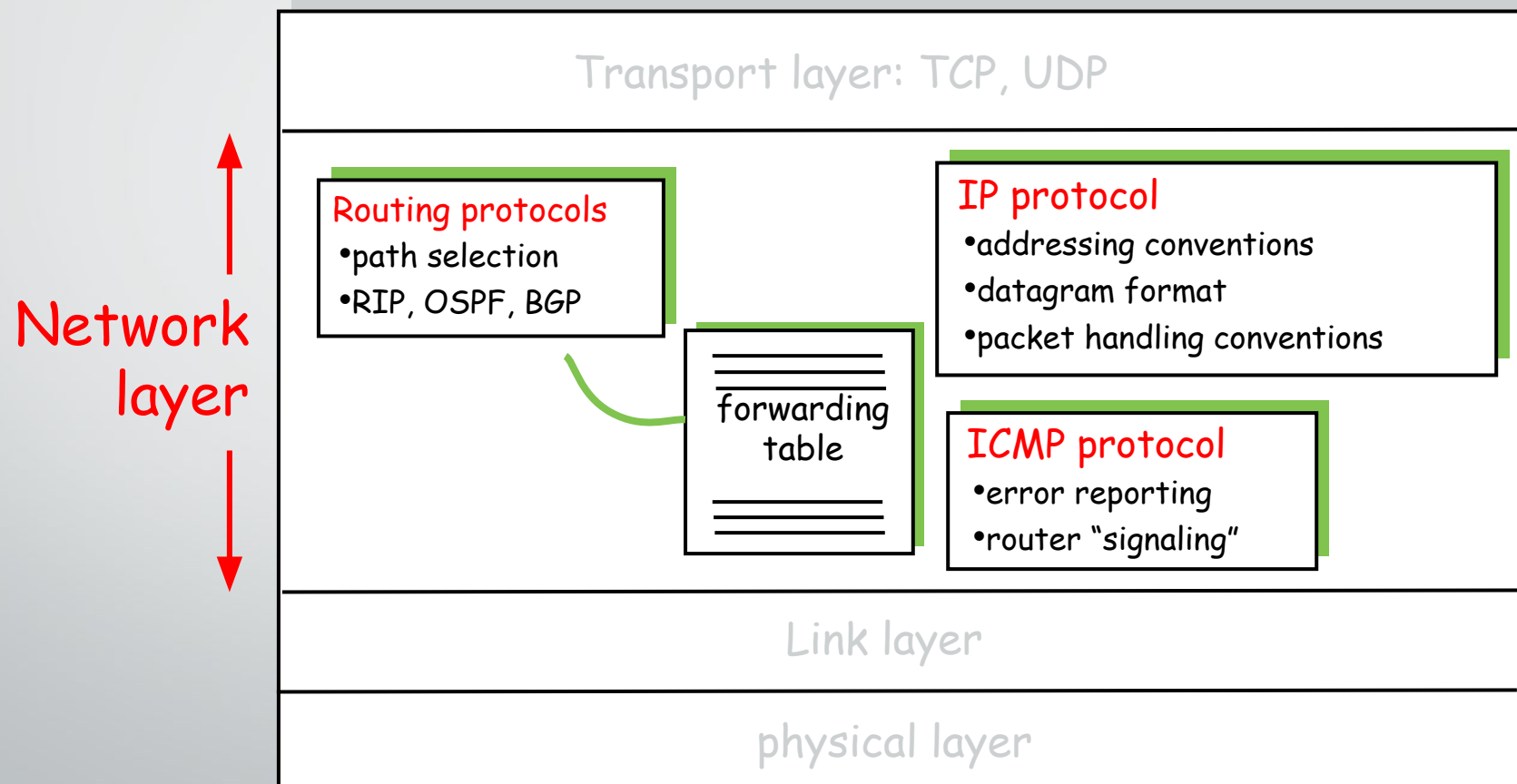
Data link layer:
e.g., Ethernet



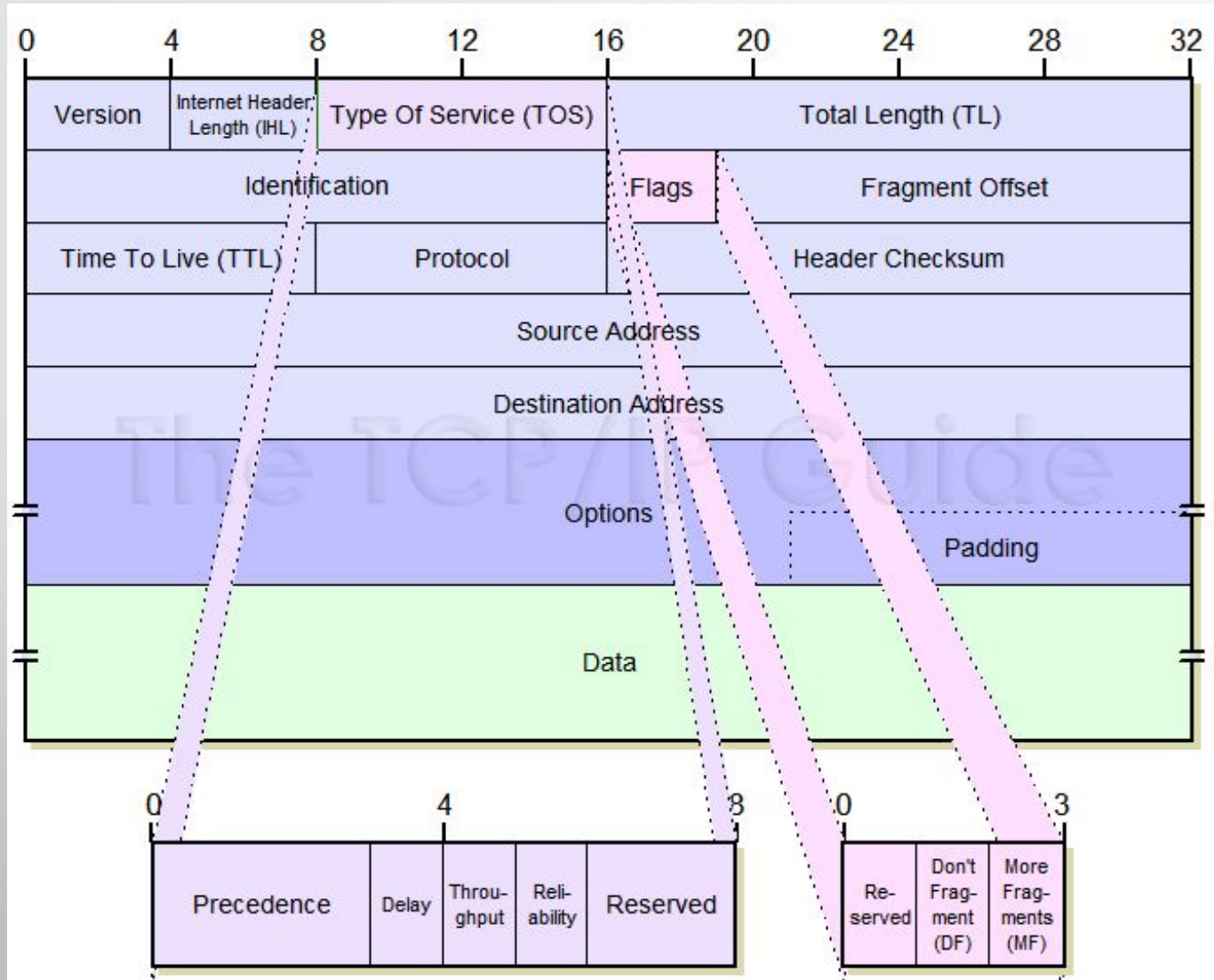
Internet Protocol

Internet Network Layer

- Host, router network layer functions:



IPv4 Datagram Format

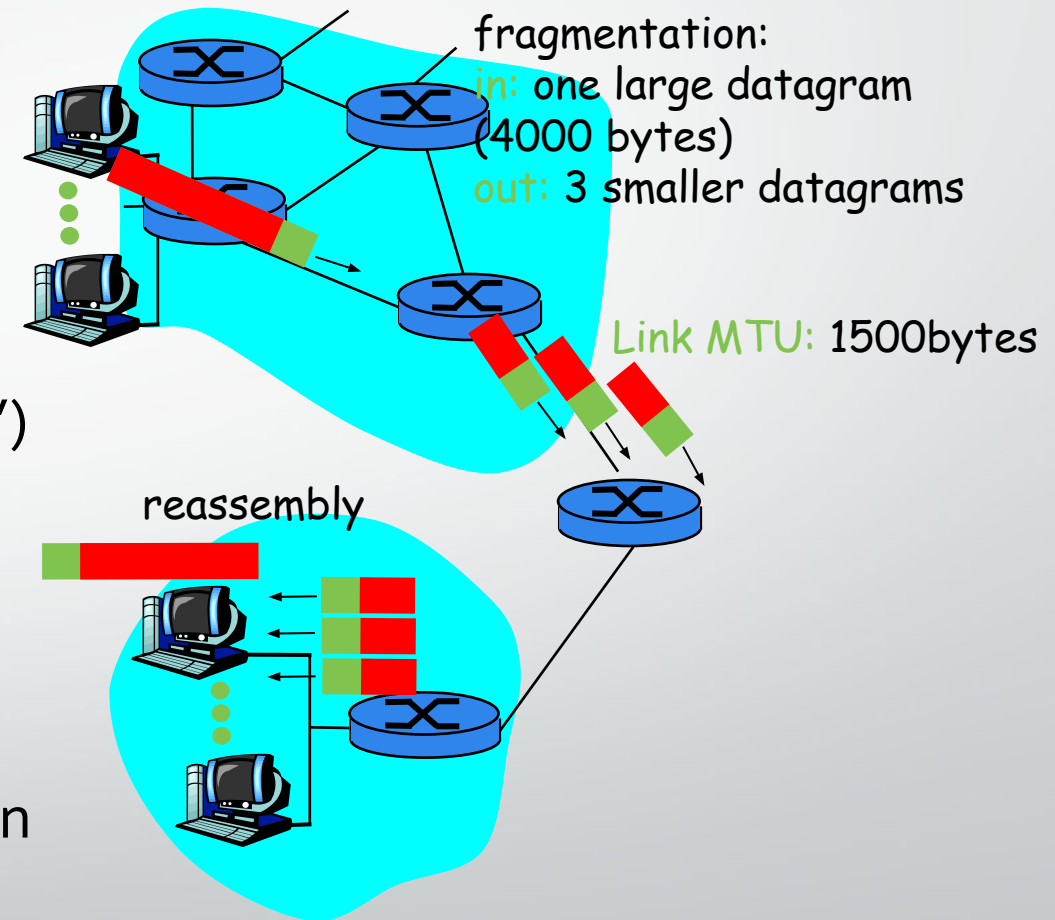


The size of an IP datagram:

- The minimum size is 20 bytes (if you have no data)
- The maximum size is 65,535 bytes

IP Fragmentation & Reassembly

- Network links have **MTU** (max. **transmission unit - max. transfer size**) - largest possible link-level frame.
 - different link types, different MTUs
- Large IP datagram divided ("fragmented") within net
 - one datagram becomes several datagrams
 - "reassembled" only at final destination
 - IP header bits used to identify, order related fragments

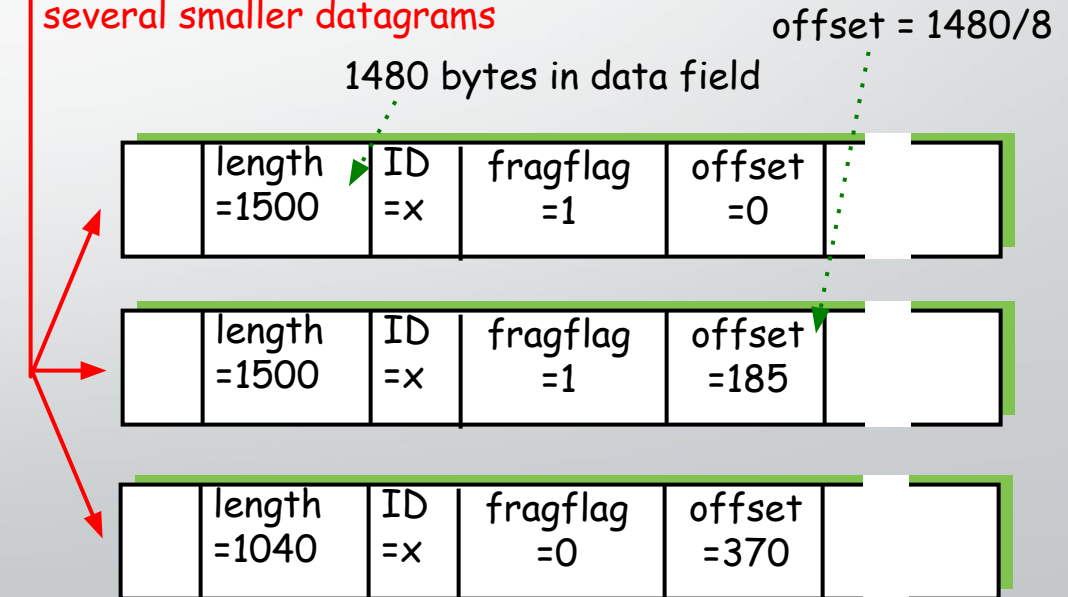


• Example: IP Fragmentation & Reassembly

- 4000 Bytes of datagram
- MTU = 1500 Bytes
 - Header + Data
 - Header size is usually 20 bytes
 - It can differ
- Offset:
 - The value of the offset is measured in units of 8 bytes.
 - This is done because the length of the offset field is only 13 bits long and cannot represent a sequence of bytes greater than 8191.
 - This forces hosts or routers that fragment datagrams to choose the size of each fragment so that the first byte number is divisible by 8.

	length =4000	ID =x	fragflag =0	offset =0	
--	-----------------	----------	----------------	--------------	--

One large datagram becomes several smaller datagrams



IP Fragmentation & Reassembly

$MTU = 20(H) + 1480(D)$
 $5140 = 20(H) + 5120(D)$
 $5120 - 1480 = 3640$ (1st)
 $3640 - 1480 = 2160$ (2nd)
 $2160 - 1480 = 680$ (3rd)

$680 + 20 = 700$

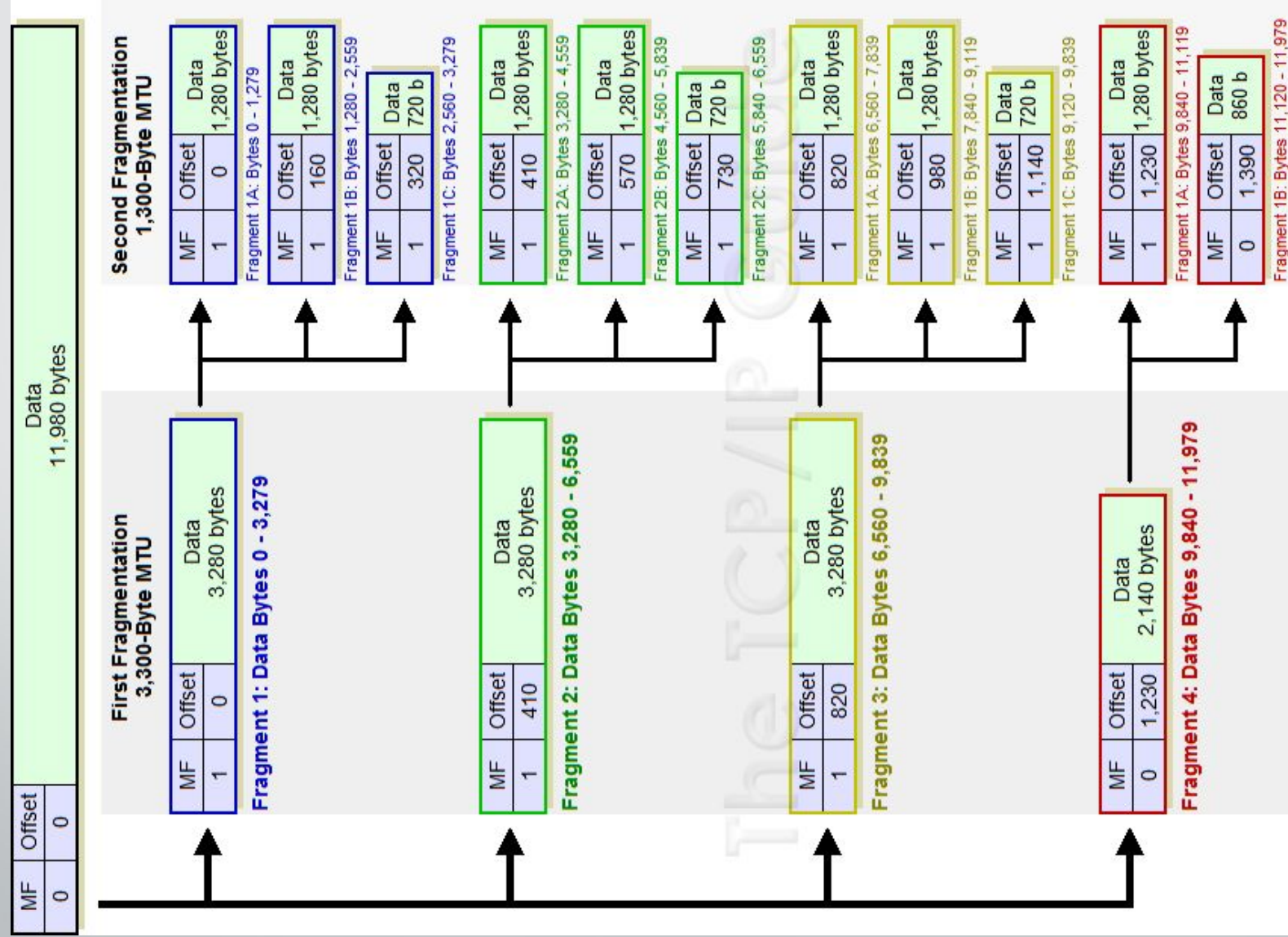
Original IP Datagram

Sequence	Identifier	Total Length	DF May / Don't	MF Last / More	Fragment Offset
0	345	5140	0	0	0

IP Fragments (Ethernet)

Sequence	Identifier	Total Length	DF May / Don't	MF Last / More	Fragment Offset	Data Bytes	Fragment Offset
0-0	345	1500	0	1	0	0 -1479	$0/8=0$
0-1	345	1500	0	1	185	1480-2959	$1480/8=185$
0-2	345	1500	0	1	370	2960-4439	$2960/8=370$
0-3	345	700	0	0	555	4440-5119	$4440/8=555$

IP Fragmentation Example



ICMP

ICMP

- Also known as **Internet Control Message Protocol**
- It is mainly used by the **operating systems** in **IP** network management and administration.
- Used for
 - errors in the underlying communications of network applications
 - availability of remote hosts
 - network congestion
- It does not carry application data, but rather information about the status of the network itself.
- Example of ICMP in practice

PING

Ping

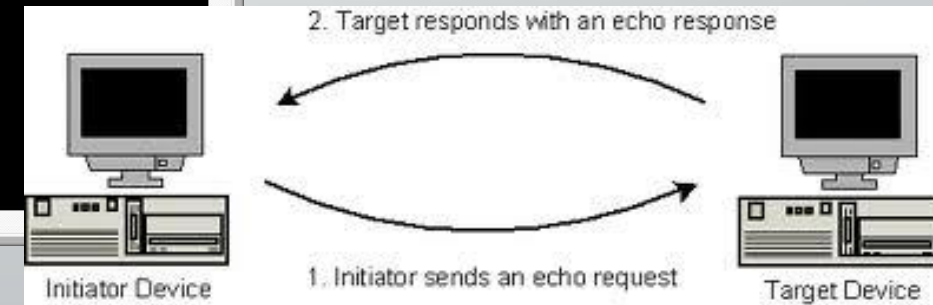
- Utility used to test the reachability of a host.
- Sends Internet Control Message Protocol (ICMP) echo request packets to the target host and waiting for an ICMP response.
- In the process it measures the time from transmission to reception

```
C:\Windows\system32\command.com
C:\USERS\LARRYP~1>ping www.pepperdine.edu

Pinging www.pepperdine.edu [137.159.8.186] with 32 bytes of data:
Reply from 137.159.8.186: bytes=32 time=37ms TTL=114
Reply from 137.159.8.186: bytes=32 time=38ms TTL=114
Reply from 137.159.8.186: bytes=32 time=36ms TTL=114
Reply from 137.159.8.186: bytes=32 time=38ms TTL=114

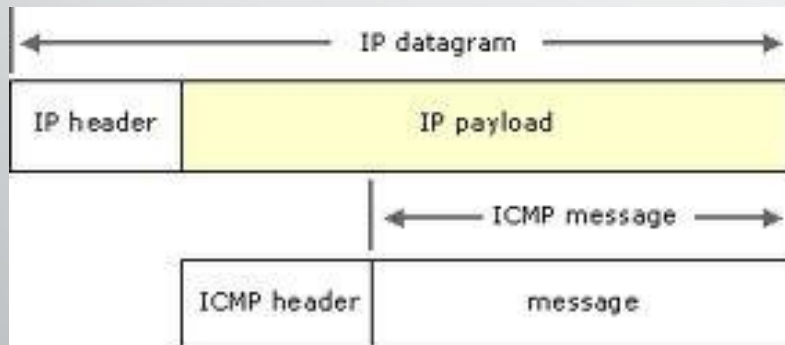
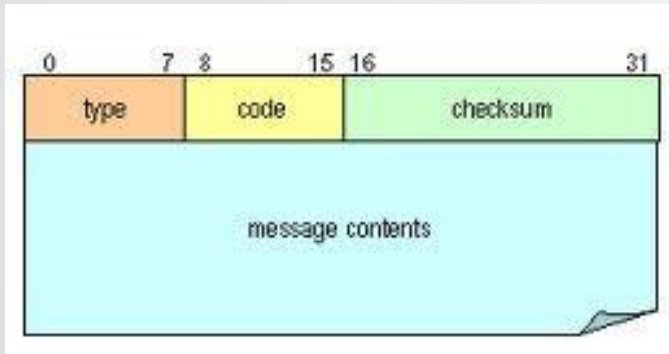
Ping statistics for 137.159.8.186:
    Packets: Sent = 4, Received = 4, Lost = 0 (0% loss),
    Approximate round trip times in milli-seconds:
        Minimum = 36ms, Maximum = 38ms, Average = 37ms

C:\USERS\LARRYP~1>
```



ICMP Message

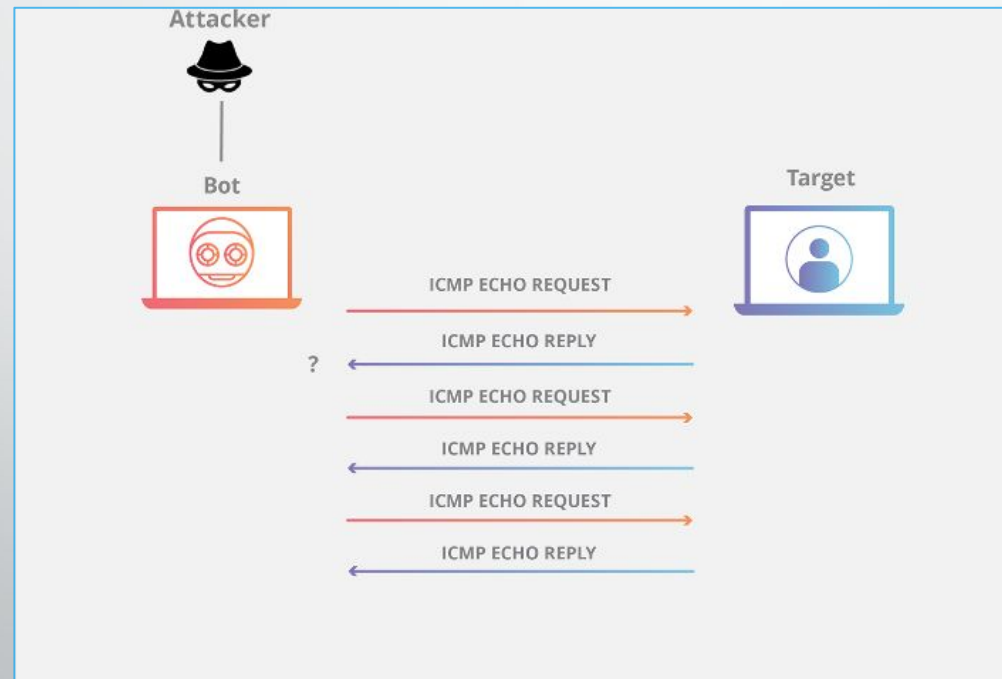
- ICMP message: type, code plus first 8 bytes of IP datagram causing error



Type	Code	description
0	0	echo reply (ping)
3	0	dest. network unreachable
3	1	dest host unreachable
3	2	dest protocol unreachable
3	3	dest port unreachable
3	6	dest network unknown
3	7	dest host unknown
4	0	source quench (congestion control - not used)
8	0	echo request (ping)
9	0	route advertisement
10	0	router discovery
11	0	TTL expired
12	0	bad IP header

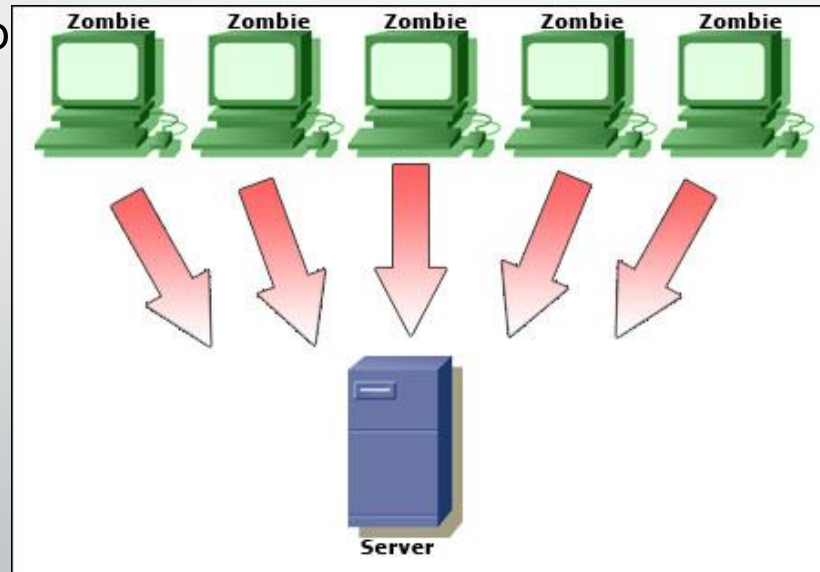
Ping Attacks

- ICMP PING flood attack /DOS Attack :
 - It uses the ICMP echo command to flood large amounts of data packets to the victim's computer in an attempt to overload it.
 - Another type of DOS Attack : Deny to give service by replying with false message.



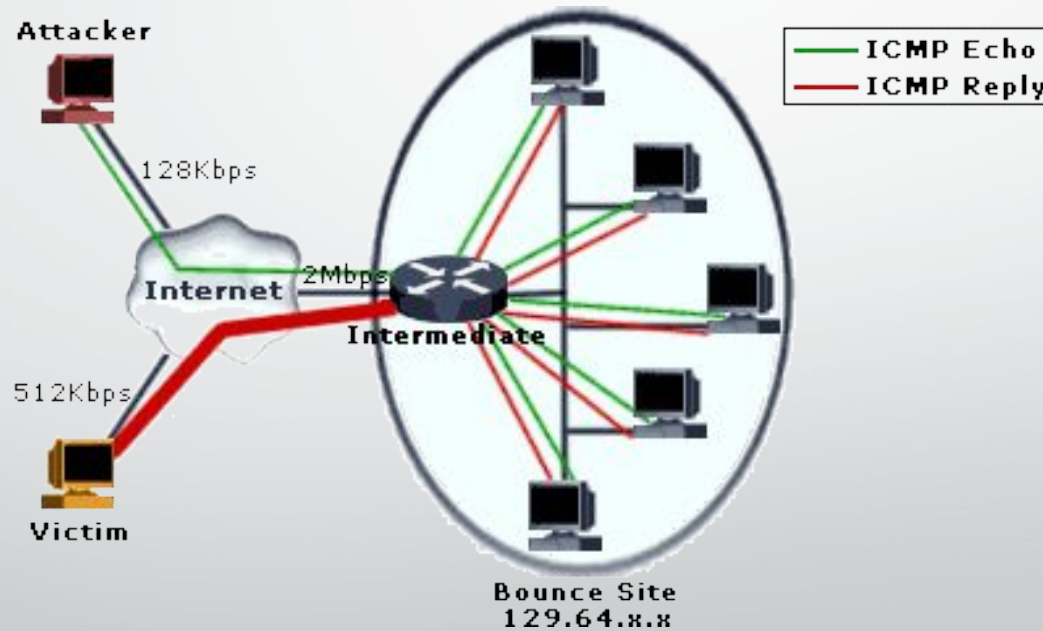
Ping Attacks

- ICMP DDOS attack – Zombie Attack:
 - Much like the ping flood method, only multiple computers are being used.
 - In this instance, the computers that are being used may or may not be aware of the fact that they are attacking a website or network.
 - Trojans and viruses commonly give the hacker control of a computer, and thus, the ability to use them for attacks. These computers are called zombies.



Ping Attacks

- ICMP DDOS attack – Packet magnification (or ICMP Smurf):
 - An attacker sends forged ICMP echo packets to vulnerable networks' broadcast addresses. All the systems on those networks send ICMP echo replies to the victim.



Traceroute

- Tool used to trace path from source to destination host.
- The IP address and domain name (if there is one) of each router is returned to the client.
- Commands:
 - Unix: **traceroute**
 - Cisco IOS: **traceroute** (t)
 - DOS: **tracert**

Hop 1: User LAN router

Hops 2-7: Verizon (ISP) network

Hops 8-10: the Yahoo LAN

```
C:\Windows\system32\command.com
C:\USERS\LARRYP~1>tracert www.yahoo.com

Tracing route to www-real.wa1.b.yahoo.com [209.131.36.158]
over a maximum of 30 hops:

  1      1 ms      1 ms      4 ms  192.168.1.1
  2     37 ms     36 ms     39 ms  L100.LSANCA-DSL-14.verizon-gni.net [71.105.96.1]
  3     35 ms     35 ms     35 ms  P15-2.LSANCA-LCR-03.verizon-gni.net [130.81.44.32]
  4     39 ms     39 ms     38 ms  so-6-1-2-0.LAX01-BB-RTR1.verizon-gni.net [130.81.28.225]
  5     47 ms     47 ms     47 ms  so-5-3-0-0.SJC01-BB-RTR1.verizon-gni.net [130.81.19.10]
  6     46 ms     47 ms     46 ms  130.81.17.229
  7     54 ms     47 ms     49 ms  130.81.14.90
  8     48 ms    129 ms     50 ms  ae0-p170.msr2.sp1.yahoo.com [216.115.107.81]
  9     90 ms     48 ms    112 ms  te-8-1.bas-a1.sp1.yahoo.com [209.131.32.17]
 10     48 ms     50 ms     49 ms  f1.www.vip.sp1.yahoo.com [209.131.36.158]

Trace complete.

C:\USERS\LARRYP~1>
```


Traceroute: Another example

Hop 1: User LAN router

Hops 2-4: Verizon network (a backbone ISP)

Hops 5-6: Altnet (a backbone ISP)

Hops 7-11: Level 3 (a backbone ISP)

Hops 12-14: the Google LAN

```
C:\Windows\system32\COMMAND.com
C:\USERS\LARRYP~1>tracert www.google.com

Tracing route to www.l.google.com [74.125.19.147]
over a maximum of 30 hops:

  1    3 ms    1 ms    1 ms  192.168.1.1
  2   38 ms   37 ms   37 ms  L100.LSANCA-DSL-14.verizon-gni.net [71.105.96.11]
  3   38 ms   34 ms   36 ms  P1-3.LSANCA-LCR-03.verizon-gni.net [130.81.35.81]
  4   34 ms   37 ms   34 ms  so-6-1-2-0.LAX01-BB-RTR1.verizon-gni.net [130.81.28.225]
  5   37 ms   35 ms   38 ms  0.so-1-3-0.XL3.LAX15.ALTER.NET [152.63.114.145]
  6   36 ms   36 ms   40 ms  0.ge-6-0-0.BR2.LAX15.ALTER.NET [152.63.116.149]
  7   38 ms   40 ms   40 ms  xe-11-0-0.edge1.SanJose3.level3.net [4.68.111.249]
  8   46 ms   38 ms   49 ms  ae-73-70.ebr3.LosAngeles1.Level3.net [4.69.144.116]
  9   47 ms   55 ms   52 ms  ae-2.ebr3.SanJose1.Level3.net [4.69.132.91]
 10   68 ms   54 ms  126 ms  ae-63-63.csw1.SanJose1.Level3.net [4.69.134.226]
 11   72 ms   45 ms  115 ms  ae-1-69.edge1.SanJose1.Level3.net [4.68.18.14]
 12  137 ms   51 ms   49 ms  GOOGLE-INC.edge1.SanJose1.Level3.net [4.79.43.146]
 13   49 ms   49 ms   54 ms  209.85.251.98
 14   47 ms   47 ms   46 ms  nuq04s01-in-f147.1e100.net [74.125.19.147]

Trace complete.
```

Traceroute: Request Timed Out

This message indicates that the router security settings keep it from revealing its identity or the router and connection are slow.

```
* * * Request timed out.
```

The End