# APPLIED MECHANISM DESIGN FOR SOCIAL GOOD

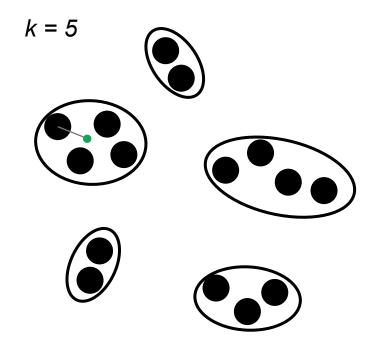
**JOHN P DICKERSON & MARINA KNITTEL** 

Lecture #23 - 04/18/2022 Lecture #24 - 04/20/2022

CMSC498T Mondays & Wednesdays 2:00pm – 3:15pm



**Clustering** is the problem of grouping data based off of the distance (or similarity score) between points.

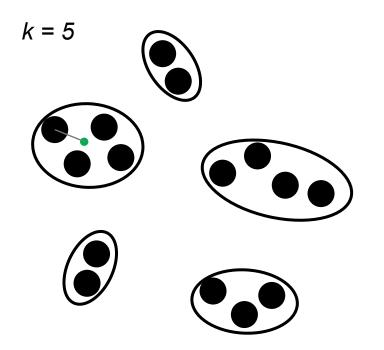


*k-center:* minimize the maximum distance from a point to its cluster center.

*k-median:* minimize the sum of distances from points to cluster centers.

*k-means:* minimize the sum of squares of distances from points to cluster centers.

**Clustering** is the problem of grouping data based off of the distance (or similarity score) between points.



Generally, we minimize:

$$\sum_{x \in X} ||\varphi(x) - x||^p$$

X is our point set.

 $\varphi$  tells us the center for a point.

p defines what we maximize.

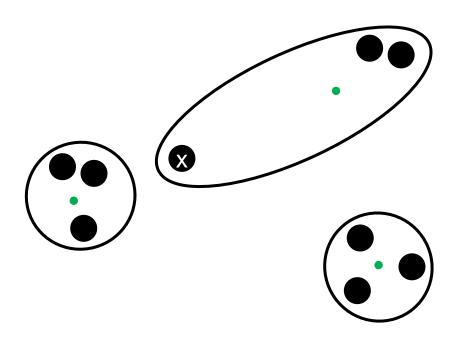
k-center:  $p = \infty$ 

k-median: p = 1

k-center: p = 2

. . .

Clustering *x* poorly is only one of many costs in *k-median* since we sum over many points. But for *k-center*, since it is the furthest point from the cluster center, it is very costly, in fact it defines the *k-center* cost.



We can write an integer linear program to assign points to clusters.

Let  $C = c_1, c_2, ..., c_k$  be the cluster centers.

Let  $x_i$  denote that x is assigned to center i.

This may look nonlinear, but it's a constant!

Objective: minimize 
$$\sum_{c_i \in C} \sum_{x \in X} x_i \times ||c_i - x||^p$$

Constraints: 
$$\sum_{c_i \in C} x_i = 1$$
 for all  $x \in X$  (only assign to 1 center)  $x_i \in \{0,1\}$  for all  $x \in X$  and  $c_i \in C$ 

For *k-centers*, we do it a bit differently. We do *not* put anything in the objective. We guess *R* as an upper bound on the distance to centers. *R* is the cost of the solution. We use binary search to find the smallest *R* with a feasible solution.

Objective: None!

Constraints:  $\sum_{c_i \in C \cap B_R(x)} x_i = 1$  for all  $x \in X$  (only assign to 1 center)  $x_i \in \{0,1\}$  for all  $x \in X$  and  $c_i \in C$ 

Note:  $B_R(x)$  is the ball of radius R around x (i.e., points within R from x).

### **DISPARATE IMPACT**

#### Griggs vs Duke Power Co:

- North Carolina, 1970
- Required high school diploma and standardized testing for promotion
- Sued for discrimination
- Ruled discriminatory by SC because it had "a disproportionate and adverse impact on certain individuals"





## DISPARATE IMPACT IN CLUSTERING

**Applying to ML:** Ensure the impact of a system across protected groups is proportionate. "Group fairness."

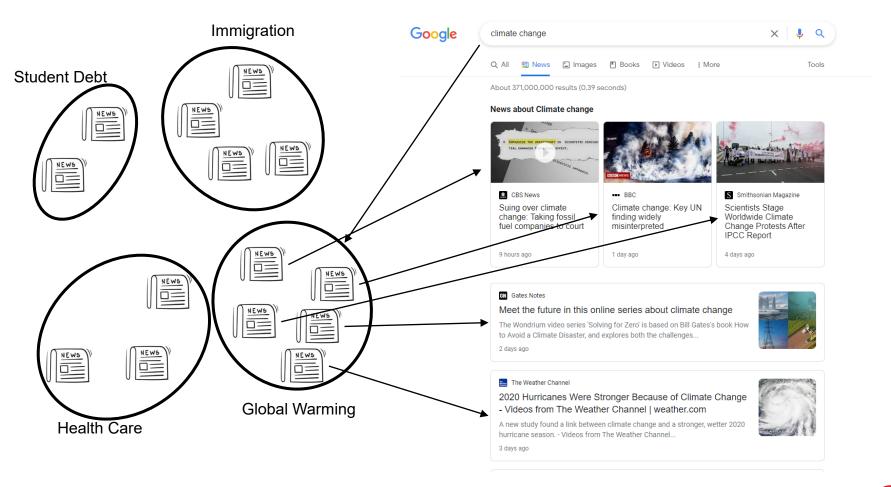




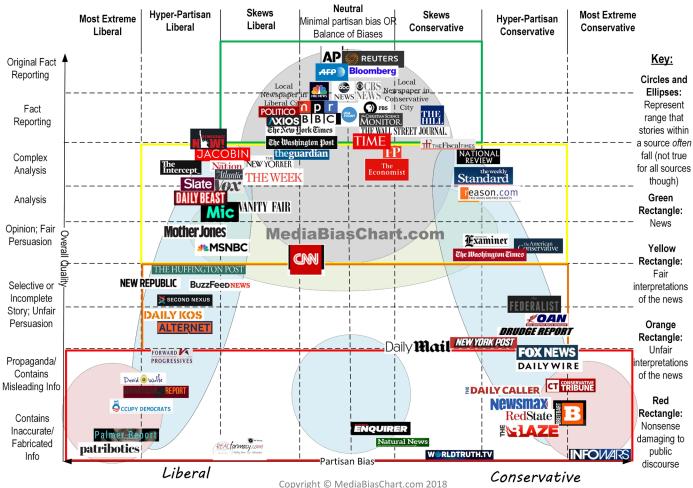
#### Applying to clustering:

- How do we measure the impact of a system on a protected group?
  - How many individuals are in a cluster
- How do we prevent disparate impact?
  - Ensure the number of individuals from any group in any cluster is proportionate to group size.

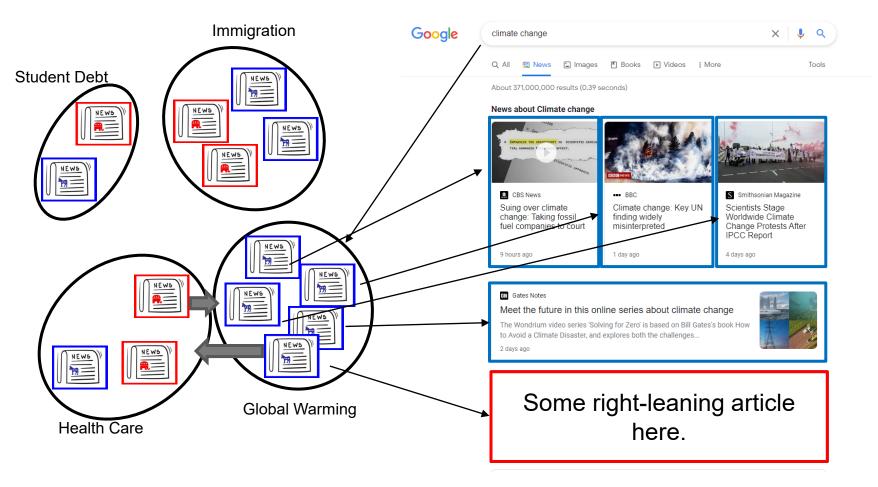
# CLUSTERING NEWS ARTICLES



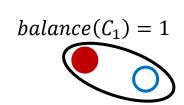
### RUNNING EXAMPLE: NEWS SEARCH

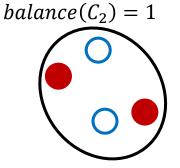


# CLUSTERING NEWS ARTICLES



### **GROUP FAIRNESS IN CLUSTERING: FORMAL**



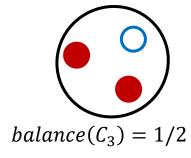


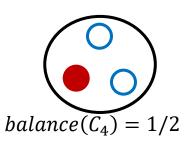
#### For a cluster C:

$$balance(C) := \min\left(\frac{\#red(C)}{\#blue(C)}, \frac{\#blue(C)}{\#red(C)}\right)$$

#### For a clustering S:

$$balance(S) \coloneqq \min_{C \in S} balance(C)$$





We want balance to be high (close to 1).

KEY:

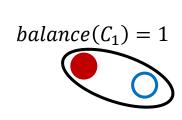
= Right-leaning article

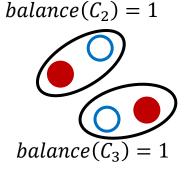
= Left-leaning article

Before: balance(S) = 1/2

After:

### **GROUP FAIRNESS IN CLUSTERING: FORMAL**



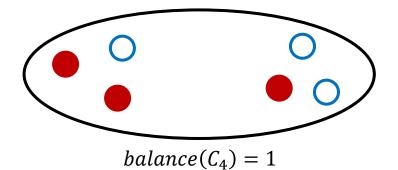


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#### For a clustering S:

$$balance(S) := \min_{C \in S} balance(C)$$



We want balance to be high (close to 1).

KEY:



Before: balance(S) = 1/2

After: balance(S) = 1

#### Let:

This is the best balance we can achieve

b/r = 1/3 is the minimum ratio between reds and blues

R = number of remaining reds = 10

B = number of remaining blues = 5

#### Iteratively...

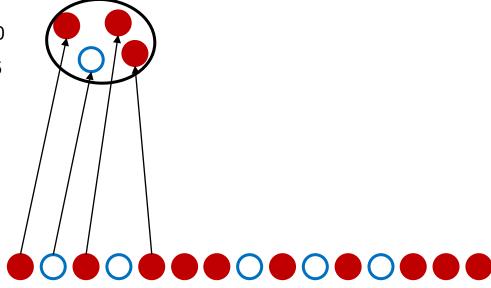


IF (R-B) ≥ (r-b): make a cluster of r reds and b blues.

ELSE: use (R-B)+b red and b blue points.

When r=b: simply match

points



#### Let:

This is the best balance we can achieve

b/r = 1/3 is the minimum ratio between reds and blues

R = number of remaining reds = 7

B = number of remaining blues = 4

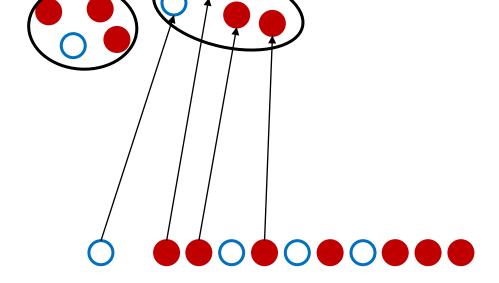
#### Iteratively...



IF (R-B) ≥ (r-b): make a cluster of r reds and b blues.

ELSE: use (R-B)+b red and b blue points.

When r=b: simply match points



#### Let:

This is the best balance we can achieve

b/r = 1/3 is the minimum ratio between reds and blues

R = number of remaining reds = 4

B = number of remaining blues = 3

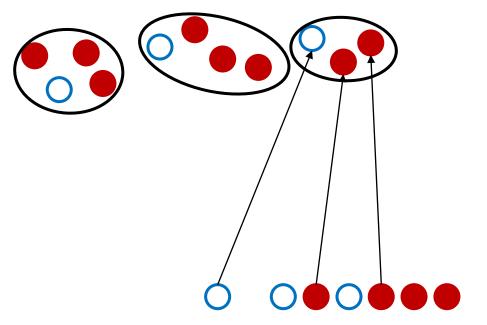
#### Iteratively...

IF (R-B) ≥ (r-b): make a cluster of r reds and b blues.



ELSE: use (R-B)+b red and b blue points.

When r=b: simply match points



#### Let:

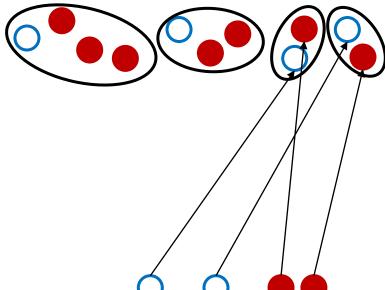
This is the best balance we can achieve

b/r = 1/3 is the minimum ratio between reds and blues

R = number of remaining reds = 2

B = number of remaining blues = 2





#### Iteratively...

IF (R-B) ≥ (r-b): make a cluster of r reds and b blues.

ELSE: use (R-B)+b red and b blue points.



When r=b: simply match points

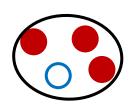
#### Let:

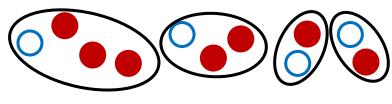
This is the best balance we can achieve

b/r = 1/3 is the minimum ratio between reds and blues

R = number of remaining reds = 0

B = number of remaining blues = 0





#### Iteratively...

IF  $(R-B) \ge (r-b)$ : make a cluster of r reds and b blues.

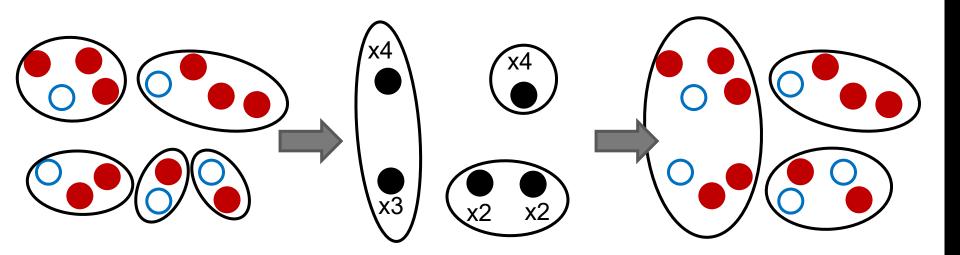
ELSE: use (R-B)+b red and b blue points.

When r=b: simply match points

**Idea:** Find a "fairlet decomposition" with small fairlets. Find a way to merge them into a good, appropriately-sized clustering.

# USING FAIRLETS FOR FAIR CLUSTERING

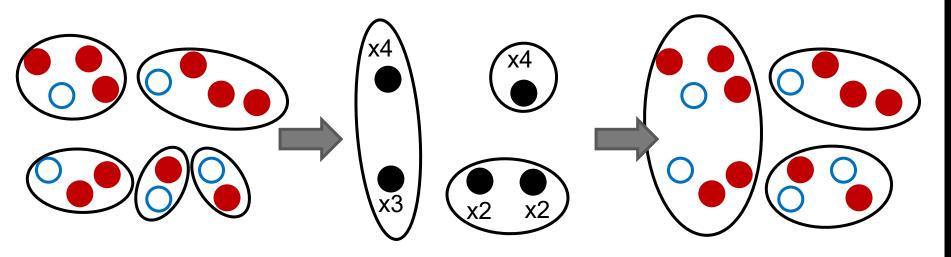
- 1. Find a *good* fairlet decomposition.
- 2. Replace each fairlet with a single point. Duplicate that point according to the fairlet size.
- 3. Run a "vanilla" clustering on these points.
- 4. Apply this clustering to the original points.



# USING FAIRLETS FOR FAIR CLUSTERING

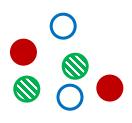
**Let:** Y be our fairlet decomposition, Y' be our transformed point set, and S be our final clustering.

**Theorem:** for k-median and k-center: cost(S) = cost(Y) + cost(Y')



# GENERALIZING TO MORE COLORS





For each color i, we get bounds  $\alpha_i$ ,  $\beta_i$  to bound the proportional representation of i.

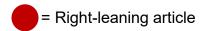




A cluster *C* is fair if for all colors *i*:

$$\alpha_i \times |C| \le i(C) \le \beta_i \times |C|$$

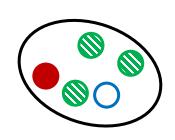
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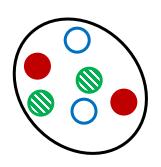


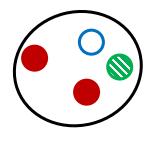
= Green party-leaning article

Where i(C) is the number of i colored points in C.

## GENERALIZING TO MORE COLORS









KEY:



Is this clustering fair for the following values?

$$\alpha_{red} = 1/4 \times \beta_{red} = 1/2 \checkmark$$

$$\alpha_{green}$$
= 1/3  $\times$   $\beta_{green}$  = 1/2  $\times$ 

$$\alpha_{blue} = 1/4 \checkmark \quad \beta_{blue} = 1/2 \checkmark$$

What about...

$$\alpha_{red} = 1/5 \checkmark \quad \beta_{red} = 1/2 \checkmark$$

$$\alpha_{green}$$
= 1/4 $\checkmark$   $\beta_{green}$  = 3/5 $\checkmark$ 

$$\alpha_{blue} = 1/4 \checkmark \quad \beta_{blue} = 1/2 \checkmark$$

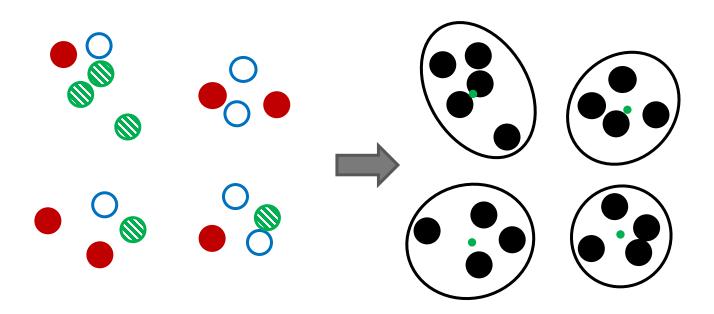
# ILP FOR FAIR CLUSTERING

For this approach, we will use a linear program as an intermediate step, but it will not be used to solve the whole problem.

- 1. Ignore fairness. Find a good vanilla clustering. This gives cluster centers.
- 2. Use a LP relaxation of an ILP to fairly assign points to centers.
- 3. Round the LP to an integer solution.

# ILP FOR FAIR CLUSTERING STEP 1

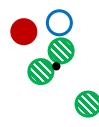
Step 1: Ignore fairness. Find a good vanilla clustering.



### ILP FOR FAIR CLUSTERING STEP 2

**Step 2:** Use a LP relaxation of an ILP to fairly assign points to centers.

Let  $C = c_1, c_2, ..., c_k$  be our given cluster centers. Let  $x_i$  denote that x is assigned to center i. Let R be our distance guess.





Objective: None!





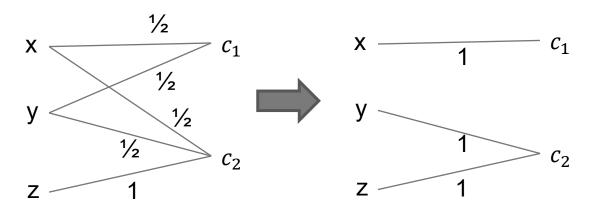
Constraints:  $\sum_{c_i \in C \cap B_R(x)} x_i = 1$  for all  $x \in X$   $\alpha_i \sum_{x \in X} x_i \leq \sum_{j \ colored \ x \in X} x_i \leq \beta_i \sum_{x \in X} x_i$  for all  $c_i \in C$  and colors j  $0 \leq x_i \leq 1$  for all  $x \in X$  and  $c_i \in C$ 

### ILP FOR FAIR CLUSTERING STEP 3

Now we have a set of cluster centers and a fair *fractional* assignment of points to centers.

In other words, x can be  $\frac{1}{2}$  assigned to one center and  $\frac{1}{2}$  assigned to another!

**Step 3:** Round these to whole numbers (i.e., a real assignment).



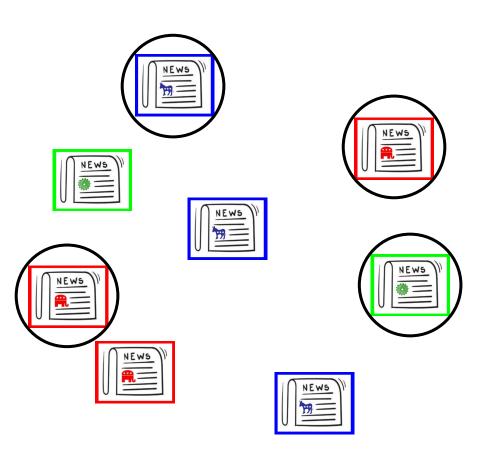
# FAIR DATA SUMMARIZATION

**Data summarization:** Select a subset of points that "represent" your data.

**Method:** Do a clustering, and select the cluster centers.

**Fairness:** for every color i, there must be at least  $k_i$  representatives of that color.

Say 
$$k_{red} = k_{blue} = k_{green} = 1$$



FAIRNESS THROUGH PROPORTIONALITY

Say you want to build parks 3 parks. You have two dense cities and 1 rural area. There are 6 park location options (blue-rimmed circles).

**Blocking coalition:** a set of *n/3* people such that there is 1 park location that they *all* prefer to their given location.

**Proportional clustering:** one with no blocking coalition.

