

Introduction to Sorting

- References:
 - Text book : Chapter 8, Skip Chapter 9: Faster Sorting Methods
 - Previous CSC313 class notes

- Java interface Comparable:

Define a method compareTo() return integer // recall String class

Example: x.compareTo(y) returns

negative ($x < y$) , 0 ($x = y$) or positive ($x > y$)

- Generic type static method

Recall generic class definition: `class A < T > { }`

If you have a static method, this is not working:

```
class A <T> {
    static public void f (T a) { .... } // not working!
}
```

Generic type static method should be:

```
class A {
    static public <T> void f (T a) { .... }
}
```

May call this method:

`A.f("string");`

Bounded Generic Types: If you want only generic types which has implemented certain interface, such as interface Comparable (which contains compareTo() method), you can use bounded type parameter, syntax is :

```
class A {
    static public <T extends Comparable<T>> void f (T a)
    {....}
}
```

This means compareTo() only work for T, but not subclasses or super classes of T. For example:

```
class Num implements Comparable<Num>
{
    ... // implement compareTo() ....
}

class Num2 extends Num
{
    // do not implement compareTo(), inherit compareTo()....
}
```

Then,

```
Num myNumObj;
...
A.f(MyNumObj); // it is OK

Num2 myNum2Obj;
..
A.f(MyNum2Obj); // it is error, no compareTo()
```

Wildcards: To ensure compareTo() work for type T (as long as it or its superclass implements compareTo), use “? Super” wild card as follows:

`<T extends Comparable<? Super T>>`

// now, both A.f(MyNum2Obj) and A.f(MyNumObj); are OK
// using compareTo() in Num class

- Java static methods that sort an array of “Comparable” objects

To ensure that generic class T (or its superclasses) must implement Comparable, use

```
public class SortArray
{
    public static <T extends Comparable<? Super T>>
        void sort(T[] a, int n)
    { ... }
    ...
}
```

Selection sort

Idea: For each loop i (from 0 to n-1),
search the smallest remaining object (between i to n)
place it in proper location, i.e. i

- Example :

Index	Initial	1 st loop	2 nd loop	3 rd loop	4 th loop
0	13	-10	-10	-10	-10
1	1	1	-5	-5	-5
2	2	2	2	1	1
3	-5	-5	1	2	2
4	-10	13	13	13	13

IterativeAlgorithm:

```

for (index = 0 ; index < n - 1 ; index++)
{
    Find indexOfNextSmalles,
    i.e. the index of the smallest value among a[index], . . . , a [n - 1]

    Interchange the values of a[index] and a[indexOfNextSmallest]
}

```

- Java program:

```

/*****
 * Class for sorting an array of Comparable objects from smallest to
 * largest.
 *****/

public class SortArray
{
    /** Task: Sorts the first n objects in an array into ascending order.
     * @param a  an array of Comparable objects
     * @param n  an integer > 0 */
    public static <T extends Comparable<? super T>>
        void selectionSort(T[] a, int n)
    {
        for (int index = 0; index < n - 1; index++)
        {
            int indexOfNextSmallest = getIndexOfSmallest(a, index, n - 1);
            swap(a, index, indexOfNextSmallest);
            // Assertion: a[0] <= a[1] <= . . . <= a[index] <= all other a[i]
        } // end for
    } // end selectionSort
}

```

```

/** Task: Finds the index of the smallest value in a portion of an
 *     array.
 * @param a    an array of Comparable objects
 * @param first an integer  $\geq 0$  and  $< a.length$  that is the index of
 *             the first array element to consider
 * @param last  an integer  $\geq first$  and  $< a.length$  that is the index
 *             of the last array element to consider
 * @return the index of the smallest value among
 *         a[first], a[first + 1], . . . , a[last] */
private static <T extends Comparable<? super T>>
    int getIndexOfSmallest(T[] a, int first, int last)
{
    T min = a[first];
    int indexOfMin = first;
    for (int index = first + 1; index <= last; index++)
    {
        if (a[index].compareTo(min) < 0)
        {
            min = a[index];
            indexOfMin = index;
            // Assertion: min is the smallest of a[first] through a[index].
        } // end if
    } // end for
    return indexOfMin;
} // end getIndexOfSmallest

```

```

/** Task: Swaps the array elements a[i] and a[j].
 * @param a an array of objects
 * @param i an integer >= 0 and < a.length
 * @param j an integer >= 0 and < a.length */
private static void swap(Object[] a, int i, int j)
{
    Object temp = a[i];
    a[i] = a[j];
    a[j] = temp;
} // end swap
} // end SortArray

```

- The efficiency of Selection sort

main for loop executes $n - 1$ times

For loop i , inner loop executes, `getIndexOdSmallest()`,
 $n - i - 1$ times

$$(n - 1) + (n - 2) + \dots + 1 = n(n - 1)/2 = \mathbf{O(n^2)}$$

- **Insertion sort**

Idea: For each loop i (from 1 to $n-1$),
 Objects from 0 to $i-1$ already sorted
 Find proper location within 0.. i to insert object i

- Example :

Index	Initial	1 st loop	2 nd loop	3 rd loop	4 th loop
0	13	1	1	-5	-10
1	1	13	2	1	-5
2	2	2	13	2	1
3	-5	-5	-5	13	2
4	-10	-10	-10	-10	13

Algorithm insertionSort (a, first, last)

// Sorts the array elements a[first] through a[last] iteratively.

```

for (unsorted = first + 1 through last)
{
    firstUnsorted = a [unsorted]
    insertInOrder (firstUnsorted, a, first, unsorted - 1)
}

```

Algorithm insertInOrder (element, a, begin, end)

// Inserts element into the sorted array elements a[begin] through a[end].

```

index = end
while ((index >= begin) and (element < a [index]))
{
    a [index + 1] = a [index] // make room
    index--
}
// Assertion: a[index + 1] is available.
a [index + 1] = element // insert

```

- The efficiency of insertion sort

main for loop executes $n - 1$ times

For loop i , inner loop executes, `insertInOrder()`,
At most i times

$$1 + 2 + \dots + (n - 2) + (n - 1) = n(n - 1)/2 = \mathbf{O(n^2)}$$

Best case running time : $O(n)$ // input a sorted list

Average case running time: $O(n^2)$

Worst Case Running time : $O(n^2)$

- **Bubble sort**

Algorithm

// compare adjacent keys from the remaining keys in $A[0, n-j-1]$

// and exchange them if they are out of order

// one record, $A[n-pass]$, is in place in every pass

```
sorted = false
```

```
for (pass=1; (pass < n) && !sorted; pass++)
```

```
{
```

```
    sorted = true
```

```
    for (k=0; k<n-pass; k++)
```

```
        if (A[k].key > A[k+1].key)
```

```
        {
```

```
            SWAP(A[k],A[k+1])
```

```
            sorted = false
```

```
        }
```

```
}
```


- Example :

Index	Initial	1 st loop	2 nd loop	3 rd loop	4 th loop
0	13	1	1	-5	-10
1	1	2	-5	-10	-5
2	2	-5	-10	1	1
3	-5	-10	2	2	2
4	-10	13	13	13	13

- Analysis

Best case running time : $O(n)$ // input a sorted list

Average Case Running time : $O(n^2)$ // sort half of the list

We skip Chapter 9 Faster Sorting Methods in this course.
Those methods will be covered in CSC340 or CSC510

Faster methods include : Merge Sort, Quick Sort and Radix Sort

	Average Case	Best Case	Worst Case
Radix sort	$O(n)$	$O(n)$	$O(n)$
Merge sort	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
Quick sort	$O(n \log n)$	$O(n \log n)$	$O(n^2)$
Shell sort	$O(n^{1.5})$	$O(n)$	$O(n^2)$ or $O(n^{1.5})$
Insertion sort	$O(n^2)$	$O(n)$	$O(n^2)$
Selection sort	$O(n^2)$	$O(n^2)$	$O(n^2)$