## SAN FRANCISCO STATE UNIVERSITY COMPUTER SCIENCE DEPARTMENT

# SORTING REVIEW (CSC220) SELECTION SORT, INSERTION SORT, SHELL SORT

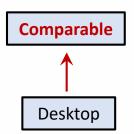
**DUC TA** 

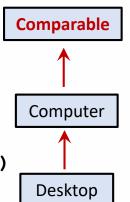
### Generic Types?

```
// Selection Sort public static T extends Comparable T super T void selection T a, int n) {
```

### Bounded Type Parameters

- Objects in the passed array be strings, Integer Objects, or any comparable objects.
- The class of these objects must implement the interface Comparable.
- T is bounded to represent class type that provides the method compareTo().
- The green (1<sup>st</sup>) T is the type parameter's name
- The blue (2<sup>nd</sup>) T is the upper bound of the green (1<sup>st</sup>) T
- The brown (3<sup>rd</sup>) T is the return type (if used or if not void)
- The black (4<sup>th</sup>) **T** is the parameter type
- The extends keyword used in general sense to mean either extends
  (as in classes) or implements (as in interfaces). "implements" in this case.
- Problems: it is more restrictive than necessary
  - compareTo only works for T, but not subclasses or super class of T.
  - Desktop smallestDesktop = MyClass.arrayMinimum(myArray)
  - Insisting comparing a Desktop vs. a Desktop, not even a Computer

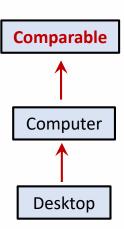




#### Bounded WildCard

```
// CORRECT public static \ensuremath{^{T}} extends Comparable\ensuremath{^{?}} super \ensuremath{^{T}>>} T arrayMinimum(\ensuremath{^{T}[]} anArray) {
```

- Problems: it is more restrictive than necessary
  - Desktop smallestDesktop = MyClass.arrayMinimum(myArray)
- <? super T> allows comparing Desktop vs. Computer. That is Desktop vs. an object of Desktop's superclass.
- <? super T> means any superclass of T
- <? super T> T is the lower bound of the wild card "?"
- ? extends T T is the upper bound of the wild card "?"



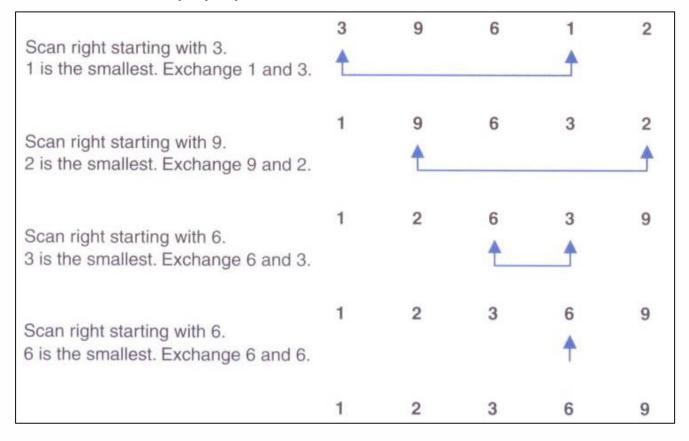
- T[] a: an array of Comparable objects
- int n: the first n objects in the array

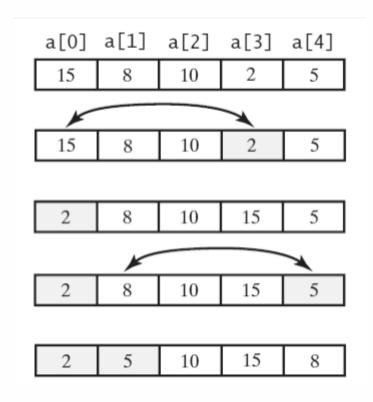


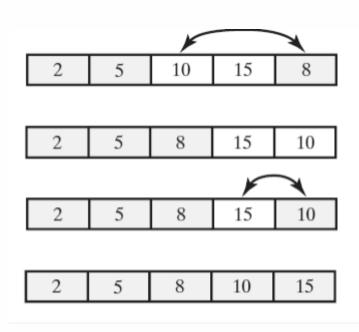
- Selection sort orders a list of values by repetitively putting a particular value into its final position
- More specifically:
  - 1. Find the smallest value in the list
  - 2. Switch it with the value in the first position
  - 3. Find the next smallest value in the list
  - 4. Switch it with the value in the second position
  - 5. Repeat until all values are in their proper places

### Smallest → Final Position

- 1. Find the smallest value in the list
- 2. Switch it with the value in the first position
- 3. Find the next smallest value in the list
- 4. Switch it with the value in the second position
- 5. Repeat until all values are in their proper places

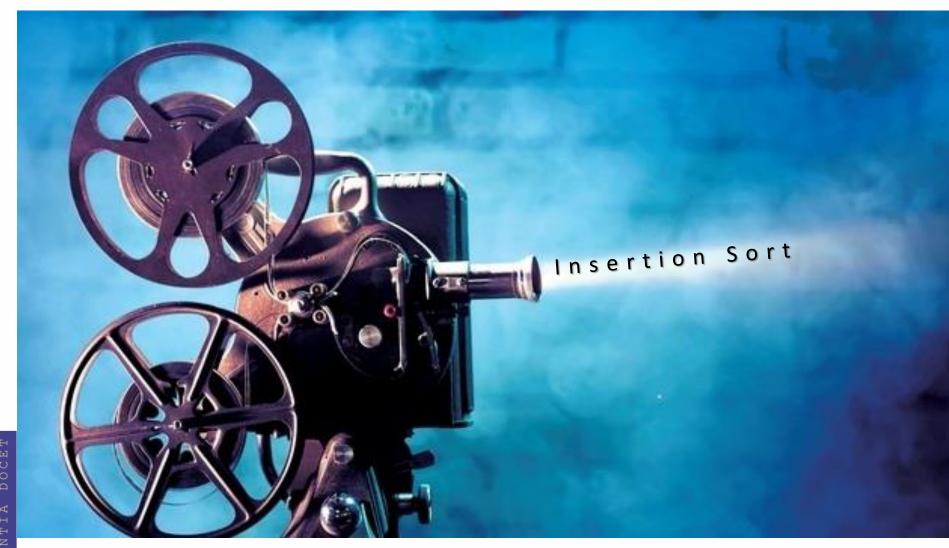






```
2
       // Finds the index of the smallest value in a portion of an array a.
       // Returns the index of the smallest value among
       private static <T extends Comparable<? super T>>
                int getIndexOfSmallest(T[] a, int first, int last) {
           T min = a[first];
           int indexOfMin = first;
10
11
           for (int index = first + 1; index <= last; index++) {</pre>
12
13
14
                if (a[index].compareTo(min) < 0) {</pre>
                    min = a[index];
15
                    indexOfMin = index;
16
17
                }
18
19
           return indexOfMin;
20
21
       }
22
23
```

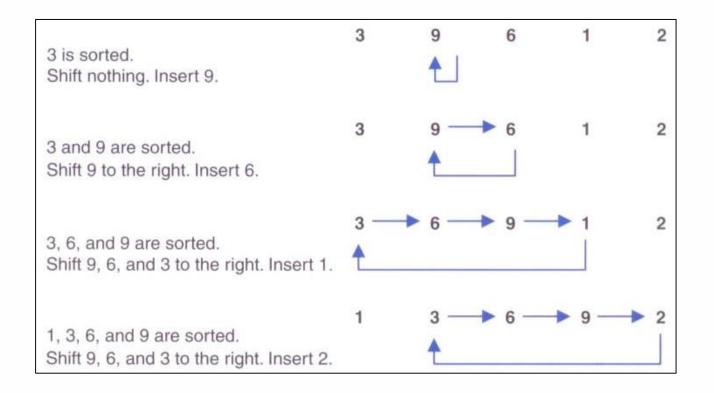
```
2
3
    public static <T extends Comparable<? super T>> void selectionSort(T[] a, int n) {
4
5
           for (int index = 0; index < n - 1; index++) {
              int indexOfNextSmallest = qetIndexOfSmallest(a, index, n - 1);
              swap(a, index, indexOfNextSmallest);
10
11
12
 2
 3
       // Swaps the array entries a[i] and a[j].
       private static void swap(Object[] a, int i, int j) {
            Object temp = a[i];
            a[i] = a[i];
            a[i] = temp;
       }
```

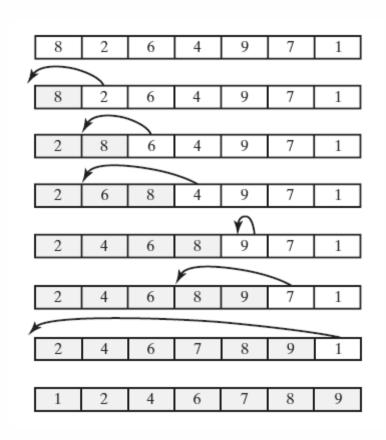


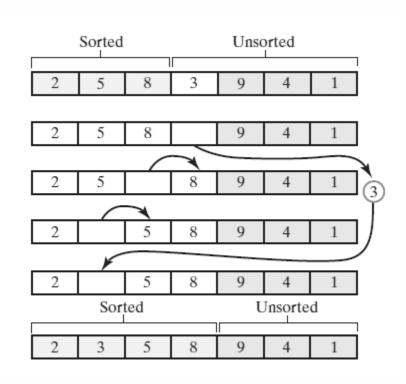
- Insertion sort orders values by repetitively inserting a particular value into a sorted subset of the list
- More specifically:
  - 1. Consider the **first** item to be a **sorted sublist** of length 1
  - 2. Insert the **second** item into the **sorted sublist**, shifting the first item if needed
  - 3. Insert the **third** item into the **sorted sublist**, shifting the other items as needed
  - 4. Repeat until all values have been inserted into their proper positions

### **Sorted Sublist**

- 1. Consider the **first** item to be a **sorted sublist** of length 1
- 2. Insert the **second** item into the **sorted sublist**, shifting the first item if needed
- 3. Insert the **third** item into the **sorted sublist**, shifting the other items as needed
- 4. Repeat until all values have been inserted into their proper positions







```
3
       // Insert anEntry, the next unsorted entry, in the sorted sublist
 4
       private static <T extends Comparable<? super T>>
 б
               void insertInOrder(T anEntry, T[] a, int begin, int end) {
 7
 8
           int index = end; // End of the sublist
 9
10
           // Keep scanning the sublist from end to beginning
11
           // Compare anEntry vs. last entry (a[index])
12
           // Shift a[index] towards the end to make room for anEntry
13
14
           while ((index >= begin) && (anEntry.compareTo(a[index]) < 0)) {</pre>
               a[index + 1] = a[index];
15
                                                      // Make room
               index--:
16
                                                      // Next left entru
17
           } // end for
18
19
           a[index + 1] = anEntry; // Insert
20
       }
21
```

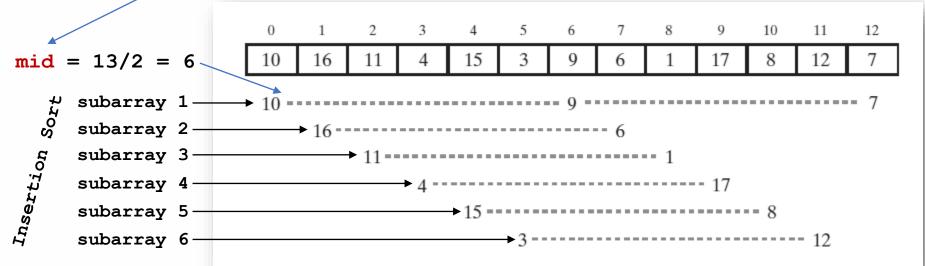
```
2
       public static <T extends Comparable<? super T>>
                void insertionSort(T[] a, int n) {
 5
           insertionSort(a, 0, n - 1);
 7
       }
 8
       public static <T extends Comparable<? super T>>
                void insertionSort(T[] a, int first, int last) {
10
11
12
           // Start at position 1, not 0
13
           11
           for (int unsorted = first + 1; unsorted <= last; unsorted++) {</pre>
14
15
16
                T firstUnsorted = a[unsorted]; // first of Unsorted
17
18
                insertInOrder(firstUnsorted, a, first, unsorted - 1);
19
20
21
```

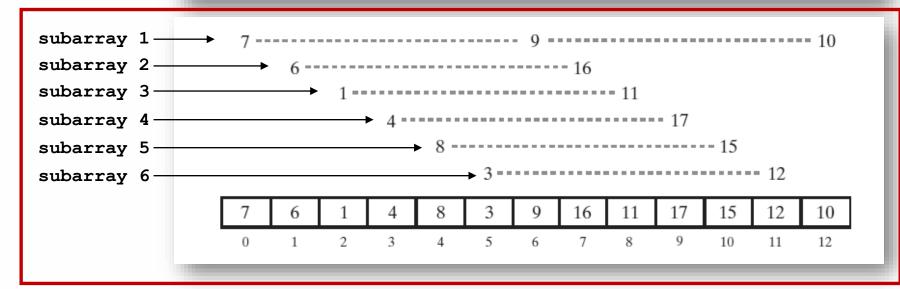


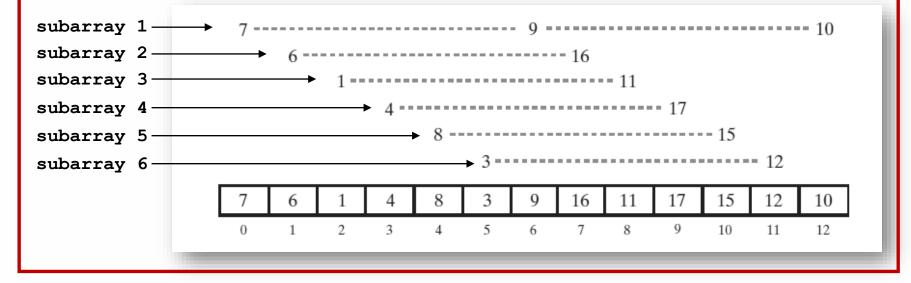
- Shell Sort is a variation of the Insertion Sort which is faster than O(n²)
- Observation during Insertion Sort:
  - When an entry is far from its correct sorted position → many moves
  - When an array is completely scrambled → Insertion Sort: **extremely inefficient**
  - When an array is almost sorted → Insertion Sort = Shell Sort: more efficient
  - Donald Shell in 1959:
    - (Insertion) sort subarrays of entries at equally spaced indices.
      - Initial separation between indices be n/2
      - Sort each subarrays separately using insertion sort
      - Halve this value each pass
      - Sort each subarrays separately using insertion sort
      - Repeat until the separation is 1
      - Final step: Insertion sort the entire array.

ort
Almost Sorted nl?

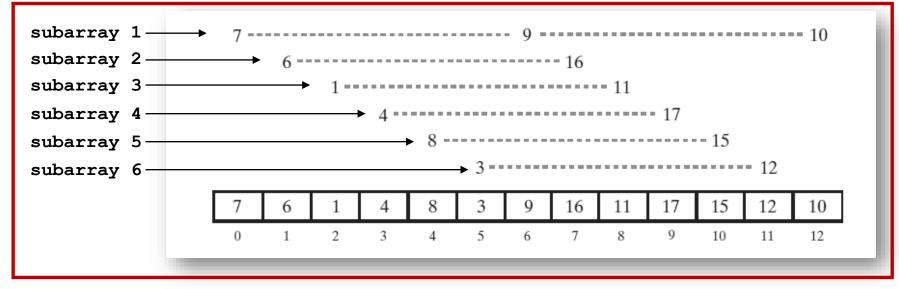
Initial separation
between indices

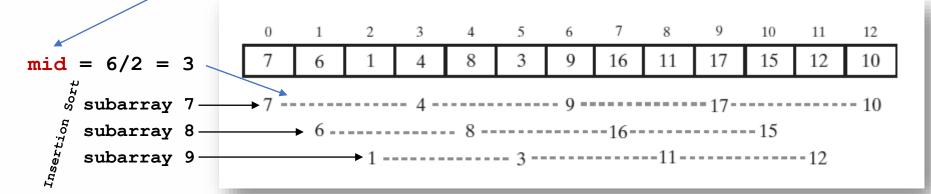


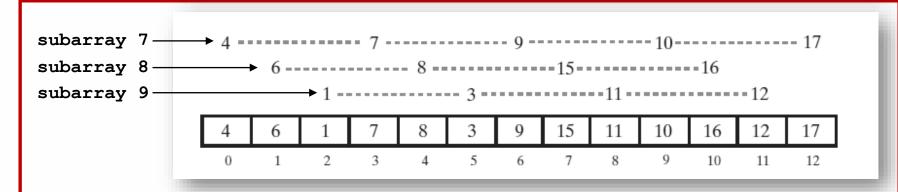




	0	1	2	3	4	5	6	7	8	9	10	11	12
mid = 6/2 = 3	7	6	1	4	8	3	9	16	11	17	15	12	10
subarray 7													
subarray 8——		<b>6</b>			- 8			16			15		
subarray 9			<b>→</b> 1			3			11			-12	



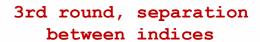




$$mid = 3/2 = 1$$

 $I_{nsertion}$  Sort

4	1	6	1	7	8	3	9	15	11	10	16	12	17
(	0	1	2	3	4	5	6	7	8	9	10	11	12



mid = 3/2 = 1  $\int_{So_{t}}^{I_{0}} S_{o_{t}}^{o_{t}}$ 

4	6	1	7	8	3	9	15	11	10	16	12	17
0	1	2	3	4	5	6	7	8	9	10	11	12

### **Sorted Sublist**

1 | 3 | 4 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 15 | 16 | 17

```
2
 3
           Sorts equally spaced elements of an array into ascending order.
           Parameters:
       //
 4
 5
                      An array of Comparable objects.
       //
              first The integer index of the first array entry to
       //
 6
                      consider; first >= 0 and < a.length.
 7
       //
 8
       //
              last
                     The integer index of the last array entry to
       //
                      consider; last >= first and < a.length.
 9
10
       //
              space The difference between the indices of the
11
       //
                      entries to sort.
12
       private static <T extends Comparable<? super T>>
13
                void incrementalInsertionSort(T[] a, int first, int last, int space) {
14
15
           int unsorted, index;
16
17
           for (unsorted = first + space; unsorted <= last;</pre>
                    unsorted = unsorted + space) {
18
19
20
               T nextToInsert = a[unsorted];
21
               index = unsorted - space;
22
23
                while ((index >= first) && (nextToInsert.compareTo(a[index]) < 0)) {</pre>
24
                    a[index + space] = a[index];
25
                    index = index - space;
26
                }
27
28
                a[index + space] = nextToInsert;
29
           }
30
       }
31
```

```
2
       public static <T extends Comparable<? super T>>
 3
 4
               void shellSort(T[] a, int n) {
 5
 6
           shellSort(a, 0, n - 1);
 7
 8
       } // end shellSort
 9
10
       /**
11
        * Sorts equally spaced elements of an array into ascending order.
12
13
        * @param a An array of Comparable objects.
        * @param first An integer >= 0 that is the index of the first array element
14
        * to consider.
15
16
        * @param last An integer ≥= first and < a.length that is the index of the
        * last array element to consider. @param space The difference between the
17
        * indices of the eleme nts to sort.
18
19
        */
20
       public static <T extends Comparable<? super T>>
21
               void shellSort(T[] a, int first, int last) {
22
23
           int n = last - first + 1; // Number of array entries
24
           int space = n / 2;
25
           while (space > 0) {
26
27
               for (int begin = first; begin < first + space; begin++) {</pre>
28
                   incrementalInsertionSort(a, begin, last, space);
29
30
               space = space / 2;
31
           }
32
       }
33
```

### More Examples

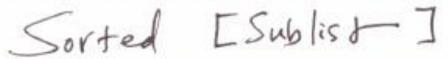
### 

Show the contents of the array above each time an Selection Sort changes it while sorting the array into ascending order.

Smaller 1 > Final Position 1. 

### 752911436

Show the contents of the array above each time an Insertion Sort changes it while sorting the array into ascending order.



171	5	2	9	1	1	4	3	6	
[7.	5.]								
5.	7	2.]							
2	5	7	91						
2.	5	7	9	1.7					
1	2.	5	7	9	1.]				
1	1	2	5.	7	9	4.]			
1	1	2	4.	5	7	9	3.]		
11	1	2	3	4	5	7.	9	6.]	
1	1	2	3	4	5	6	7	9	

8 5 3 2 2 9 3 7 4 6 A most Sorted n/2
Show the contents of the array above each time a Shell Sort changes it while sorting the array into ascending order.

										rough	
8	5	3	2	2	9	3	7	4	6	10/2=5	
8.					9.						
	3.					5.					
		3.					7.				
			2.					4.			
				2.					6.		
8	3	3	2	2	9	5	7	4	6	5/2=2	
8 2.		3.		4.		5.		8.			
	2.		3.		6.		7.		9.		
[2]	2	3	3	4	6	5	7	8	9	2/2=1	*
. 2	2]										
2	2	3]									
2	2	3	3								
2	2	3	3	4]							
2	2	3	3	4	6]						
. 2	2	3	3	4	6.	5.]					
2	2	3	3	4	5	6	7]				
- 2	2	3	3	4	5	6	子	8	1		
-2	2	3	3	4	7	6	7	D	97		
2	2	3	3	4	5	6	7	8	9		V

8 5 3 2 2 9 3 7 4 Al most Sorted M/Z
Show the contents of the array above each time a Shell Sort changes it while sorting the array into ascending order.

										nua
	8	5	3	2	2	9	3	7	4	9/2=4
	2.				4.				8.	
		5.				9.				
			3.				3.			
				2.				7.		
	2	5	3	2	4	9	3	7	8	4/2=2
	2.		3.		3.		4.		8.	
		2.		5.		7.		9.		
K	[2]	2	3	5	3	7	4	9	8	2/2 = 1
	[2	2]								
	[2	2	3 ]							
	12	2	3	5]						
	[2	2	3	5.	3.					
	t2	2	3	3	5	7]				
	[2	2	3	3	5.	7	4.]			
	T2	2	3	3	4	5	7	9]		
	[2	2	3	3	4	5	7	9.	8.]	
/	2	2	3	3	4	5	7	8	9	

9 5 6 4 3 9 3 8 2 Htmos Sorted M/2
Show the contents of the array above each time a Shell Sort changes it while sorting the array into ascending order.

									1006-1	
9	5	6	4	3	9	3	8	2	9/2=4	
2.				3.				9.	,	
	5.				9.					
		3.				6.				
			4.				8.			
2	5	3	4	3	9	6	8	9	4/2 = 2	
2.		3.		3.		6.		9.		
	4.		5.		8.		9.			
2]	4	3	5	3	8	6	9	9	2/2 = 1	*
,2	4]									
2	4.	3.]								
2	3	4	7]							
_ 2	3 3 3	4.	5	3.7						
2	3	3	4	5	8]					
- 2	3	3	4	5	8.	6.]				
_ 2	3	3	4	5	6	8	9			
2	3	3	4	5	6	8	9	9]		
2		7	4	5	6	8	9	9		レ
			•				•			
3 3 4 5 6	3 4 5 6	4 5 6	5 6	6			9			レ

### COMPARE THE SORTING ALGORITHMS

```
public static <T extends Comparable<? super T>> void selectionSort(T[] a, int n) {
        for (int index = 0; index < n - 1; index++) {
                                                          n – 1, for-loop
           int indexOfNextSmallest = getIndexOfSmallest(a, index, n - 1);
                                                                              Worst case, O(n)
           swap(a, index, indexOfNextSmallest);
                                                                              Always 3 ops, O(1)
       }
}
                           (n-1) * (O(n) + O(1)) = (n-1) * O(n) = O(n^2)
                                       simplify -> simplify ->
```

swap():

```
Always 3 steps → Number of Operations is independent of number
    of entries in an array → O(1)
    Data movement, O(n) swap, little movement.

private static void swap(Object[] a, int i, int j) {
    Object temp = a[i];
    a[i] = a[j];
    a[j] = temp;
}
```

```
private static <T extends Comparable<? super T>>
          int getIndexOfSmallest(T[] a, int first, int last) {
     T min = a[first];
     int indexOfMin = first;
     for (int index = first + 1; index <= last; index++) {</pre>
          if (a[index].compareTo(min) < 6) {</pre>
               min = a[index];
                indexOfMin = index;
     }
     return indexOfMin;
}
getIndexOfSmallest():
     last = n - 1
     first ranges from 0 to n-2
      Each time the 2 loops execute last – first. Worst (n-1) - 0 = n - 1. Best (n-1) - (n-2) = 1.
     Then total: (n-1) + (n-2) + ... + 1 = n (n-1) / 2 = \frac{1}{2} n^2 - \frac{1}{2} n
     The selection sort then is O(n^2)
 or
     The worst: O (n) for n comparisons, 1 basic operation compareTo() each time.
      Outer for-loop runs n -1 so the selection sort then is (n-1) * O(n) = O(n^2)
```

```
public static <T extends Comparable<? super T>>
        void insertionSort(T[] a, int n) {
    insertionSort(a, 0, n - 1);
public static <T extends Comparable<? super T>>
        void insertionSort(T[] a, int first, int last) {
    // Start at position 1, not 0
    //
    for (int unsorted = first + 1; unsorted <= last; unsorted++) { n − 1, for loop</pre>
        T firstUnsorted = a[unsorted]; // first of Unsorted
        insertInOrder(firstUnsorted, a, first, unsorted - 1);
                                                                   Worst case, O(n)
    }
```

```
Worst Case: (n-1) * O(n) = O(n^2)
Best Case: (n-1) * O(1) = O(n)
```

```
private static <T extends Comparable<? super T>>
         void insertInOrder(T anEntry, T[] a, int begin, int end) {
    int index = end; // End of the sublist
    // Keep scanning the sublist from end to beginning
    // Compare anEntry vs. last entry (a[index])
    // Shift a[index] towards the end to make room for anEntry
    //
    while ((index >= begin) && (anEntry.compareTo(a[index]) < 0)) {</pre>
         a[index + 1] = a[index];
                                                  // Make room
         index--;
                                                  // Next left entru
     } // end for
    a[index + 1] = anEntry; // Insert
insertInOrder():
     begin is 0
     end ranges from 0 to n-2
     1 basic operation compareTo() each time.
     Worst case O(n). Best case O(1).
     Total: 1 + 2 + ... + (n-1) = n(n-1)/2
     Thus, insertion sort is O(n^2) \leftarrow worst case
     When array already sorted, O(n) \leftarrow best case
```

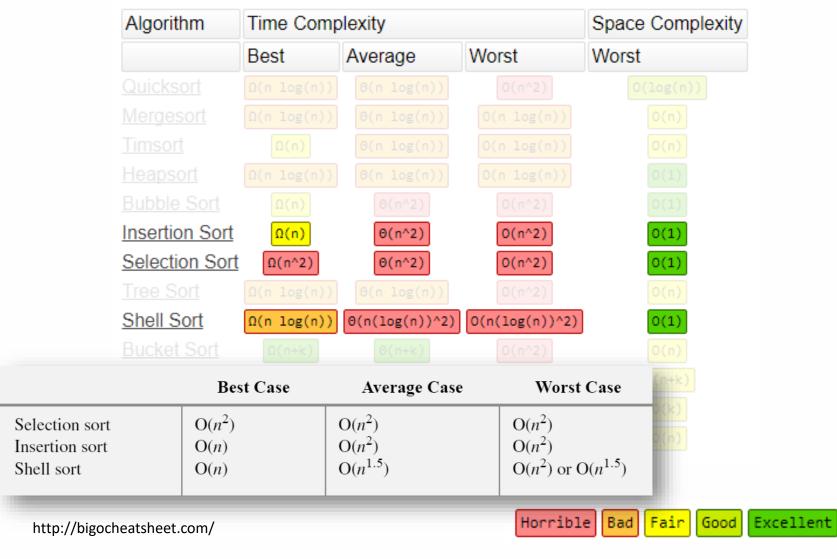
```
public static <T extends Comparable<? super T>>
         void shellSort(T[] a, int first, int last) {
    int n = last - first + 1; // Number of array entries
    int space = n / 2;
    while (space > 0) {
         for (int begin = first; begin < first + space; begin++) {</pre>
              incrementalInsertionSort(a, begin, last, space);
         space = space / 2;
private static <T extends Comparable<? super T>>
       void incrementalInsertionSort(T[] a, int first, int last, int space) {
    int unsorted, index:
                                                    O(n³), 3 nested loops
                                                    but space n/2 each time
   for (unsorted = first + space; unsorted <= last;</pre>
           unsorted = unsorted + space) {
                                                         O(n<sup>1.5</sup>) is the worst-case of Shell
                                                         O(n) is the best, an ordinary insertion sort
       T nextToInsert = a[unsorted];
       index = unsorted - space;
       ➡ile ((index >= first) && (nextToInsert.compareTo(a[index]) < 6)) {</p>
           a[index + space] = a[index];
           index = index - space;
       a[index + space] = nextToInsert;
    }
```

### **COMPARISONS**

Т	ime efficiencies i	in Big Oh notatio	n
	Best	Average	Worst
Selection Sort	O(n²)	O(n²)	O(n²)
Insertion Sort	O(n)	O(n²)	O(n²)
Shell Sort	O(n)	O(n <sup>1.5</sup> )	O(n <sup>1.5</sup> )

### COMPARING THE ALGORITHMS

### **Array Sorting Algorithms**



### **COMPARING THE ALGORITHMS**

### **Array Sorting Algorithms**

$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	Algorithm	Time Comp	lexity		Space Complexity
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		Best	Average	Worst	Worst
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	Quicksort	$\Omega(\text{n log(n)})$	Θ(n log(n))	O(n^2)	O(log(n))
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	Mergesort	$\Omega(\text{n log(n)})$	Θ(n log(n))	O(n log(n))	O(n)
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	<u>Timsort</u>	$\Omega(n)$	Θ(n log(n))	O(n log(n))	O(n)
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	<u>Heapsort</u>	$\Omega(\text{n log(n)})$	Θ(n log(n))	O(n log(n))	0(1)
	Bubble Sort	$\Omega(n)$	Θ(n^2)	O(n^2)	0(1)
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	Insertion Sort	$\Omega(n)$	Θ(n^2)	O(n^2)	0(1)
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	Selection Sort	Ω(n^2)	Θ(n^2)	O(n^2)	0(1)
Bucket Sort $\Omega(n+k)$ $\Theta(n+k)$ $O(n^2)$ $O(n)$ Radix Sort $\Omega(nk)$ $O(nk)$ $O(nk)$	Tree Sort	$\Omega(\text{n log(n)})$	Θ(n log(n))	O(n^2)	O(n)
Radix Sort Ω(nk) Θ(nk) Ο(n+k)	Shell Sort	$\Omega(\text{n log(n)})$	Θ(n(log(n))^2)	O(n(log(n))^2)	0(1)
	Bucket Sort	$\Omega(n+k)$	Θ(n+k)	O(n^2)	0(n)
Counting Sort $O(n+k)$ $O(n+k)$ $O(n+k)$	Radix Sort	$\Omega(nk)$	Θ(nk)	O(nk)	0(n+k)
Southing Soft Little State Sta	Counting Sort	$\Omega(n+k)$	0(n+k)	0(n+k)	0(k)
$\underline{\text{Cubesort}} \qquad \boxed{\text{O(n log(n))}} \qquad \boxed{\text{O(n log(n))}}$	Cubesort	Ω(n)	Θ(n log(n))	O(n log(n))	0(n)

Horrible Bad Fair Good Excellent

# See you next class!