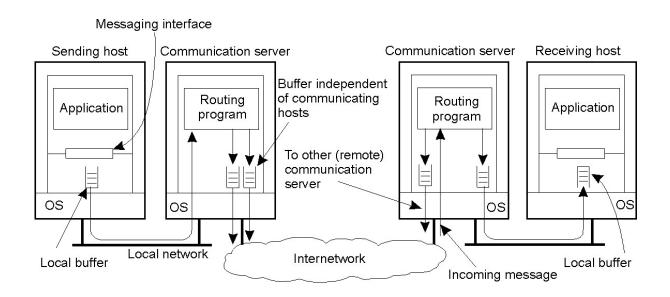
#### Last Class: RPCs and RMI

- Case Study: Sun RPC
- Lightweight RPCs
- Remote Method Invocation (RMI)
  - Design issues

#### **Communication Issues**

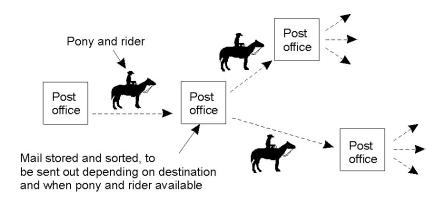
- Message-oriented communication
  - Persistence and synchronicity
- Stream-oriented communication

# Module 1: Persistence and Synchronicity in Communication



#### **Persistence**

- Persistent communication
  - Messages are stored until (next) receiver is ready
  - Examples: email, pony express



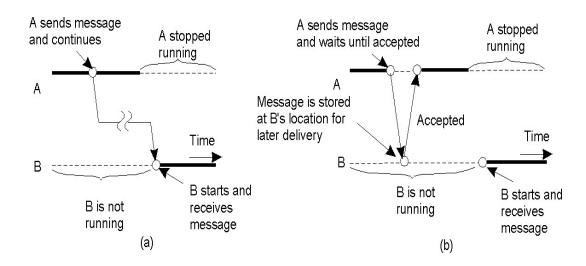
#### **Transient Communication**

- Transient communication
  - Message is stored only so long as sending/receiving application are executing
  - Discard message if it can't be delivered to next server/receiver
  - Example: transport-level communication services offer transient communication
  - Example: Typical network router discard message if it can't be delivered next router or destination

# Synchronicity

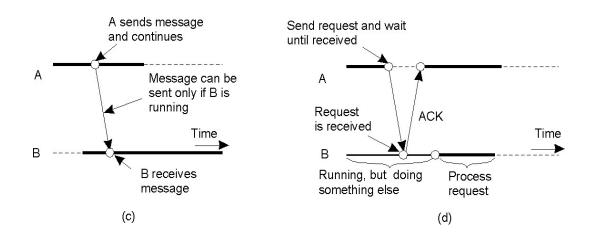
- Asynchronous communication
  - Sender continues immediately after it has submitted the message
  - Need a local buffer at the sending host
- Synchronous communication
  - Sender blocks until message is stored in a local buffer at the receiving host or actually delivered to sending
  - Variant: block until receiver processes the message
- Six combinations of persistence and synchronicity

#### Persistence and Synchronicity Combinations



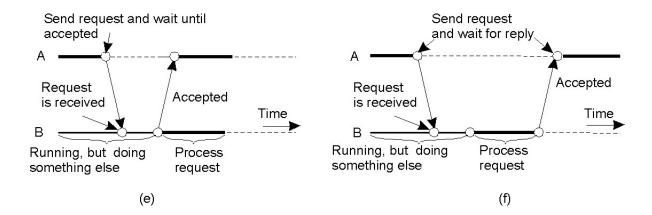
- a) Persistent asynchronous communication (e.g., email)
- b) Persistent synchronous communication

#### Persistence and Synchronicity Combinations



- c) Transient asynchronous communication (e.g., UDP)
- d) Receipt-based transient synchronous communication

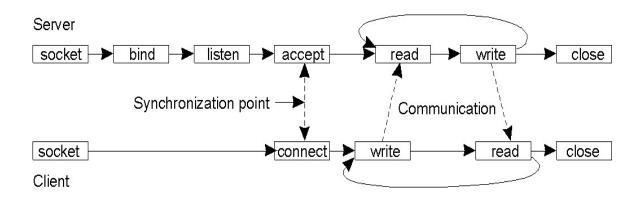
#### Persistence and Synchronicity Combinations



- e) Delivery-based transient synchronous communication at message delivery (e.g., asynchronous RPC)
- f) Response-based transient synchronous communication (RPC)

#### Message-oriented Transient Communication

- Many distributed systems built on top of simple message-oriented model
  - Example: Berkeley sockets



#### **Berkeley Socket Primitives**

Primitive	Meaning	
Socket	Create a new communication endpoint	
Bind	Attach a local address to a socket	
Listen	Announce willingness to accept connections	
Accept	Block caller until a connection request arrives	
Connect	Actively attempt to establish a connection	
Send	Send some data over the connection	
Receive	Receive some data over the connection	
Close	Release the connection	

## Message-Passing Interface (MPI)

- Sockets designed for network communication (e.g., TCP/IP)
  - Support simple send/receive primitives
- Abstraction not suitable for other protocols in clusters of workstations or massively parallel systems
  - Need an interface with more advanced primitives
- Large number of incompatible proprietary libraries and protocols
  - Need for a standard interface
- Message-passing interface (MPI)
  - Hardware independent
  - Designed for parallel applications (uses *transient communication*)
- Key idea: communication between groups of processes
  - Each endpoint is a (groupID, processID) pair

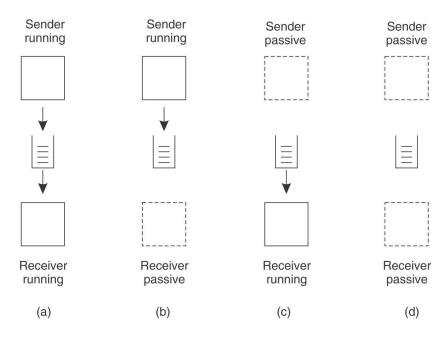
# **MPI** Primitives

Primitive	Meaning
MPI_bsend	Append outgoing message to a local send buffer
MPI_send	Send a message and wait until copied to local or remote buffer
MPI_ssend	Send a message and wait until receipt starts
MPI_sendrecv	Send a message and wait for reply
MPI_isend	Pass reference to outgoing message, and continue
MPI_issend	Pass reference to outgoing message, and wait until receipt starts
MPI_recv	Receive a message; block if there are none
MPI_irecv	Check if there is an incoming message, but do not block

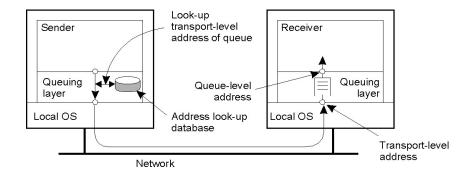
# Module 2 : Message-oriented Persistent Communication

- Message queuing systems
  - Support asynchronous persistent communication
  - Intermediate storage for message while sender/receiver are inactive
  - Example application: email
- Communicate by inserting messages in queues
- Sender is only guaranteed that message will be eventually inserted in recipient's queue
  - No guarantees on when or if the message will be read
  - "Loosely coupled communication"

# Message-Queuing Model (1)

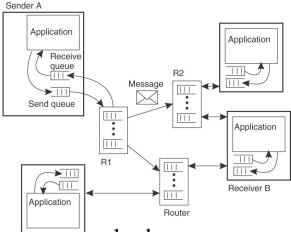


#### Message-Queuing Model



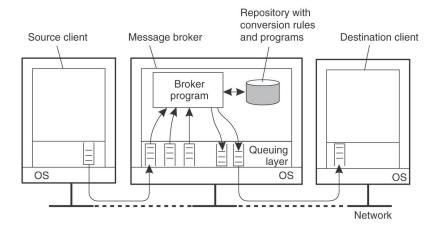
Primitive	Meaning
Put	Append a message to a specified queue
Get	Block until the specified queue is nonempty, and remove the first message
Poll	Check a specified queue for messages, and remove the first. Never block.
Notify	Install a handler to be called when a message is put into the specified queue.

# General Architecture of a Message-Queuing System (2)



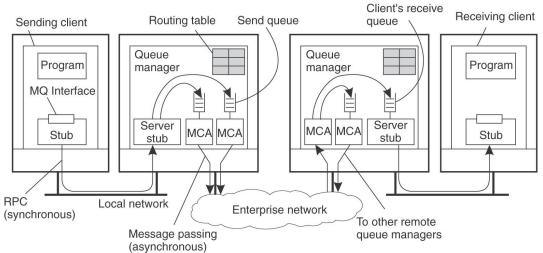
- Queue manager and relays
  - Relays use an overlay network
  - Relays know about the network topology and how to route

#### Message Brokers



- Message broker: application level gateway in MQS
  - Convert incoming messages so that they can be understood by destination (format conversion)
  - Also used for pub-sub systems

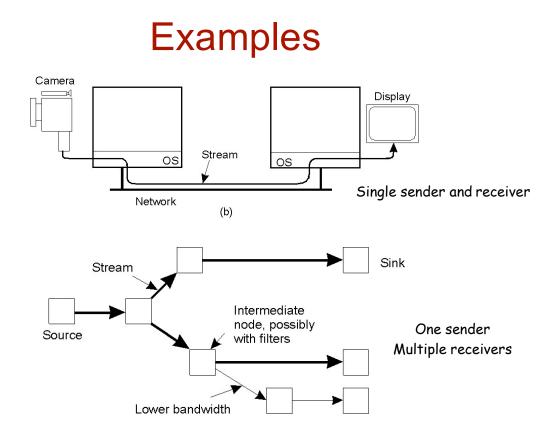
#### IBM's WebSphere MQ



- Queue managers manage queues
  - Connected through message channels
- Message channel agent (MCA)
  - Checks queue, wraps into TCP packet, send to receiving MCA

# Module 3: Stream Oriented Communication

- Message-oriented communication: request-response
  - When communication occurs and speed do not affect correctness
- Timing is crucial in certain forms of communication
  - Examples: audio and video ("continuous media")
  - 30 frames/s video => receive and display a frame every 33ms
- Characteristics
  - Isochronous communication
    - Data transfers have a maximum bound on end-end delay and jitter
  - Push mode: no explicit requests for individual data units beyond the first "play" request



#### Streams and Quality of Service

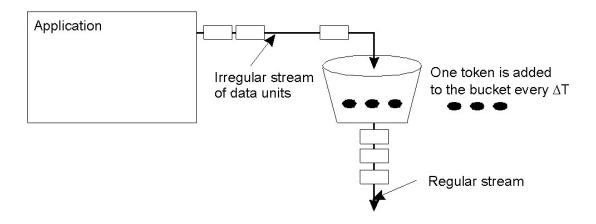
- Properties for Quality of Service:
- The required bit rate at which data should be transported.
- The maximum delay until a session has been set up
- The maximum end-to-end delay.
- The maximum delay variance, or jitter.
- The maximum round-trip delay.

## Quality of Service (QoS)

- Time-dependent and other requirements are specified as quality of service (QoS)
  - Requirements/desired guarantees from the underlying systems
  - Application specifies workload and requests a certain service quality
  - Contract between the application and the system

Characteristics of the Input	Service Required
<ul> <li>maximum data unit size (bytes)</li> <li>Token bucket rate (bytes/sec)</li> <li>Toke bucket size (bytes)</li> <li>Maximum transmission rate (bytes/sec)</li> </ul>	<ul> <li>Loss sensitivity (bytes)</li> <li>Loss interval (μsec)</li> <li>Burst loss sensitivity (data units)</li> <li>Minimum delay noticed (μsec)</li> <li>Maximum delay variation (μsec)</li> <li>Quality of guarantee</li> </ul>

## Specifying QoS: Token bucket

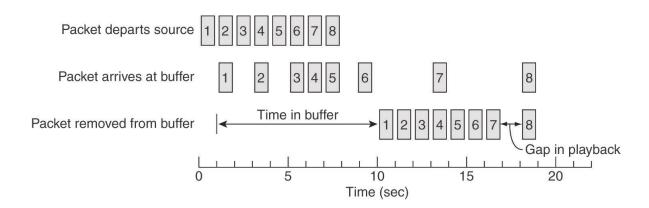


- The principle of a token bucket algorithm
  - Parameters (rate r, burst b)
  - Rate is the average rate, burst is the maximum number of packets that can arrive simultaneously

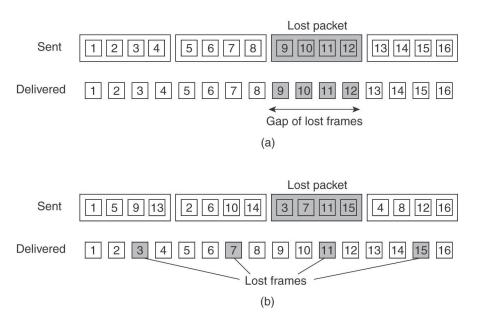
# Enforcing QoS Stream synchronization Client Stream decoder decoder QoS control

- Eniorce at ena-points (e.g., Token bucket)
  - No network support needed
- Mark packets and use router support
  - Differentiated services: expedited & assured forwarding
- Use buffers at receiver to mask jitter
- Packet losses
  - Handle using forward error correction
  - Use interleaving to reduce impact

## Enforcing QoS (1)



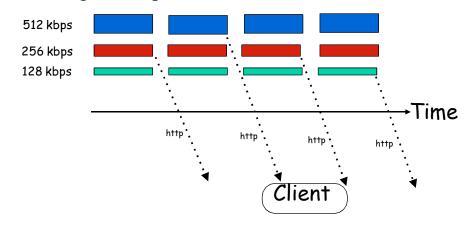
## Enforcing QoS (2)



Can also use forward error correction (FEC)

#### **HTTP Streaming**

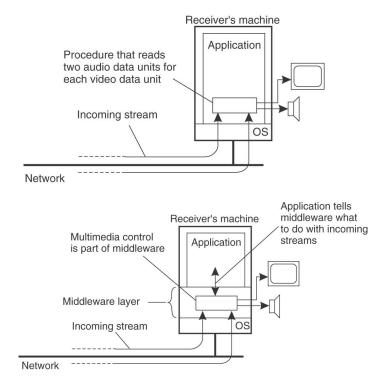
- UDP is inherently better suited for streaming
  - Adaptive streaming, specialized streaming protocols
- Yet, almost all streaming occurs over HTTP (and TCP)
  - Universal availability of HTTP, no special protocol needed
- Direct Adaptive Streaming over HTTP (DASH)
  - Intelligence is placed at the client



#### Stream synchronization

- Multiple streams:
  - Audio and video; layered video
- Need to sync prior to playback
  - Timestamp each stream and sync up data units prior to playback
- Sender or receiver?
- App does low-level sync
  - 30 fps: image every 33ms, lip-sync with audio
- Use middleware and specify playback rates

#### Synchronization Mechanism



#### Multicasting

- Group communication
  - IP multicast versus application-level multicast
  - Construct an overlay multicast tree rooted at the sender
  - Send packet down each link in the tree
- Issues: tree construction, dynamic joins and leaves

# **Overlay Construction**

