Alogorithm:-Gtep1: - Get MIN to location o. step 2: - search येव येव विवा येथ येथ येथ येथ the minimum element in the list. => Here no swap steps: - swap with Value at location MIN. 22 22 22 (22) 22 22 22. stepy: Increment MIN to point to next => Here no swap element ये येथे येथे येथे थिये थेथे थेथे. steps:- Repeat until list is stored. =) Here no Gwap २२ २२ २२ २२ २२

-Here, no bwap = 2) In the above list all the elements are same so, there are no swappings at all.

output: - 22 22 22 22 22 22 22.

2. cost-the following list using Insertion cost varun Amar karthik Ramesh Bhuvan pinesh Hroz Ganesh.

A: Insertion list: — It is also a sorting algorithm. But it is more efficient because it replaces sorting swapping with shifting.

Here every element is compared to its previous element If we found only bigger element before the key, then we shift their places.

Vorun Amar karthik Ramesh Bhuvan Dinesh firot Ganesh

Temp

co, shift varun right and insert Amar at oth position

Amar varun karthik Ramesh Bhuvan pinesh Hrot Ganch

тетр

varun > Karthik

Chift varun right and insert karthik at 1st position

Amar karthik varun Ramesh Bhuvan Dinesh Herot Ganesh

Temp

shift varun Right Ginsert Ramesh at 2nd position

				3			
hmay	Karthi	k Ramesh	varung	pinesh	19101	Ganesh	
Va Chill k	run > Bh karthik, R	awan , Rame amesh , yaru	chy Bhuvan in Lo right	mp ,karthi & inser	K) Bhuva t Bhuvar	n at i positio	
Amay	Bhuvan	Karthik R	amesh vari	un pine	sh) 18102	Ganesh	
shiff	karlhik,	Ramesh, Va	run to tigh	t and	insert pi	nesh at	
Amar	Bhuyan		rthik Rame	F	1	The black	
GhiH	Karthik,	Ramesh, var position		t and 1	insert fa	not at and	
nmar	BhuVan	pinesh fly	02 Karthik	Romest	varun	Ganesh	
shift k	rarthik, n	Pamesh, vari	into Right	and in	sert Gar	remp nesh at	
nmar e		pinesh floor		karthik	Romest	h yarun	



```
public class sortstringarrayusing Insertion sorté
public static void main (string [ Jargs ) {
string []arr = { "varun, "Amar", "karthik", "Ramesh", "Bhuvan",
                  "pinesh", "frot", "Ganesh"};
int count = 0;
stringsorted Array[] = sort - sub (arr, arr. length);
for (inti=0; ix sorted array length; i+1);
 Gystem. out. print in (corted Array[i]);
public static string[] sort-sub (string array[].int+) {
 string temp = "1"
for (int =0; ixt; i++)}
for (intj=i+1;j<f;j++);
it (array[i]. comparto Ingnore case (array[i],0)
 temp = array[i];
  array[i] zarray[i]:
  array[i] =temp;
 return array;
```

Output :-Amax Bhuvan Dinesh fr02 Ganesh Karthik. Ramesh Varun. cort the following numbers using Quick cort: 67,54,9,21, 12, 65, 56, 43, 34, 79, 70 and 45. Quick cort: -A: procedure: -I. We take a list of elements. R. We identify the first element on the Key [PNOT] element 3. comparsion starts from right to left (smaller element) 4. latter on from left to Right (bigger element) 5. on pemand, we devide-this list into & halfs. 6. We repeat steps 3,5 on the left and right sublists. Given numbers? -67), 54, 9, 21, 12, 65, 56, 43, 34, 79, 70, 45, Key 45,54,9,21,12,65,56,43,34,79,70(67) 45, (54) 9, 21, 12, 65, 56, 43, 34, 79, 70, 67

45, 34, 9, 21, 12, 65, 56, 43, 54, 67, (70, 79 45, 34, 9, 21, 12, 65, 56, 43, 54, 67, 70, 79 43,34,9,21,12,66,56,45) 43, 84, 9, 21, 12, 65, 56, 45, 54 43,12,9,21,34,145/,66,65,54,67,60, 43,17,9,21,34,145/, (54) key, (5), (6), (67), (79) (43),12,9,21,34,1451,1541,65 Key,56,1671,70,79. (ai), 12,9,1341,43,45,54,56,66,67, 40,79. 9,12,21,134,43,45,54,56,65,67,70,79 Given, list is corted using Quick cort public class Quickcorts public static void main (string [Jargs) { inti;

```
int[] arr = 867, 54, 9, 21, 12, 65, 56, 43, 34, 79, 70, 45};
   quicksort(arr, 0,9);
   System. Out. println ("In The sorted array is: In");
  for (1=0; 1<10; 1++)
   cystem out println (arr[i]).
 public static int partition (intal7, int beg, int end) &
   int left, right, temp, loc, flag;
    loc=left=beg;
     right zend?
    flag =0;
    While (flag!=1)
       While [(a[10c] <= a[right] (10gc, = right))
       right -- ?
       it (loge = = right)
       -flag=1;
       else it (a[loc]>a[right])
       temp = a sloc]:
        aloc]-afright];
        afright]=temp?
         loc = right;
       if (flag! =1)
        While ((a[loc])=a[left](loc!=left))
```

```
flag=13
  else it (a [loc] (a [left])
    temp=a[loc];
     a [loc] za [left];
     asiett]=temp;
       loe = left;
 return loc;
static void quicksort (intaII, int beg, intend)
  int loe;
   it (begrend)
    loc=partition (a, beg, end);
quicksort (a, beg, loc-1);
     quicksort (a, loc+il) ?
output: -
   12
  34
```

```
43
     54
     56
     65
      67
     79
4. Implement linear search and Binary search using Recursion.
     linear cearch: -
A:
             public class Test &
                    Gtate int arr[] = {12, 34, 54, 2,3};
                   /* Recursive method to search x in arr[1.7]*/
                   Gtatic int research (intarr [], int I, intr, intx)
                  it (x(1)
                    return-1; it (arr[2] == x)
                    return ?;
                   it (arr[r]==x)
                    return;
                   return research (arr, 1+1, r-1, x);
                11 priver method
                public static void main[string[Jargs] ?
                int x = 3;
               1/ Method call to Hnd x
              int index = research (arr, o, arr. length-1,x);
              îf (index! =-1)
```

```
(10)
```

```
ogstem.out.println ("clement"+x+"is present at index"+
   ese
      cystem.out.println ("flement + x +" is not present");
output: -
Element 3 is present at index 4.
Binary Gearch: -
public class Binary Gearch &
   1/ Returns index of xit it is present in arr 12, 7, else
   11 return-1
   int binary search (int arr [], int 2, intr, intx)
   { (+ (x>=I) {
      int mid = 1+ (r-1)/2;
     1/2+ the element is present at the middle itself
      it (arr[mid] == x)
       return mid
  1124 element is smaller-than mid, then it can only
  11 be present in left subarray
    it(arr[mid]>x)
        return binary search (arr, 1, mid-1, x)
  l'else the element can only be present in right
   1/subarray
   return binary search (arr, mid +1, r, x);
```

```
1/ We reach here when element is not present in array
      return-1;
    11 oriver method to test above
    public static void main (string args []) {
       Binary search ob = new Binary search();
       intari[]={2,3,4,10,40};
       int n = arrolength;
        int x =10;
        int result = ob. Sinarysearch (arr, o, n-1, x);
        if (result ==-1)
         system.out.println["Flement not present"];
        else
         Cystem. out println (" Element tound at index " + result);
    output:
     Element found at index 3
   Explain in brief, the various factors that determine the selec-
   tion of an algorithm to solve a computational problem.
A:- The performance of an algorithm can be measured by two
     properties.
     1. Time complexity.
    2. space complexity.
    1. time complexity: — of an algorithm quantifies the amount
```

Of time taken by an algorithm to run as a function of the length of the input.

2. Cimilarly, space complexity: - of an algorithm quantifies the amount of space (or) memory taken by an algorithm to run as a

function of the length of the input-

space complexity: - An algorithm represents the amount of memory space needed the algorithm in its life cycle.

space needed by an algorithm is equal to sum of the follow-

ing two components:

Spaces(p) = fixed part (+) + Variable part sp(2).

A tixed part that is a space required to store certain data and variables (i.e. simple variables and constants, program size etc.), that are not dependent of the size of the problem. A variable part is a space required by variables, whose size is totally dependent on the size of the problem.

For example, recursion stack space, dynamic memory alloca-

tion etc....

In other words: space can be calculated based on amount of space required.

1) to store program instructions.

2) to store constant values.

3) to store variable values.

4) and for tew other things like function calls, jumping statements etc.

```
Ex =1:-
intoum (intacz, inta)
 Int sum=0,13
 for (=0; kn; i++)
   cum = cum +A[i]
  return sum;
In the above piece of code it requires
> n*4' bytes of memory to store array variable 'al]
> 4 bytes of memory for integer parameter n.
=> 8 bytes of memory for Local integer variables 'sum' and'
   (4 bytes each)
=> 4 bytes of memory for return value.
In case of Time complexity, the running time of an algorithm
depends upon the following.
-> Whether it is a 32 bit machine (or) 64 bit machine.
- Read and write speed of the machine.
 =) the amount of time required by an algorithm to perform
Arithmetic operators, logical operators, return value and assi-
 gnment operations etc.
 => Input data.
```