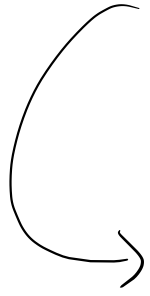


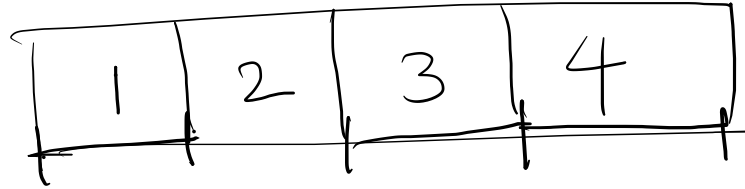
LINKED LIST



Linear data structure

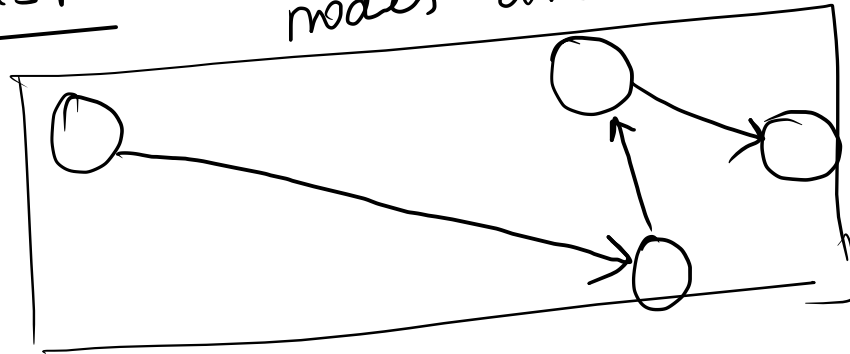
elements → sequentially

ARRAY



contiguous memory allocation

linked list

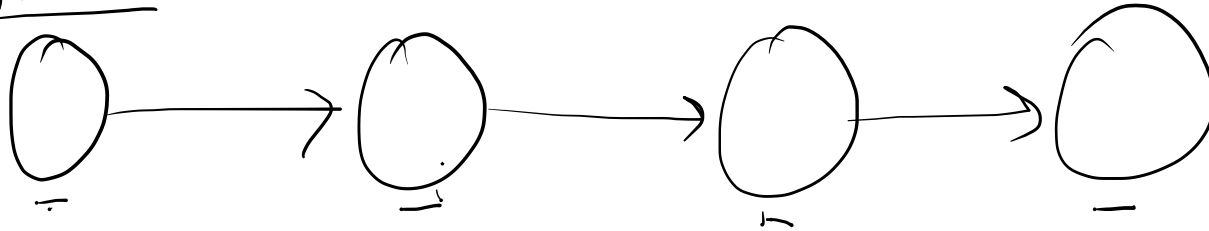


nodes are connected

non-contiguous memory allocation

Head

forward

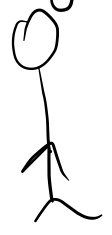
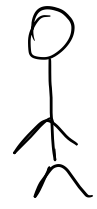


Singly linked list

↳ Unidirectional

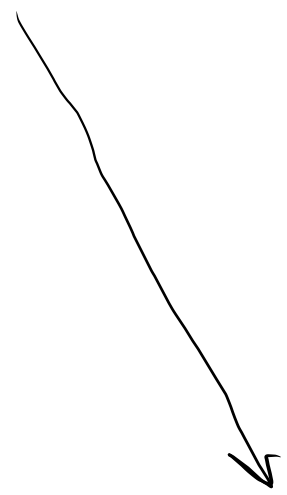
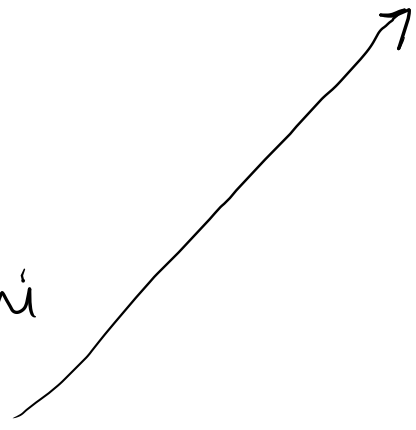
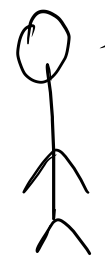
Mumbai

Bangalore

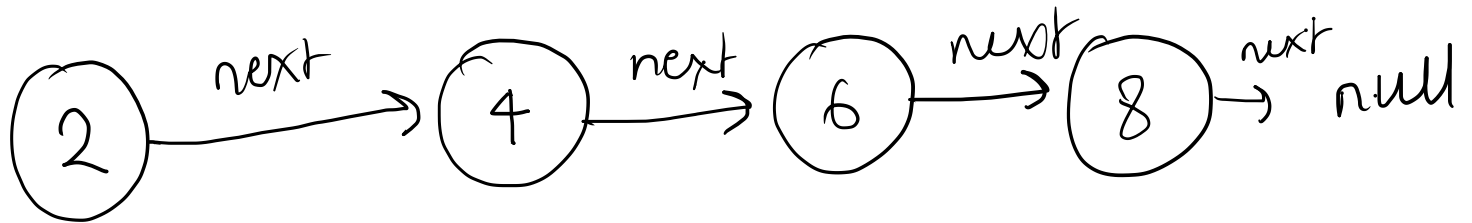


Delhi

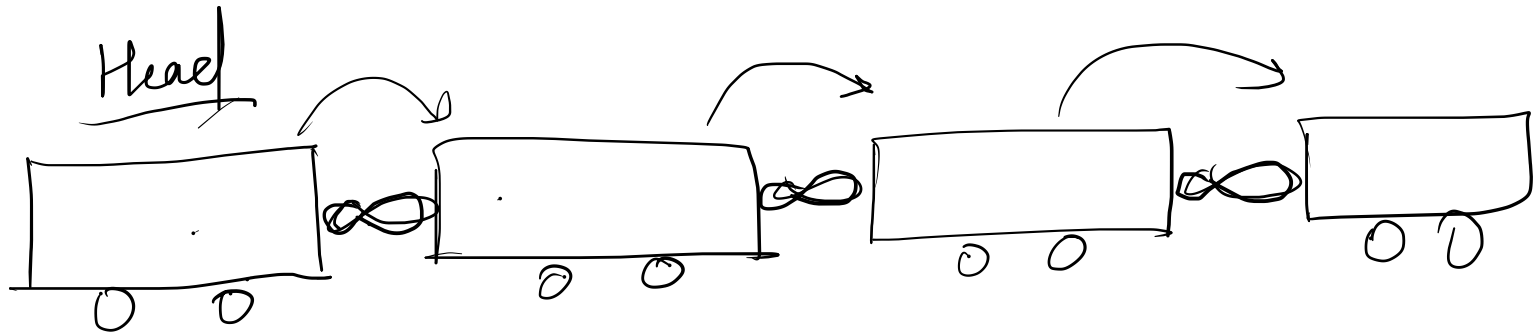
Pune



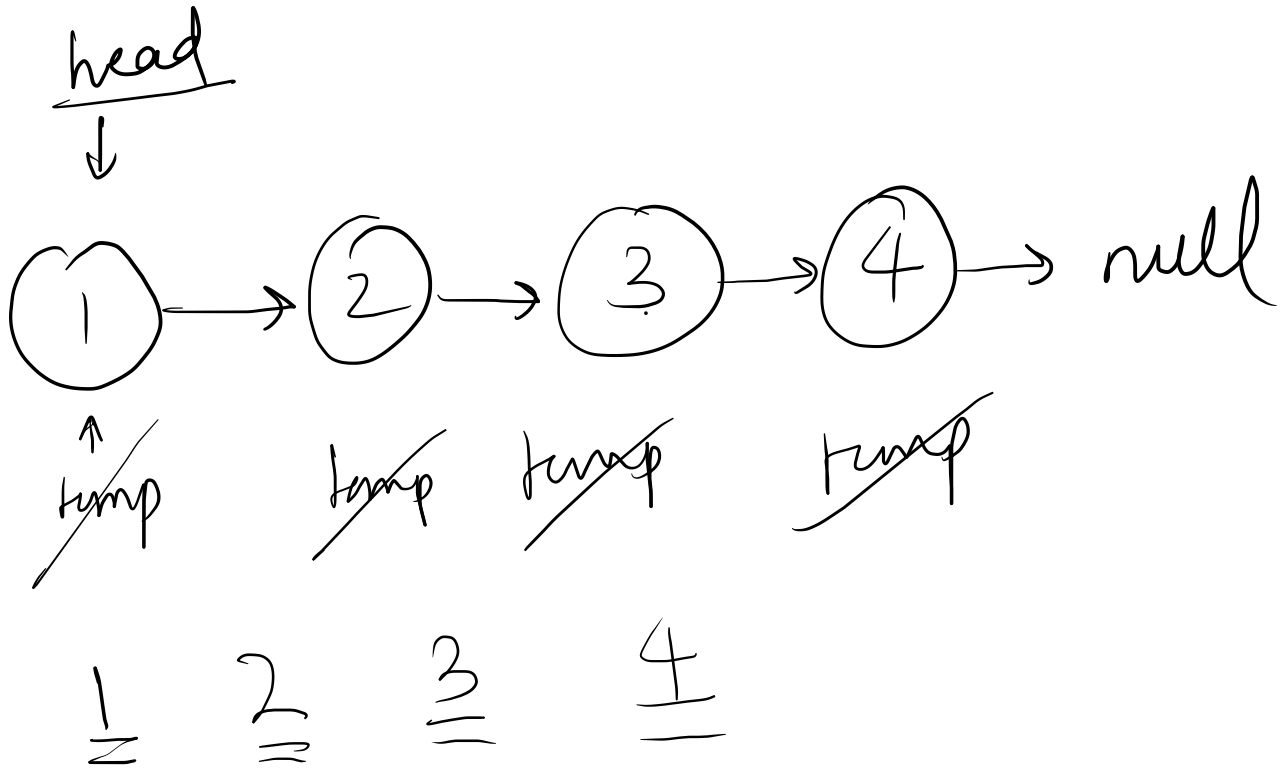
Node $\begin{cases} \rightarrow \text{data} \\ \rightarrow \text{address of next} \end{cases}$



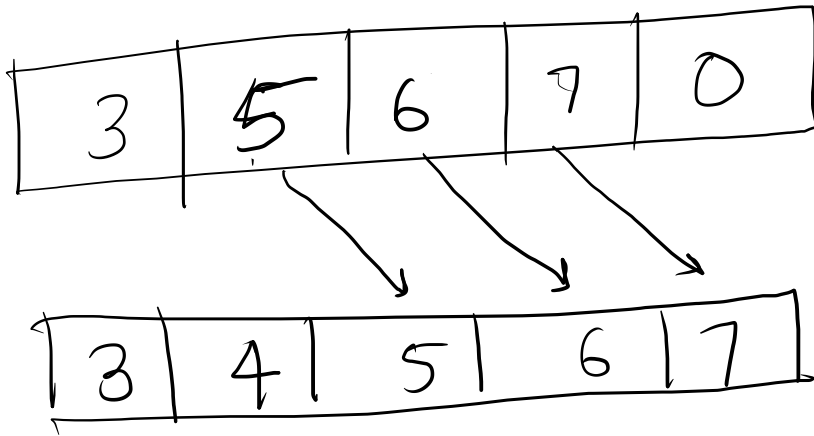
lastNode $\xrightarrow{\text{next}}$ null

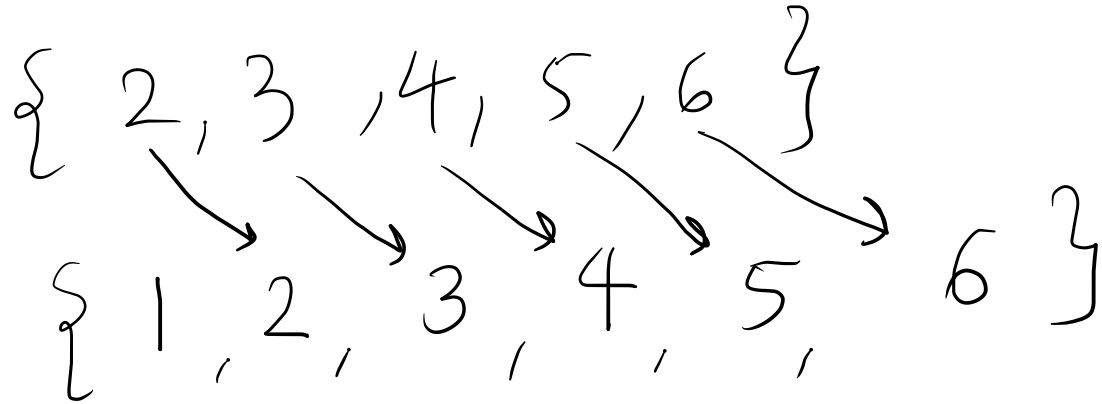


Nodes are connected

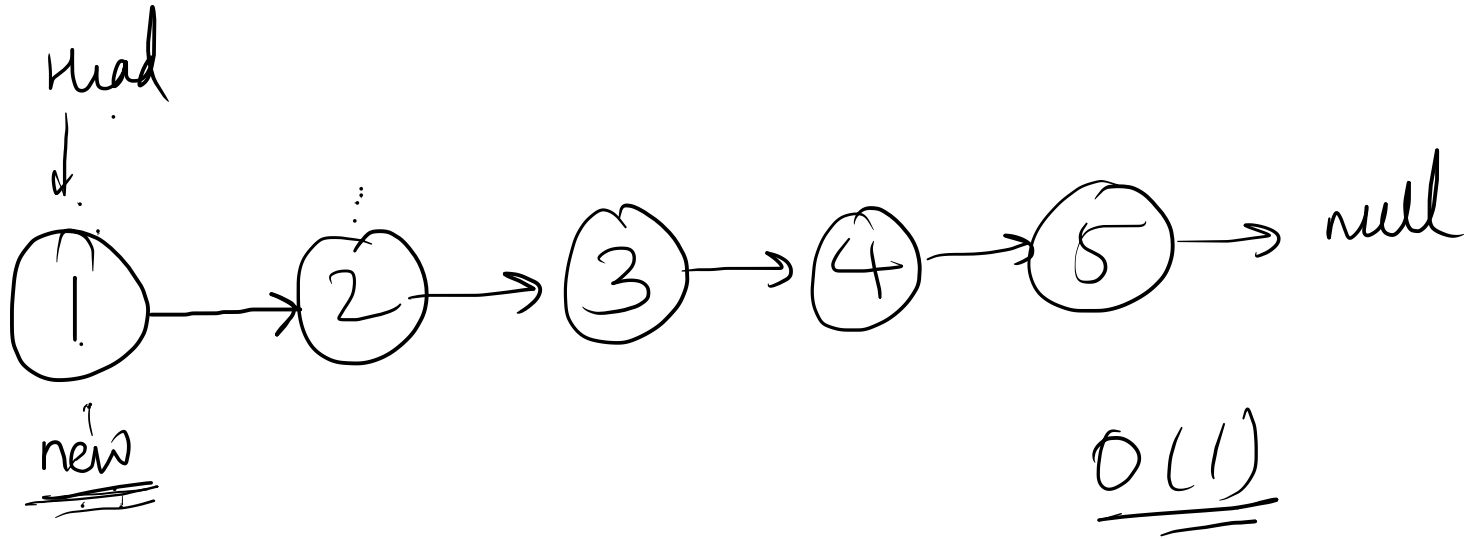


Advantage of linked list





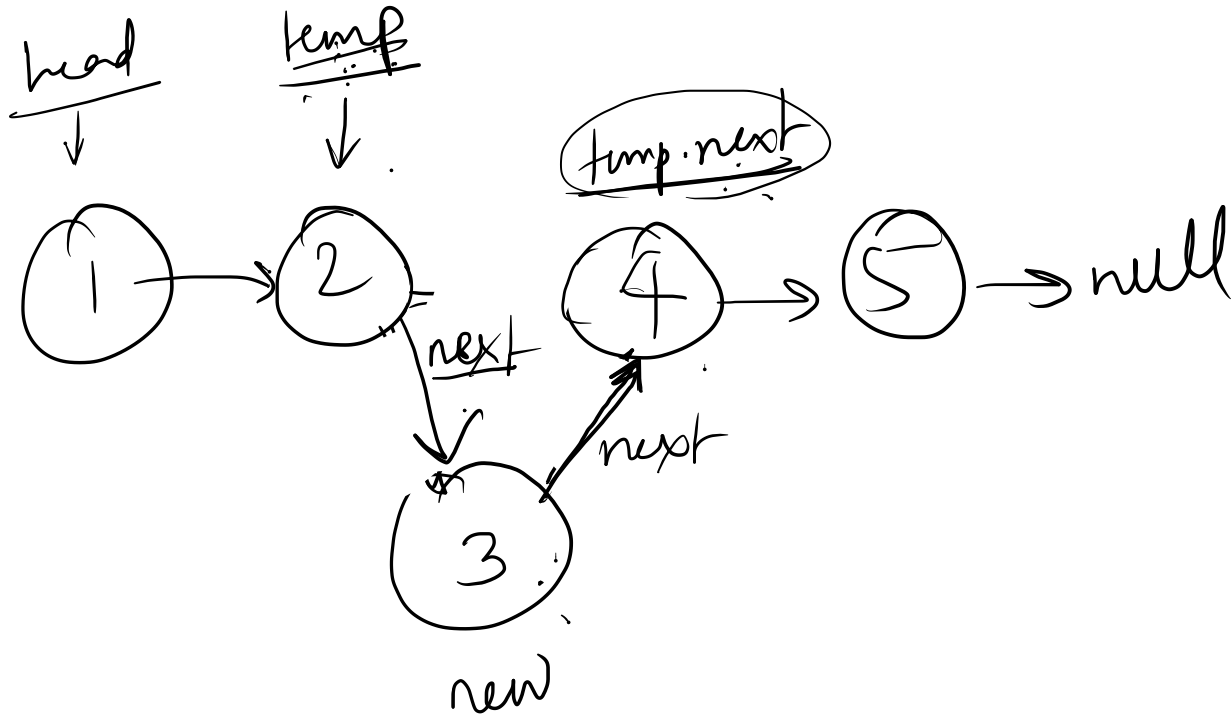
Insertion in Array
 $O(n)$



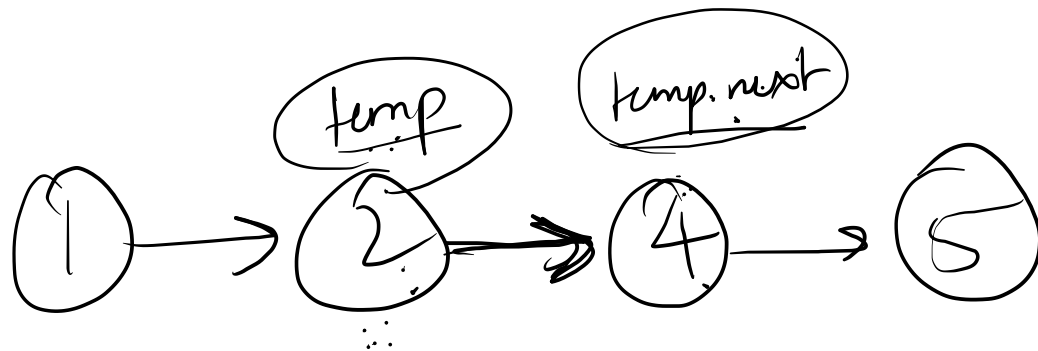
$new.next = head$
 $head = new$

node.next ✓

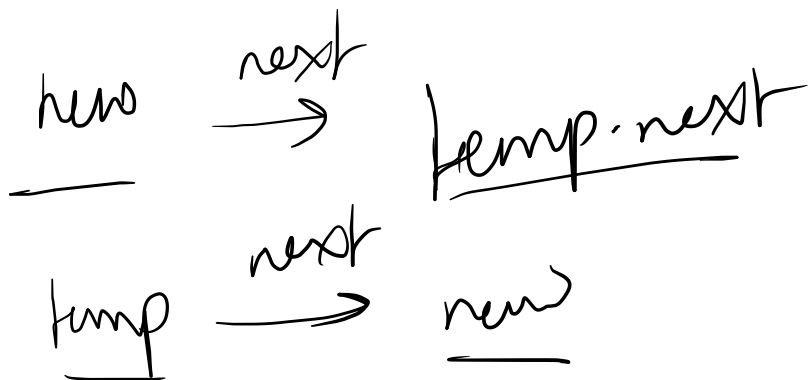
node.prev ✗

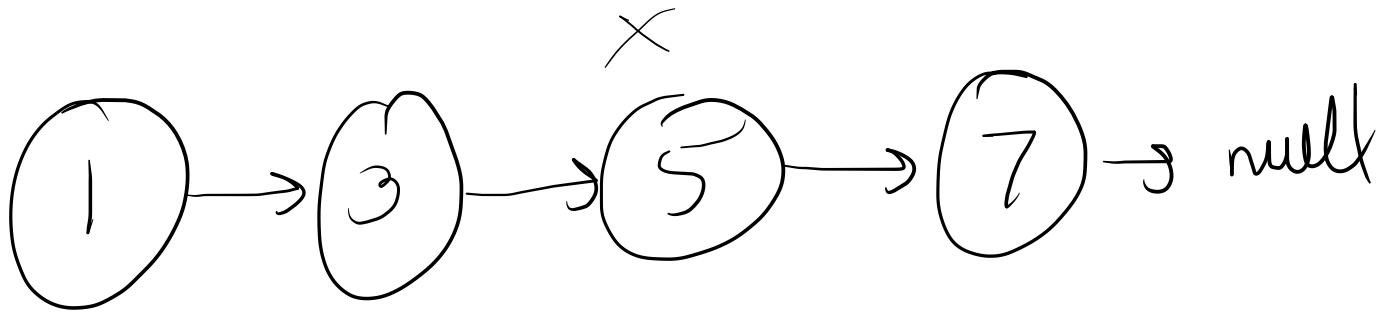


$new.next = temp.next$
 $temp.next = new$

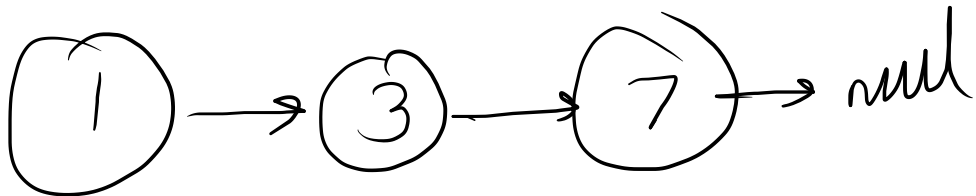


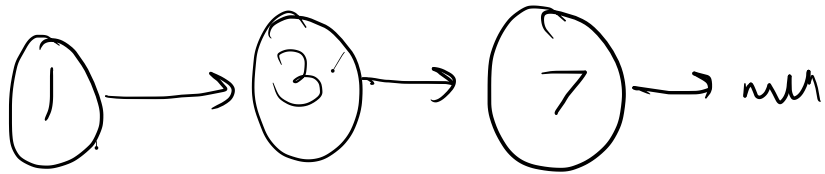
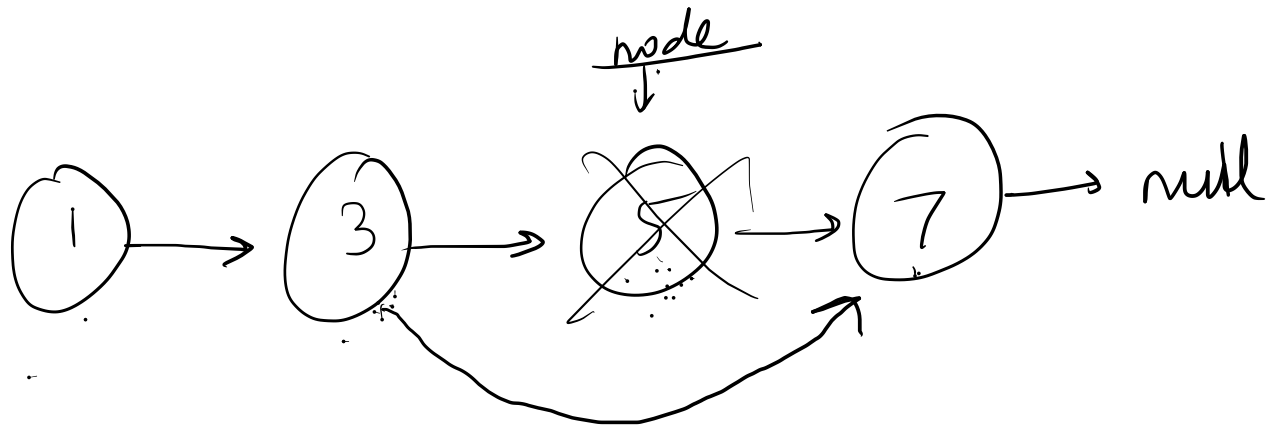
O(1)

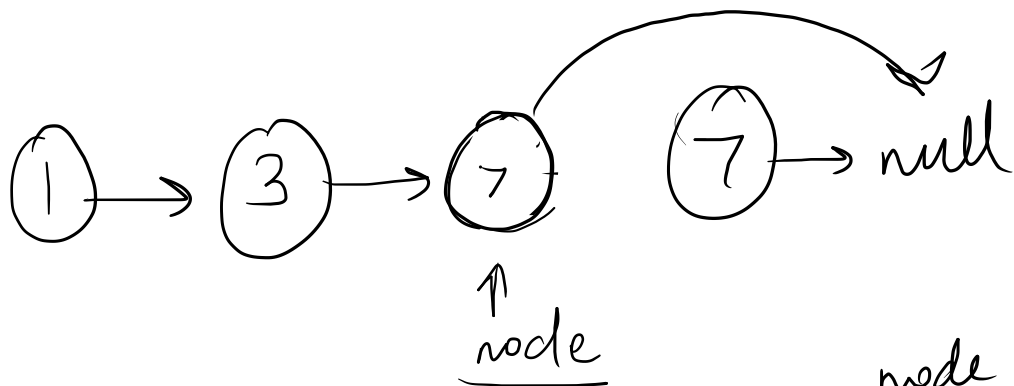




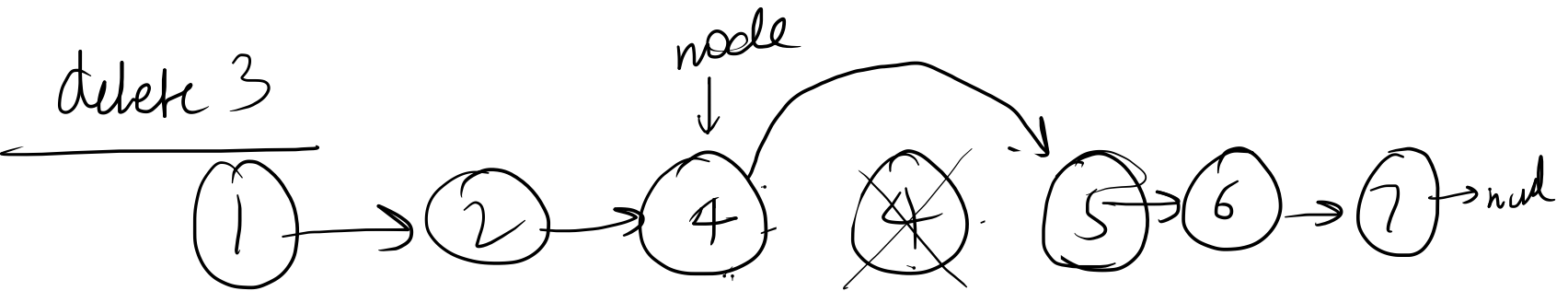
delete node 5





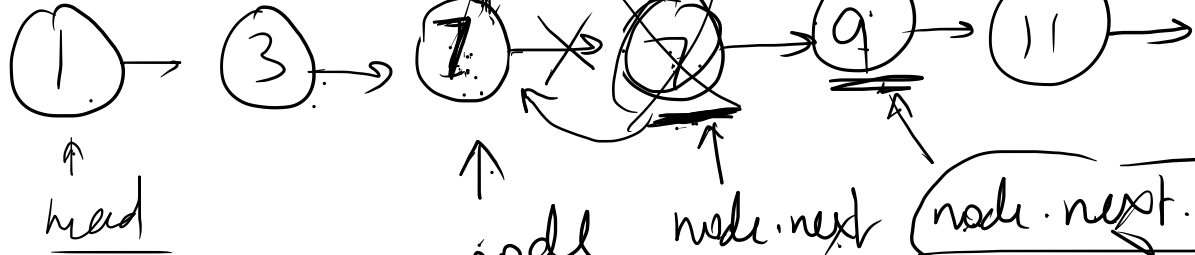


node.value == node.next

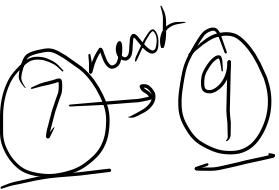


① Replace (3) with 4

delete 5



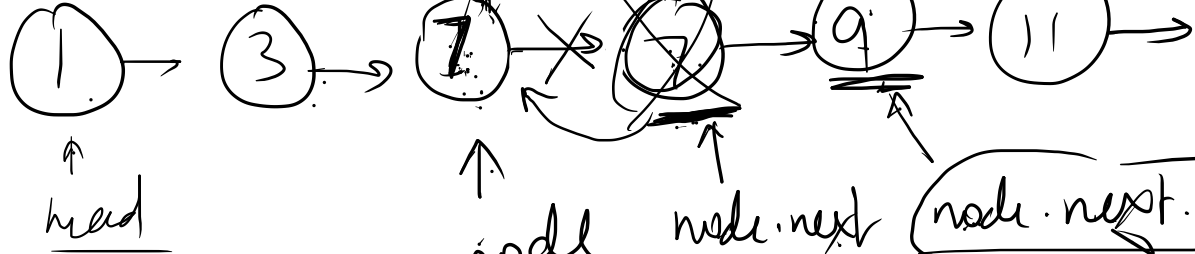
1 3 7 9 11



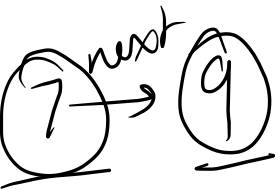
node.val = node.next.val

node.next = node.next.next

delete 5



1 3 7 9 11



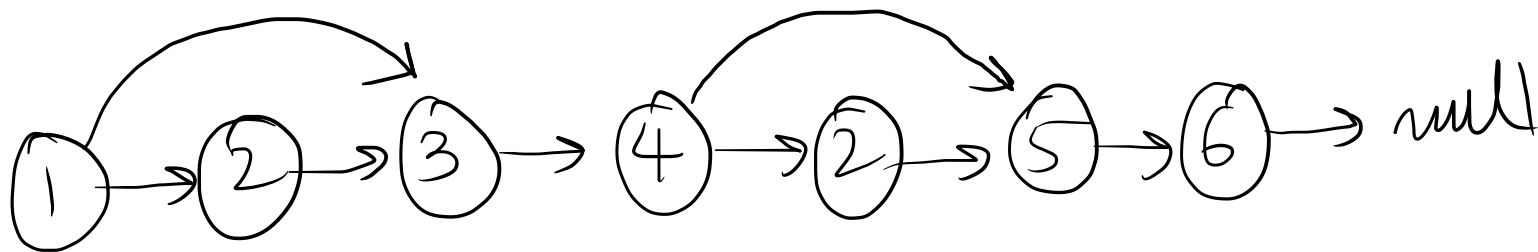
node.val = node.next.val

node.next = node.next.next

$$\underline{\underline{T.C}} = \frac{\text{constant time}}{\underline{\underline{O(1)}}}$$

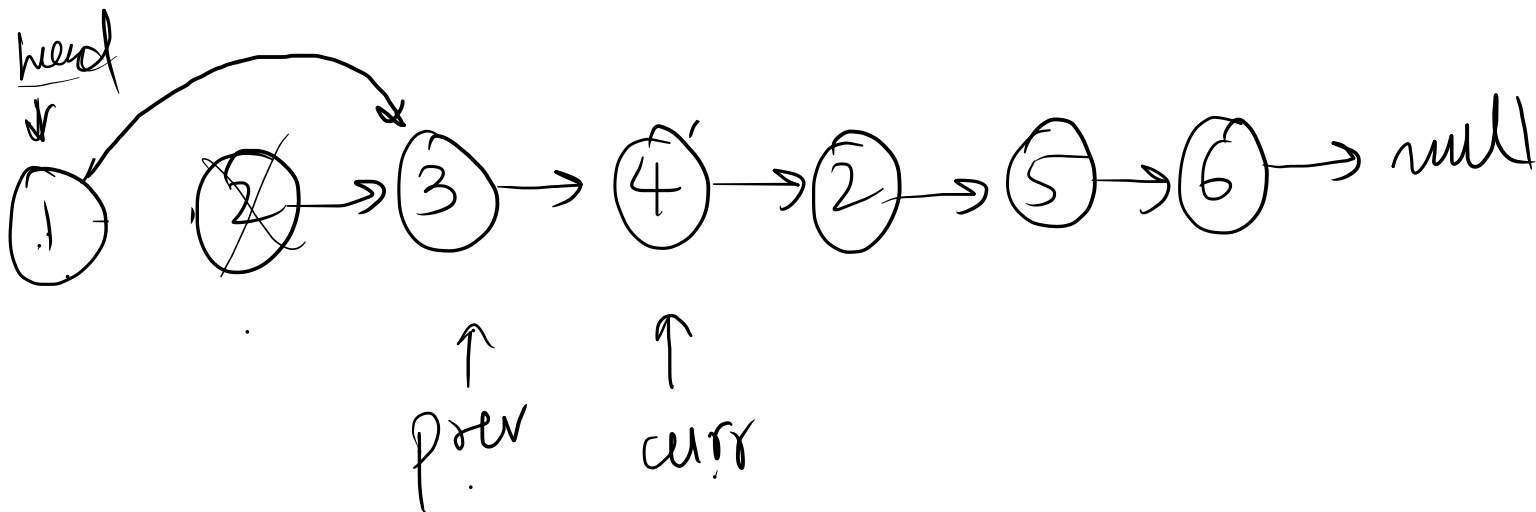
Auxiliary space: $O(1)$

2 mins Break



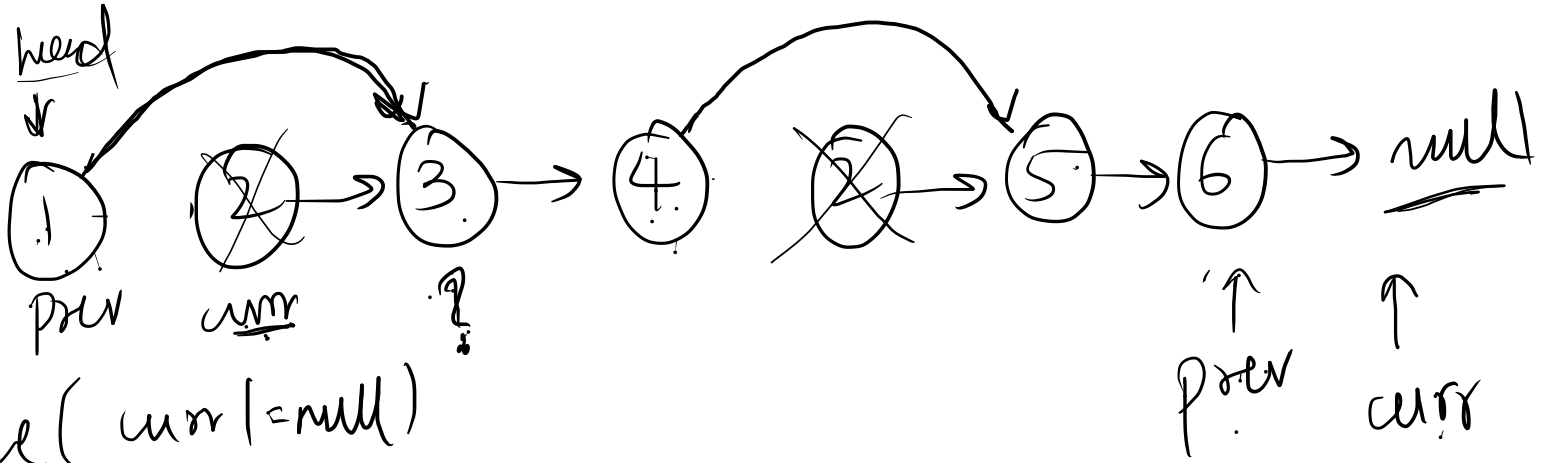
Remove 2





val = 2

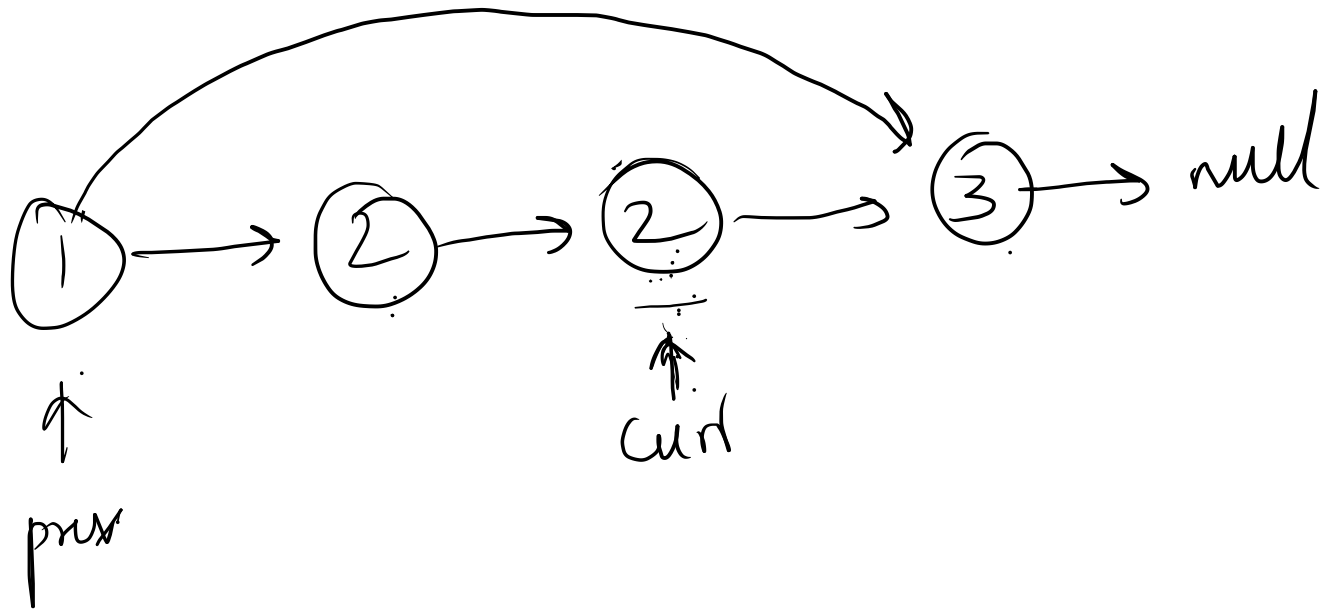
prev $\xrightarrow{\text{next}}$



```
while (curr != null)
{
```

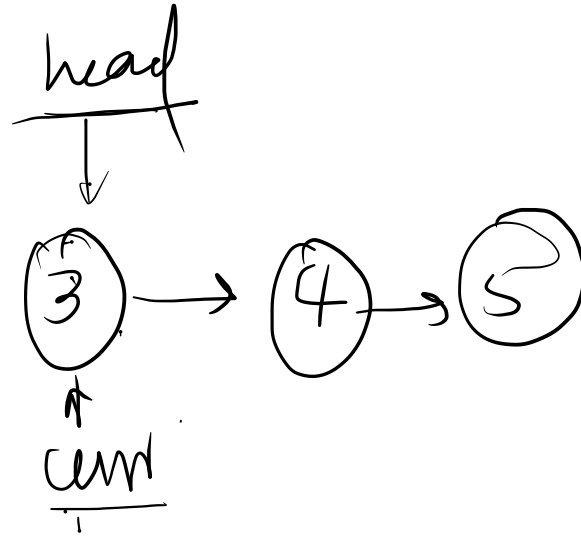
val = 2

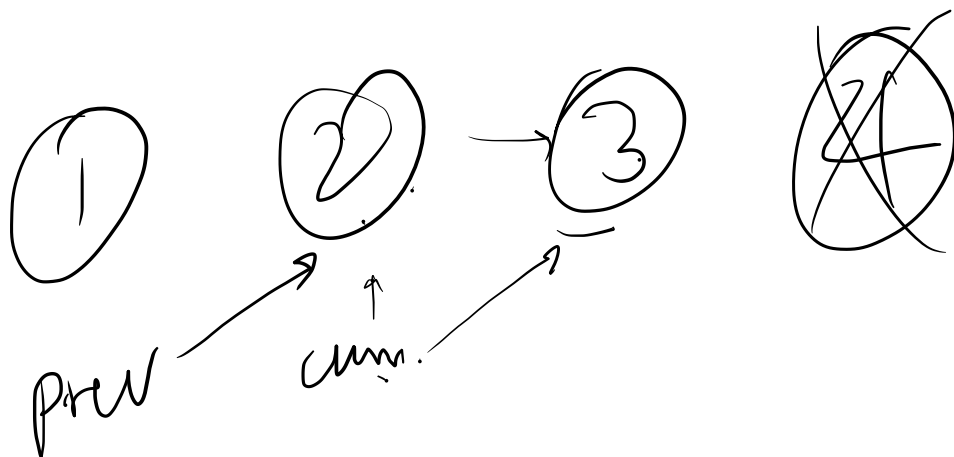
prev → next → curr → next

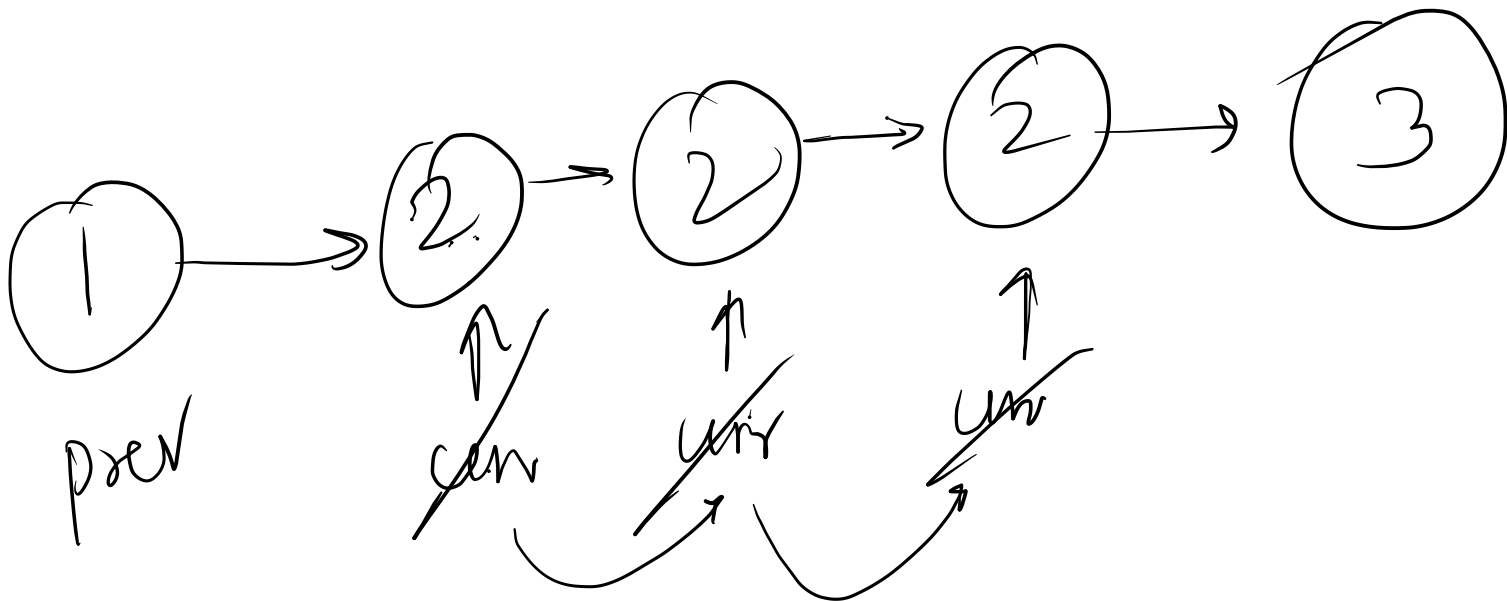


val = 2

mid
prev



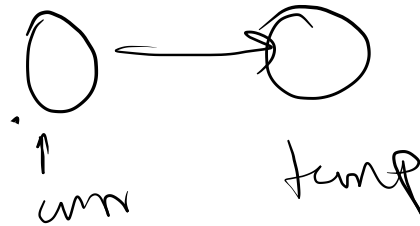
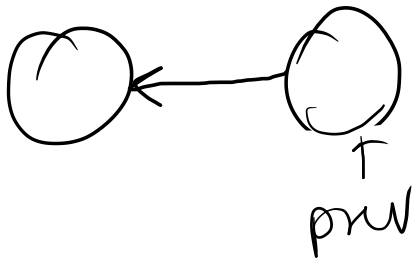
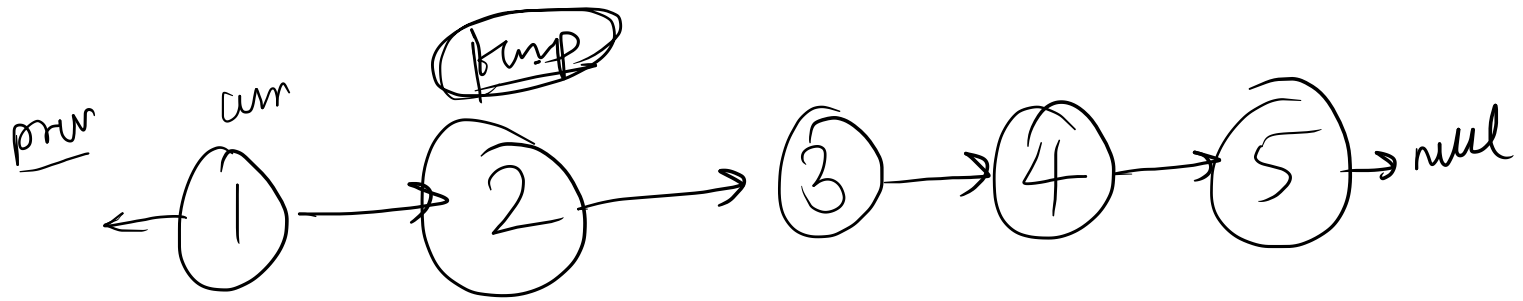




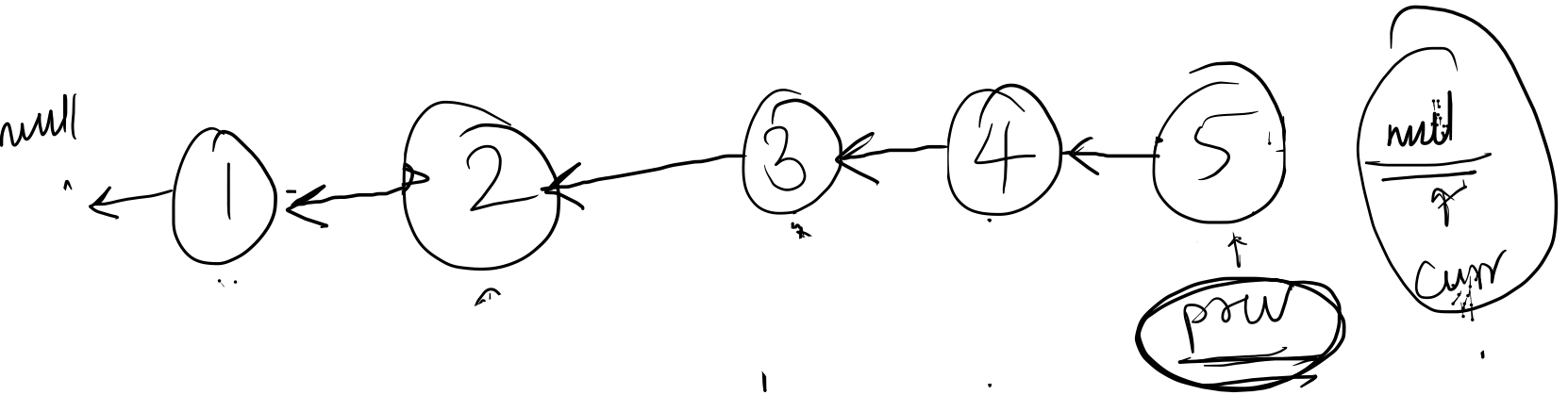
2

$$\underline{TC = O(n)}$$

$$\underline{\text{Auxiliary space} = O(1)}$$



temp = curr.next



curr.next = prev

while(curr != null)
{
}

$$\underline{TC = O(n)}$$

$$\underline{\text{Aux. space} = O(1)}$$

