

2D array

Matrix in Maths

	0	1	2	3
0	00	01	02	03
1	10	11	12	13
2	20	21	22	23
3	30	31	32	33

4x4

[i][j] i → row
j → column

3x4

```
int arr = new int[3][4];
```

$4 \times 3 \times 4 = 48$ bytes

No. of elements * size of 1 element = Memory allocated

arr[2][3]

Square Matrix

→ rows = columns

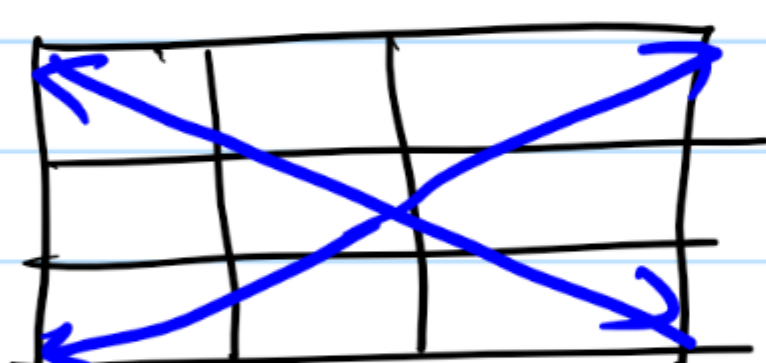
4x4

3x3 ... nxn

Rectangular Matrix

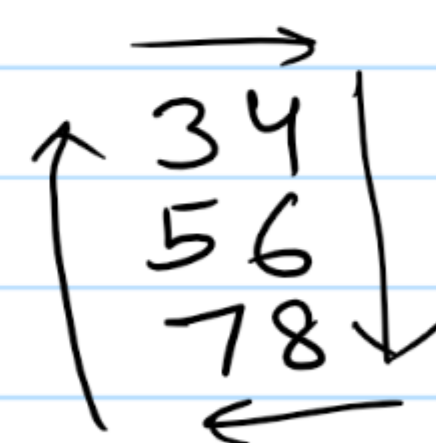
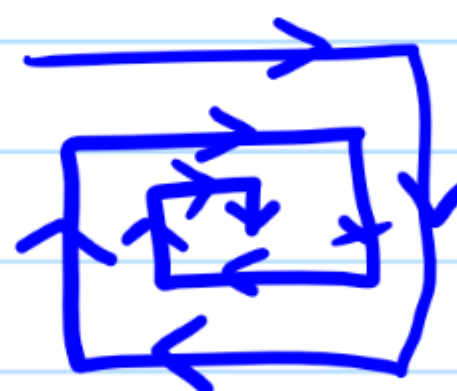
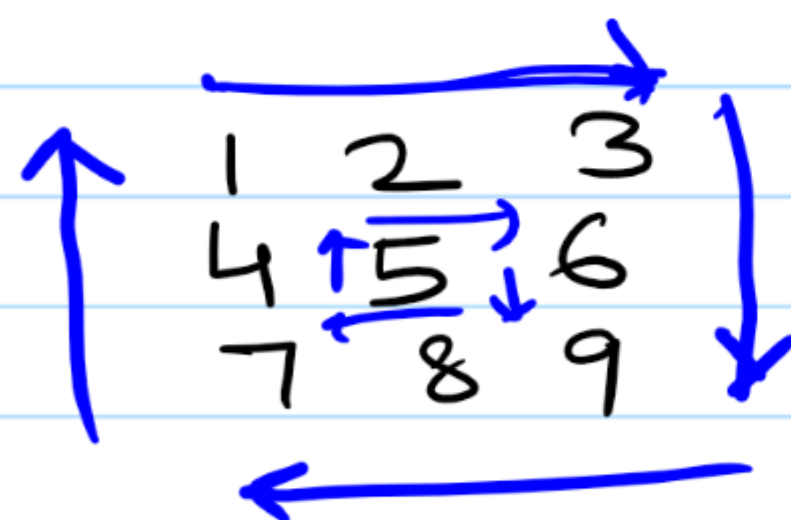
nxm

$n \neq m$



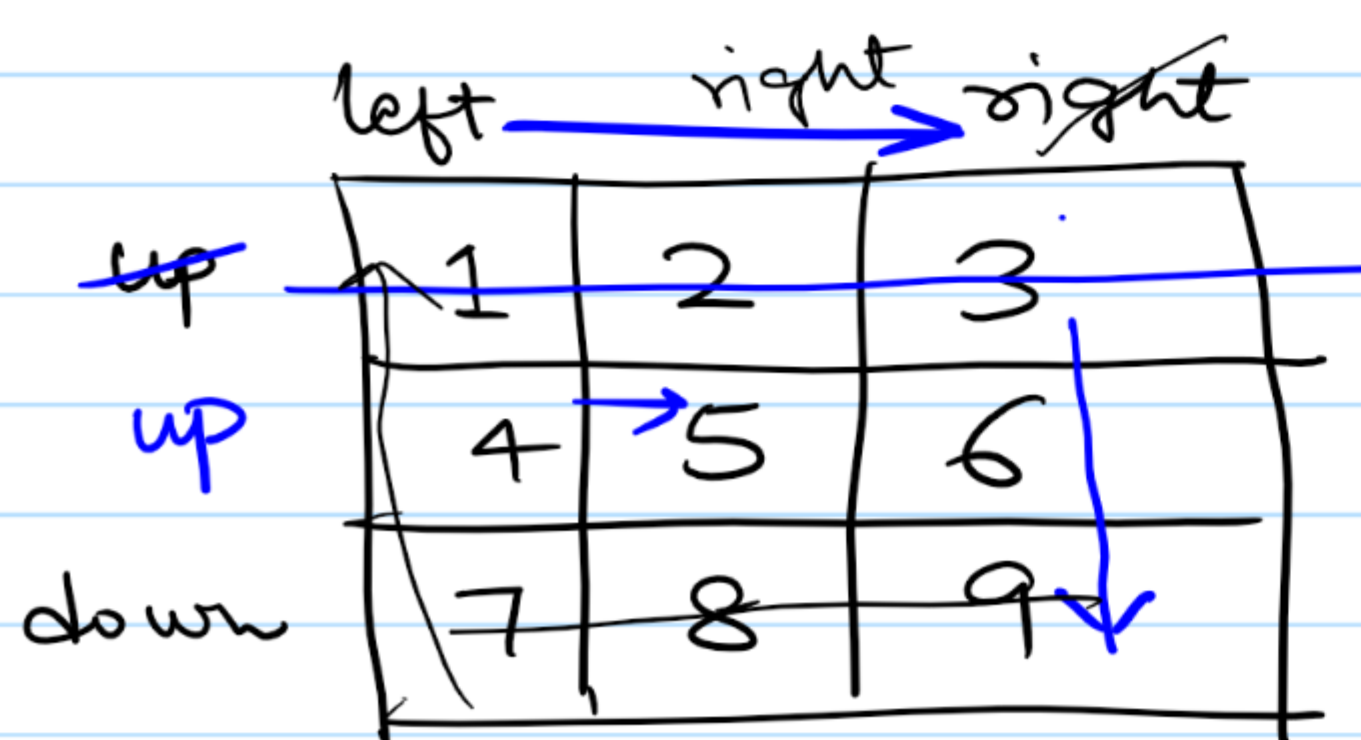
Q1

Print matrix in spiral order



3 4 6 8 7 5

1 2 3 6 9 8 7 4 5



1 2 3 6 9 8 7 4 5

L2R

```
for (int col = left; col <= right; col++)
    print arr[up][col];
    up++;
```

T2B

```
for (int row = up; row <= down; row++)
    print (arr[row][right]);
    right--;
```

left → right
up → down
1 2 3 4 = 1 2 3 4 4 3 2 1

right = 3
left = 0

if (left <= right) {

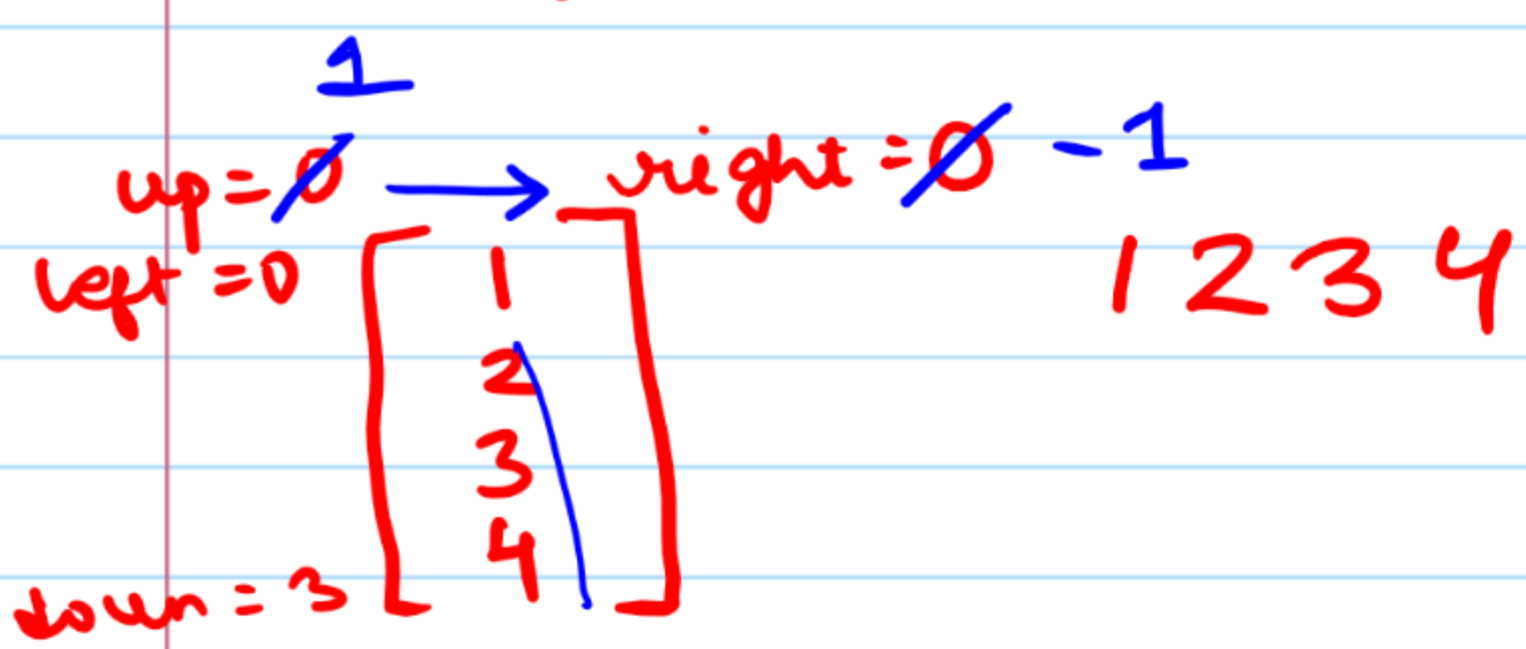
up = 1
down = 0


```
for (int row = down; row > up; row--)
```

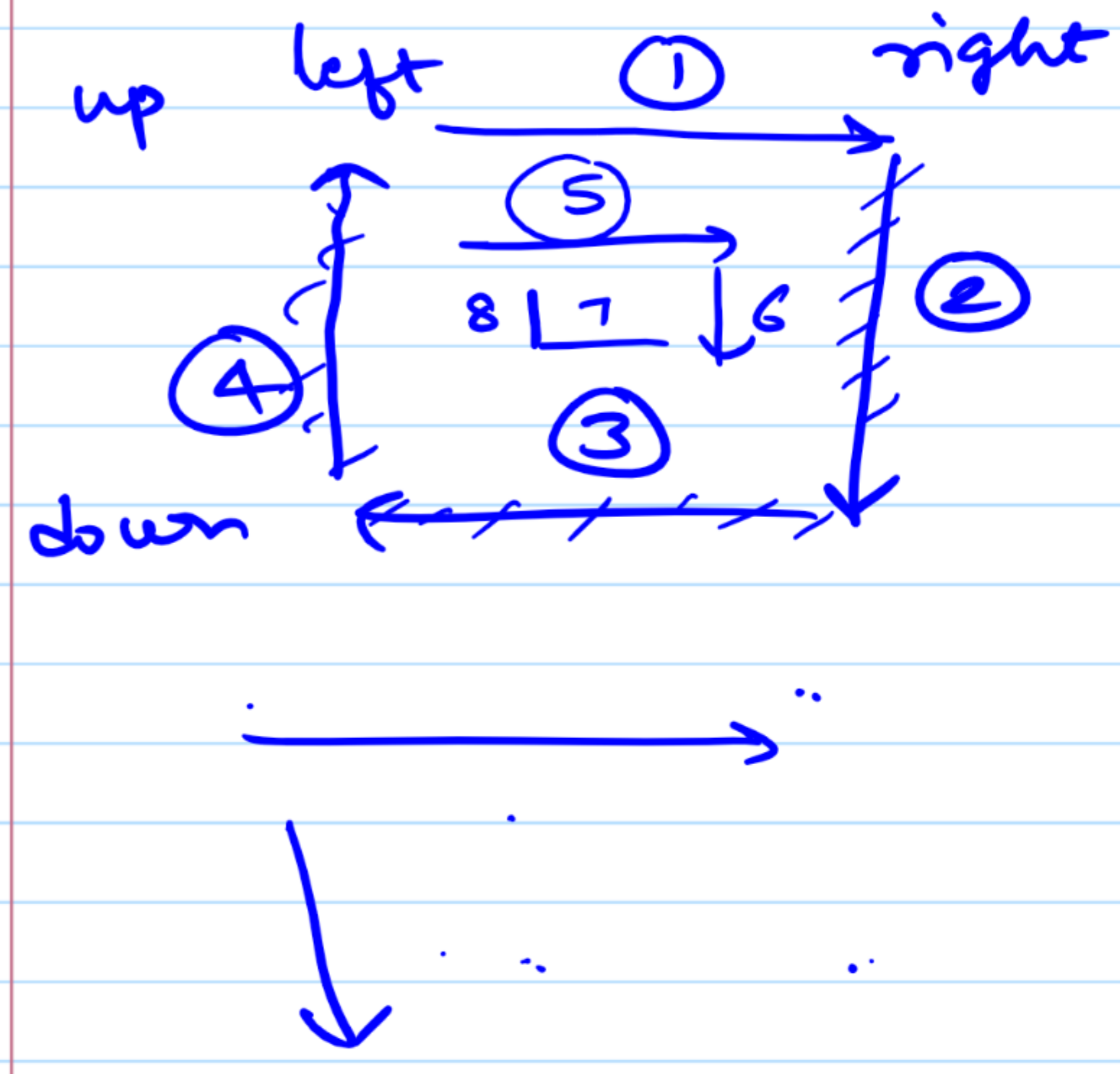
```
{
    print (arr[row][left]);
    left++;
}
```

```
right--;
if (up <= down)
```

```
{
    for (int col = right; col >= left; col--)
        print (arr[down][col]);
    down--;
}
```



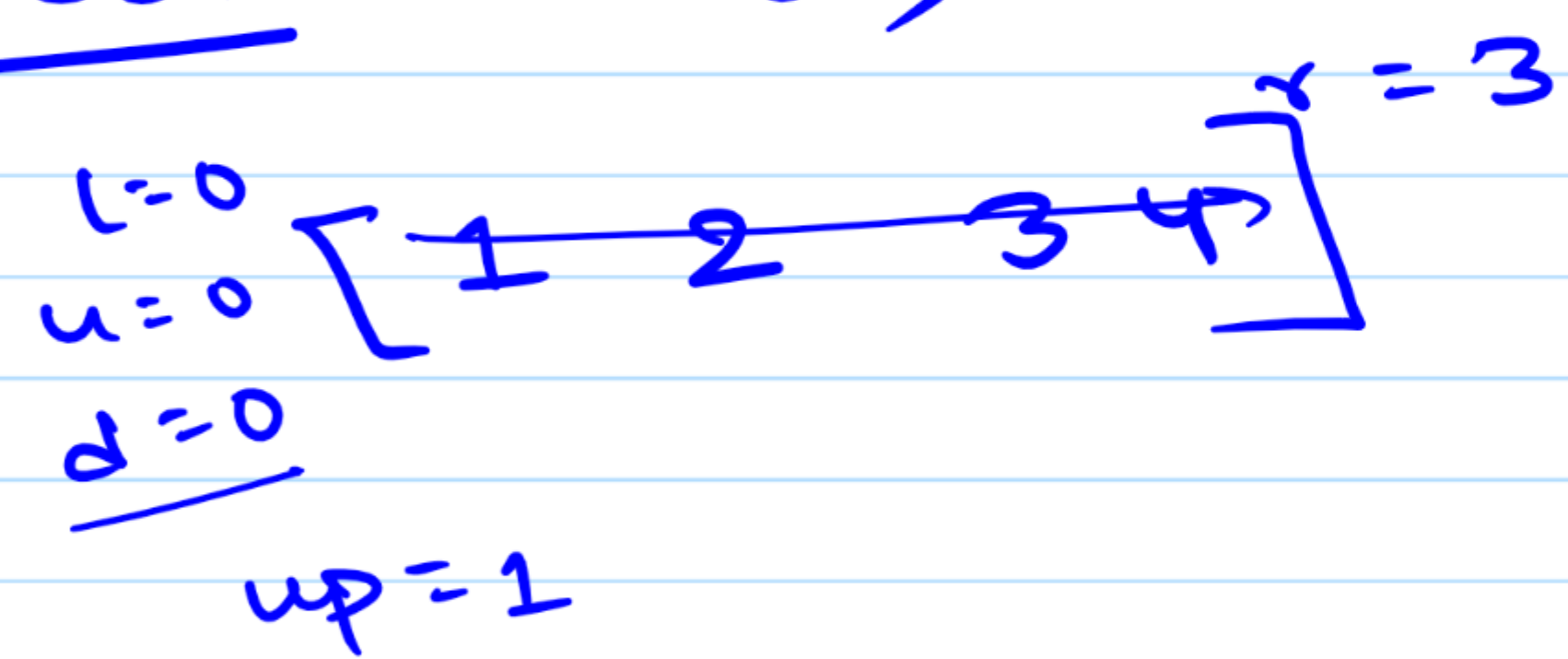
1 2 3 4



left to right
up++
top to bottom
right--;
if (right < left)
right to left;
down--;
if (down < up)
down to up
left++;

TC: M rows N col $O(M \times N)$

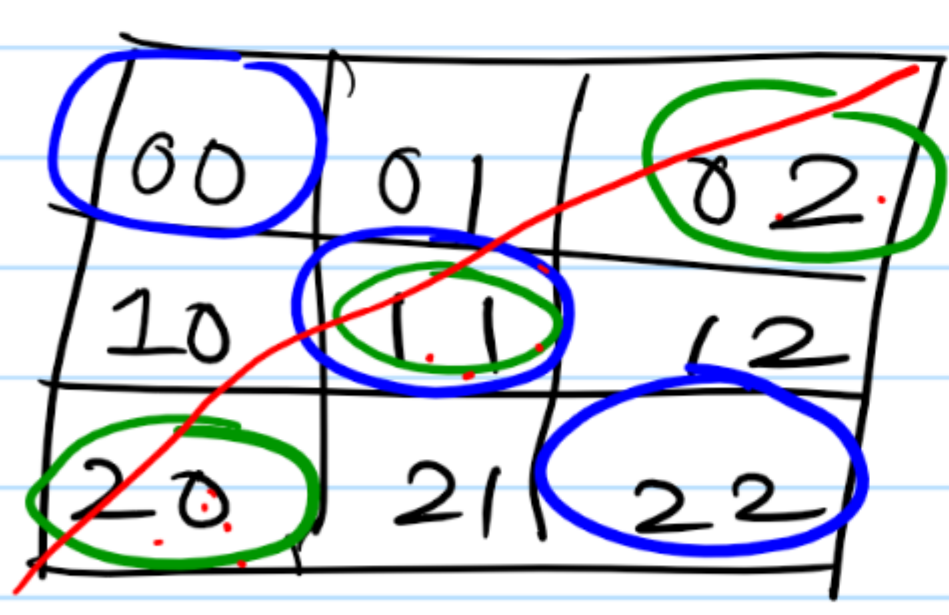
SC: $O(1)$



1 2 3 4
4 3 2 1 if (u <= d)

Q2 Matrix diagonal sum

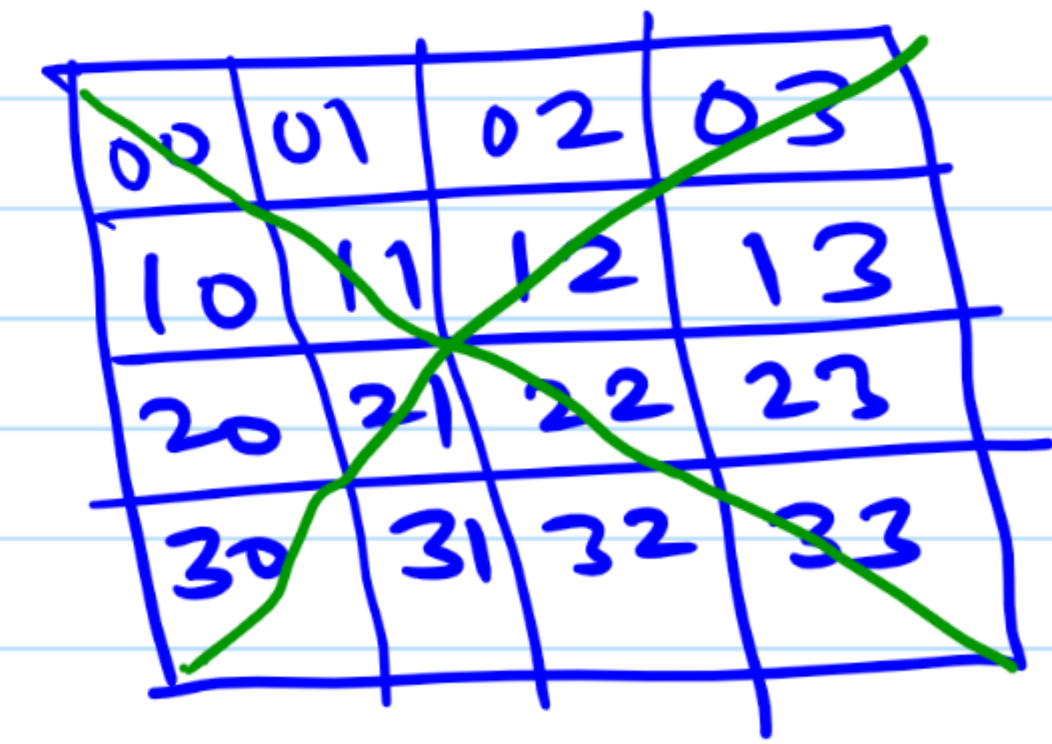
3x3



= 2

n x n

i = j
i + j = n - 1
j = n - 1 - i



mat[3/2][3/2]
mat[i][i]

```
{
    ans = 0;
    for (int i = 0; i < n; i++)
    {
```

ans += mat[i][i]; // adding primary diagonal elements

ans += mat[i][n-1-i]; // adding secondary diagonal elements

duplicate element

```
}
if (n % 2 == 1) // odd order matrix
```



```

    { ans -= mat[n/2][n/2];
    }

```

$$3/2 = 1$$

return ans;

}

TC: $O(N)$

SC: $O(1)$

Q3: count negative numbers

0	4	3	2	-1
3	2	1	-1	
1	1	-1	-2	
3	-1	-1	-2	-3

rows & columns are in decreasing order

count = 0;

Brute force:

```

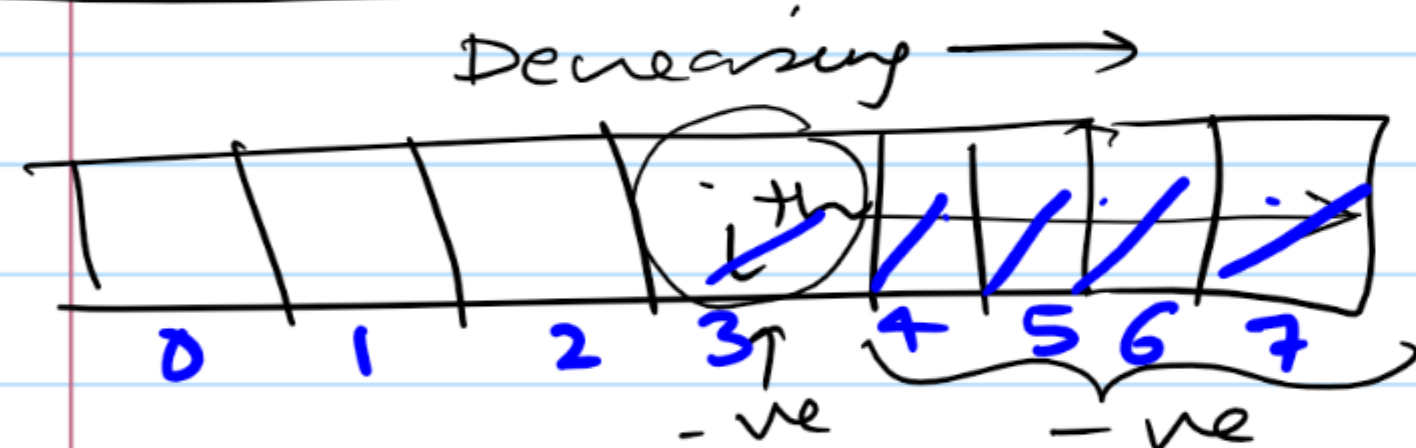
for (i = 0 ... m-1)
    for (j = 0 ... n-1)
        if mat[i][j] < 0
            count++;
return count;

```

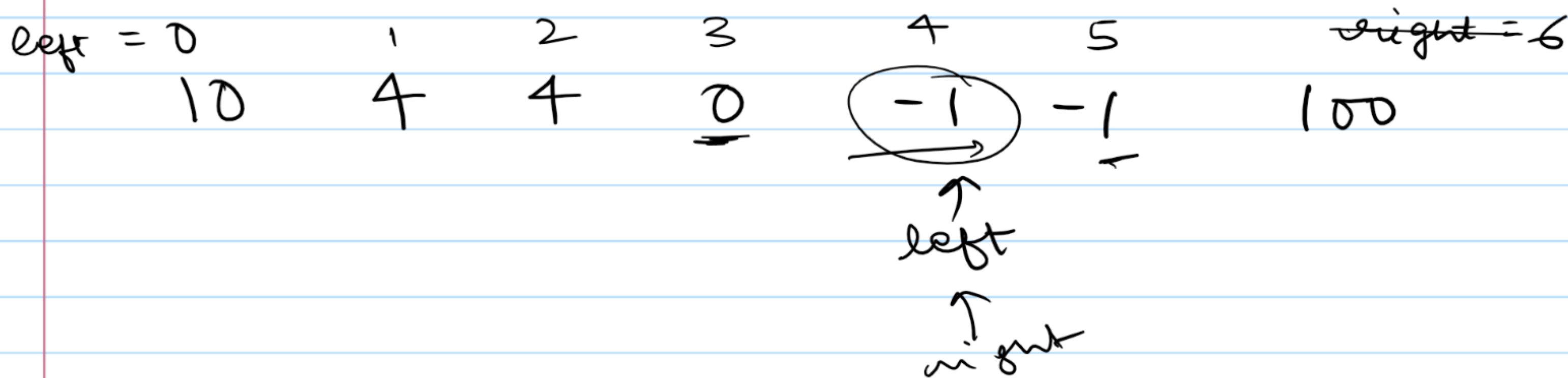
TC: $O(m \times n)$

SC: $O(1)$

Approach 2: Hunt: Binary Search



$$(n - i) = 8 - 3 = 5$$



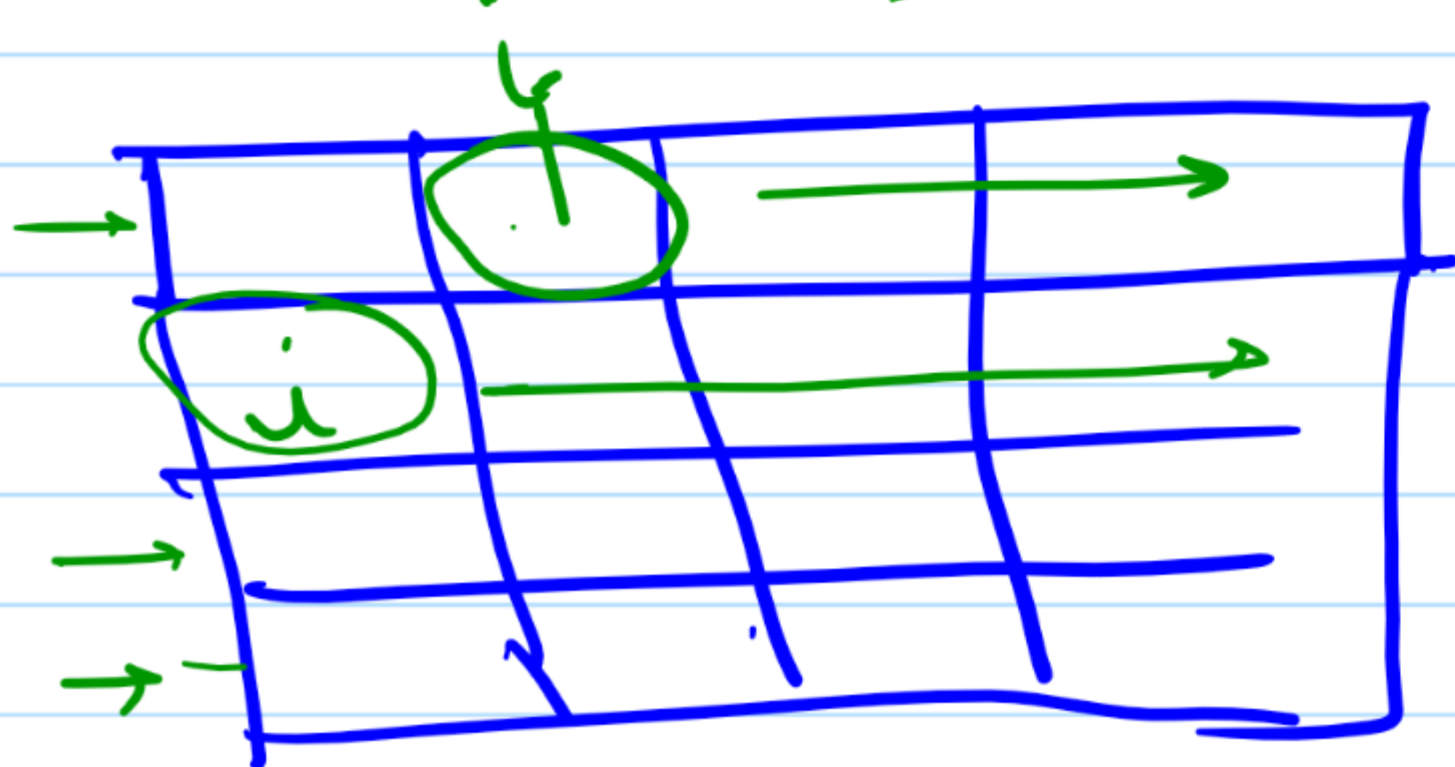
$$\text{mid} = \frac{\text{left} + \text{right}}{2} = \frac{0 + 6}{2} = 3$$

$$\text{left} = \text{mid} + 1 = 4$$

$$\text{mid} = \frac{4 + 6}{2} = 5$$

$$\text{right} = 4$$

n $n/2$... 1 $(\log n)$



Algo:

1) Initialize count = 0 ← Total -ve elements

2) Let n be no. of columns, m rows

iterate on each row of the matrix &

find the index of 1st -ve element (left).

3) all element from left to n will be
-ve \therefore row is in decreasing order.

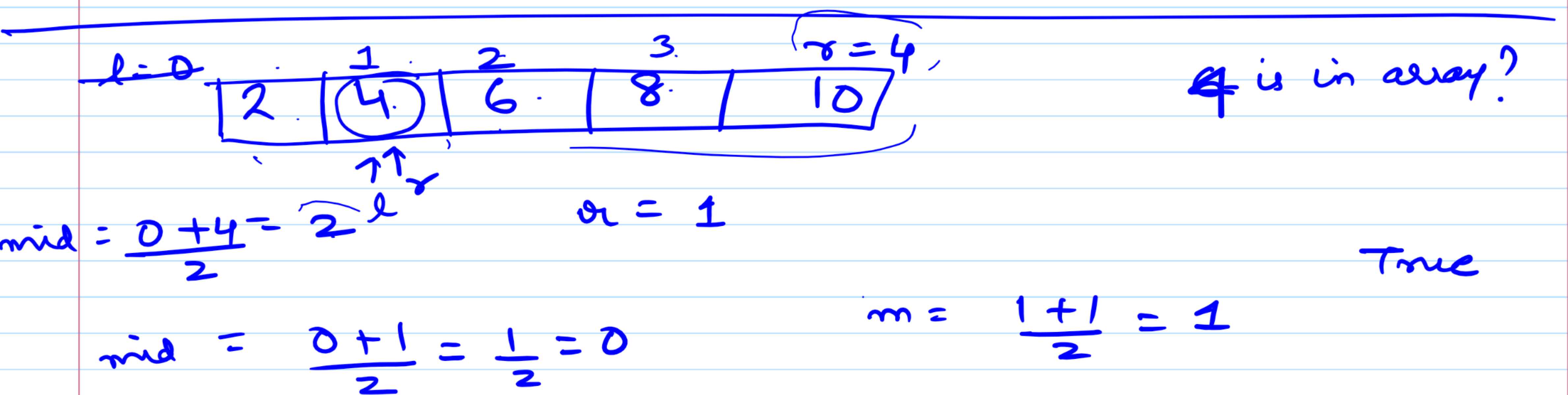
Increment count by $(n - \text{left})$

TC: n elements in each row

BS: $O(\log n)$

$$m \times \log n = O(m \log n)$$

SC: $O(1)$



Q4: Richest customer wealth (M customers N accounts)

$\begin{bmatrix} 1 & 2 & 8 \\ 4 & 1 & 2 \end{bmatrix}$

$$1^{\text{st}} = 1 + 2 + 8 = 11$$

$$2^{\text{nd}} = 4 + 1 + 2 = 7$$

1st

Algorithm: 1) Calculate row sum for each row
2) Return row with max. sum

TC: $O(M \times N)$

SC: $O(1)$