



# Functional Programming

## Lecture 2: Lambda abstraction

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# Last lecture

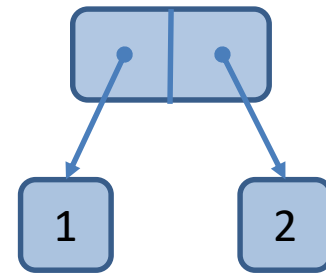
- What is (pure) functional programming
- Why do we care about it?
- Recursion is the main tool
- Scheme
  - S-expression, quote, identifiers, define, if, cond

# Pairs

Allow to construct compound data structures

```
(cons 1 2)
```

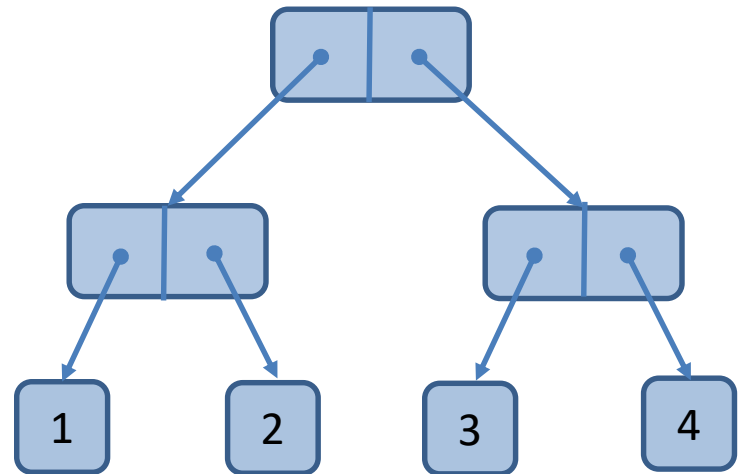
=> (1 . 2)



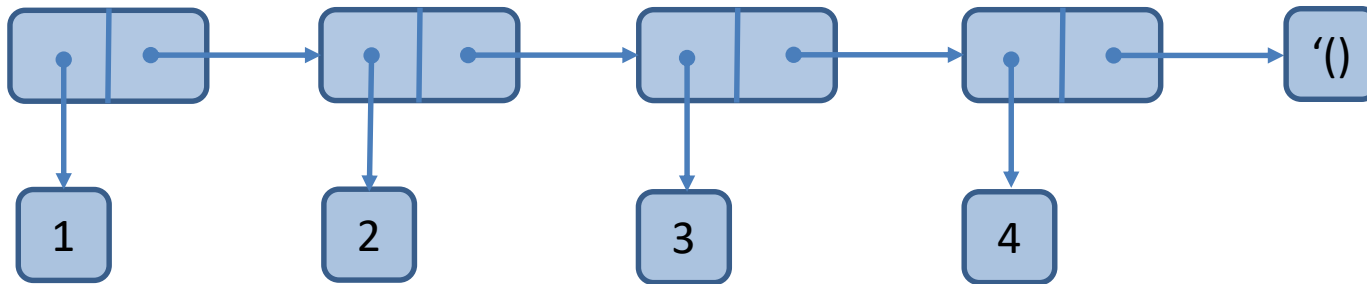
```
(cons
```

```
  (cons 1 2)
```

```
  (cons 3 4) )
```



# Lists



Lists are linked lists of pairs with '()' at the end

S-expressions are just lists

```
'(+ 1 2 3 4 5)
```

Lists can be created by a function `cons` or

```
(list item1 item2 ... itemN)
```

# Lists

Pairs forming the lists can be decomposed by

`car`    *[car]*                    first element of the pair

`cdr`    *[could-er]*                second element of the pair

`(caddr x)`    shortcut for `(car (cdr (cdr x)))`

Empty list is a null pointer

`null?`    tests whether the argument is the empty list

# Append

```
;;; Append two lists
(define (append2 a b)
  (cond
    ((null? a) b)
    (else (cons (car a)
                  (append2 (cdr a) b))))
  )
)
```

# Equality

Function = is only for numbers

Equivalence of the objects `eqv?`

```
(eqv? 1 1), (eqv? 'a 'a) ==> #t
```

```
(eqv? (list 'a) (list 'a)) ==> #f
```

More restrictive version is `eq?`

Typically the same pointer

Recursive version of `eqv?` on lists is

```
(equal? (list 'a) (list 'a)) ==> #t
```

# Debugging Basics

## Tracing function calls and returns

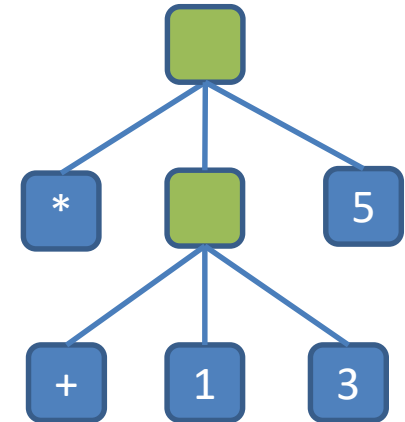
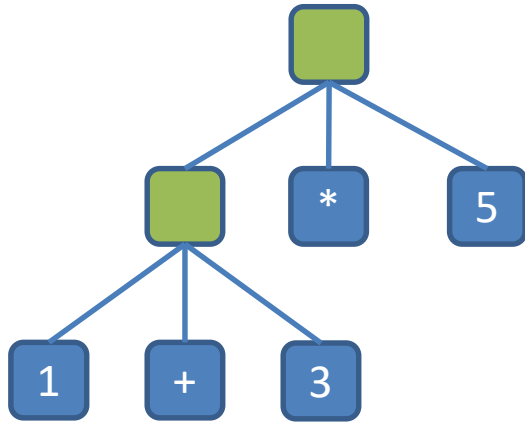
```
#lang scheme
(require racket/trace)
(trace append2)
(untrace append2)
```

## Helper print-outs

```
(begin (display x)
       (newline)
       <do-work>)
```



# Infix -> prefix



( ( 1 + 3 ) \* 5 )

( \* ( + 1 3 ) 5 )

# Evaluation strategy

Defines the order of evaluating the expressions  
influences program termination, not the result

Evaluation of scheme is eager (or strict)

left to right

evaluate all arguments before executing a function

Evaluation of some special forms is lazy

if, and, or, **lambda**

# Functions are 1st-class citizens

- They may be named by variables
- They may be passed as arguments to a function
- They may be returned by a function
- They may be included in data structures

# Lambda abstraction

A construction for creating nameless procedures

```
(lambda (arg1 ... argN) <expr>)
```

Define for functions is an abbreviation

```
(define (<var> <formals>) <body>)
```

Is the same as

```
(define <var>  
  (lambda (<formals>) <body>))
```

# Filter

```
;;; Filter a list by the given predicate
(define (my-filter pred lst)
  (cond
    ((null? lst) '())
    ((pred (car lst))
     (cons (car lst)
           (my-filter pred (cdr lst))))
    (else (my-filter pred (cdr lst))))
  )
)
```

# Derivative

$$Dg(x) \approx \frac{g(x + dx) - g(x)}{dx}$$

```
(define dx 0.00001)
(define (deriv g)
  (lambda (x)
    (/ (- (g (+ x dx)) (g x)) dx)
  )
)
```

# Functions as data

```
(define (mult-fn fns x y)
  (cond
    ((null? fns) '())
    (else
     (cons
      ((car fns) x y)
      (mult-fn (cdr fns) x y))))
  )
)
```

# Let

## Motivation

reuse of computation/result is often required

e.g., `minimum`, `roots` from the labs

```
(let ( (<var1> <exp1>)
      (<var2> <exp2>))
      <body-using-var1-var2>)
```



$$f(x, y) = x(1 + xy)^2 + y(1 - y) + (1 + xy)(1 - y)$$

```
;;; Local variables
(define (f x y)
  (let
    ((a (+ 1 (* x y)))
     (b (- 1 y)))
    (+ (* x a a)
       (* y b)
       (* a b)))
  )
)
```

# Implementing let

```
(let ((x <exp1>)
      (y <exp2>)) <body>)
```

Can be implemented as

```
((lambda (x y) <body>) <exp1> <exp2>)
```

# Let as lambda

```
(define (f2 x y)
  ((lambda (a b)
    (+ (* x a a)
      (* y b)
      (* a b)))
    (+ 1 (* x y))
    (- 1 y))
)
```

# Let\*

We might want to use the earlier definitions in the following.

```
(let ((x <exp>))  
      (let ((y <exp-with-x>)) <body-x-y>)
```

Equivalent to

```
(let* ((x <exp>)  
       (y <exp-with-x>)) <body-x-y>)
```

# Quicksort

```
(define (qsort lst cmp)
  (cond
    ((null? lst) '())
    (else (let*
              ((pivot (car lst))
               (smaller (lambda (x) (cmp x pivot)))
               (greater (lambda (x) (not (cmp x pivot)))))
              (append
                (qsort (filter smaller (cdr lst)) cmp)
                (list pivot)
                (qsort (filter greater (cdr lst)) cmp))
              )
            )
    )
  )
```

# Scheme home assignments

Three connected assignments

- Robot simulation

- Population evaluation

- Code synthesis

Why this assignment?

Work on your own

Submit by midnight of the day after your lecture

<https://cw.felk.cvut.cz/brute/> (in 2 weeks)