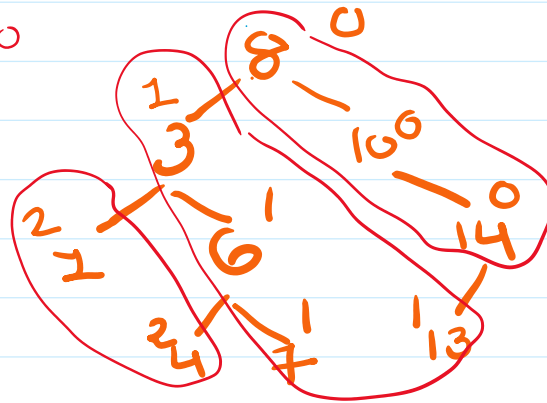


gfg

Diagonal Traversal

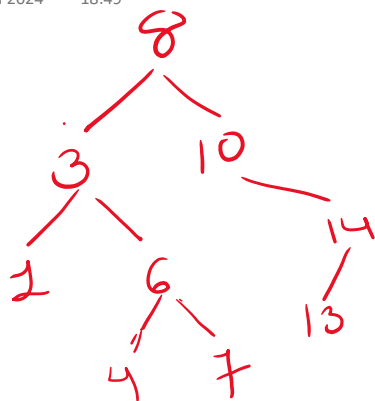
$d = 0$



=> going left +1
going right +
noting

$0 \Rightarrow 8 \quad 10 \quad 14$
 $1 \Rightarrow 3 \quad 6 \quad 7 \quad 13$
 $2 \Rightarrow 1 \quad 4$

map using

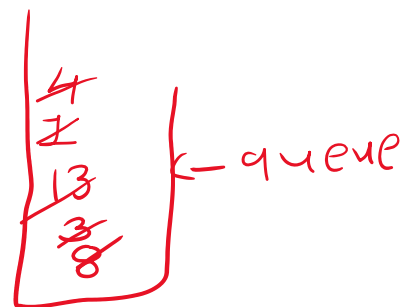


right print
left ele
push

8 10 14

3 6 7

13 1 4



temp \Rightarrow 8

while(temp != null)

{

print temp val

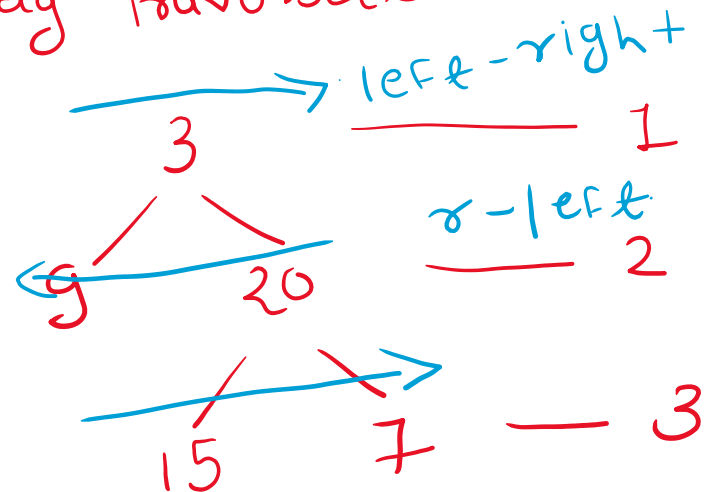
if (temp \rightarrow left)

q.push

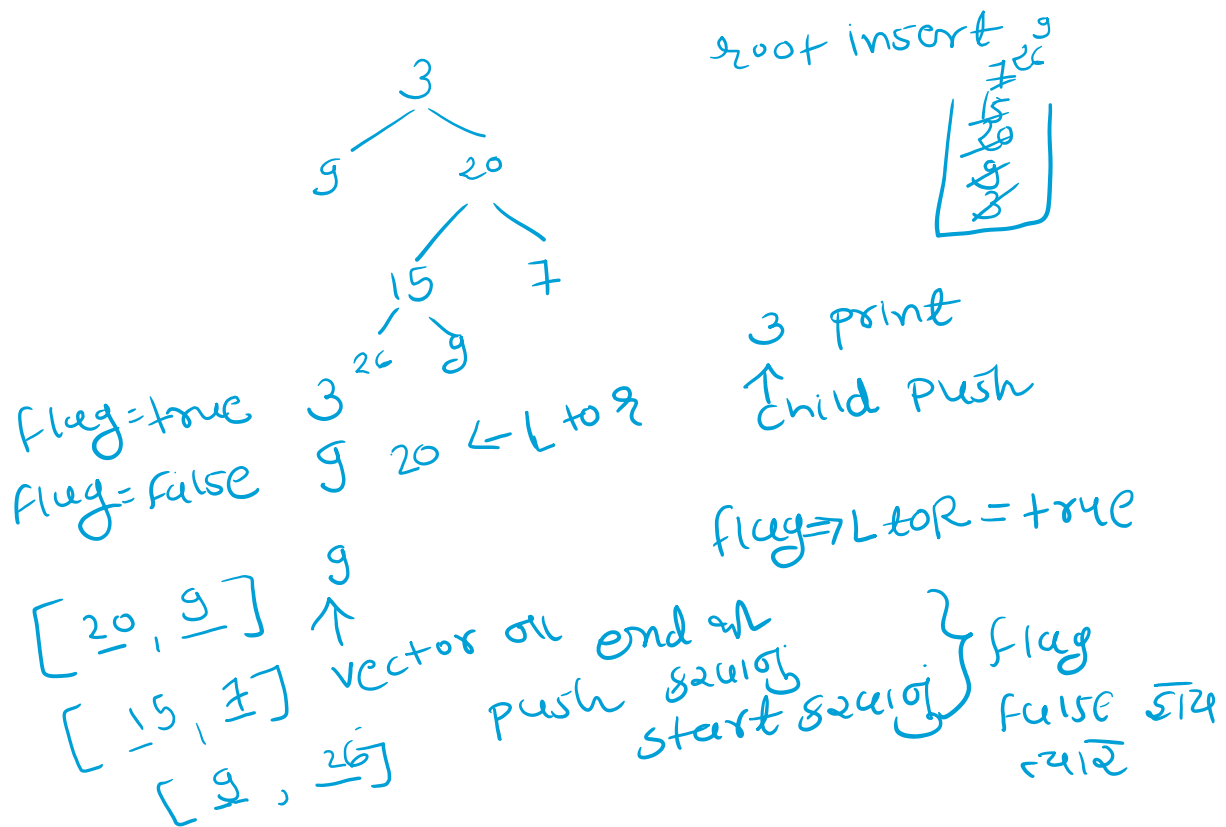
temp = temp \rightarrow right

Lc:- 103

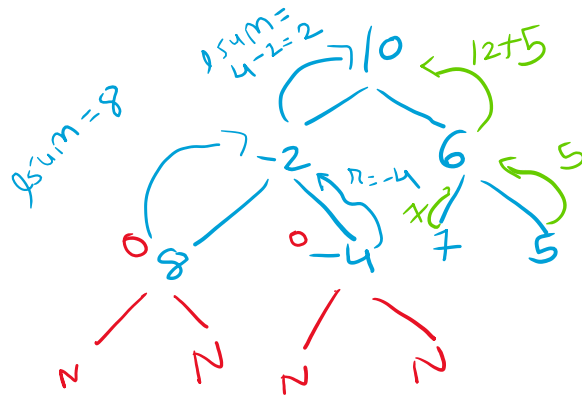
Zig-zag Traversal



[3] [20 9]
[15 7]

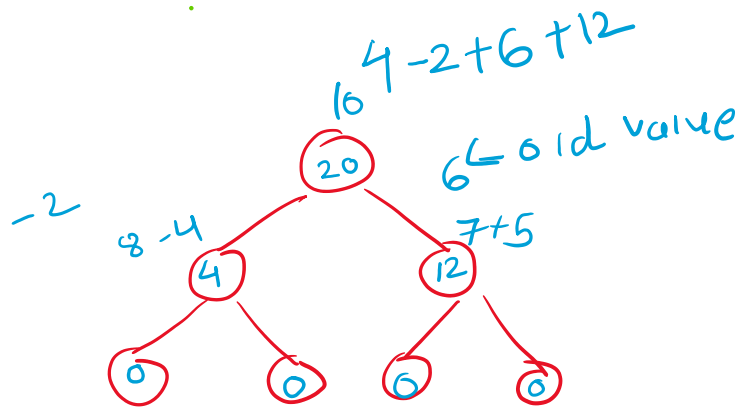


gfg: Transform to sum Tree



(1) left Node will be zero

left, right
value sum

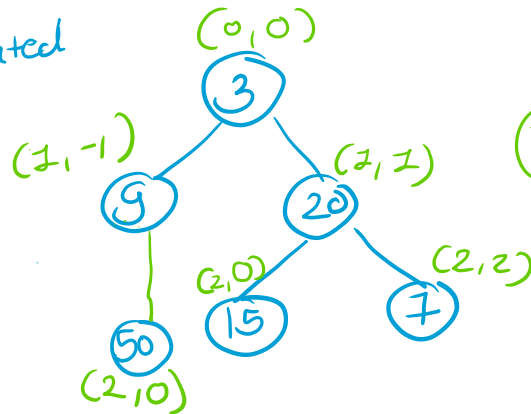


987

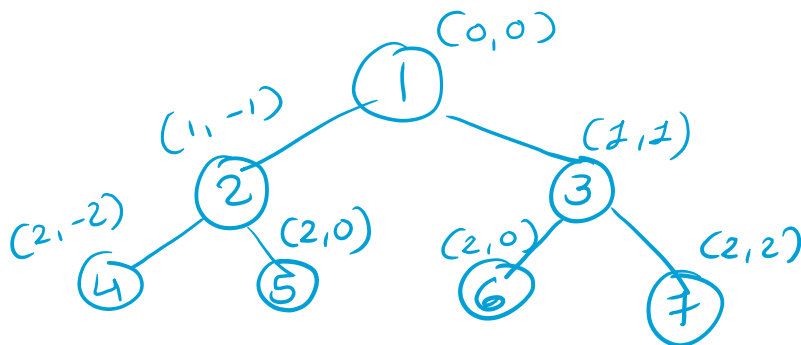
Vertical Order tra

if both node co-ordinates are same

-1 = 9
0 = 3, 5, 15
1 = 20
2 = 7



(row, col)
(row+1, col-1) (row+1, col+1)



-2 = 4
-1 = 2
0 = 1, 5, 6
1 = 3
2 = 7

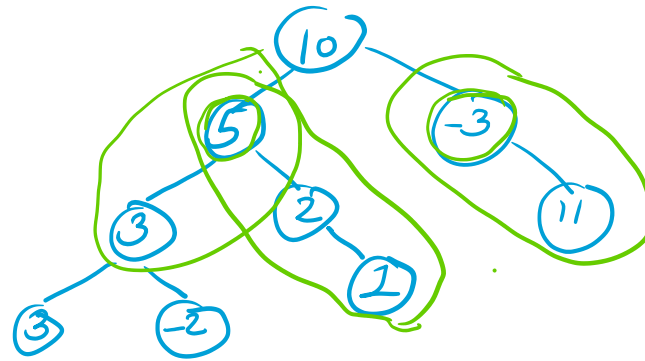
create a map
map { Col → map < row, multiset > }
sorted
sorted and multiple value are store

Col row val

0 0 → 1

437

k-sum paths

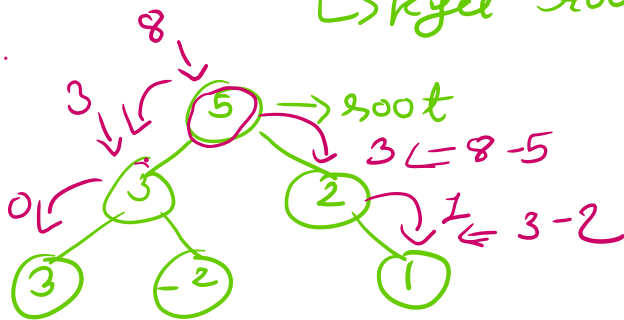


target sum = 8

path = 3

Break it down

↳ kya root se path aata hai?



gu

morris Traversal

Inorder traversal $\Rightarrow TC(O(n)) \Rightarrow O(n)$ } space
 $SC(O(n)) \Rightarrow O(1)$ }



2 1 3

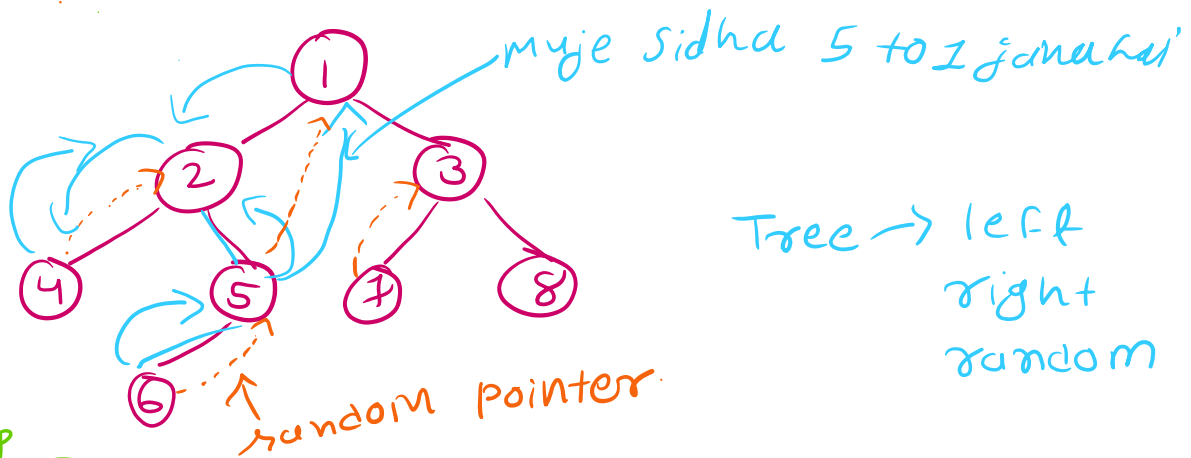
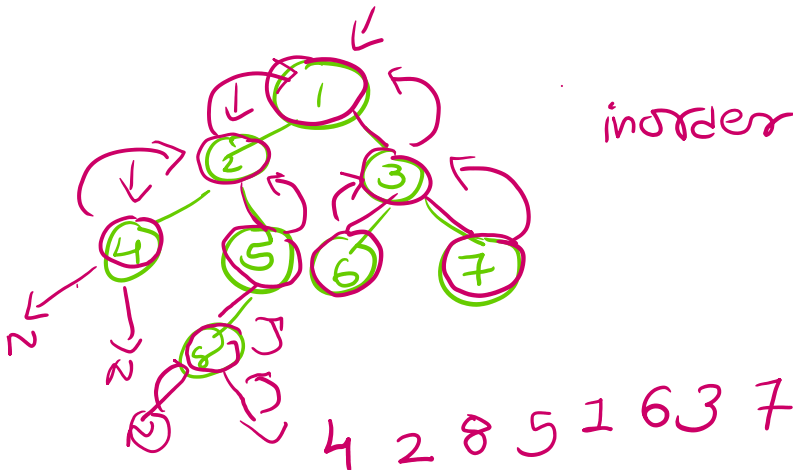
↑ this is a morris-traversal.

inorder(root->left)

cout << root->data

inorder(root->right)

morris Traversal \Rightarrow inorder $\Rightarrow O(n) = TC$
 $O(1) = SC$



inorder predecessor

L) in tree Node ki previous se Node main k value us se

1 → 5
7 → 1

while(curr) curr = root;

{
if left is null then, visit it

else { left ≠ null

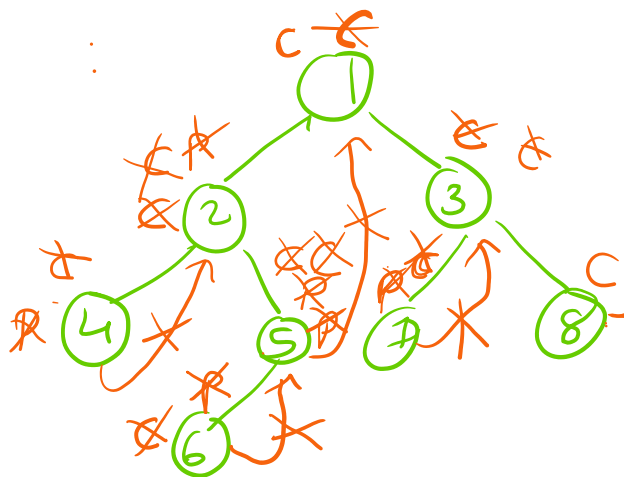
① predecessor find

② if right node null then, go left after establishing link from predecessor to curr

else

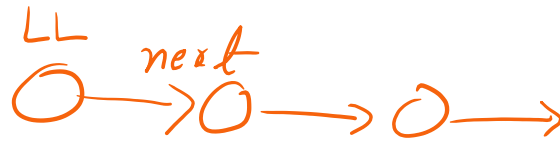
③ left is already visited, go right after visiting curr, & remaining connection

}

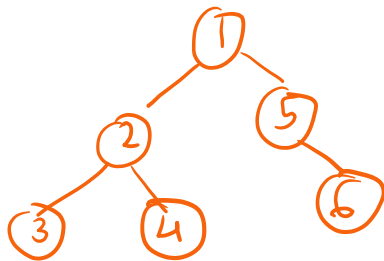


4 2 6 5 1 7 3 8

114 Flatten BT into LL (> convert



right = next
 left = null

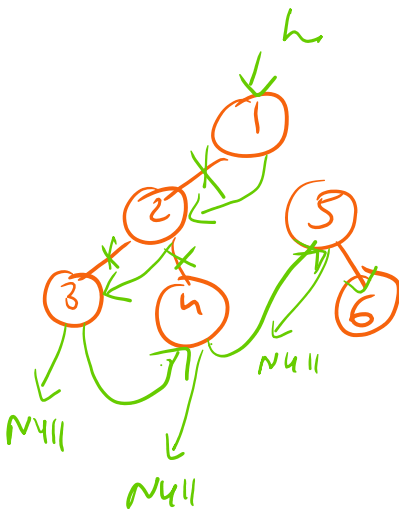


Preorder NLR

1 2 3 4 5 6

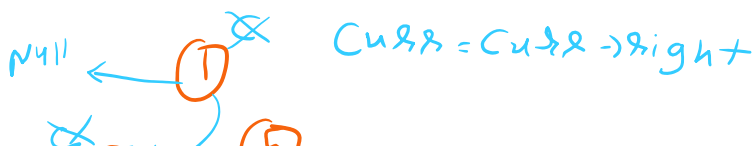
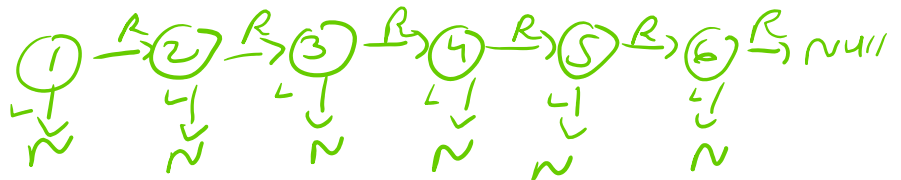
right pointer

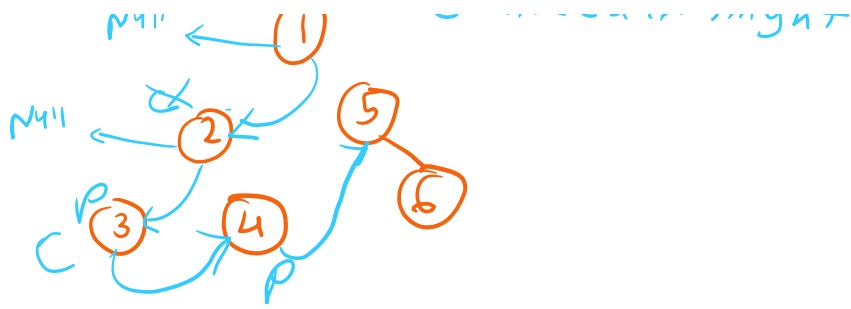
left pointer = null



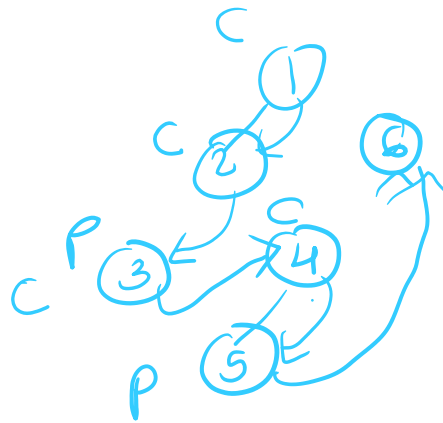
Head ① → root

1 2 3 4 5 6





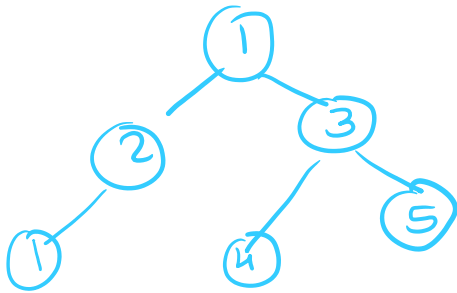
$prev \rightarrow right = curr \rightarrow right$
 $curr \rightarrow right = curr \rightarrow left$
 $curr \rightarrow left = null$
 $move\ curr = curr \rightarrow right$



1 2 3 4 5 6

gfg

maximum sum of non-adjacent nodes



if we choose 1 then that child not choose

$$C1 = 1, 1, 4, 5 = 11$$

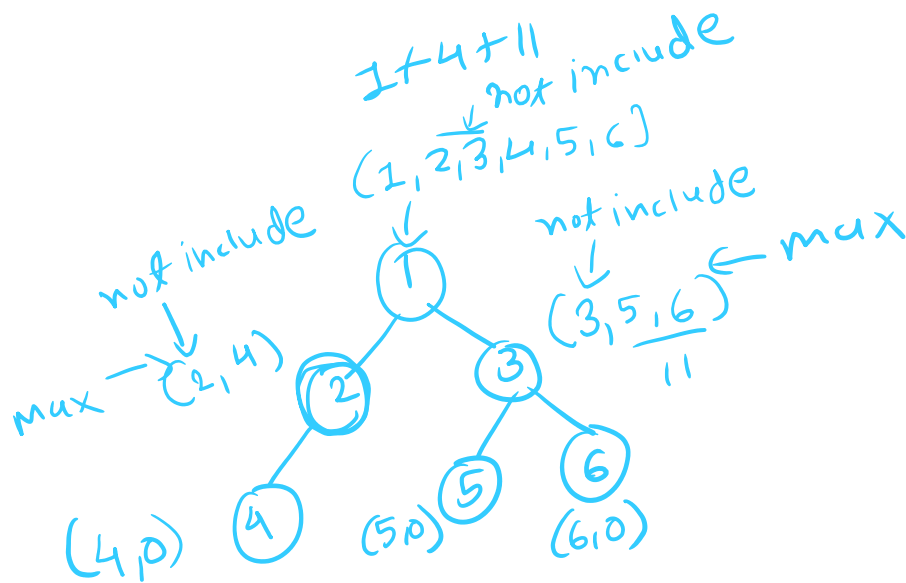
$$C2 = 2, 3 = 5$$

$$C3 = 1, 4, 5, 1 = 11$$

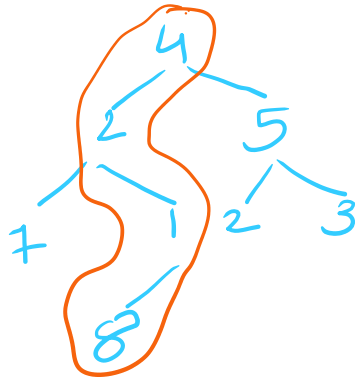
$$C4 = 4, 5, 1, 1 = 11$$

$$C5 = 5, 4, 1, 1 = 11$$

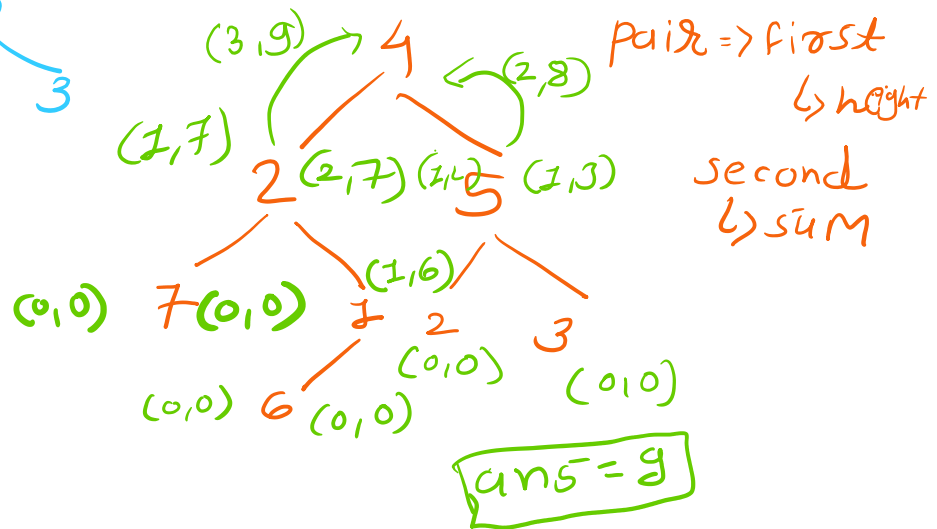
sum \Rightarrow including all nodes at that level.



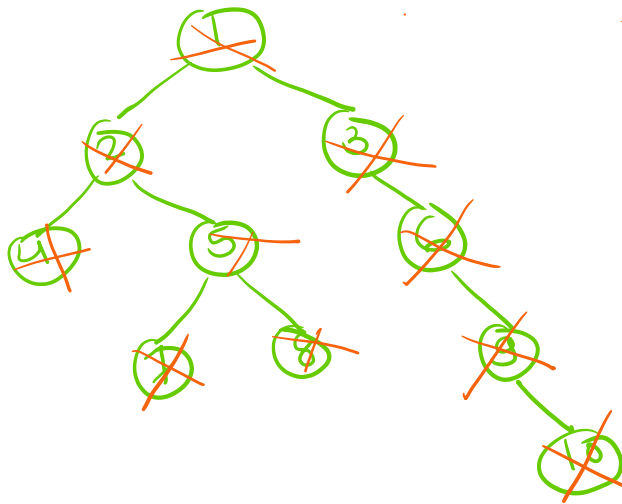
gfg Sum of the longest Bloodline
of a tree



$$ans = 4 + 2 + 1 + 8$$



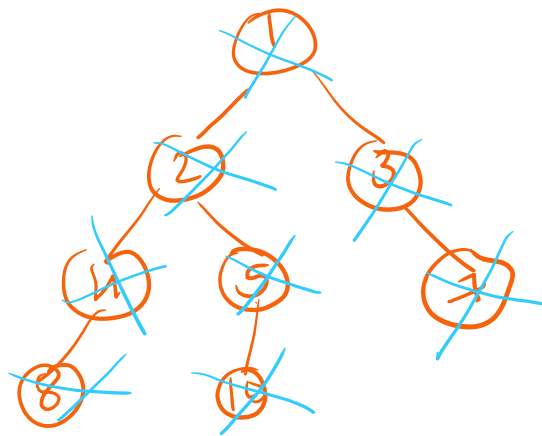
gfg Burning Tree



burn
 $t=0, 8$
 $t=1, 5$
 $t=2, 7, 2$
 $t=3, 4, 1$
 $t=4, 3$
 $t=5, 6$
 $t=6, 9$
 $t=7, 10$

7 sec burn

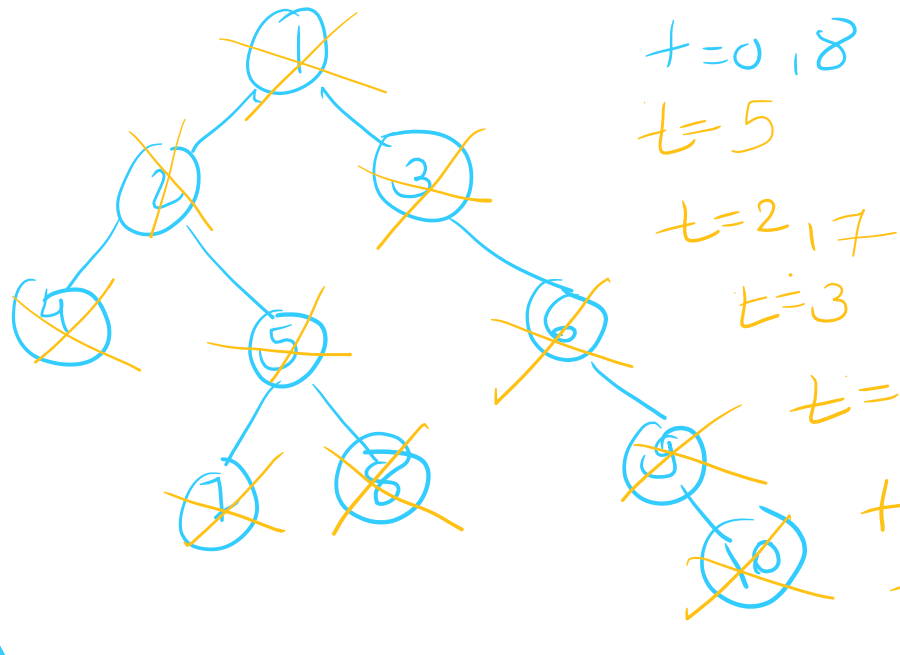
target = 10



$t=0, 10$
 $t=1, 5$
 $t=2, 2$
 $t=3, 4, 1$
 $t=4, 8, 3$
 $t=5, 7$

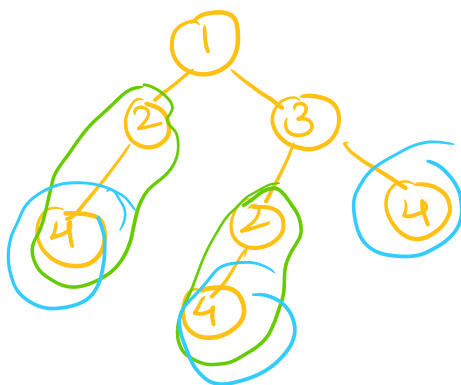
ans $t=5$

- ① Find target node
- ② make node to parent node mapping

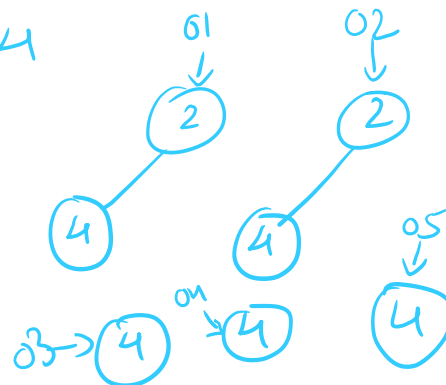


Visited
 queue $C =$ ~~1~~ ~~2~~ ~~5~~ ~~2, 7~~ ~~1, 4~~ ~~3~~ ~~6~~ ~~8~~ ~~10~~

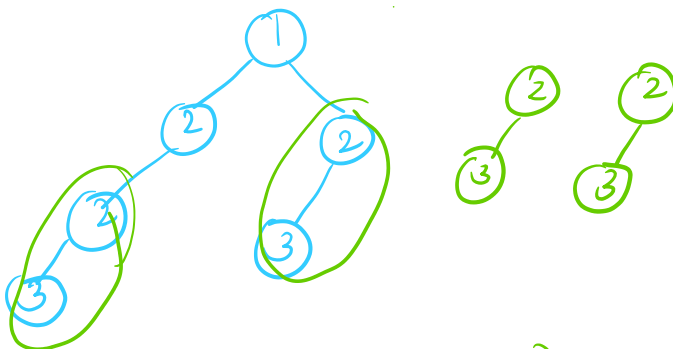
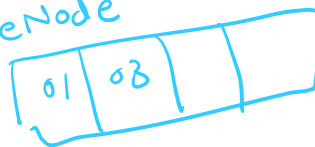
G52 Find Duplicate subtrees



2,4
4

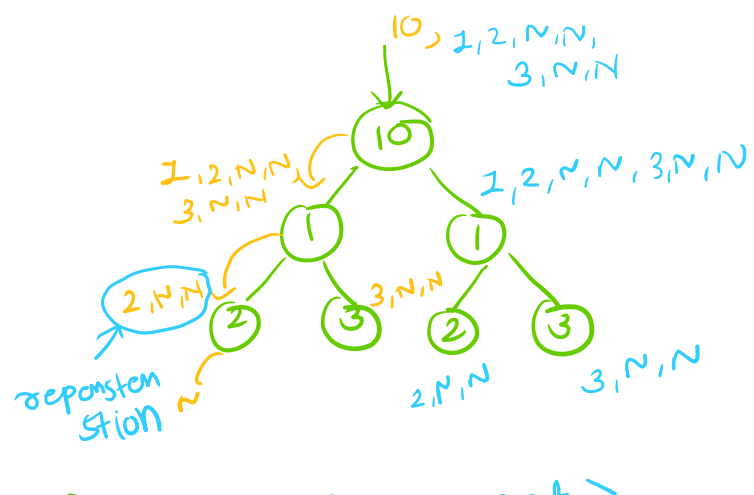
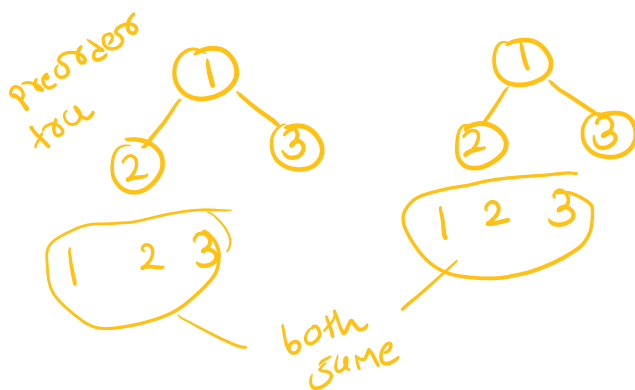


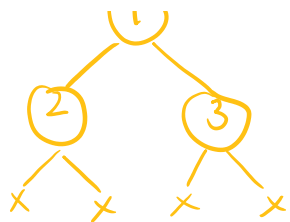
TreeNode



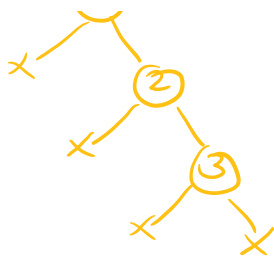
NI Boucle Forcée $O(n^2)$

↳ one node ^a then check
other node ^b





1, 2, N, N, 3, N, N



1, N, 2, N, 3, N, N

unique identify

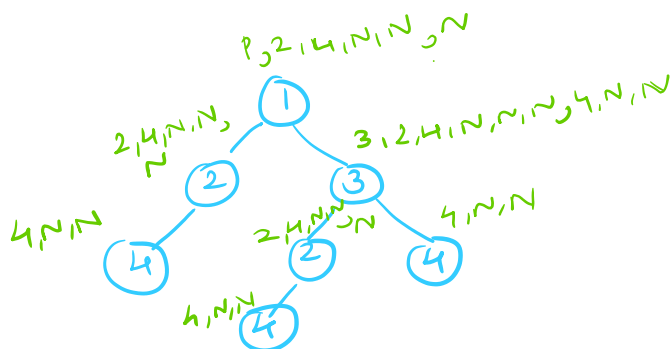
stion

Map<String, int>

2, N, N x 2
3, N, N x 2

1, 2, N, N, 1 x 2
3, N, N

if Map hai equal
storing use krni
ans = 2 + 3 = 5
node



4, N, N x 2 3
2, 4, N, N, N x 2
3, 2, 4, N, N, N, N, 1
4, N, N, N

ans = 4, 2

if 1 equal krni
8 ans hai store.