Reusing ITT Compiled Code



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Reusing Just-in-Time Compiled Code

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```
make_coffee (...) {
   Get xx% coffee
   Get yy% milk
    . . .
   Boil to zz degree
   Serve counter1;
```

A **sequence of instructions** to make coffee.

An execution takes 10 minutes

But we want to serve every order within **2 minutes**.

JIT is the principle of

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   Get xx% coffee
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JIT is the principle of

Profiling

Specialization

Amortization

```
make_coffee (...) {
   Get xx% coffee
   Get yy% milk
    . . .
   Boil to zz degree
   Serve counter1;
```

10 minutes

12g coffee, 40 ml milk, 2g sugar 12g coffee, 40 ml milk, 1sp honey 4g coffee, 60 ml milk, 0g sugar 12g coffee, 40 ml milk, 1sp honey 12g coffee, 40 ml milk, 1sp honey 4g coffee, 60 ml milk, 0g sugar

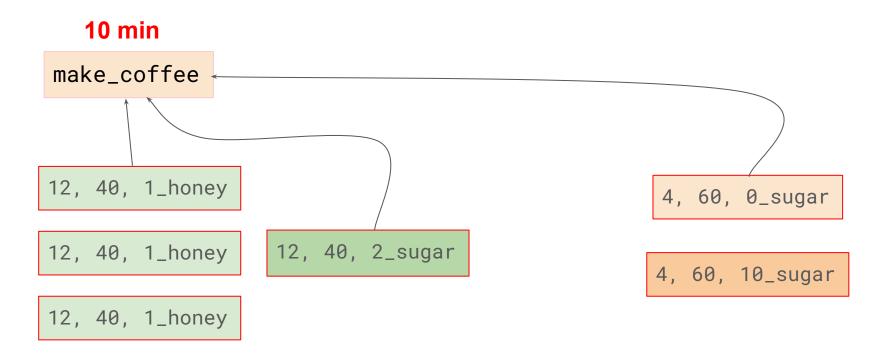
The perfect coffee, always

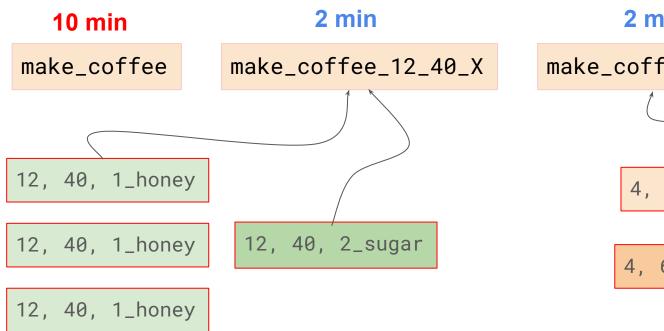
```
make_coffee (...) {
   Get xx% coffee
   Get yy% milk
   Boil to zz degree
   Serve counter1;
```

10 minutes + 10 sec profiling

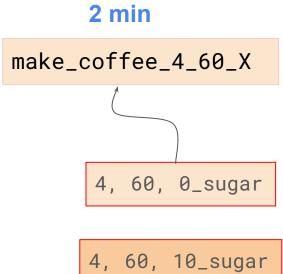
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12g coffee, 40 ml milk, 2g sugar
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12g coffee, 40 ml milk, 1sp honey
12g coffee, 40 ml milk, 1sp honey
4g coffee, 60 ml milk, 10g sugar
```

The perfect coffee, always

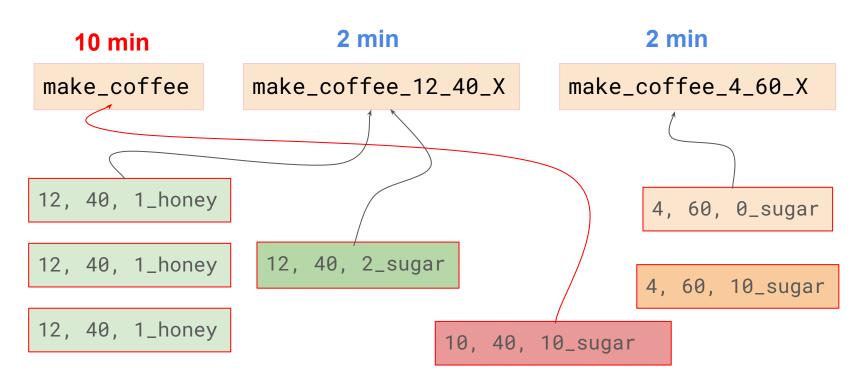




Specialization



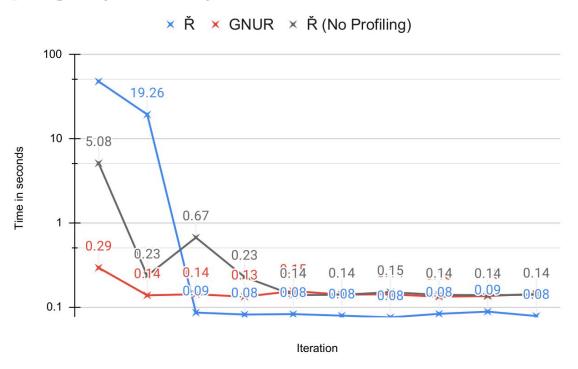
Amortization



The perfect coffee, always

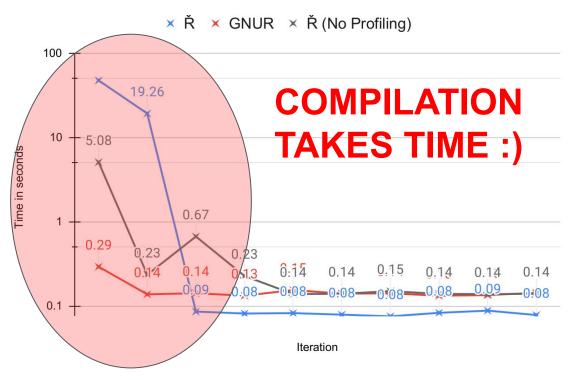
Motivation: Why reusing compiled code is useful

pidigits (shootout)



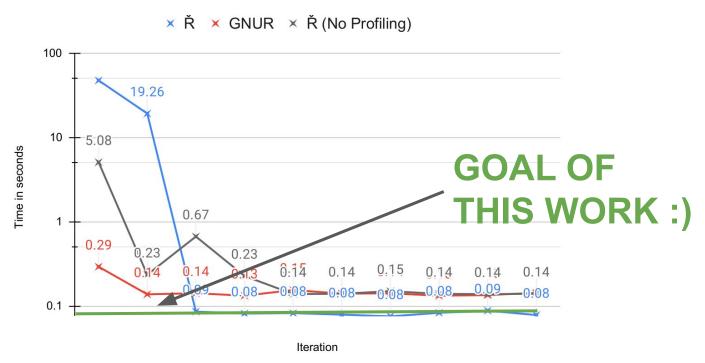
Motivation: Why reusing compiled code is useful

pidigits (shootout)



Motivation: Why reusing compiled code is useful

pidigits (shootout)



Motivation: Why reusing JIT compiled code is hard

```
foo <- function(x, y, z=TRUE) {
    x;
    if (z) res <- m1(x, y)
    else res <- m2(g, y)
    res <- bar(res, g)
    res
}</pre>
```

```
foo(□,□)

If we supply just two
arguments
```

Are these predicates true?

```
foo <- function(x, y, z=TRUE) {
    # x;
    if (z) res <- m1(x, y)
    else res <- m2(g, y)
    res <- bar(res, g)
    res
}</pre>
```

```
foo(□,□)
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If we supply just two
arguments

Are these predicates true?

We can always remove line 2

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foo <- function(x, y, z=TRUE) {
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    res <- bar(res, g)
    res
}</pre>
Are these predicates
true?

We can always remove line 2
```

Else condition is never taken

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foo <- function(x, y, z=TRUE) {
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    res
}</pre>
```

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foo function(x, y, z=TRUE) {
    x;
    if (z) res <- m1(x, y)
    else res <- m2(g, y)
    res <- bar(res, g)
    res
}</pre>
```

Lazy Evaluation

```
foo(\Box, \Box)
```

```
x <- function()
   assign("z",FALSE,
   sys.frame(-1))

foo(x(),2)</pre>
```

```
foo <- function(x, y, z=TRUE) {
                                        foo(\Box, \Box)
   X;
   if (z) res <- m1(x, y)
                           False branch is taken
   else res <- m2(q, y)
    res <- bar(res, g)
                                      x <- function()
    res
                                          assign("z", FALSE,
                                          sys.frame(-1))
                                      foo(x(),2)
```

```
foo <- function(x, y, z=TRUE) {
    # x;
    if (z) res <- m1(x, y)
    # else res <- m2(g, y)
    res <- bar(res, g)
    res
}</pre>
```

Call-Site Context

 $foo(\Box, \Box)$

We can always remove line 2, if x is non-reflective

Else condition is never taken, if x is non-reflective

```
foo' <- function(x, y, z=TRUE) {
   if (z) res <- m1(x, y)
   res <- bar(res, g)
   res
}</pre>
```

Call-Site Context

 $foo(\Box, \Box)$

We can always remove line 2, if x is non-reflective

Else condition is never taken, if x is non-reflective

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Else condition is never taken, if x is non-reflective

foo(10, 20)

g <- matrix(...)

```
foo" <- function(x, y, z=TRUE) {
    res <- x + y
    res <- ns::add(res, g)
    res
}

m1 <- function(a, b) a + b
bar <- ns::add</pre>
Callee Context
```

Type Context

```
foo" <- function(x, y, z=TRUE) {</pre>
                                          Call-Site Context
    res <- x + y
                                          foo(□,□)
    res <- ns::add(res, g)</pre>
    res
                                          Callee Context
m1 \leftarrow function(a, b) a + b
                                          m1, m2, bar
bar <- ns::add
foo(10, 20)
                                          Type Context
g <- matrix(...)
```

```
foo" <- function(x, y, z=TRUE) {</pre>
                                      Call-Site Context
   res <- x + y
   res <- ns::add(res, g)</pre>
                                      foo(□.□)
     COMPILATION HAPPENS UNDER CONTEXT
m1
bar <- ns::add
                                      IIII, III∠, Dai
foo(10, 20)
                                      Type Context
g <- matrix(...)
```

foo <- function(...)</pre>

One function

foo <- function(...)

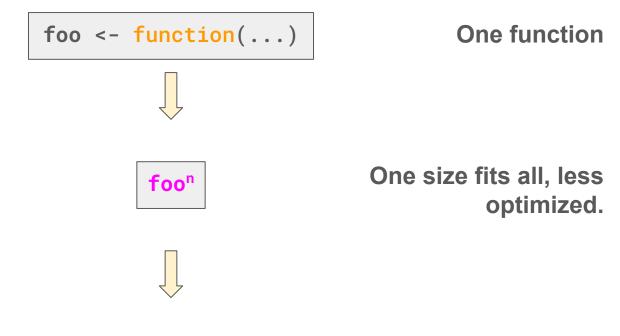
Many different versions of the same function.

Different users have different use-cases

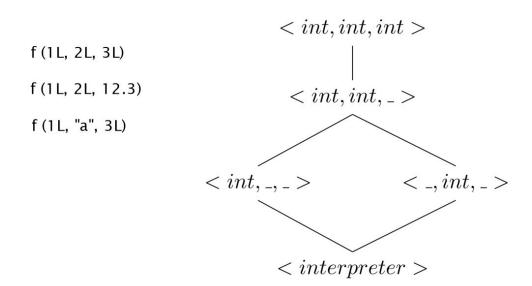


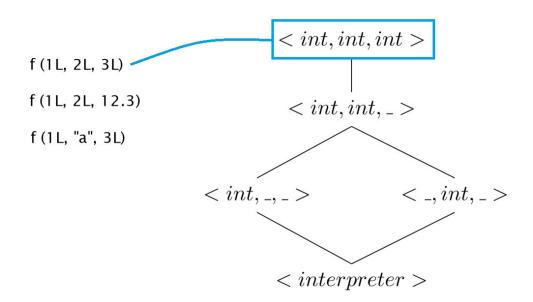
Reuse?

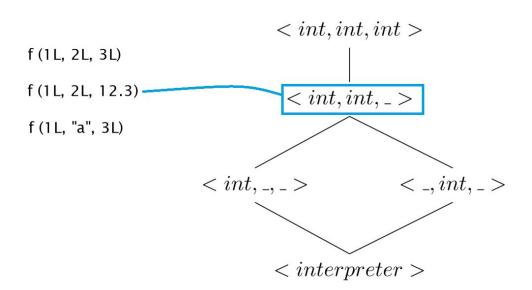
Motivation: Previous approaches

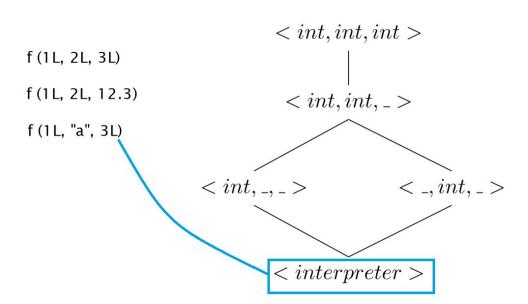


Lose peak performance









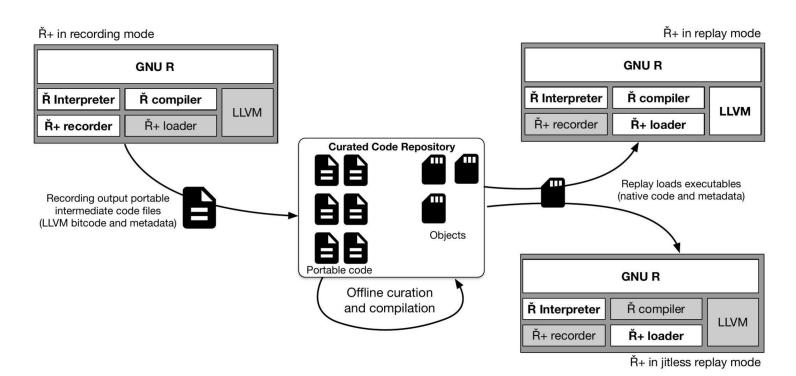
Challenges

1. Recording - How to save the compiled code?

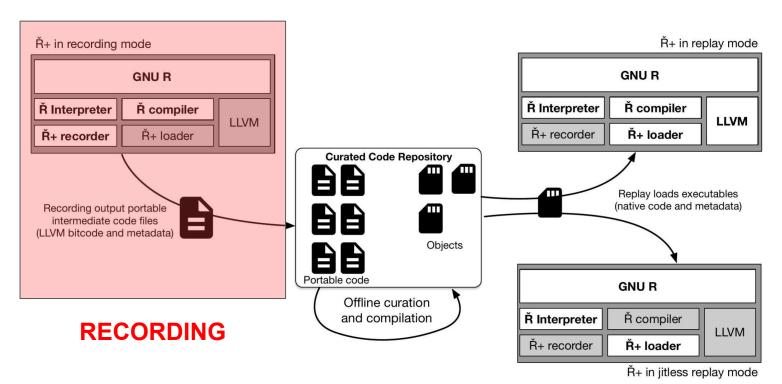
2. Duplicates - Identification and removal of duplicates

3. Reuse - Out of multiple options, which one to dispatch?

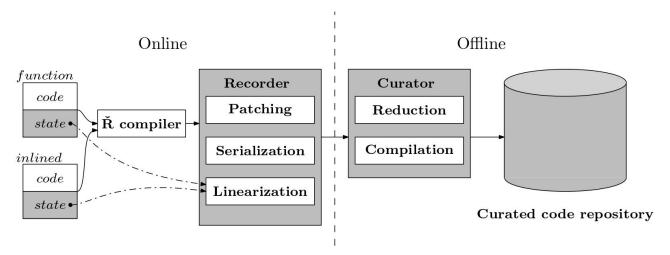
Our Approach: Ř+



Our Approach: Ř+



System Overview



Compiled code: *.bc (LLVM bitcode file)

Pool references: *.pool (encoded binary file)

Function metadata: *.meta (encoded binary file)

Speculative Contexts

Ř+ recording saves f' of a function f

```
f' = compiler(f)

What it looks like
```

Speculative Contexts

Ř+ recording saves f' of a function f

f' = compiler(⟨Code, ⟨C,F⟩⟩)

What it ic!:0

Speculative Contexts

Pair of $\langle \mathbf{C}, \mathbf{F} \rangle$

 $f(\Box, \Box, \Box)$

Set of predicates on the call-site arguments

Speculative Contexts

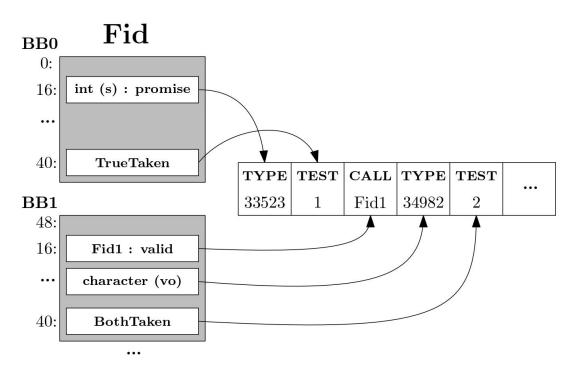
Pair of $\langle C, F \rangle$

F is a vector that holds

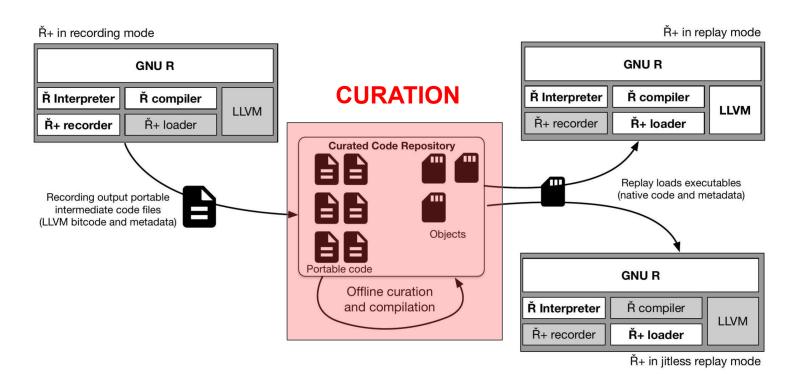
feedback context.

```
foo' <- function(x, y, z=TRUE) {</pre>
    res <- ns::add(res, g)
    res
m1 <- function(a, b) a + b
bar <- ns::add
foo(10, 20)
g <- matrix(...)
```

System Overview



Our Approach: Ř+



Curation

```
The complete compilation context contains:

Complete: <(int, int), <m1_inlined, ns::add_static, g_matrix)>
```

Curation

```
The complete compilation context contains:

Complete: \((int, int), \langle m1_inlined, ns::add_static, g_matrix\)\\
Useful(D): \((int, int), \langle m1_inlined, ns::add_static\)\\
```

We identify the useless part of the context using deopt points D.

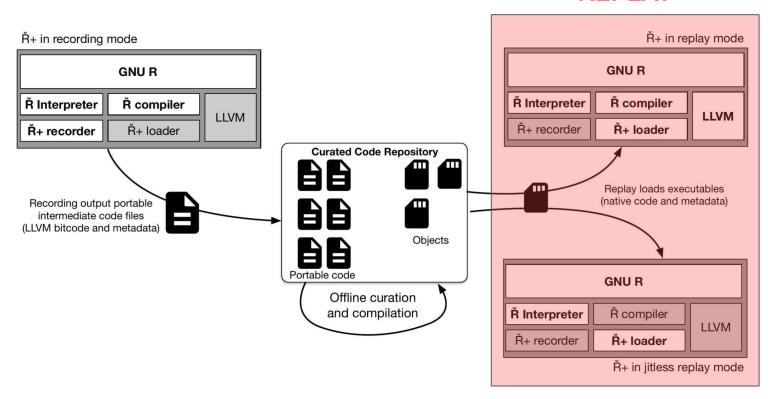
Curation

```
V': (C, D)

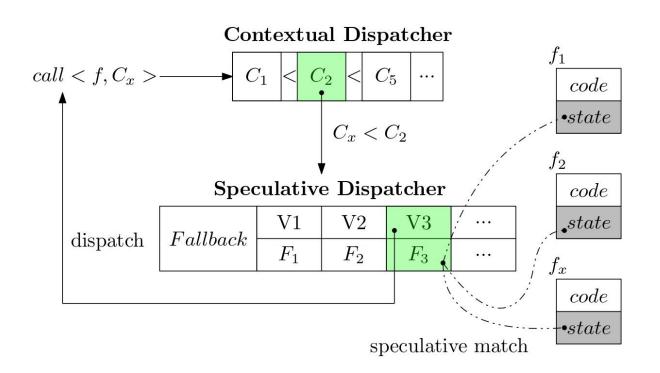
V": (C', D')
```

Our Approach: Ř+

REPLAY



The dispatcher



The dispatcher

```
Function* Dispatcher::dispatch() {
  if (cache != NULL) return cache;
  for (int i = length() - 1; i >= 0; i--) {
    auto f = getFunction(i);
    if (f->abled() && f->matchD())
      return cache = f;
  }
  return cache = getFallback();
}
```

Contextual dispatcher ensures correctness

Speculative dispatcher ensures precision

The dispatcher

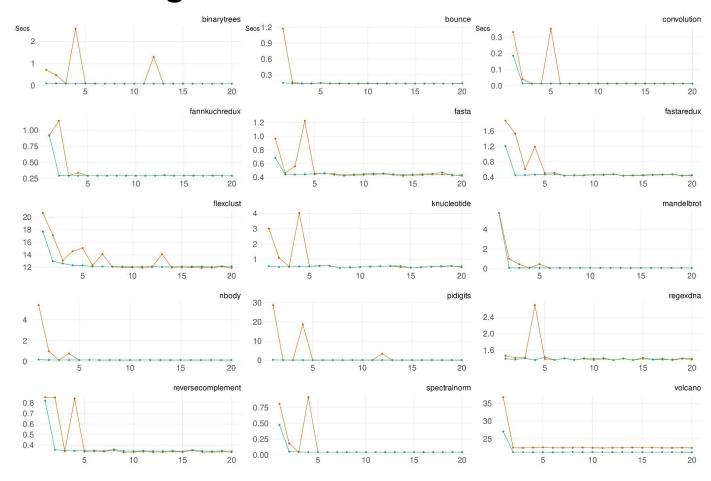
```
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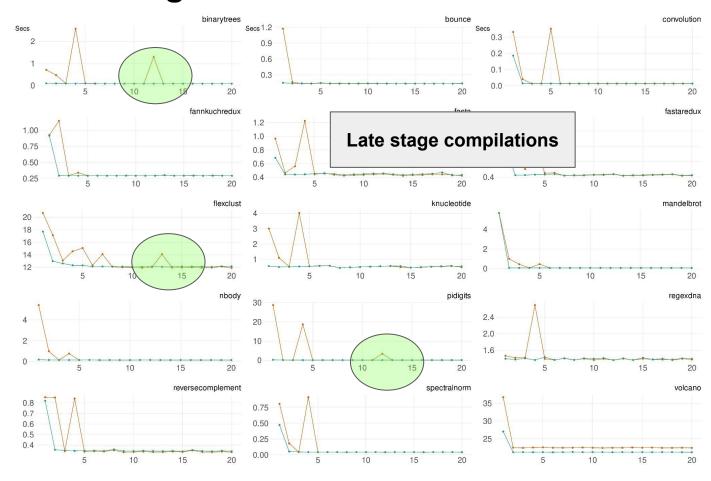
Expensive check needs to happen only when state changes :)

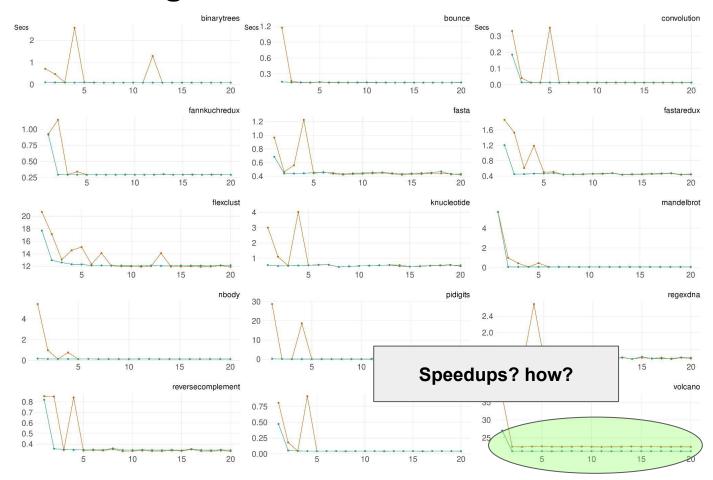
Otherwise we can keep a fast cache.

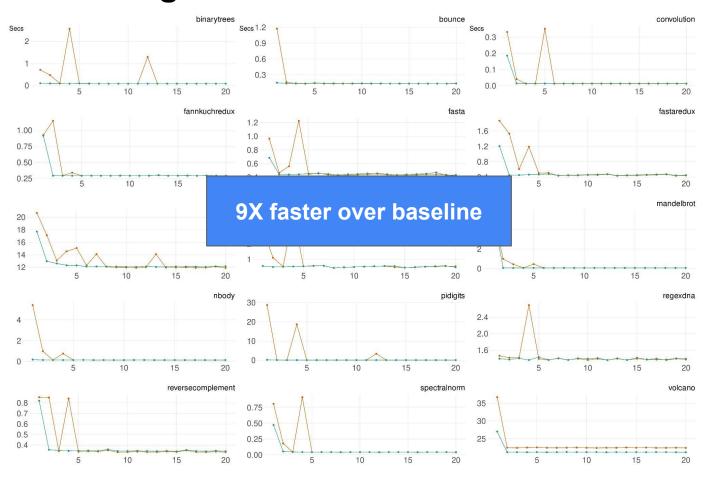
Contextual dispatcher ensures correctness

Speculative dispatcher ensures precision

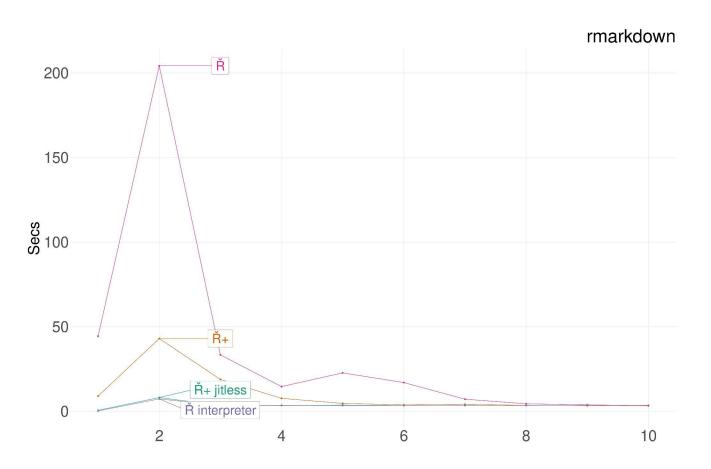




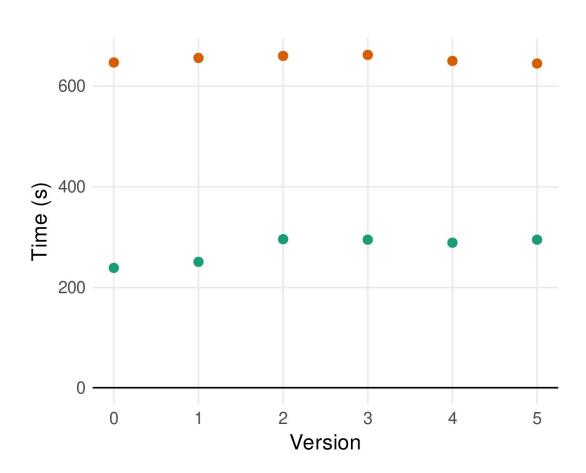




Real-World Use Case

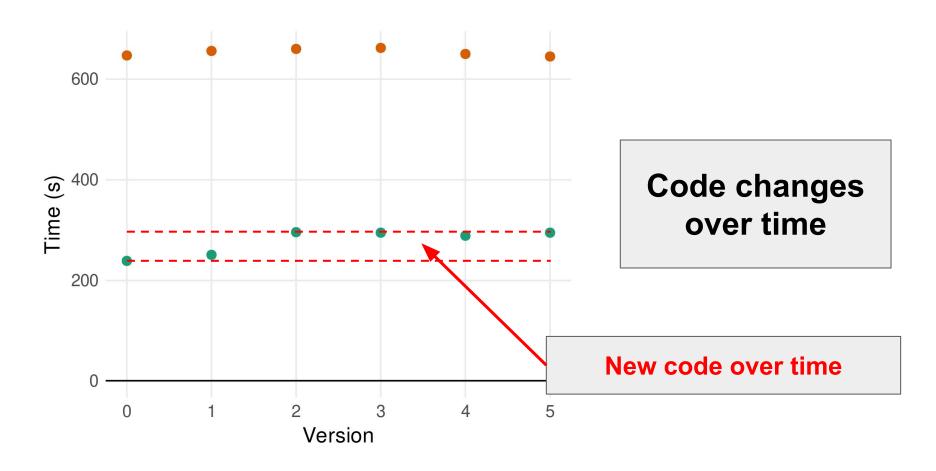


End-to-End Performance

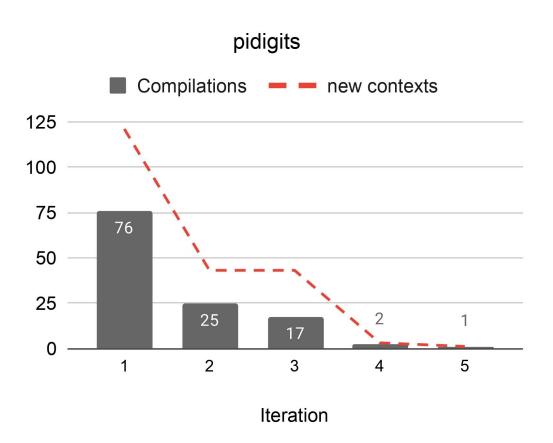


Code changes over time

End-to-End Performance



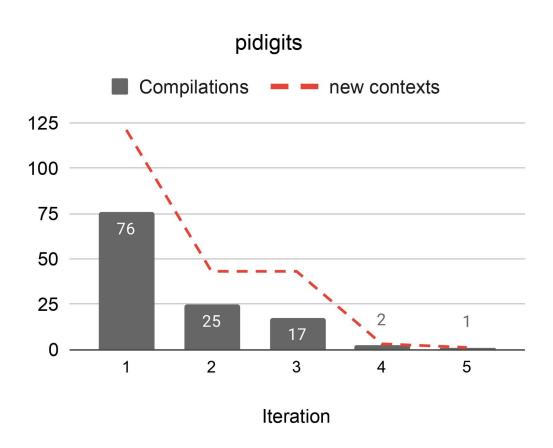
Repository Construction



Same program, Same inputs,

Iterative **loop**

Repository Construction

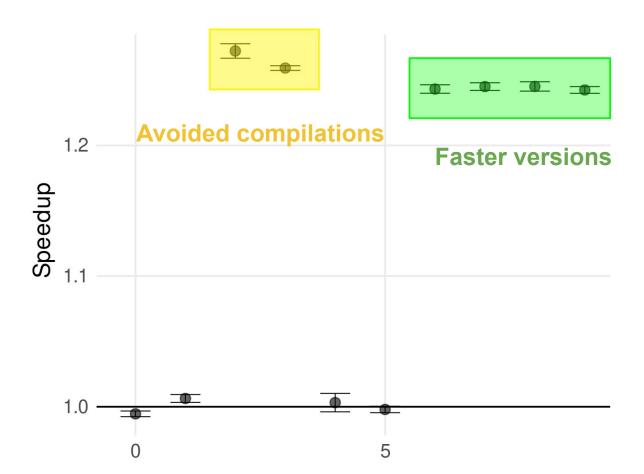


Same program, Same inputs,

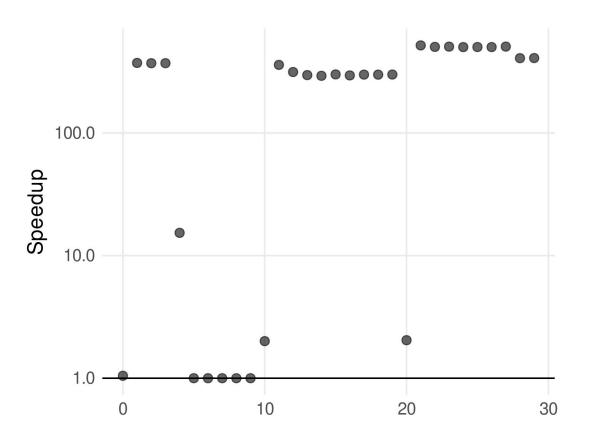
Iterative loop

Why still compilations?

Phase Change Behaviour



Phase Change Behaviour



Conclusion

Less compilations, fancy dispatcher, runs faster

- Speculative Dispatcher -> more complex operations?
- Global optimization of serialized code?
- More de-optimizations are good???
- Current work only focuses on maintaining JIT performance.
- But we show cases where we exceed it, this requires more careful exploration.

Thank you:)