

1. Write-read conflict

T1	T2
R(x)	
R(y)	
W(x)	
	R(x)
	R(y)
	W(y)
	R(x)
	R(y)
	W(x)
	R(z)
	W(z)

T1 writing x in before T2 writing x may cause a dirty read. Since T1 never commit its updates, rollback T1 would cause both T2 reading x give us a false result that does not match with The actual value written by T2

2. Read-write conflict

T1	T2
	R(x)
	R(y)
	W(y)
R(x)	
R(y)	
W(x)	
	R(x)
	R(y)
	W(x)
	R(z)
	W(z)

This might also cause a read-write conflict since T2 read x twice while T1 updated x in between. The same read query might result in different value, causing an unrepeatable read

3. Write-write conflict

T1	T2
	R(x)
	R(y)
	W(y)
	R(x)
R(x)	
R(y)	
W(x)	
	R(y)
	W(x)
	R(z)
	W(z)

The T1 write x is overwritten by T2 write x happens after T1 actions, causing a lose update.