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11. For $(n)^{(1/2)} + n$, n is of higher order, so it will take precedence as the function increases. Because the function $(n \log n)$ increases at a greater rate than the function (n) as n grows larger, it is impossible to have any C or n_0 value that would make $n \geq C(n \log n)$.

12. For $(100n^3 + n^2)$, $100n^3$ is of higher order, so it will take precedence as the function increases. If we let $C = 102$, and let n_0 be any number ≥ 1 , for example $n_0 = 1$. We can see the following:

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100(1)^3 + (1)^2 <= (102)(1)^3
100 + 1 <= 102
101 <= 102, which holds true

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At the same time, if we let $C = 100$, and let n_0 be any number ≥ 1 , for example $n_0 = 1$. We can see the following:

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100(1)^3 + (1)^2 >= (100)(1)^3
100 + 1 >= 100
101 >= 100, which holds true

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Therefore we can say that $(100n^3 + n^2)$ is both big O and big ω of (n^3) . The functions are of the same order, therefore it is true that $(100n^3 + n^2)$ is big θ of (n^3) (there is both an upper and lower bound).

13. If we simplify the left function we get:

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(1/n) + (n^2/n)
n^(-1) + n, because n is of higher order than n^(-1), n takes precedence as
the function increases

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If we let $C = 3$, and let n_0 be some number ≥ 1 , for example $n_0 = 1$, we see the following:

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(1/(1)) + ((1)^2/(1)) <= 3((1))
1 + 1 <= 3
2 <= 3, which holds true

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At the same time, if we let $C = 1$, and let n_0 be some number ≥ 1 , for example $n_0 = 1$, we see that:

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(1/(1)) + ((1)^2/(1)) >= 1((1))
1 + 1 >= 1
2 >= 1, which holds true

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Therefore we can say that $(1/n + n^2/n)$ is both big O and big ω of (n) . The functions are of the same order, so it holds true that $(1/n + n^2/n)$ is big θ of (n) (there is both an upper and lower bound).

14. For $(1/n^{100}) + \log 16$, we can rewrite this as $n^{(-100)} + (n^0) \log 16$. Because $(n^0) \log 16$ is of higher order than $n^{(-100)}$, it takes precedence as the function increases. If we let $C = (\log 16 + 2)$, and let n_0 be some number ≥ 1 , for example $n_0 = 1$, we can see that:

$1/(1)^{100} + \log 16 \leq (\log 16 + 2)(1)$
 $1 + \log 16 \leq \log 16 + 2$, which holds true

At the same time, if we let $C = \log 16$, and let n_0 be some number ≥ 1 , for example $n_0 = 1$, we can see that:

$1/(1)^{100} + \log 16 \geq (\log 16)(1)$
 $1 + \log 16 \geq \log 16$, which holds true

Therefore, we can say that $(1/n^{100} + \log 16)$ is both big O and big omega of (1) . Because both functions are of the same order, it is true that $(1/n^{100} + \log 16)$ is big theta of (1) (there is both an upper and lower bound).

15. For the left function $(n + \log(n^2))$ n grows at a faster rate and takes precedence as the function increases. If we let $C = 1$, and let n_0 equal some number ≥ 1 , for example $n_0 = 1$, we find the following:

$(1) + \log(1)^2 \leq (1)(5(1))$
 $1 + 0 \leq 5$
 $1 \leq 5$, which holds true

At the same time, if we let $C = 1/10$, and let n_0 equal some number ≥ 1 , for example $n_0 = 1$, we find the following:

$(1) + \log(1)^2 \geq (1/10)(5(1))$
 $1 + 0 \geq 1/2$
 $1 \geq 1/2$, which holds true

Therefore, we can say that it is true that $(n + \log n^2)$ is both big O and big omega of $(5n)$. The functions are of the same order, so therefore $(n + \log n^2)$ is big theta of $(5n)$ (there is both an upper and lower bound).