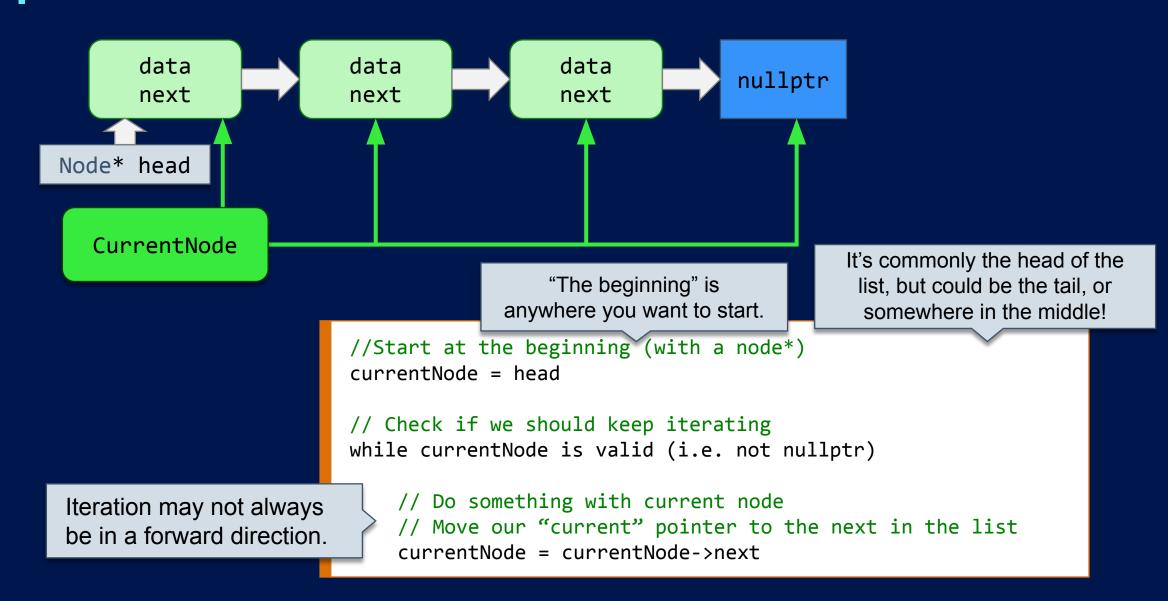
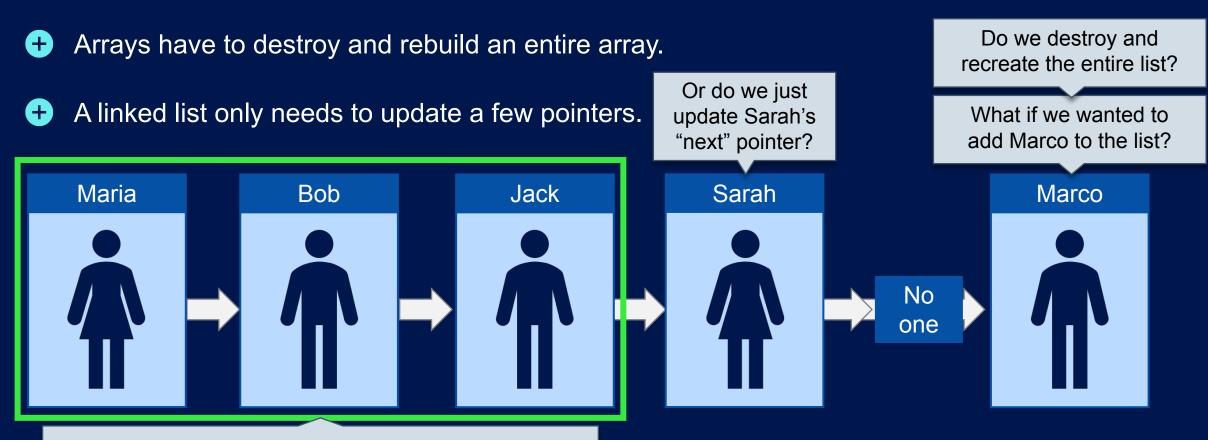


Back to Our List of Numbers



Inserting a Node Into an Existing List

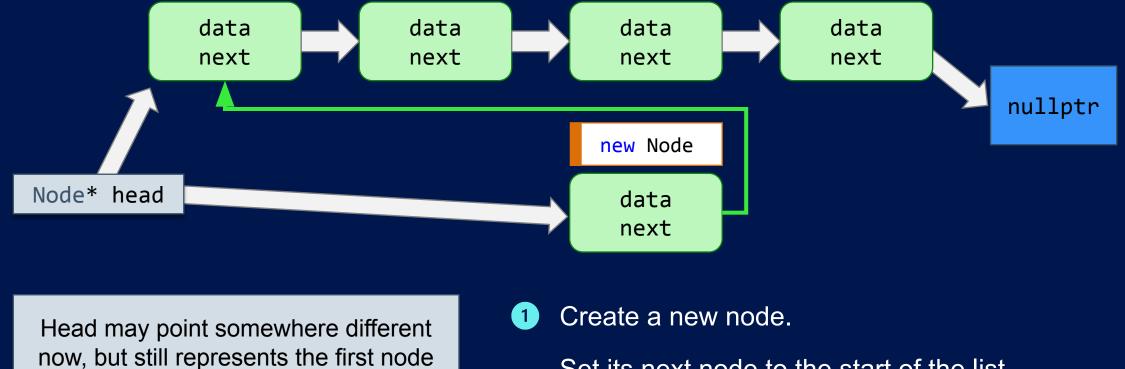
4 A big advantage a linked list has over arrays is the ability to add and remove nodes quickly.



These "nodes" have no idea the list was changed.

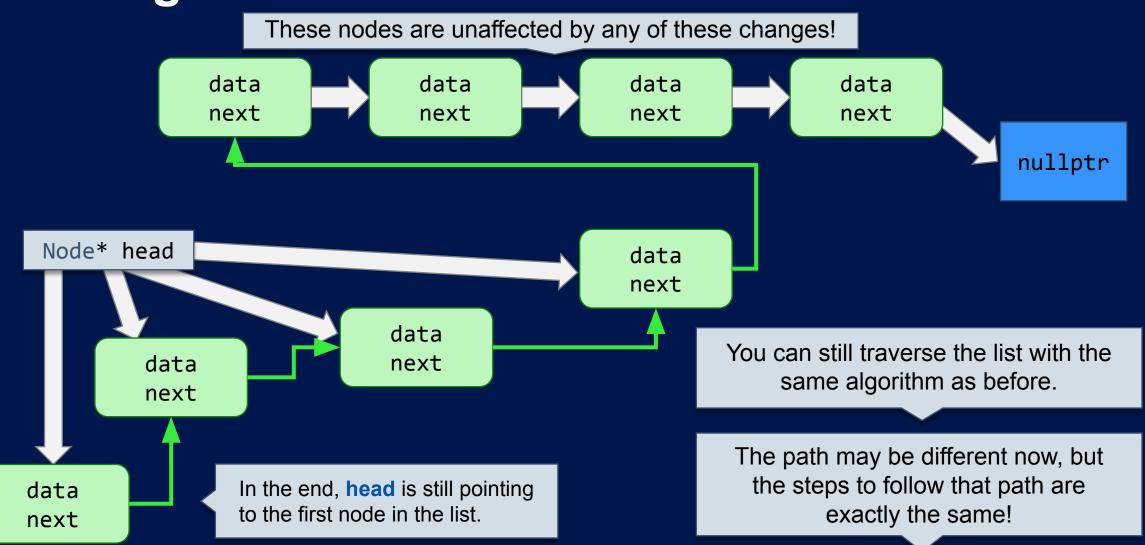
Adding to the Front of the List

in the list.



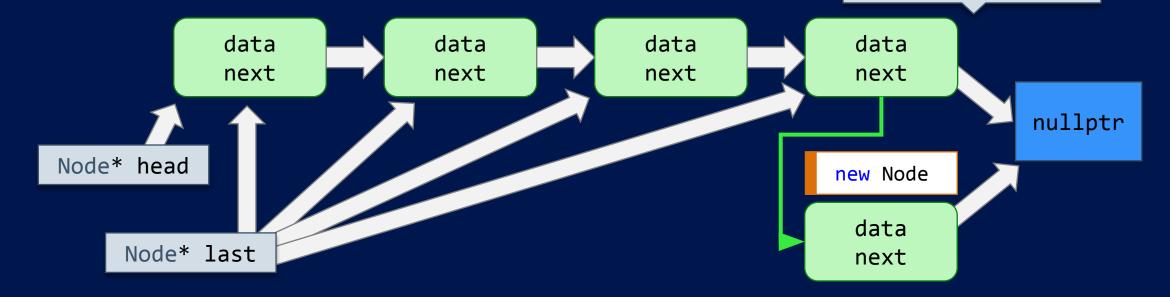
- Set its next node to the start of the list (i.e., the head)
 - Set the current head pointer to the new node.

Adding More Nodes to the Front



Adding More Nodes to the Front

This is the last node, or the "tail" of the list.



```
// Slightly different traversal
currentNode = head

// Keep going until we find node with a "next"
// pointer that is nullptr—that's the last
while currentNode->next is valid
    currentNode = currentNode->next
```

- + Similar to adding to the front of the list:
 - 1 Create a new node.
 - 2 Set the next pointer of the "tail" to this new node.
 - If you have a tail pointer, set the tail pointer itself to this new node (not applicable in this example).

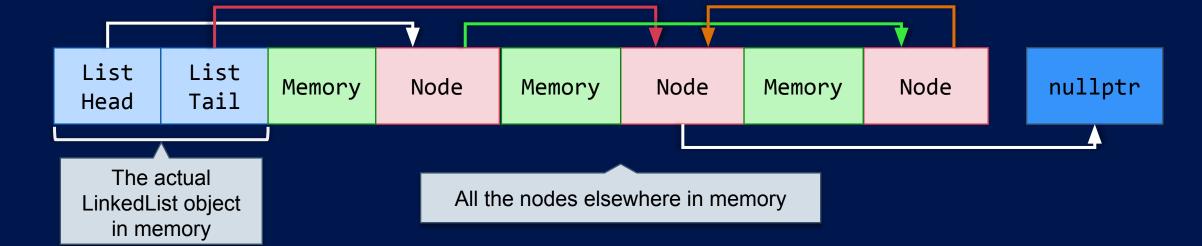
Storing Head and Tail Pointers

A basic Linked List implementation might only store the first node:

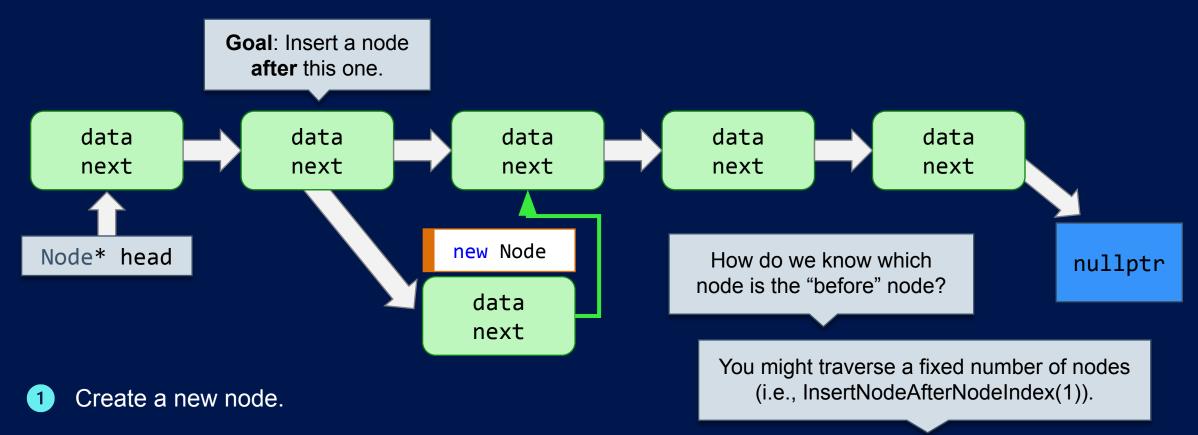
+ You could also store a tail pointer.

```
class LinkedList
{
    Node* head;
};
```

```
class LinkedList
{
    Node* head;
    Node* tail;
};
```



Inserting a Node in Between Two Nodes

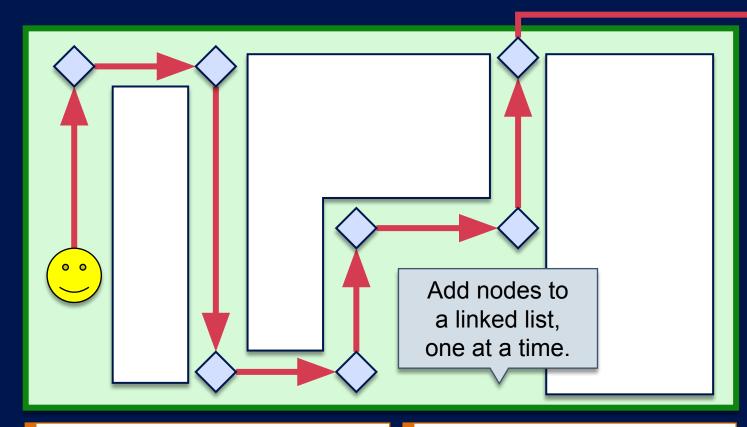


- 2 Set its "next" pointer to the node that will ultimately follow it.
- 3 Set the "previous" node's next pointer to the new node.

You might traverse and search for a node with a specific value (i.e., InsertNodeAfterValue(200)).

nullptr

Use Case: Movement / pathfinding



Last waypoint has no next.

Player Movement Algorithm:

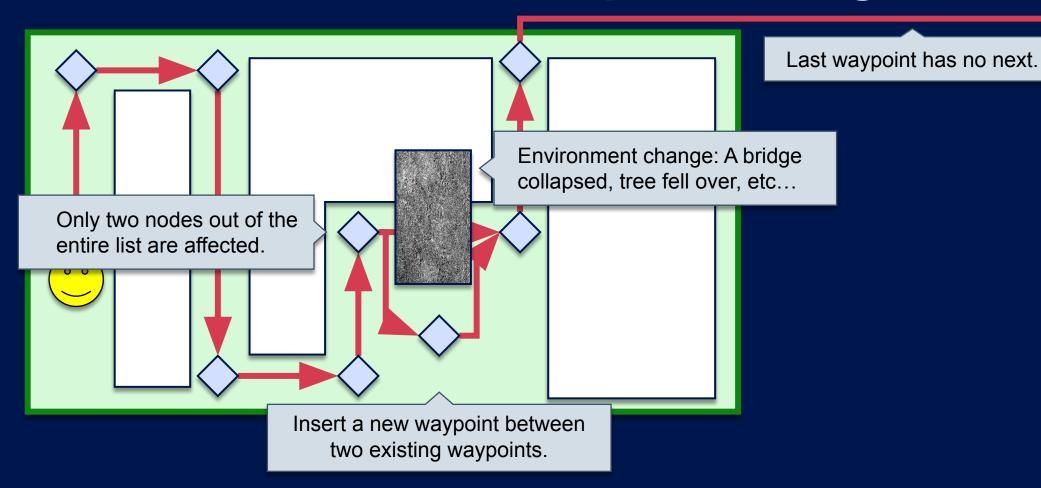
If destination is not null and player is not at destination, take a step toward destination.

Then,

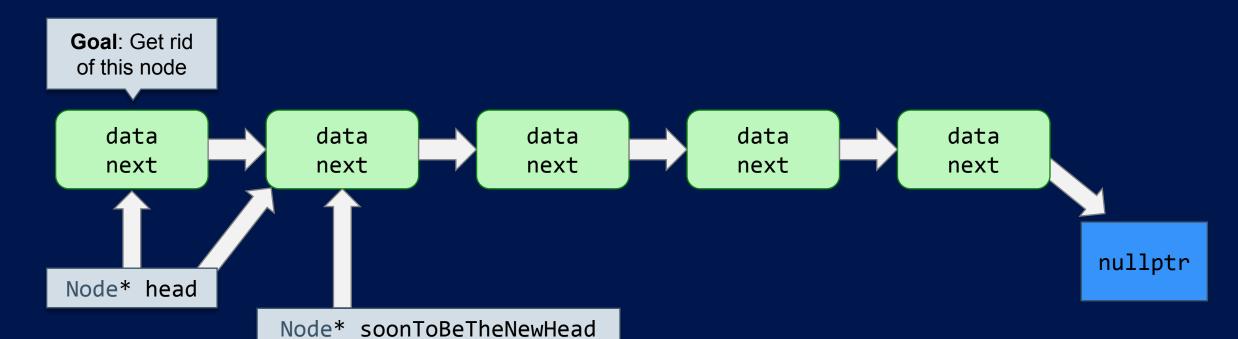
if player is at destination destination = destination->next.

What if the environment changed and blocked part of the path?

Use Case: Movement / pathfinding



Removing the Head Node

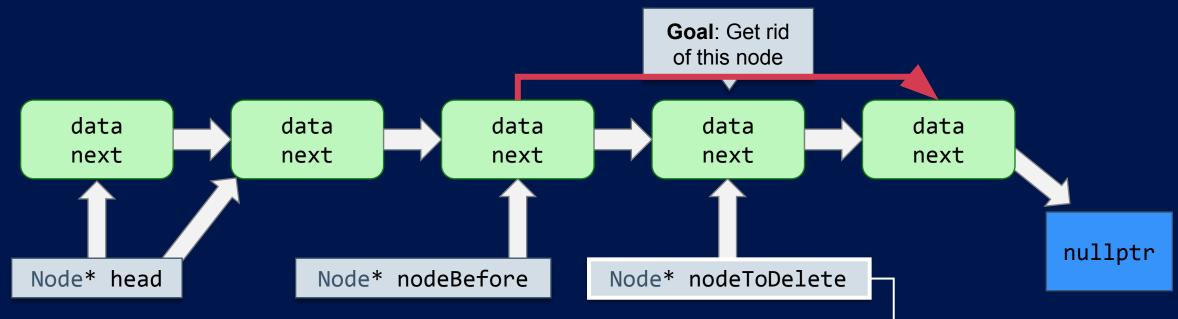


- 1 Create a pointer to the node after the head.
- 2 Delete the head node.
- 3 Set the head node to the pointer you created.

Why not just call delete Head and be done?

Once head is deleted, any data at that location (including the location of head->next) is lost!

Removing a Node



- Get a pointer to the node before the one to delete.
- Set its "next" pointer to the node that will ultimately follow it (the one **after** the deleted node).

```
nodeBefore->next = nodeToDelete->next; // OR...
nodeBefore->next = nodeBefore->next->next;
```

3 Delete the node.

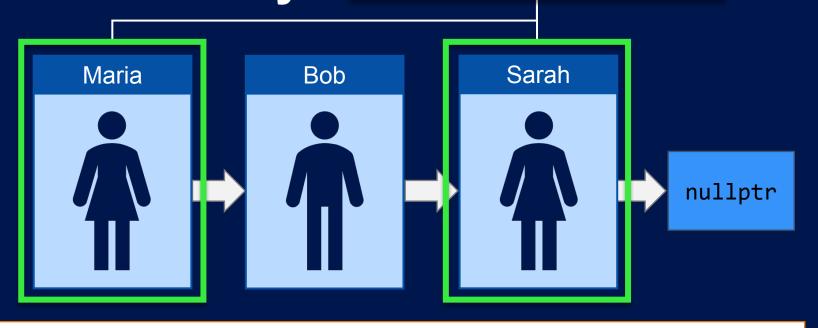
If this is a temporary variable declared in a function somewhere, it will fall out of scope eventually.

Just don't try to use it after deleting what it pointed to!

Checking for Node Validity!

What if we wanted to visit **every other** node, starting with the first?

```
struct PersonNode
{
    Person thePerson;
    PersonNode* next;
};
```



```
PersonNode* p = someList.GetFirst(); // Start at the beginning
p->thePerson.SomeOperation();

p = p->next->next;
p->next == Bob Node
(Bob Node)->next == Sarah Node

p = p->next->next;
From Sarah, p->next is nullptr

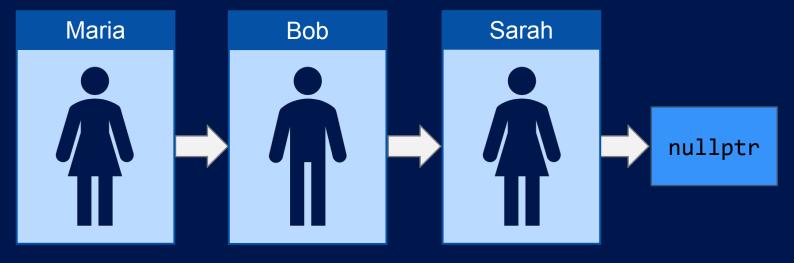
nullptr->next will likely crash your program
```

You can see this is a bad idea, looking at a diagram.

Your program doesn't have a diagram, and can't see "obviously bad" decisions.

Checking for Node Validity!

```
struct PersonNode
{
    Person thePerson;
    PersonNode* next;
};
```



```
PersonNode* p = someList.GetFirst();
if (p != nullptr)
   p->thePerson.SomeOperation();

if (p->next != nullptr) // IF someone follows me...
{
   p = p->next->next; // Move to who follows THEM

   if (p != nullptr) // If we actually have a valid pointer...
        p->thePerson.SomeOperation(); // THEN, FINALLY do something
}
```

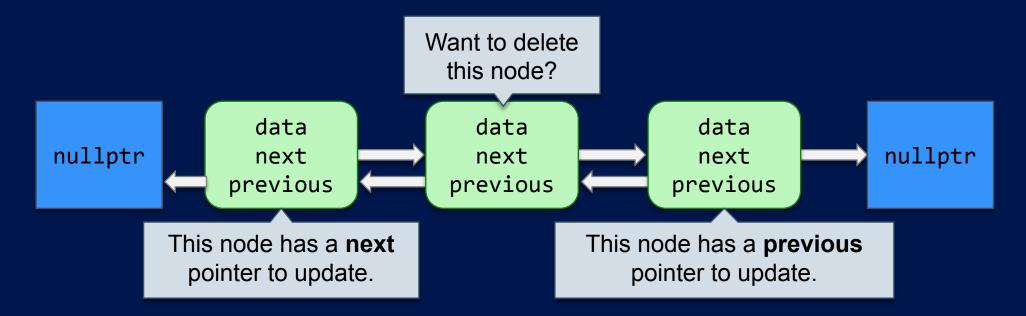
Don't make assumptions on the structure of the entire list.

Handle each node one at a time, and check to make sure you have valid pointers!

Changes With a Doubly-Linked List

- The overall algorithms for any operation is largely the same.
- You have to account for a "previous" pointer in nodes.

 Initialize them—to nullptr if nothing else.
- Some algorithms may be easier to implement if you have previous pointers.



| Edge Cases

- Hany algorithms in our code work correctly in almost all cases.
- **Edge cases** are the special circumstances that need specific code.
- Common edge cases:
 - An empty list (or a full one!)
 - An operation on the first or last node
 - If a value is at some minimum (often zero), some maximum, or nullptr.

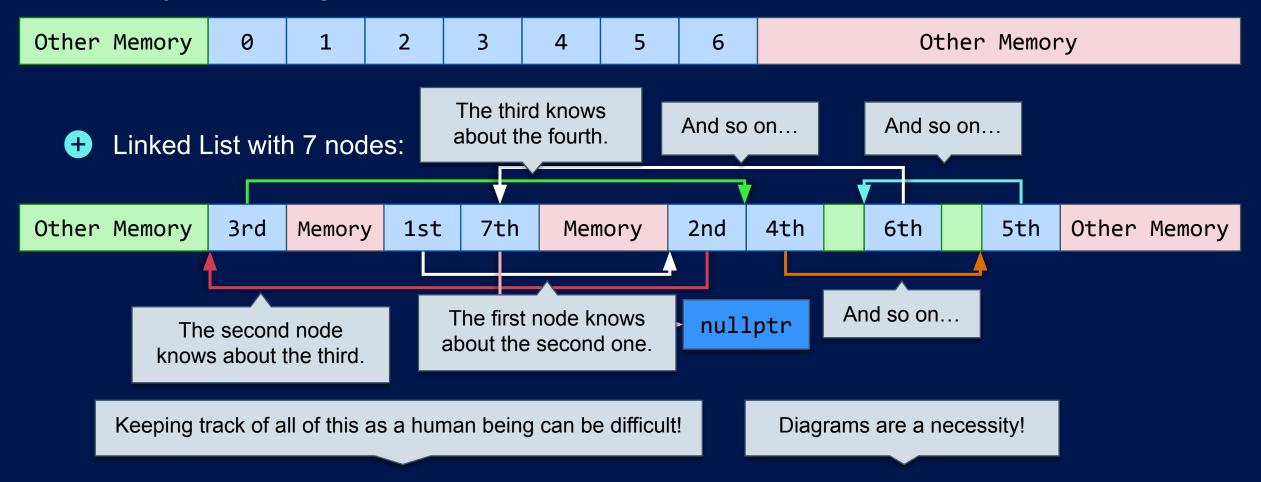
```
// Example: RemoveNode(Node* target)
if (target == head)
    // Edge case to remove the first node
else if (target == tail)
    // Edge case to remove the last node
else
    // Code for any "in-the-middle" nodes

// Any "universal" code that should always execute
```

Many algorithms will have a setup kind of like this.

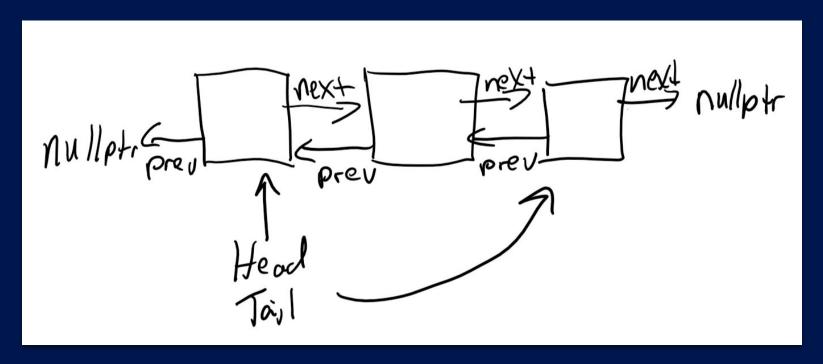
Visualizing Memory Can Be a Challenge

Array with 7 contiguous elements:



Draw Diagrams to Plan and Debug!

- You don't have to be "good enough" to remember all of it.
- Draw a diagram to represent the current state of your list.



- Use a diagram like this with every algorithm
- Walk through your code,updating this diagram each step of the way.
- If your code has a bug
 in it, you'll "break"
 your diagram.

Other Operations Beyond These

- We've looked at the most common operations.
- You could create more, depending on what your program needs:
 - Combine two lists of nodes.
 - Split a list into multiple parts.
 - Remove every other node.
 - Insert <X> nodes after/before an existing node.
- + These "advanced" operations can utilize existing functionality.
 - └── Set up the basic operations first, then build upon them.



Conclusion



Placeholder for the instructor's welcome message. Video team, please insert the instructor's video here.

