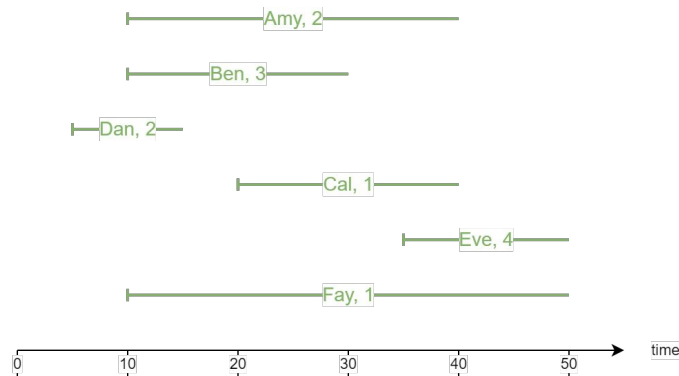


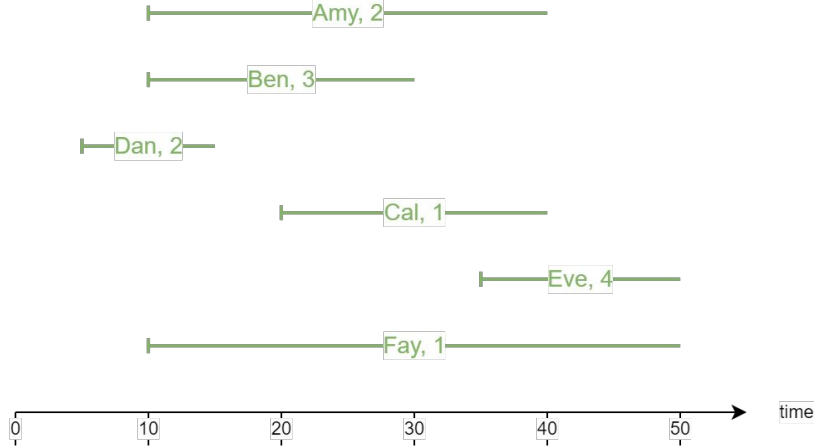
SB-Tree

Incremental computation and maintenance of temporal aggregates

Mirko Bristle, Markus Eggimann

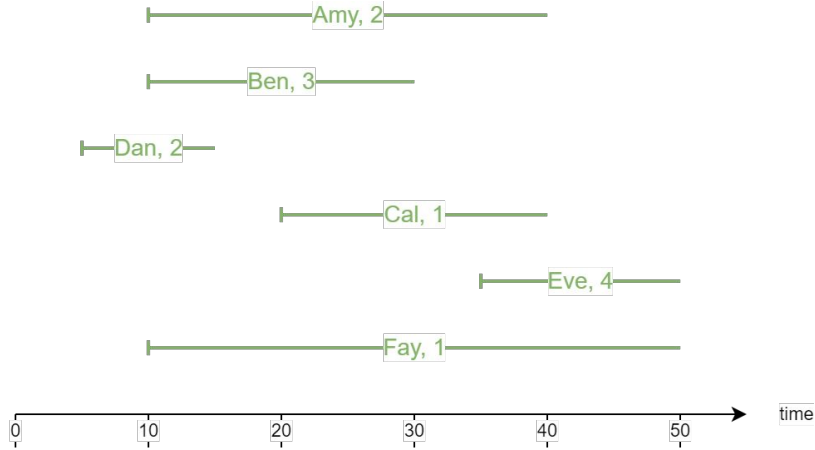


Context I: Temporal data



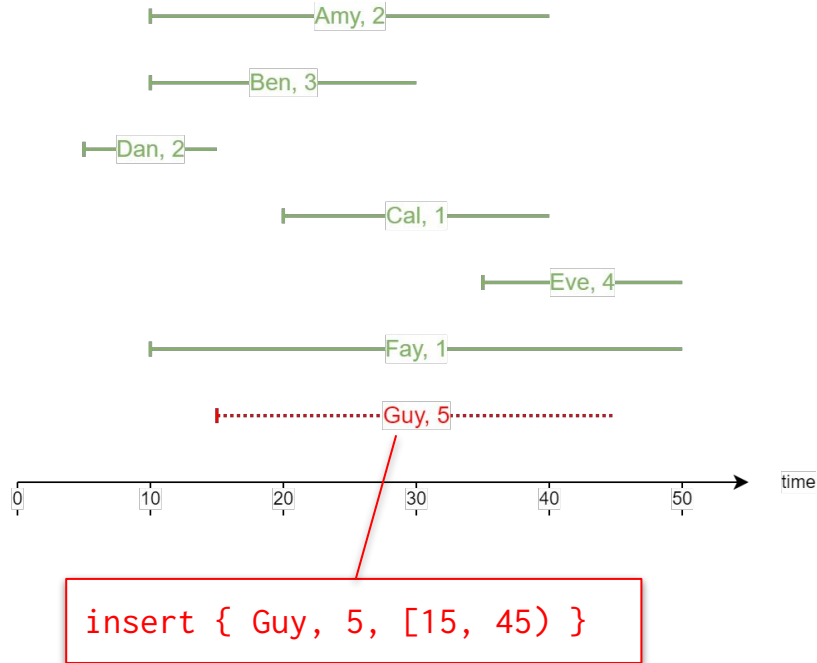
Patient	Dosage	Interval
Amy	2	[10, 40)
Ben	3	[10, 30)
Cal	1	[20, 40)
Dan	2	[5, 15)
Eve	4	[35, 45)
Fay	1	[10, 50)

Context II: Temporal aggregates



Sum dosage	Interval
0	[0, 5)
2	[5, 10)
8	[10, 15)
6	[15, 20)
7	[20, 30)
4	[30, 35)
8	[35, 40)
5	[40, 45)
1	[45, 50)
0	[50, ...)

The problem



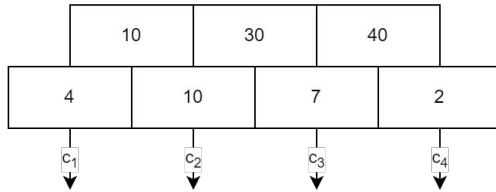
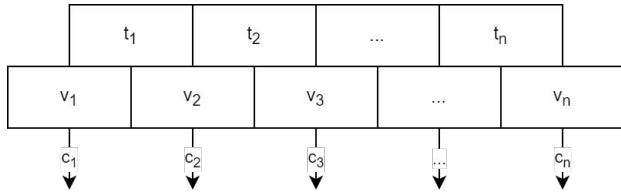
Sum dosage	Interval
0	[0, 5)
2	[5, 10)
8	[10, 15)
11	[15, 20)
12	[20, 30)
9	[30, 35)
13	[35, 40)
10	[40, 45)
1	[45, 50)
0	[50, ...)

A solution

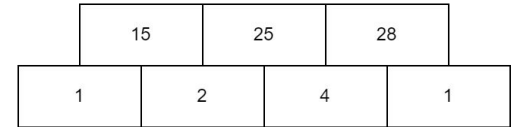
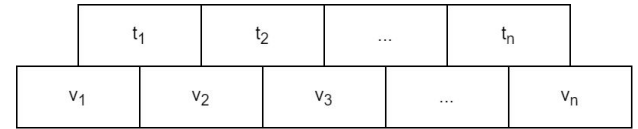
- Use a tree-based index for each aggregate
- Use a combination of B-trees and segment trees
- Allow fast lookups and efficient updates

→ Introducing the **SB-tree**!

SB-tree: Building blocks

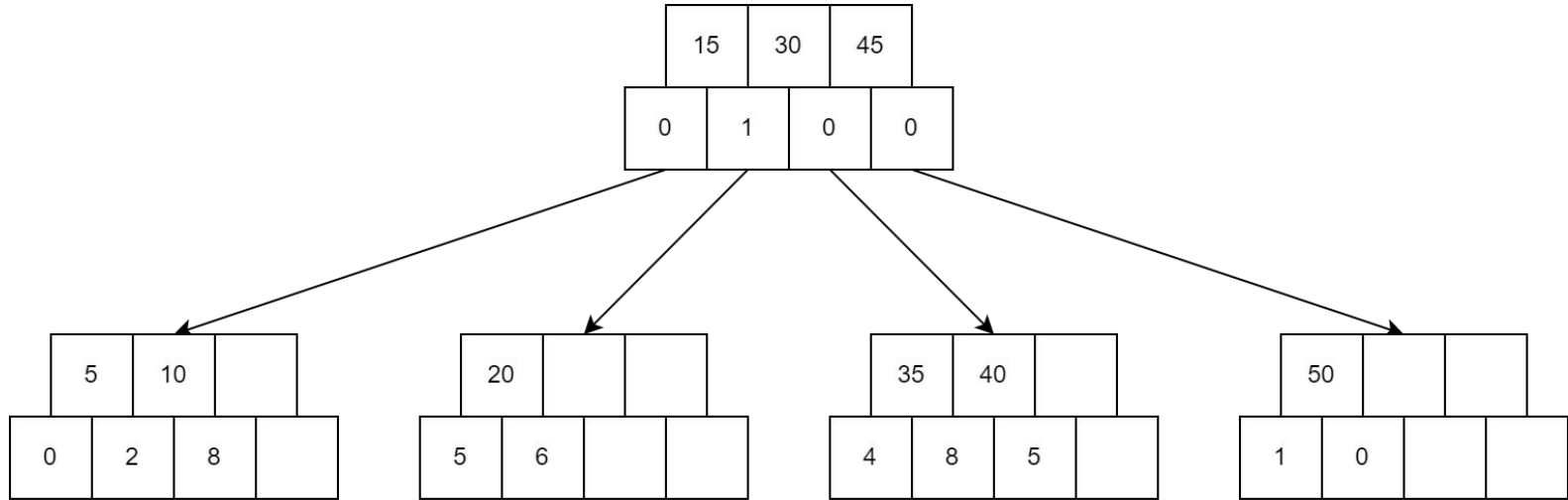


Interior node



Leaf node

SB-tree: The complete structure



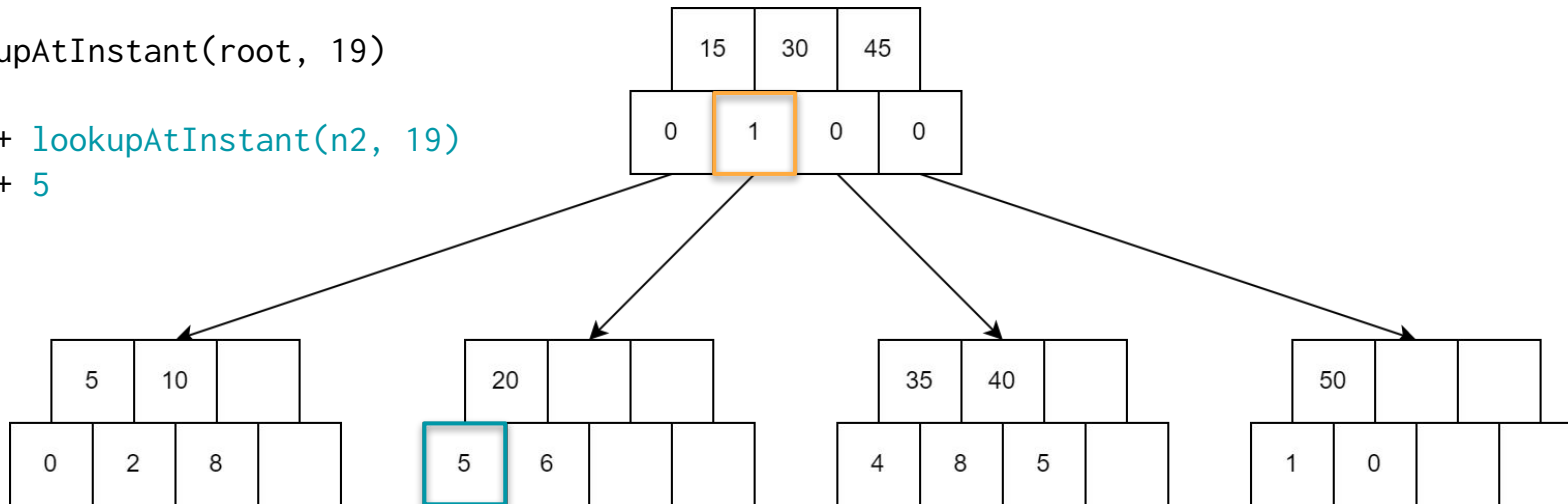
SB-tree: Lookup

lookupAtInstant(root, 19)

= 1 + lookupAtInstant(n2, 19)

= 1 + 5

= 6



SB-tree: Range query

rangeQuery([14, 28), 0)

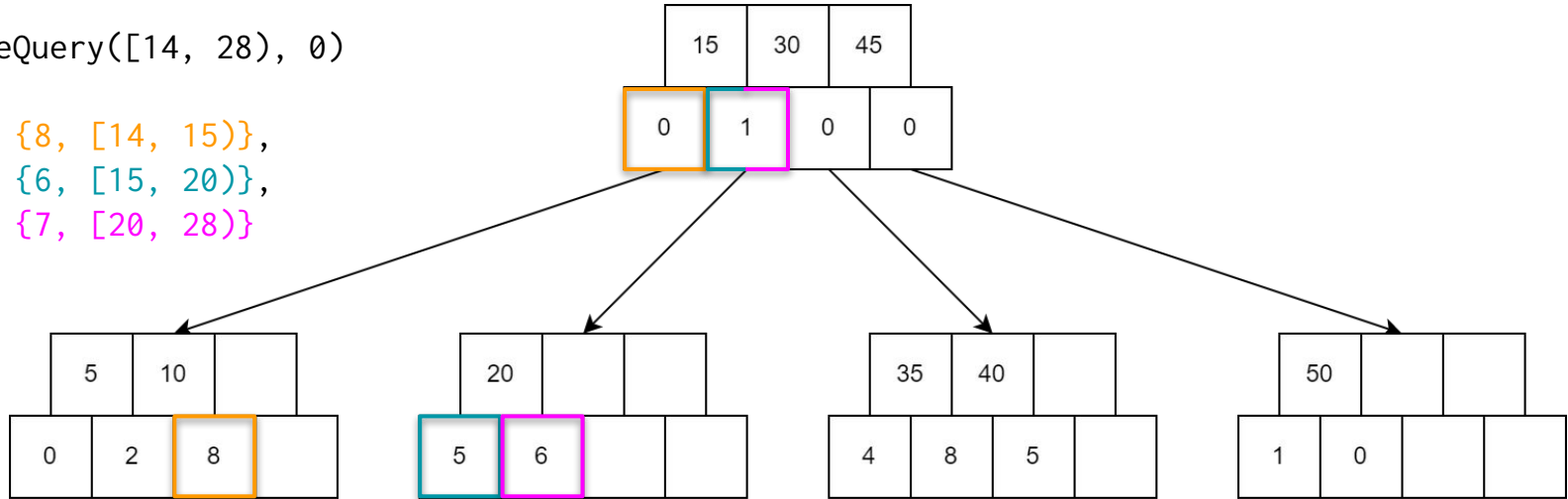
= [

{8, [14, 15)},

{6, [15, 20)},

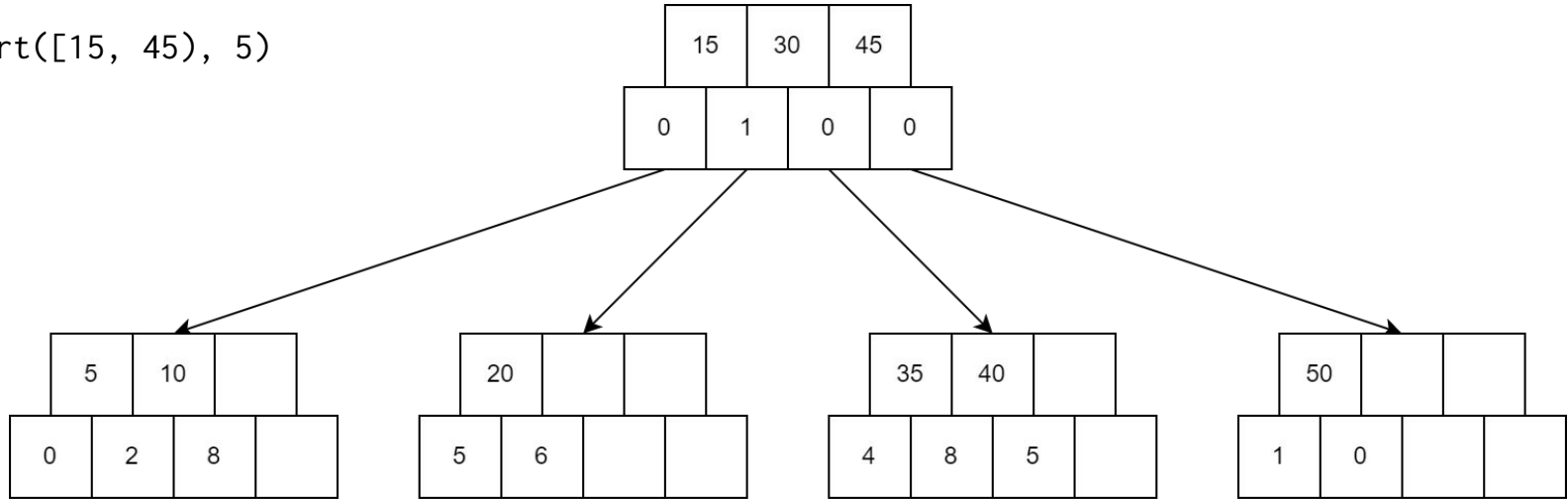
{7, [20, 28)}

]



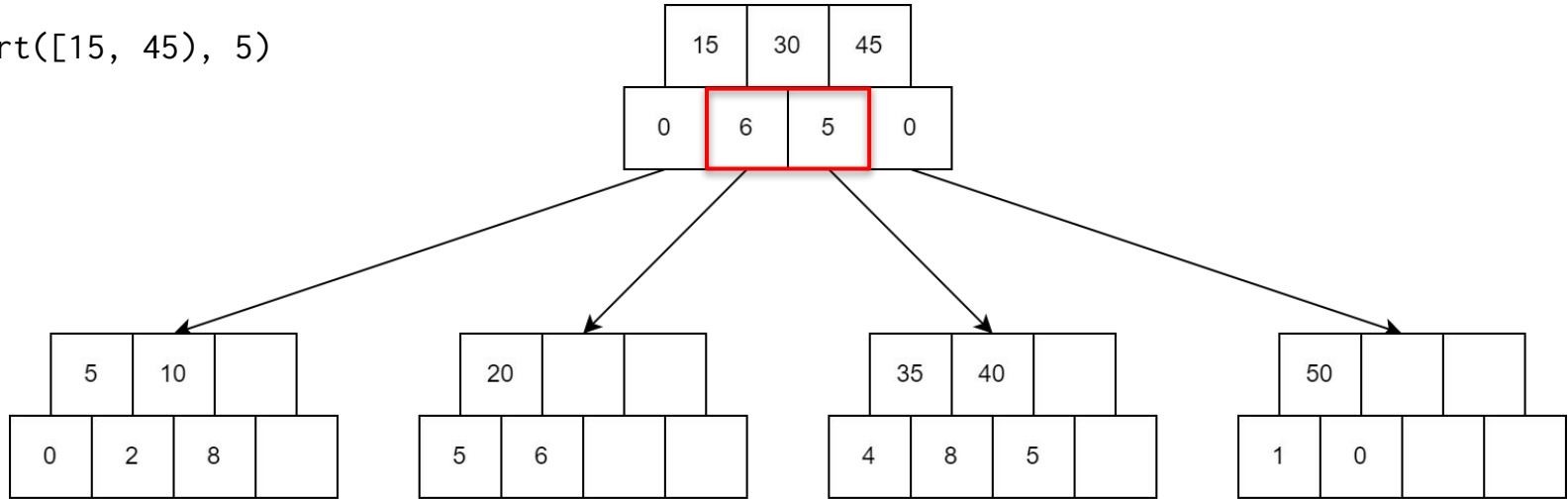
SB-tree: Insert I

insert([15, 45), 5)



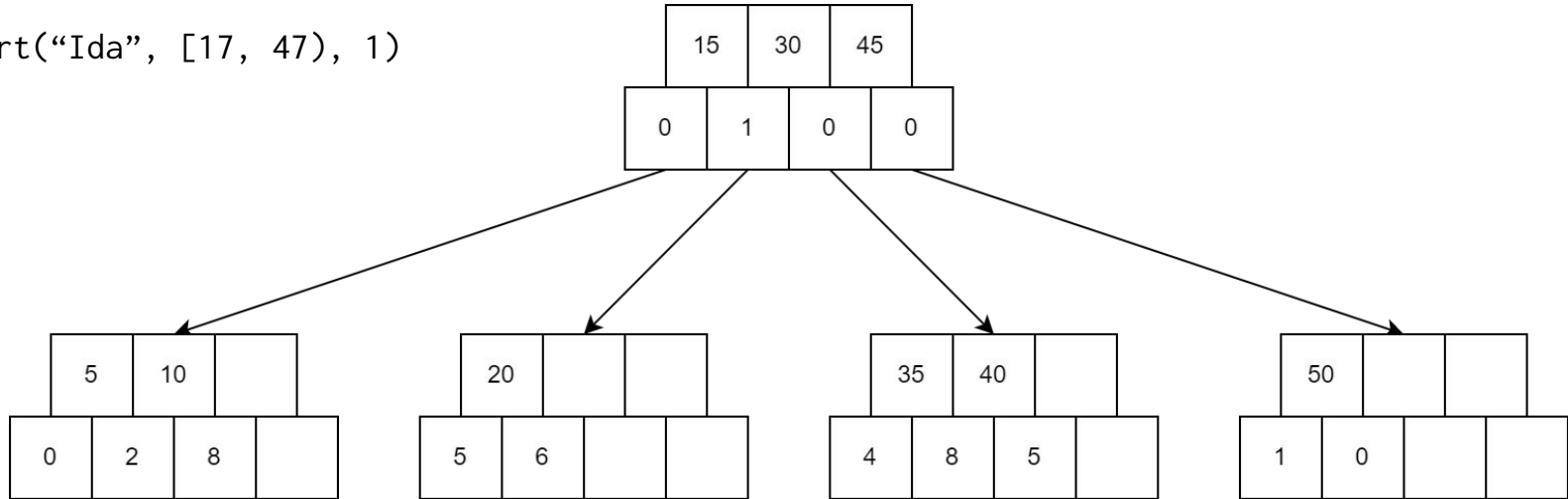
SB-tree: Insert I

insert([15, 45), 5)



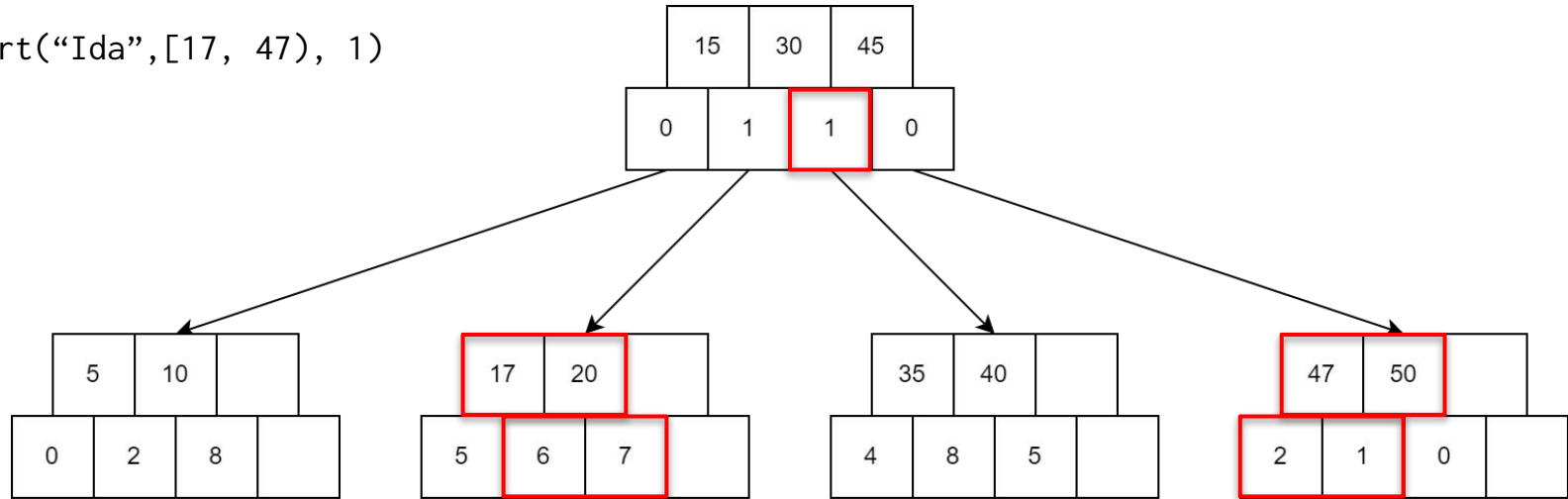
SB-tree: Insert II

insert("Ida", [17, 47], 1)



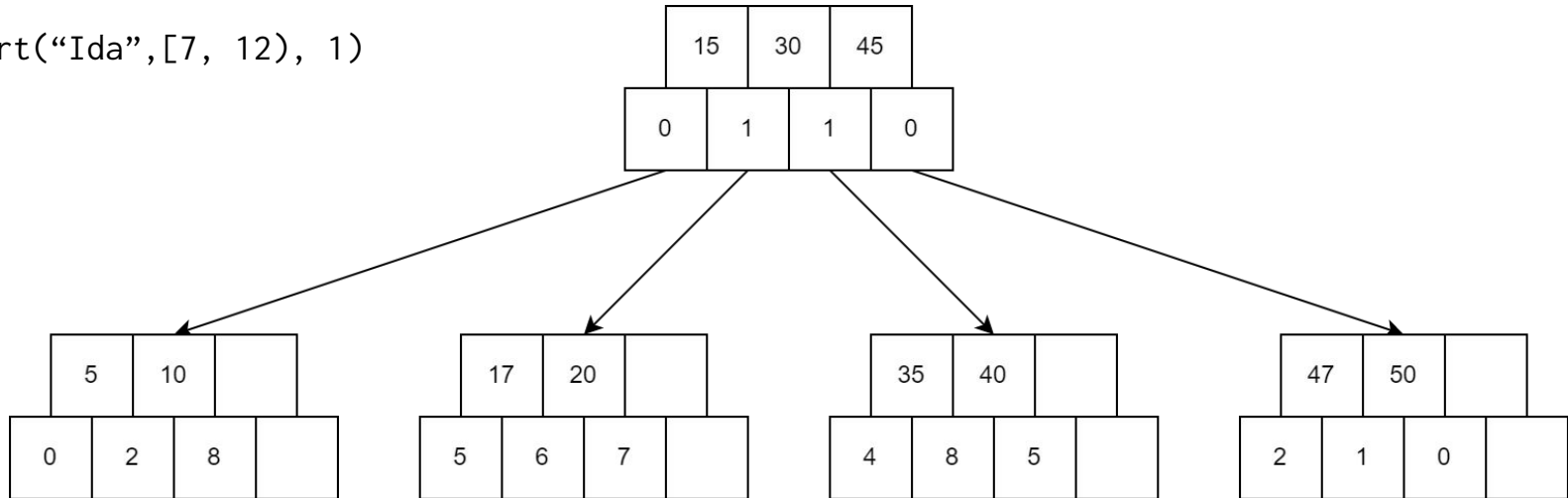
SB-tree: Insert II

insert("Ida",[17, 47), 1)



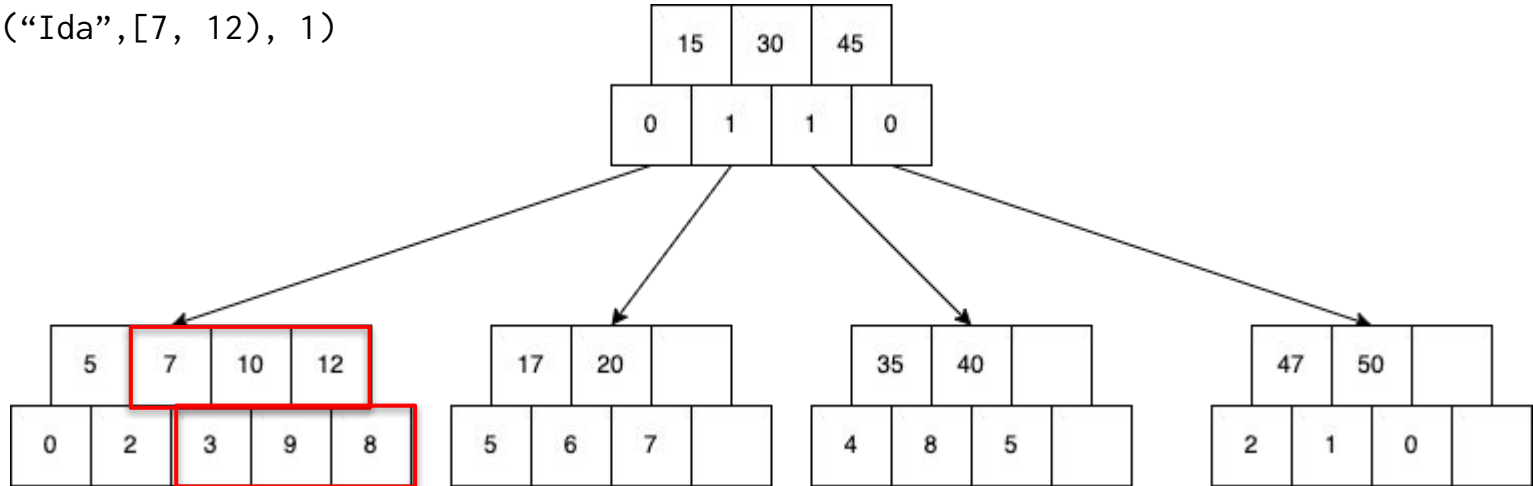
SB-tree: Node splits

insert("Ida",[7, 12], 1)



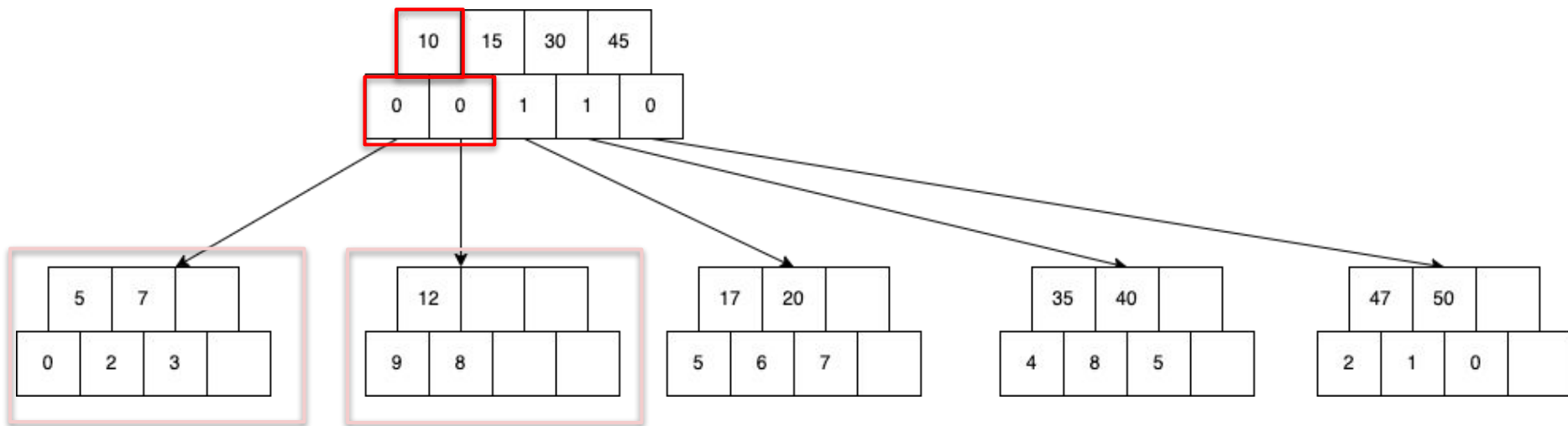
SB-tree: Node splits

insert("Ida",[7, 12], 1)



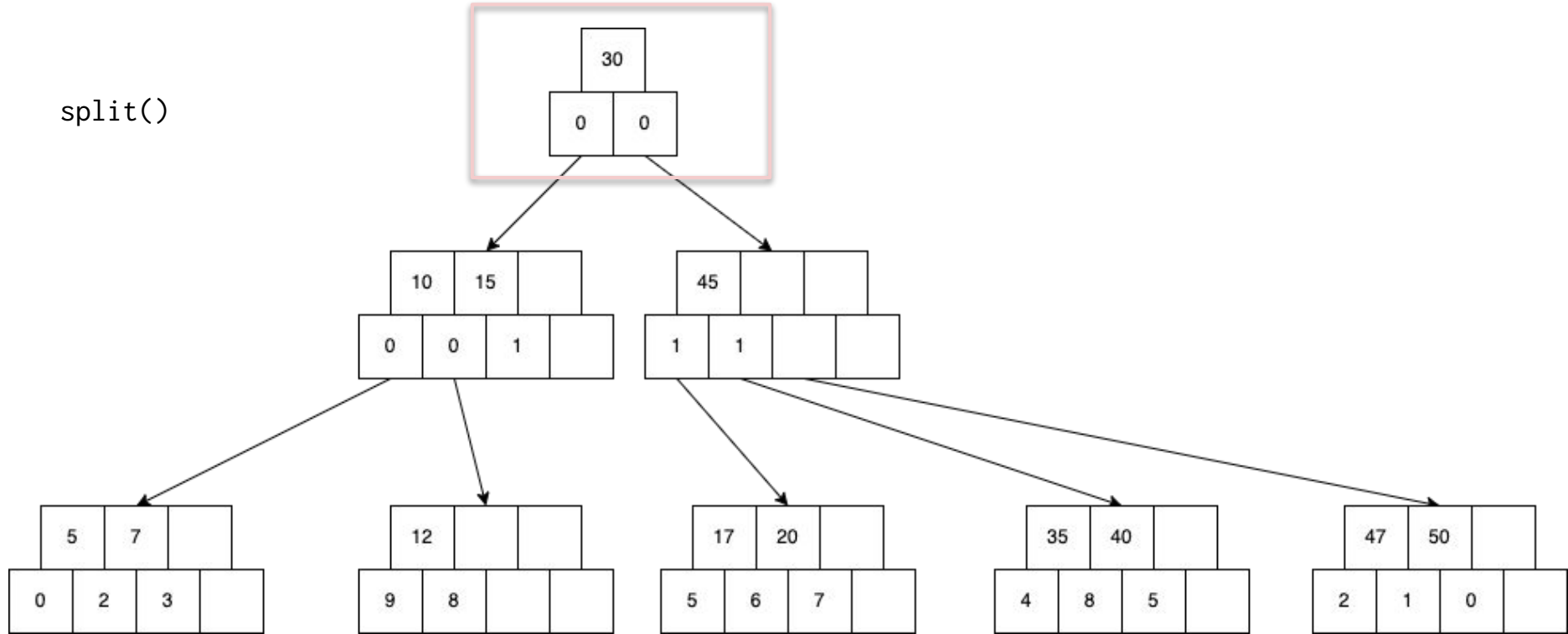
SB-tree: Node splits

split()



SB-tree: Node splits

split()

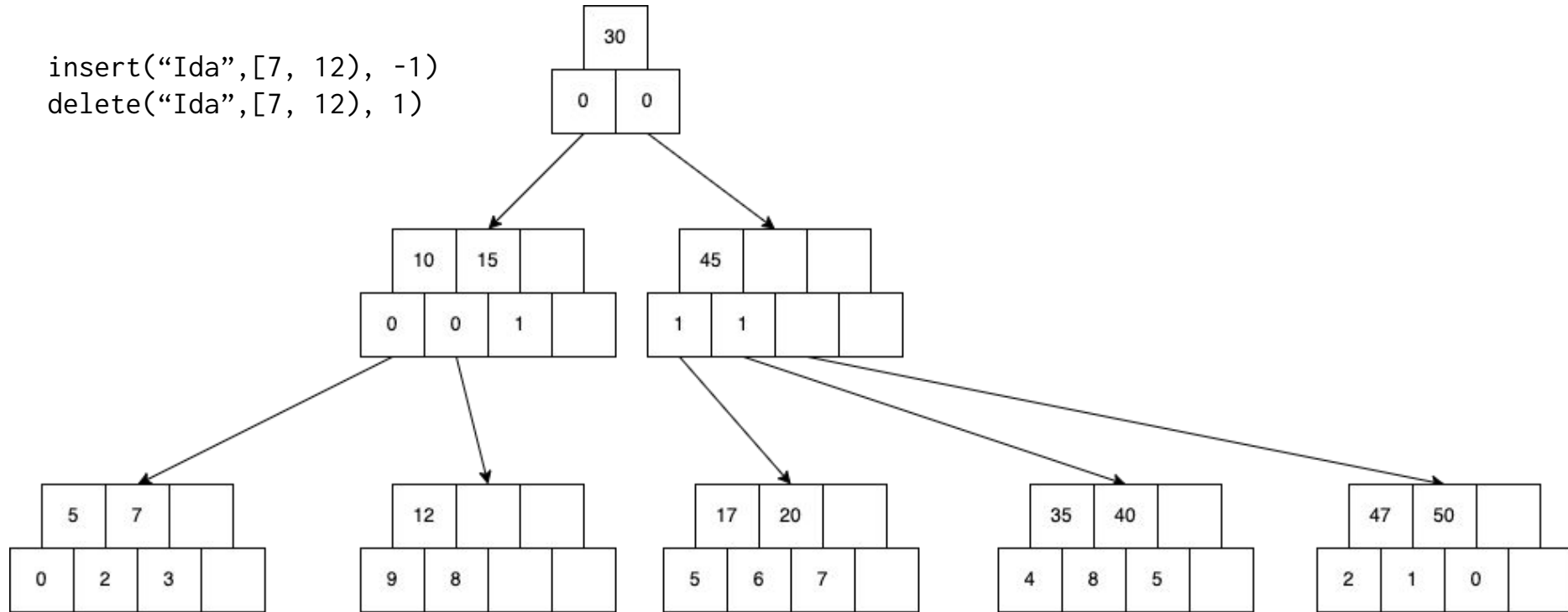


SB-tree: Delete

- Like inserts!
- Example:
 - delete { “Ida”, 1, [17, 47) }
 - is equivalent to insert { . , -1, [17, 47) }
- But underflows may occur!

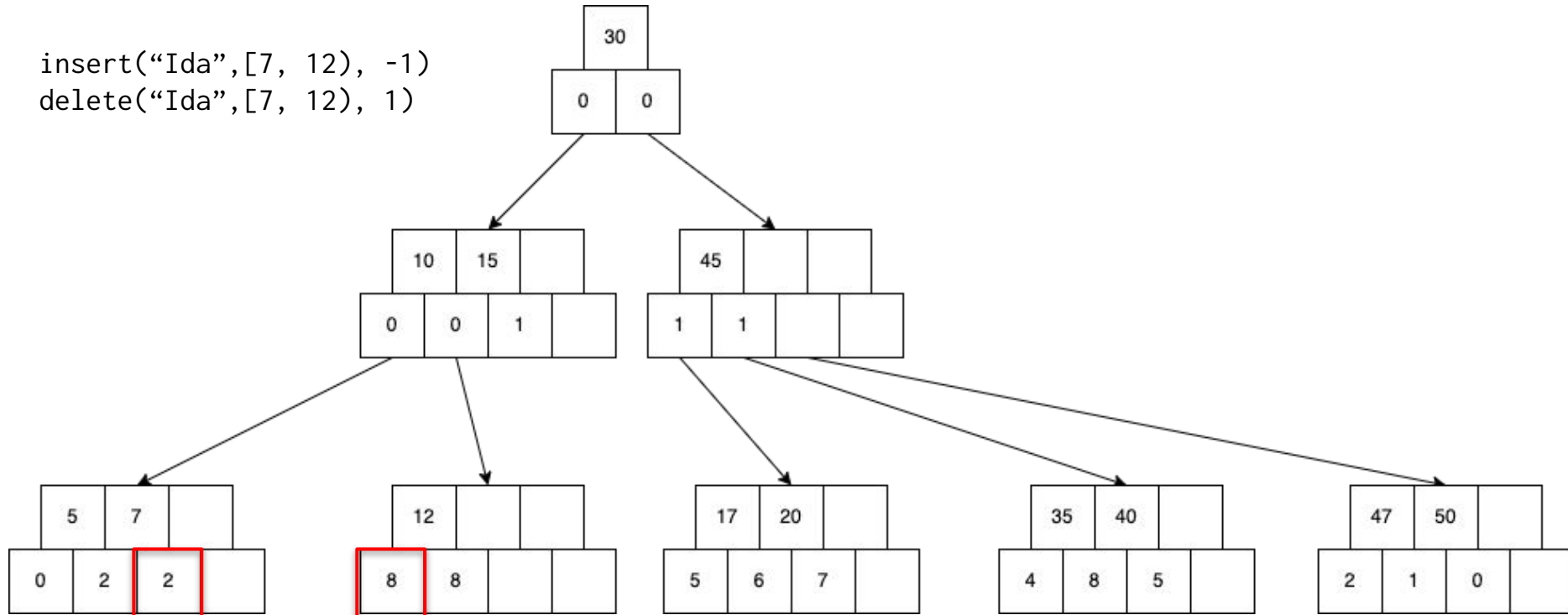
SB-tree: Node merges

insert("Ida",[7, 12), -1)
delete("Ida",[7, 12), 1)



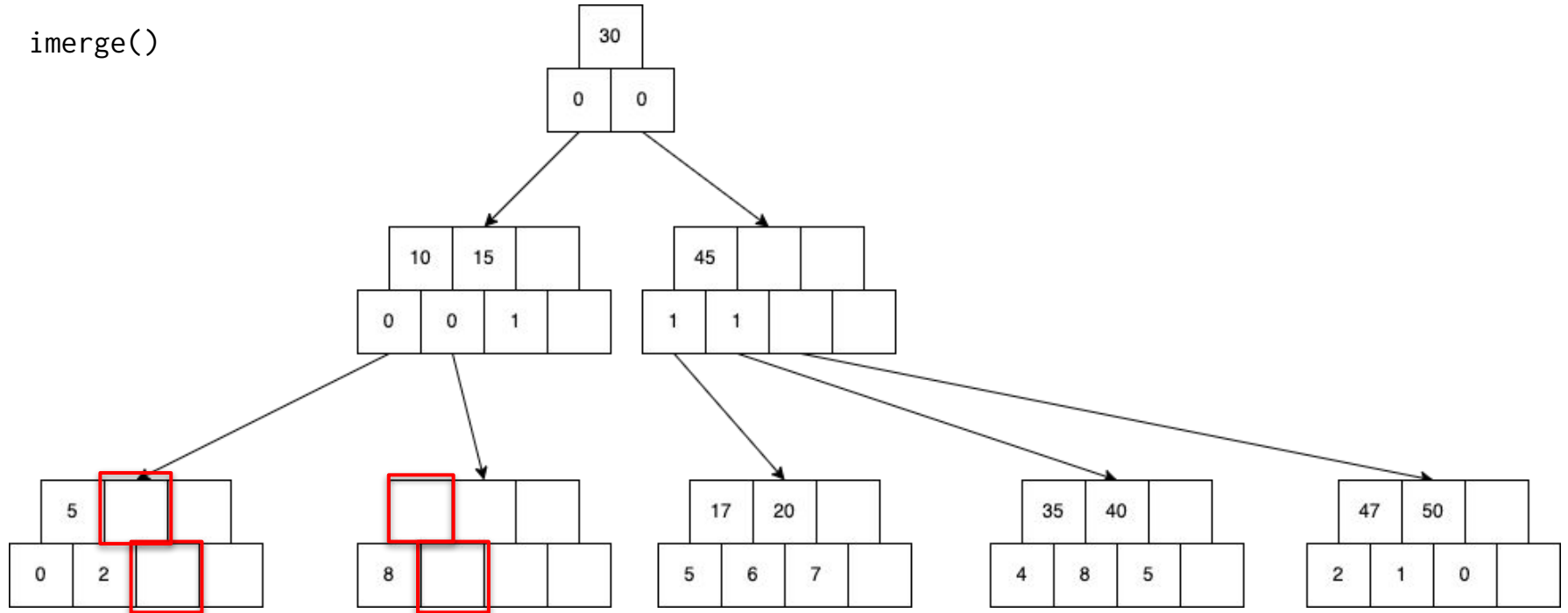
SB-tree: Node merges

insert("Ida",[7, 12), -1)
delete("Ida",[7, 12), 1)



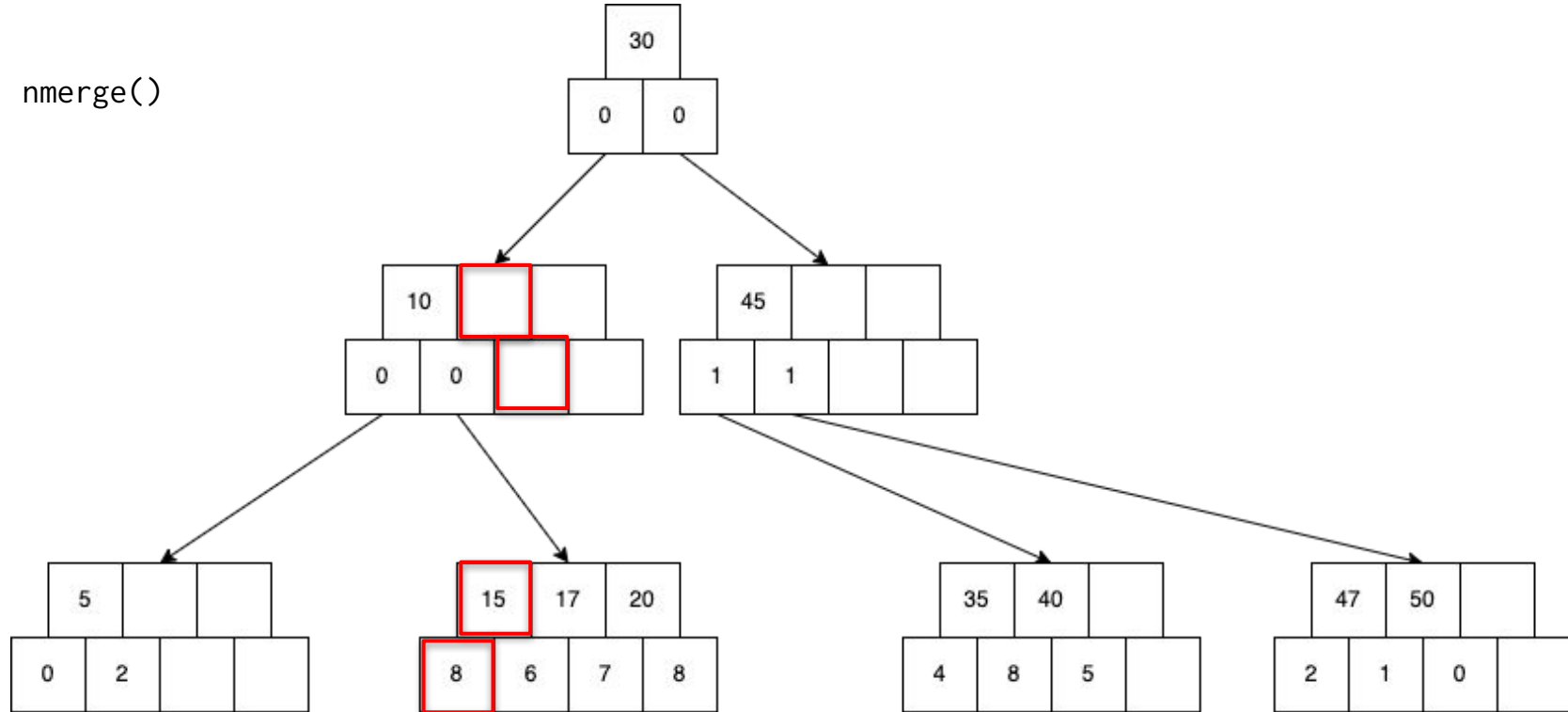
SB-tree: Node merges

imerge()



SB-tree: Node merges

nmerge()

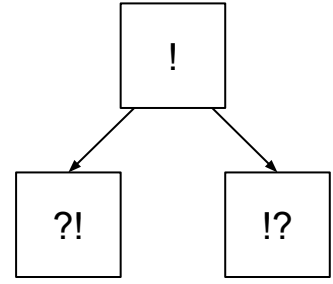


Performance

- Lookup: $O(h)$
- Range query: $O(h + r)$
- Insert: $O(h)$
- Delete: $O(h)$

h : height of the tree

r : number of leafs intersecting with the query range



Discussion

References

Incremental computation and maintenance of temporal aggregates

Jun Yang, Jennifer Widom

The VLDB Journal (2003) 12, p262-283