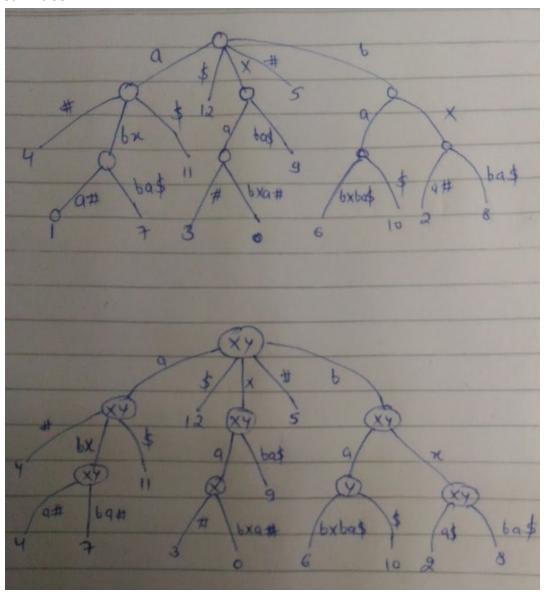
Algorithm for Circular DNA alignment:

Suppose we have two circular strings to be matched of length m and n. Now we use the suffix tree for finding the rotation of each.

1>Let's say string 1 is: xabxa and string 2 as: babxba. We make a single string out of the two by separating them via # character such that the final string to be put in suffix tree is xabxa#babxba. We then obtain the following suffix tree for both the strings to be compared. 2>Later we modify our substring with label 'XY', 'X', 'Y' representing that the string belong to Xand Y both or just Y or just X.

3>Considering this we traverse the entire suffix tree and find which all subsequences belong to both the strings.

4>Now we can generate the alignment by adding gap for the unmatched and get the final matching result of the alignment of all the possible rotations of circular strings using a single suffix tree.



Complexity:

For generating the suffix tree: O(m+n) //for all the rotations of each

For finding the subsequences: $O(max(m,n))^*(m+n)$ which is either $O(n^2 \text{ or } m^2 \text{ or } m^*n)$