

**Example 7.59:** Convert the following grammar

$$A \rightarrow BC$$

$$B \rightarrow CA \mid b$$

$$C \rightarrow AB \mid a$$

into GNF

**Solution:**

Let  $A = A_1$ ,  $B = A_2$ ,  $C = A_3$  and the resulting grammar is

$$A_1 \rightarrow A_2 A_3$$

$$A_2 \rightarrow A_3 A_1 \mid b$$

$$A_3 \rightarrow A_1 A_2 \mid a$$

First two productions are of the form:  $A_i \rightarrow A_j \alpha$  where  $i < j$ . So, we consider  $A_3$ -production.

**Consider  $A_3$ -production:** Substituting for  $A_1$  in  $A_3$ -production we get

$$A_3 \rightarrow A_1 A_2 \mid a = (A_2 A_3) A_2 \mid a = A_2 A_3 A_2 \mid a$$

Again replacing the first  $A_2$  in  $A_3$ -production we get,

$$\begin{aligned} A_3 &\rightarrow A_2 A_3 A_2 \mid a = (A_3 A_1 \mid b) A_3 A_2 \mid a \\ &= A_3 A_1 A_3 A_2 \mid b A_3 A_2 \mid a \end{aligned}$$

we get the resulting  $A_3$ -production as

$$A_3 \rightarrow A_3 A_1 A_3 A_2 \mid b A_3 A_2 \mid a$$

which is having left recursion. After eliminating left recursion, we get

$$\begin{aligned} A_3 &\rightarrow b A_3 A_2 \mid a \mid b A_3 A_2 Z \mid a Z \\ Z &\rightarrow A_1 A_3 A_2 \mid A_1 A_3 A_2 Z \end{aligned}$$

Now, all  $A_3$ -productions are in GNF.

**Consider  $A_2$ -productions:** Since all  $A_3$ -productions are in GNF, substituting  $A_3$ -productions in  $A_2$ -production we get

$$\begin{aligned} A_2 &\rightarrow (b A_3 A_2 \mid a \mid b A_3 A_2 Z \mid a Z) A_1 \mid b \\ &= b A_3 A_2 A_1 \mid a A_1 \mid b A_3 A_2 Z A_1 \mid a Z A_1 \mid b \end{aligned}$$

which is in GNF. Now, all  $A_2$ -productions are in GNF.

**Consider  $Z$ -productions:** Since  $A_1$ -productions are in GNF, substituting  $A_1$ -production in  $Z$ -production we get  $Z$ -productions also in GNF as shown below:

$$\begin{aligned} Z &\rightarrow A_1 A_3 A_2 \mid A_1 A_3 A_2 Z \\ &= (b A_3 A_2 A_1 A_3 \mid a A_1 A_3 \mid b A_3 A_2 Z A_1 A_3 \mid a Z A_1 A_3 \mid b A_3) A_3 A_2 \\ &\quad (b A_3 A_2 A_1 A_3 \mid a A_1 A_3 \mid b A_3 A_2 Z A_1 A_3 \mid a Z A_1 A_3 \mid b A_3) A_3 A_2 Z \end{aligned}$$

which can be written as

$$Z \rightarrow bA_3A_2A_1A_3A_3A_2 \mid aA_1A_3A_3A_2 \mid bA_3A_2ZA_1A_3A_3A_2 \mid \\ aZA_1A_3A_3A_2 \mid bA_3A_3A_2$$

$$Z \rightarrow bA_3A_2A_1A_3A_3A_2Z \mid aA_1A_3A_3A_2Z \mid bA_3A_2ZA_1A_3A_3A_2Z \mid \\ aZA_1A_3A_3A_2Z \mid bA_3A_3A_2Z$$

Since all productions are in GNF, the resulting grammar is also in GNF. So, the final grammar obtained in GNF notation is

$$G = (V, T, P, S)$$

where

$$V = \{A_1, A_2, A_3, Z\}$$

$$T = \{a, b\}$$

$$P = \{$$

$$A_1 \rightarrow bA_3A_2A_1A_3 \mid aA_1A_3 \mid bA_3A_2ZA_1A_3 \mid aZA_1A_3 \mid bA_3$$

$$A_2 \rightarrow bA_3A_2A_1 \mid aA_1 \mid bA_3A_2ZA_1 \mid aZA_1 \mid b$$

$$A_3 \rightarrow bA_3A_2 \mid a \mid bA_3A_2Z \mid aZ$$

$$Z \rightarrow bA_3A_2A_1A_3A_3A_2 \mid aA_1A_3A_3A_2 \mid bA_3A_2ZA_1A_3A_3A_2 \mid \\ aZA_1A_3A_3A_2 \mid bA_3A_3A_2$$

$$Z \rightarrow bA_3A_2A_1A_3A_3A_2Z \mid aA_1A_3A_3A_2Z \mid bA_3A_2ZA_1A_3A_3A_2Z \mid \\ aZA_1A_3A_3A_2Z \mid bA_3A_3A_2Z$$

}

$A_1$  is the start symbol

**Example 7.60:** Convert the following grammar

$$S \rightarrow AA \mid 0$$

$$A \rightarrow SS \mid 1$$

into GNF

**Solution:**

Let  $S = A_1$  and  $A = A_2$ . After substitution, the resulting grammar obtained is shown below:

$$A_1 \rightarrow A_2A_2 \mid 0$$

$$A_2 \rightarrow A_1A_1 \mid 1$$

First production is of the form:  $A_i \rightarrow A_j\alpha$  where  $i < j$ . So, we consider  $A_2$ -production.

**Consider  $A_2$ -production:** Substituting for  $A_1$  in  $A_2$ -production we get

$$A_2 \rightarrow A_2A_2A_1 \mid 0A_1 \mid 1$$

The above grammar is having left recursion. After eliminating left recursion, we get

$$A_2 \rightarrow 0A_1 \mid 1 \mid 0A_1Z \mid 1Z$$

$$Z \rightarrow A_2A_1 \mid A_2A_1Z$$

Now, all  $A_2$ -productions are in GNF.

**Consider  $A_1$ -productions:** Since all  $A_2$ -productions are in GNF, substituting  $A_2$ -productions in  $A_1$ -production we get

$$A_1 \rightarrow 0A_1A_2 \mid 1A_2 \mid 0A_1ZA_2 \mid 1ZA_2 \mid 0$$

which is in GNF. Now, all  $A_1$ -productions are in GNF.

**Consider  $Z$ -productions:** Since  $A_2$ -productions are in GNF, substituting  $A_2$ -production in  $Z$ -production we get  $Z$ -productions also in GNF as shown below:

$$Z \rightarrow 0A_1A_1 \mid 1A_1 \mid 0A_1ZA_1 \mid 1ZA_1$$

$$Z \rightarrow 0A_1A_1Z \mid 1A_1Z \mid 0A_1ZA_1Z \mid 1ZA_1Z$$

So, the final grammar obtained which is in GNF is shown below:

$$A_1 \rightarrow 0A_1A_2 \mid 1A_2 \mid 0A_1ZA_2 \mid 1ZA_2 \mid 0$$

$$A_2 \rightarrow 0A_1 \mid 1 \mid 0A_1Z \mid 1Z$$

$$Z \rightarrow 0A_1A_1 \mid 1A_1 \mid 0A_1ZA_1 \mid 1ZA_1$$

$$Z \rightarrow 0A_1A_1Z \mid 1A_1Z \mid 0A_1ZA_1Z \mid 1ZA_1Z$$

## Exercises

- 1) Define the following:      i) Context-free grammar      ii) Derivation
- 2) What is a sentence or sentential form? Explain with example.
- 3) Consider the following grammar

$$S \rightarrow aCa$$

$$C \rightarrow aCa \mid b$$

What is the language generated by this grammar?

- 4) Obtain a grammar to generate the following languages:
  - a.  $L = \{a^n b^n : n \geq 0\}$
  - b.  $L = \{a^{n+1} b^n : n \geq 0\}$
  - c.  $L = \{a^n b^{n+1} : n \geq 0\}$
  - d.  $L = \{a^n b^{n+2} : n \geq 0\}$
  - e.  $L = \{a^n b^{2n} : n \geq 0\}$
- 5) Obtain a grammar to generate set of all palindromes over  $\Sigma = \{a, b\}$ .
- 6) Obtain a grammar to generate  $L = \{ww^R \text{ where } w \in \{a, b\}^*\}$ .
- 7) Obtain a grammar to generate strings of all non-palindromes over  $\{a, b\}$ .

Obtain the grammar to generate following languages:

a)  $L = \{ 0^m 1^n 2^n \mid m \geq 1 \text{ and } n \geq 0 \}$

b)  $L = \{ w \mid n_a(w) = n_b(w) \}$

What is the language generated by the following grammar

$$S \rightarrow 0A \mid \epsilon$$

$$A \rightarrow 1S$$

Obtain a CFG to generate a string of balanced parentheses.

Obtain a grammar to generate the following languages:

a)  $L = \{ 0^i 1^j \mid i \neq j, i \geq 0 \text{ and } j \geq 0 \}$

b)  $L = \{ a^{n+2} b^m \mid n \geq 0 \text{ and } m > n \}$

c)  $L = \{ a^n b^m \mid n \geq 0, m > n \}$

d)  $L = \{ a^n b^{n-3} \mid n \geq 3 \}$

e)  $L = L_1 L_2$  where  $L_1 = \{ a^n b^m \mid n \geq 0, m > n \}$   $L_2 = \{ 0^n 1^{2n} \mid n \geq 0 \}$

f)  $L = L_1 \cup L_2$  where  $L_1 = \{ a^n b^m \mid n \geq 0, m > n \}$   $L_2 = \{ 0^n 1^{2n} \mid n \geq 0 \}$

g)  $L = \{ w : |w| \bmod 3 \neq |w| \bmod 2 \}$  on  $\Sigma = \{a\}$

h)  $L = \{ w : |w| \bmod 3 \geq |w| \bmod 2 \}$  on  $\Sigma = \{a\}$

Obtain a grammar to generate set of all strings with exactly one  $a$  when  $\Sigma = \{a, b\}$ .

Obtain a grammar generating all strings with at least one  $a$  if  $\Sigma = \{a, b\}$ .

Obtain a grammar to generate the set of all strings with no more than three  $a$ 's when  $\Sigma = \{a, b\}$ .

Obtain a grammar to generate the following languages:

a)  $L = \{ w \mid n_a(w) = n_b(w) + 1 \}$

b)  $L = \{ w \mid n_a(w) > n_b(w) \}$

c)  $L = \{ a^n b^m c^k \mid n + 2m = k \text{ for } n \geq 0, m \geq 0 \}$

d) Strings of 0's and 1's having a substring 000

e)  $L = \{ a^{n_1} b^{n_1} a^{n_2} b^{n_2} \dots a^{n_k} b^{n_k} : n, k \geq 0 \}$

Obtain a grammar to generate integers. Show the derivation for the unsigned number 1965 and signed number +1965.

Is the following grammar ambiguous?

$$S \rightarrow aSb \mid SS \mid \epsilon$$

Show that the following grammar is ambiguous.

$$S \rightarrow SbS \mid a$$

Given the following grammar:

$$E \rightarrow I \mid E+E \mid E * E \mid (E)$$

$$I \rightarrow a \mid b \mid c \mid d$$



Obtain the derivations for the strings  $(a+b)^*c^*d$  and  $a+b^*c$  and the parse tree for each derivation.

20. Obtain a reduced grammar for the grammar shown below:

$$\begin{aligned} S &\rightarrow aAa \\ A &\rightarrow Sb \mid bCC \mid aDA \\ C &\rightarrow ab \mid aD \mid \\ E &\rightarrow aC \\ D &\rightarrow aAD \end{aligned}$$

21. Find an equivalent grammar without  $\epsilon$ -productions for the grammar shown below.

$$\begin{aligned} S &\rightarrow aSa \mid bSb \mid A \\ A &\rightarrow aBb \mid bBa \\ B &\rightarrow aB \mid bB \mid \epsilon \end{aligned}$$

22. Remove the unit productions from the grammar

$$\begin{aligned} S &\rightarrow A \mid B \mid Cc \\ A &\rightarrow aBb \mid B \\ B &\rightarrow aB \mid bb \\ C &\rightarrow Cc \mid B \end{aligned}$$

23. Eliminate useless productions from the grammar

$$\begin{aligned} S &\rightarrow aA \mid a \mid B \mid C \\ A &\rightarrow aB \mid \epsilon \\ B &\rightarrow aA \\ C &\rightarrow CcD \\ D &\rightarrow abd \end{aligned}$$

**Note:** First remove  $\epsilon$  and unit productions, then simplify CFG.

24. Obtain the following grammar in CNF

$$\begin{aligned} E &\rightarrow E + T \mid T \\ T &\rightarrow T * F \mid F \\ F &\rightarrow (E) \mid I \\ I &\rightarrow a \mid b \mid c \mid Ia \mid Ib \mid Ic \end{aligned}$$

25. Obtain the following grammar in CNF and GNF

$$\begin{aligned} S &\rightarrow aA \mid a \mid B \mid C \\ A &\rightarrow aB \mid \epsilon \\ B &\rightarrow aA \mid \\ C &\rightarrow cCD \\ D &\rightarrow abd \end{aligned}$$

26. Simplify the following CFG and convert it into CNF

$$S \rightarrow AaB \mid aaB$$

$$A \rightarrow \epsilon$$

$$B \rightarrow bbA \mid \epsilon$$

27. Convert the following grammar into GNF

$$S \rightarrow abAB$$

$$A \rightarrow bAB \mid \epsilon$$

$$B \rightarrow Aa \mid \epsilon$$

28. What is left recursion? How it can be eliminated?

29. Eliminate left recursion from the following grammar

$$E \rightarrow E + T \mid T$$

$$T \rightarrow T * F \mid F$$

$$F \rightarrow (E) \mid id$$

30. Eliminate the useless symbols in the grammar

$$S \rightarrow aA \mid bB$$

$$A \rightarrow aA \mid a$$

$$B \rightarrow bB$$

$$D \rightarrow ab \mid Ea$$

$$E \rightarrow aC \mid d$$

31. Eliminate left recursion from the following grammar

$$S \rightarrow Ab \mid a$$

$$A \rightarrow Ab \mid Sa$$

32. What is the need for simplifying a grammar?

33. What is a useless symbol? Explain with example.

34. Simplify the following grammar

$$S \rightarrow aA \mid a \mid Bb \mid cC$$

$$A \rightarrow aB$$

$$B \rightarrow a \mid Aa$$

$$C \rightarrow cCD$$

$$D \rightarrow ddd$$

35. What is an  $\epsilon$  - production? What is a Nullable variable? Explain with example.

36. Eliminate all  $\epsilon$ - productions from the grammar

$$S \rightarrow ABCa \mid bD$$

$$A \rightarrow BC \mid b$$

$$B \rightarrow b \mid \epsilon$$

$$C \rightarrow c \mid \epsilon$$

$$D \rightarrow d$$

37. Eliminate all  $\epsilon$ - productions from the grammar

$$S \rightarrow BAAB$$

$$A \rightarrow 0A2 \mid 2A0 \mid \epsilon$$

$$B \rightarrow AB \mid 1B \mid \epsilon$$

38. Obtain a grammar to generate an arithmetic expression using the operators  $+$ ,  $-$ ,  $*$ ,  $/$  and  $^$  (indicating power). An identifier can start with any of the letters from  $\{a, b, c\}$  and can be followed by zero or more symbols from  $\{a, b, c\}$

39. Define the following terms:

a) Derivation tree

b) yield of a tree

c) ambiguous grammar

40. Is the following grammar ambiguous?

$$S \rightarrow iCtS \mid iCtSeS \mid a$$

$$C \rightarrow b$$

41. What is dangling else problem? How dangling else problem can be solved?"

42. Eliminate ambiguity from the following ambiguous grammar:

$$S \rightarrow iCtS \mid iCtSeS \mid a$$

$$C \rightarrow b$$

43. Convert the following ambiguous grammar into unambiguous grammar

$$E \rightarrow E * E \mid E - E$$

$$E \rightarrow E \wedge E \mid E / E$$

$$E \rightarrow E + E$$

$$E \rightarrow (E) \mid id$$

44. What is inherently ambiguous grammar?"

45. Obtain the inherently ambiguous grammar for the following inherently ambiguous languages:

$$L = \{a^n b^n c^m d^m \mid m \geq 1, n \geq 1\} \cup \{a^n b^m c^m d^n \mid m \geq 1, n \geq 1\}$$

$$L = \{a^i b^j c^k \mid i, j, k, \geq 0 \text{ and } i = j \text{ or } j = k\}$$

46. What is left recursion? How to eliminate left recursion?

47. What is useless variable? Explain with example.