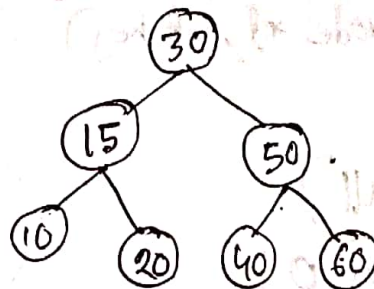


BST → Binary search tree

Lecture 1



Left side element
is smaller than
root

Right > root
left < root

this type Binary
tree for searching

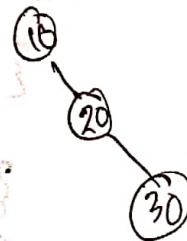
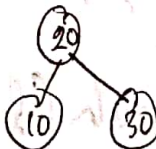
① No Duplicates

② Inorder gives sort order

③ No. of BST for "n" nodes

④ remains
and search.

Inorder : 10, 15, 20, 30, 40, 50, 60



⑤ shape
are possible

$$\frac{2^n C_n}{n+1}$$

Catalan number
of tree generation

BST → represent using
Link list most of the
time

time taking in BST searching
lecture 2 $\log n \leq h \leq n$ $O(h) = O(\log n)$

Node *

Rsearch (Node *t, int key)

height

t

root

if (t == Null)

return 0;

if (key == t->data)

{ return t;

}

else if (key < t->data)

{

Rsearch(t->lchild, key);

}

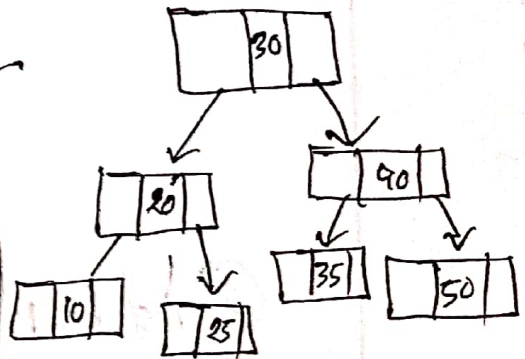
else if (key > t->data)

{

Rsearch(t->rchild, key);

}

}



key = 25

① t first root

② t < key → lchild
else → rchild

③ repeat ②

④ if t == key
return t

⑤ last null

Return

Rsearch(t->lchild, key)

⇒ iterative function using loop

Node * search(Node *t, int key)

{

while (t != null)

{

if (key == t->data)

{ return t;

}

else if (key < t->data)

{ t = t->lchild;

else if (key > t->data)

{ t = t->rchild;

}

}

return 0;

}

$O(\log n)$

Inserting in Binary Search tree

Lecture 3

```
void Insert(Node *f, int key)
```

```
{
    Node *r = Null, *p;
```

```
    while(f != Null)
```

```
    {
        r = f; // follow (f)
        if (key == f->data)
        {
            return;
        }
    }
```

```
    else if (key < f->data)
    {
        f = f->lchild;
    }
```

```
    else if (key > f->data)
    {
        f = f->rchild;
    }
}
```

creating new node

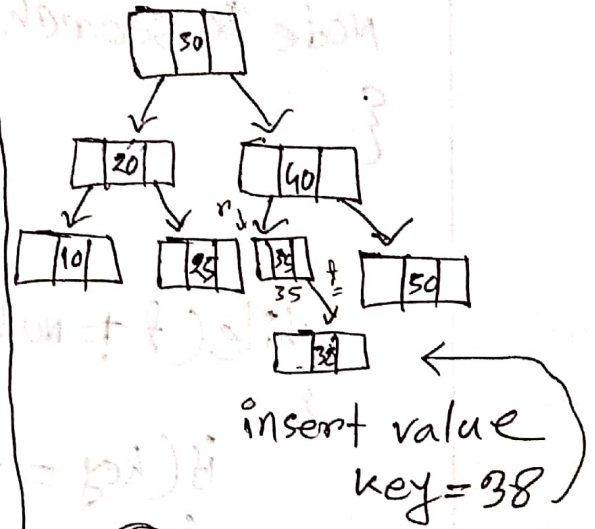
```
p = New Node;
```

```
if (key < r->data)
{
    r->lchild = p;
}
```

```
else
{
    r->rchild = p;
}
```

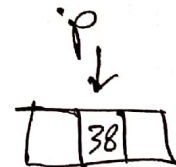
```
p->data = key;
p->lchild = p->rchild = Null;
```

compare for left or Right New Node Insert



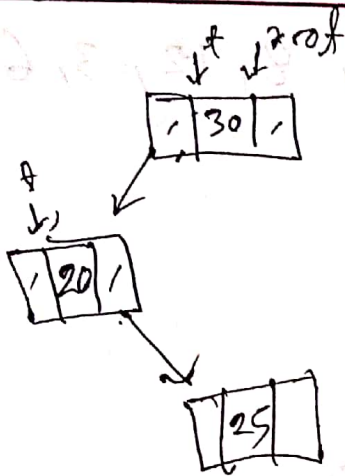
① Searching and insert the key

② using trail pointer
 (r) follow the (f)



Recursive Insert

Lecture 4



```
Node* insert(Node* p, int key)
{
    Node *t;
```

```
    if(p == Null)
```

```
    {
        t = new Node;
        t->data = key;
        t->lchild = t->rchild = null;
        return t;
    }
```

```
    if(key < p->data)
```

```
    {
        p->lchild = insert(p->lchild, key);
```

```
    }
    else if(key > p->data)
```

```
    {
        p->rchild = insert(p->rchild, key);
```

```
    }
    return p;
```

When return
that time
Link up

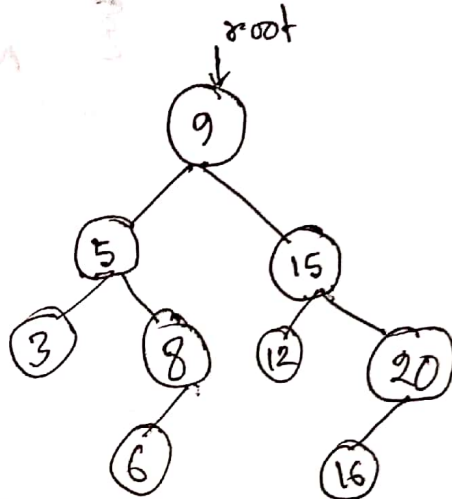
```
int main()
```

```
{
    Node *root = null;
    root = insert(root, 30);
    insert(root, 20);
    insert(root, 25);
}
```

Lecture 5

Creating Binary search tree. and Delete

key: 9, 15, 5, 20, 16, 8, 12, 3, 6

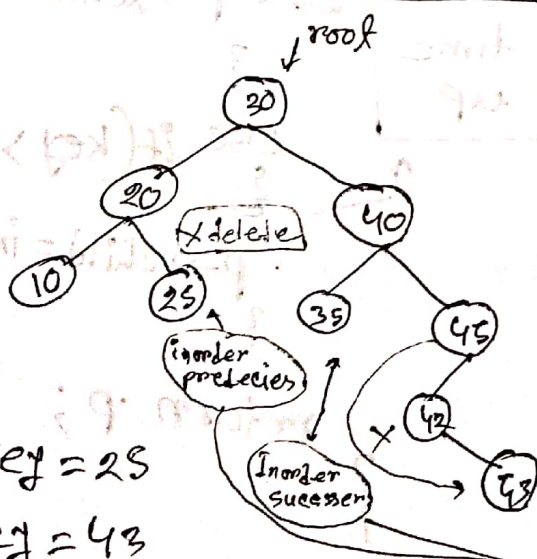


Lecture 6

⇒ code in pdf

Lecture 7

Deleting Binary Search tree



key = 25

1. search key
2. deleting the node. if success

key = 25

key = 43

key = 42

key = 30

here is a problem
sub tree are there
we have to modife
the link

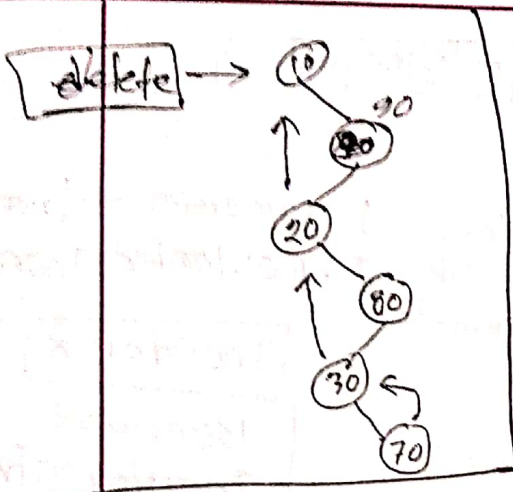
root
befor and after

(like)

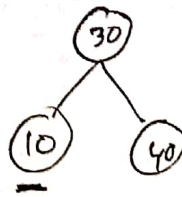
In: 10, 20, 25, 30, 35, 40...

Find out Inorder Predecess // who are before 30 root

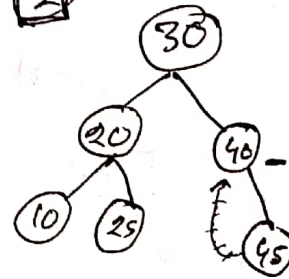
number of modification $O(\log n)$
 deleting $O(\log n)$



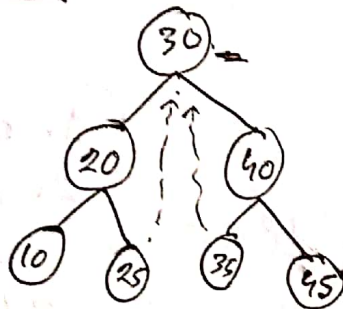
(1)



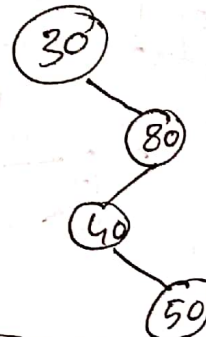
(2)



(3)



(4)



Between (25, 35) one of them, predices or succesor

Node * Delete(Node *p, int key)

if (height(p->lchild) > height(p->rchild))

{ q = Inpre(p->lchild)

p->data = q->data;

p->lchild = Delete(p->lchild, q->data)

}

else { if (with rchild)

}

if (key < p->data)

{ p->lchild = Delete(p->lchild, key)

}

else if (key > p->data)

{ p->rchild = Delete(p->rchild, key)

}

else

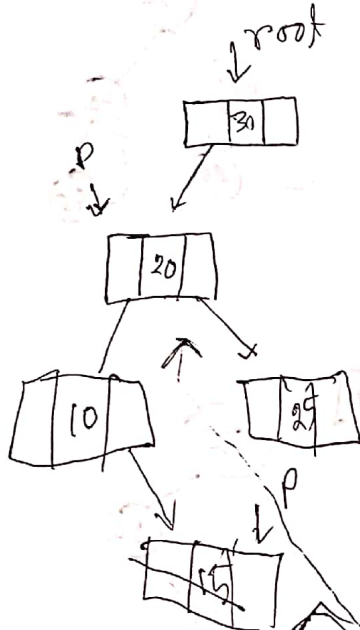
{

}

Generating BST from Preorder

Lecture 9 :

Pre [30 | 20 | 10 | 15 | 25 | 40 | 50 | 45 |]
i ----- ↑



(p) this is smaller than root add left and Right

1. preorder + Inorder
2. postorder + Inorder

Inorder X

Because Inorder gives sorted order

(BST)

1. Need stack
2. Left element and push in stack
3. (i) check for left and Right

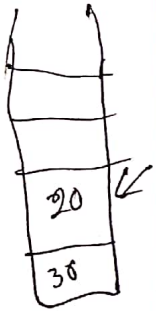
(9) when left child then push address in stack

if Right X push

When right child and root not lying value pop in the stack

Simply think that creating tree by human being what condition follow check root value and

Left and Right this type follow and code



stack

Lecture 10

Create code.

Bst \rightarrow preorder

```
void create(int pre[], int n)
{
    stack stk;
    Node * f;
    int i = 0;
    root = new Node;
    root->data = pre[i++];
    root->lchild = root->rchild = null;
    p = root;
```

```
while (i < n)
```

```
{ if (pre[i] < p->data)
```

```
{ f = new Node;
  f->data = pre[i++];
  f->lchild = f->rchild = null;
  p->lchild = f;
  push(&stk, p);
  p = f;
}
```

```
else (pre[i] > p->data)
```

```
{ if (pre[i] > p->data && pre[i] < stack_top(stack)->data)
```

```
{ f = new Node; f->data = pre[i++];
  f->lchild = f->rchild = null;
  p->rchild = f; p = f;
}
```

```
else { p = pop(&stk);
```

```
}
```

if data is
pre
getter from
root value
each
then it is
Range [root, last
value]
then one case
or another
case

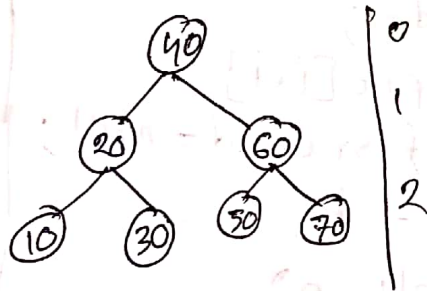
pop from
stack

Lecture 11

Drawback of Binary search tree

- * we can not control height of Binary Search tree :
- * we are expecting that height $O(\log n)$ But not
- * we need another method to control of height BST.
- * we can not control user Insertion :

key : 40, 20, 30, 60, 50, 10, 70

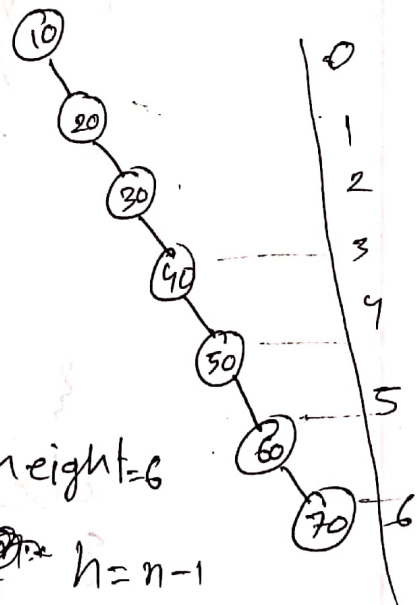


height = 2

$$\log_2(n+1) - 1$$

$$\approx O(\log n)$$

key : 10, 20, 30, 40, 50, 60, 70



height = 6

$$h = n - 1$$

$$O(n)$$

We need AVL tree
tree itself control height