

# Lab Report-2

**Dijkstra:** Dijkstra's Algorithm is a method used in computer science to find the **shortest path** from a starting node (point) to all other nodes in a network (graph). Its primary goal is to determine the most efficient route by minimizing the total "cost" accumulated along the path. This algorithm is specifically designed to work on **weighted graphs**.

## Graph Search Example:

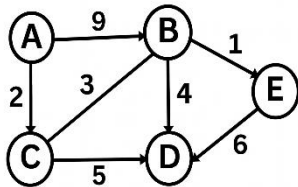
**Relaxation Formula:** It's the mechanism that checks if a shorter path to a neighbor has been found.

**Formula:**  $\text{if}(d[u] + c(u,v) < d[v]) \Rightarrow d[v] = d[u] + c(u,v)$

## Meaning of Each Symbol

- **u** : The *current node* (the node we are exploring now).
- **v** : A *neighbor node* of u.
- **d[u]** : The shortest distance from the start node to u (known so far).
- **d[v]** : The shortest distance from the start node to v (recorded so far).
- **c(u,v)** : The cost/weight of the edge from u to v.

## Graph:



Find the path from A to E, where A is the source Node.

## Initialization:

	A	B	C	D	E
A	0	$\infty$	$\infty$	$\infty$	$\infty$

Node	A	B	C	D	E
Parent	-1	-1	-1	-1	-1

**Step-1:**

	A	B	C	D	E
A	0	$\infty$	$\infty$	$\infty$	$\infty$
C	0	9	2	$\infty$	$\infty$

**Check Condition for B:**

$$d[u] + c(u,v) < d[v] \Rightarrow 0 + 9 < \infty$$

Condition is **true**, so we **relax**:

$$d[v] = d[u] + c(u,v) = 0 + 9 = 9$$

**Check Condition for C:**

$$d[u] + c(u,v) < d[v] \Rightarrow 0 + 2 < \infty$$

Condition is **true**, so we **relax**:

$$d[v] = d[u] + c(u,v) = 0 + 2 = 2$$

Node	A	B	C	D	E
Parent	-1	A	A	-1	-1

**Step-2:**

	A	B	C	D	E
A	0	$\infty$	$\infty$	$\infty$	$\infty$
C	0	9	2	$\infty$	$\infty$
B	0	5	2	7	$\infty$

**Check Condition for B:**

$$d[u] + c(u,v) < d[v] \Rightarrow 2 + 3 < 9$$

Condition is **true**, so we **relax**:

$$d[v] = d[u] + c(u,v) = 2 + 3 = 5$$

**Check Condition for D:**

$$d[u] + c(u,v) < d[v] \Rightarrow 2 + 5 < \infty$$

Condition is **true**, so we **relax**:

$$d[v] = d[u] + c(u,v) = 2 + 5 = 7$$

Node	A	B	C	D	E
Parent	-1	C	A	B	-1

**Step-3:**

	A	B	C	D	E
A	0	$\infty$	$\infty$	$\infty$	$\infty$
C	0	9	2	$\infty$	$\infty$
B	0	5	2	7	$\infty$
E	0	5	2	7	6

**Check Condition for D:**

$$d[u] + c(u,v) < d[v] \Rightarrow 5+4 > 7$$

Condition is **false**, so:

$$d[v] = 7$$

**Check Condition for E:**

$$d[u] + c(u,v) < d[v] \Rightarrow 5+1 < \infty$$

Condition is **true**, so we **relax**:

$$d[v] = d[u] + c(u,v) = 5 + 1 = 6$$

Node	A	B	C	D	E
Parent	-1	C	A	7	B

**Step-4:**

	A	B	C	D	E
A	0	$\infty$	$\infty$	$\infty$	$\infty$
C	0	9	2	$\infty$	$\infty$
B	0	5	2	7	$\infty$
E	0	5	2	7	6
D	0	5	2	7	6

**Check Condition for D:**

$$d[u] + c(u,v) < d[v] \Rightarrow 6+6 > 7$$

Condition is **false**, so:

$$d[v] = 7$$

Node	A	B	C	D	E
Parent	-1	C	A	7	B

➤ **The Shortest path from A to E Node:**

A → C → B → E

**Pseudocode:**

```
Dijkstra(G, source) {  
    for each vertex u in G {           // initialization  
        dist[u] = infinity  
        parent[u] = null  
    }  
  
    dist[source] = 0                    // initialize source  
  
    PQ = priority_queue()              // min-priority queue  
    PQ.push( (0, source) )            // (distance, vertex)  
  
    while (PQ is not empty) {  
        (du, u) = PQ.pop()             // extract node with min distance  
        for each v in Adj[u] {         // explore all neighbors  
            w = weight(u, v)  
            if (dist[u] + w < dist[v]) { // relaxation  
                dist[v] = dist[u] + w  
                parent[v] = u  
                PQ.push( (dist[v], v) ) // push updated distance  
            }  
        }  
    }  
  
    for each vertex u in G {  
        print dist[u]  
    } }
```