



YILDIZ TECHNICAL UNIVERSITY

Department of Computer Engineering

Algorithm Analysis
Homework-1

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20011011 Mehmet Emin Aydın Homework 1

1) a) $c_1 \cdot g(n) \leq (n^2+1)^{10}$ ve $c_2 \cdot g(n) \geq (n^2+1)^{10}$

$$g(n) = n^{20} \Rightarrow c_1 \cdot n^{20} \leq (n^2+1)^{10}$$

$$n_0 = 1 \Rightarrow n_0 \leq n \Rightarrow c_1 = 1,$$

$$c_2 \cdot n^{20} \geq (n^2+1)^{10} \Rightarrow c_2 \cdot (n^2)^{10} \geq (n^2+1)^{10}$$

$$\sqrt[10]{c_2} \cdot n^2 \geq n^2+1 \Rightarrow n_0 = 1$$

$$n_0 \leq n \text{ için } \sqrt[10]{c_2} = 2 \Rightarrow c_2 = 2^{10}$$

$$(n^2+1)^{10} \in \Theta(n^{20})$$

b) $c_1 \cdot g(n) \leq \sqrt{10n^2+7n+3} \leq c_2 \cdot g(n)$

$$c_1 \cdot g(n) \leq \sqrt{10n^2+7n+3} \Rightarrow g(n) = \sqrt{n^2}$$

$$n_0 = 1 \text{ ve } n_0 \leq n \text{ için } c_1 = \sqrt{10}$$

$$c_2 \cdot g(n) \geq \sqrt{10n^2+7n+3} \quad n_0 = 1 \text{ için}$$

$$n_0 \leq n \text{ için } c_2 = \sqrt{20}$$

$$\sqrt{10n^2+7n+3} \in \Theta(n)$$

2) $\sum_{i=0}^{n-1} \sum_{j=0}^{i-1} (i+j) \rightarrow \sum_{j=0}^{i-1} (i) + \sum_{j=0}^{i-1} (j) = i^2 + \frac{(i-1) \cdot i}{2}$

$$\sum_{i=0}^{n-1} \left(\frac{3i^2}{2} - \frac{i}{2} \right) = \frac{3}{2} \sum_{i=0}^{n-1} (i^2) - \frac{1}{2} \sum_{i=0}^{n-1} (i)$$

$$\frac{3}{2} \cdot \frac{(n-1) \cdot n \cdot (2n-1)}{6} - \frac{1}{2} \cdot \frac{(n-1) \cdot n}{2}$$

$$\frac{2n^3 - 3n^2 + n}{4} - \frac{(n^2 - n)}{4} = \frac{2n^3 - 4n^2 + 2n}{4} = \frac{n^3 - 2n^2 + n}{2}$$

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3) a) Mathematical Notation of the code $\Rightarrow \sum_{k=i}^n \sum_{j=i+1}^{n-1} \sum_{i=0}^{n-2} 1$

$$\sum_{i=0}^{n-2} 1 = n-2-0+1 = n-1$$

$$\begin{aligned} \sum_{j=i+1}^{n-1} (n-1) &= \sum_{j=i+1}^{n-1} (n) - \sum_{j=i+1}^{n-1} 1 = (n-1-(i+1)+1) \cdot n - (n-1-(i+1)+1) \cdot 1 \\ &= (n-i-1) \cdot n - (n-i-1) \\ &= n^2 - i \cdot n + n - n + i + 1 \\ &= n^2 - i \cdot n + i + 1 \end{aligned}$$

$$\sum_{k=i}^n (n^2) - n \cdot \sum_{k=i}^n i + \sum_{k=i}^n i + \sum_{k=i}^n 1 = (n-i+1) \cdot n^2 - \left(\frac{n(n+1)}{2} - \frac{(i-1)i}{2} \right) \cdot (n+1) + (n-i)$$

$$= n^3 - i \cdot n^2 + n^2 - \left(\frac{n^2 + n - i^2 + i}{2} \right) \cdot (n+1) + (n-i)$$

The fastest growing term n^3 olduğundan

$$Big-O = O(n^3)$$

b) $A[j][i]$ 'nin 0 olduğu değerler için en içteki for döngüsü hiç işlenmeyeceği için while döngüsüne çevrildi.

for $i \leftarrow 0$ to $n-2$ do

for $j \leftarrow i+1$ to $n-1$ do

$k \leftarrow i$

while $(k < n$ and $A[j][i] \neq 0)$

$A[j][k] = A[j][k] - A[i][k] * A[j][i] / A[i][i];$

4)

$$T(n) = T(n/3) + 1$$

$$T(n/3) = T(n/9) + 1$$

$$T(n/9) = T(n/27) + 1$$

$$T(n) = T\left(\frac{n}{3^k}\right) + k$$

$$n = 3^k \Rightarrow k = \log_3 n$$

$$T(n) = \frac{T(1)}{0} + \log_3 n \Rightarrow T(n) = \log_3 n$$