

🔑 master ▾

system-design / system-design-master 2 /
typeahead.md

Go to file

...



Vineet Sagar Self Prep System D...

Latest commit b89d37c on Apr 22, 2020

🕒 History

👤 0 contributors

Typeahead

- [Scenario](#)
- [Initial design](#)
- [Storage](#)
 - [Query service DB](#)
 - [Word count table](#)
 - [Prefix table](#)
 - [Trie](#)
 - [Data collections service](#)
- [Scale](#)
 - [How to reduce response time](#)
 - [What if the trie too large for one machine](#)
 - [How to reduce the size of log file](#)

Scenario

- [Google suggestion](#)

- Prefix -> top n hot key words
- DAU: 500M
- Search: 66500M = 18b (Every one search for 6 words, each word has 6 characters)
- QPS = 18b / 86400 ~ 200k
- Peak QPS = QPS * 2 ~ 400k
- Twitter typeahead

Initial design

- Querv service

☰ 105 lines (87 sloc) | 4.01 KB

Raw

Blame



- Data collection service

Storage

Query service DB

Word count table

- How to query on the db
- Query SQL: Select * from hit_stats where keyword like \${key}% order by hitCount DESC Limit 10
 - Like operation is expensive. It is a range query.
 - where keyword like 'abc%' is equivalent to where keyword >= 'abc' AND keyword < 'abd'

keyword	hitCount
Amazon	20b
Apple	15b
Adidas	7b
Airbnb	3b

Prefix table

- Convert a keyword table to a prefix table, put into memory

prefix	keywords
a	"amazon","apple"
am	"amazon","amc"
ad	"adidas","adobe"
don	"don't have", "donald trump"

Trie

- Trie (in memory) + Serialized Trie (on disk).
 - Trie is must faster than DB because
 - All in-memory vs DB cache miss
- Store word count at node, but it's slow
 - e.g. TopK. Always need to traverse the entire trie. Exponential complexity.
- Instead, we can store the top n hot key words and their frequencies at each node, search becomes $O(\text{len})$.

prefix	keywords
a	"amazon","apple"
am	"amazon","amc"
ad	"adidas","adobe"
don	"don't have", "donald trump"

- How do we add a new record {abd: 3b} to the trie
 - Insert the record into all nodes along its path in the trie.
 - If a node along the path is already full, then need to loop through all records inside the node and compared with the node to be inserted.

Data collections service

- How frequently do you aggregate data
 - Real-time not impractical. Read QPS 200K + Write QPS 200K. Will slow down query service.
 - Once per week. Each week data collection service will fetch all the data within the most recent one week and aggregate them.
- How does data collection service update query service? Offline update and works online.
 - All in-memory trie must have already been serialized. Read QPS already really high. Do not write to in-memory trie directly.
 - Use another machine. Data collection service updates query service.

Scale

How to reduce response time

- Cache result
 - Front-end browser cache the results
- Pre-fetch
 - Fetch the latest 1000 results

What if the trie too large for one machine

- Use consistent hashing to decide which machine a particular string belongs to.
 - A record can exist only in one machine. Sharding according to char will not distribute the resource evenly. Instead, calculate consistent hashing code
 - a, am, ama, amax stored in different machines.

How to reduce the size of log file

- Probabilistic logging.
 - Too slow to calculate and too large amount of data to store.
 - Log with 1/10,000 probability
 - Say over the past two weeks "amazon" was searched 1 billion

times, with $1/1000$ probability we will only log 1 million times.

- For a term that's searched 1000 times, we might end up logging only once or even zero times.