



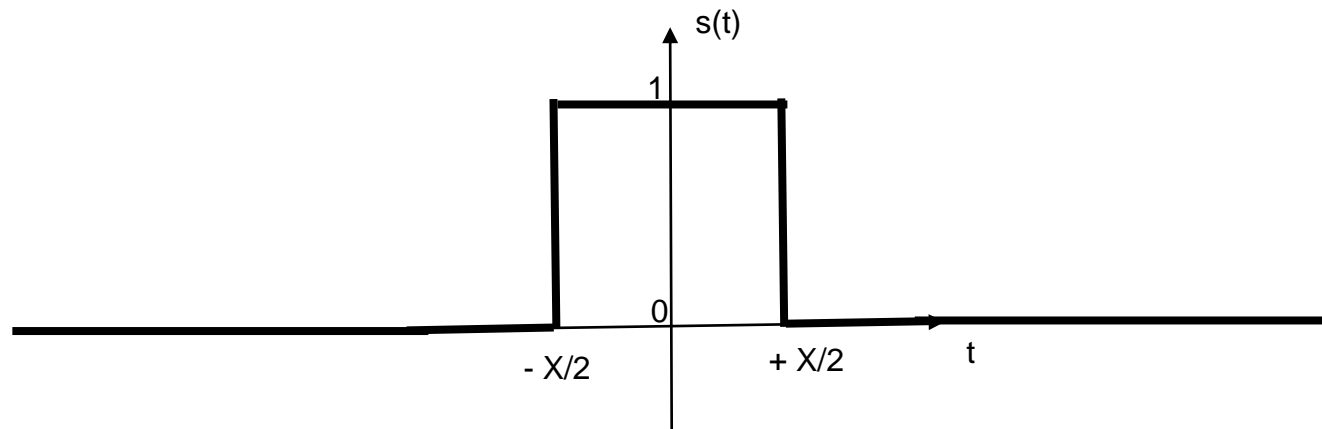
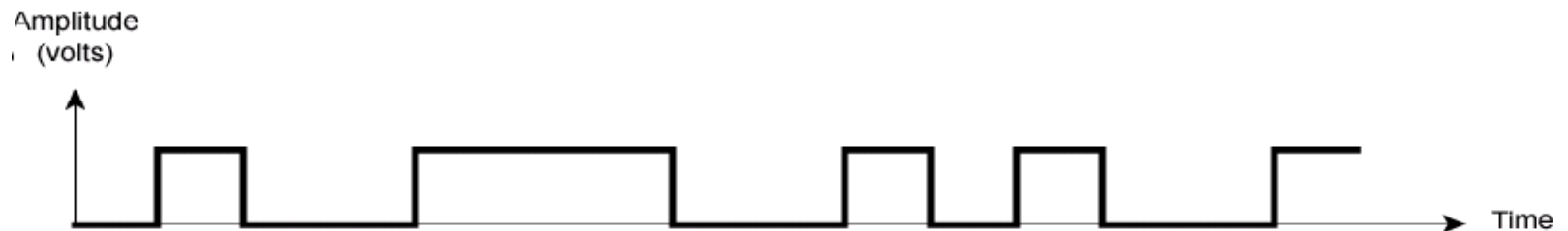
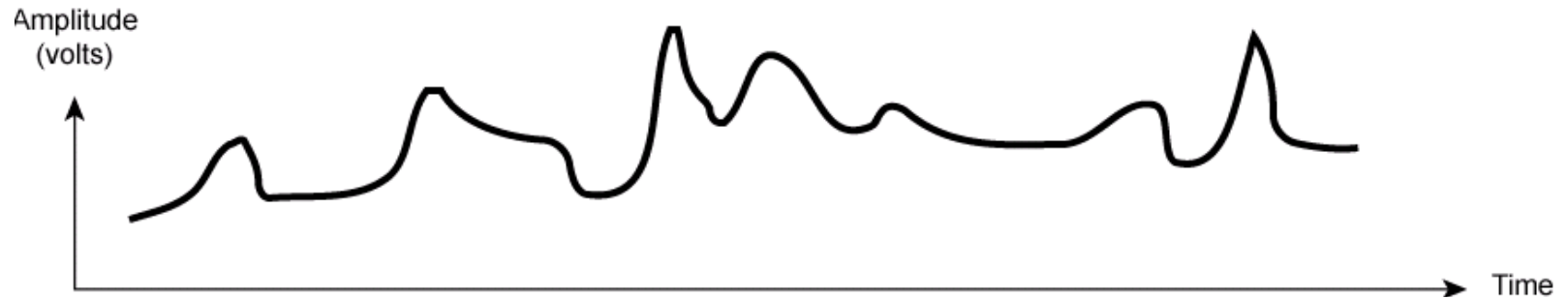
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Computer Networks 1

BANDWIDTH AND CAPACITY

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APERIODIC SIGNALS IN TIME



PERIODIC SIGNALS

For any periodic wave:

$$S(t + mT) = S(t); \quad 0 \leq t \leq T$$

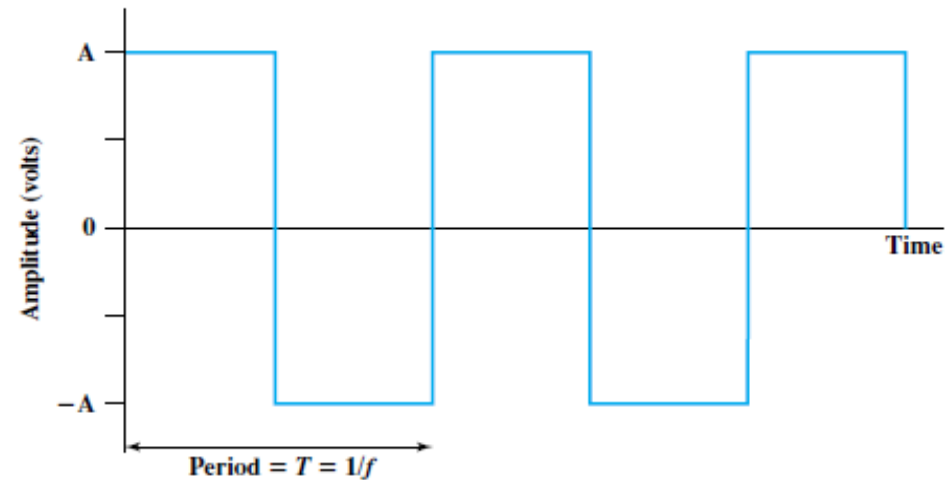
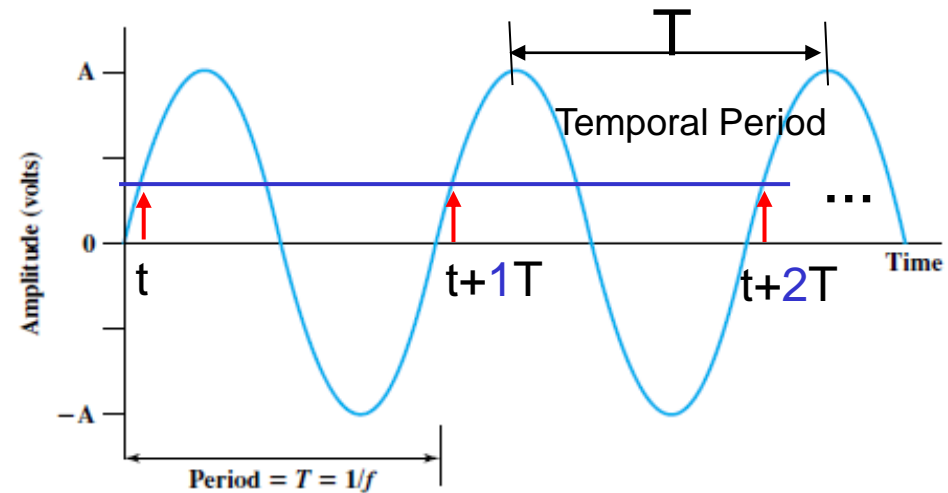
Where:

t is time over first period

T is the waveform period

m is an integer

Signal behavior over **one period** describes behavior at all times



SINE WAVE

□ $s(t) = A \sin(2\pi ft + \phi)$

□ A: Peak Amplitude

- Peak strength of signal, volts

□ T: Temporal (time) Period, (s)

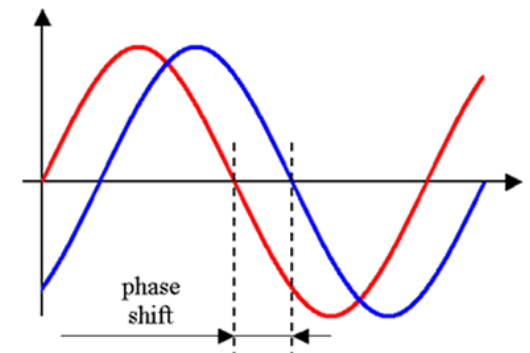
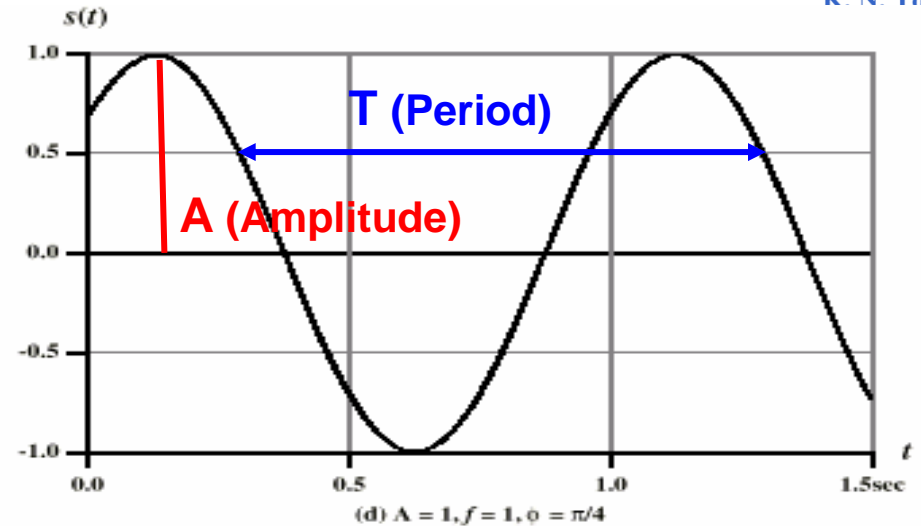
□ f: Repetition Frequency, (Cycles/s) (Hz)

- Measures how fast the signal varies with time
- $f = 1 / T$

□ ω : Angular Frequency, (Radians/s)

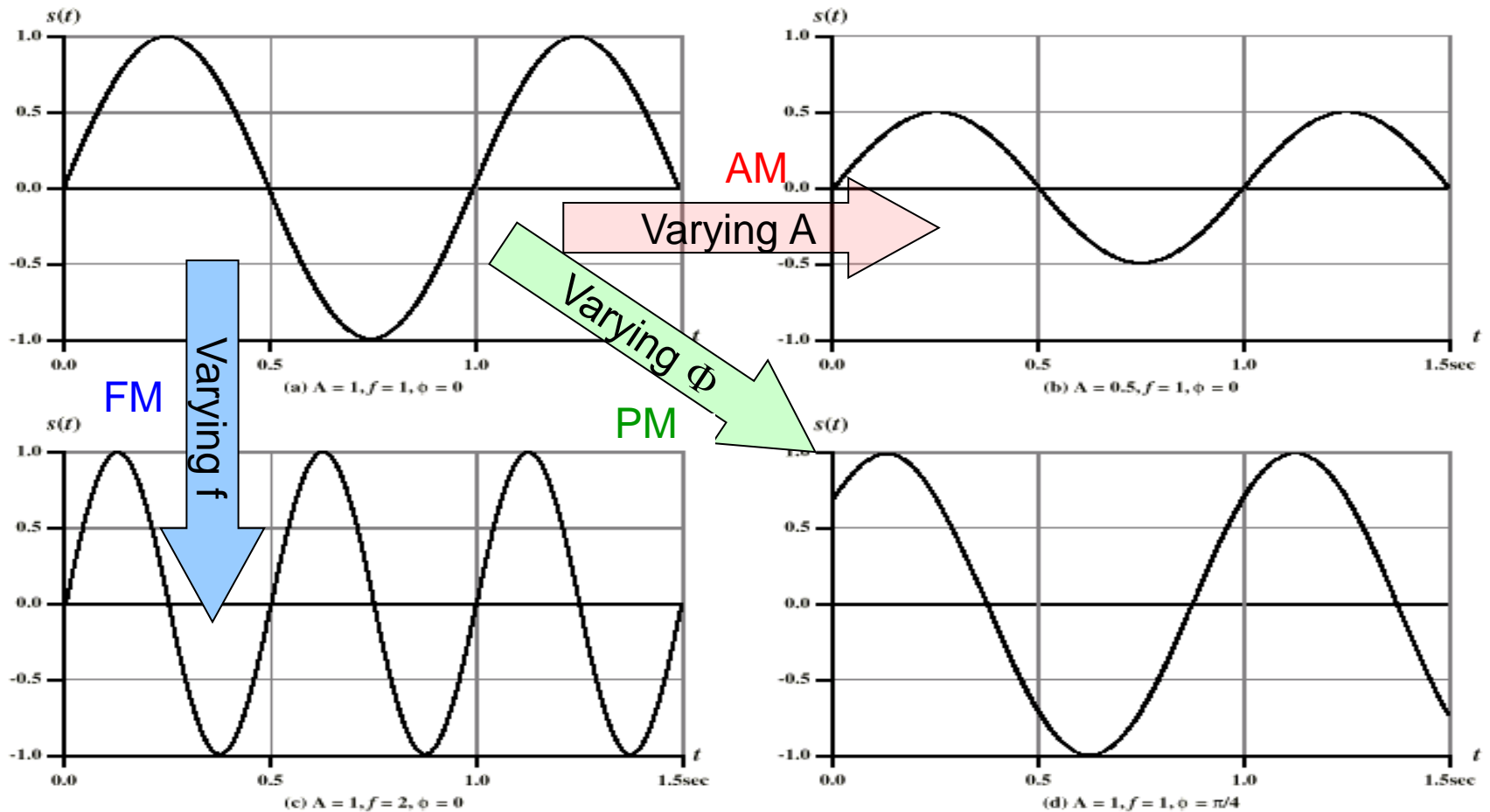
$$\omega = 2\pi f = 2\pi / T$$

□ ϕ : Phase Angle, (Radians)



CONVEYING INFORMATION BY A SINE WAVE CARRIER

$$s(t) = A \sin(2\pi ft + \Phi)$$



WAVELENGTH

□ λ (meters)

□ Spatial period of the wave:

- distance between two points in space on the wave propagation path where the wave has the same total phase

□ Distance traveled by the wave during one temporal (time) cycle

$$d_T = v T$$

$$v = \lambda f$$

$$d_T = (\lambda f) T = \lambda$$

FREQUENCY DOMAIN CONCEPTS

- ❑ Response of systems to a **sine waves** is easy to analyze
- ❑ **Fourier analysis** shows that signals can be treated as
 - Sum of many sine wave components
 - having different frequencies, amplitudes, and phases
- ❑ Basis for **frequency domain analysis**
- ❑ Dealing with functions in the frequency domain is simpler

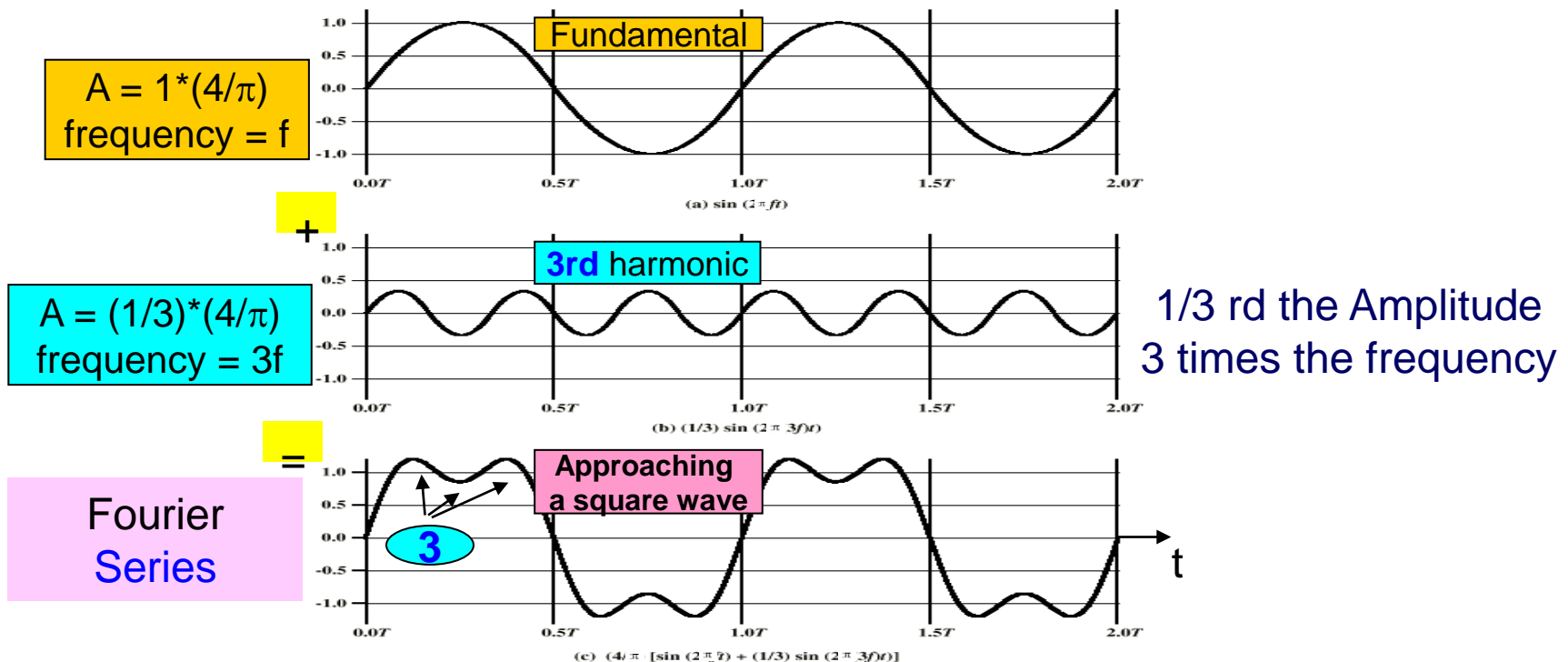
FREQUENCY DOMAIN REPRESENTATION

$$S(t) = (4/\pi f) [(\sin(2\pi f t) + (1/3)(\sin(2\pi(3f)t))]$$

$3f$:integer multiple of the first frequency , f , known as **fundamental frequency**

The **period** of the **total signal** is equal to the period of the **fundamental frequency**

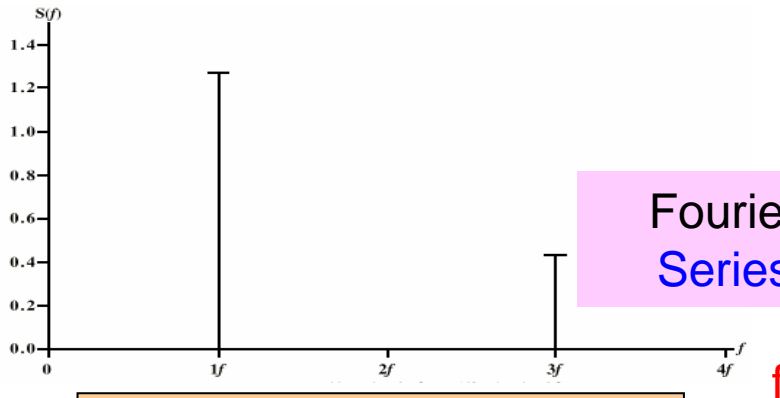
The period of the component $\sin(2\pi f t)$ is T and the period of $s(t)$ is also T .



1/3 rd the Amplitude
3 times the frequency

SIGNAL SPECTRUM

Frequency Spectrum



Frequency Domain: $S(f)$ vs f

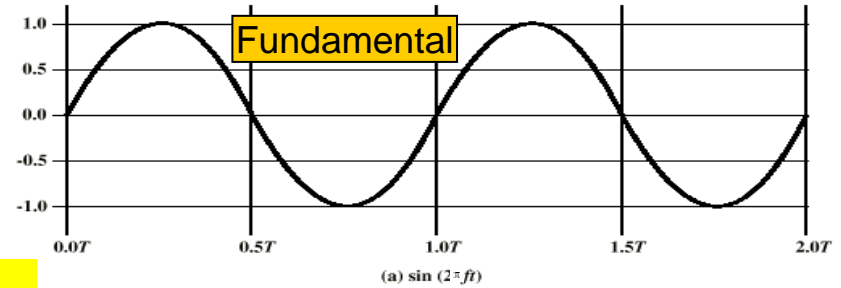
Discrete Function in f

Fourier Series

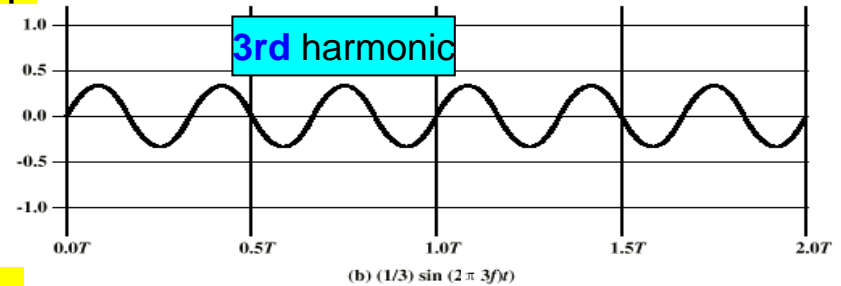


Time Domain: $s(t)$ vs t

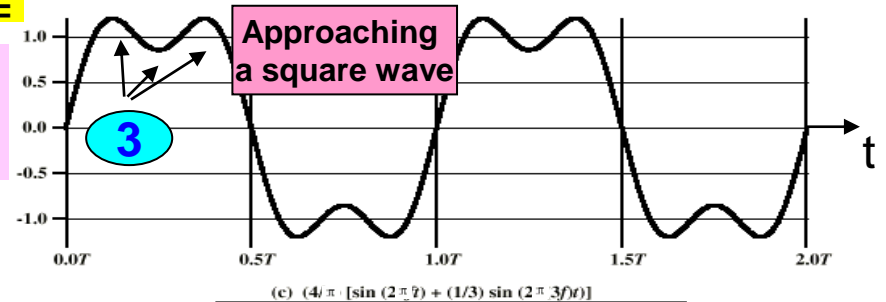
Periodic function in t



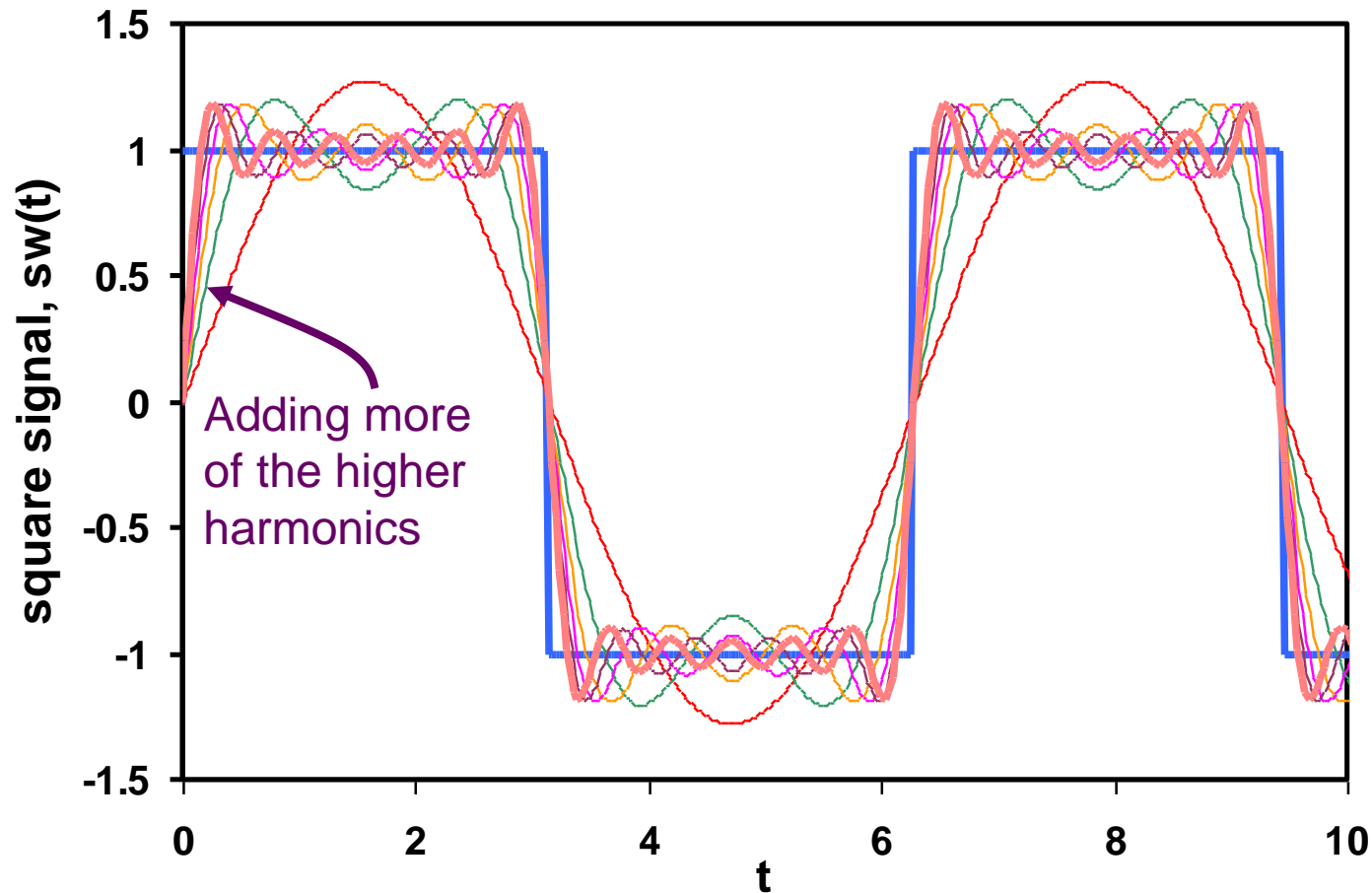
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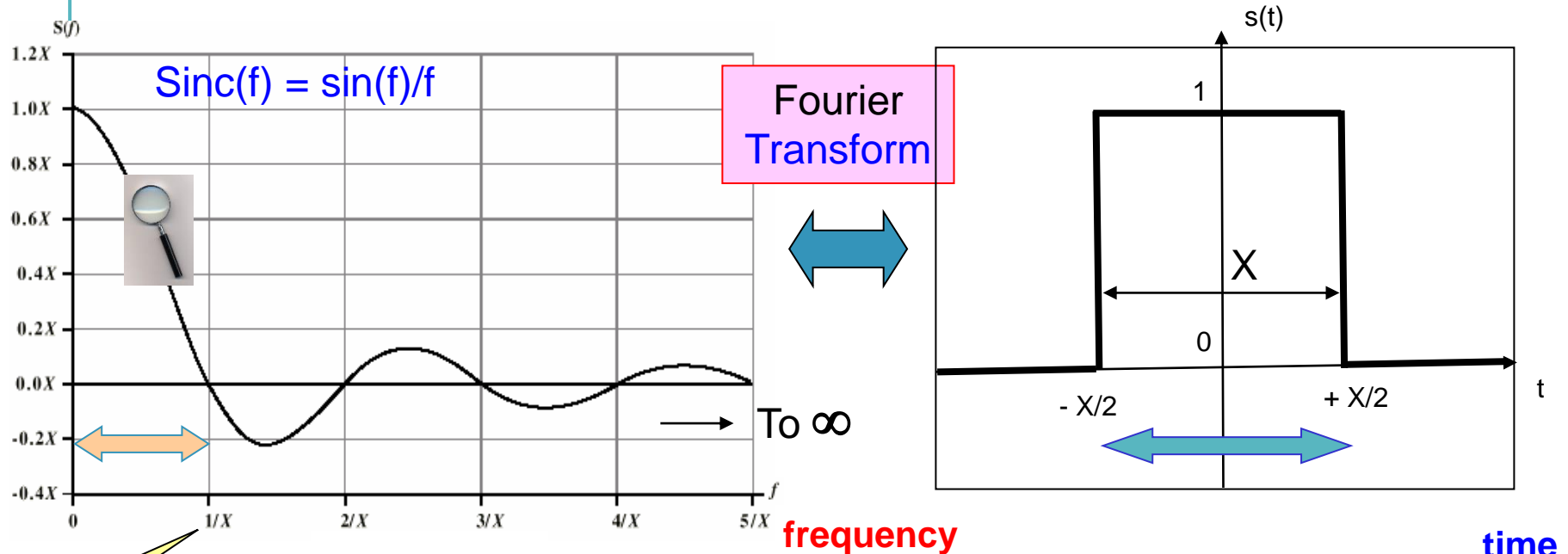


ASYMPTOTICALLY APPROACHING A SQUARE WAVE



Combining the fundamental + an infinite number of **odd** harmonics at proper amplitudes

SPECTRUM OF A SINGLE SQUARE PULSE



$1/X$

Frequency Domain: $S(f)$ vs f

Continuous Function in f

Fourier Transform

Time Domain: $s(t)$ vs t

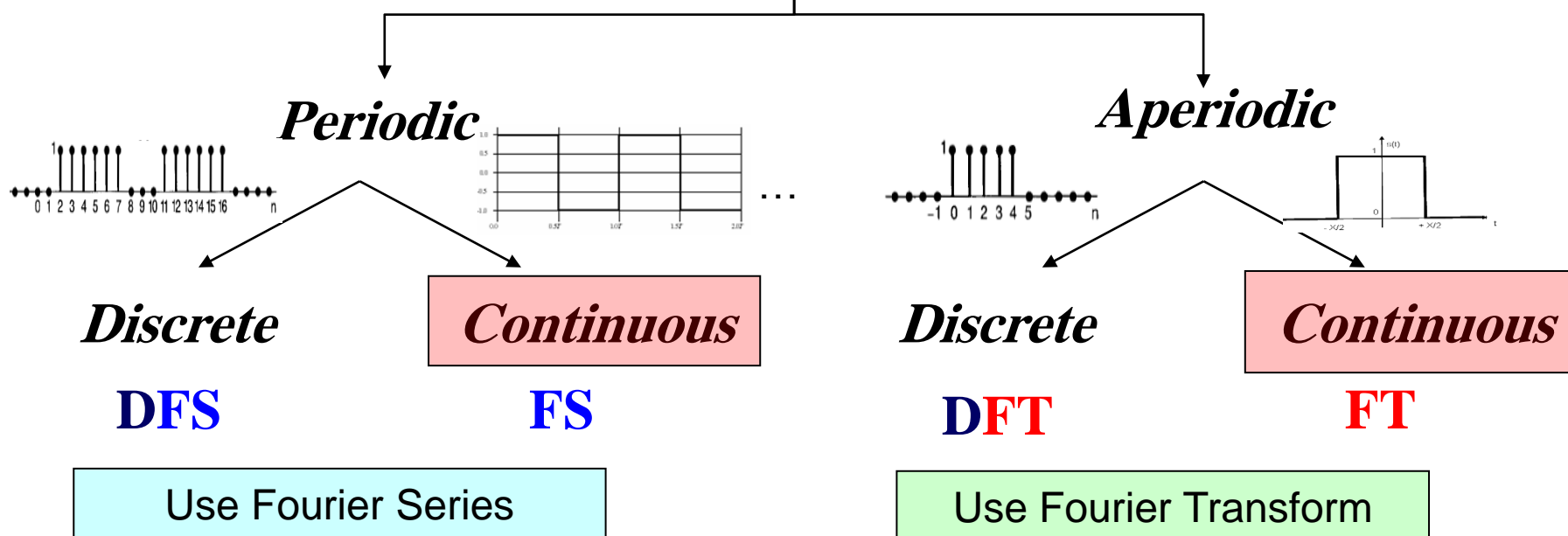
Aperiodic function in t

• What happens to the spectrum as the pulse gets broader ... \rightarrow DC ?

• What happens to the spectrum as the pulse gets narrower ... \rightarrow spike ?

FOURIER ANALYSIS

Signals in Time



- FS** : Fourier Series
- DFS** : Discrete Fourier Series
- FT** : Fourier Transform
- DFT** : Discrete Fourier Transform

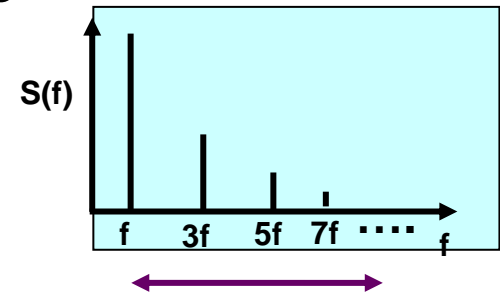
SPECTRUM & BANDWIDTH OF A SIGNAL

□ Spectrum of a signal

- Range of frequencies contained in a signal

□ Absolute (theoretical) Bandwidth (BW):

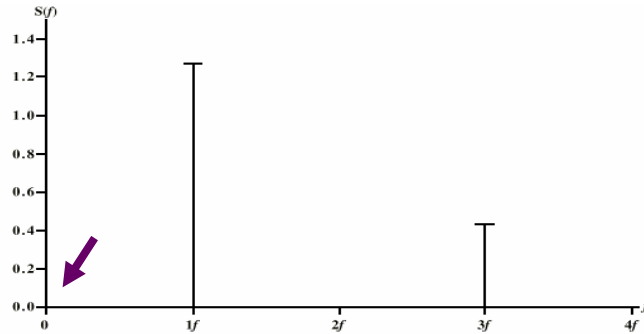
- Full width of spectrum = $f_{\max} - f_{\min}$
- But in many situations, $f_{\max} = \infty$!



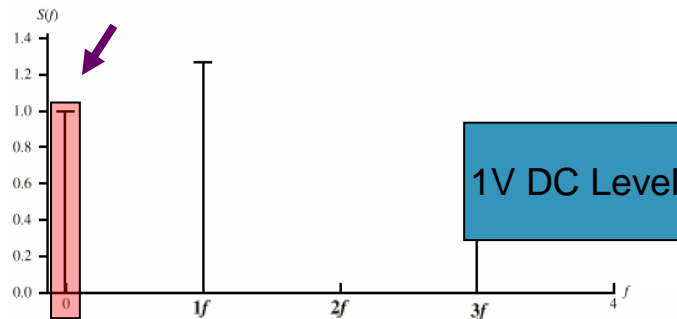
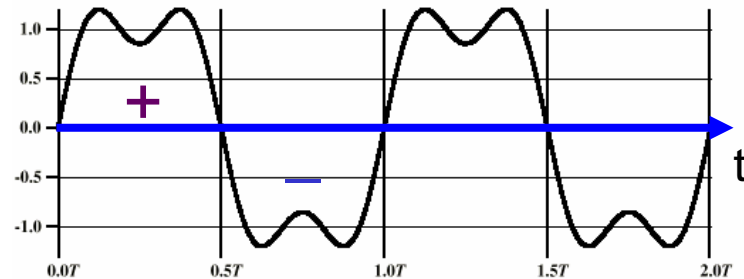
□ Effective Bandwidth

- Often called just *bandwidth*
- Narrow band of frequencies containing *most* of the signal energy
- e.g. that contains say 95% of the energy of the signal

SIGNALS WITH A DC COMPONENT

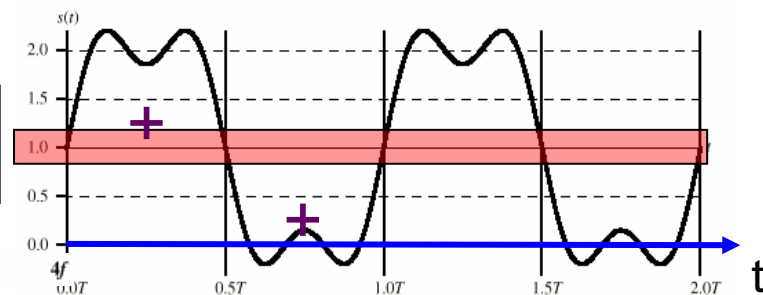


NO DC Component,
Signal average over a period = 0



1V DC
Component

1V DC Level

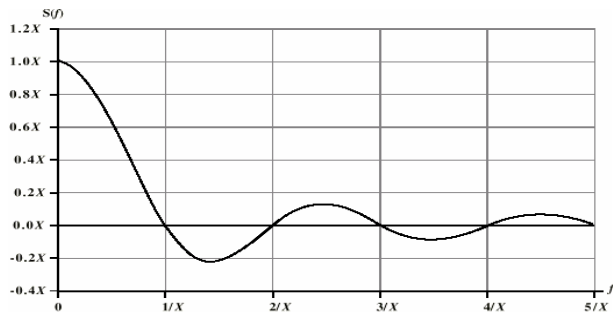
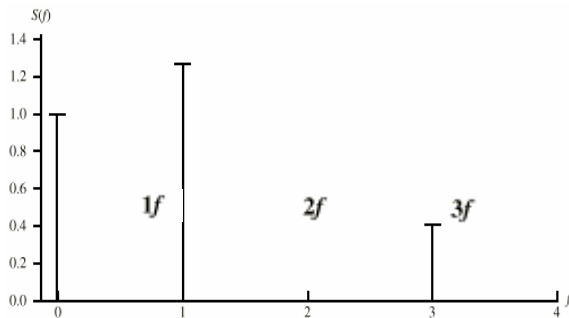
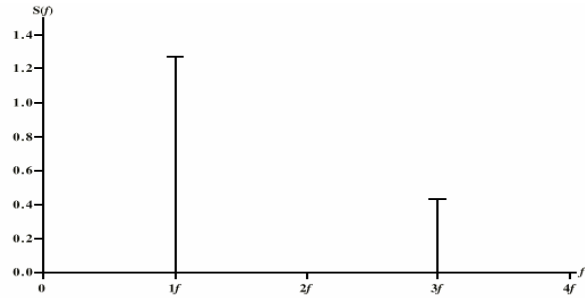


DC Component:

Component at zero frequency

Determines if $f_{min} = 0$ or not

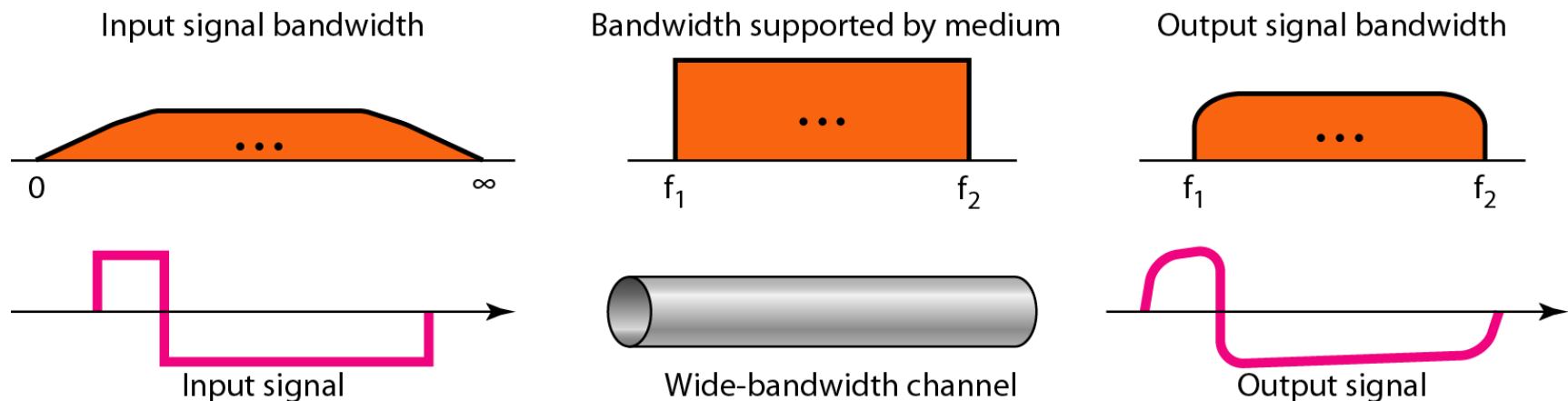
BANDWIDTH FOR THESE SIGNALS



f_{min}	f_{max}	Absolute BW	Effective BW
$1f$	$3f$	$2f$	$2f$
0	$3f$	$3f$	$3f$
0	∞	∞	$1/X ?$

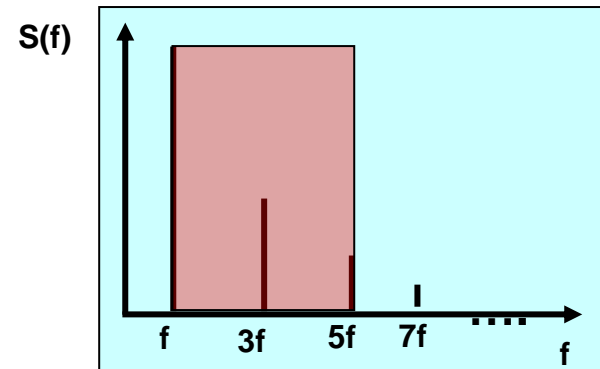
BANDWIDTH OF A TRANSMISSION SYSTEM

- ❑ Range of signal frequencies that are adequately passed by the system
- ❑ Transmission system (TX, medium, RX) acts as a filter
 - Poor transmission media, e.g. twisted pairs, have a narrow bandwidth



BANDWIDTH OF A TRANSMISSION SYSTEM

- Narrow media bandwidth effectively cuts off higher frequency signal components
 - poor signal quality at receiver
 - limiting the signal frequencies (Hz) that can be used for transmission
 - limiting the data rates that can be used (bps), examples:
 - Twisted pair: 4 KHz BW → 100 Kbps
 - Optical fiber: 4 THz BW → 10 Gbps



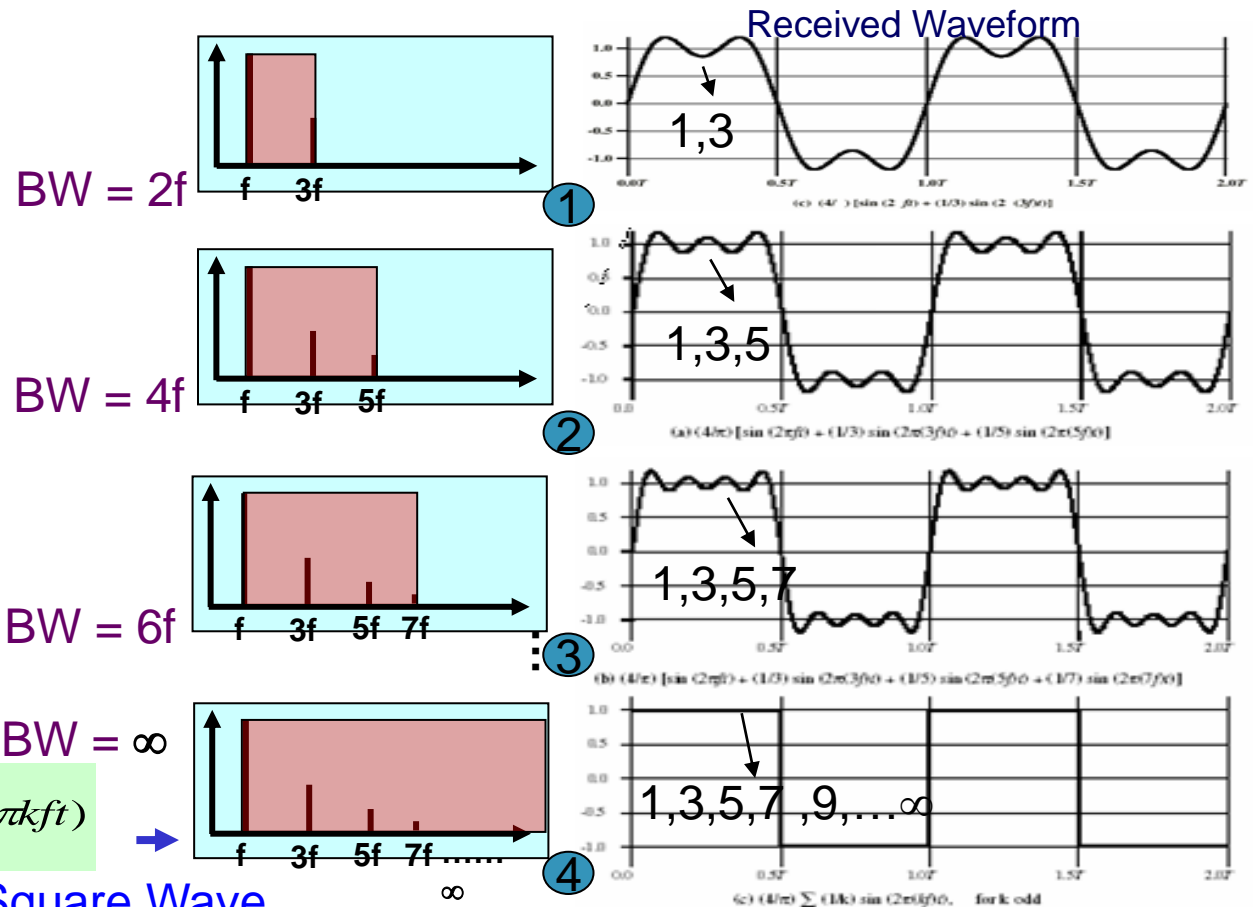
LIMITING EFFECT OF SYSTEM BANDWIDTH

Varying System BW

Better reception requires larger BW

$$s(t) = \frac{4}{\pi} \sum_{k \text{ odd}, k=1}^{\infty} \frac{1}{k} \sin(2\pi kft)$$

Fourier Series for a Square Wave



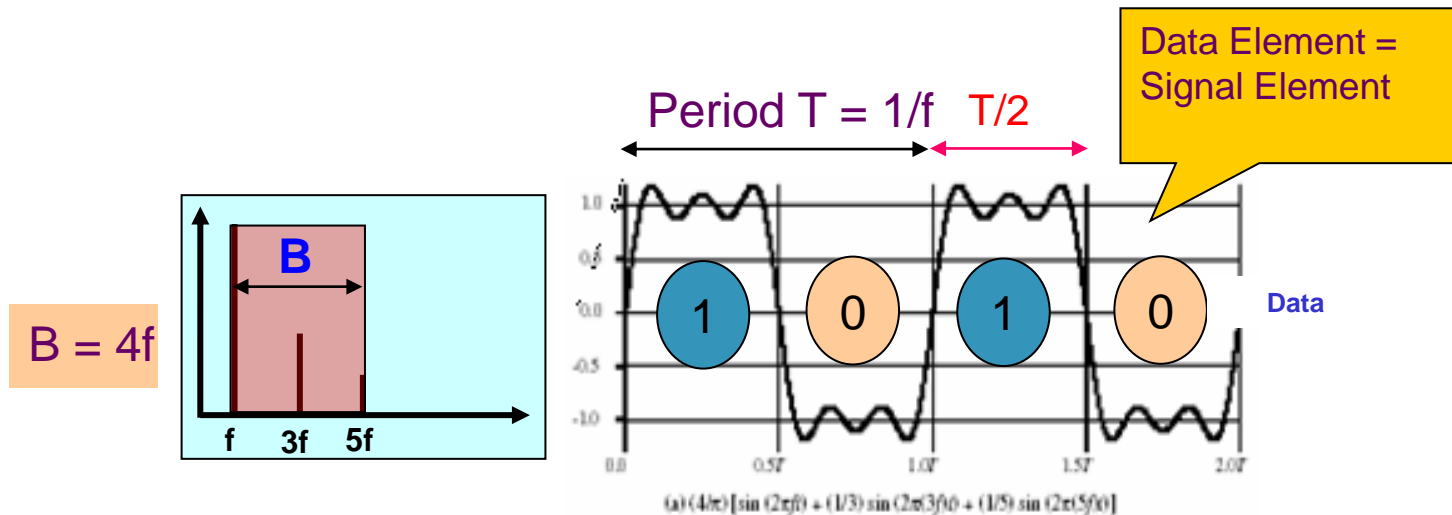
SYSTEM BANDWIDTH AND ACHIEVABLE DATA RATES

- ❑ Any transmission system supports only a limited range of frequencies (bandwidth) for satisfactory transmission
- ❑ For example, this bandwidth is largest for expensive optical fibers and smallest for cheap twisted pair wires
- ❑ Bandwidth is money → Economize in its use

SYSTEM BANDWIDTH AND ACHIEVABLE DATA RATES

- ❑ Limited system bandwidth degrades higher frequency components of the signal transmitted
 - ⇒ poorer received waveforms
 - ⇒ more difficult to interpret the signal at the receiver (especially with noise)
 - ⇒ Data Errors
- ❑ More degradation occurs when higher data rates are used (signal will have higher frequency components)

BANDWIDTH AND DATA RATES



1- Data rate relates to f

$$\text{Data rate} = 1/(T/2) = (2/T) \text{ (bps)} = 2f \text{ (bps)}$$

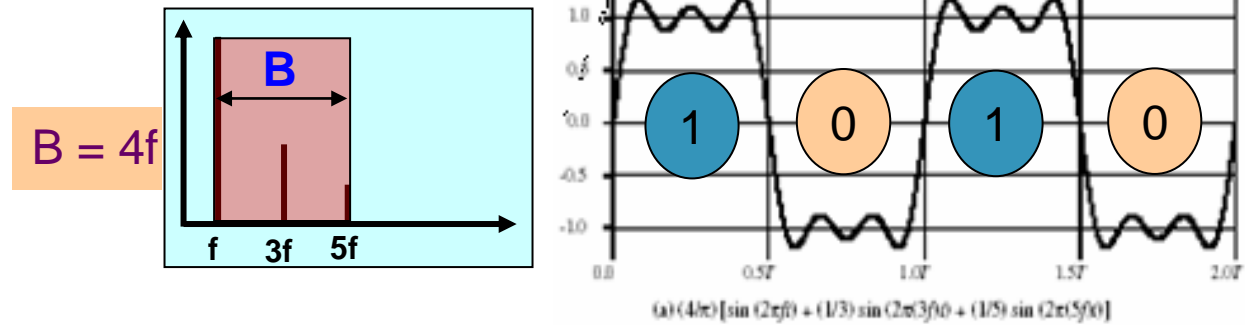
2- Bandwidth relates to quality of Fourier harmonics in the signal

e.g. harmonics $1f, 3f, 5f$: $B=4f$

From 1 and 2:

$$\text{Given a bandwidth } B, \text{ Data rate} = 2f = B/2$$

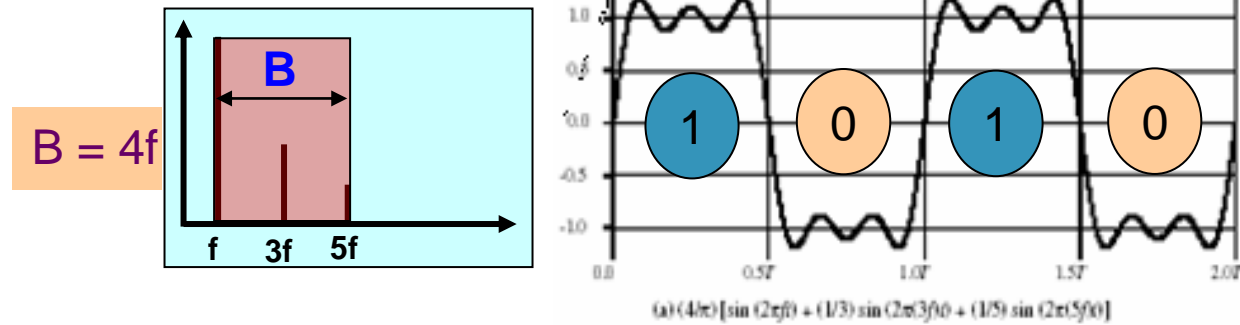
BANDWIDTH AND DATA RATES EXAMPLE



$f=1\text{Mhz}$, Data rate=?

$f=2\text{MHz}$, Data rate= ?

BANDWIDTH AND DATA RATES EXAMPLE



$f=1\text{MHz}$, bit time =.5 micro sec, Data rate=2Mbps

$f=2\text{MHz}$, , bit time =.25 micro sec, Data rate= 4Mbps

DOUBLING THE DATA RATE

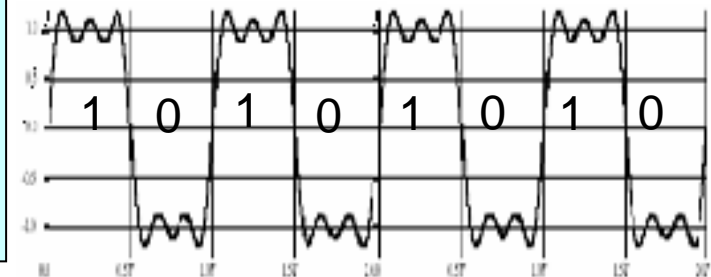
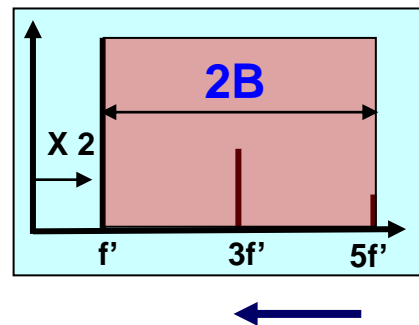
To double the data rate you need to double f : $f' = 2f$
Two ways to do this...

1. Double the bandwidth,
same received waveform quality (same RX conditions & error rate)

$$\text{New BW} = B' = 4f' = 8f = 2B$$

$$\text{Data rate} = 2f' = 2(2f) = 4f = B$$

e.g.
 $f' = 2\text{MHz}$
 $2B = 8\text{MHz}$
Data rate = 4Mbps



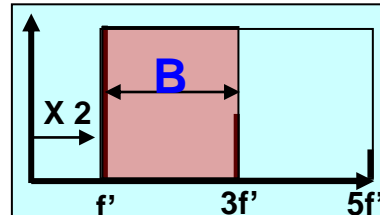
DOUBLING THE DATA RATE

2. **Same** bandwidth, B , but tolerate poorer received waveform
(needs better receiver, higher S/N ratio, or tolerating more errors in data)

Bandwidth : $B = 4f = 2f'$

Data rate = $2f' = 2(2f) = 4f = B$

$B = 2f'$

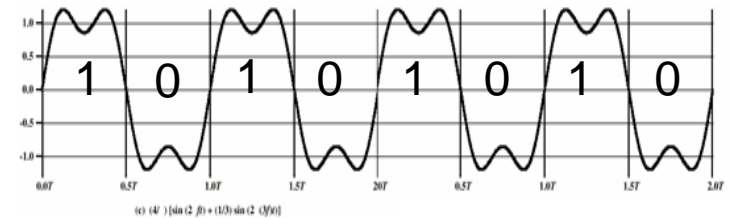


e.g.

$f' = 2\text{MHz}$,

$B = 4\text{MHz}$

Data rate = 4Mbps



MAXIMUM DATA RATE (LINK CHANNEL CAPACITY)

- ☐ Bandwidth of transmission system
- ☐ Signal to noise ratio (SNR)
- ☐ Receiver type
- ☐ Acceptable error performance

BANDWIDTH & DATA RATES: TRADEOFFS...

- ❑ Increasing the data rate (bps) while keeping BW the same (to economize)
- ❑ Means working with inferior (poorer) waveforms at the receiver, which may require:
 - Ensuring higher signal to noise ratio (SNR) at RX
 - More sensitive (& costly!) receiver
 - Suffering from higher bit error rates
 - Tolerate them?
 - Add more efficient means for error detection and correction

CHANNEL CAPACITY

❑ **Channel capacity:** Maximum data rate usable under a given set of communication conditions

- Max rate at which data can be communicated on the channel, bits per second (bps)

❑ How channel BW (B), signal level, noise and impairments, and the amount of data error that can be tolerated limit the channel capacity?

❑ In general, Max possible data rate, C , on a given channel

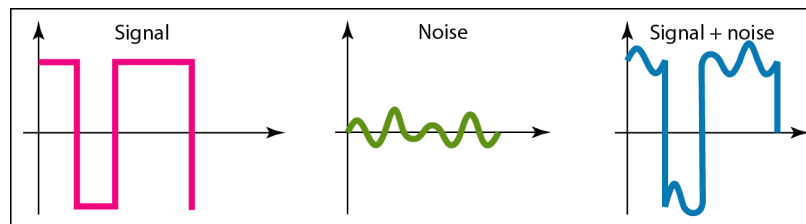
Function (B , Signal to noise, Bit error rate allowed)

CHANNEL CAPACITY

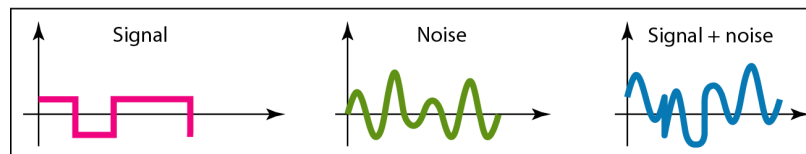
❑ **Bandwidth:** BW of the transmitted signal as constrained by the transmission system, (Hz)

❑ **Signal relative to Noise,**

- $\text{SNR} = \text{signal power} / \text{noise power ratio}$
- (Higher SNR \rightarrow better communication conditions \rightarrow higher C)



a. Large SNR



b. Small SNR

CHANNEL CAPACITY

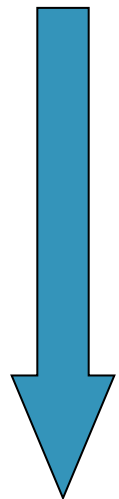
- ❑ **Bandwidth:** BW of the transmitted signal **as constrained** by the transmission system, (Hz)
- ❑ **Signal relative to Noise,**
 - $\text{SNR} = \text{signal power} / \text{noise power ratio}$
 - (Higher SNR \rightarrow better communication conditions \rightarrow higher C)
- ❑ **Bit error rate (BER) allowed:**
 - in $(\text{bits received in error}) / (\text{total bits transmitted})$
 - Equal to the **bit error probability**
 - Higher allowed \rightarrow higher usable data rates \rightarrow higher C

CHANNEL CAPACITY, C:

$$C \text{ (bps)} = F(B, \text{SNR}, \text{BER})$$

Three Formulations under different assumptions:

Idealistic



Assumptions	Formulation
Ideal: Noise-free, Error-free: $C = F(B)$	Nyquist
Noisy, Error-free: $C = F(B, \text{SNR})$	Shannon
Practical: Noisy, Error: $C = F(B, \text{SNR}, \text{BER})$	E_b/N_0 Vs Error Rate

Realistic

BANDWIDTH (OR SPECTRAL) EFFICIENCY (BE):

- ❑ Measures how well we are utilizing a given bandwidth to send data at a high rate....
- ❑ Can be greater than 1 (not like engineering efficiencies)
- ❑ The larger the better

$$BE = \frac{\text{Channel Capacity } C}{\text{Bandwidth } B}, \quad \text{bps} / \text{Hz}$$

1. NYQUIST CHANNEL CAPACITY: (NOISE-FREE, ERROR-FREE)

- ❑ Idealized, theoretical
 - Assumes a noise-free → error-free channel
- ❑ If rate of **signal** transmission is $2B$ then a signal with frequency components up to B Hz is sufficient to carry that **signalling rate**
- ❑ Given bandwidth B , **highest signalling rate** possible is $2B$ **signal elements/s**
- ❑ How much data rate does this represent?
(depends on how many bits are represented by each signal element!)
 - Given a binary signal (1,0), data rate is **same as** signal rate → Data rate supported by a BW of B Hz is $2B$ bps → **$C = 2B$**
 - For the same B , data rate can be increased by sending one of M different signals (symbols): as each signal level now represents $\log_2 M$ bits
- ❑ Generalized **Nyquist Channel Capacity**, **$C = 2B \log_2 M$ (bps)**
- ❑ Bandwidth efficiency = $C/B = 2 \log_2 M$ (bits/s)/Hz : Dimensionless quantity

NYQUIST BANDWIDTH: EXAMPLE

□ $C = 2B \log_2 M$ bits/s

- C = Nyquist Channel Capacity
- B = Bandwidth
- M = Number of discrete signal levels (symbols) used

□ Data on telephone Channel:
 $3400 - 300 = 3100$ Hz

□ With a **binary signal** ($M = 2$ symbols, e.g. 2 amplitudes):

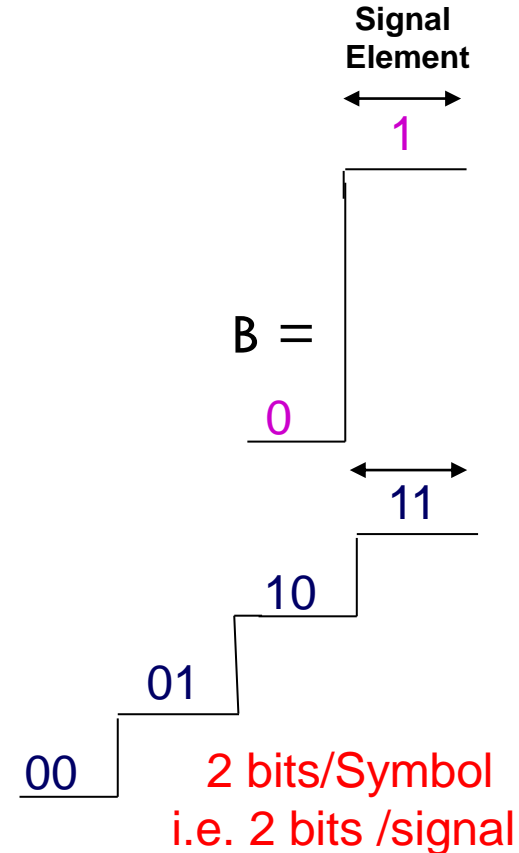
$$C = 2B \log_2 2 = 2B \times 1 = 6200 \text{ bps}$$

□ With a **quadrary signal** ($M = 4$ symbols):

$$C = 2B \log_2 4 = 2B \times 2 = 4B = 12,400 \text{ bps}$$

□ Channel capacity increased, but

disadvantage: Larger number of signal levels (M) makes it more difficult for the receiver to determine data correctly in the presence of noise



2. SHANNON CAPACITY FORMULA: (NOISY, ERROR-FREE)

- Highest **error-free** data rate in the presence of noise
- Signal to noise ratio $SNR = \text{signal} / \text{noise levels}$
 $SNR_{dB} = 10 \log_{10} (SNR \text{ ratio})$
- Errors are less likely with lower noise (larger SNR ratios). This allows higher **error-free** data rates i.e. larger Shannon channel capacities
- Shannon Capacity $C = B \log_2(1+SNR)$:

Highest data rate transmitted **error-free** with a given noise level

- For a given BW, the larger the SNR the higher the data rate I can use **without introducing errors**
- C/B : Spectral (bandwidth) efficiency, BE, (bps/Hz) (> 1)
- Larger BEs mean better utilization of a given bandwidth B for transmitting data fast.

SHANNON CAPACITY FORMULA: COMMENTS

❑ Formula says: for data rates \leq calculated C , it is theoretically possible to find an **encoding scheme** that achieves error-free transmission at the given SNR... **But it does not say how!**

Also:

❑ It is a theoretical approach based on **thermal (white) noise only**. But in practice, we also have impulse noise, attenuation and delay distortions, etc...

- So, maximum error-free data rates **measured** in practice are expected to be **lower** than the C predicted by the Shannon formula due to the greater noise

❑ However, maximum error-free data rates can be used to compare practical systems: The higher that rate the better the system...

SHANNON CAPACITY FORMULA: COMMENTS CONTD.

□ Formula suggests that changes in B and SNR can be done **arbitrarily** and **independently**... but

→ In practice, this may not be the case!

- Higher SNR obtained through excessive amplification may also introduce nonlinearities
→ increased distortion and inter-modulation **noise** ... which **reduces SNR**!
- High Bandwidth B opens the system up for more thermal noise (kTB), and therefore **reduces SNR**!

SHANNON CAPACITY FORMULA: EXAMPLE

- Spectrum of communication channel extends from 3 MHz to 4 MHz
- $\text{SNR} = 24\text{dB}$
- Then $B = 4\text{MHz} - 3\text{MHz} = 1\text{MHz}$
 $\text{SNR}_{\text{dB}} = 24\text{dB} = 10 \log_{10} (\text{SNR})$
 $\text{SNR (ratio)} = \log_{10}^{-1} (24/10) = 10^{24/10} = 251$
- Using Shannon's formula: $C = B \log_2 (1 + \text{SNR})$
 $C = 10^6 * \log_2 (1 + 251) \sim 10^6 * 8 = 8 \text{ Mbps}$
- Based on Nyquist's formula, determine M that gives the above channel capacity:
 $C = 2B \log_2 M$
 $8 * 10^6 = 2 * (10^6) * \log_2 M$
 $4 = \log_2 M$
 $M = 16$

3. E_b/N_0 VS ERROR RATE FORMULATION (NOISE AND ERROR ARE BOTH SPECIFIED **TOGETHER**)

- Handling both noise and a quantified error rate **simultaneously**
- We introduce E_b/N_0 : A standard **quality measure** of three channel parameters (B, SNR, R) and can **also** be independently related to the **error rate**

3. E_b/N_0 VS ERROR RATE FORMULATION (NOISE AND ERROR ARE BOTH SPECIFIED **TOGETHER**)

□ It expresses SNR in a manner related to the data rate, R

▪ E_b = Signal **energy** in one bit interval (Joules)

= Signal **power** (Watts) x bit interval T_b (second)

= $S \times (1/R) = S/R$

▪ N_0 = Noise power (watts) in 1 Hz = kT . **Two formulations:**

$$\frac{E_b}{N_0} = \frac{ST_b}{N_0} = \frac{S/R}{kT} = \frac{S}{kTR} \quad T_b = 1/R$$

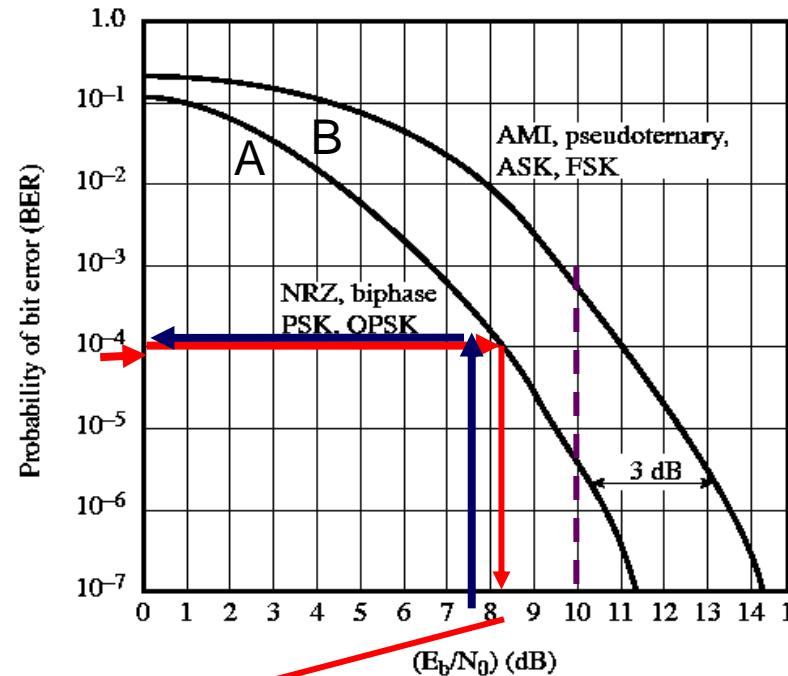
$$\frac{E_b}{N_0} = \frac{S/R}{N_0} = \frac{S}{N} \frac{B}{R} = SNR \left(\frac{B}{R} \right) \longrightarrow = \text{SNR/BE}$$

BER vs E_b/N_0 curve for a given encoding scheme

BER CURVE

Which encoding scheme is better: A or B?

Lower Error Rate: larger E_b/N_0



$$\left(\frac{E_b}{N_0} \right)_{dB} = S_{dBW} - 10 \log R - 10 \log k - 10 \log T$$

$$= S_{dBW} - 10 \log R + 228.6 \text{ dBW} - 10 \log T$$

$$\frac{E_b}{N_0} = \frac{S / R}{N_0} = \frac{S}{N} \frac{B}{R} = SNR \left(\frac{B}{R} \right) = \frac{SNR}{BE}$$

$$\text{Max } R = C, BE = C/B$$

EXAMPLE:

$$\left(\frac{E_b}{N_0} \right)_{dB} = S_{dBW} - 10 \log R - 10 \log k - 10 \log T$$

$$= S_{dBW} - 10 \log R + 228.6 \text{ dBW} - 10 \log T$$

- Given:
- The effective noise temperature, T , is 290°K
- The data rate, R , is 2400 bps
- Would like to operate with a bit error rate of 10^{-4} (e.g. 1 error in 10^4 bits)

What is the minimum signal level required for the received signal?

- From curve,
a minimum E_b/N_0 needed to achieve a bit error rate of $10^{-4} = 8.4 \text{ dB}$

$$8.4 = S(\text{dBW}) - 10 \log 2400 + 228.6 \text{ dBW} - 10 \log 290$$

$$= S(\text{dBW}) - (10)(3.38) + 228.6 - (10)(2.46)$$

$$S = -161.8 \text{ dBW}$$

E_b/N_0 IN TERMS OF BE, ASSUMING SHANNON CHANNEL CAPACITY

From Shannon's formula:

$$C = B \log_2(1 + \text{SNR})$$

We have:

$$\text{SNR} = (2^{C/B} - 1) = (2^{BE} - 1)$$

From the E_b/N_0 formula:

$$\frac{E_b}{N_0} = \frac{\text{SNR}}{BE} = \frac{1}{BE} (2^{BE} - 1)$$

C/B (bps/Hz) is the spectral (bandwidth) efficiency BE based on Shannon channel capacity

EXAMPLE

Find the minimum E_b/N_0 required to achieve a Shannon bandwidth efficiency ($BE = C_{\text{Shannon}}/B$) of 6 bps/Hz:

$$\frac{E_b}{N_0} = \frac{1}{BE} (2^{BE} - 1)$$

Substituting in the equation above:

$$E_b/N_0 = (1/6) (2^6 - 1) = 10.5 = 10.21 \text{ dB}$$

SNR is mostly applicable to analog signal and E_b/N_0 digital signal, where we are dealing with bits