

Problem 5.21. Let $AMBIG_{CFG} = \{\langle G \rangle \mid G \text{ is an ambiguous CFG}\}$. Show that $AMBIG_{CFG}$ is undecidable.

Proof. To show that $AMBIG_{CFG}$ is undecidable, we give a reduction from PCP to $AMBIG_{CFG}$. Given an instance

$$P = \left\{ \left[\frac{t_1}{b_1} \right], \left[\frac{t_2}{b_2} \right], \dots, \left[\frac{t_k}{b_k} \right] \right\}$$

of the Post Correspondence Problem, construct a CFG G with the rules

$$\begin{aligned} S &\rightarrow T \mid B \\ T &\rightarrow t_1 T a_1 \mid \dots \mid t_k T a_k \mid t_1 a_1 \mid \dots \mid t_k a_k \\ B &\rightarrow b_1 B a_1 \mid \dots \mid b_k B a_k \mid b_1 a_1 \mid \dots \mid b_k a_k, \end{aligned}$$

where a_1, a_2, \dots, a_k are new terminal symbols¹. Next we show that this reduction works.

If P is an instance of Post Correspondence Problem, then for every possible arrangement of the dominoes i_1, i_2, \dots, i_l , where $t_{i_1} t_{i_2} \dots t_{i_l}$ is the top string and $b_{i_1} b_{i_2} \dots b_{i_l}$ the bottom string, there exists exactly one derivation for the top and bottom strings in G . The terminal symbols a_1, a_2, \dots, a_k , make sure that every string generated using the variable T has at most one left-most derivation even if some t_i and t_j are same. Same argument can be said for the variable B .

If $P \in PCP$, then there exists a match i_1, i_2, \dots, i_l , where $t_{i_1} t_{i_2} \dots t_{i_l} = b_{i_1} b_{i_2} \dots b_{i_l}$. In this case, the CFG G is guaranteed to have following two left-most derivations that produce the same string.

$$\begin{aligned} S &\rightarrow T \\ &\rightarrow t_{i_1} T a_{i_1} \\ &\rightarrow t_{i_1} t_{i_2} T a_{i_2} a_{i_1} \\ &\rightarrow t_{i_1} t_{i_2} \dots T \dots a_{i_2} a_{i_1} \\ &\rightarrow t_{i_1} t_{i_2} \dots t_{i_l} a_{i_l} \dots a_{i_2} a_{i_1} \end{aligned}$$

$$\begin{aligned} S &\rightarrow B \\ &\rightarrow b_{i_1} B a_{i_1} \\ &\rightarrow b_{i_1} b_{i_2} B a_{i_2} a_{i_1} \\ &\rightarrow b_{i_1} b_{i_2} \dots B \dots a_{i_2} a_{i_1} \\ &\rightarrow b_{i_1} b_{i_2} \dots b_{i_l} a_{i_l} \dots a_{i_2} a_{i_1} \end{aligned}$$

□

¹This is a hint mentioned in the book.