

Problem 5.14. Consider the problem of determining whether a Turing machine M on an input w ever attempts to move its head left when its head is on the left-most tape cell. Formulate this problem as a language and show that it is undecidable.

Proof. Let $T = \{\langle M, w \rangle \mid M \text{ is a } \mathbf{TM}, \text{ which on input } w \text{ moves its head left when its head is on the left-most tape cell}\}$.

Show that A_{TM} reduces to T , where $A_{TM} = \{\langle M, w \rangle \mid M \text{ is a } \mathbf{TM} \text{ and } M \text{ accepts } w\}$. Assume for the sake of contradiction that $\mathbf{TM} R$ decides T . Then construct a $\mathbf{TM} S$ that uses R to decide A_{TM} . The idea is to construct a new $\mathbf{TM} M_w$, which only moves its head left when its head is on the left-most tape cell if M accepts w . We construct $\mathbf{TM} M_w$ using 2 tapes as follows.

$M_w =$ “On input w :

1. Simulate M on w using the second tape.
2. If the simulation shows that M accepts, move the head of the first tape left.”

Next, we construct $\mathbf{TM} S$ to decide A_{TM} .

$S =$ “On input $\langle M, w \rangle$, where M is a \mathbf{TM} and w is a string:

1. Use M and w to construct $\mathbf{TM} M_w$ as discussed above.
2. Convert the 2 tape $\mathbf{TM} M_w$ to equivalent single tape $\mathbf{TM} M'$.
3. Run R on $\langle M', w \rangle$.
4. If R accepts, M accepts w , so *accept*. Otherwise, *reject*.”

Thus, if $\mathbf{TM} R$ exists, we can decide A_{TM} , but we know that A_{TM} is undecidable¹. By virtue of this contradiction, we can conclude that R does not exist. Therefore, T is undecidable. \square

¹Theorem 4.11 A_{TM} is undecidable.