**Problem 7.22.** Let  $DOUBLE\text{-}SAT = \{\langle \phi \rangle \mid \phi \text{ has at least two satisfying assignments}\}$ . Show that DOUBLE-SAT is NP-complete.

*Proof.* The proof is in two parts. First, we show that  $DOUBLE\text{-}SAT \in NP$ . Second, we show that every  $A \in P$  is polynomial time reducable to DOUBLE-SAT.

**Part a.** Show that DOUBLE- $SAT \in NP$ .

A certificate is simply two satisfying assignments.

**Part b.** Show that every  $A \in P$  is polynomial time reducable to DOUBLE-SAT.

To show that every  $A \in P$  is polynomial time reducable to DOUBLE-SAT, we give a polynomial time reduction F from SAT to DOUBLE-SAT.

F = "On input  $\langle \phi \rangle$ , where  $\phi$  is a satisfiable Boolean formula:

- 1. Let  $\alpha$  be a new Boolean variable, which is not in  $\phi$ .
- 2. Let  $\phi_2 = \phi \wedge (\alpha \vee \overline{\alpha})$ .
- 3. Output  $\langle \phi_2 \rangle$ ."

If  $\phi$  is satisfiable, then  $\phi_2$  has at least two satisfiable assignments. First, where  $\alpha = T$ , and the second where  $\alpha = F$ . If  $\phi_2 = \phi \wedge (\alpha \vee \overline{\alpha})$  has at least two satisfiable assignments, then  $\phi$  is satisfiable.