

Problem 4.22. Let $PREFIX-FREE_{REG} = \{\langle R \rangle \mid R \text{ is a regular expression and } L(R) \text{ is prefix-free}\}$. Show that $PREFIX-FREE_{REG}$ is decidable. Why does a similar approach fail to show that $PREFIX-FREE_{CFG}$ is decidable?

Proof. We present a **TM** I that decides $PREFIX-FREE_{REG}$.

$I =$ “On input $\langle R \rangle$, where R is a regular expression:

1. Convert R to equivalent NFA N .
2. Construct an NFA M , such that $L(M) = NOPREFIX(N)$ by following the construction given in solution to Problem 1.40 a.
3. Convert M to equivalent DFA D .
4. Test $L(D) = \emptyset$ using the E_{DFA} decider T from Theorem 4.4.
5. If T accepts, *accept*; if T rejects, *reject*.”

□

The class of context-free languages is not closed under $NOPREFIX$ ¹. Therefore, the approach used to construct a decider for $PREFIX-FREE_{REG}$ cannot be used to show that $PREFIX-FREE_{CFG}$ is decidable.

¹Refer to the proof given in solution to Problem 2.41 a.