Problem 1.32. Let

$$\Sigma_3 = \left\{ \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix}, \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix}, \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix}, \dots, \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix} \right\}$$

 $\Sigma_3$  contains all size 3 columns of 0s and 1s. A string of symbols in  $\Sigma_3$  gives three rows of 0s and 1s. Consider each row to be a binary number and let

 $B = \{w \in \Sigma_3^* \mid the \ bottom \ row \ of \ w \ is \ the \ sum \ of \ the \ top \ two \ rows\}$ 

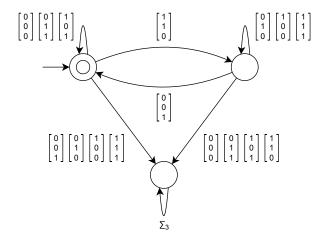
For example,

$$\begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \\ 0 \end{bmatrix} \in B, but \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \\ 1 \end{bmatrix} \not \in B.$$

Show that B is regular.

*Proof Idea.* To prove B is regular, we show that B's reverse  $B^R$  is regular. According to the proof given in Problem 1.31, if  $B^R$  is regular, then its reverse B is also regular.

*Proof.* The proof is by construction. The following state diagram shows the construction of a DFA that recognizes  $B^R$ .



State diagram of a DFA that recognizes  $B^R$ .

Problem 1.33. Let

$$\Sigma_2 = \left\{ \begin{bmatrix} 0 \\ 0 \end{bmatrix}, \begin{bmatrix} 0 \\ 1 \end{bmatrix}, \begin{bmatrix} 1 \\ 0 \end{bmatrix}, \begin{bmatrix} 1 \\ 1 \end{bmatrix} \right\}$$

Here,  $\Sigma_2$  contains all columns of 0s and 1s of height two. A string of symbols in  $\Sigma_2$  gives two rows of 0s and 1s. Consider each row to be a binary number and let

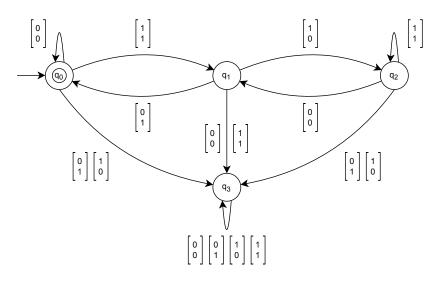
$$C = \{w \, \in \, \Sigma_2^* \, | \, the \, bottom \, row \, of \, w \, is \, three \, times \, the \, top \, row \}$$

For example,

$$\begin{bmatrix} 0 \\ 0 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \end{bmatrix} \begin{bmatrix} 0 \\ 0 \end{bmatrix} \in C, but \begin{bmatrix} 0 \\ 1 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} \not\in C.$$

Show that C is regular.

*Proof Idea.* To prove C is regular, we show that C's reverse  $C^R$  is regular. According to the proof given in Problem 1.31, if  $C^R$  is regular, then its reverse C is also regular.



State diagram of a DFA that recognizes  $C^R$ .

*Proof.* The proof is by construction. The above state diagram shows the construction of a DFA that recognizes  $C^R$ .

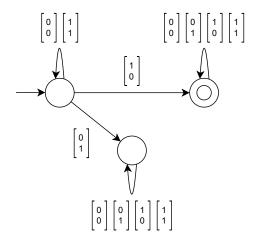
**Problem 1.34.** Let  $\Sigma_2$  be the same as in Problem 1.33. Consider each row to be a binary number and let

$$D = \{ w \in \Sigma_2^* \mid the \ top \ row \ of \ w \ is \ a \ larger \ number \ than \ is \ the \ bottom \ row \}$$

For example,

$$\begin{bmatrix} 0 \\ 0 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \end{bmatrix} \begin{bmatrix} 0 \\ 0 \end{bmatrix} \in D, \ but \begin{bmatrix} 0 \\ 0 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \end{bmatrix} \begin{bmatrix} 0 \\ 0 \end{bmatrix} \not\in D$$

Show that D is regular.



State diagram of a DFA that recognizes D.

*Proof.* The proof is by construction. The given state diagram shows the construction of a DFA that recognizes D.

**Problem 1.35.** Let  $\Sigma_2$  be the same as in Problem 1.33. Consider the top and bottom rows to be strings of 0s and 1s, and let

 $E = \{w \in \Sigma_2^* \mid the \ bottom \ row \ of \ w \ is \ the \ reverse \ of \ the \ top \ row\}$ 

Show that E is not regular.

*Proof.* The proof is by contradiction. Assume that E is regular. Let p be the pumping length given by the pumping lemma. Choose s to be the string:

$$s = \begin{pmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} \end{pmatrix}^p \begin{pmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} \end{pmatrix}^p$$

Therefore,  $s \in E$  and  $|s| \ge p$ , so the pumping lemma guarantees that s can be split into three pieces, s = xyz, where for any  $i \ge 0$ ,  $xy^iz \in E$ . According to condition 3 (i.e.  $|xy| \le p$ ) of the pumping lemma, y can only contain  $\begin{bmatrix} 1 \\ 0 \end{bmatrix}$ s. The string xyyz has more  $\begin{bmatrix} 1 \\ 0 \end{bmatrix}$ s than  $\begin{bmatrix} 0 \\ 1 \end{bmatrix}$ s, so its bottom row is not the reverse of its top row, which is a contradiction. Therefore, E is not regular.

**Problem 1.35.** Let  $\Sigma_2$  be the same as in Problem 1.33. Consider the top and bottom rows to be strings of 0s and 1s, and let

 $E = \{w \in \Sigma_2^* | \text{ the bottom row of } w \text{ is the reverse of the top row of } w\}.$ 

Proof.