Problem 2.27. Let $G = (V, \Sigma, R, \langle STMT \rangle)$ be the following grammar.

$$\langle \text{STMT} \rangle \rightarrow \langle \text{ASSIGN} \rangle \mid \langle \text{IF-THEN} \rangle \mid \langle \text{IF-THEN-ELSE} \rangle$$

$$\langle \text{IF-THEN} \rangle \rightarrow \text{if condition then } \langle \text{STMT} \rangle$$

$$\langle \text{IF-THEN-ELSE} \rangle \rightarrow \text{if condition then } \langle \text{STMT} \rangle \text{ else } \langle \text{STMT} \rangle$$

$$\langle \text{ASSIGN} \rangle \rightarrow \text{a:=1}$$

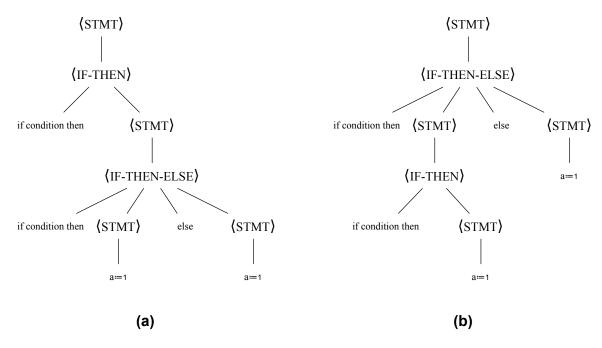
$$\Sigma = \{ \text{if, condition, then, else, a:=1} \}$$

$$V = \{ \langle \text{STMT} \rangle, \langle \text{IF-THEN} \rangle, \langle \text{IF-THEN-ELSE} \rangle, \langle \text{ASSIGN} \rangle \}$$

G is a natural-looking grammar for a fragment of a programming language, but G is ambiguous.

Part a. Show that G is ambiguous.

The string "if condition then if condition then a:=1 else a:=1" has two different parse trees in G.



Two different parse trees of the string "if condition then if condition then a:=1 else a:=1".

Part b. Give a new unambiguous grammar for the same language.

$$\begin{split} \langle \text{STMT} \rangle &\to \langle \text{ASSIGN} \rangle \mid \langle \text{IF-THEN} \rangle \\ \langle \text{IF-THEN} \rangle &\to \text{if condition then } \langle \text{IF-BODY} \rangle \\ \langle \text{IF-BODY} \rangle &\to \langle \text{ASSIGN} \rangle \mid \langle \text{IF-THEN} \rangle \mid \langle \text{ASSIGN} \rangle \; \langle \text{ELSE} \rangle \\ \langle \text{ELSE} \rangle &\to \text{else } \langle \text{ELSE-BODY} \rangle \\ \langle \text{ELSE-BODY} \rangle &\to \langle \text{ASSIGN} \rangle \mid \langle \text{IF-THEN} \rangle \\ \langle \text{ASSIGN} \rangle &\to \text{a}{:}{=}1 \end{split}$$