

**Problem 7.45.** Modify the algorithm for context-free language recognition in the proof of Theorem 7.16 to give a polynomial time algorithm that produces a parse tree for a string, given the string and a CFG, if that grammar generates the string.

*Solution.* Let  $G$  be a CFG in Chomsky normal form generating the CFL  $L$ . Assume that  $S$  is the start variable. The following algorithm  $D'$  produces a parse tree for a string in polynomial time, given the string and a CFG.

$D' =$  “On input  $w = w_1 \cdots w_n$ :

1. For  $w = \varepsilon$ , if  $S \rightarrow \varepsilon$  is a rule, *accept*; else, *reject*. [[ $w = \varepsilon$  case]]
2. For  $i = 1$  to  $n$ : [[ examine each substring of length 1 ]]
3.     For each variable  $A$ :
4.         Test whether  $A \rightarrow b$  is a rule, where  $b = w_i$ .
5.         If so, place  $A$  in  $table(i, i)$ .
6. For  $l = 2$  to  $n$ : [[  $l$  is the length of the substring ]]
7.     For  $i = 1$  to  $n - l + 1$ : [[  $i$  is the start position of the substring ]]
8.         Let  $j = i + l - 1$ . [[  $j$  is the end position of the substring ]]
9.         For  $k = i$  to  $j - 1$ : [[  $k$  is the split position ]]
10.         For each rule  $A \rightarrow BC$ :
11.             If  $table(i, k)$  contains  $B$  and  $table(k + 1, j)$  contains  $C$ ,  
               put  $A$  in  $table(i, j)$ .
12. If  $S$  is in  $table(1, n)$ , then:
- 13.
14. Otherwise, output  $\varepsilon$ .”

□