

**Problem 7.45.** Modify the algorithm for context-free language recognition in the proof of Theorem 7.16 to give a polynomial time algorithm that produces a parse tree for a string, given the string and a CFG, if that grammar generates the string.

*Solution.* Let  $G$  be a CFG in Chomsky normal form generating the CFL  $L$ . Assume that  $S$  is the start variable. The following algorithm  $D'$  produces a parse tree for a string in polynomial time, given the string and a CFG.

$D'$  = “On input  $w = w_1 \cdots w_n$ :

1. For  $w = \varepsilon$ , if  $S \rightarrow \varepsilon$  is a rule, *accept*; else, *reject*. [[ $w = \varepsilon$  case]]
2. For  $i = 1$  to  $n$ : [[ examine each substring of length 1 ]]
3. For each variable  $A$ :
4. Test whether  $A \rightarrow b$  is a rule, where  $b = w_i$ .
5. If so, place  $A$  in  $\text{table}(i, i)$ .
6. For  $l = 2$  to  $n$ : [[ 1 is the length of the substring ]]
7. For  $i = 1$  to  $n - l + 1$ : [[ i is the start position of the substring ]]
8. Let  $j = i + l - 1$ . [[ j is the end position of the substring ]]
9. For  $k = i$  to  $j - 1$ : [[ k is the split position ]]
10. For each rule  $A \rightarrow BC$ :
11. If  $\text{table}(i, k)$  contains  $B$  and  $\text{table}(k + 1, j)$  contains  $C$ , put  $A$  in  $\text{table}(i, j)$ .
12. If  $S$  is in  $\text{table}(1, n)$ , then:
- 13.
14. Otherwise, output  $\varepsilon$ .”

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