

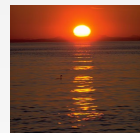


Chapter 3: SQL

Database System Concepts, 5th Ed.

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Chapter 3: SQL

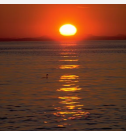
- Data Definition
- Basic Query Structure
- Set Operations
- Aggregate Functions
- Null Values
- Nested Subqueries
- Complex Queries
- Views
- Modification of the Database
- Joined Relations**





History

- IBM Sequel language developed as part of System R project at the IBM San Jose Research Laboratory
- Renamed Structured Query Language (SQL)
- ANSI and ISO standard SQL:
 - SQL-86
 - SQL-89
 - SQL-92
 - SQL:1999 (language name became Y2K compliant!)
 - SQL:2003
- Commercial systems offer most, if not all, SQL-92 features, plus varying feature sets from later standards and special proprietary features.
 - Not all examples here may work on your particular system.

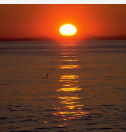




Data Definition Language

Allows the specification of not only a set of relations but also information about each relation, including:

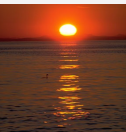
- The schema for each relation.
- The domain of values associated with each attribute.
- Integrity constraints
- The set of indices to be maintained for each relations.
- Security and authorization information for each relation.
- The physical storage structure of each relation on disk.





Domain Types in SQL

- **char(n).** Fixed length character string, with user-specified length n .
- **varchar(n).** Variable length character strings, with user-specified maximum length n .
- **int.** Integer (a finite subset of the integers that is machine-dependent).
- **smallint.** Small integer (a machine-dependent subset of the integer domain type).
- **numeric(p,d).** Fixed point number, with user-specified precision of p digits, with n digits to the right of decimal point.
- **real, double precision.** Floating point and double-precision floating point numbers, with machine-dependent precision.
- **float(n).** Floating point number, with user-specified precision of at least n digits.
- More are covered in Chapter 4.





Create Table Construct

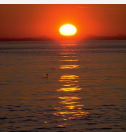
- An SQL relation is defined using the **create table** command:

```
create table  $r$  ( $A_1 D_1, A_2 D_2, \dots, A_n D_n,$   
                (integrity-constraint1),  
                ...,  
                (integrity-constraintk))
```

- r is the name of the relation
- each A_i is an attribute name in the schema of relation r
- D_i is the data type of values in the domain of attribute A_i

- Example:

```
create table branch  
  (branch_name  char(15) not null,  
   branch_city char(30),  
   assets       integer)
```





Integrity Constraints in Create Table

- not null
- primary key (A_1, \dots, A_n)

Example: Declare *branch_name* as the primary key for *branch*

```
create table branch
    (branch_name char(15),
     branch_city  char(30),
     assets       integer,
     primary key (branch_name))
```

primary key declaration on an attribute automatically ensures **not null** in SQL-92 onwards, needs to be explicitly stated in SQL-89





Drop and Alter Table Constructs

- The **drop table** command deletes all information about the dropped relation from the database.
- The **alter table** command is used to add attributes to an existing relation:

alter table r add A D

where A is the name of the attribute to be added to relation r and D is the domain of A .

- All tuples in the relation are assigned *null* as the value for the new attribute.
- The **alter table** command can also be used to drop attributes of a relation:

alter table r drop A

where A is the name of an attribute of relation r

- Dropping of attributes not supported by many databases



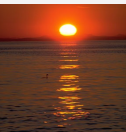


Basic Query Structure

- SQL is based on set and relational operations with certain modifications and enhancements
- A typical SQL query has the form:

select A_1, A_2, \dots, A_n
from r_1, r_2, \dots, r_m
where P

- A_i represents an attribute
 - R_i represents a relation
 - P is a predicate.
- This query is equivalent to the relational algebra expression.
$$\Pi_{A_1, A_2, \dots, A_n} (\sigma_P(r_1 \times r_2 \times \dots \times r_m))$$
 - The result of an SQL query is a relation.





The select Clause

- The **select** clause list the attributes desired in the result of a query
 - corresponds to the projection operation of the relational algebra

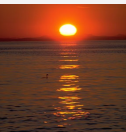
- Example: find the names of all branches in the *loan* relation:

select *branch_name*
from *loan*

- In the relational algebra, the query would be:

$\Pi_{branch_name}(loan)$

- NOTE: SQL names are case insensitive (i.e., you may use upper- or lower-case letters.)
 - E.g. *Branch_Name* *BRANCH_NAME* *branch_name*
 - Some people use upper case wherever we use bold font.





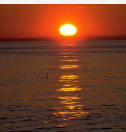
The select Clause (Cont.)

- SQL allows duplicates in relations as well as in query results.
- To force the elimination of duplicates, insert the keyword **distinct** after select.
- Find the names of all branches in the *loan* relations, and remove duplicates

```
select distinct branch_name  
from loan
```

- The keyword **all** specifies that duplicates not be removed.

```
select all branch_name  
from loan
```





The select Clause (Cont.)

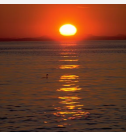
- An asterisk in the select clause denotes “all attributes”

```
select *  
from loan
```

- The **select** clause can contain arithmetic expressions involving the operation, +, −, *, and /, and operating on constants or attributes of tuples.
- The query:

```
select loan_number, branch_name, amount * 100  
from loan
```

would return a relation that is the same as the *loan* relation, except that the value of the attribute *amount* is multiplied by 100.



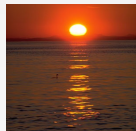


The where Clause

- The **where** clause specifies conditions that the result must satisfy
 - Corresponds to the selection predicate of the relational algebra.
- To find all loan number for loans made at the Perryridge branch with loan amounts greater than \$1200.

```
select loan_number  
from loan  
where branch_name = 'Perryridge' and amount > 1200
```

- Comparison results can be combined using the logical connectives **and**, **or**, and **not**.
- Comparisons can be applied to results of arithmetic expressions.

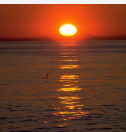




The where Clause (Cont.)

- SQL includes a **between** comparison operator
- Example: Find the loan number of those loans with loan amounts between \$90,000 and \$100,000 (that is, $\geq \$90,000$ and $\leq \$100,000$)

```
select loan_number  
      from loan  
      where amount between 90000 and 100000
```





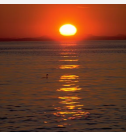
The from Clause

- The **from** clause lists the relations involved in the query
 - Corresponds to the Cartesian product operation of the relational algebra.
- Find the Cartesian product *borrower X loan*

```
select *  
from borrower, loan
```

- Find the name, loan number and loan amount of all customers having a loan at the Perryridge branch.

```
select customer_name, borrower.loan_number, amount  
from borrower, loan  
where borrower.loan_number = loan.loan_number and  
branch_name = 'Perryridge'
```





The Rename Operation

- The SQL allows renaming relations and attributes using the **as** clause:
old-name as new-name
- Find the name, loan number and loan amount of all customers; rename the column name *loan_number* as *loan_id*.

```
select customer_name, borrower.loan_number as loan_id, amount  
from borrower, loan  
where borrower.loan_number = loan.loan_number
```





Tuple Variables

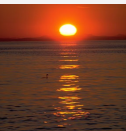
- Tuple variables are defined in the **from** clause via the use of the **as** clause.
- Find the customer names and their loan numbers for all customers having a loan at some branch.

```
select customer_name, T.loan_number, S.amount  
from borrower as T, loan as S  
where T.loan_number = S.loan_number
```

- Find the names of all branches that have greater assets than some branch located in Brooklyn.

```
select distinct T.branch_name  
from branch as T, branch as S  
where T.assets > S.assets and S.branch_city = 'Brooklyn'
```

- Keyword **as** is optional and may be omitted
borrower as T borrower T





String Operations

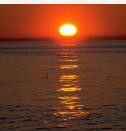
- SQL includes a string-matching operator for comparisons on character strings. The operator “like” uses patterns that are described using two special characters:
 - percent (%). The % character matches any substring.
 - underscore (_). The _ character matches any character.
- Find the names of all customers whose street includes the substring “Main”.

```
select customer_name  
from customer  
where customer_street like '% Main%'
```

- Match the name “Main%”

```
like 'Main\%' escape '\'
```

- SQL supports a variety of string operations such as
 - concatenation (using “||”)
 - converting from upper to lower case (and vice versa)
 - finding string length, extracting substrings, etc.





Ordering the Display of Tuples

- List in alphabetic order the names of all customers having a loan in Perryridge branch

```
select distinct customer_name  
from    borrower, loan  
where borrower loan_number = loan.loan_number and  
        branch_name = 'Perryridge'  
order by customer_name
```

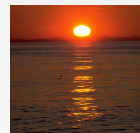
- We may specify **desc** for descending order or **asc** for ascending order, for each attribute; ascending order is the default.
 - Example: **order by** *customer_name* **desc**





Duplicates

- In relations with duplicates, SQL can define how many copies of tuples appear in the result.
- **Multiset** versions of some of the relational algebra operators – given multiset relations r_1 and r_2 :
 1. $\sigma_{\theta}(r_1)$: If there are c_1 copies of tuple t_1 in r_1 , and t_1 satisfies selections σ_{θ} , then there are c_1 copies of t_1 in $\sigma_{\theta}(r_1)$.
 2. $\Pi_A(r)$: For each copy of tuple t_i in r_1 , there is a copy of tuple $\Pi_A(t_i)$ in $\Pi_A(r_1)$ where $\Pi_A(t_i)$ denotes the projection of the single tuple t_i .
 3. $r_1 \times r_2$: If there are c_1 copies of tuple t_1 in r_1 and c_2 copies of tuple t_2 in r_2 , there are $c_1 \times c_2$ copies of the tuple $t_1 \cdot t_2$ in $r_1 \times r_2$





Duplicates (Cont.)

- Example: Suppose multiset relations $r_1 (A, B)$ and $r_2 (C)$ are as follows:

$$r_1 = \{(1, a) (2, a)\} \quad r_2 = \{(2), (3), (3)\}$$

- Then $\Pi_B(r_1)$ would be $\{(a), (a)\}$, while $\Pi_B(r_1) \times r_2$ would be

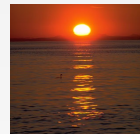
$$\{(a, 2), (a, 2), (a, 3), (a, 3), (a, 3), (a, 3)\}$$

- SQL duplicate semantics:

select A_1, A_2, \dots, A_n
from r_1, r_2, \dots, r_m
where P

is equivalent to the *multiset* version of the expression:

$$\Pi_{A_1, A_2, \dots, A_n} (\sigma_P(r_1 \times r_2 \times \dots \times r_m))$$



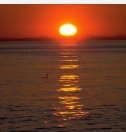


Set Operations

- The set operations **union**, **intersect**, and **except** operate on relations and correspond to the relational algebra operations \cup , \cap , $-$.
- Each of the above operations automatically eliminates duplicates; to retain all duplicates use the corresponding multiset versions **union all**, **intersect all** and **except all**.

Suppose a tuple occurs m times in r and n times in s , then, it occurs:

- $m + n$ times in r **union all** s
- $\min(m, n)$ times in r **intersect all** s
- $\max(0, m - n)$ times in r **except all** s





Set Operations

- Find all customers who have a loan, an account, or both:

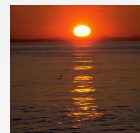
```
(select customer_name from depositor)  
union  
(select customer_name from borrower)
```

- Find all customers who have both a loan and an account.

```
(select customer_name from depositor)  
intersect  
(select customer_name from borrower)
```

- Find all customers who have an account but no loan.

```
(select customer_name from depositor)  
except  
(select customer_name from borrower)
```





Aggregate Functions

- These functions operate on the multiset of values of a column of a relation, and return a value

avg: average value

min: minimum value

max: maximum value

sum: sum of values

count: number of values





Aggregate Functions (Cont.)

- Find the average account balance at the Perryridge branch.

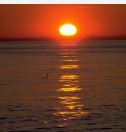
```
select avg (balance)  
  from account  
 where branch_name = 'Perryridge'
```

- Find the number of tuples in the *customer* relation.

```
select count (*)  
  from customer
```

- Find the number of depositors in the bank.

```
select count (distinct customer_name)  
  from depositor
```



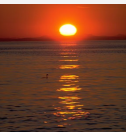


Aggregate Functions – Group By

- Find the number of depositors for each branch.

```
select branch_name, count (distinct customer_name)  
  from depositor, account  
  where depositor.account_number = account.account_number  
 group by branch_name
```

Note: Attributes in **select** clause outside of aggregate functions must appear in **group by** list





Aggregate Functions – Having Clause

- Find the names of all branches where the average account balance is more than \$1,200.

```
select branch_name, avg (balance)  
      from account  
      group by branch_name  
      having avg (balance) > 1200
```

Note: predicates in the **having** clause are applied after the formation of groups whereas predicates in the **where** clause are applied before forming groups





Null Values

- It is possible for tuples to have a null value, denoted by *null*, for some of their attributes
- *null* signifies an unknown value or that a value does not exist.
- The predicate **is null** can be used to check for null values.
 - Example: Find all loan number which appear in the *loan* relation with null values for *amount*.

```
select loan_number  
from loan  
where amount is null
```

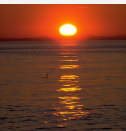
- The result of any arithmetic expression involving *null* is *null*
 - Example: $5 + \text{null}$ returns null
- However, aggregate functions simply ignore nulls
 - More on next slide





Null Values and Three Valued Logic

- Any comparison with *null* returns *unknown*
 - Example: $5 < null$ or $null <> null$ or $null = null$
- Three-valued logic using the truth value *unknown*:
 - OR: $(unknown \text{ or } true) = true$,
 $(unknown \text{ or } false) = unknown$
 $(unknown \text{ or } unknown) = unknown$
 - AND: $(true \text{ and } unknown) = unknown$,
 $(false \text{ and } unknown) = false$,
 $(unknown \text{ and } unknown) = unknown$
 - NOT: $(\text{not } unknown) = unknown$
 - “*P* is unknown” evaluates to true if predicate *P* evaluates to *unknown*
- Result of **where** clause predicate is treated as *false* if it evaluates to *unknown*





Null Values and Aggregates

- Total all loan amounts

```
select sum (amount )  
from loan
```

- Above statement ignores null amounts
- Result is *null* if there is no non-null amount
- All aggregate operations except **count(*)** ignore tuples with null values on the aggregated attributes.





Nested Subqueries

- SQL provides a mechanism for the nesting of subqueries.
- A **subquery** is a **select-from-where** expression that is nested within another query.
- A common use of subqueries is to perform tests for set membership, set comparisons, and set cardinality.





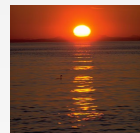
Example Query

- Find all customers who have both an account and a loan at the bank.

```
select distinct customer_name  
from borrower  
where customer_name in (select customer_name  
                        from depositor )
```

- Find all customers who have a loan at the bank but do not have an account at the bank

```
select distinct customer_name  
from borrower  
where customer_name not in (select customer_name  
                        from depositor )
```



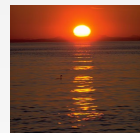


Example Query

- Find all customers who have both an account and a loan at the Perryridge branch

```
select distinct customer_name
from borrower, loan
where borrower.loan_number = loan.loan_number and
       branch_name = 'Perryridge' and
       (branch_name, customer_name) in
       (select branch_name, customer_name
        from depositor, account
        where depositor.account_number =
              account.account_number)
```

- **Note:** Above query can be written in a much simpler manner. The formulation above is simply to illustrate SQL features.





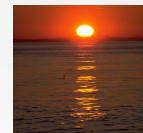
Set Comparison

- Find all branches that have greater assets than some branch located in Brooklyn.

```
select distinct T.branch_name  
from branch as T, branch as S  
where T.assets > S.assets and  
S.branch_city = 'Brooklyn'
```

- Same query using > **some** clause

```
select branch_name  
from branch  
where assets > some  
(select assets  
from branch  
where branch_city = 'Brooklyn')
```





Definition of Some Clause

- $F \text{ <comp> some } r \Leftrightarrow \exists t \in r \text{ such that } (F \text{ <comp> } t)$

Where <comp> can be: <, ≤, >, =, ≠

$$(5 < \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline 6 \\ \hline \end{array}) = \text{true} \quad (\text{read: } 5 < \text{some tuple in the relation})$$

$$(5 < \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline \end{array}) = \text{false}$$

$$(5 = \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline \end{array}) = \text{true}$$

$$(5 \neq \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline \end{array}) = \text{true (since } 0 \neq 5)$$

$(= \text{some}) \equiv \text{in}$

However, $(\neq \text{some}) \not\equiv \text{not in}$





Example Query

- Find the names of all branches that have greater assets than all branches located in Brooklyn.

```
select branch_name
from branch
where assets > all
      (select assets
from branch
where branch_city = 'Brooklyn')
```





Definition of all Clause

- $F \text{ <comp> all } r \Leftrightarrow \forall t \in r (F \text{ <comp> } t)$

$$(5 < \text{all } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline 6 \\ \hline \end{array}) = \text{false}$$

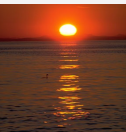
$$(5 < \text{all } \begin{array}{|c|} \hline 6 \\ \hline 10 \\ \hline \end{array}) = \text{true}$$

$$(5 = \text{all } \begin{array}{|c|} \hline 4 \\ \hline 5 \\ \hline \end{array}) = \text{false}$$

$$(5 \neq \text{all } \begin{array}{|c|} \hline 4 \\ \hline 6 \\ \hline \end{array}) = \text{true (since } 5 \neq 4 \text{ and } 5 \neq 6)$$

$(\neq \text{all}) \equiv \text{not in}$

However, $(= \text{all}) \not\equiv \text{in}$





Test for Empty Relations

- The **exists** construct returns the value **true** if the argument subquery is nonempty.
- **exists** $r \Leftrightarrow r \neq \emptyset$
- **not exists** $r \Leftrightarrow r = \emptyset$



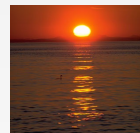


Example Query

- Find all customers who have an account at all branches located in Brooklyn.

```
select distinct S.customer_name
from depositor as S
where not exists (
    (select branch_name
from branch
where branch_city = 'Brooklyn')
except
    (select R.branch_name
from depositor as T, account as R
where T.account_number = R.account_number and
        S.customer_name = T.customer_name ))
```

- Note that $X - Y = \emptyset \Leftrightarrow X \subseteq Y$
- *Note:* Cannot write this query using = **all** and its variants

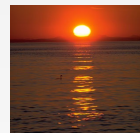




Test for Absence of Duplicate Tuples

- The **unique** construct tests whether a subquery has any duplicate tuples in its result.
- Find all customers who have at most one account at the Perryridge branch.

```
select T.customer_name
from depositor as T
where unique (
    select R.customer_name
from account, depositor as R
where T.customer_name = R.customer_name and
       R.account_number = account.account_number and
       account.branch_name = 'Perryridge')
```



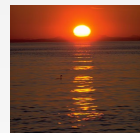


Example Query

- Find all customers who have at least two accounts at the Perryridge branch.

```
select distinct T.customer_name  
from depositor as T  
where not unique (  
    select R.customer_name  
    from account, depositor as R  
    where T.customer_name = R.customer_name and  
          R.account_number = account.account_number and  
          account.branch_name = 'Perryridge')
```

- Variable from outer level is known as a **correlation variable**



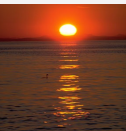


Derived Relations

- SQL allows a subquery expression to be used in the **from** clause
- Find the average account balance of those branches where the average account balance is greater than \$1200.

```
select branch_name, avg_balance
from (select branch_name, avg (balance)
      from account
      group by branch_name )
as branch_avg ( branch_name, avg_balance )
where avg_balance > 1200
```

Note that we do not need to use the **having** clause, since we compute the temporary (view) relation *branch_avg* in the **from** clause, and the attributes of *branch_avg* can be used directly in the **where** clause.

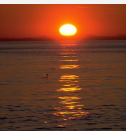




With Clause

- The **with** clause provides a way of defining a temporary view whose definition is available only to the query in which the **with** clause occurs.
- Find all accounts with the maximum balance

```
with max_balance (value) as  
    select max (balance)  
    from account  
select account_number  
from account, max_balance  
where account.balance = max_balance.value
```

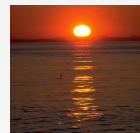




Complex Queries using With Clause

- Find all branches where the total account deposit is greater than the average of the total account deposits at all branches.

```
with branch_total (branch_name, value) as  
    select branch_name, sum (balance)  
    from account  
    group by branch_name  
with branch_total_avg (value) as  
    select avg (value)  
    from branch_total  
select branch_name  
from branch_total, branch_total_avg  
where branch_total.value >= branch_total_avg.value
```



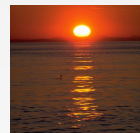


Views

- In some cases, it is not desirable for all users to see the entire logical model (that is, all the actual relations stored in the database.)
- Consider a person who needs to know a customer's name, loan number and branch name, but has no need to see the loan amount. This person should see a relation described, in SQL, by

```
(select customer_name, borrower.loan_number, branch_name  
      from borrower, loan  
      where borrower.loan_number = loan.loan_number )
```

- A **view** provides a mechanism to hide certain data from the view of certain users.
- Any relation that is not of the conceptual model but is made visible to a user as a “virtual relation” is called a **view**.





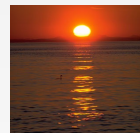
View Definition

- A view is defined using the **create view** statement which has the form

create view *v* **as** < query expression >

where <query expression> is any legal SQL expression. The view name is represented by *v*.

- Once a view is defined, the view name can be used to refer to the virtual relation that the view generates.
- When a view is created, the query expression is stored in the database; the expression is substituted into queries using the view.





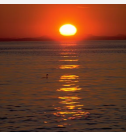
Example Queries

- A view consisting of branches and their customers

```
create view all_customer as  
    (select branch_name, customer_name  
     from depositor, account  
     where depositor.account_number =  
           account.account_number )  
union  
    (select branch_name, customer_name  
     from borrower, loan  
     where borrower.loan_number = loan.loan_number )
```

- Find all customers of the Perryridge branch

```
select customer_name  
     from all_customer  
     where branch_name = 'Perryridge'
```





Views Defined Using Other Views

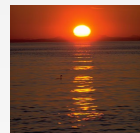
- One view may be used in the expression defining another view
- A view relation v_1 is said to *depend directly* on a view relation v_2 if v_2 is used in the expression defining v_1
- A view relation v_1 is said to *depend on* view relation v_2 if either v_1 depends directly to v_2 or there is a path of dependencies from v_1 to v_2
- A view relation v is said to be *recursive* if it depends on itself.





View Expansion

- A way to define the meaning of views defined in terms of other views.
- Let view v_1 be defined by an expression e_1 that may itself contain uses of view relations.
- View expansion of an expression repeats the following replacement step:
 - repeat**
 - Find any view relation v_i in e_1
 - Replace the view relation v_i by the expression defining v_i
 - until** no more view relations are present in e_1
- As long as the view definitions are not recursive, this loop will terminate





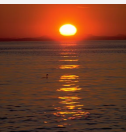
Modification of the Database – Deletion

- Delete all account tuples at the Perryridge branch

```
delete from account  
where branch_name = 'Perryridge'
```

- Delete all accounts at every branch located in the city 'Needham'.

```
delete from account  
where branch_name in (select branch_name  
                        from branch  
                        where branch_city = 'Needham')
```



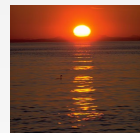


Example Query

- Delete the record of all accounts with balances below the average at the bank.

```
delete from account  
  where balance < (select avg (balance)  
                    from account)
```

- Problem: as we delete tuples from deposit, the average balance changes
- Solution used in SQL:
 1. First, compute **avg** balance and find all tuples to delete
 2. Next, delete all tuples found above (without recomputing **avg** or retesting the tuples)





Modification of the Database – Insertion

- Add a new tuple to *account*

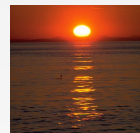
```
insert into account  
values ('A-9732', 'Perryridge', 1200)
```

or equivalently

```
insert into account (branch_name, balance, account_number)  
values ('Perryridge', 1200, 'A-9732')
```

- Add a new tuple to *account* with *balance* set to null

```
insert into account  
values ('A-777', 'Perryridge', null )
```





Modification of the Database – Insertion

- Provide as a gift for all loan customers of the Perryridge branch, a \$200 savings account. Let the loan number serve as the account number for the new savings account

insert into *account*

select *loan_number, branch_name, 200*

from *loan*

where *branch_name* = 'Perryridge'

insert into *depositor*

select *customer_name, loan_number*

from *loan, borrower*

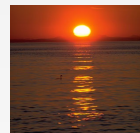
where *branch_name* = 'Perryridge'

and *loan.account_number = borrower.account_number*

- The **select from where** statement is evaluated fully before any of its results are inserted into the relation (otherwise queries like

insert into table1 select * from table1

would cause problems)





Modification of the Database – Updates

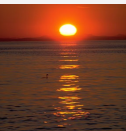
- Increase all accounts with balances over \$10,000 by 6%, all other accounts receive 5%.

- Write two **update** statements:

```
update account  
set balance = balance * 1.06  
where balance > 10000
```

```
update account  
set balance = balance * 1.05  
where balance ≤ 10000
```

- The order is important
- Can be done better using the **case** statement (next slide)





Case Statement for Conditional Updates

- Same query as before: Increase all accounts with balances over \$10,000 by 6%, all other accounts receive 5%.

```
update account  
set balance = case  
    when balance <= 10000 then balance * 1.05  
    else balance * 1.06  
end
```





Update of a View

- Create a view of all loan data in the *loan* relation, hiding the *amount* attribute

```
create view loan_branch as  
    select loan_number, branch_name  
    from loan
```

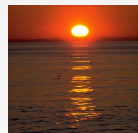
- Add a new tuple to *branch_loan*

```
insert into branch_loan  
    values ('L-37', 'Perryridge')
```

This insertion must be represented by the insertion of the tuple

('L-37', 'Perryridge', *null*)

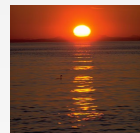
into the *loan* relation





Updates Through Views (Cont.)

- Some updates through views are impossible to translate into updates on the database relations
 - **create view *v* as**
select *loan_number, branch_name, amount*
from *loan*
where *branch_name* = 'Perryridge'
insert into *v* values ('L-99', 'Downtown', '23')
- Others cannot be translated uniquely
 - **insert into *all_customer* values** ('Perryridge', 'John')
 - ▶ Have to choose loan or account, and create a new loan/account number!
- Most SQL implementations allow updates only on simple views (without aggregates) defined on a single relation

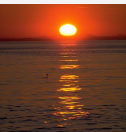




Joined Relations**

- **Join operations** take two relations and return as a result another relation.
- These additional operations are typically used as subquery expressions in the **from** clause
- **Join condition** – defines which tuples in the two relations match, and what attributes are present in the result of the join.
- **Join type** – defines how tuples in each relation that do not match any tuple in the other relation (based on the join condition) are treated.

<i>Join types</i>		<i>Join Conditions</i>
inner join left outer join right outer join full outer join		natural on <predicate> using (A_1, A_1, \dots, A_n)



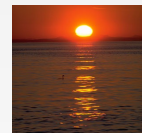


Joined Relations – Datasets for Examples

- Relation *loan*
- Relation *borrower*

<i>loan_number</i>	<i>branch_name</i>	<i>amount</i>	<i>customer_name</i>	<i>loan_number</i>
L-170	Downtown	3000	Jones	L-170
L-230	Redwood	4000	Smith	L-230
L-260	Perryridge	1700	Hayes	L-155
<i>loan</i>			<i>borrower</i>	

- Note: borrower information missing for L-260 and loan information missing for L-155





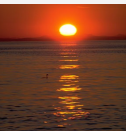
Joined Relations – Examples

- **loan inner join borrower on**
loan.loan_number = borrower.loan_number

<i>loan_number</i>	<i>branch_name</i>	<i>amount</i>	<i>customer_name</i>	<i>loan_number</i>
L-170	Downtown	3000	Jones	L-170
L-230	Redwood	4000	Smith	L-230

- **loan left outer join borrower on**
loan.loan_number = borrower.loan_number

<i>loan_number</i>	<i>branch_name</i>	<i>amount</i>	<i>customer_name</i>	<i>loan_number</i>
L-170	Downtown	3000	Jones	L-170
L-230	Redwood	4000	Smith	L-230
L-260	Perryridge	1700	<i>null</i>	<i>null</i>





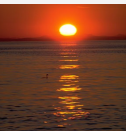
Joined Relations – Examples

- *loan natural inner join borrower*

<i>loan_number</i>	<i>branch_name</i>	<i>amount</i>	<i>customer_name</i>	<i>loan_number</i>
L-170	Downtown	3000	Jones	L-170
L-230	Redwood	4000	Smith	L-230

- *loan natural right outer join borrower*

<i>loan_number</i>	<i>branch_name</i>	<i>amount</i>	<i>customer_name</i>
L-170	Downtown	3000	Jones
L-230	Redwood	4000	Smith
L-155	<i>null</i>	<i>null</i>	Hayes





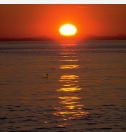
Joined Relations – Examples

- *loan full outer join borrower using (loan_number)*

<i>loan_number</i>	<i>branch_name</i>	<i>amount</i>	<i>customer_name</i>
L-170	Downtown	3000	Jones
L-230	Redwood	4000	Smith
L-260	Perryridge	1700	<i>null</i>
L-155	<i>null</i>	<i>null</i>	Hayes

- Find all customers who have either an account or a loan (but not both) at the bank.

```
select customer_name  
      from (depositor natural full outer join borrower)  
      where account_number is null or loan_number is null
```





End of Chapter 3

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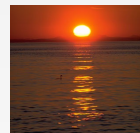




Figure 3.1: Database Schema

branch (*branch_name*, *branch_city*, *assets*)

customer (*customer_name*, *customer_street*, *customer_city*)

loan (*loan_number*, *branch_name*, *amount*)

borrower (*customer_name*, *loan_number*)

account (*account_number*, *branch_name*, *balance*)

depositor (*customer_name*, *account_number*)

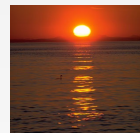




Figure 3.3: Tuples inserted into *loan* and *borrower*

<i>loan_number</i>	<i>branch_name</i>	<i>amount</i>	<i>customer_name</i>	<i>loan_number</i>
L-11	Round Hill	900	Adams	L-16
L-14	Downtown	1500	Curry	L-93
L-15	Perryridge	1500	Hayes	L-15
L-16	Perryridge	1300	Jackson	L-14
L-17	Downtown	1000	Jones	L-17
L-23	Redwood	2000	Smith	L-11
L-93	Mianus	500	Smith	L-23
<i>null</i>	<i>null</i>	1900	Williams	L-17
<i>loan</i>			Johnson	<i>null</i>
			<i>borrower</i>	

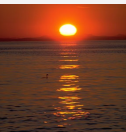




Figure 3.4: The *loan* and *borrower* relations

<i>loan_number</i>	<i>branch_name</i>	<i>amount</i>	<i>customer_name</i>	<i>loan_number</i>
L-170	Downtown	3000	Jones	L-170
L-230	Redwood	4000	Smith	L-230
L-260	Perryridge	1700	Hayes	L-155
<i>loan</i>			<i>borrower</i>	

