***ROUND ROBIN ALGORITHM***

A round-robin is a CPU scheduling algorithm that shares equal portions of resources in circular orders to each process and handles all processes without prioritization. In the round-robin, each process gets a fixed time interval of the slice to utilize the resources or execute its task called time **quantum or time slice**. Some of the round-robin processes are pre-empted if it executed in a given time slot, while the rest of the processes go back to the ready queue and wait to run in a circular order with the scheduled time slot until they complete their task. It removes the starvation for each process to achieve [CPU.](https://www.javatpoint.com/cpu-full-form)

***ROUND ROBIN CPU SCHEDULING ALGORITHM***

**Step 1:** Organize all processes according to their arrival time in the ready queue. The queue structure of the ready queue is based on the FIFO structure to execute all CPU processes.

**Step 2:** Now, we push the first process from the ready queue to execute its task for a fixed time, allocated by each process that arrives in the queue.

**Step 3:** If the process cannot complete their task within defined time interval or slots because it is stopped by another process that pushes from the ready queue to execute their task due to arrival time of the next process is reached. Therefore, CPU saved the previous state of the process, which helps to resume from the point where it is interrupted. (If the burst time of the process is left, push the process end of the ready queue).

**Step 4:** Similarly, the scheduler selects another process from the ready queue to execute its tasks. When a process finishes its task within time slots, the process will not go for further execution because the process's burst time is finished.

**Step 5:** Similarly, we repeat all the steps to execute the process until the work has finished.

***CHARACTERISTICS OF ROUND ROBIN ALGORITHM***

1. It is a pre-emptive algorithm.
2. It shares an equal time interval between all processes to complete their task.
3. It is a starvation free CPU scheduling algorithm. Hence it is known as the fairest and simple algorithm.

***ADVANTAGES***

1. It does not face any starvation issues or convoy effect.
2. Each process gets equal priority to the fair allocation of CPU.
3. It is easy to implement the CPU Scheduling algorithm.
4. Each new process is added to the end of the ready queue as the next process's arrival time is reached.
5. Each process is executed in circular order that shares a fixed time slot or quantum.
6. Every process gets an opportunity in the round-robin scheduling algorithm to reschedule after a given quantum period.

***DISADVANTAGES***

1. If the time quantum is lower, it takes more time on context switching between the processes.
2. It does not provide any special priority to execute the most important process.
3. The waiting time of a large process is higher due to the short time slot.
4. The performance of the algorithm depends on the time quantum.
5. The response time of the process is higher due to large slices to time quantum.
6. Getting a correct time slot or quantum is quite difficult for all processes in the round-robin algorithm.

***FLOWCHART***

***Diagram

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***Round Robin CPU Scheduling Example:***

Let's understand the concepts of Round Robin with an example. Suppose we have five processes P1, P2, P3, P4 and P5. The arrival and burst time of each process are mentioned in the following table, as shown below. The time quantum is three units.

Graphical user interface, application

Description automatically generated

Now we have to create the **ready queue** and the **Gantt chart** for Round Robin CPU Scheduler.

Ready queue: P1, P3, P1, P2, P4, P3, P5, P1, P3, P5

Here is the Gantt chart:

Table

Description automatically generated with medium confidence

**Step 1:** At time 0, process P1 enters the ready queue and starts its execution for the defined time slot 3. During 3 units of the time slice, another process, P3, arrives in the ready queue because its arrival time is 1.

**Step 2:** Now, process P3 starts its execution with a time slot of 3 units while process P1 has to wait. After execution of process P3, process P1 again resumes its execution for time slot 3.

**Step 3:** During the execution of process P1, two more processes P2 and P4, arrive in the ready queue to begin their execution. Since process P2 has come first, it will be executed for time quantum 2 units after that P4 is executed.

**Step 4:** Here, P4 has executed for time slot 3 units, and its task is completed because BT (Burst Time) is 2. Hence it will not go to the ready queue.

**Step 5:** After that, process P3 is executed from the ready queue for time slot 3 and then process P5 arrives for time slot 3.

**Step 6:** Meanwhile, process P5 is executed, process P1 and P3 have to wait in the ready queue.

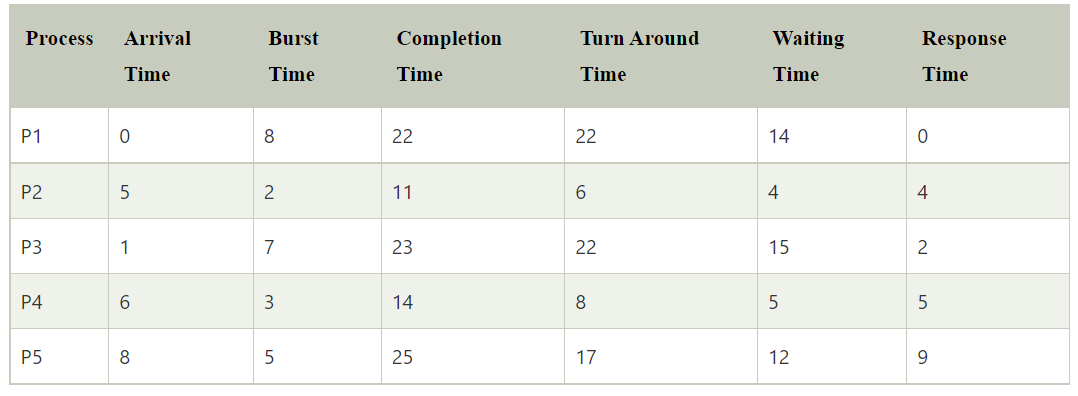
**Step 7:** Now process P1 is fetched from the ready queue and starts their execution for time slot 2 as it requires only 2 BT to finish its tasks. Hence it will not go to the ready queue for further execution.

**Step 8:** Now, the process P3 is executed for time slot 1 as it requires only 1 BT to complete its tasks.

**Step 9:** The final process P5 is executed for time slot 2 because it requires only 2 BT to complete its tasks.

The following are the important terms to find the Completion time, Turn Around Time (TAT), Response Time (RT) and Waiting Time (WT).

1. **Completion Time**: It defines the time when processes complete their execution.
2. **Turn Around Time**: It defines the time difference between the completion time (CT) and the arrival time (AT).  
   Turn Around Time (TAT) = Completion Time (CT) - Arrival Time (AT)
3. **Waiting Time**: It defines the total time between requesting action and acquiring the resource.  
   Waiting Time (WT) = Turn Around Time (TAT) - Burst Time (BT)
4. **Response Time**: It is the time that defines at which time the **system** response to a process.



Completion time for process P1 = 22, P2 = 11, P3 = 23, P4 = 14 and P5 = 25.

Turn Around Time for P1 = Completion Time (CT) - Arrival Time (AT)  
                          22 - 0 = 22

Turn Around Time for P2 = 11 - 5 = 6

Turn Around Time for P3 = 23 - 1 = 22

Turn Around Time for P4 = 14 - 6 = 8

Turn Around Time for P5 = 25 - 8 = 17

Waiting Time for P1 = Turn Around Time (TAT) - Burst Time (BT)  
                        22 - 8 = 14

Waiting Time for P2 = 6 - 2 = 4

Waiting Time for P3 = 22 - 7 = 15

Waiting Time for P4 = 8 - 3 = 5

Waiting Time for P5 = 17 - 5 = 12

Average Waiting Time = (14 + 4 + 15 + 5 + 12)/5 = 50/5 = 10

Average Turn Around Time = (22+6+22+8+17)/5= 75/5 = 15

### *Round Robin CPU Scheduling example with sequential arrival time*

Let's understand another example of Round Robin with sequential arrival time. Here we have four processes P1, P2, P3, and P4. The arrival and burst time of each process are mentioned in the following table, as shown below. The time quantum is **6** units.

### *Table Description automatically generated with medium confidence*

### Here is the Gantt chart:

### *Table Description automatically generated*

**Step 1:** At time 0, process P1 arrives in the ready queue and executes its tasks for time quantum 6 units. During 6 units of the time slice, another process P2, P3 and P4 arrive in the ready queue. After that, process P1 will return to the end of the ready queue and await their execution.

**Step 2:** Now, process P2 starts their execution for time slot 5 units because the Burst Time (BT) is 5, and it does not go to the ready queue for further execution.

**Step 3:** After that process, P3 will start its execution, which has 10 Burst time, but the time quantum is 6 units. Therefore, it executes its tasks for a defined time limit and is added to the ready queue's end.

**Step 4:** After that, the process P4 starts their execution, which has burst time 11, but the time quantum is 6 units. It executes its tasks for only 6 seconds and then adds to the end of the ready queue.

**Step 5:** After the execution of P4, now P1 will start its execution again for 2 units or second, and the process P1 terminate or end. Similarly, the complete execution of process P1, then P3, starts its remaining execution for Burst Time 4, and the process is completed.

**Step 6:** After the complete execution of process P3, now process P4 executes for the remaining time slot, which is 5, and the process is finished.

### *Table Description automatically generated*

Now we find the **completion, Turn around time, waiting time** and the average **TAT** and waiting time.

The completion time of P1 is: 25  
The completion time of P2 is: 11  
The completion time of P3 is: 29  
The completion time of P4 is: 34

**Turn Around Time: Completion Time (CT) - Arrival Time (AT)**  
For process P1: 25 - 0 = 25  
For process P2: 11 -1 = 10  
For process P3: 29 - 2 = 27  
For process P4: 34 - 3 = 31  
Average Turn Around Time is: (25+10+27+31)/4 = 23.25

**Process Waiting Time:**

P1 = 25-8 =17  
P2 = 10-5 =5  
P3 = 27-10= 17  
P4 = 31-11=20  
**Average Waiting Time is:** (17+5+17+20)/4 = 59/4 = 14.75

***CODE***

#include<stdio.h>

#include<conio.h>

void main()

{

// initlialize the variable name

int i, NOP, sum=0,count=0, y, quant, wt=0, tat=0, at[10], bt[10], temp[10];

float avg\_wt, avg\_tat;

printf(" Total number of process in the system: ");

scanf("%d", &NOP);

y = NOP; // Assign the number of process to variable y

// Use for loop to enter the details of the process like Arrival time and the Burst Time

for(i=0; i<NOP; i++)

{

printf("\n Enter the Arrival and Burst time of the Process[%d]\n", i+1);

printf(" Arrival time is: \t"); // Accept arrival time

scanf("%d", &at[i]);

printf(" \nBurst time is: \t"); // Accept the Burst time

scanf("%d", &bt[i]);

temp[i] = bt[i]; // store the burst time in temp array

}

// Accept the Time qunat

printf("Enter the Time Quantum for the process: \t");

scanf("%d", &quant);

// Display the process No, burst time, Turn Around Time and the waiting time

printf("\n Process No \t\t Burst Time \t\t TAT \t\t Waiting Time ");

for(sum=0, i = 0; y!=0; )

{

if(temp[i] <= quant && temp[i] > 0) // define the conditions

{

sum = sum + temp[i];

temp[i] = 0;

count=1;

}

else if(temp[i] > 0)

{

temp[i] = temp[i] - quant;

sum = sum + quant;

}

if(temp[i]==0 && count==1)

{

y--; //decrement the process no.

printf("\nProcess No[%d] \t\t %d\t\t\t\t %d\t\t\t %d", i+1, bt[i], sum-at[i], sum-at[i]-bt[i]);

wt = wt+sum-at[i]-bt[i];

tat = tat+sum-at[i];

count =0;

}

if(i==NOP-1)

{

i=0;

}

else if(at[i+1]<=sum)

{

i++;

}

else

{

i=0;

}

}

// represents the average waiting time and Turn Around time

avg\_wt = wt \* 1.0/NOP;

avg\_tat = tat \* 1.0/NOP;

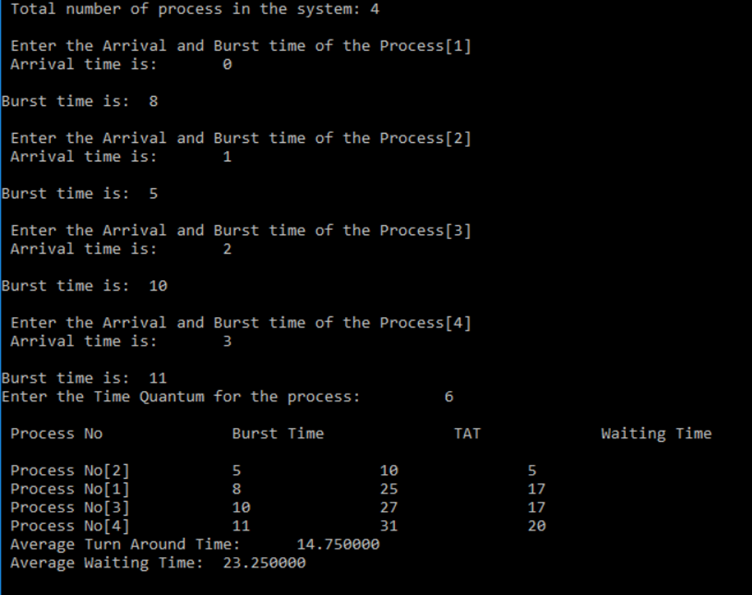
printf("\n Average Turn Around Time: \t%f", avg\_wt);

printf("\n Average Waiting Time: \t%f", avg\_tat);

getch();

}

***OUTPUT***

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