Haxl: Haskell at Facebook

Simon Marlow Jon Coens Louis Brandy Bartosz Nitka Jon Purdy Aaron Roth & others



What's in this talk

- The Haxl project: ~12 months later, where are we
- Haxl published at ICFP'14!
- Haxl open source release!
 - walking through an example data source

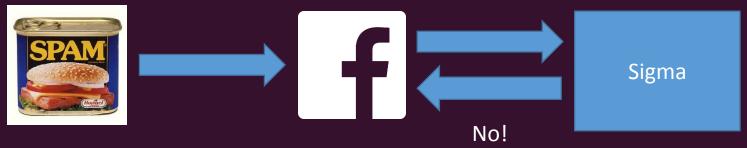


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You can't post this because it has a blocked link.

The content you're trying to share includes a link that's been blocked for being spammy or unsafe:

http://snopes.com/images/template/snopes.gif

For more information, visit the Help Center. If you think you're seeing this by mistake, please let us know.

Okay





Sigma

How does Sigma know what's spam?

• FXL

```
JaffaCakeSpam =
   MessageContains("Jaffa Cakes") &&
   Let
    LikesJaffaCakes(X) = Likes(JaffaCakes, X)
   In
   Length(Filter(LikesJaffaCakes, FriendsOf(SourceId))) < 3</pre>
```

+ machine learning.

```
SpamMessage = SpamScore > 0.99
```

- Where do the inputs to the ML come from?
 - FXL expressions.

What can you do in FXL?

Fetch data from the Facebook graph:

```
FriendCount(uid) = AssocCountByType(uid, AssocFriends)
```

- Fetch data from any of the other 18 data sources
- Run machine learning classifiers
- Perform simple computations

What's good about FXL?

- Clean syntax
 - SI engineers concentrate on fighting spam, not the language
- Static typing
 - cannot push type-incorrect code
- We can push changes fast
 - a couple of minutes from commit to production

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 - No user-defined types
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 Static typing is limited
 - Only a few types: int, double, string, vector, map, JSON
 - No user-defined types
 - Type system doesn't catch as many errors as it could
- Slow (it's an interpreter)

Why are we switching to Haskell?

- Expressivity
- Learning resources available
- Lots of libraries
- Faster (it's compiled)
- Better implementation (error messages etc.)
- Chance to redesign the whole system
 - guaranteed replayability

Technical challenges

- 1. Implicit concurrency
- 2. Implement all the FXL functionality in Haskell
- Translate all the FXL code
- Figure out how to compile+push all the Haskell code to all of the machines in a few minutes

- Status summary: we're mostly done with 1,2,3 and experimenting with a solution for 4.
- Now: testing, bug fixing and optimisation.

Implicit concurrency

In FXL you can write this:

```
NumCommonFriends(x, y) =
  Length(Intersect(FriendsOf(x), FriendsOf(y)));
```

- And Sigma automatically batches the two requests together.
- With existing languages & frameworks you have to specify the concurrency explicitly...

e.g using Haskell asyncs:

```
do
  ax <- async (friendsOf x)
  ay <- async (friendsOf y)
  fx <- wait ax
  fy <- wait ay
  return (length (intersect fx fy))</pre>
```

- Too verbose
- Prone to false dependencies

Larger example: a blog server

My Blog

Popular posts

- Post 2
- Post 1
- Post 3

Categories

- Topic A
 - Post 3
- Topic B
 - Post 2
- Topic C
 - Post 1

Post 1

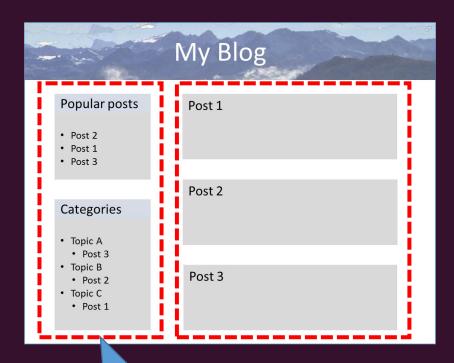
Post 2

Post 3

Another example: a blog server

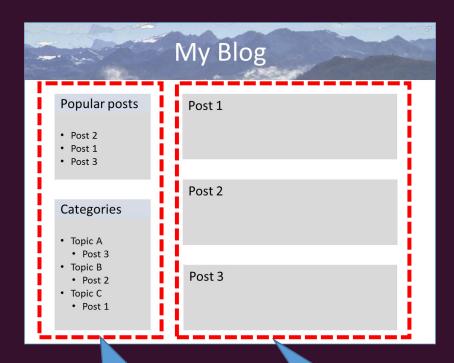


Another example: a blog server



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We want modular code – code each pane independently

We want it to execute efficiently.

- Concurrent data fetches
- No repeated data fetches

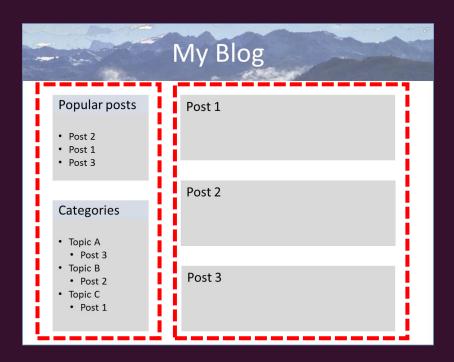
```
data Date -- a calendar date
data PostContent -- the content of a post
data PostInfo = PostInfo
 { postId :: PostId
 , postDate :: Date
 , postTopic :: String
-- data-fetching operations
getPostIds :: Haxl [PostId]
getPostInfo :: PostId -> Haxl PostInfo
getPostContent :: PostId -> Hax1 PostContent
getPostViews :: PostId -> Haxl Int
-- rendering functions
renderPosts :: [(PostInfo,PostContent)] -> Html
renderPage :: Html -> Html -> Html
```

```
blog :: Haxl Html
blog = renderPage <$> leftPane <*> mainPane
getAllPostsInfo :: Haxl [PostInfo]
getAllPostsInfo = mapM getPostInfo =<< getPostIds</pre>
mainPane :: Haxl Html
mainPane = do
  posts <- getAllPostsInfo</pre>
  let ordered =
        take 5 $
        sortBy (flip (comparing postDate)) posts
  content <- mapM (getPostContent . postId) ordered</pre>
  return $ renderPosts (zip ordered content)
```

```
leftPane :: Haxl Html
leftPane = renderSidePane <$> popularPosts <*> topics
getPostDetails :: PostId -> Haxl (PostInfo, PostContent)
getPostDetails pid =
 (,) <$> getPostInfo pid <*> getPostContent pid
popularPosts :: Haxl Html
popularPosts = do
  pids <- getPostIds</pre>
  views <- mapM getPostViews pids</pre>
  let ordered =
        take 5 $ map fst $
        sortBy (flip (comparing snd)) (zip pids views)
  content <- mapM getPostDetails ordered</pre>
  return $ renderPostList content
```

```
topics :: Haxl Html
topics = do
  posts <- getAllPostsInfo
  let topiccounts =
         Map.fromListWith (+) [ (postTopic p, 1) | p <- posts ]
  return $ renderTopics topiccounts</pre>
```

So how did we do?



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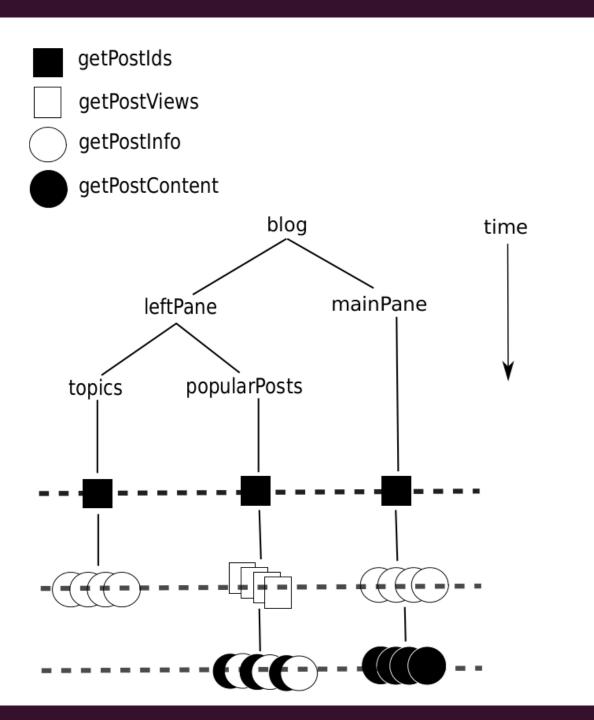
Code is clean and modular

So how did we do?



Code is clean and modular

But does it execute efficiently?



Demo

Implementation

Start with a concurrency monad

Start with a concurrency monad

```
data Result a
 = Done a
  | Blocked (Seq BlockedRequest) (Haxl a)
newtype Haxl a = Haxl { unHaxl :: IO (Result a) }
instance Monad Haxl where
 return a = Haxl $ return (Done a)
 Haxl m >>= k = Haxl $ do
   r < - m
   case r of
     Done a -> unHaxl (k a)
     Blocked br c -> return (Blocked br (c >>= k))
```

Add an Applicative instance

Fetching Data

dataFetch :: Request a -> Haxl a

Type parameter is the result type of the request

Fetching Data

dataFetch :: Request a -> Haxl a

GADT

Type parameter is the result type of the request

data Request a where
FetchPosts :: Request [PostId]
FetchPostInfo :: PostId -> Request PostInfo
FetchPostContent :: PostId -> Request PostContent
FetchPostViews :: PostId -> Request Int

Fetching Data

```
dataFetch :: Request a -> Haxl a
                                              Type parameter is
                     GADT
                                              the result type of
                                                 the request
data Request a where
 FetchPosts :: Request [PostId]
 FetchPostInfo :: PostId -> Request PostInfo
 FetchPostContent :: PostId -> Request PostContent
 FetchPostViews :: PostId -> Request Int
data FetchStatus a = NotFetched | FetchSuccess a
data BlockedRequest =
 forall a . BlockedRequest (Request a) (IORef (FetchStatus a))
```

We can implement dataFetch:

```
dataFetch :: Request a -> Haxl a
dataFetch request = Haxl $ do
  box <- newIORef NotFetched
let br = BlockedRequest request box
let cont = Haxl $ do
    FetchSuccess a <- readIORef box
    return (Done a)
return (Blocked (singleton br) cont)</pre>
```

To fetch data, we need

```
fetch :: [BlockedRequest] -> IO ()
```

Application-specific datafetching function.

Batches multiple requests, uses concurrency, etc.

 To run a computation to completion, we need a loop:

```
runHaxl :: Haxl a -> IO a
runHaxl (Haxl h) = do
  r <- h
  case r of
   Done a -> return a
   Blocked br cont -> do
    fetch (toList br)
    runFetch cont
```

Done!

We also want caching

- Reader monad passes an IORef DataCache around
- Complication:
 - cache maps Request a to a
 - can't do this with Data. Map alone

```
newtype DataCache =
  DataCache (forall a . Map (Request a) (IORef (FetchStatus a)))
```

There is no Fork: an Abstraction for Efficient, Concurrent, and Concise Data Access

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Abstract

We describe a new programming idiom for concurrency, based on Applicative Functors, where concurrency is implicit in the Applicative <*> operator. The result is that concurrent programs can be written in a natural applicative style, and they retain a high degree of clarity and modularity while executing with maximal concurrency. This idiom is particularly useful for programming against external data sources, where the application code is written without the use of explicit concurrency constructs, while the implementation is able to batch together multiple requests for data from the same source, and fetch data from multiple sources concurrently. Our abstraction uses a cache to ensure that multiple requests for the same data return the same result, which frees the programmer from having to arrange to fetch data only once, which in turn leads to greater modularity.

While it is generally applicable, our technique was designed with a particular application in mind: an internal service at Facebook that identifies particular types of content and takes actions based on it. Our application has a large body of business logic that fetches data from as many as 15 different external sources. The framework described in this paper enables the business logic to execute efficiently by automatically fetching data concurrently; we present some preliminary results.

efficiency in this setting: accessing multiple remote data sources efficiently requires *concurrency*, and that normally requires the programmer to intervene and program the concurrency explicitly.

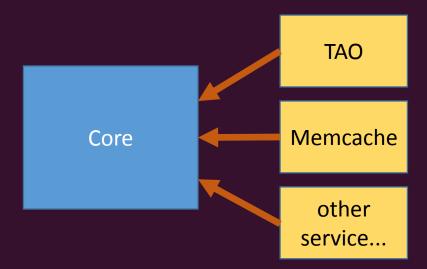
When the business logic is only concerned with *reading* data from external sources and not *writing*, the programmer doesn't care about the order in which data accesses happen, since there are no side-effects that could make the result different when the order changes. So in this case the programmer would be entirely happy with not having to specify either ordering or concurrency, and letting the system perform data access in the most efficient way possible. In this paper we present an embedded domain-specific language (EDSL), written in Haskell, that facilitates this style of programming, while automatically extracting and exploiting any concurrency inherent in the program.

Our contributions can be summarised as follows:

- We present an Applicative abstraction that allows implicit concurrency to be extracted from computations written with a combination of Monad and Applicative. This is an extension of the idea of concurrency monads [9], using Applicative <*> as a way to introduce concurrency (Section [4]). We then develop the idea into an abstraction that supports concurrent access to remote data (Section [5]), and failure (Section [8]).
- We show how to add a *cache* to the framework (Section 6).

But...

- The Request type was wired into the monad
- How can we make the monad independent of the data source(s)?



- Core code includes the monad, caching support etc.
- Core is generic: no data sources built-in

dataFetch :: Request a -> Haxl a

dataFetch :: Request a -> Haxl a

DataSource class: every data source implements this

```
dataFetch :: (DataSource u req, Request req a)
```

- => req a
- -> GenHaxl u a

DataSource class: every data source implements this

Request class: just Eq, Hashable, Typeable, Show

dataFetch :: (DataSource u req, Request req a)

- => req a
- -> GenHaxl u a

DataSource class: every data source implements this

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User state – passed around, can be accessed by data sources

DataSource walk-through

- We'll walk through constructing a complete data source
- We'll make a data source for the Facebook Graph API
 - web API for querying the Facebook Graph
 - using Felipe Lessa's fb package to do the real work
 - our data source will perform requests concurrently up to a maximum number of threads

Start with the request type:

```
data FacebookReq a where
   GetObject :: Id -> FacebookReq Object
   GetUser :: UserId -> FacebookReq User
   GetUserFriends :: UserId -> FacebookReq [Friend]
   deriving Typeable
```

GADT, as before

We also need some boilerplate:

```
deriving instance Eq (FacebookReq a)
deriving instance Show (FacebookReq a)
instance Show1 FacebookReq where show1 = show
instance Hashable (FacebookReq a) where ...
```

• A data source has some state:

```
instance StateKey FacebookReq where
                                                 API keys
  data State FacebookReq =
    FacebookState
       { credentials :: Credentials
       , userAccessToken :: UserAccessToken
         manager :: Manager
         numThreads :: Int
                                         HTTP connection
                                            manager
                                   Concurrency
                                     control
```

Initialise the state:

```
initGlobalState
  :: Int
  -> Credentials
  -> UserAccessToken
  -> IO (State FacebookReq)
initGlobalState threads creds token = do
  manager <- newManager tlsManagerSettings</pre>
  return FacebookState
    { credentials = creds
    , manager = manager
    , userAccessToken = token
    , numThreads = threads
```

nothing surprising there.

Make an instance of DataSource

```
instance DataSourceName FacebookReq where
  dataSourceName _ = "Facebook"

instance DataSource u FacebookReq where
  fetch = facebookFetch
```

Implement fetch

```
facebookFetch
                                  IO to do while
  :: State FacebookReq
                                   the requests
  -> Flags
                                  are in progress
  -> ()
  -> [BlockedFetch FacebookReq]
  -> PerformFetch
facebookFetch FacebookState{..} flags user bfs =
  AsyncFetch $ \inner -> do
    sem <- newQSem numThreads</pre>
    asyncs <- mapM (async . fetchAsync credentials manager
                                 userAccessToken sem) bfs
    inner
    mapM wait asyncs
```

Implement fetchAsync

```
fetchAsync
  :: Credentials -> Manager -> UserAccessToken -> QSem
  -> BlockedFetch FacebookReq
  -> IO ()
fetchAsync creds manager tok sem (BlockedFetch req rvar) =
  bracket (waitQSem sem) (signalQSem sem) $ do
    e <- Control.Exception.try $
         runResourceT $
         runFacebookT creds manager $
         fetchReq tok req
    case e of
      Left ex -> putFailure rvar (ex :: SomeException)
      Right a -> putSuccess rvar a
```

fetchReq maps FacebookReq to FacebookT computations

```
fetchReq
  :: UserAccessToken
  -> FacebookReq a
  -> FacebookT Auth (ResourceT IO) a
fetchReq tok (GetObject (Id id)) =
  getObject ("/" <> id) [] (Just tok)
fetchReq _tok (GetUser id) =
 getUser id [] Nothing
fetchReq tok (GetUserFriends id) = do
  f <- getUserFriends id [] tok</pre>
  source <- fetchAllNextPages f</pre>
  source $$ consume
```

Example

```
main :: IO ()
main = do
   (creds, access_token) <- getCredentials
   facebookState <- initGlobalState 10 creds access_token
   env <- initEnv (stateSet facebookState stateEmpty) ()
   r <- runHaxl env $ do
        likes <- getObject "me/likes"
        mapM getObject (likeIds likes)
   print r</pre>
```

Many requests, performed concurrently.

Back to our Haxl project...

But do people have to learn <\$>, <*>, etc?

```
numCommonFriends x y =
  length <$> (intersect <$> friendsOf x <*> friendsOf y)
```

No, because this

```
numCommonFriends x y = do
  fx <- friendsOf x
  fy <- friendsOf y
  return (length (intersect fx fy))</pre>
```

- can be silently translated to the Applicative form in the compiler
 - (not implemented yet)

 Going further, we could write a pre-processor from this:

```
numCommonFriends :: Haxl Int
numCommonFriends =
  length (intersect (friendsOf sourceId) (friendsOf targetId))
```

• To this:

```
numCommonFriends :: Haxl Int
numCommonFriends
= length <$>
    (intersect <$>
        (join (friendsOf <$> sourceId)) <*>
        (join (friendsOf <$> targetId)))
```

straightforward with haskell-src-exts

Dilemna: monads or no monads?

- Using a preprocessor
 - Advantages
 - Everything is monadic, but looks pure to the programmer
 - Easier to understand
 - Disadvantages
 - Can't write pure code
 - Hard to interpret error messages
 - Two languages adds complexity
- We decided not to go this route (for now)

Technical challenges

- 1. Implicit concurrency
- 2. Implement all the FXL functionality in Haskell
- Translate all the FXL code
- Figure out how to compile+push all the Haskell code to all of the machines in a few minutes

Implement all the FXL functionality in Haskell

- several data sources
- Each needs a Haskell/C++ FFI layer
- ~450 built-in functions
 - Ranging from easy (StrCmp) to really annoying (ParseActivityLog)
- We created lots of tasks
 - some done by the team
 - (we're working on our 4th iteration of the TAO layer)
 - others grabbed by interested people around Facebook: hack-a-month projects and bootcamp

• As of two weeks ago, we have everything implemented!

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Translate all the FXL code

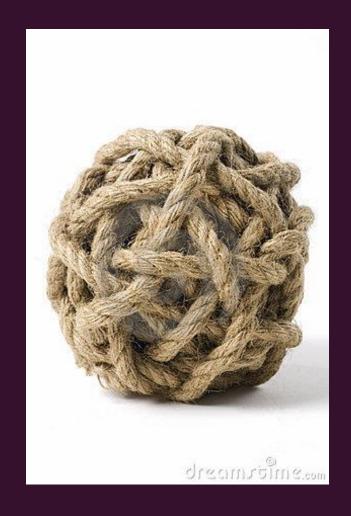
- We have a lot of FXL code
 - impractical to translate it all by hand
- Wrote a translation tool
 - tricky bit is converting to do-syntax or Applicative where necessary, while keeping as much code as possible pure
- Auto-translated code will become the source
 - Try to produce readable code

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Compile time

- At first, compiling the whole translated codebase took ~30 mins
- (FXL push currently takes ~2 mins)

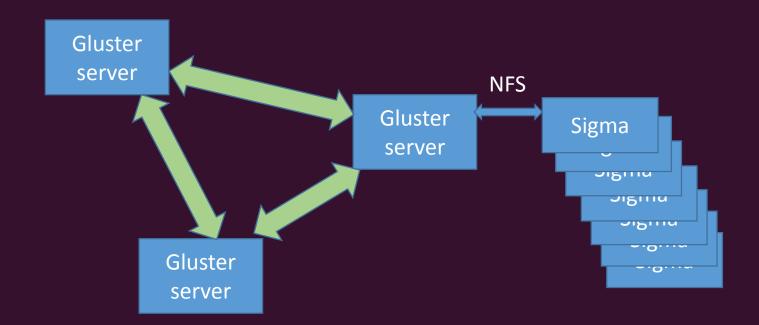


Reducing compile time

- Long laborious process to impose a sensible module structure
 - FXL source files now form a DAG
 - Compilation has some parallelism now
- We have full compile down to ~5 mins
 - incremental compile usually much faster (~2 mins)

How to push to Sigma machines?

- We're experimenting with GlusterFS for distribution.
 - seems good: all machines get a new object in <1 min



Hot code swapping

server code (C++)

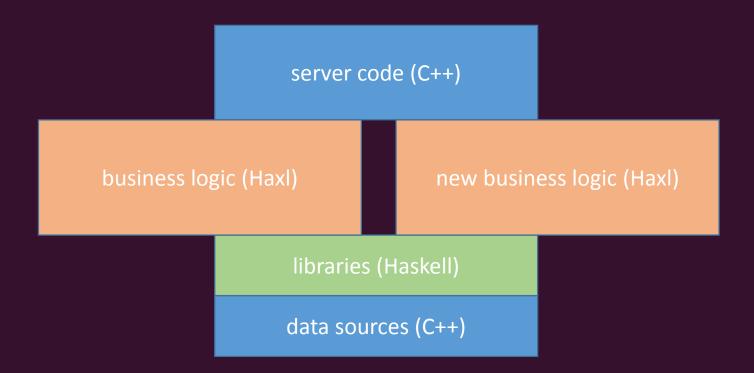
business logic (Haxl)

libraries (Haskell)

data sources (C++)

new business l

Hot code swapping



Hot code swapping

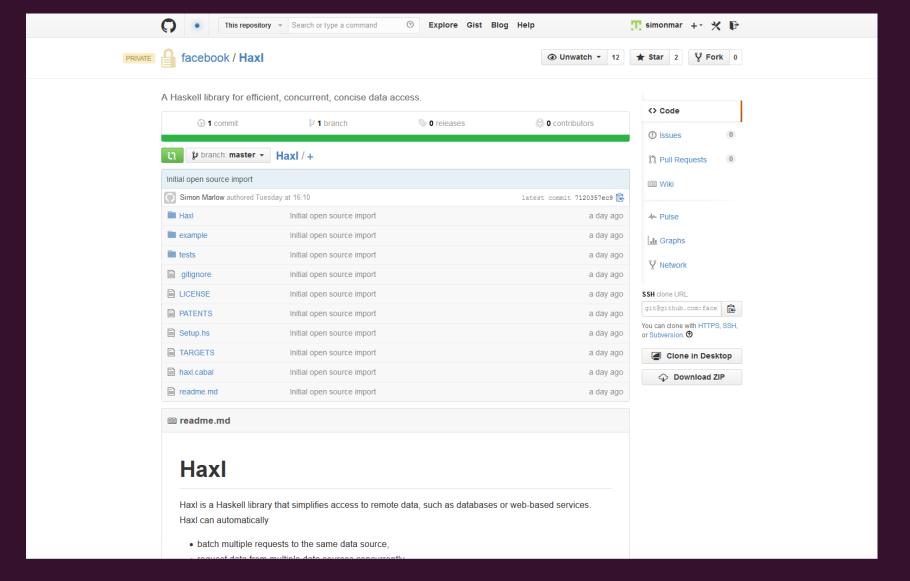
Keep serving requests while we load new code

- Use GHC's built-in linker
 - Had to modify it to unload code (shipped in GHC 7.8)
 - GC detects when it is safe to release old code
 - We can have multiple copies of the code running while existing requests drain

Status

- Call graph complete
- Full FXL codebase translated
- Next goals:
 - achieve 99% correctness (100% hard due to random effects)
 - get performance up par with FXL
 - experiment with running production traffic
 - open source...

Open Source! (coming next week)



Haxl: Haskell at Facebook

Simon Marlow
Jon Coens
Louis Brandy
Bartosz Nitka
Jon Purdy
Aaron Roth
& others
<your name here>



Questions?