

# COMP3401 Computer Organization

Fall 2023– HW#3 [CO2,3]

DUE: 7.12.2023

1) If A, B, C are 4 bit, write the result of the operations. State if the operation results an overflow

a) A=0100, B=0101  
integers

A and B are signed 2's complement

**A+B=?**

A is 4 in decimal

B is 5 in decimal

A + B = 1001 is 9 in decimal overflow occurs

b) A=1010, B=0101  
integers

A and B are signed 2's complement

**A+B=?**

A is -6 in decimal Why? We accept the first digit of the number A as the number that determines the sign. Since the first number of A is 1, the number A is negative. We use the 2's complement formula.

A -> 1010 then 0101 then  $0101 + 1 = 0110$

B is 5 in decimal

A+B = 1111 is -1 in decimal no overflow

c) A=0111, B=1111  
integers

A and B are signed 2's complement

**A-B=?**

A = 0111 is 7 in decimal

B = 1111 is -1 in decimal (2's complement formula)

A - B = 1000 is 8 in decimal overflow occurs

d) A=0111, B=1101

A and B are **unsigned** integers

**A+B = ?**

A = 0111 is 7 in decimal

B = 1101 is 13 in decimal

A+B = 10100 is 20 in decimal Overflow occurs

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2. Manually create a 8 bit fp point representation -similar to IEEE754-, that has 1 sign bit, 3 bits of exponent bits, 5 fraction bits.

- Show the range of (min, max) fp numbers that can be represented by this encoding.
- Represent 3.14 in binary using your own representation.

a. (We created these values as if there was no reserve for the exponent.)

We create manually 1 111 11111

1->sign bit so our value is negative

111-> 7 exponent bit value in decimal

11111-> fraction part is =  $(1/2+1/4+1/8+1/16+1/32)=0,921875$

$(-1)^{\text{sign}} \times (1+\text{fraction}) \times 2^{(\text{exponent}-\text{bias})}$  according to that instruction :

We reach that :

$(-1)^1 \times (1+0,921875) \times 2^{(7-3)} = - 30,75$  minimum range

If we set the sign bit to 0 instead of 1, we can find the max range:

We create manually 0 111 11111

0->sign bit so our value is negative

111-> 7 exponent bit value in decimal

11111-> fraction part is =  $(1/2+1/4+1/8+1/16+1/32)=0,921875$

$(-1)^{\text{sign}} \times (1+\text{fraction}) \times 2^{(\text{exponent}-\text{bias})}$  according to that instruction :

We reach that :

$(-1)^0 \times (1+0,921875) \times 2^{(7-3)} = + 30,75$  minimum range

b. Our value is positive so that our sign bit must be 0

Firstly, we convert 3.14 to binary: 11.00100

Secondly, we normalize the binary representation:  $1.100100 \times 2^1$ .

Thirdly, we adjust the exponent:  $1.100100 \times 2^{(1+3)}=1100.100$

Then we reach exponent bit's value 4 in decimal so now we placed it in the table in binary form -> 100

We obtained the binary equivalent of the fraction part from the binary equivalent of the number: 10010

So, the representation of 3.14 in this 8-bit format is:

0 100 10010

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3. Write the MIPS assembly code of the following C/JAVA code.

```
If (x>=5.0)
    z=2.0* (y-2.0)/(x-5.0);
else
    z= y;
```

where x, y, z are all double precision floating point numbers. Assume x, y, z are in memory locations \$sp, \$sp+8, \$sp+16, respectively. You can use the following and other MIPS instructions in the reference card.

```
ldc1 $f1, $0($sp) # Load a double precision value into register $f1 from the stack at offset 0.
ldc1 $f2, $8($sp) # Load a double precision value into register $f2 from the stack at offset 8.
ldc1 $f3, $16($sp) # Load a double precision value into register $f3 from the stack at offset 16.
ldc1 $f4, const2($gp) #Load a double precision value into register $f4 from global memory at address const2 + $gp.
ldc1 $f5, const5($gp) #Load a double precision value into register $f5 from global memory at location const5 + $gp.
```

```
c.le.d $f5, $f1 #Compare to see if $f5 and $f1 are equal or less.
bc1t Label1 #If the previous comparison holds true, branch to Label 1.
bc1f Label2 #If the previous comparison is not true, branch to Label 2.
```

```
Label1:
sub.d $f2, $f2, $f4 #Subtract $f4 from $f2.
sub.d $f1, $f1, $f5 #Subtract $f5 from $f1.
div.d $f3, $f2, $f1 #Divide $f2 by $f1 and store the result in $f3.
mul.d $f3, $f3, $f4 #Multiply $f3 by $f4.
j endLabel #Jump to endLabel
```

```
Label2:
sdc1 $f2, 8($sp) #At offset 8, place the value of $f2 in the stack.
ldc1 $f3, 8($sp) #A double precision value at offset 8 should be loaded into $f3 from the stack.
j endLabel #Jump to endLabel
```

```
endLabel:
sdc1 $f3, 16($sp) #At offset 16, place the value of $f3 in the stack.
```