# CPSC 202 PSET 9

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## A.9.1: Quadrangle Closure

- 1. Show that every graph G has a unique quadrangle closure.
  - a. Define n s.t. there are n paths with 4 vertices:

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a_{00}a_{01}a_{02}a_{03} \rightarrow \text{add } a_{03}a_{00}
a_{10}a_{11}a_{12}a_{13} \rightarrow \text{add } a_{13}a_{10}
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. . .

 $a_{n0}a_{n1}a_{n2}a_{n3} \to \text{add } a_{n3}a_{n0}$ 

- b. Let's define H:  $H = G \cup a_{01}, a_{02} \cup a_{11}, a_{12} \cup ... \cup a_{n0}, a_{n3}$
- c. Proof of uniqueness:
  - $\exists H' : \forall i \forall H' : a_{io}a_{i3} \in H'$ , therefore, H is either a supergraph or equal to H'. Proof by contradiction:

 $G \subseteq H'$  s.t. H' is a supergraph of G, then H' would have to be the closer of G and  $H \neq H'$ , so  $H \neq H' \cap H \nsubseteq H'$ 

For this to be true, then  $\exists a_{i0}a_{i3} \in H \text{ s.t. } a_{i0}a_{i3} \neq H'$ 

However, then  $a_{i0}a_{i1}a_{i2}a_{i3}$  isn't closed in H' is not closure. Therefore H has to be unique.

2. Show that the quadrangle closure of a bipartite graph is bipartite.

 $a_0 \in S, a_1 \in T, a_2 \in S, a_3 \in T$ . The edges will always run from  $S \to T$  or  $T \to S$ , therefore, the position of the first vertex doesn't matter.

#### A.9.2 Cycles

Show that G is a cycle.

a. Show G has the cycle C

 $C = V_0 V_1 V_2 ... V_n V_0$ , since G was defined as any two vertices with simple paths with no edges in common.

If C weren't contained in G then there wouldn't be two paths between  $V_0$  and  $V_1$ 

b. G doesn't have any vertices or edges in C. Proof G=C by contradiction: If  $\exists w_0$  s.t.  $w_0 \in G$  and  $w_0 \notin C$ , then we would have 2 paths that go from any  $V_i$  in C to  $w_0$  not in C, but then those paths would share an edge.  $w_0$  cannot exist and C=G.

## A.9.3 Deleting a Graph

Show that you can reduce a finite graph  $G_1$  to the empty graph with no vertices by this process  $\iff G_0$  is acrylic.

Prove every vertex of the cycle has degree of at least 2.

- a. If acrylic, it can be deleted:  $\forall v \in C : d(v) \geq 2$
- b. If there are edge points out of the cycle they can be deleted: Proof by induction on |v|:

Base case: |v| = 1

It has a degree of 0 so it can be deleted. For an acrylic graph |v| = n - 1, there exists one vertex of degree 1.

Induction step:  $|v| = 1 \rightarrow |v| = n$  |E| = |v| - 1 for an acrylic graph and  $2|E| = \sum_{v \in V} d(v)$  If  $\forall d(v) \geq 2$ , then  $2|E| = \sum_{v \in V} d(v) \geq \sum_{v \in v} 2$ , so  $2|E| \geq 2|v|$ , and  $|E| \geq |v|$ . Therefore,  $\exists$  at least one vertex of degree 1.