

Melocoton

A Program Logic for Verified Interoperability Between OCaml and C

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OOPSLA 2023, Cascais

27 October, 2023



MAX PLANCK INSTITUTE
FOR SOFTWARE SYSTEMS



AARHUS
UNIVERSITY
DEPARTMENT OF COMPUTER SCIENCE

Multi-Language Programs Are Everywhere



NumPy



Firefox



Python

C++

C

C

Rust

Bindings for:

Fortran

JavaScript

- Rust
- Python
- OCaml
- Go
- ...

The Goal: Verifying Multi-Language Programs

How do we

verify functional correctness

of programs written in

different languages?

Single-Language Functional Correctness

Hoare Logic for simple imperative languages.
Separation Logic for modularity and aliasing.

Multi-Language Functional Correctness

Multi-Language Functional Correctness

Existing work on Semantics and Logical Relations.

How do we prove functional correctness of
individual, potentially unsafe programs?

A Multi-Language Program in OCaml and C

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C business logic

```
void hash_ptr(int * x) {  
    // Implemented in OpenSSL  
    // tedious to port to OCaml  
}
```

A Multi-Language Program in OCaml and C

OCaml business logic

```
let main () =
  let r = ref 42 in
  hash_ref r; (*written in C*)
  print_int !r
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C glue code

```
value caml_hash_ref(value r) {
  int x = Int_val(Field(r, 0));
  hash_ptr(&x);
  Store_field(r, 0, Val_int(x));
  return Val_unit;
}
```

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OCaml glue code

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external hash_ref
  : int ref -> unit
= "caml_hash_ref"
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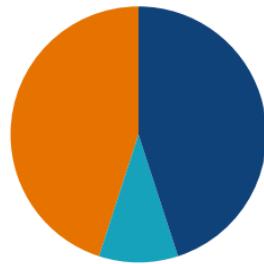
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A Schematic Multi-Language Program

Most multi-language programs look like this:

OCaml business logic
oblivious of C

C business logic
oblivious of OCaml



glue code

where the languages actually interact

We Need to Reason Language-Locally!

Our Contribution: Melocoton

$\lambda_{\text{ML+C}}$ **Program Logic**

Glue Code Verification

$\lambda_{\text{ML+C}}$ **Semantics**

Glue Code Semantics

Common Approach: program logic on top of semantics, **but**

- **Language Interaction:** new semantics and logic for glue code

Our Contribution: Melocoton

OCaml* Program Logic

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Glue Code Verification

C* Program Logic

OCaml* Semantics

$\lambda_{\text{ML+C}}$ Semantics

Glue Code Semantics

C* Semantics

Common Approach: program logic on top of semantics, **but**

- **Language Interaction:** new semantics and logic for glue code
- **Language Locality:** embed existing semantics and logics

* simplified/idealized versions of OCaml and C

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Language Interaction: Different Views of the Same Data

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How is **OCaml data** accessed from **C glue code**?

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How is **OCaml** data accessed from **C** glue code?

High-level **OCaml** values are accessed..
..through a **low-level block representation**.

Language Interaction: Semantics

High-level OCaml value \sim_{ML} Low-level block representation

Language Interaction: Semantics

High-level OCaml value	\sim_{ML}	Low-level block representation
integers	\sim_{ML}	integers
booleans	\sim_{ML}	integers (0 or 1)

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Language Interaction: Semantics

High-level OCaml value	\sim_{ML}	Low-level block representation
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λ_{ML+C} Semantics

$\sigma : Heap_{ML}$

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Language Interaction: Program Logic, Take 1

$\lambda_{\text{ML+C}}$ Semantics

$\sigma : \text{Heap}_{\text{ML}}$



$\zeta : \text{BlockHeap}$

Language Interaction: Program Logic, Take 1

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$\lambda_{\text{ML+C}}$ **Semantics**

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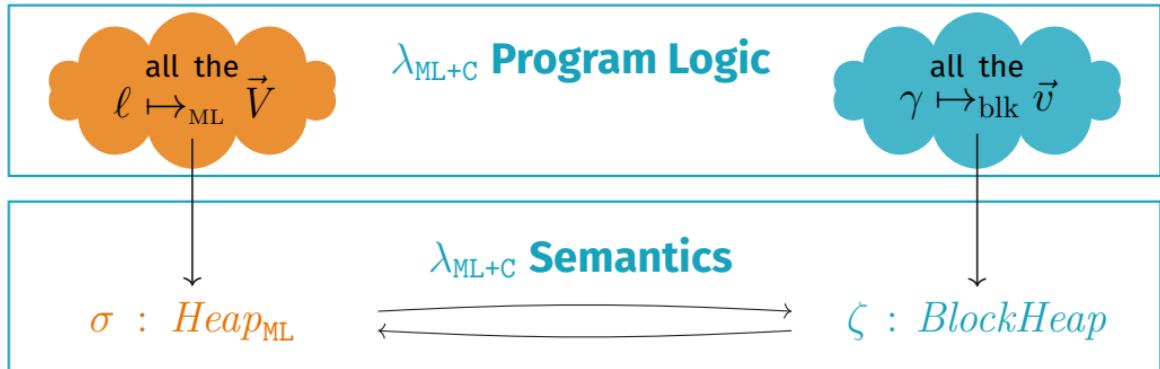


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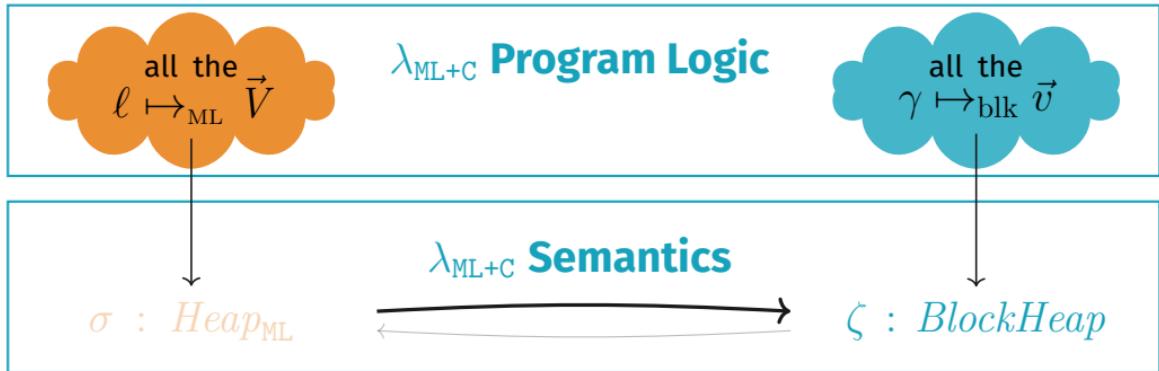
Language Interaction: Program Logic, Take 1

 $\ell \mapsto_{\text{ML}} \vec{V}$ $\lambda_{\text{ML+C}}$ **Program Logic** $\gamma \mapsto_{\text{blk}} \vec{v}$ $\sigma : \text{Heap}_{\text{ML}}$ $\lambda_{\text{ML+C}}$ **Semantics** $\zeta : \text{BlockHeap}$

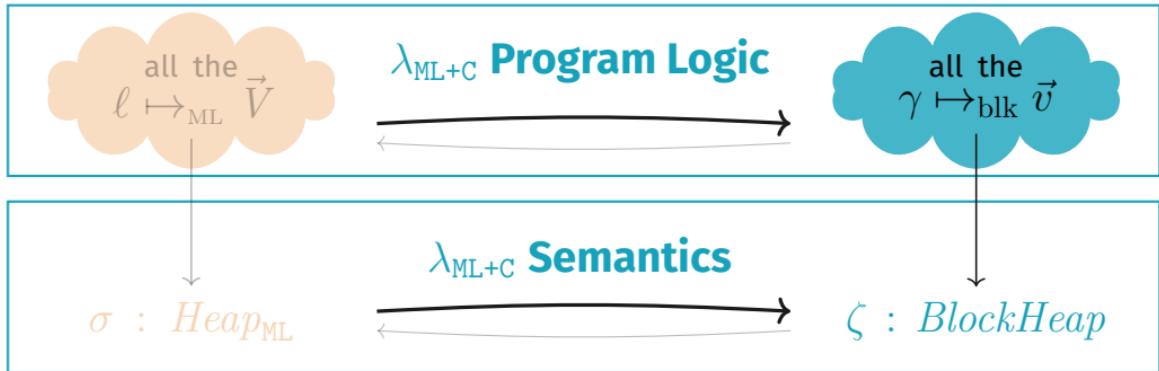
Language Interaction: Program Logic, Take 1



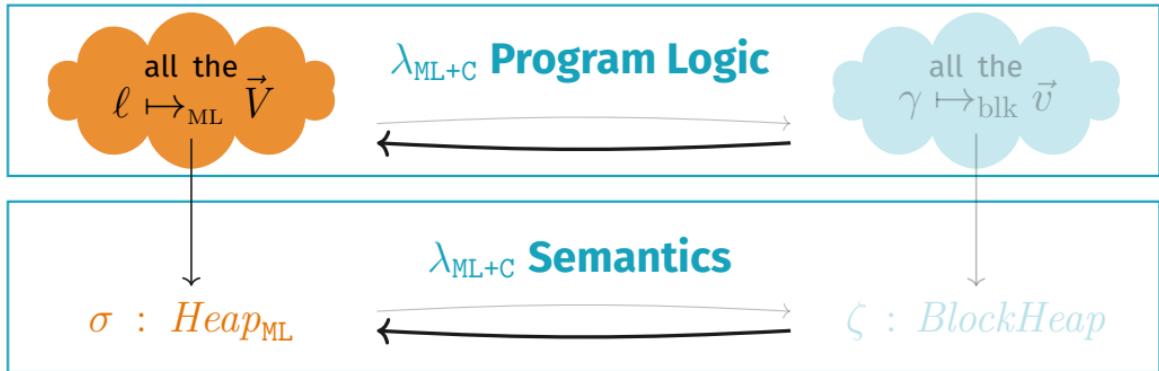
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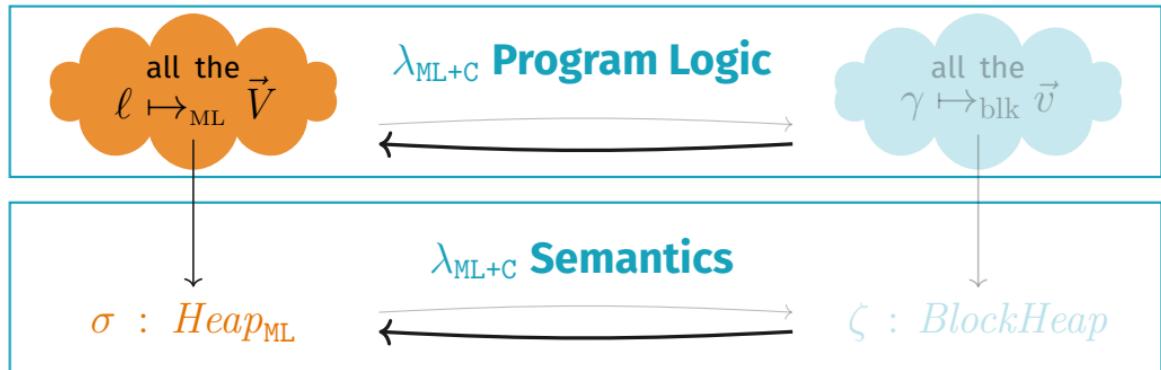
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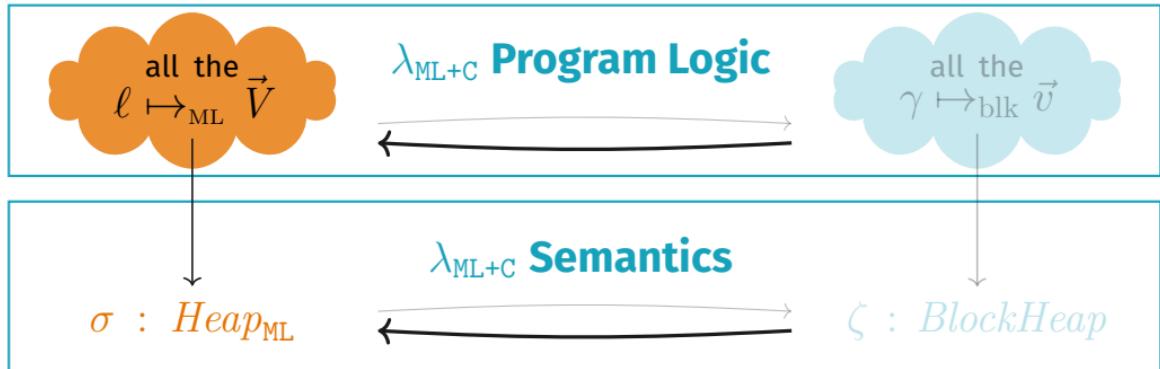


EXTCALL

{
 all
} C function body {
 all
}

{
 all
} call into C {
 all
}

Language Interaction: Program Logic, Take 1



EXTCALL

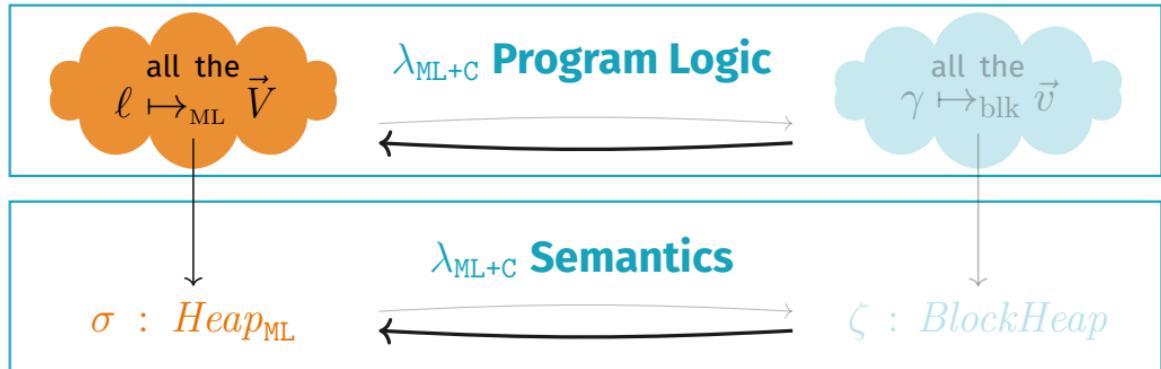
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$$\frac{\{P\} e \{Q\}}{\{R * P\} e \{Q * R\}}$$

Language Interaction: Program Logic, Take 1



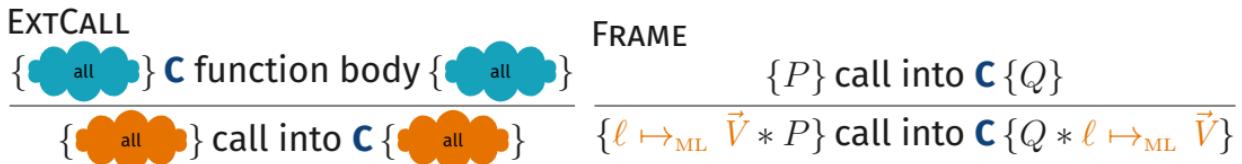
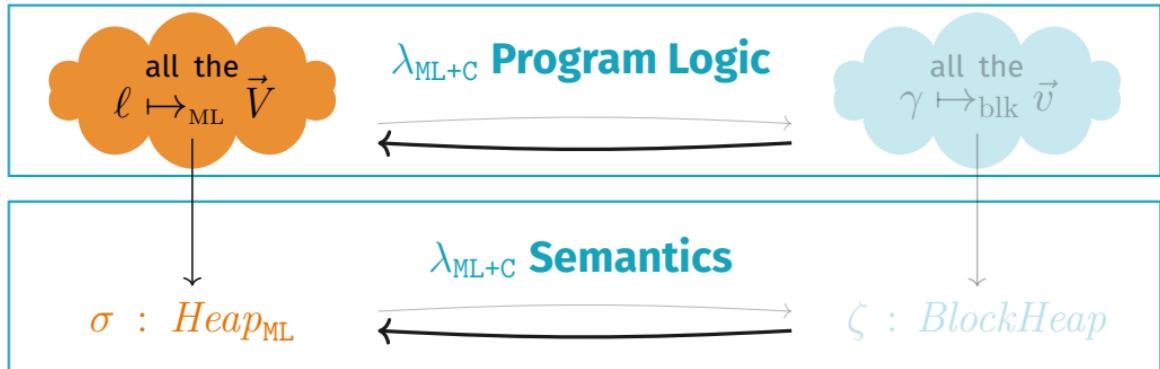
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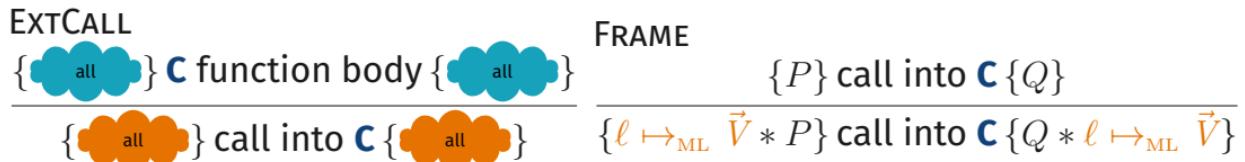
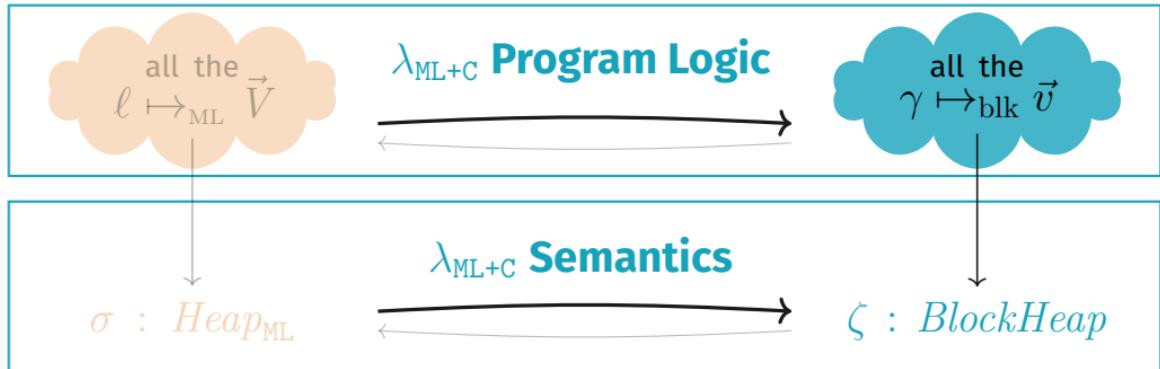
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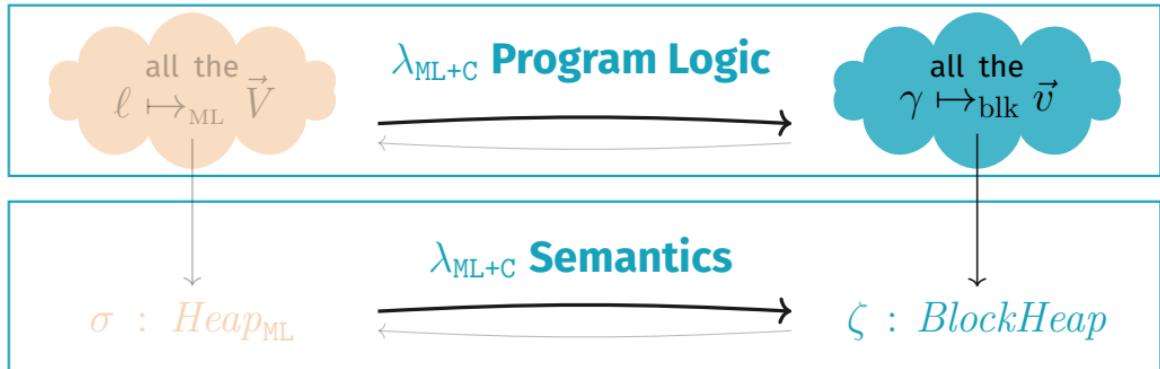
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Language Interaction: Program Logic, Take 1



Language Interaction: More Gradual Rules

OCaml points-toes remain valid when switching to C!

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$$\ell \mapsto_{\text{ML}} \vec{V}$$

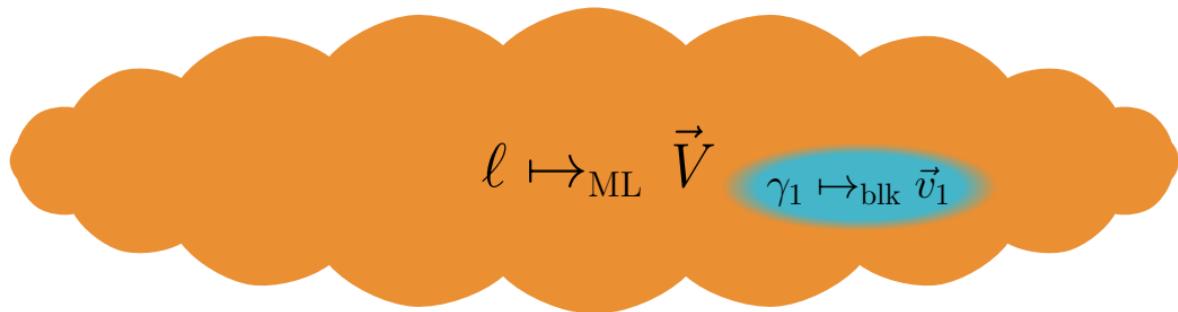
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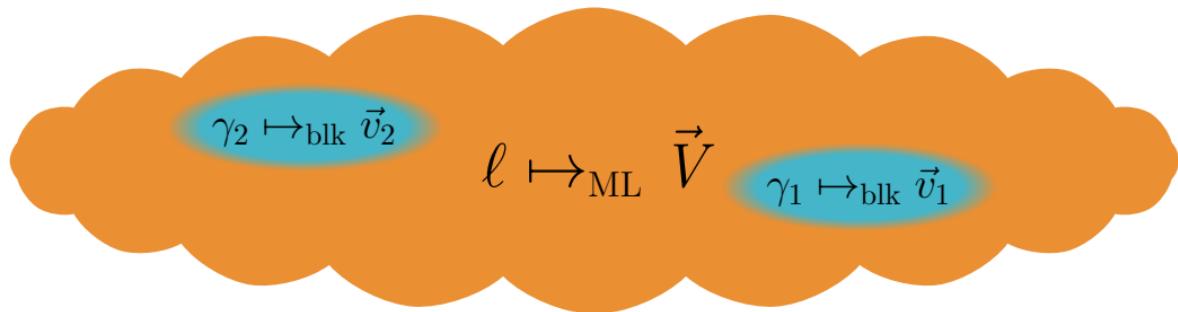
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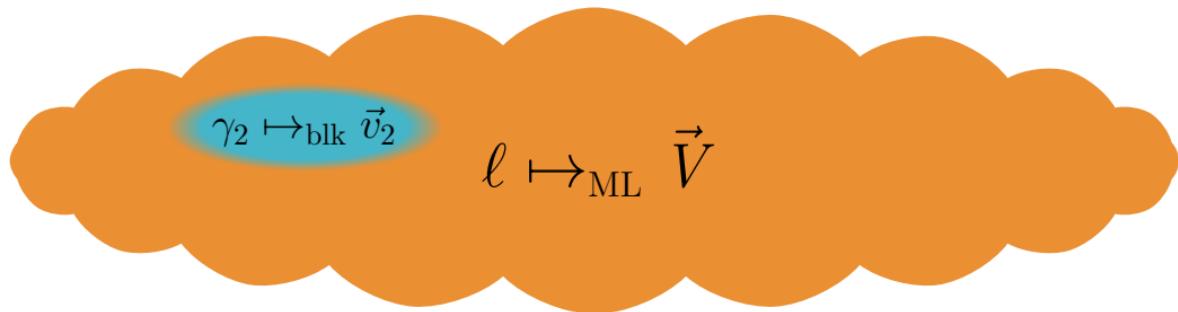
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View Reconciliation Rules for Converting On-Demand:

$$\ell \mapsto_{\text{ML}} \vec{V} \Rightarrow \exists \gamma \vec{v}. \gamma \mapsto_{\text{blk}} \vec{v} * \ell \sim_{\text{ML}} \gamma * \vec{V} \sim_{\text{ML}} \vec{v}$$

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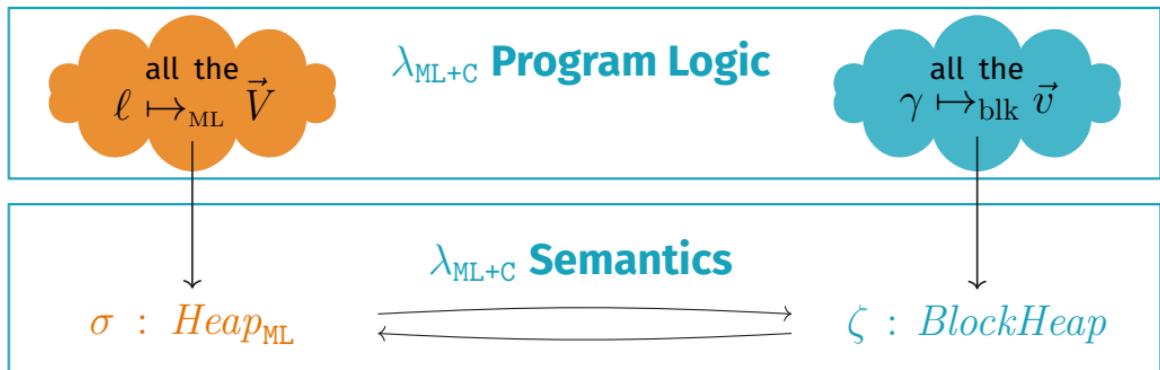
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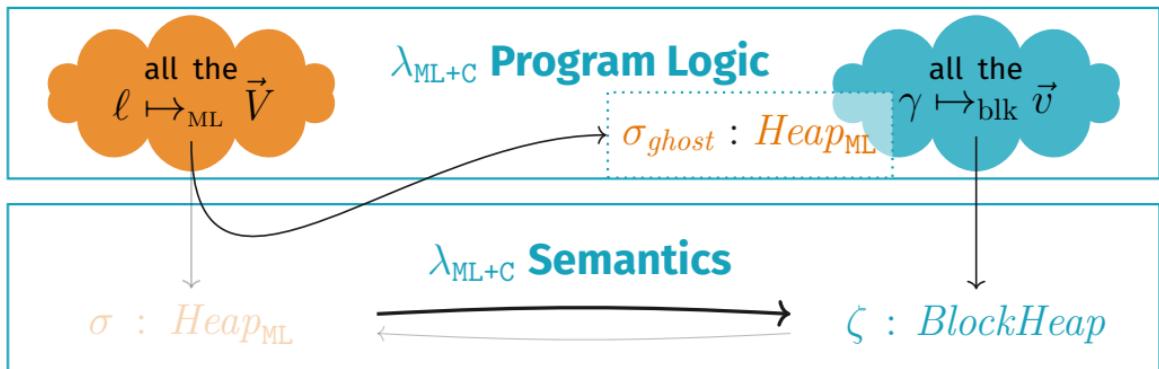
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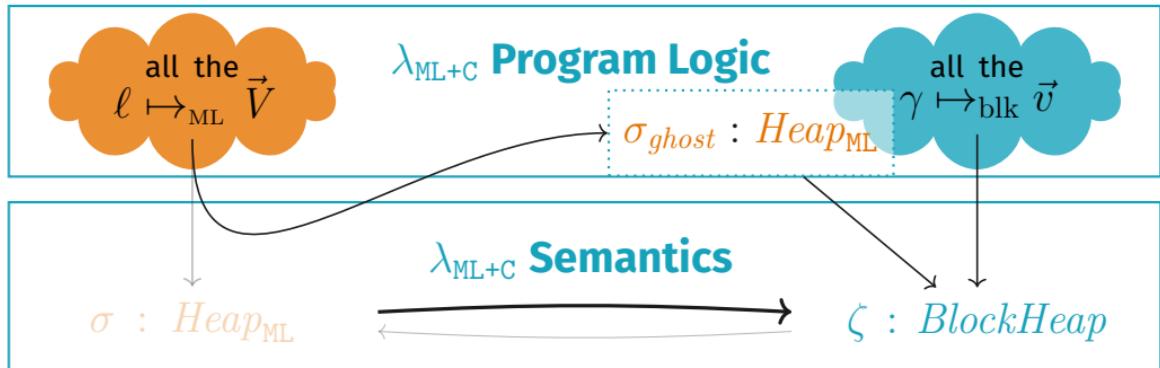


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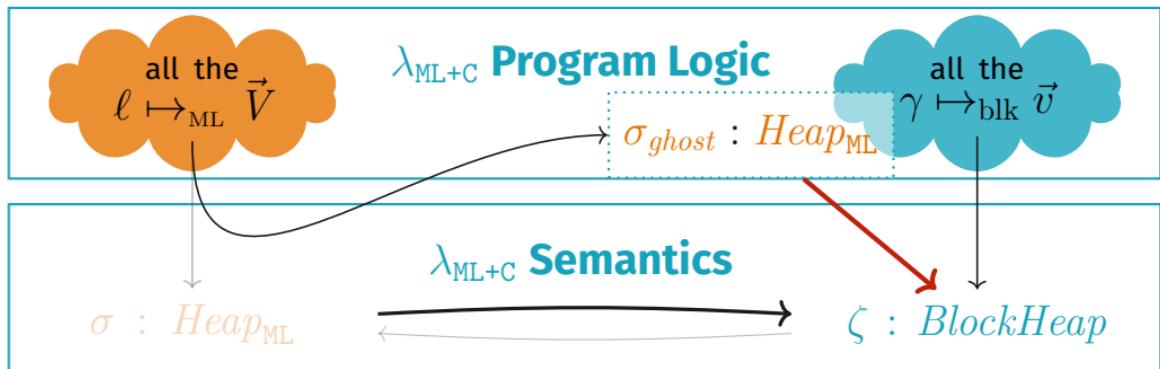


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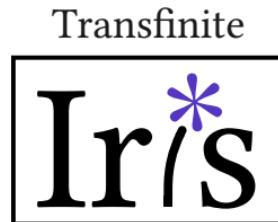
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More in the paper ...

- Language-local reasoning for **external calls**.
- Additional **OCaml FFI features**: garbage collection, registering roots, custom blocks, callbacks, etc.
- **Case studies** utilising all of these features.
- **Step-indexed logical relation** to prove OCaml type safety of external C functions.



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Language Locality: Embed Existing Languages

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$\lambda_{\text{ML+C}}$ Program Logic

C Program Logic

Glue Code Verification

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<https://melocoton-project.github.io>