

# Decentralized Finance

## DeFi Security

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 Stanford  
University



Imperial College  
London



 UNIVERSITY OF  
ILLINOIS  
URBANA-CHAMPAIGN



Berkeley  
UNIVERSITY OF CALIFORNIA



# DeFi Security Affects Multiple Layer

Third party Layer

**UI**

Wallet, Website, APIs

**Other**

Oracle data feed, Centralized governance

DeFi Protocol + Application Layer

**Asset**

Fungible, Non-Fungible

**Atomic Composable DeFi**

Exchange, Loan, Mixer, Liquidity incentive

Smart contract Layer

**Data**

Block, Transaction, Contract

**Virtual Machine**

Contract execution, State transition

Blockchain Layer

**Consensus**

Proof-of-Work, Proof-of-Stake

**Incentive Protocol**

Block reward, MEV reward, TX fee

Network Layer

**Network Services**

DNS, IP, BGP

**Network Protocols**

P2P overlay, Peer discovery, Data propagation

 I can attack any layer!





# Network Layer Security

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# Network Layer

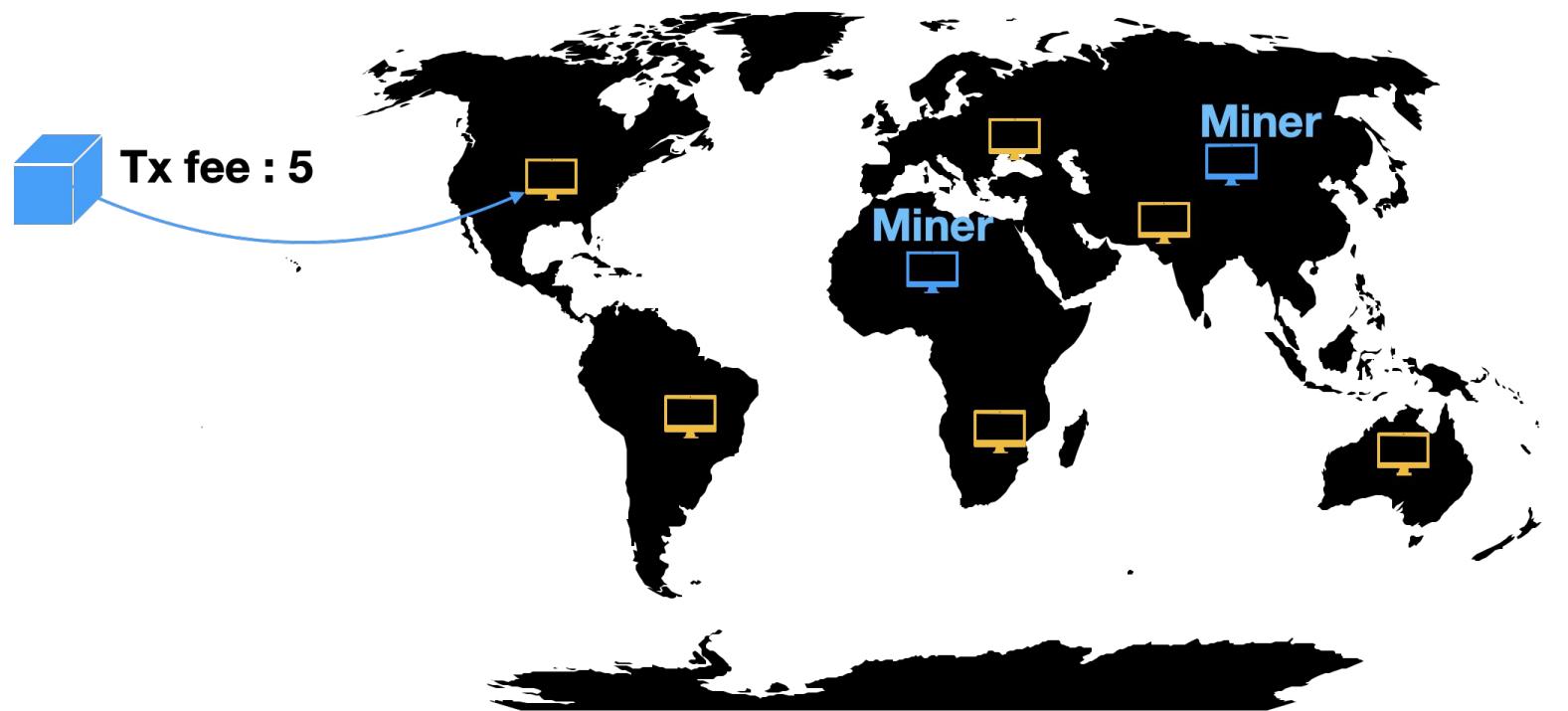
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- Why Network Layer?
  - Information dissemination and propagation.
  - Latency matters!
- How many nodes?
  - Bitcoin: about 10'000 reachable full nodes (TCP/8333)
  - Ethereum:
  - Dogecoin:
- What type of nodes exist?
  - Full nodes
  - Light nodes

# Exchange Transaction Propagation

Trader

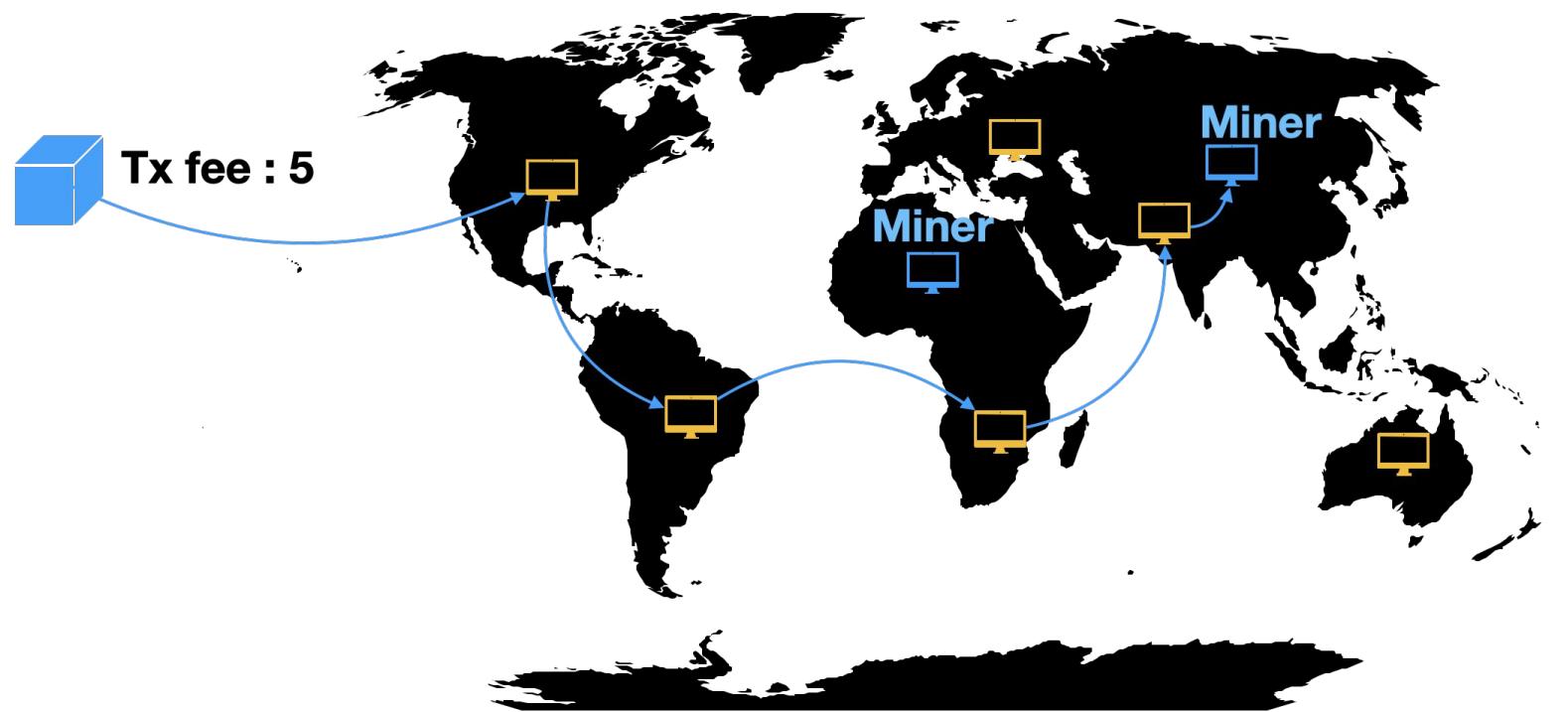
P2P Network



# Exchange Transaction Propagation

Trader

P2P Network

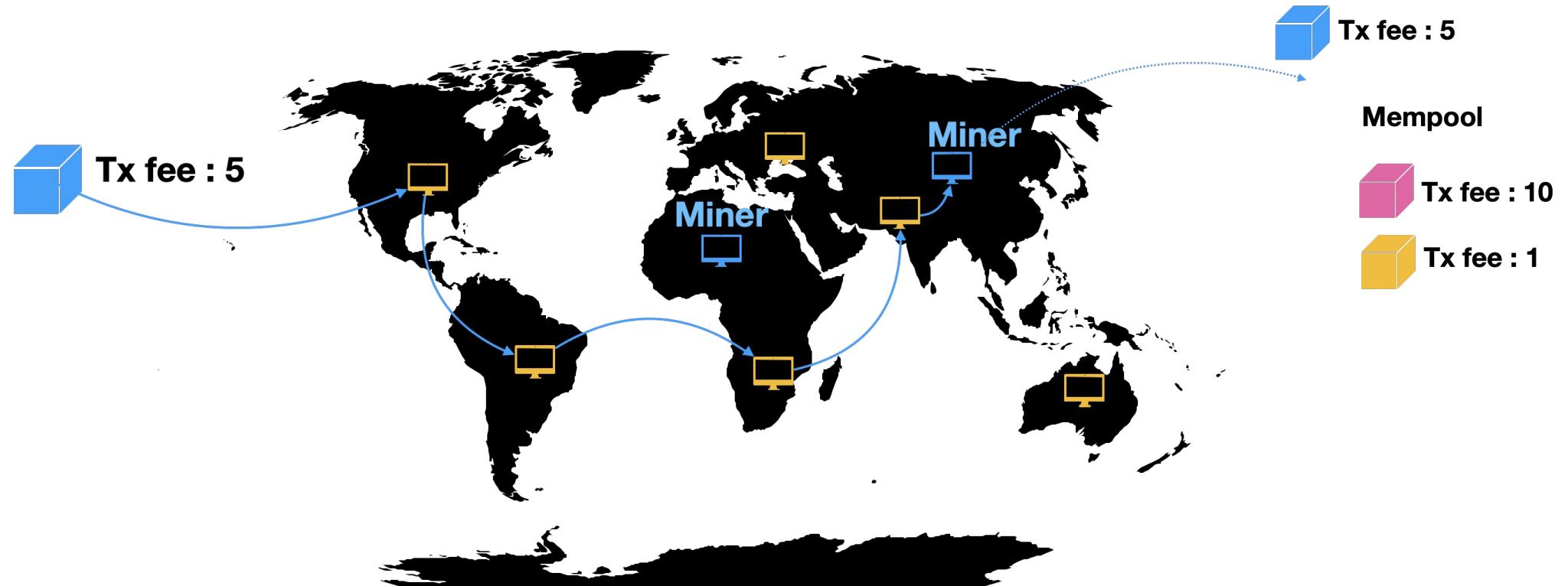


# Exchange Transaction Propagation

Trader

P2P Network

Elected Leader/Miner

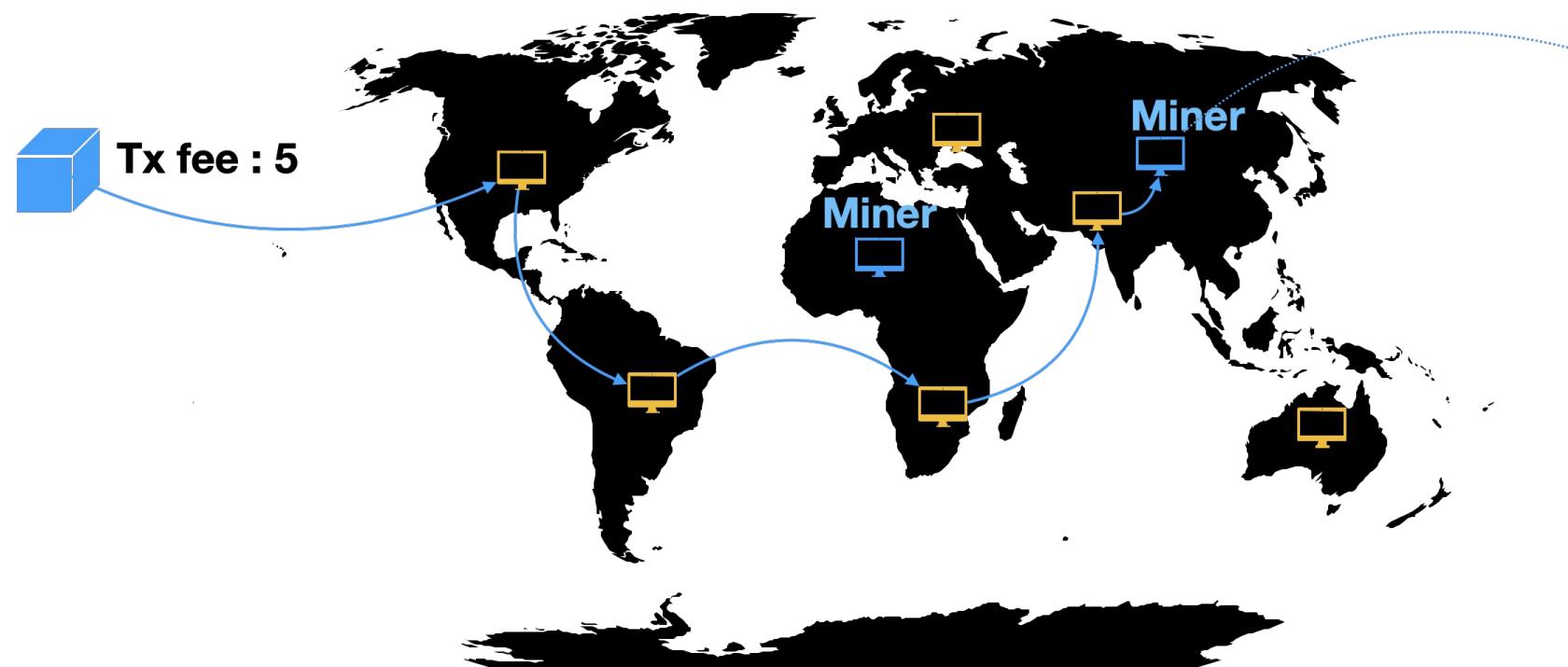


# Exchange Transaction Propagation

Trader

P2P Network

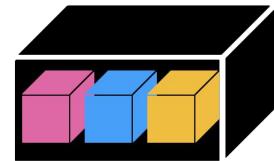
Elected Leader/Miner



Mempool

- Tx fee : 10
- Tx fee : 5
- Tx fee : 1

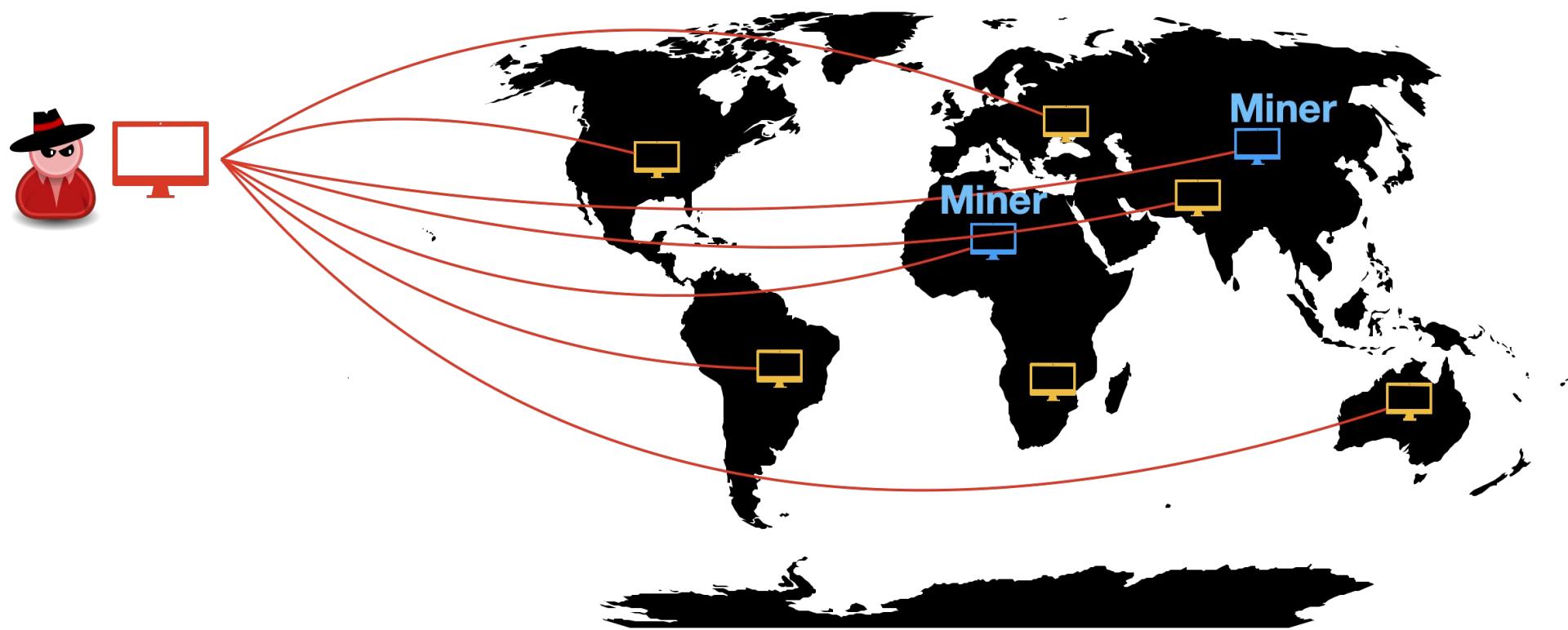
Final Block



# Network Layer – Spy Node

Attacker (Spy Node)

P2P Network

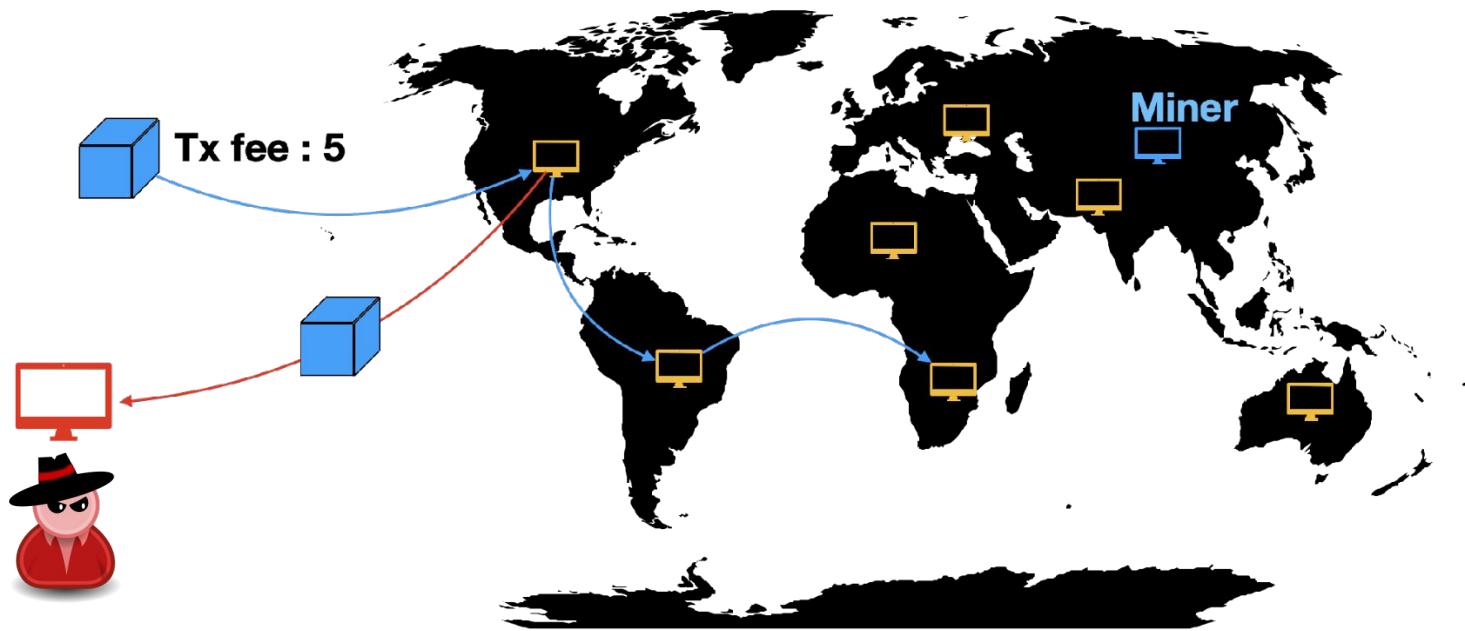


# Network Layer – Spy Node

Trader

P2P Network

Elected Leader/Miner



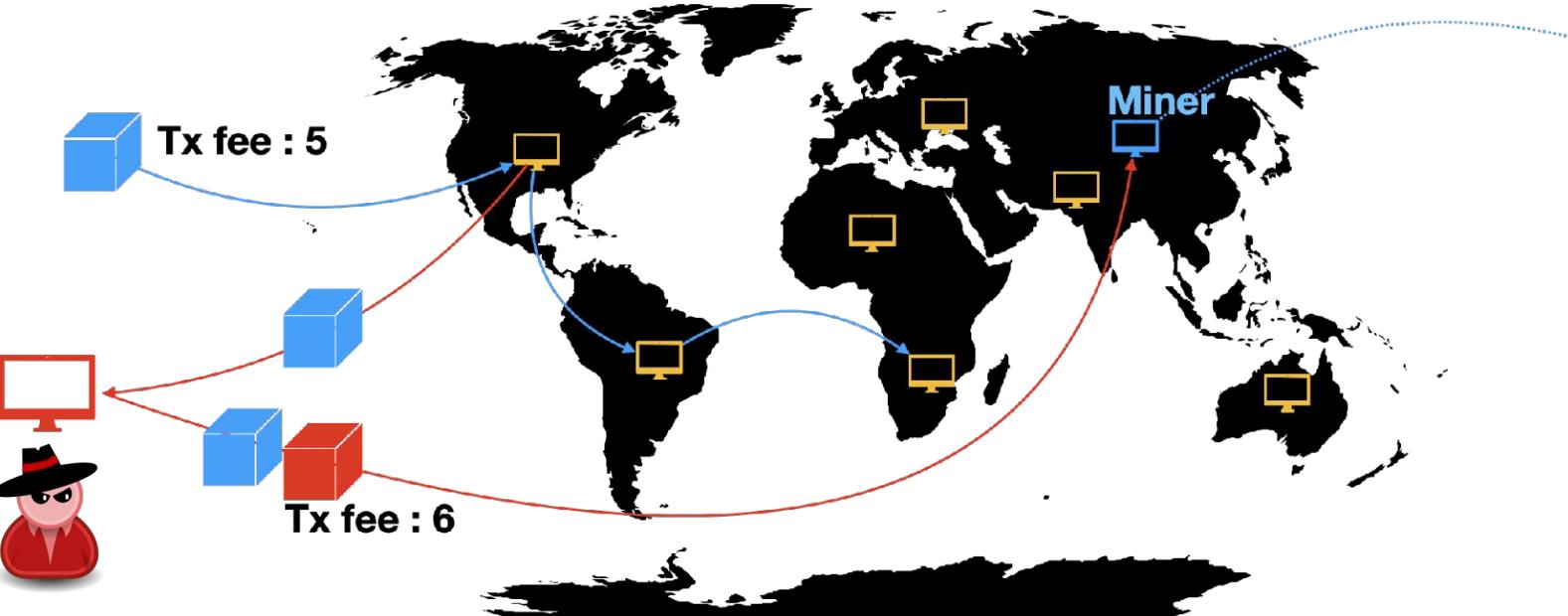
Mempool

Tx fee : 10

Tx fee : 1

# Front-running

Trader



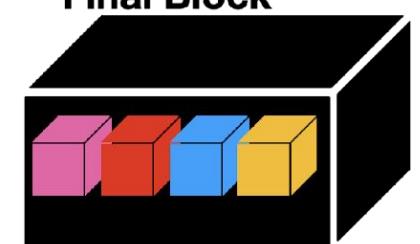
P2P Network

Elected Leader/Miner

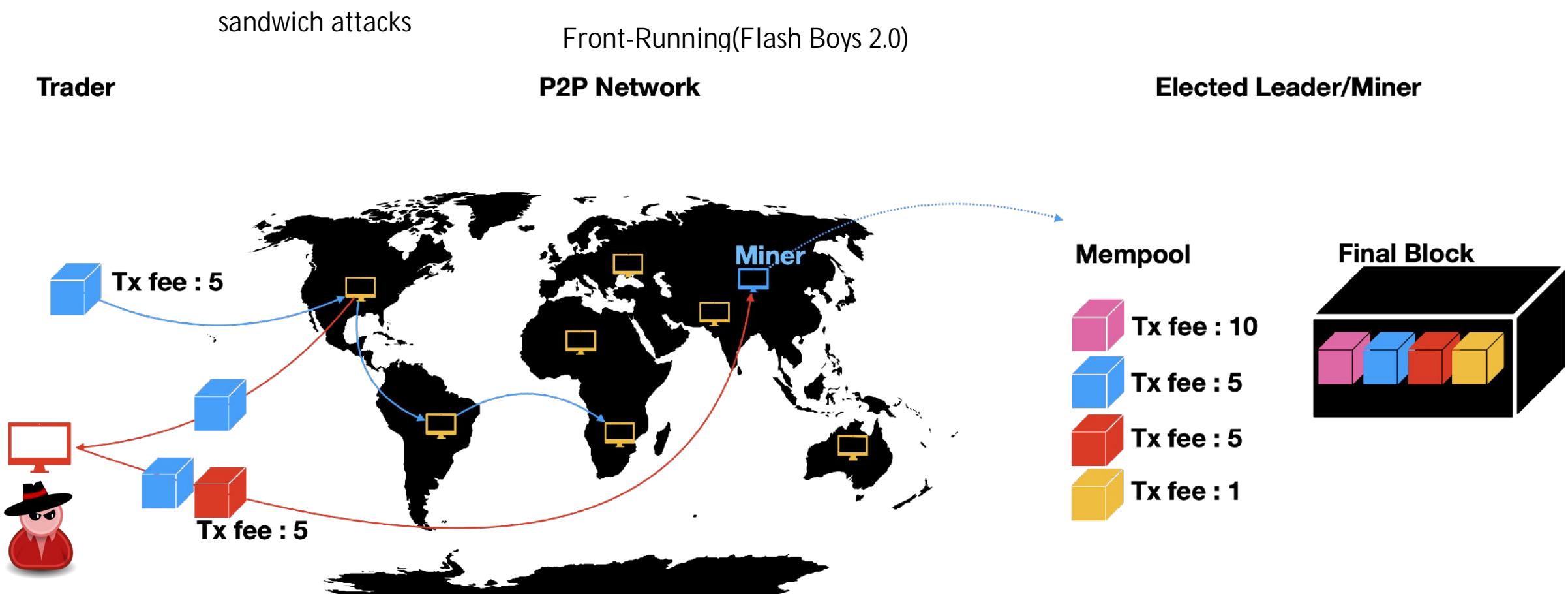
Mempool

	Tx fee : 10
	Tx fee : 6
	Tx fee : 5
	Tx fee : 1

Final Block



# Back-running

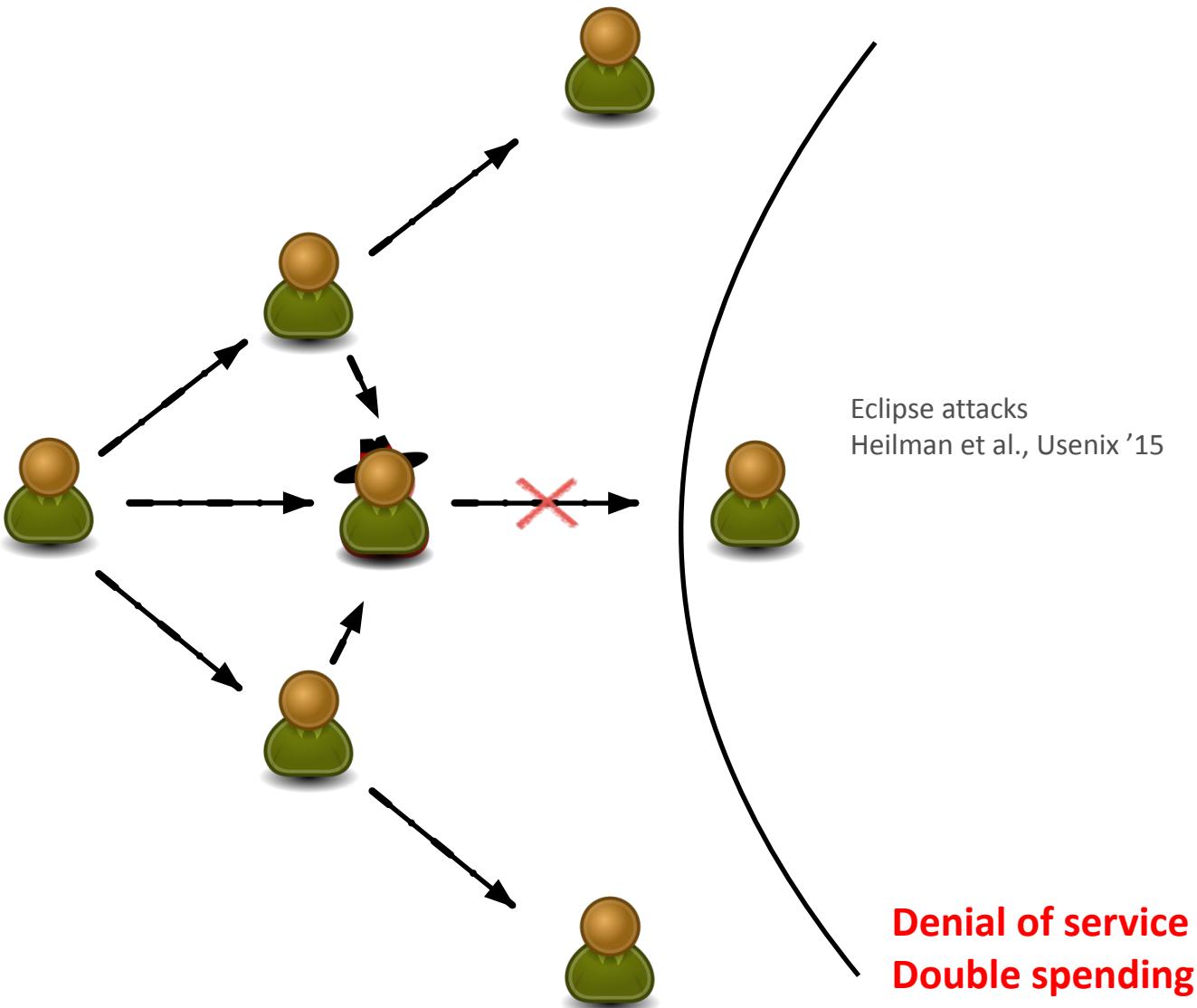




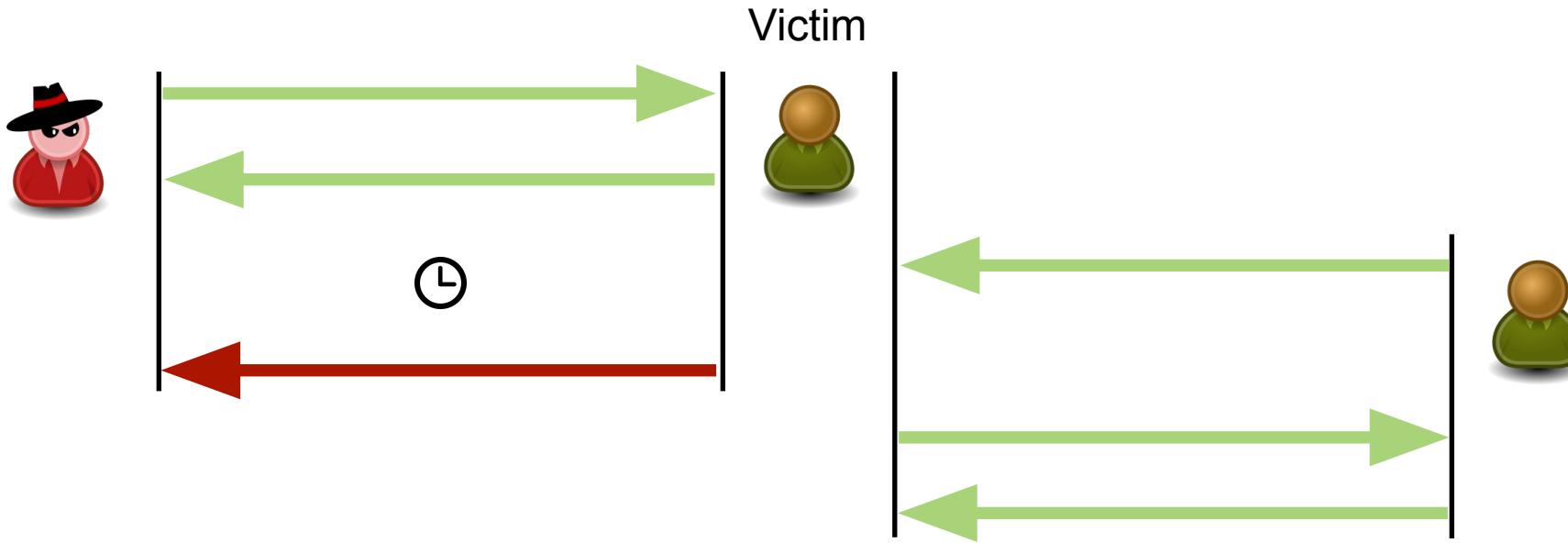
# Eclipse Attacks

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# Eclipse Attacks



# Request timeouts



Block timeout: 20 minutes

Transaction timeout: 2 minutes

# Security Implications

- Adversary

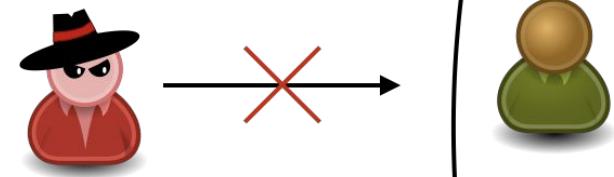
- Blinds victim from blocks and transaction > 20 min
- Experimental validation

- Impact

- Double spend transactions
- Aggravated selfish mining
- Network wide Denial of Service

- Mitigations

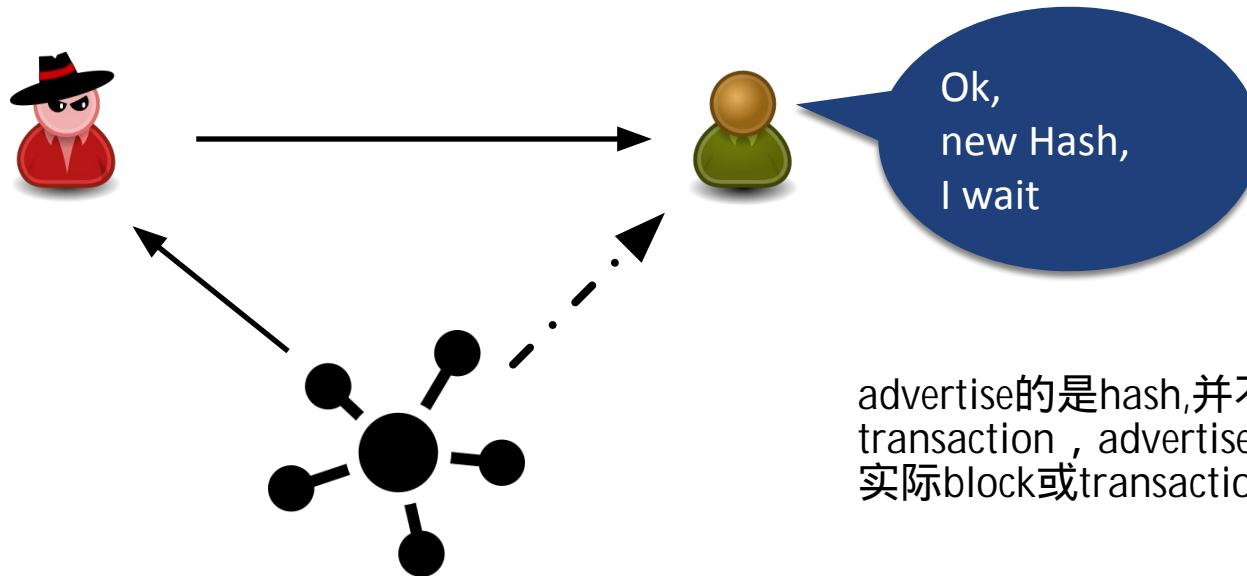
- Hardening measures
- Estimate waiting time for secure transactions



# Eclipse Requirements

所必须得条件

1. Must be **first** peer to advertise Transaction/Block



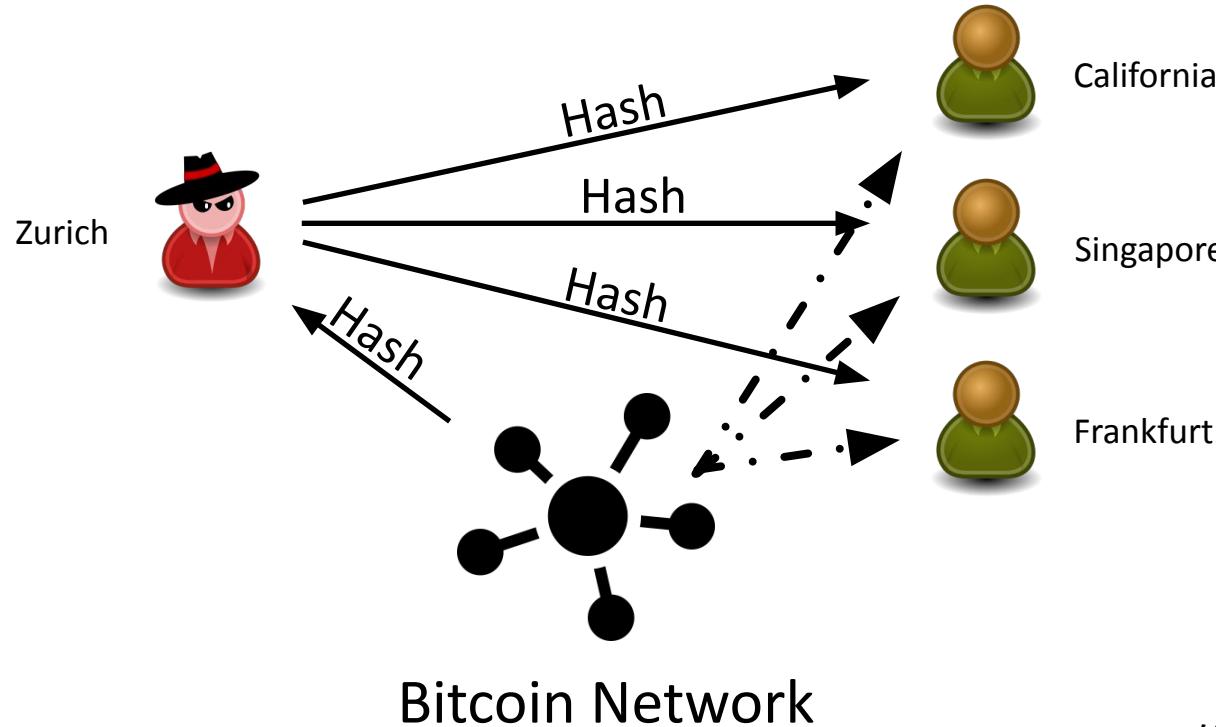
advertise的是hash,并不是实际的block或者transaction , advertise的hash确认后会进行后续实际block或transaction的传递

后续其他节点的advertise , 就会被接收但并不接受数据传输

2. Victim should wait

- Block timeout: 20 minutes
- Transaction timeout: 2 minutes

# Being First on the Network Layer



一般来说，建立的攻击conn越多，那么就会有越高的成功率成为第一个建立的连接

Connections of Adversary	40	80	200	800
Connections of Victim	40	40	40	40
Average success in being first	0.44± 0.14	0.57± 0.20	0.80± 0.14	<b>0.89± 0.07</b>

# Network Layer Timeouts

- Transactions

- After 2 minutes request from other peer (FIFO)

FIFO queue



面对TX和blocks有不同的  
具体阻塞策略

- Blocks (older Bitcoin version)

- After 20 minutes disconnect and do nothing
  - If received header, disconnect and request block from another peer



# Blockchain Layer Security

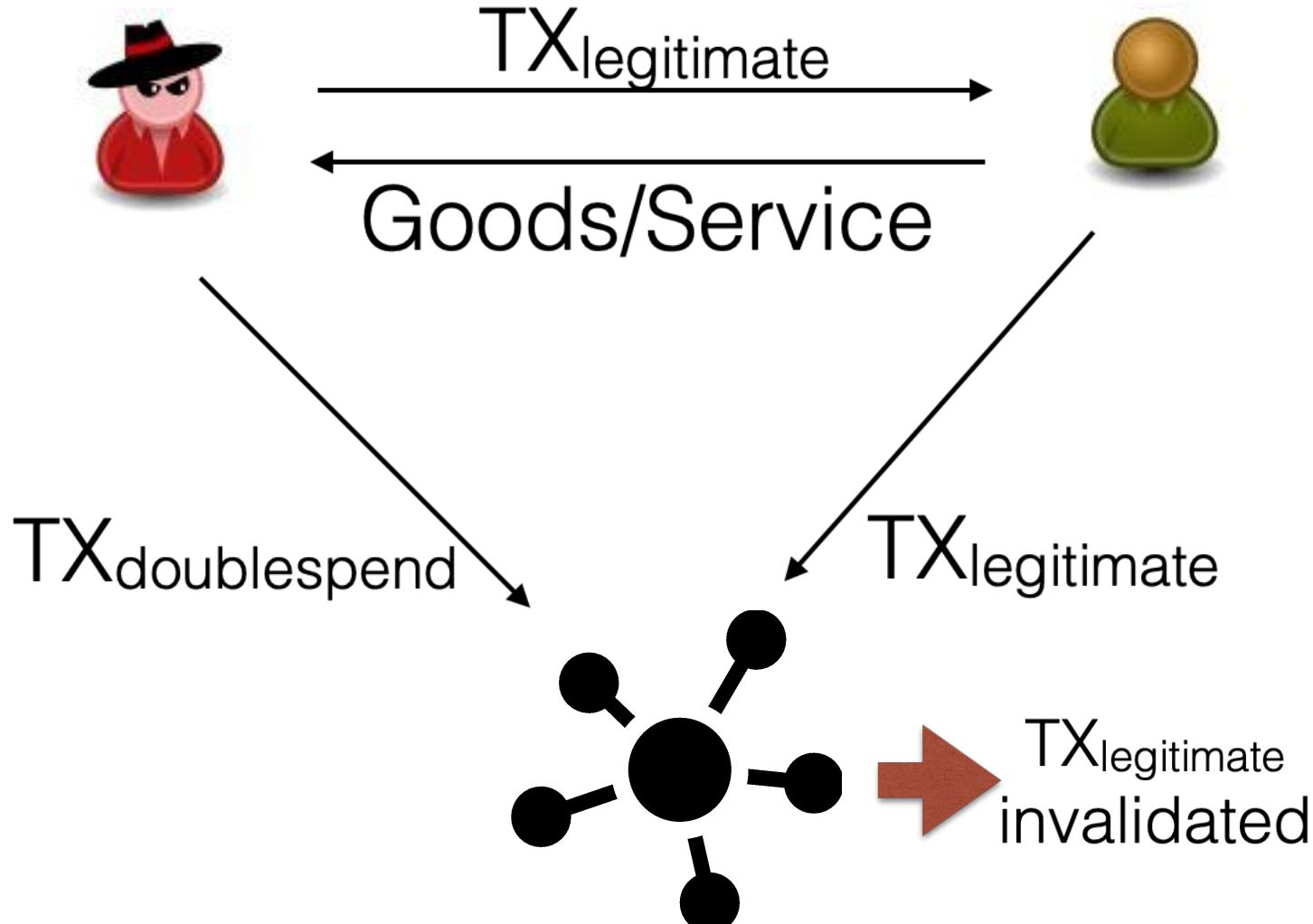
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# Why Blockchain Layer?

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- Double-Spending
- Selfish Mining
- Undercutting
- Bribery

# Double-Spending

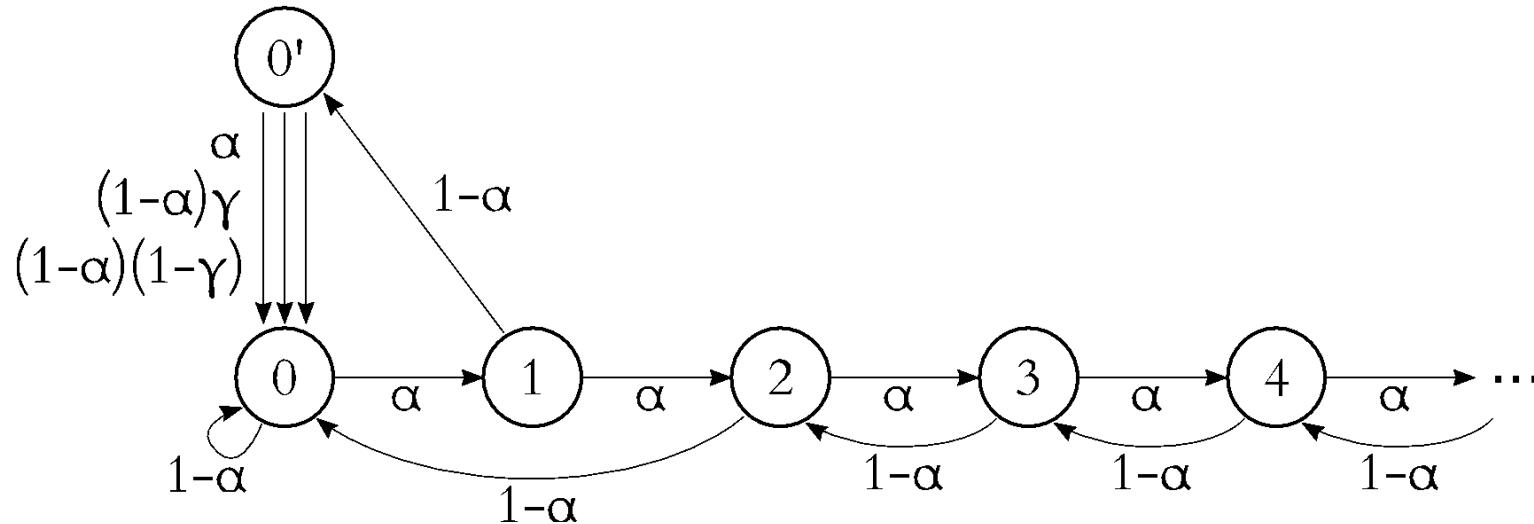


# Increasing Mining Advantage with an Eclipse

- Idea from Eyal et. al:

暂时不去深纠 selfish-mining 和eclipse的混合体

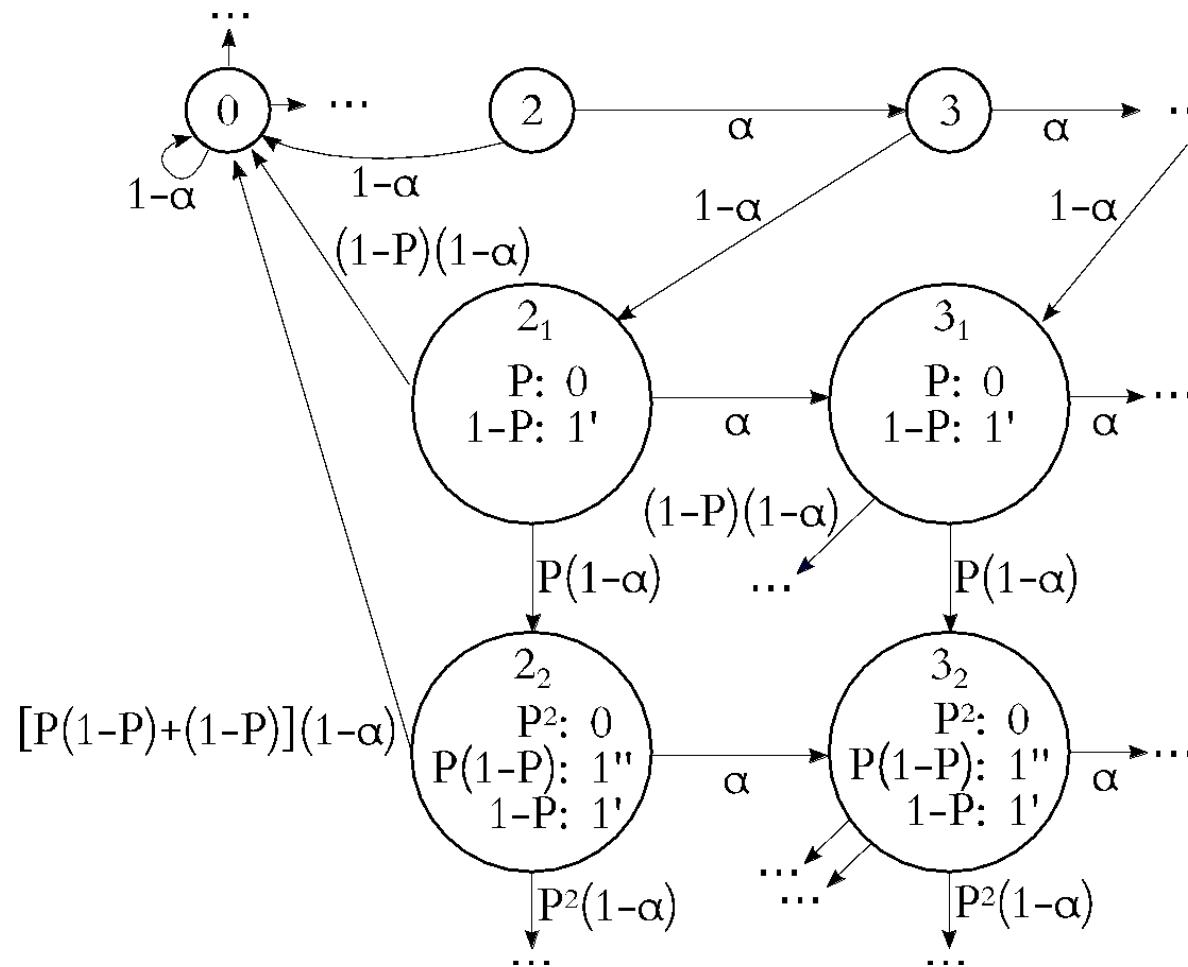
- Instead of publishing, keep a block private
- Other miners will perform wasteful computations



$\alpha$  : hashing power of adversary

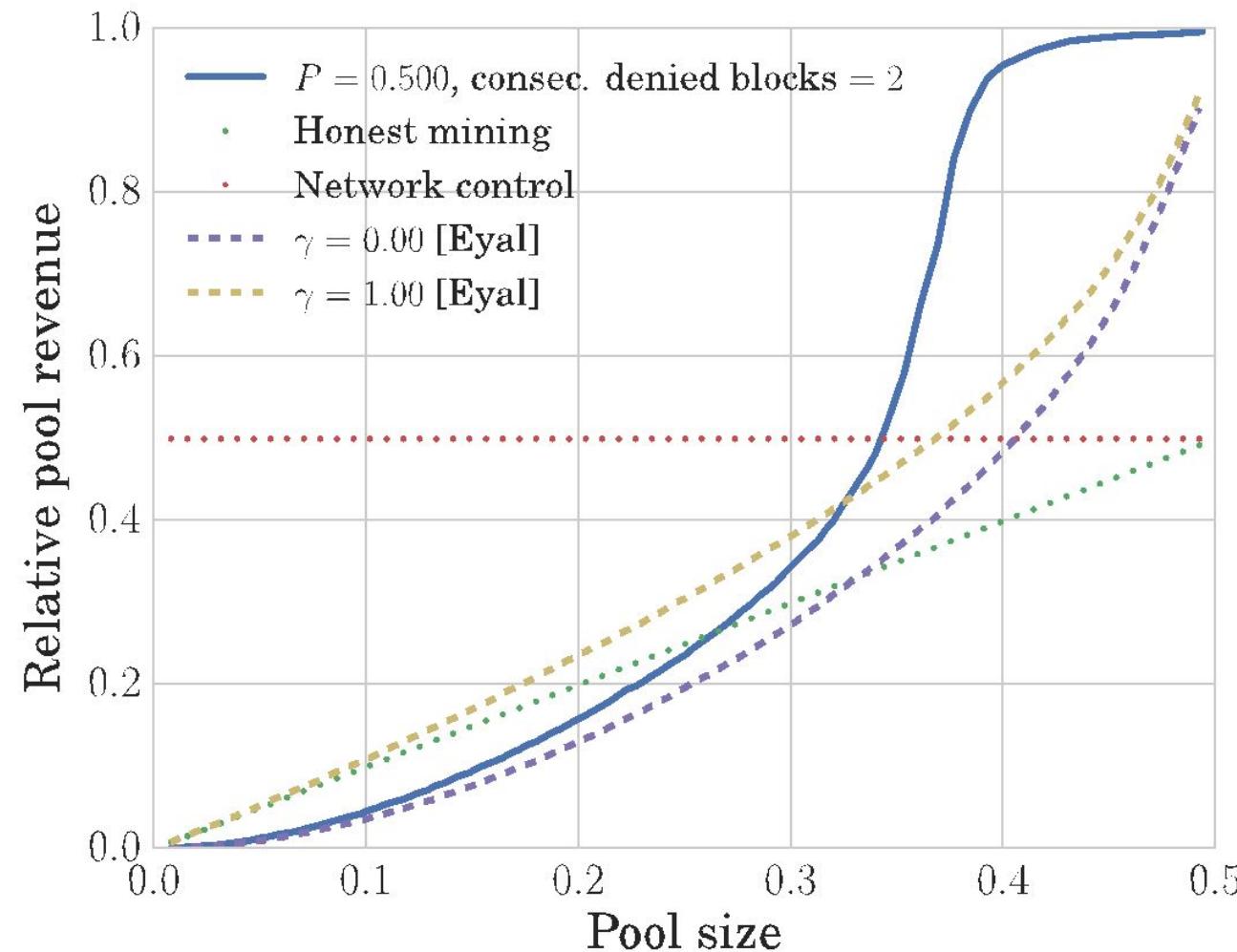
$\gamma$  : propagation parameter

# Increasing Mining Advantage with an Eclipse



P: probability to eclipse a block to a miner

# Increasing Mining Advantage with an Eclipse





# Smart Contract Layer Security

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# Smart Contract Layer

看到了这里

```
contract Wallet {  
    uint balance = 10;  
  
    function withdraw(){  
        if(balance > 0)  
            msg.sender.call.value(balance)();  
        balance = 0;  
    } }
```

Transfer \$\$\$  
to the caller

- Programs that handle money
  - Executed on a blockchain, written in a high-level language, compiled to VM code
- No patching after release
- What can go wrong?

# The DAO attack

Funds Stolen From the DAO One Year Ago Would be Worth \$1.35bn Today

JP Buntinx June 18, 2017 Crypto.News

The screenshot shows a website for "Etherdice". On the left, there's a large Ethereum logo and a smaller icon with a die and the number 33. A message box contains the text: "Etherdice is down for maintenance. We are having troubles with our smart contract and will probably need to invoke the fallback mechanism." Below this, a section titled "King of the Ether Throne" describes the DApp as granting riches and immortalizing names. An "Important Notice" section highlights a "SERIOUS ISSUE" where monarch compensation payments might not be sent. It advises users not to send payments to the contract and that refunds will not be made after the issue was identified on 2016-02-07.

Etherdice is down for maintenance. We are having troubles with our smart contract and will probably need to invoke the fallback mechanism.

## King of the Ether Throne

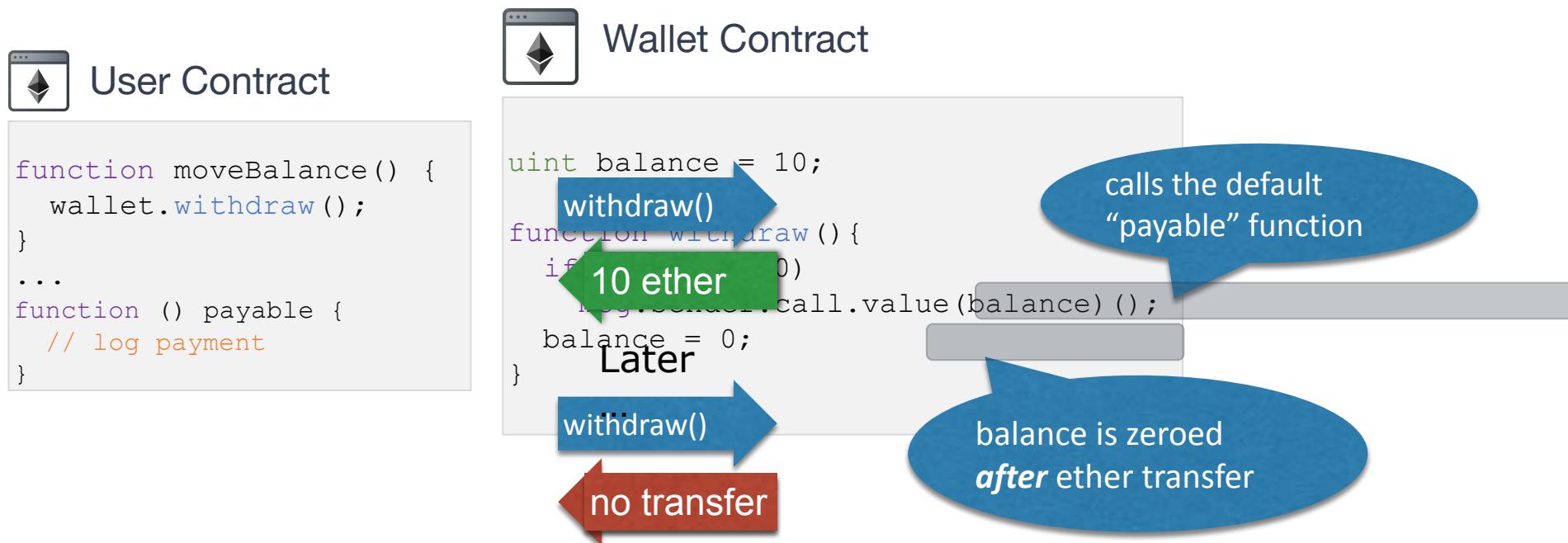
An Ethereum DApp (a "contract"), living on the blockchain, that will make you a King or Queen, might grant you riches, and will immortalize your name.

### Important Notice

A SERIOUS ISSUE has been identified that can cause monarch compensation payments to not be sent.

DO NOT send payments to the contract previously referenced on this page, or attempt to claim the throne. Refunds will CERTAINLY NOT be made for any payments made after this issue was identified on 2016-02-07.

# Security Bug #1: Reentrancy



Can the user contract withdraw more than 10 ether?

# Security Bug #1: Reentrancy



```
function moveBalance() {  
    wallet.withdraw();  
}  
  
function () payable {  
    wallet.withdraw();  
}
```



Calls withdraw()  
before balance  
is set to 0



Wallet Contract

```
uint balance = 10;  
  
function withdraw() {  
    if(balance > 0)  
        msg.sender.call.value(balance)();  
    balance = 0;  
}
```

balance is zeroed  
*after* ether transfer

An adversary stole 3.6M Ether !

# Security Bug #2: Unprivileged write to storage

Any user may change the wallet's owner



## Wallet Contract

```
address owner = ...;

function initWallet(address _owner) {
    owner = _owner;
}

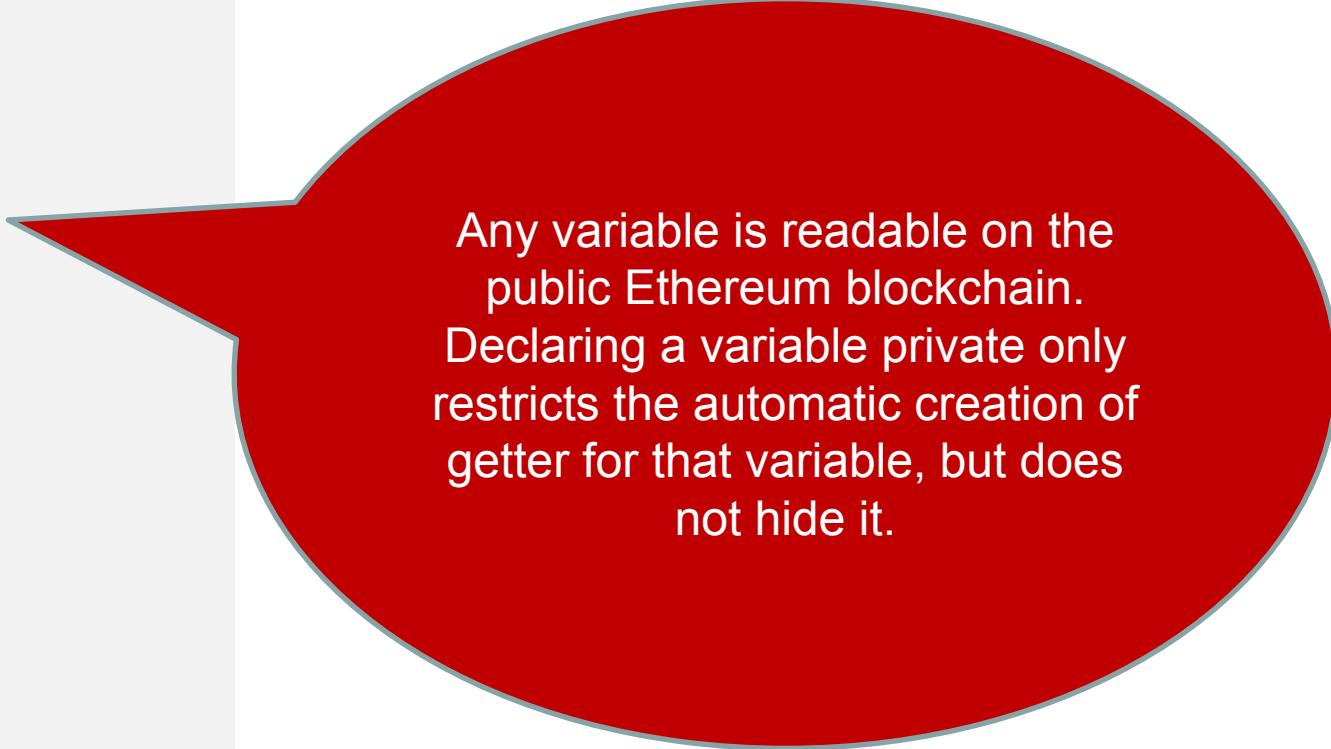
function withdraw(uint amount) {
    if (msg.sender == owner) {
        owner.send(amount);
    }
}
```

Only owner can send ether

An attacker used a similar bug to *steal \$32M*

# Smart Contract Bug Exercise 1

```
contract Example {  
  
    address public owner;  
    string private mySecret;  
  
    constructor {  
        owner = msg.sender;  
    }  
  
    function setSecret(string _secret) public {  
        require(msg.sender == owner);  
        mySecret = _secret;  
    }  
  
    function getSecret() public returns (string) {  
        require(msg.sender == owner);  
        return mySecret;  
    }  
}
```



Any variable is readable on the public Ethereum blockchain. Declaring a variable `private` only restricts the automatic creation of getter for that variable, but does not hide it.

Hint: who would be able to read `mySecret`?

# Smart Contract Bug Exercise 2

```
contract Vulnerable {  
  
    mapping(address => bool) authorized;  
    mapping(address => uint) balances;  
  
    function refund(uint amount) public {  
        require(authorized[msg.sender]);  
        require(amount <= balances[msg.sender]);  
  
        msg.sender.call.value(amount)("");  
        balances[msg.sender] -= amount;  
    }  
}
```

The code is vulnerable to a **reentrancy attack**.

The balance of the *msg.sender* is only updated after a transfer is made. If the *msg.sender* is a contract and has a fallback function that calls into the contract again, the *msg.sender* can deplete the contract of the funds.

Hint: who can be *msg.sender*?

# Smart Contract Bug Exercise 2

```
contract Vulnerable {  
    ... // vulnerable as the previous example  
}
```

```
contract Exploit {
```

```
    Vulnerable v;  
  
    function register(address contract) public {  
        v = Vulnerable(contract);  
    }
```

```
    function exploit() public {  
        // your code here  
    }
```

```
    // your code here  
}
```

Hint: check the previous example

# Smart Contract Bug Exercise 2 - Solution

```
contract Vulnerable {  
    ... // vulnerable as the previous example  
}
```

```
contract Exploit {  
  
    Vulnerable v;  
  
    function register(address contract) public {  
        v = Vulnerable(contract);  
    }  
  
    function exploit() public {  
        v.refund(1);  
    }  
  
    function () public {  
        v.refund(1);  
    }  
}
```

# More smart contract security bugs

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Unexpected ether  
flows



Insecure coding, such as unprivileged writes (e.g., *Multisig Parity bug*)



Use of unsafe inputs (e.g., reflection, hashing, ...)



Reentrant method calls (e.g., *DAO bug*)

# More smart contract security bugs

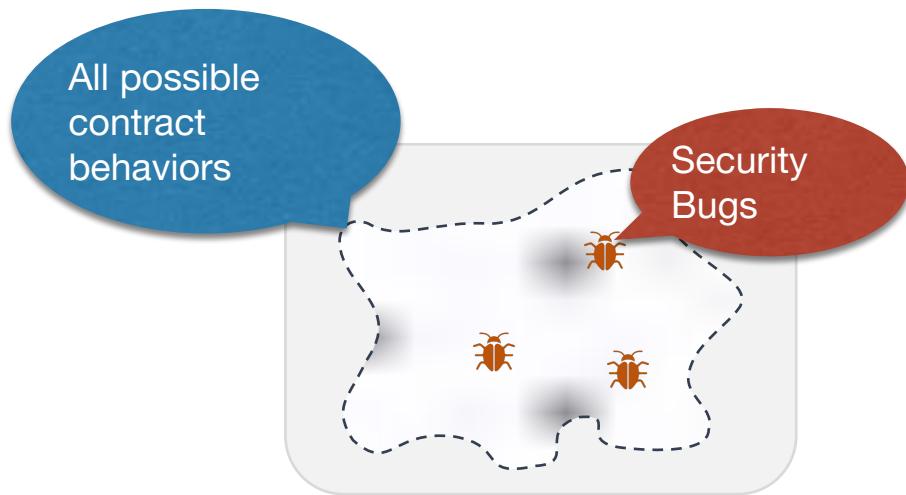
The screenshot shows a web browser window displaying the "Known Attacks" page from the Ethereum Smart Contract Best Practices GitHub repository. The URL is [https://consensys.github.io/smart-contract-best-practices/known\\_attacks/](https://consensys.github.io/smart-contract-best-practices/known_attacks/). The page has a dark blue header with the title "Ethereum Smart Contract Best Practices" and a search bar. On the left, there's a sidebar with navigation links like Home, General Philosophy, Secure Development Recommendations, Known Attacks (which is the current page), Software Engineering Techniques, Token specific recommendations, Documentation and Procedures, Security Tools, Bug Bounty Programs, and About. The main content area is titled "Known Attacks" and contains a section on Reentrancy. It explains that reentrancy occurs when a function can be called again before its first invocation has completed, leading to data race conditions. Below this is a code snippet illustrating an insecure withdrawal function:

```
// INSECURE
mapping (address => uint) private userBalances;

function withdrawBalance() public {
    uint amountToWithdraw = userBalances[msg.sender];
    (bool success, ) = msg.sender.call.value(amountToWithdraw)("");
    require(success);
    userBalances[msg.sender] = 0;
}
```

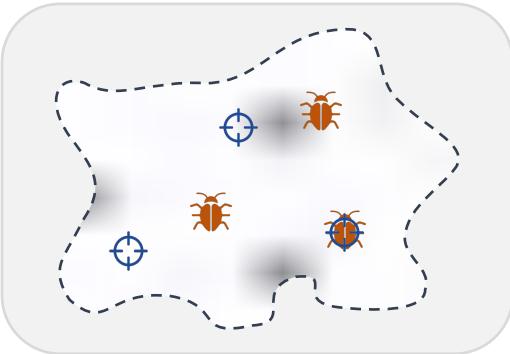
A note at the bottom states: "Since the user's balance is not set to 0 until the very end of the function, the second (and later) invocations will succeed". To the right of the main content is a "Table of contents" sidebar listing various attack categories: Reentrancy, Reentrancy on a Single Function, Cross-function Reentrancy, Pitfalls in Reentrancy Solutions, Front-Running, Taxonomy, Displacement, Insertion, Suppression, Mitigations, Timestamp Dependence, Integer Overflow and Underflow, DoS with (Unexpected) revert, DoS with Block Gas Limit, Gas Limit DoS on a Contract via Unbounded Operations, Gas Limit DoS on the Network via Block Stuffing, Insufficient gas griefing, Forcibly Sending Ether to a Contract, Deprecated/historical attacks, Call Depth Attack (deprecated), Constantinople Reentrancy Attack, and Other Vulnerabilities.

# Automated security analysis

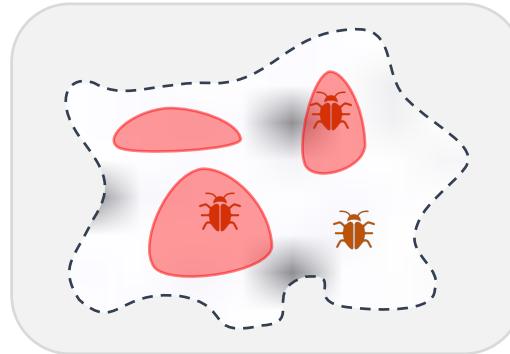


Problem: Cannot enumerate all possible contract behaviors...

# Automated security analysis – Existing solutions

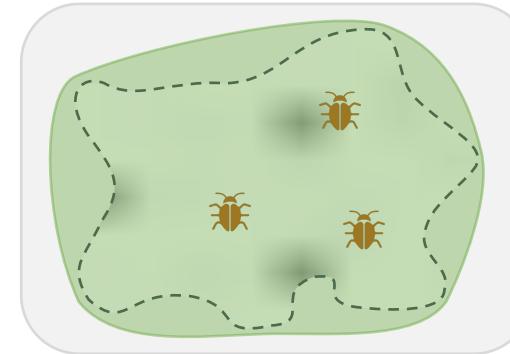


Testing



Dynamic analysis  
Symbolic execution

Easy to implement, but  
very limited guarantees



Static analysis  
Formal verification

Strong guarantees, but many  
false positives



# DeFi Flash Loan „Attacks“

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# Flash Loan Attacks

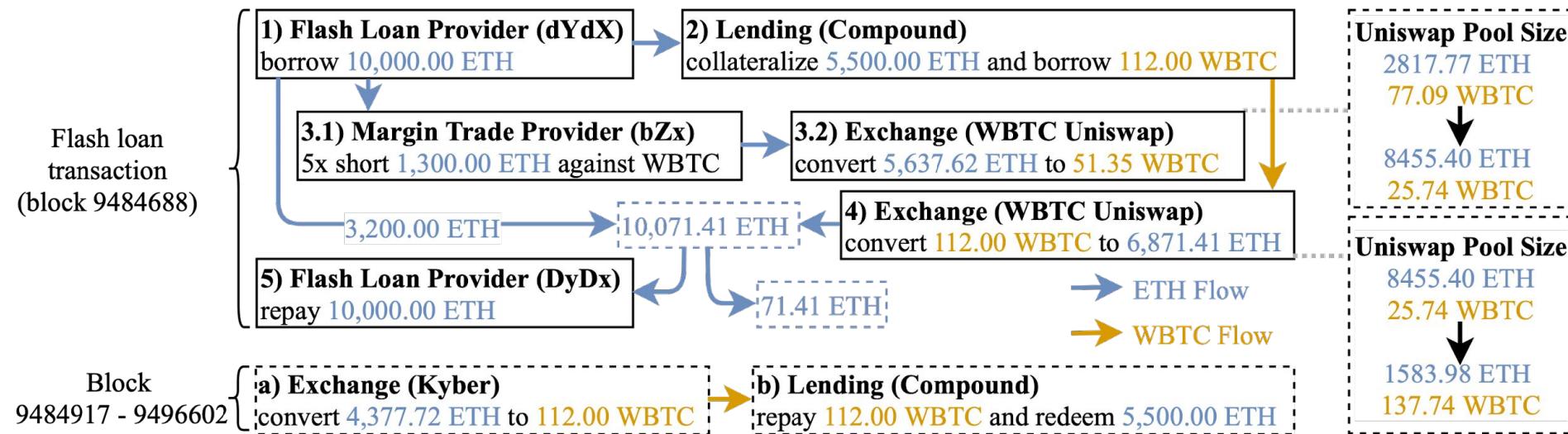
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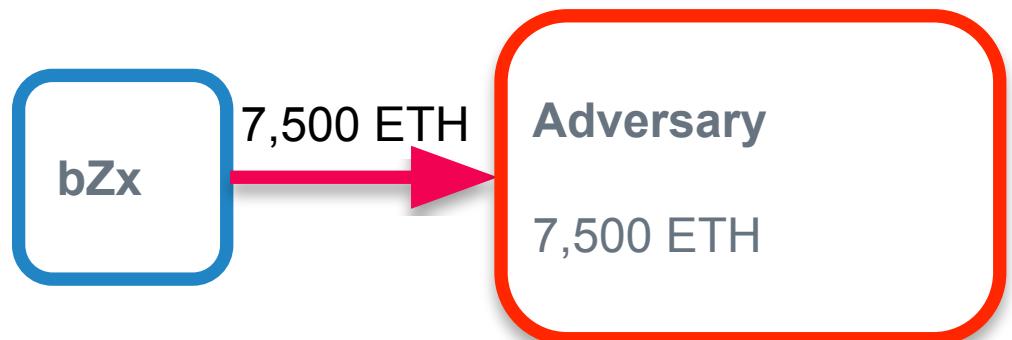
# bZx - Pump and Arbitrage Attack – February 2020



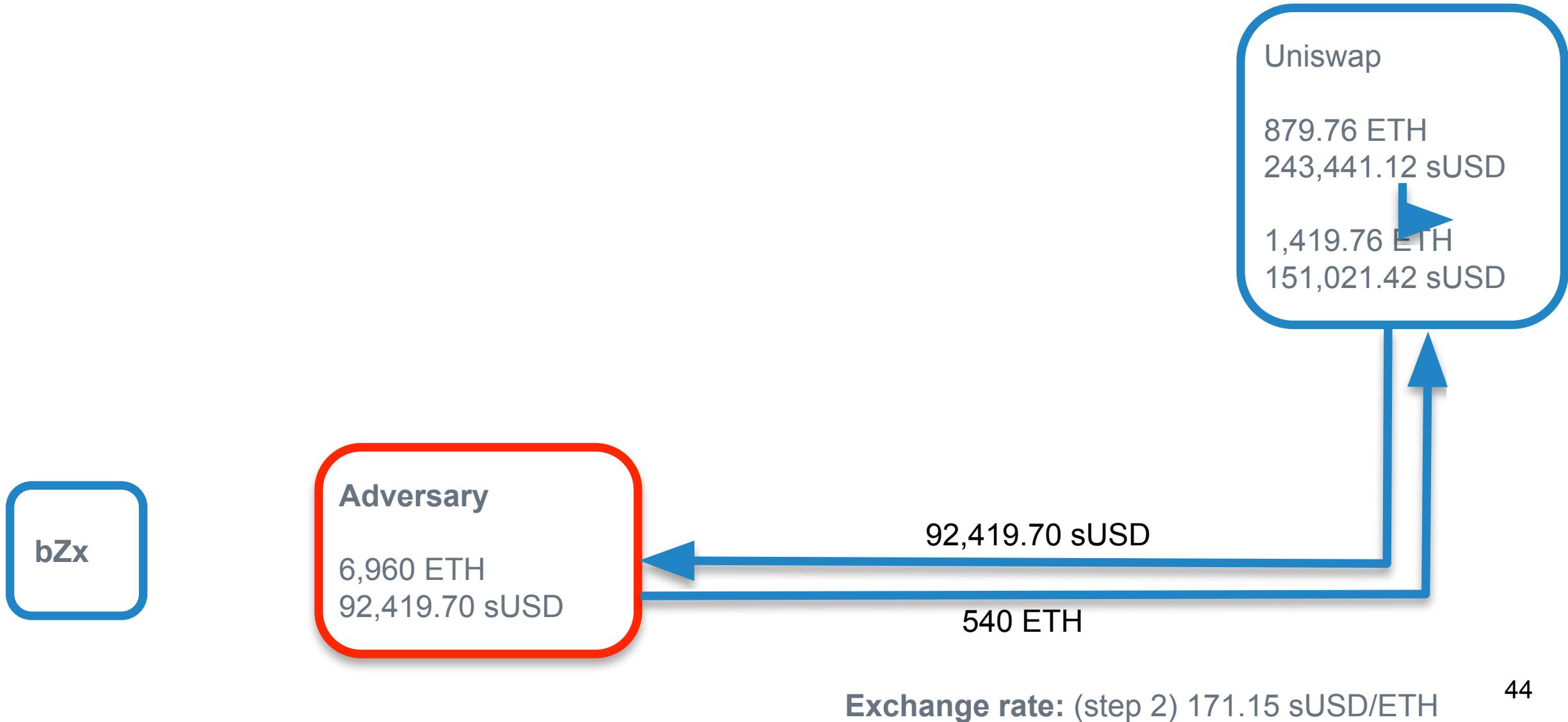
**Input: 130 USD gas**  
**Output: 350,000 USD**  
**Optimal: 830,000 USD**

# bZx – Oracle manipulation – February 2020

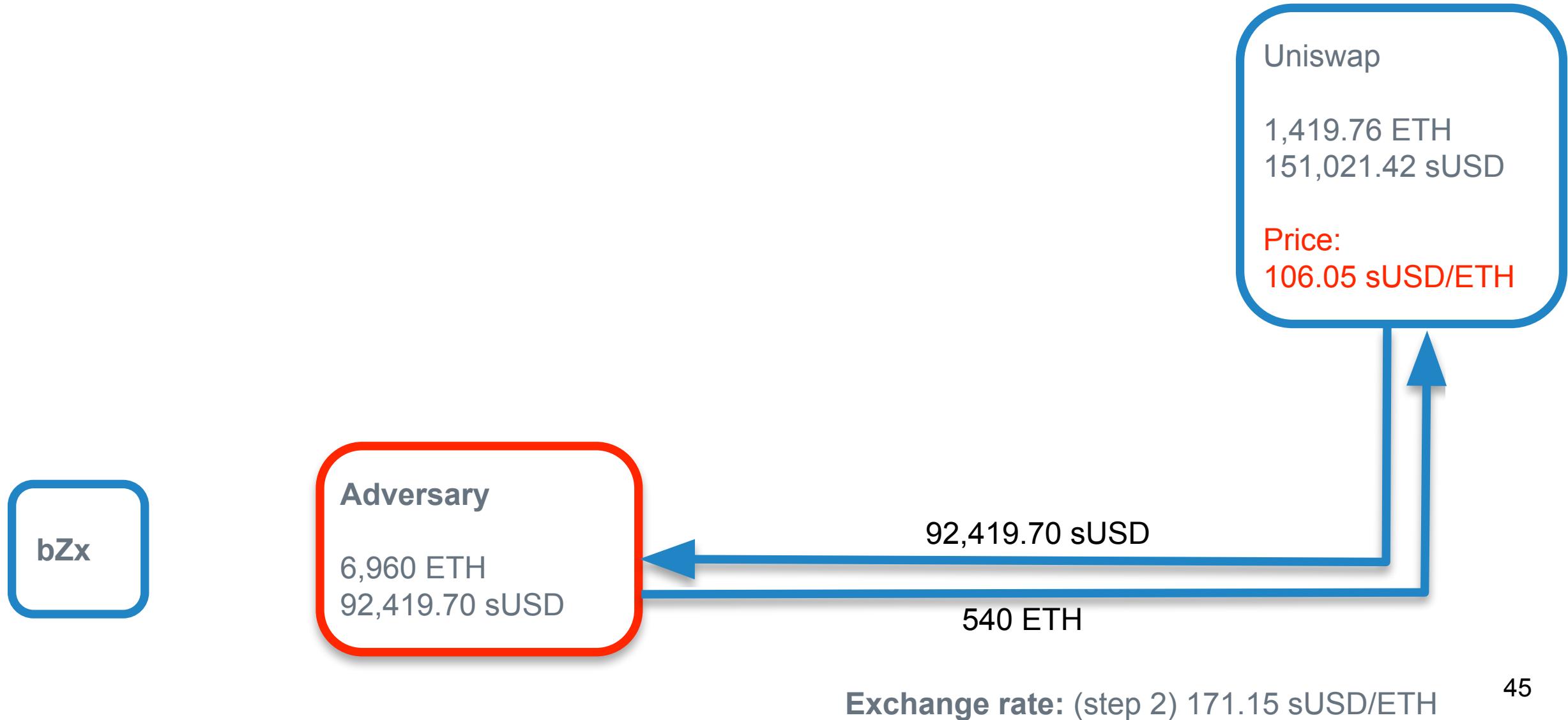
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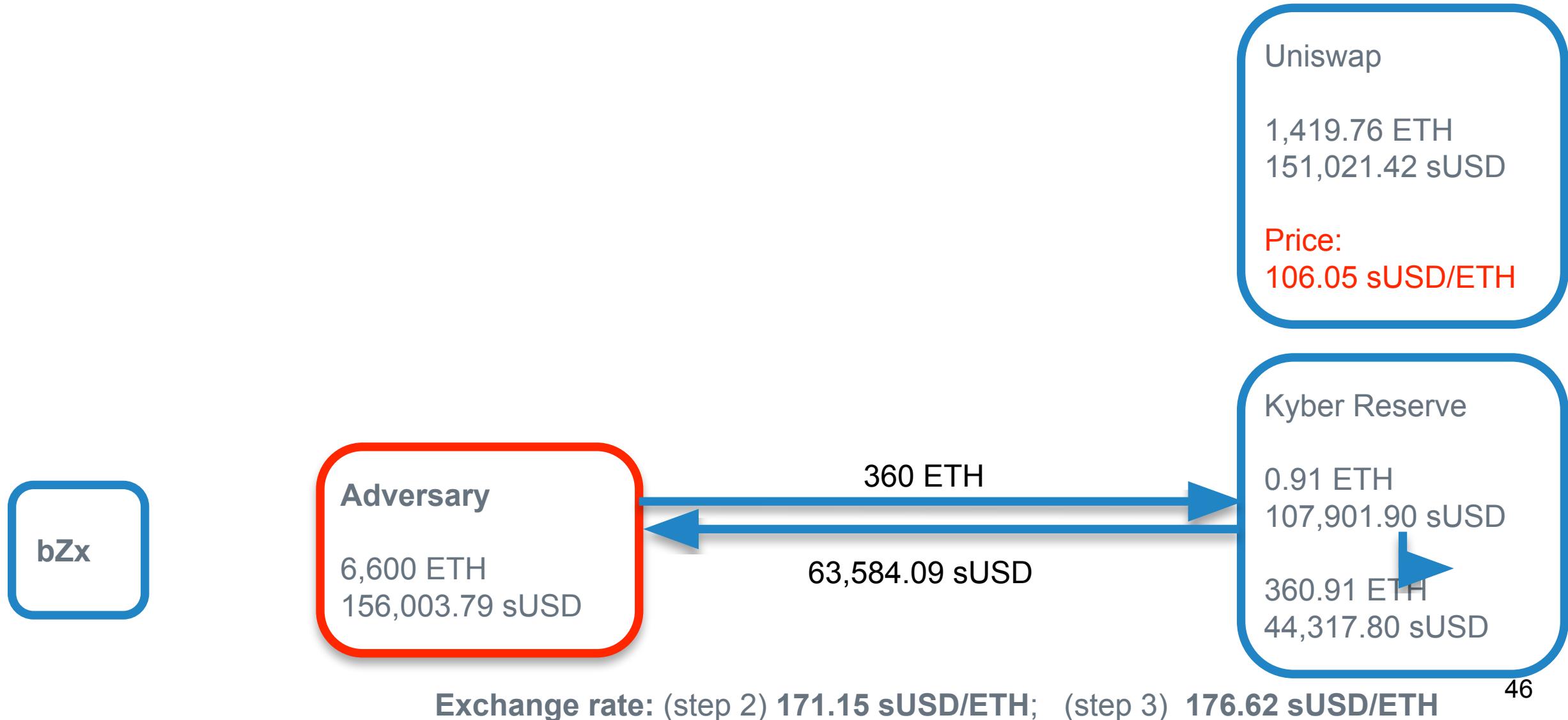
# bZx – Oracle manipulation – February 2020



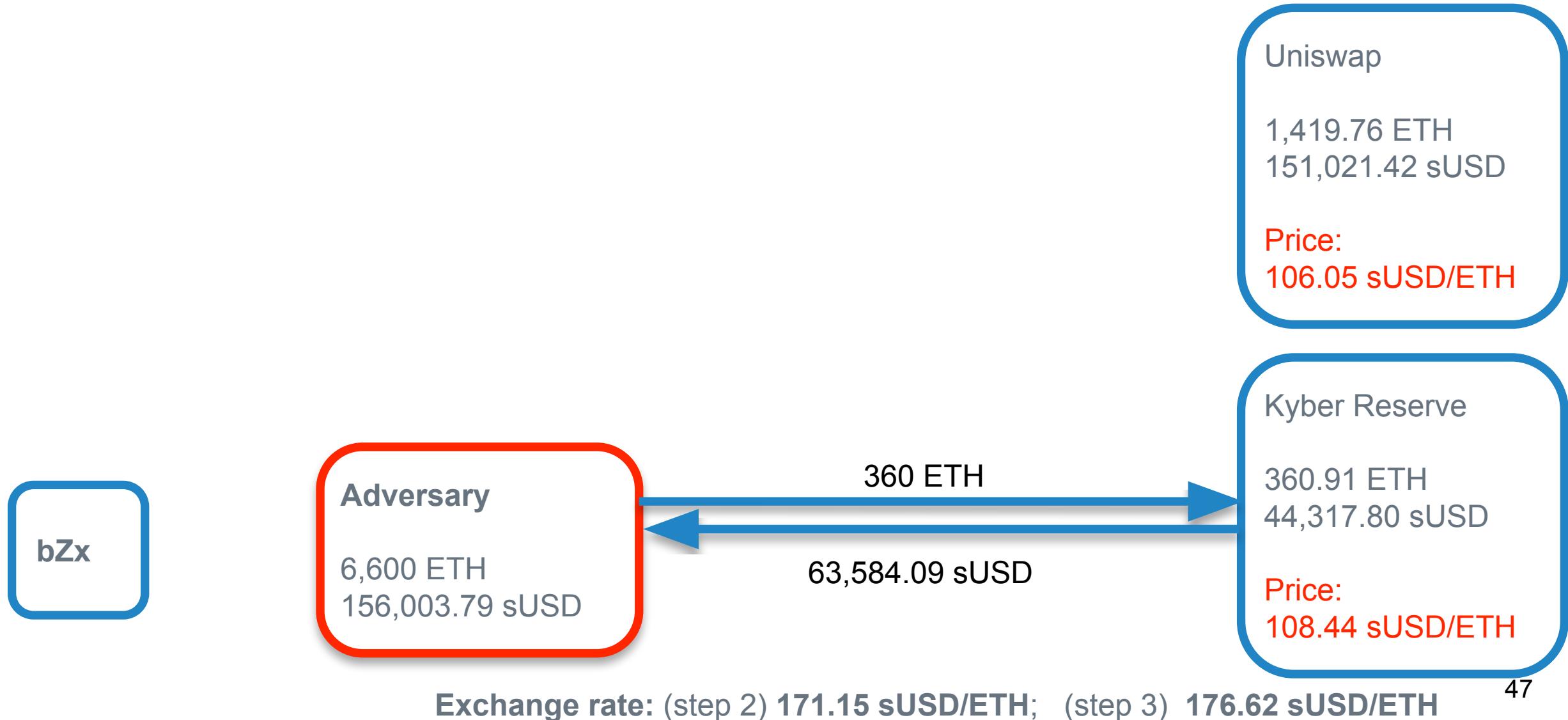
# bZx – Oracle manipulation – February 2020



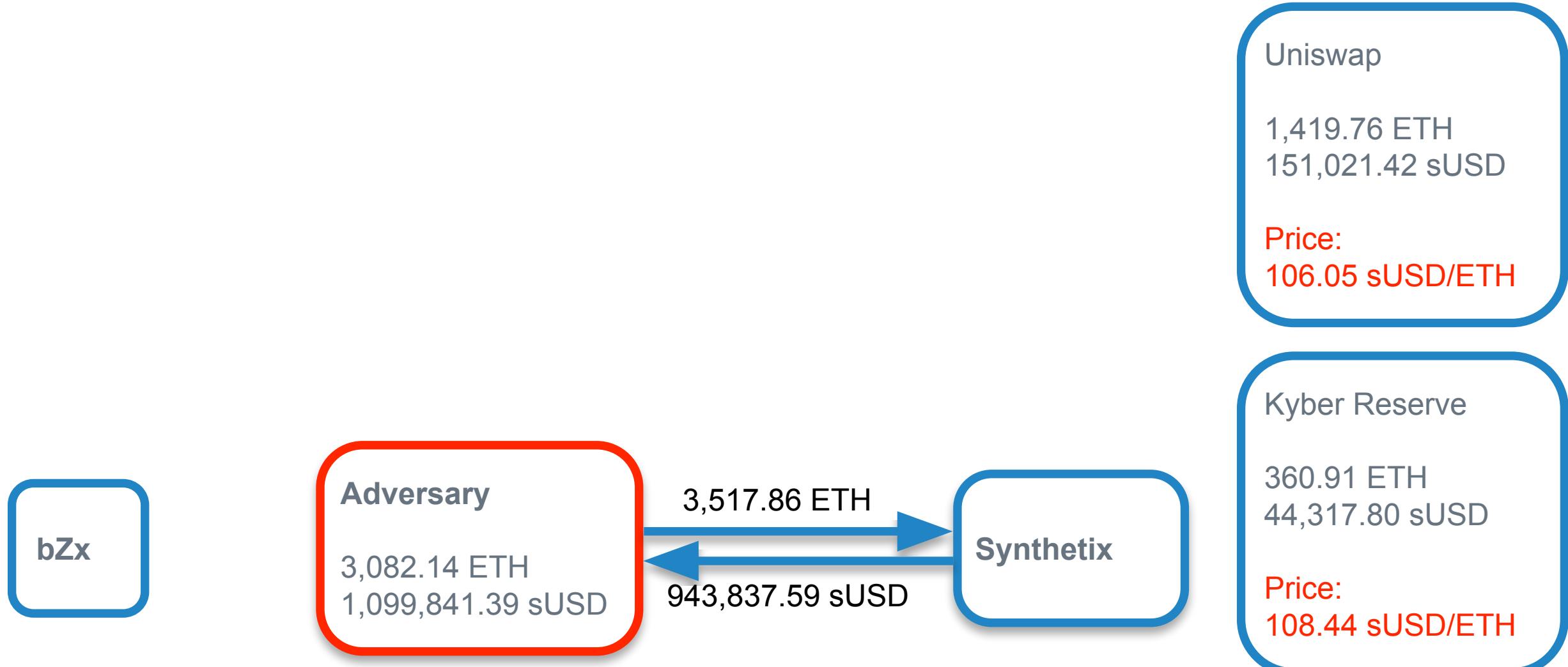
# bZx – Oracle manipulation – February 2020



# bZx – Oracle manipulation – February 2020

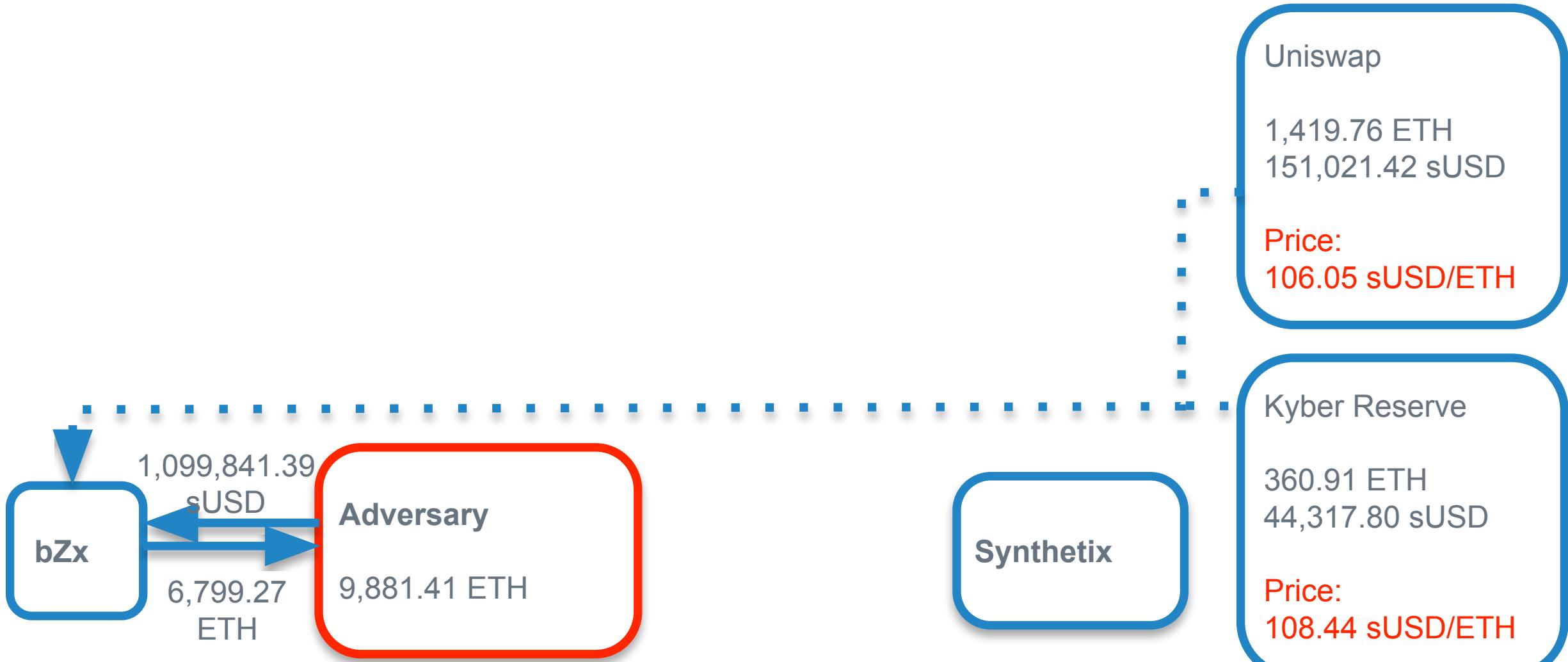


# bZx – Oracle manipulation – February 2020



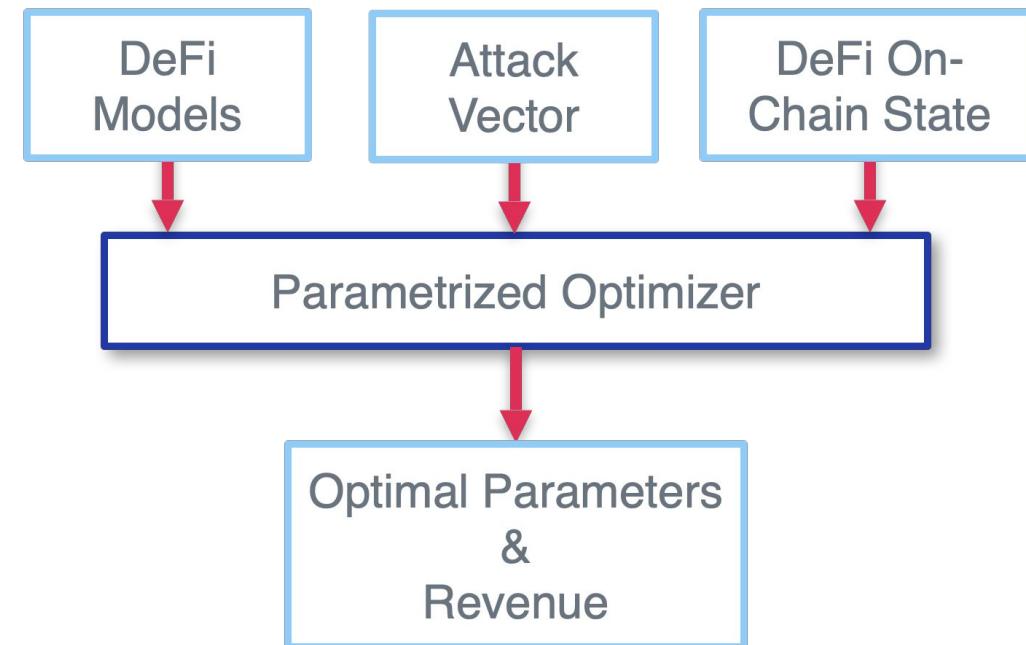
Exchange rate: (step 2) 171.15 sUSD/ETH; (step 3) 176.62 sUSD/ETH; (step 4) 268.30 sUSD/ETH

# bZx – Oracle manipulation – February 2020



# Constrained Optimization Framework

- Formulate DeFi actions in models
  - Constant product AMM:  $\Delta y = y - \frac{xy}{x+\Delta x}$
- Construct a constrained optimization problem based on the attack vector
  - Objective function: outcome profit
- Fetch the on-chain state that the



# Optimizing the bZx attack 2

---

- Borrow  $X$  ETH (bZx flash loan)
  - Convert  $p_1$  ETH to  $f_1(p_1)$  sUSD (Uniswap)
  - Convert  $p_2$  ETH to  $f_2(p_2)$  sUSD (Kyber)
  - Deposit  $p_3$  ETH for  $f_3(p_3)$  sUSD (Synthetix)
  - Collateralize  $z$  sUSD to borrow  $g(z)$  ETH
  - $z=f_1(p_1)+f_2(p_2)+f_3(p_3)$
- Repay  $X$  ETH (bZx flash loan)
- Objective:  $o=g(f_1(p_1)+f_2(p_2)+f_3(p_3))-X$ 
  - s.t.  $p_1+p_2+p_3 < X$

# Optimizing the bZx attack 2

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- Sequential Least Squares Programming (SLSQP)
  - SciPy
- Ubuntu 18.04.2, 16 CPU cores, 32 GB RAM
- Validation by concrete execution
  - Execution on the real blockchain state



# Sandwich Attacks

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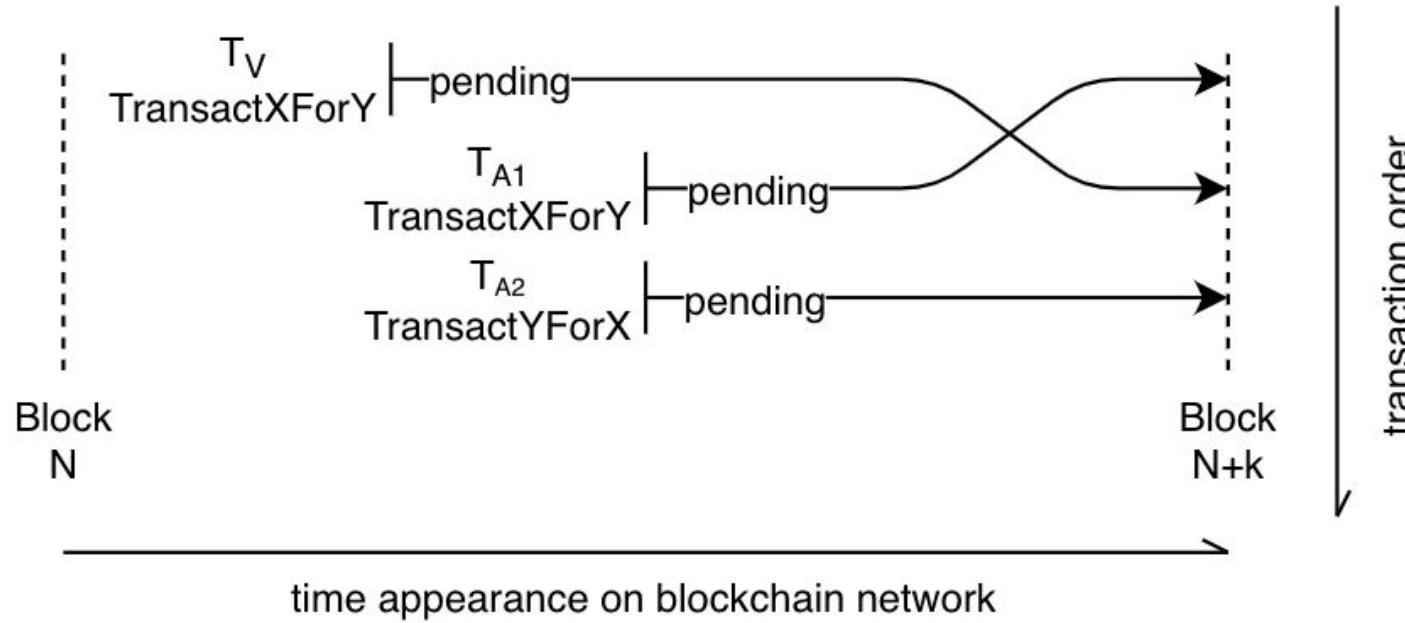
# AMM – Automated Market Maker

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$$x \times y = k$$

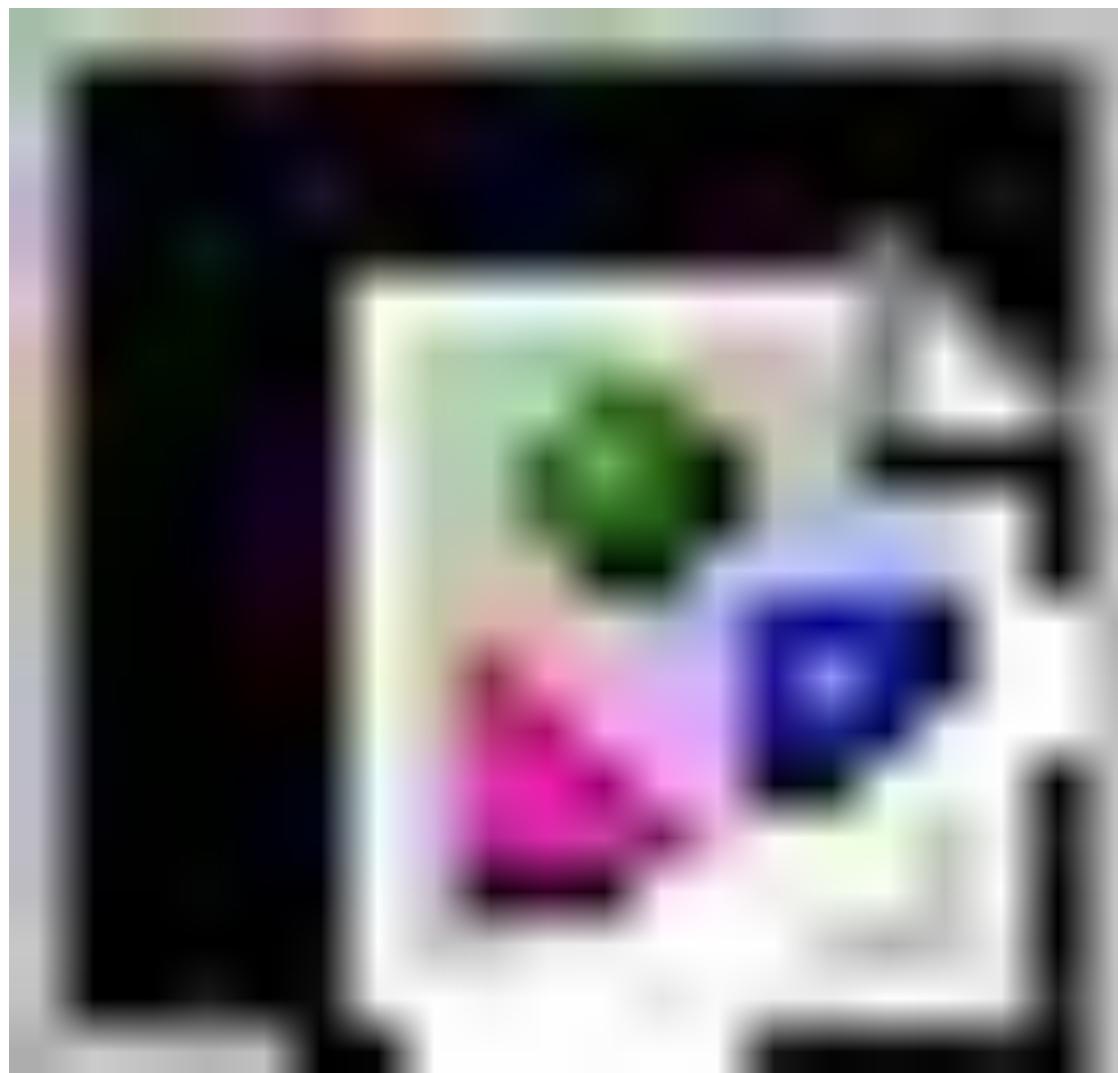
The diagram illustrates the components of the formula  $x \times y = k$ . The variable  $x$  is associated with "Asset X quantity", the variable  $y$  is associated with "Asset Y quantity", and the product  $x \times y$  is equated to the constant  $k$ .

# Sandwich Attack

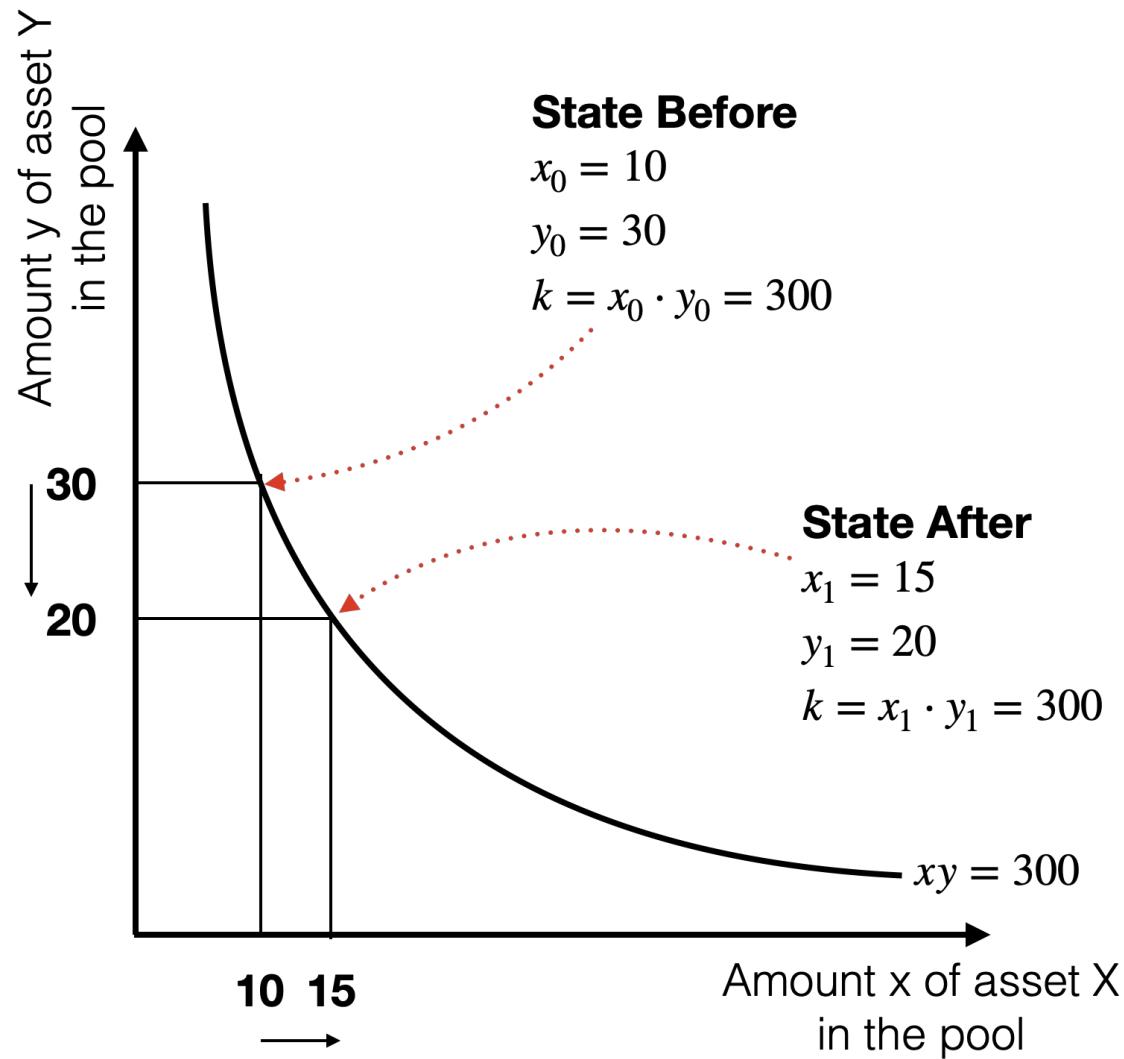


# AMM – Constant product formula

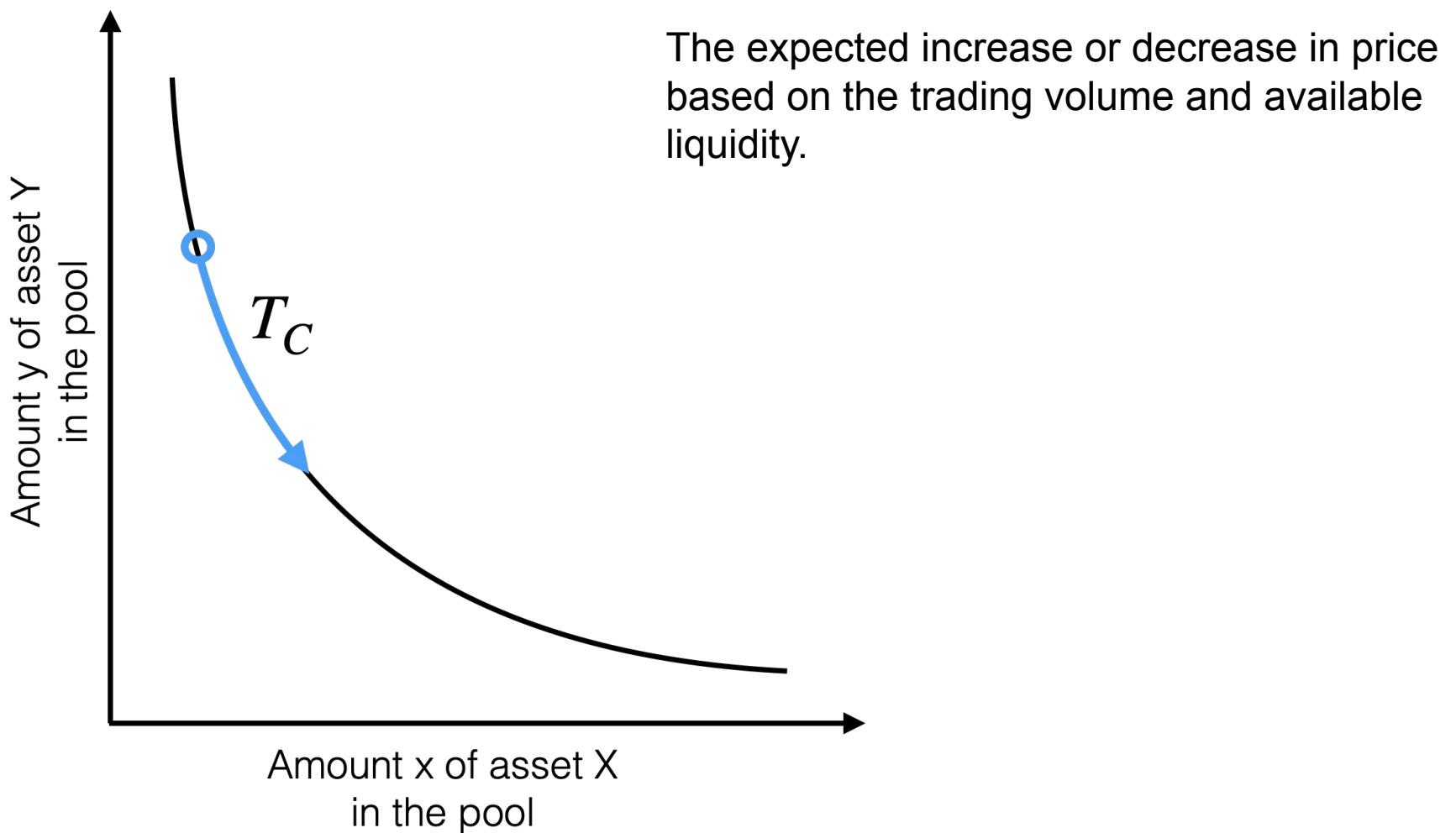
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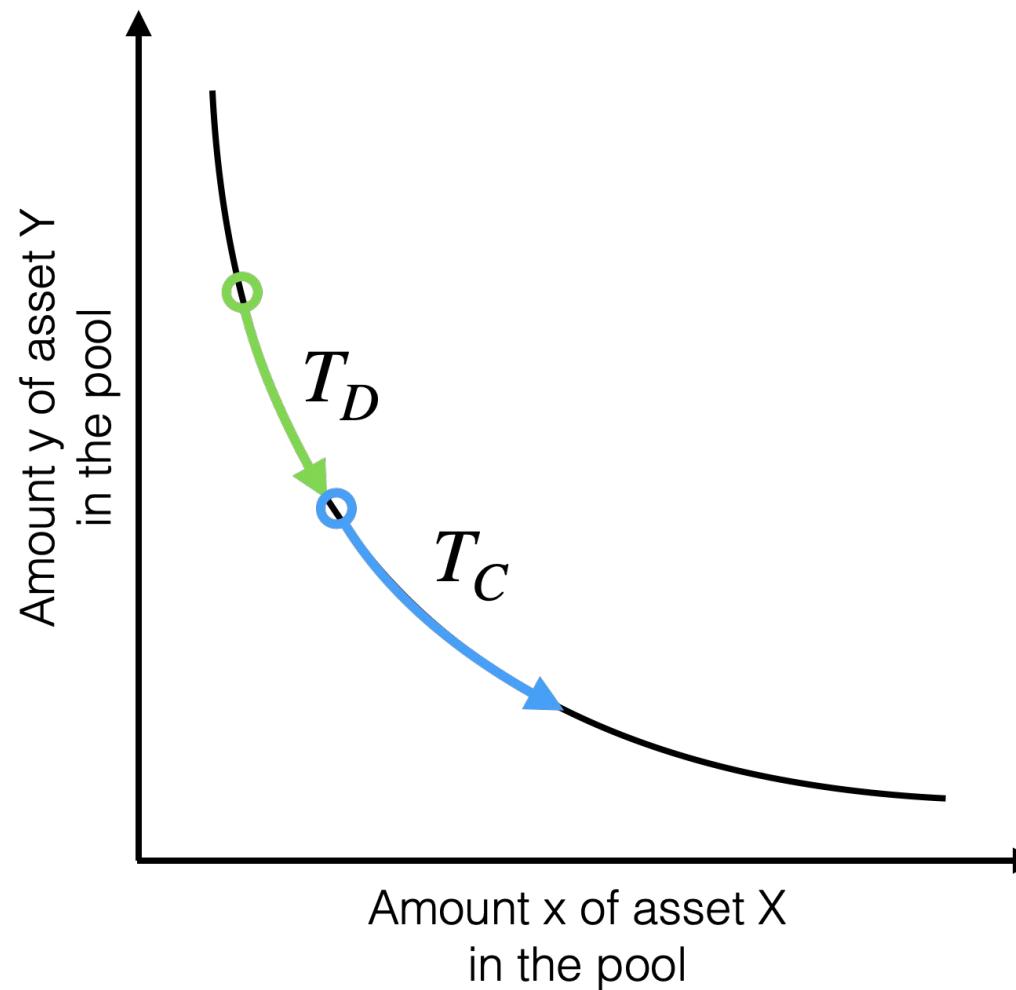
# AMM – Constant product formula



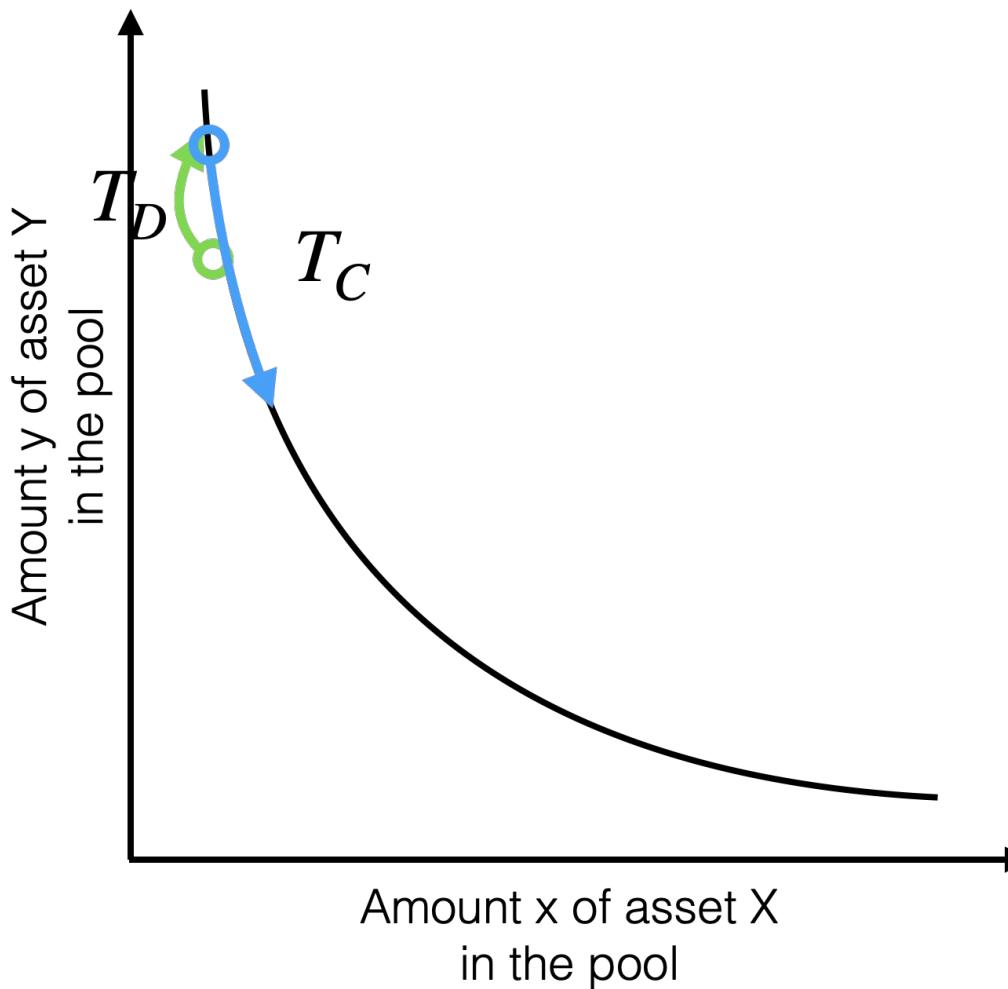
# Expected Slippage



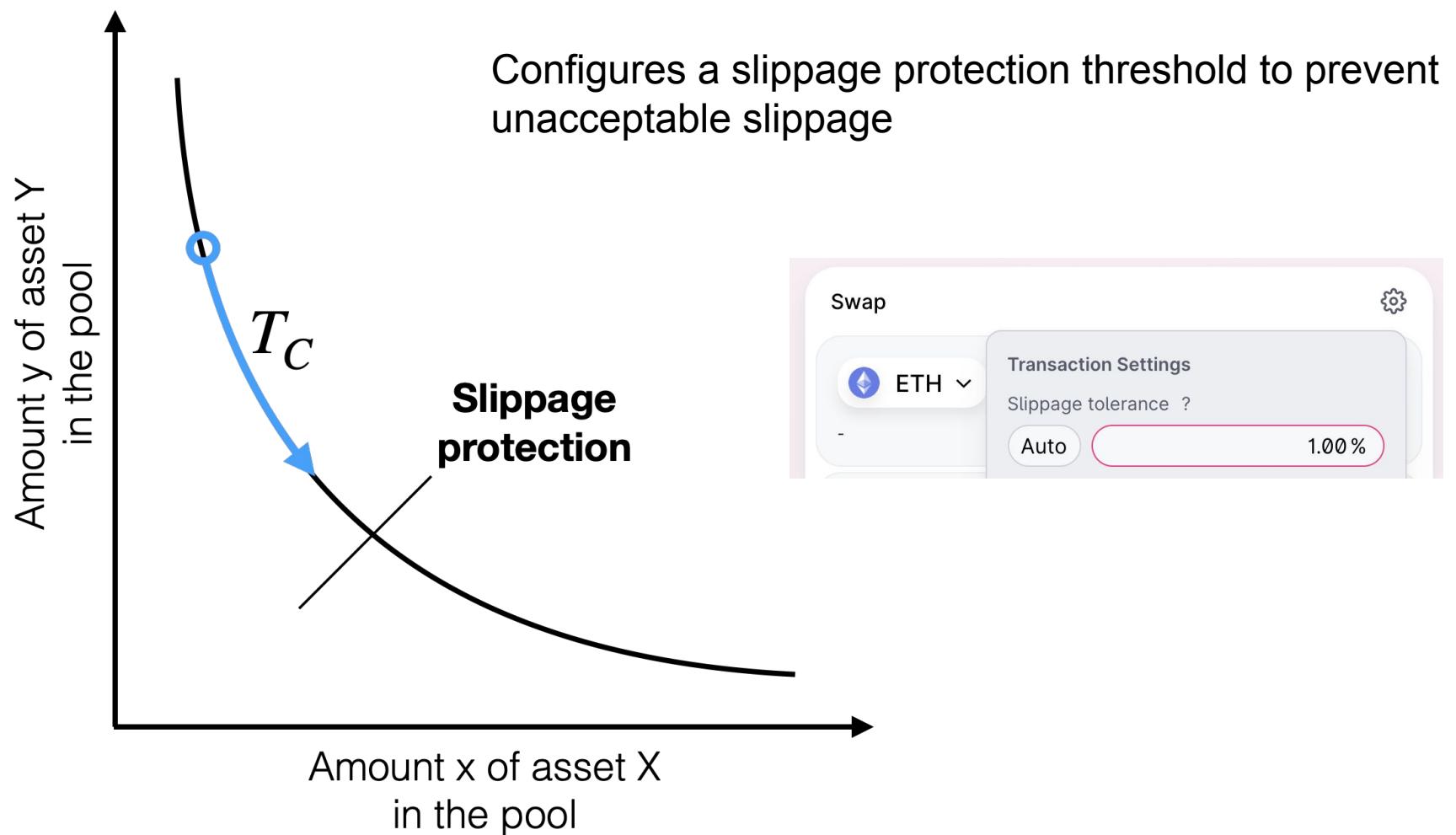
# Unexpected Slippage -> Worse Execution Price



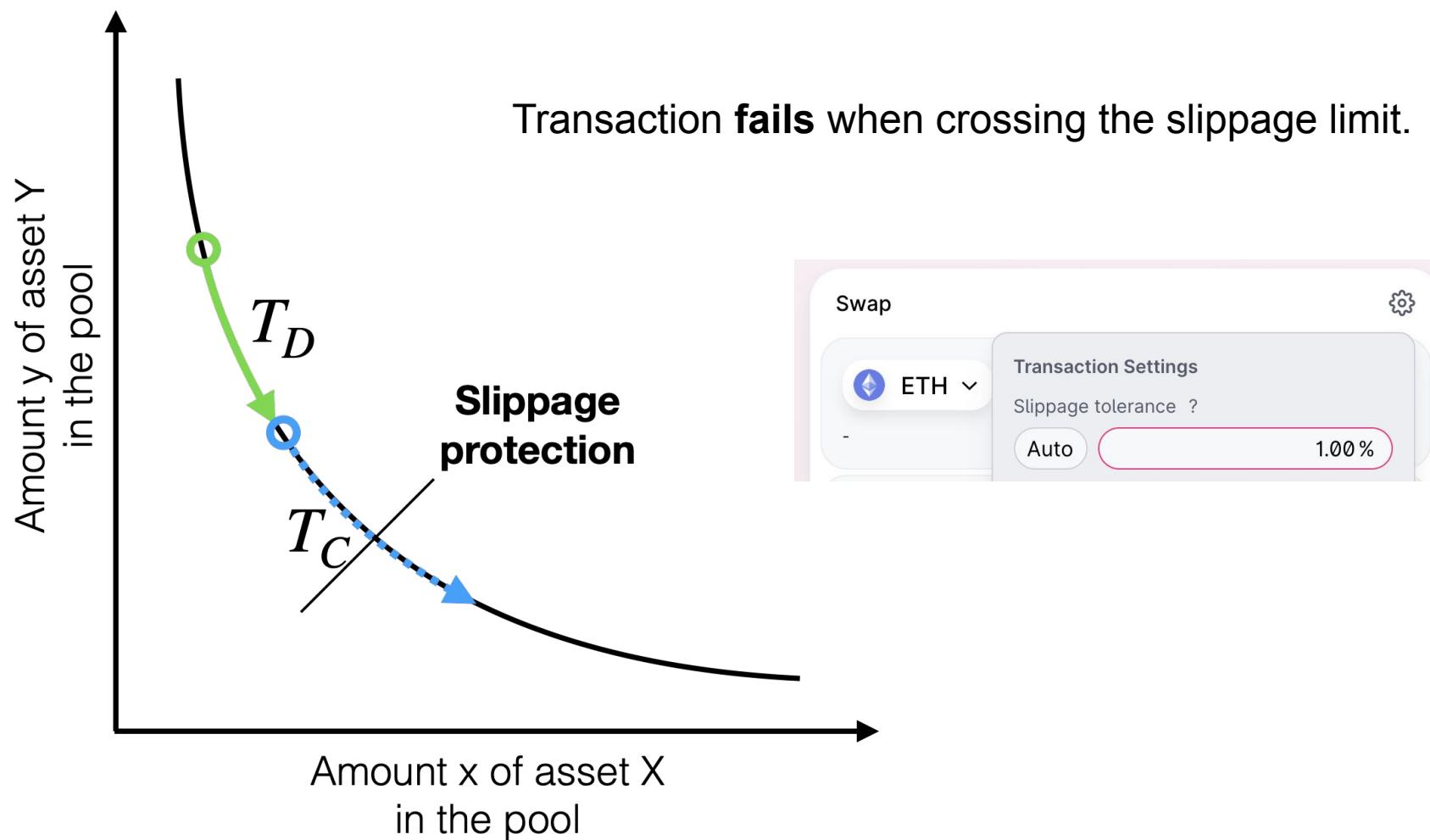
# Unexpected Slippage -> Better Execution Price



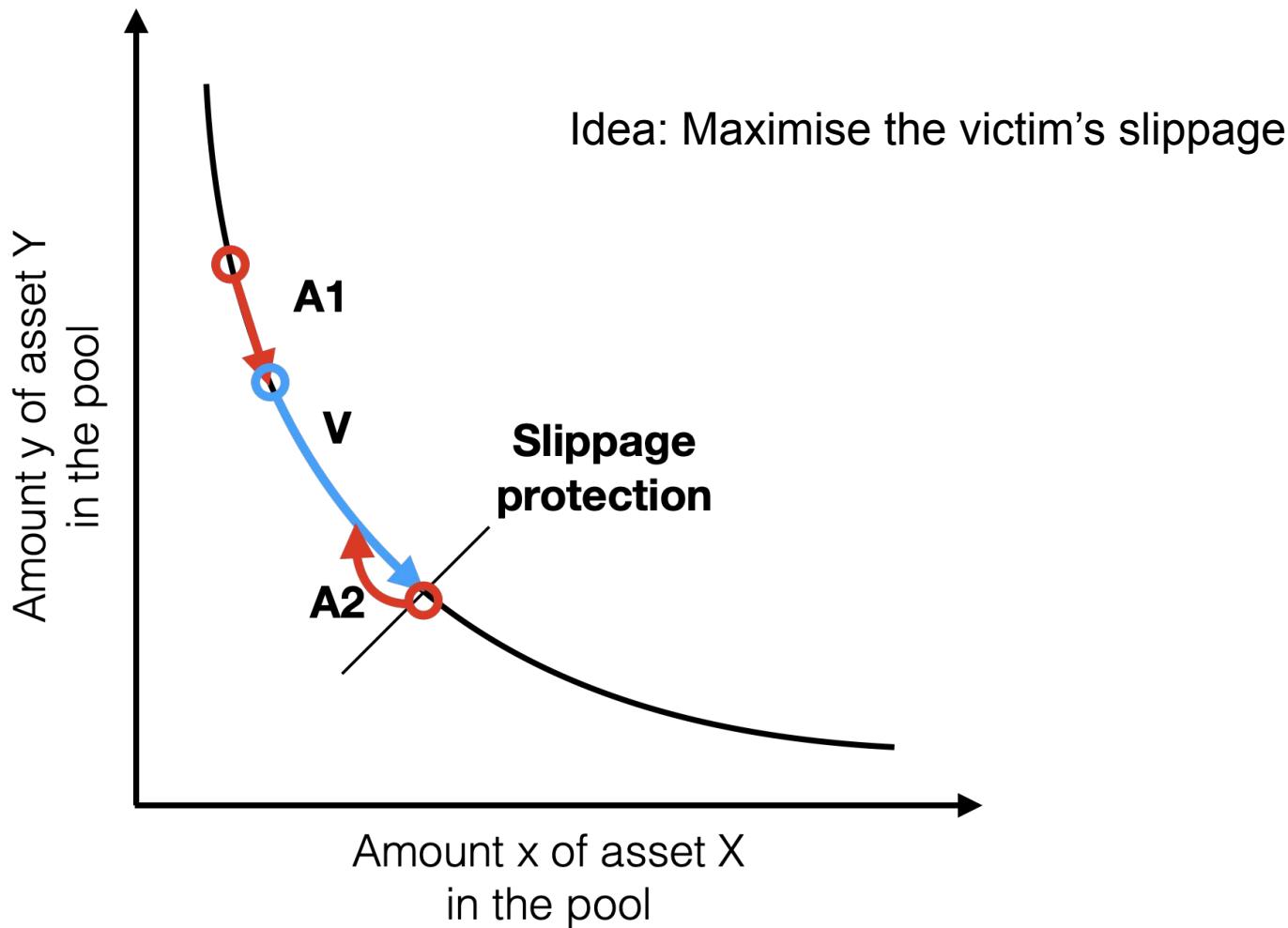
# Slippage Protection



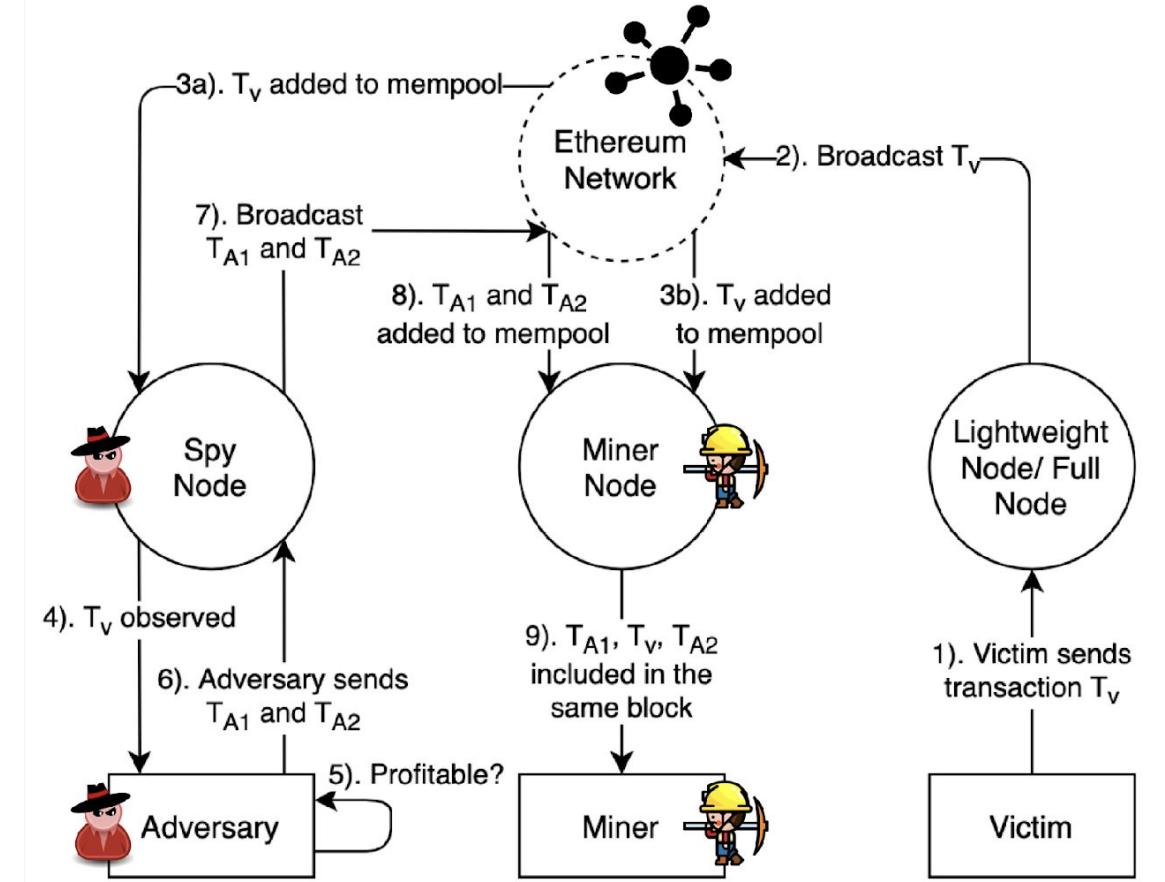
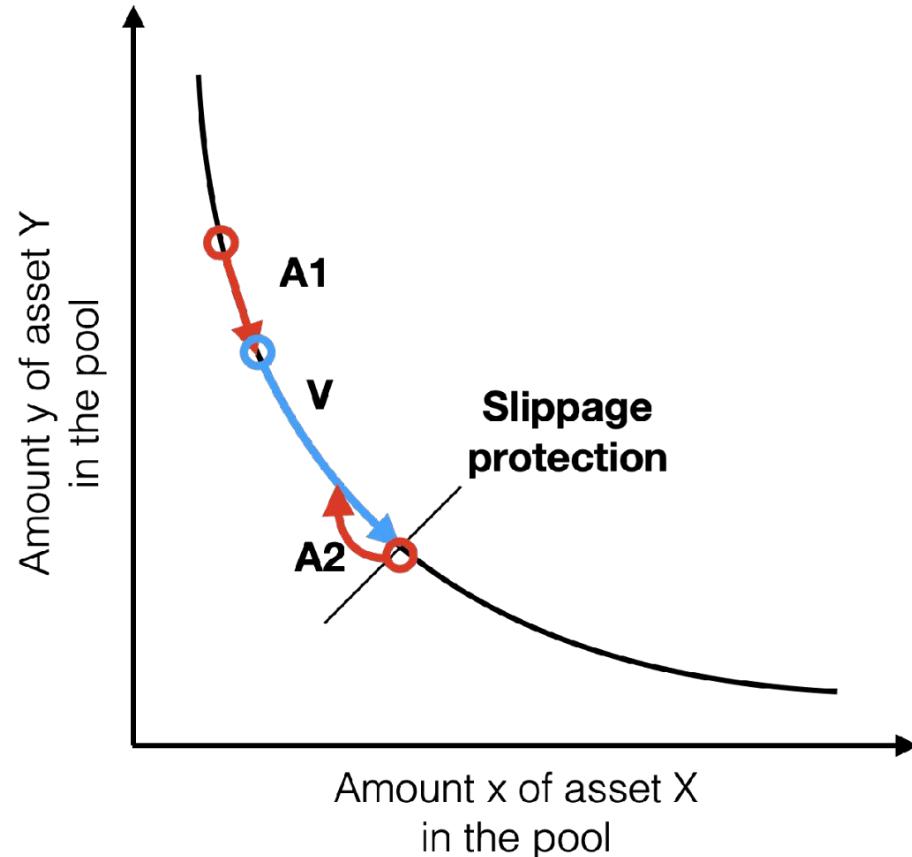
# Slippage Protection



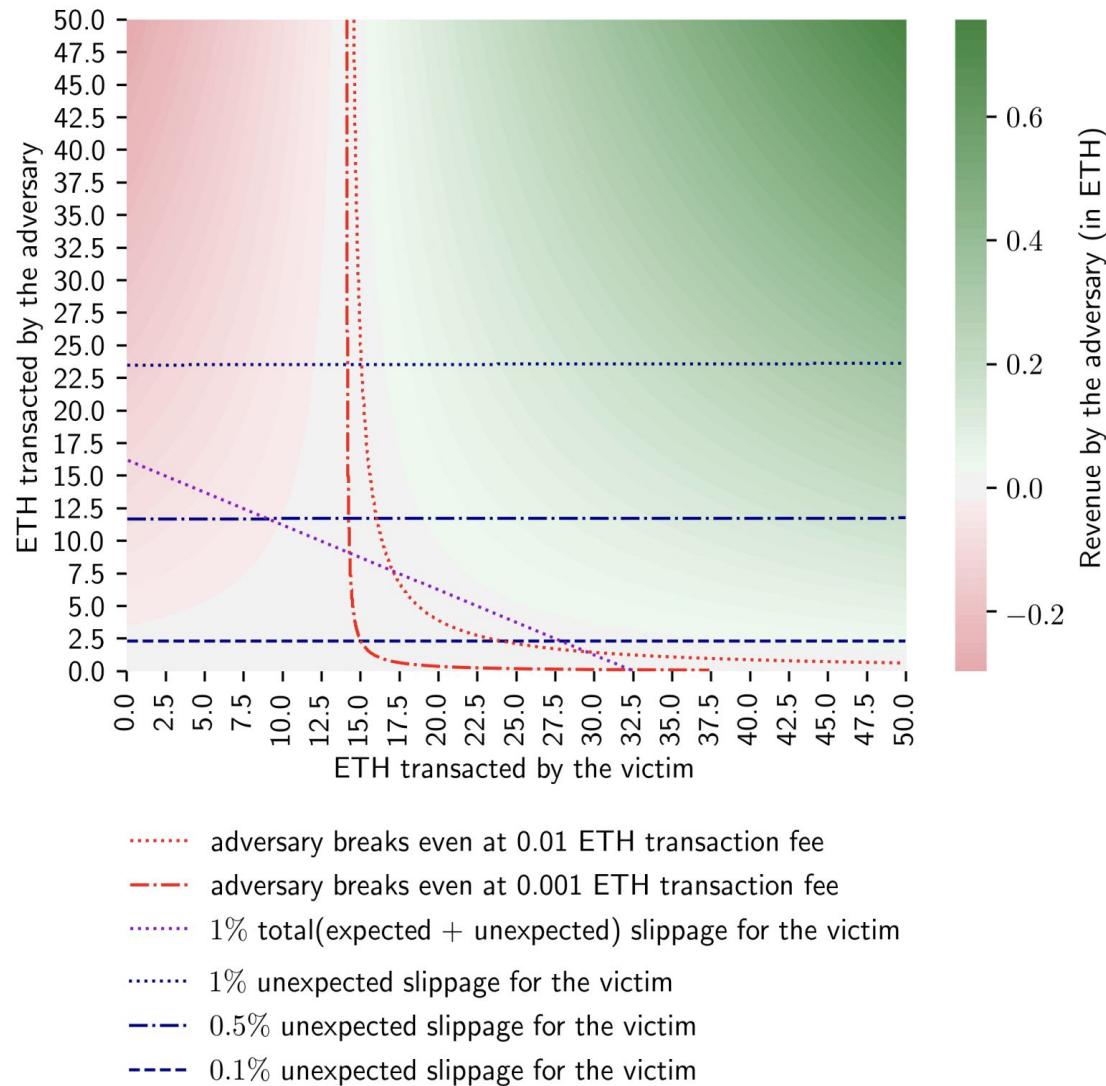
# Sandwich Attack Against Taker



# Network layer + DeFi protocol layer

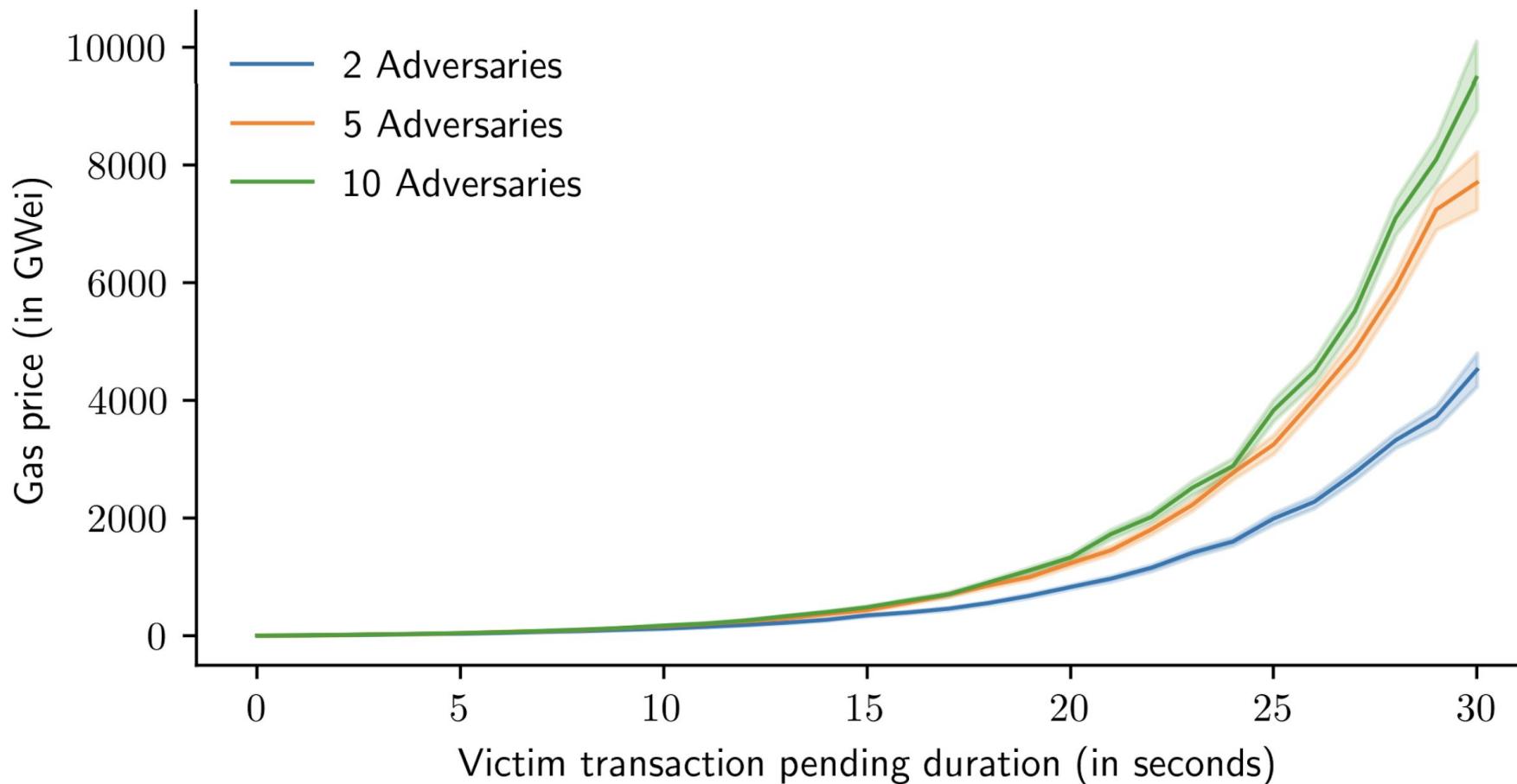


# Sandwich attack profitability



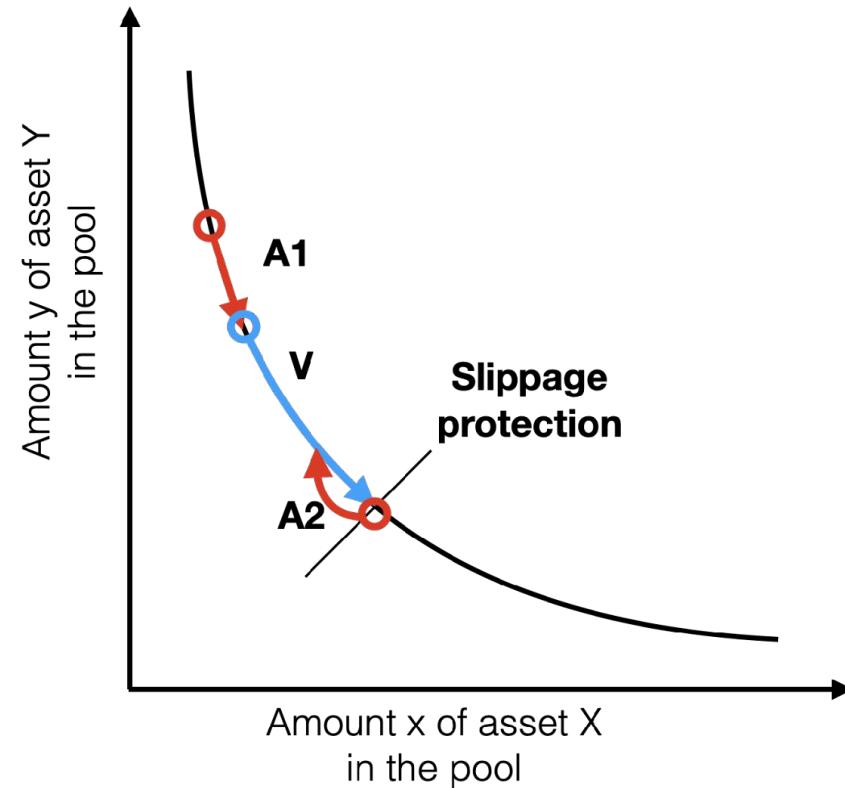
# Multiple Adversaries

Break-even of the attacker becomes harder to attain

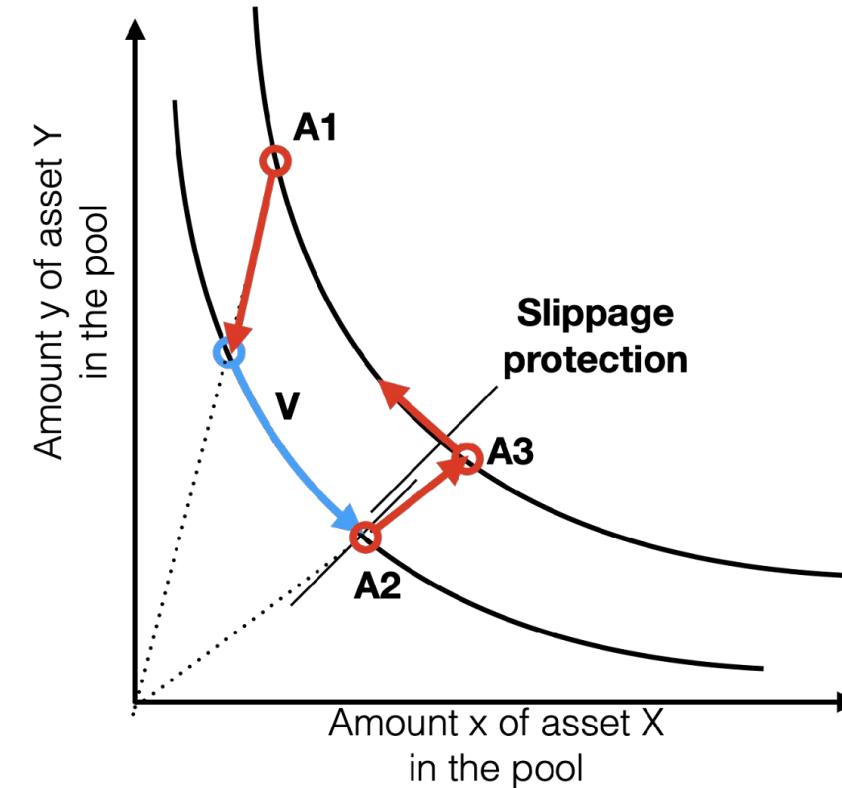


# Advanced Sandwich Attack

Taker attacks Taker



Provider attacks Taker



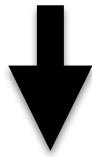


# Blockchain Extractable Value

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# What is Blockchain (or Miner) Extractable Value?

Price of collateral drops below health factor



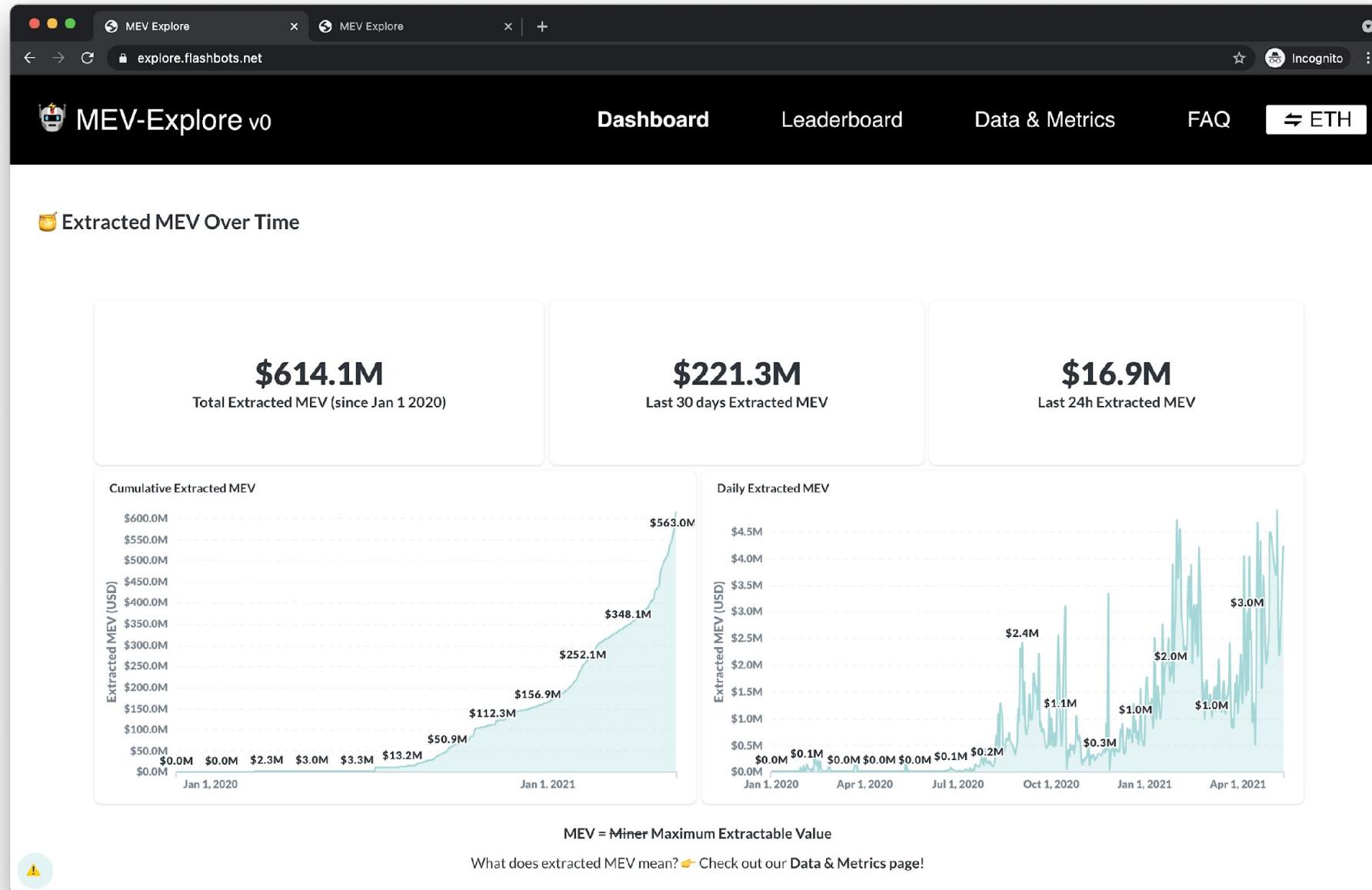
Liquidation! 



Who will liquidate?



# How much MEV?



# How much MEV? – Sandwich Attacks

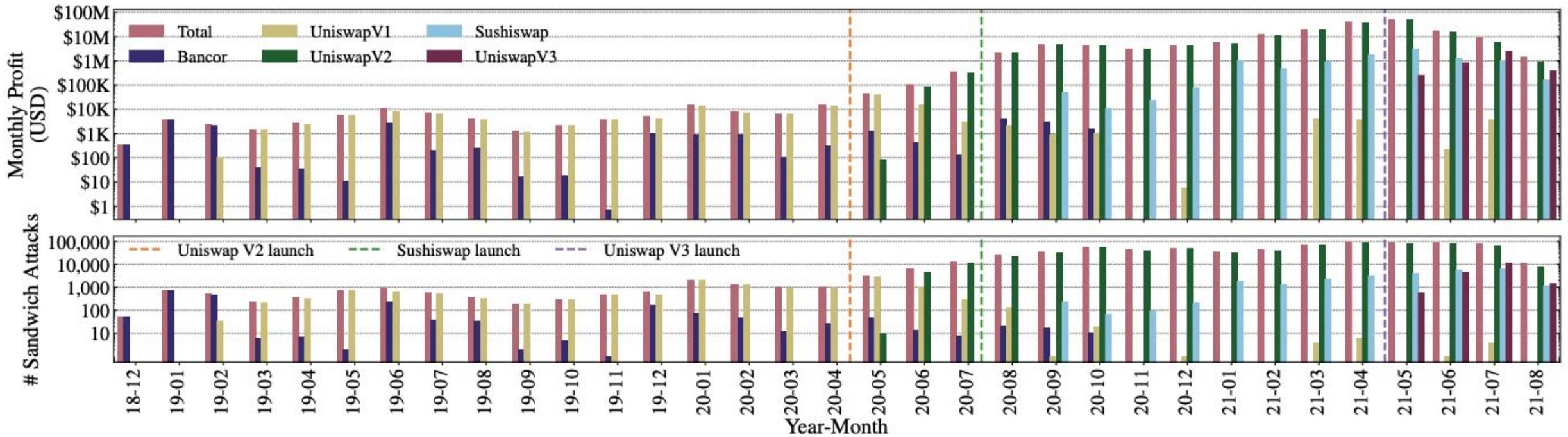
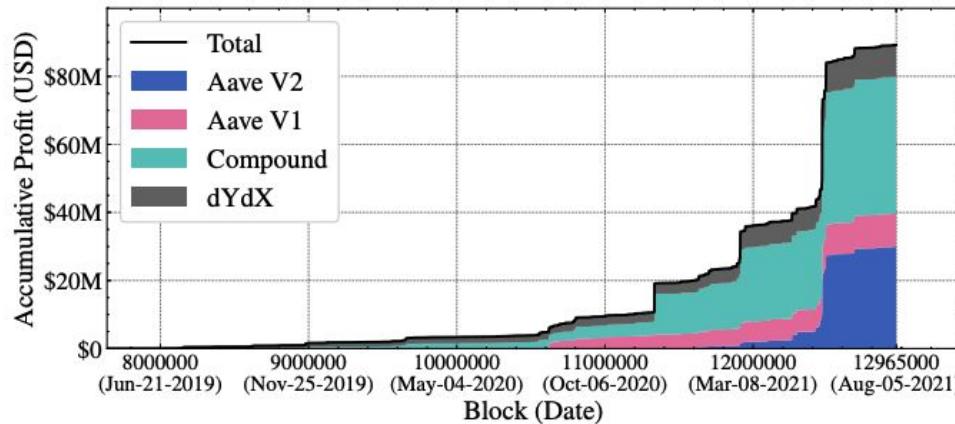
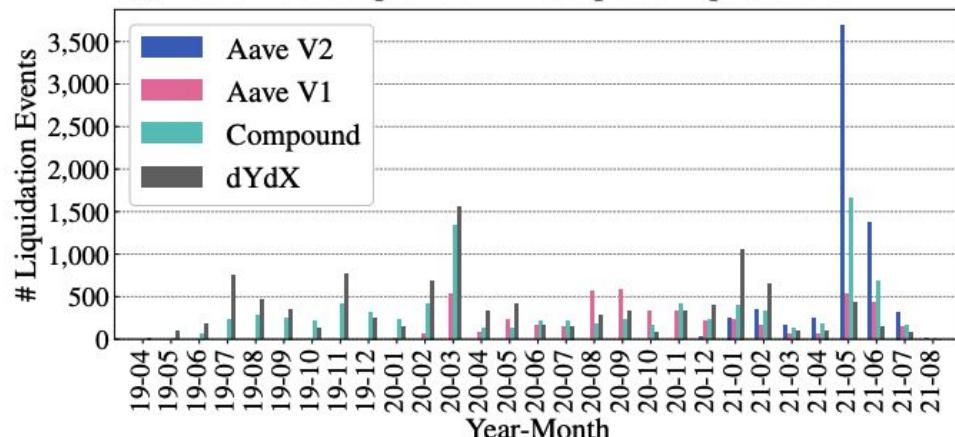


Fig. 3: Sandwich attacks, from block 6803256 (1st of December, 2018) to 12965000 (5th of August, 2021).

# How much MEV? – Liquidations



(a) Accumulative profit of fixed spread liquidations.



(b) The monthly number of fixed spread liquidation events.

Fig. 5: The number of liquidations increase in months where the ETH price collapses, e.g., in March, 2020 and May, 2021.

# How much MEV? – Arbitrage

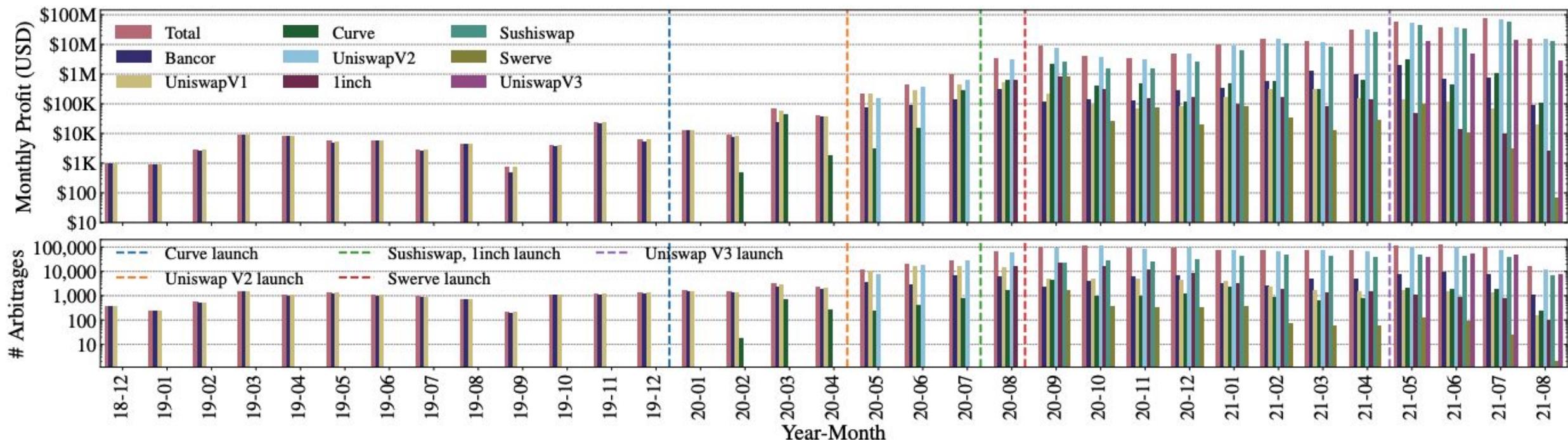


Fig. 7: Monthly arbitrage statistics from block [6803256](#) (1st of December, 2018) to block [12965000](#) (5th of August, 2021).



# Transaction Replay Attacks

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# Generalized Front-Running

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- “Copy Cat” or “Replay”
  - Observe transaction on the network layer
  - Replace certain data, sign, and broadcast copy
- Potential Profit
  - 35M USD over 32 months
  - 188,365 profitable transactions (0.02%)
  - Real-time algorithm ( $0.18s \pm 0.29$ )

# Generalized Front-Running Algorithm & Results

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**Algorithm 1:** Transaction Replay Algorithm.

**Input:** The current highest block  $B_i$ ; the potential victim transaction  $T_V$ ; the adversarial account address  $\mathcal{A}$ .

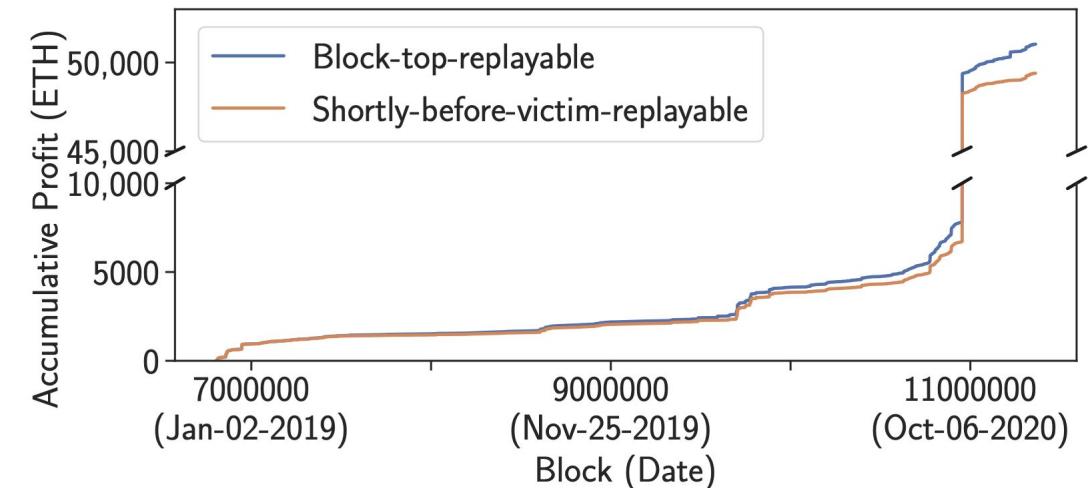
**Function**  $\text{ConstructReplay}(T_V, \mathcal{A})$ :

```
 $T.sender \leftarrow \mathcal{A}$ 
 $T.value \leftarrow T_V.value$ 
 $T.input \leftarrow \text{substituting } T_V.sender \text{ in } T_V.input \text{ with } \mathcal{A}$ 
return  $T$ 
end
```

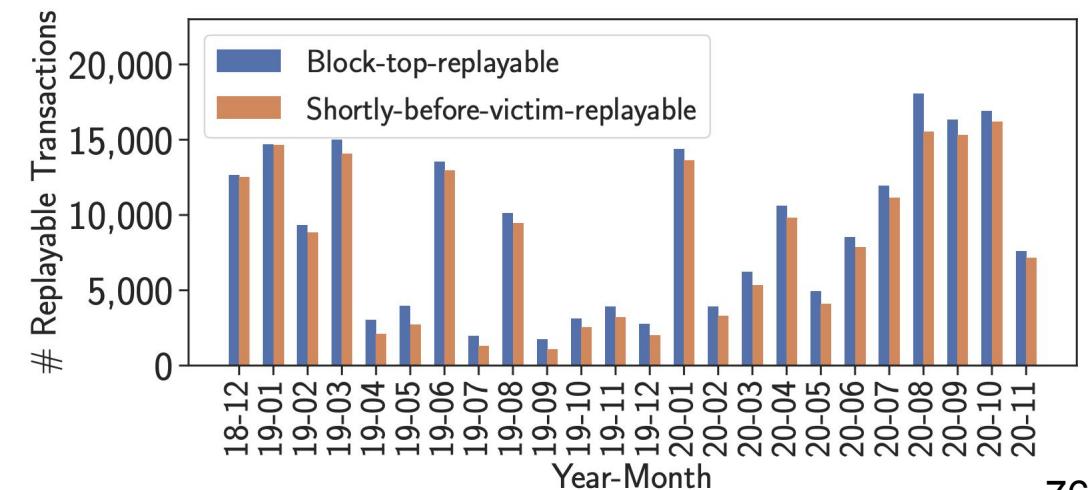
**Algorithm**  $\text{TransactionReplay}(T_V, \mathcal{A})$ :

```
 $T_{replay} \leftarrow \text{ConstructReplay}(T_V, \mathcal{A})$ 
Concretely Execute  $T_{replay}$  upon block  $B_i$ 
if  $T_{replay}$  is profitable then
| Front-run  $T_V$  with  $T_{replay}$ 
end
end
```

---



(a) Accumulative profit that can be extracted by replay attacks.

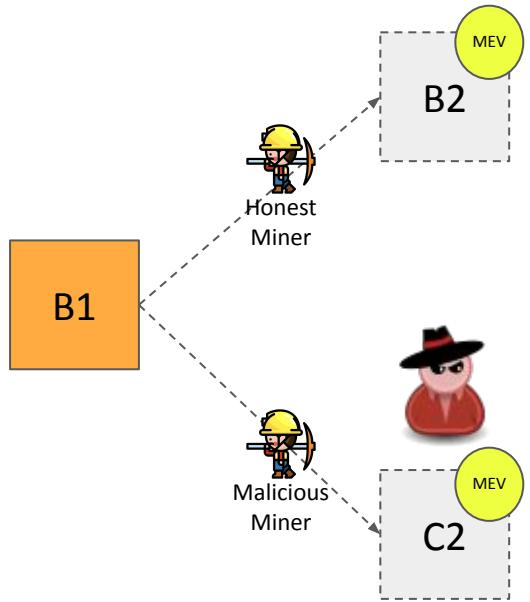




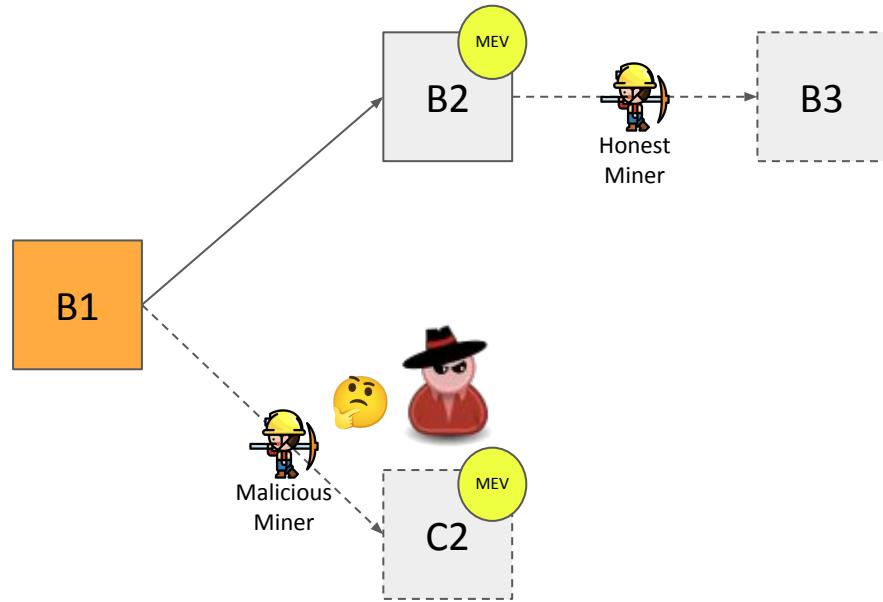
# BEV Forking and Chain Reorganisation

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# The dangers of naively maximizing MEV

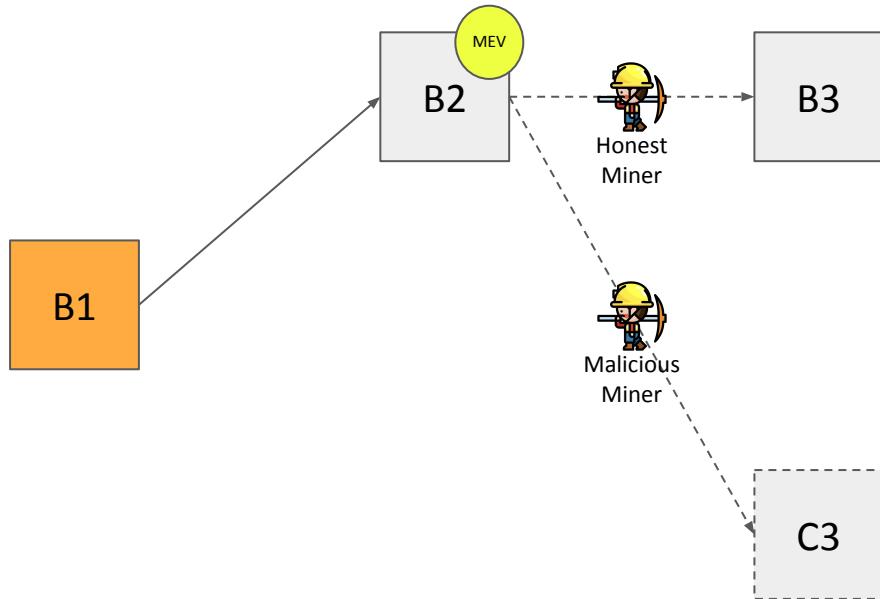


# The dangers of naively maximizing MEV



# The dangers of naively maximizing MEV

Case 1:

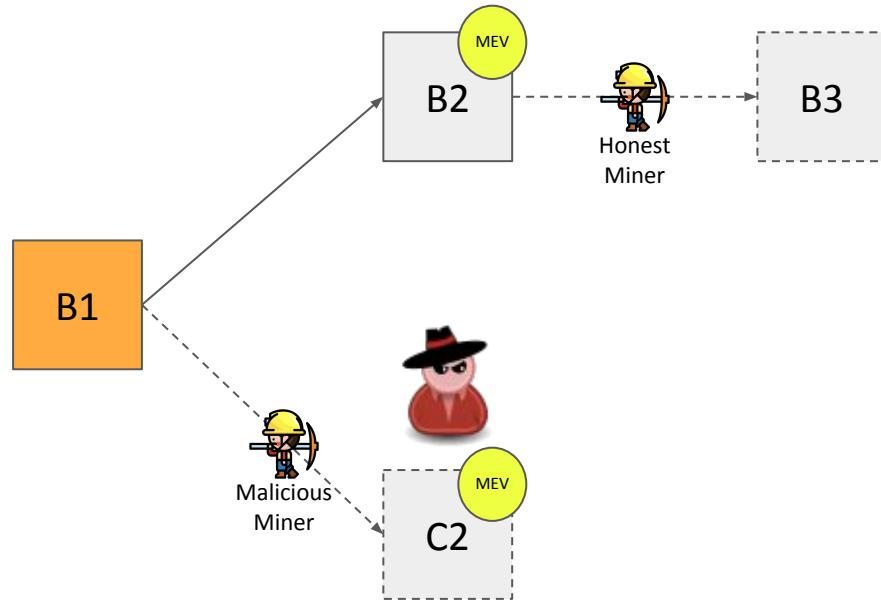


**Case 1:**

Malicious miner forfeits MEV opportunity

# The dangers of naively maximizing MEV

Case 2:



**Case 1:**

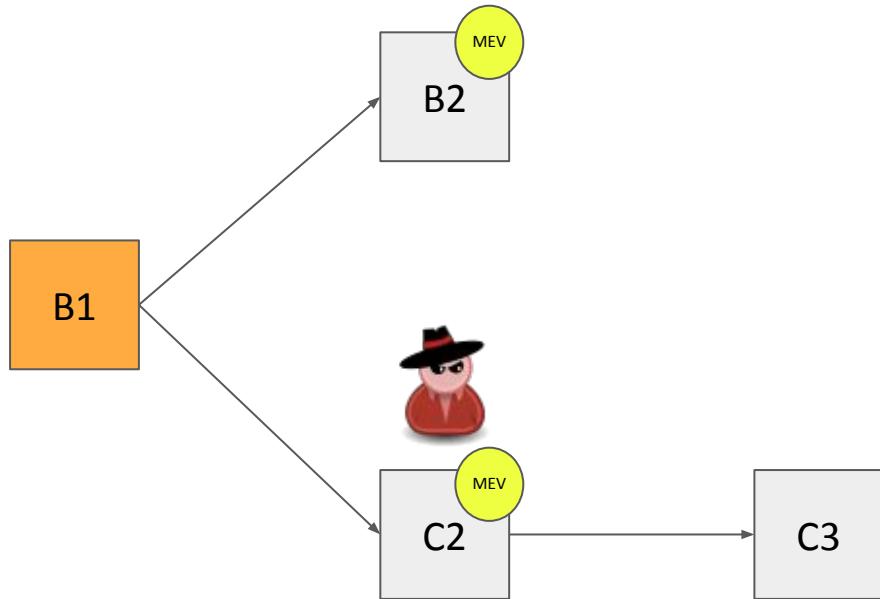
Malicious miner forfeits MEV opportunity

**Case 2:**

Keeps mining block C2

# The dangers of naively maximizing MEV

Case 2:



**Case 1:**

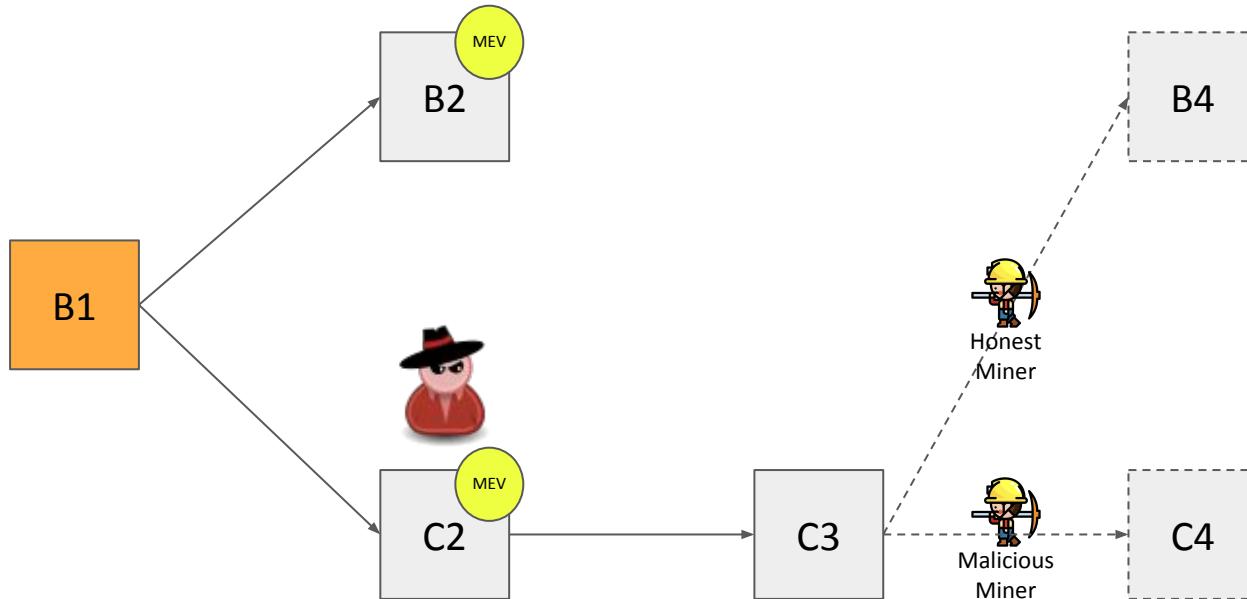
Malicious miner forfeits MEV opportunity

**Case 2:**

Keeps mining block C2

# The dangers of naively maximizing MEV

Case 2:



Case 1:

Malicious miner forfeits MEV opportunity

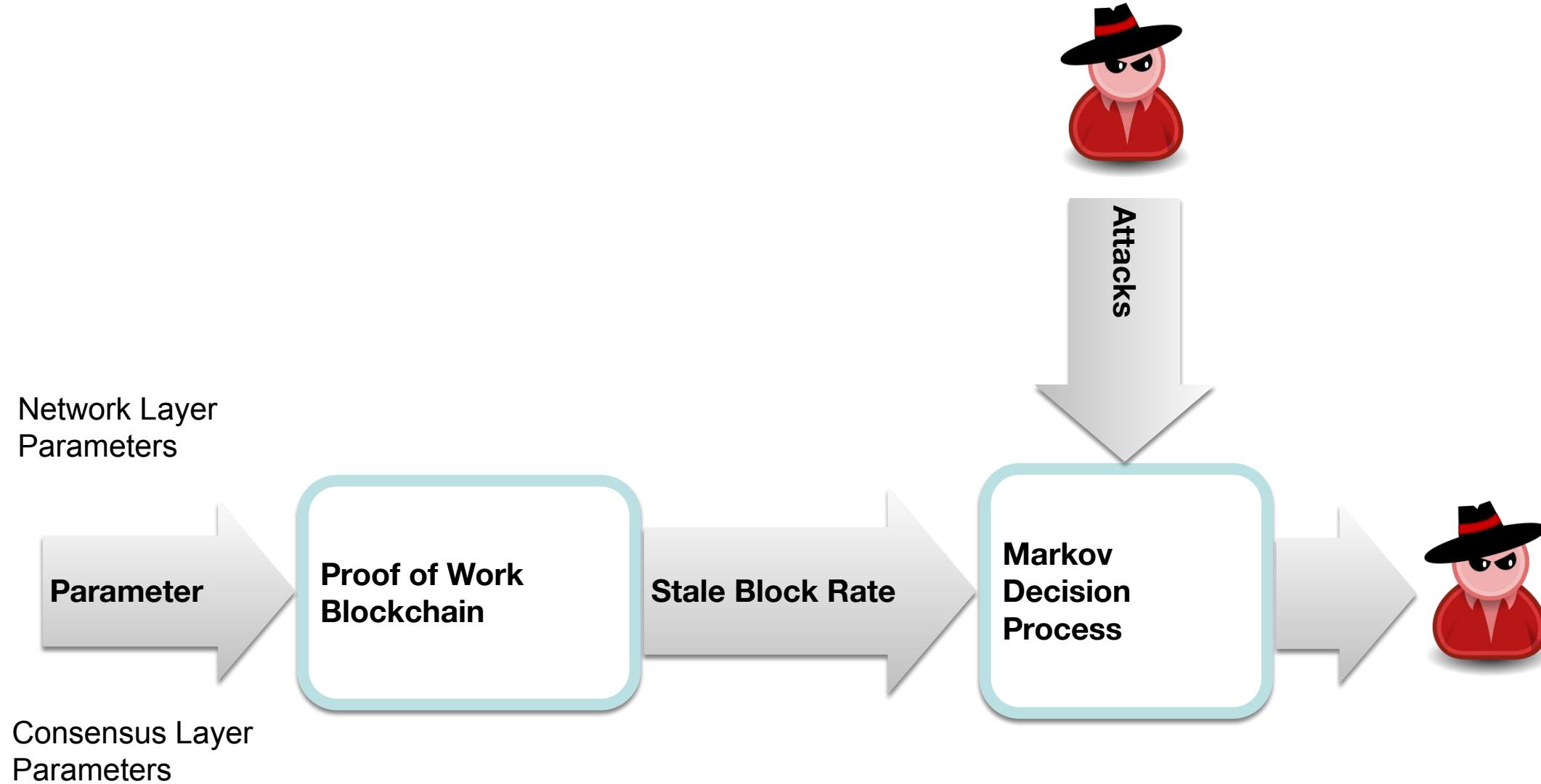
Case 2:

Keeps mining on block C2

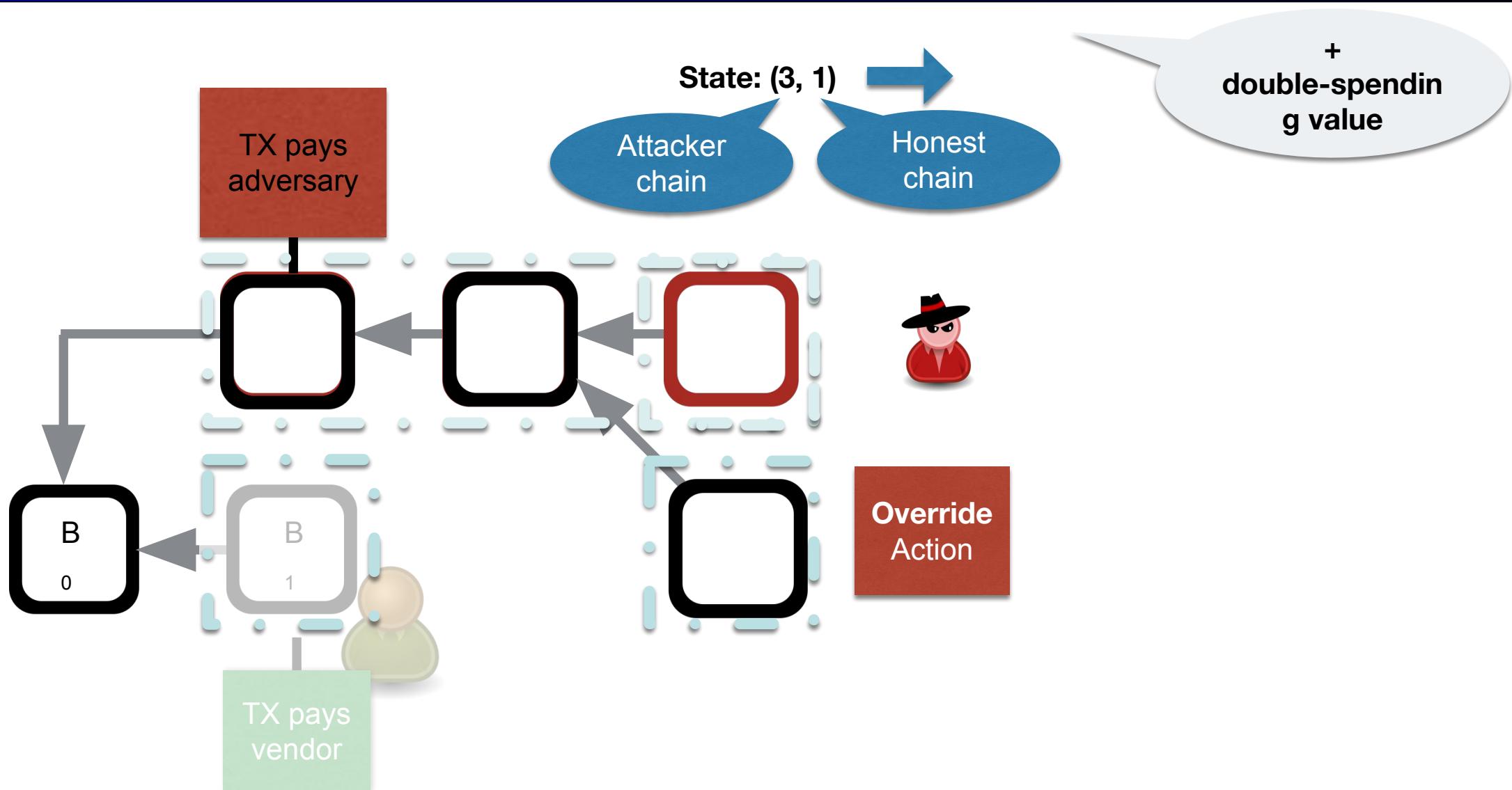
- Waste computational power
- Increase stale block rates and risks for:
  - Double spending
  - Selfish mining



# Markov Decision Process (MDP)



# Markov Decision Process (MDP)



# Reducing MEV is the key to security (example)



10%  
miner

+



MEV, 4x average  
block reward

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[“On the just-in-time discovery of profit-generating transactions in defi protocols.”](#) peer-reviewed at S&P’21

# Reducing MEV is the key to security (example)

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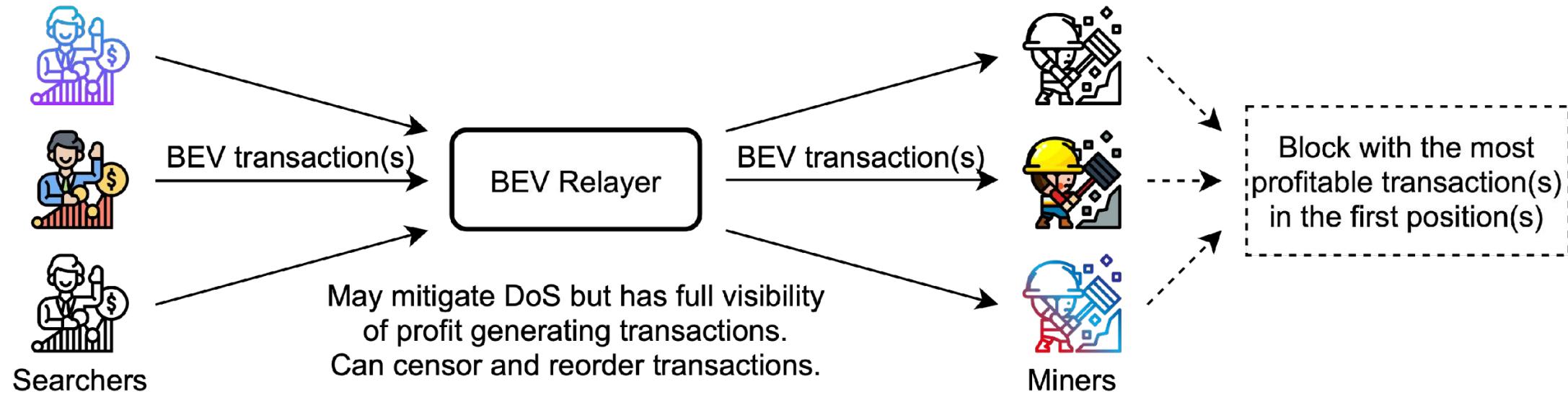
874x



# BEV Relayer & How to Mitigate BEV?

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# BEV Relay Architecture



# BEV Relayer Concerns

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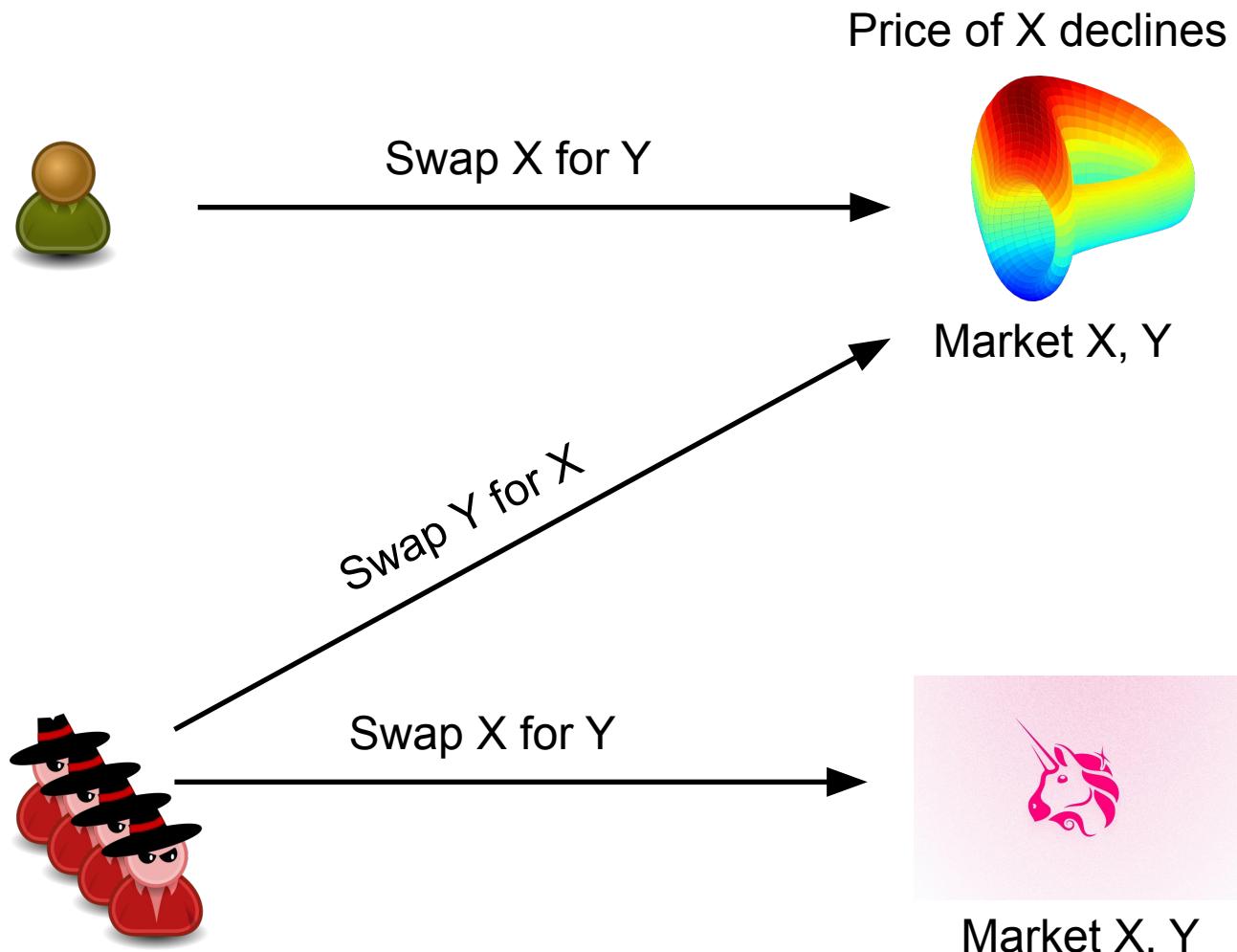
- BEV provably incentivises miners to fork (cf. S&P'21)
- BEV relayer centralise the P2P Network
- The relayer may resell/profit from searcher strategies
- The relay system doesn't necessarily reduce P2P overhead
- A for profit company distributes the geth client to >50% of the miners
- Innocent users are being stolen from systematically

# Anti-MEV Solution Space

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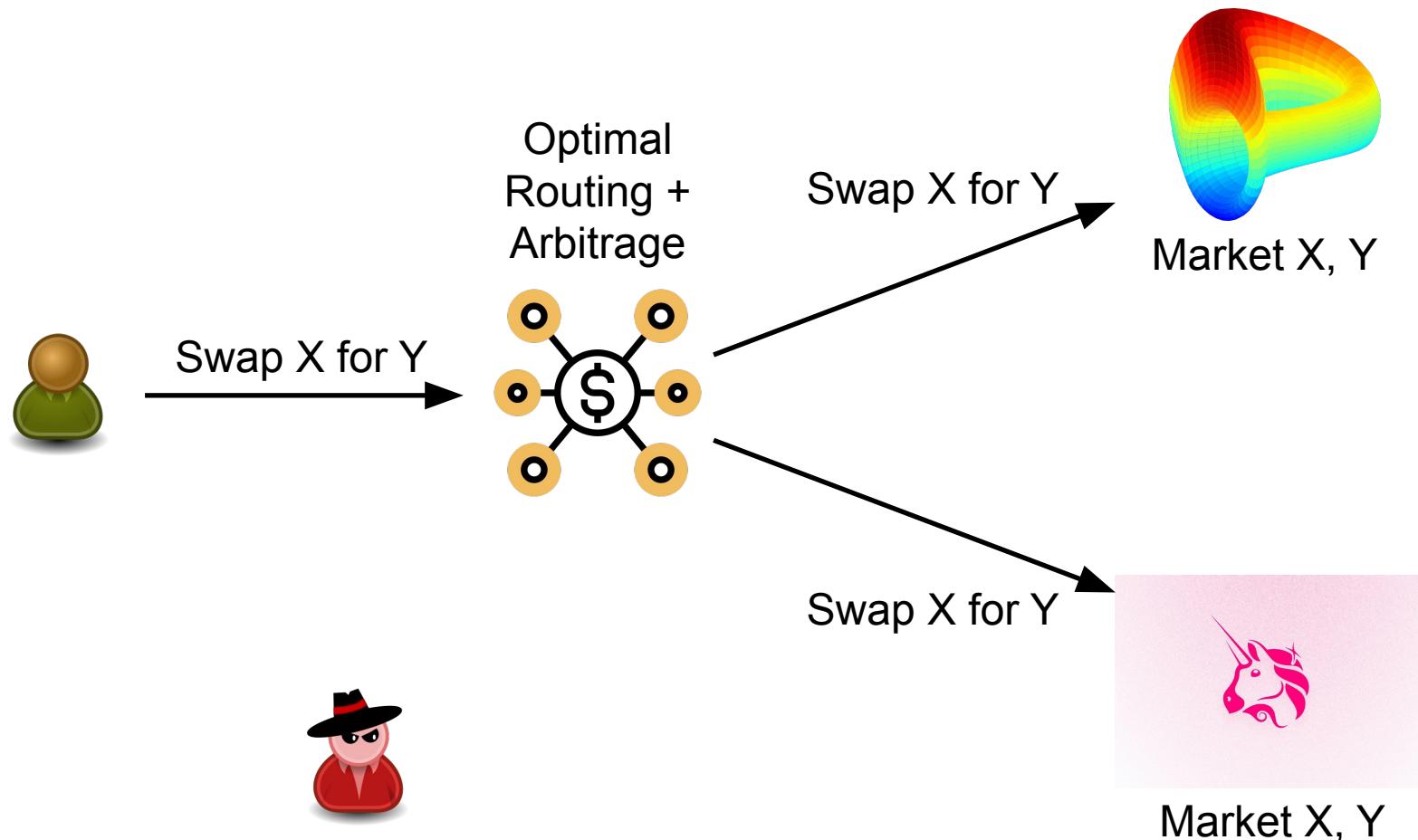
- Fair-Ordering on the Blockchain Layer
  - e.g., Aequitas Protocol Family
- Fixing MEV of existing dApps
  - Merging AMM DEX into one
    - On-chain aggregators such as A2MM (see DEX lecture)
- Designing MEV-Mindful dApps
  - Avoiding MEV by design
  - e.g., a price oracle update immediate performs a liquidation
- Might not fix cross-chain MEV..

# Application-Specific MEV Mitigation



- Causes
  - Back-run Flooding
  - Network Congestions
  - Price Gas Auctions
  - Transaction Fee Increase
- The user forgoes an arbitrage opportunity.

# Application-Specific MEV Mitigation



- Cons
  - Higher Gas Fees
- Pros
  - Better ex rate
  - Arbitrage profit
  - MEV reduction
  - Healthier chain