

# Decentralized Finance

## Introduction to Blockchain technology

Instructors: **Dan Boneh**, Arthur Gervais, Andrew Miller, Christine Parlour, Dawn Song



# What is a blockchain?

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Abstract answer: a blockchain provides coordination between many parties, when there is no single trusted party

if trusted party exists  $\Rightarrow$  no need for a blockchain

[financial systems: often no trusted party]

# What is a blockchain?

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**user facing tools** (cloud servers)

**applications** (DAPPs, smart contracts)

**compute layer** (blockchain computer)

**consensus layer**

# Consensus layer (informal – not the topic of this course)

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A public append-only data structure:

- **Persistence:** once added, data can never be removed\*
- **Consensus:** all honest participants have the same data\*\*
- **Liveness:** honest participants can add new transactions
- **Open(?)**: anyone can add data

achieved by replication

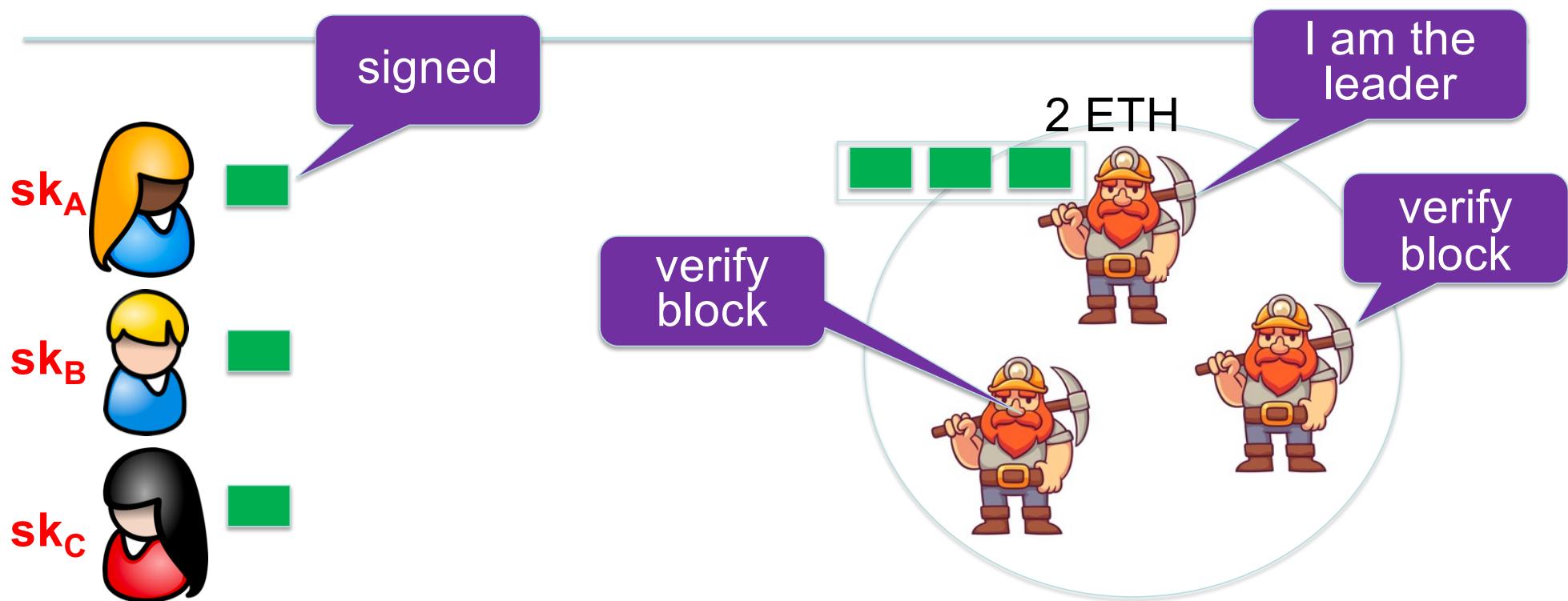


Layer 1:

consensus layer

# How are blocks added to chain?

blockchain

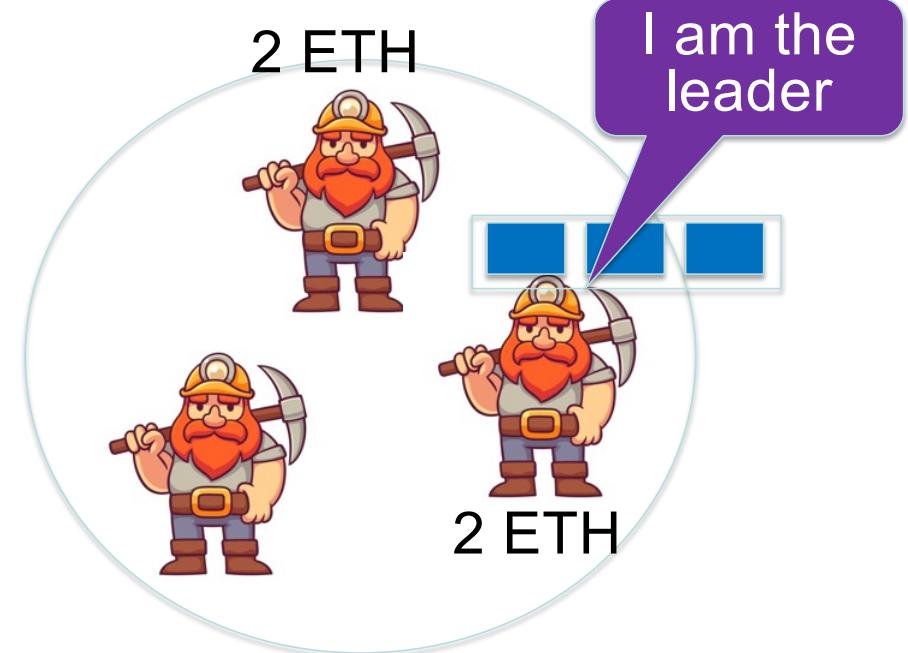
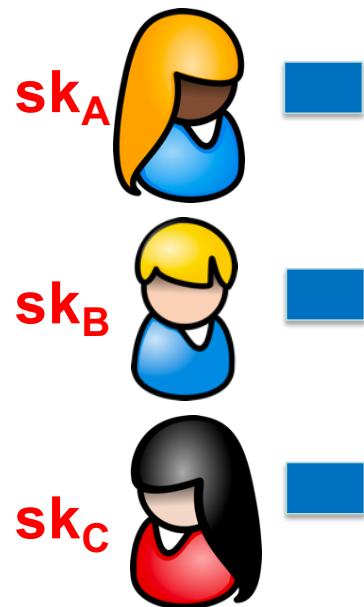


# How are blocks added to chain?

blockchain



...



# Compute layer: The blockchain computer

**DAPP logic is encoded in a program that runs on blockchain**

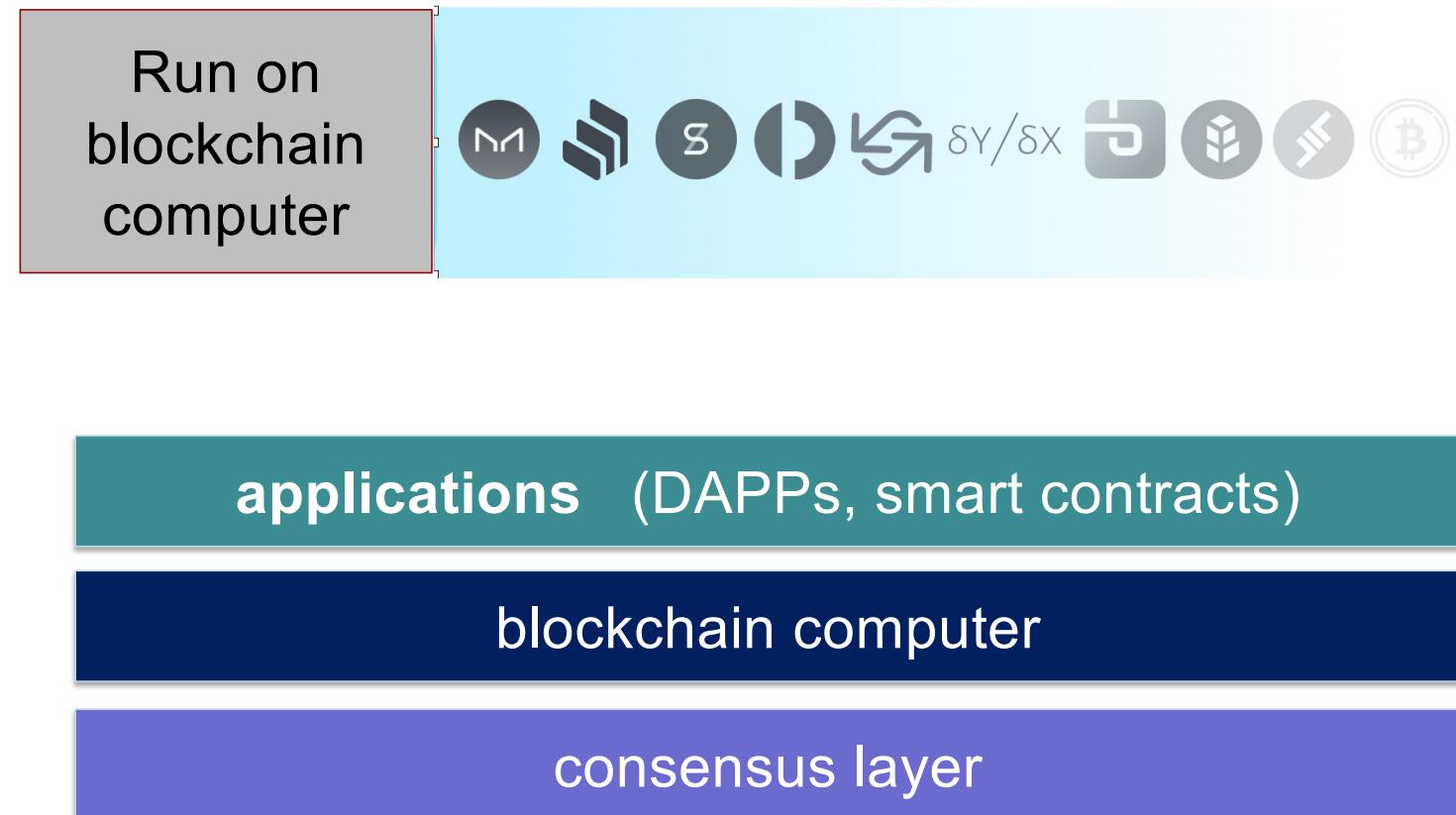
- Rules are enforced by a public program (public source code)
  - ⇒ transparency: no single trusted 3<sup>rd</sup> party
- The DAPP program is executed by parties who create new blocks
  - ⇒ public verifiability: everyone can verify state transitions

compute layer

consensus layer

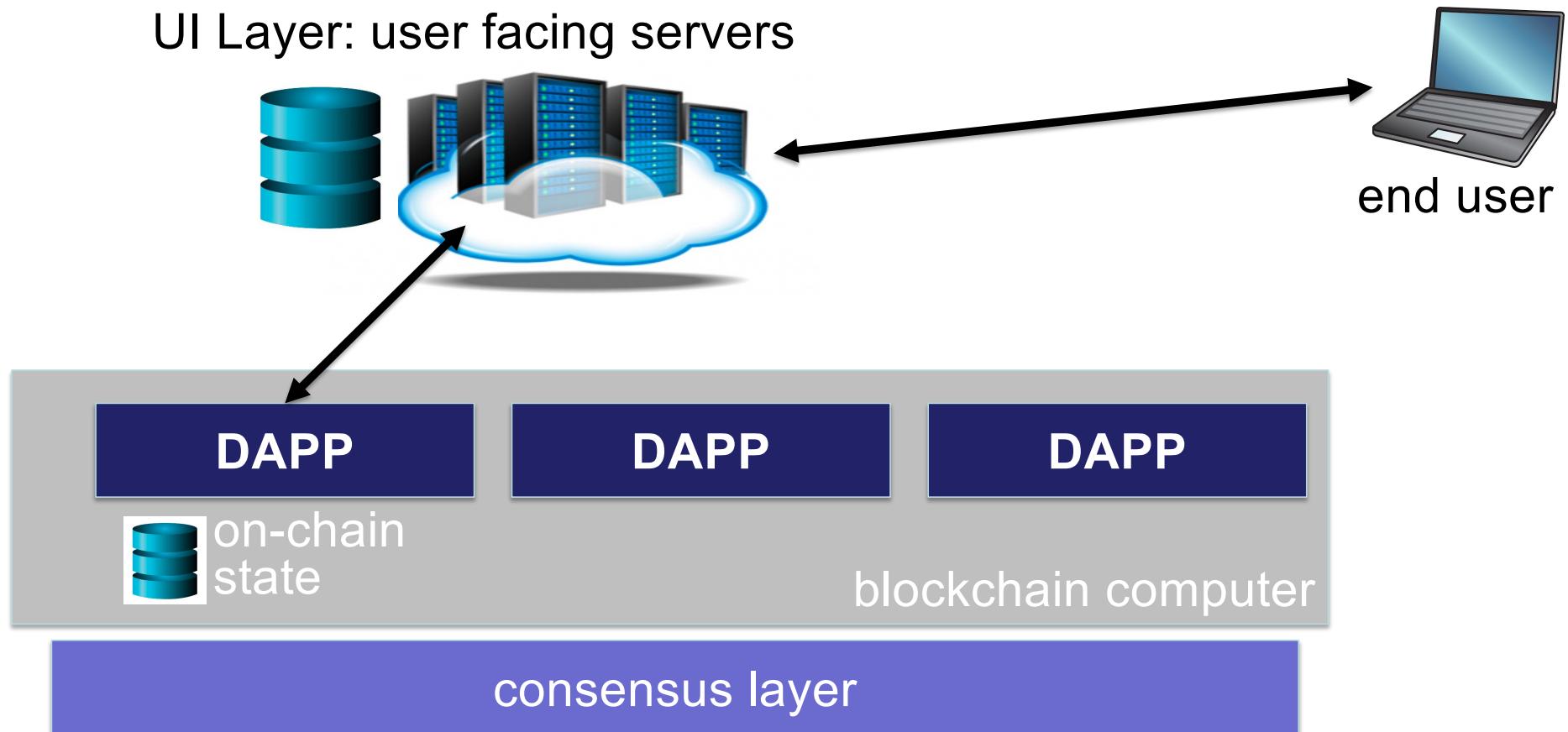
# Apps layer: Decentralized applications (DAPPS)

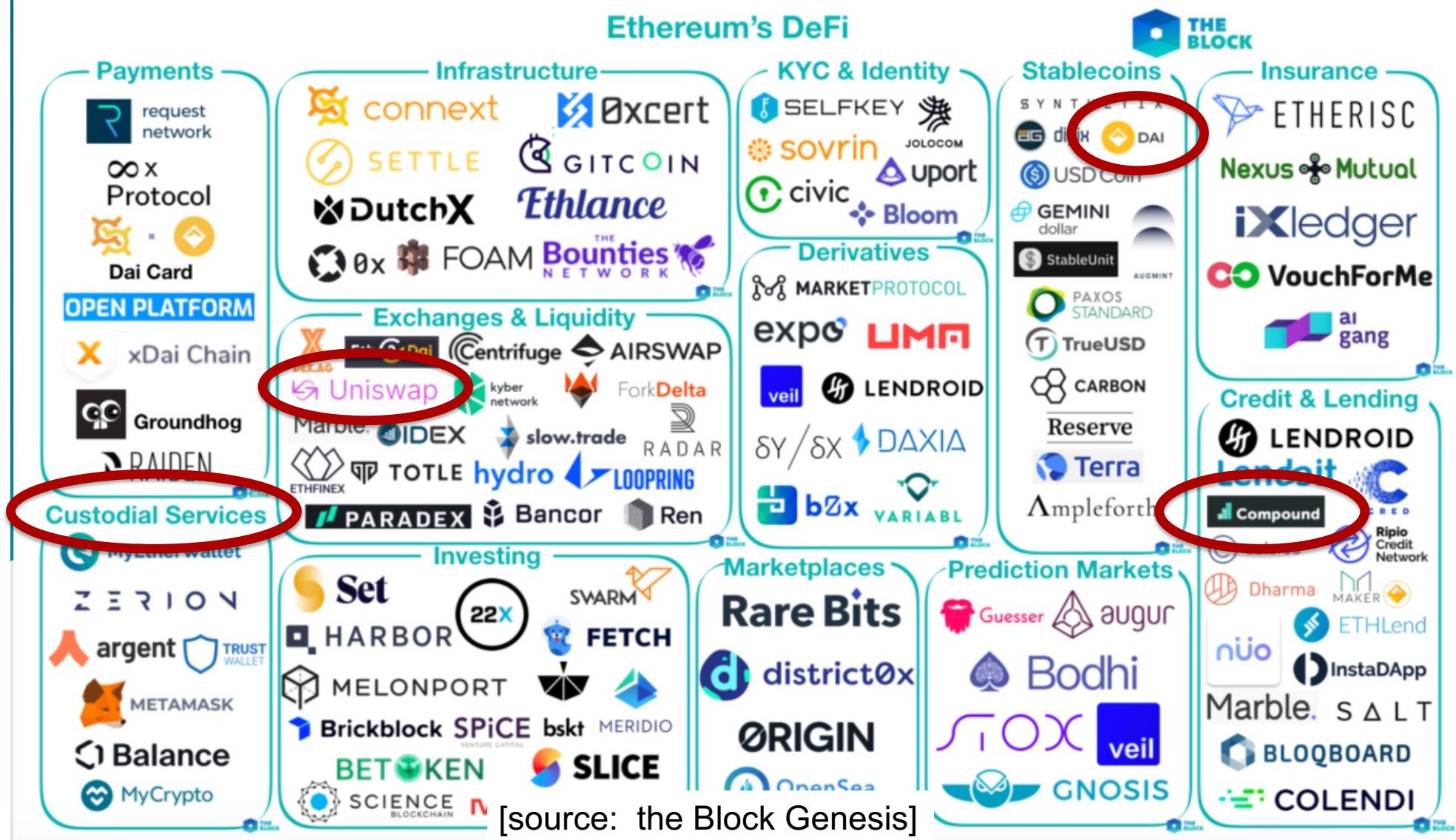
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# UI Layer: Common DAPP architecture

UI Layer: user facing servers





# lots of experiments ...

DEFI PULSE	Name	Chain	Category	Locked (USD) ▼
🏆 1.	Aave	Multichain	Lending	\$15.63B
🥈 2.	Curve Finance	Multichain	DEXes	\$10.71B
🥉 3.	InstaDApp	Ethereum	Lending	\$10.66B
4.	Compound	Ethereum	Lending	\$10.43B
5.	Maker	Ethereum	Lending	\$9.11B
6.	Uniswap	Ethereum	DEXes	\$7.29B
7.	Convex Finance	Ethereum	Assets	\$5.59B

Let's get started ...

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Next segment: cryptographic background

See you there

# Cryptographic Background: hash functions

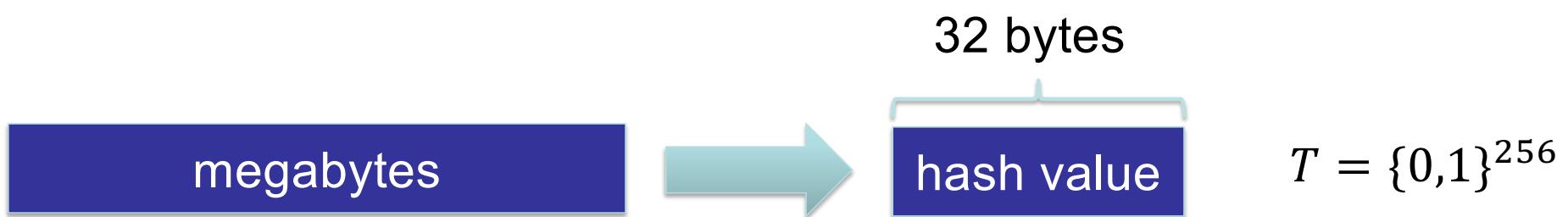
<https://defi-learning.org/>

# Cryptography Background

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## (1) cryptographic hash functions

An efficiently computable function  $H: M \rightarrow T$   
where  $|M| \gg |T|$



# Collision resistance

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Def: a collision for  $H: M \rightarrow T$  is pair  $x \neq y \in M$  s.t.  $H(x) = H(y)$

$|M| \gg |T|$  implies that many collisions exist

Def: a function  $H: M \rightarrow T$  is collision resistant if it is “hard” to find even a single collision for  $H$  (we say  $H$  is a CRH)

Example: **SHA256:**  $\{x : \text{len}(x) < 2^{64} \text{ bytes}\} \rightarrow \{0,1\}^{256}$

details in crypto MOOC

## An application: committing to data

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Alice has a large file  $m$ . She publishes  $h = H(m)$  (32 bytes)

Bob has  $h$ . Later Alice sends  $m'$  s.t.  $H(m') = h$

$H$  is a CRH  $\Rightarrow$  Bob is convinced that  $m' = m$   
(otherwise,  $m$  and  $m'$  are a collision for  $H$ )

We say that  $h = H(m)$  is a **binding commitment** to  $m$

(note: not hiding,  $h$  may leak information about  $m$ )

# Committing to a list (of transactions)

Alice has  $S = (m_1, m_2, \dots, m_n)$

32 bytes

## Goal:

- Alice publishes a short binding commitment to  $S$ ,  $h = \text{commit}(S)$
- Bob has  $h$ . Given  $(m_i, \text{proof } \pi_i)$  can check that  $S[i] = m_i$   
Bob runs  $\text{verify}(h, i, m_i, \pi_i) \rightarrow \text{accept/reject}$

security: adv. cannot find  $(S, i, m, \pi)$  s.t.  $m \neq S[i]$  and

$\text{verify}(h, i, m, \pi) = \text{accept}$  where  $h = \text{commit}(S)$

# Merkle tree (Merkle 1989)

commitment



$h$

Merkle  
tree  
commitment

$m_1 \ m_2 \ m_3 \ m_4 \ m_5 \ m_6 \ m_7 \ m_8$



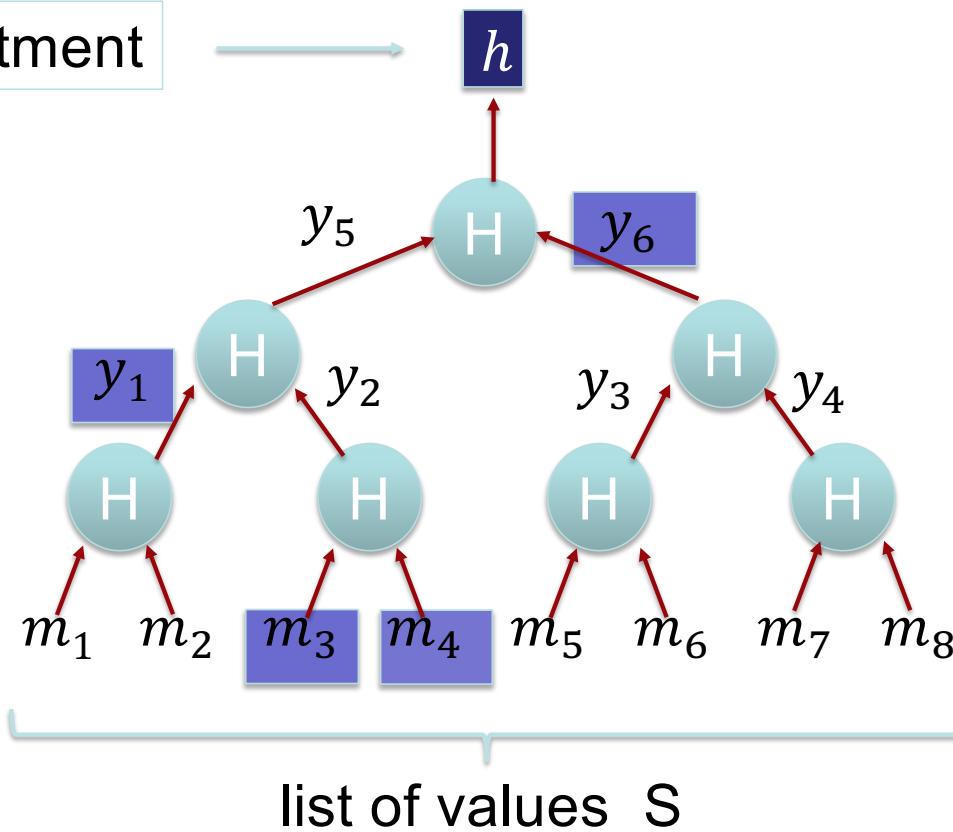
list of values  $S$

Goal:

- commit to list  $S$  of size  $n$
- Later prove  $S[i] = m_i$

# Merkle tree (Merkle 1989)

commitment



Goal:

- commit to list  $S$  of size  $n$
- Later prove  $S[i] = m_i$

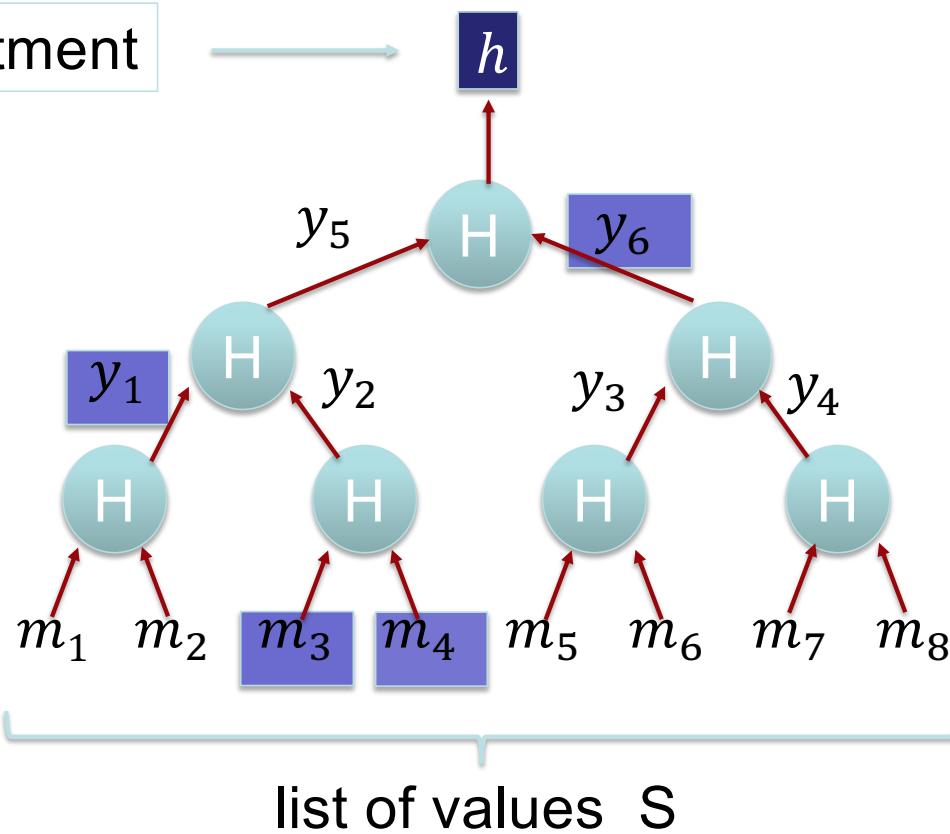
To prove  $S[4] = m_4$

proof  $\pi = (m_3, y_1, y_6)$

length of  $\pi$ :  $\log_2 n$

# Merkle tree (Merkle 1989)

commitment



To prove  $S[4] = m_4$

proof  $\pi = (m_3, y_1, y_6)$

Bob does:

$$y_2 \leftarrow H(m_3, m_4)$$

$$y_5 \leftarrow H(y_1, y_2)$$

$$h' \leftarrow H(y_5, y_6)$$

accept if  $h = h'$

## Merkle tree (Merkle 1989)

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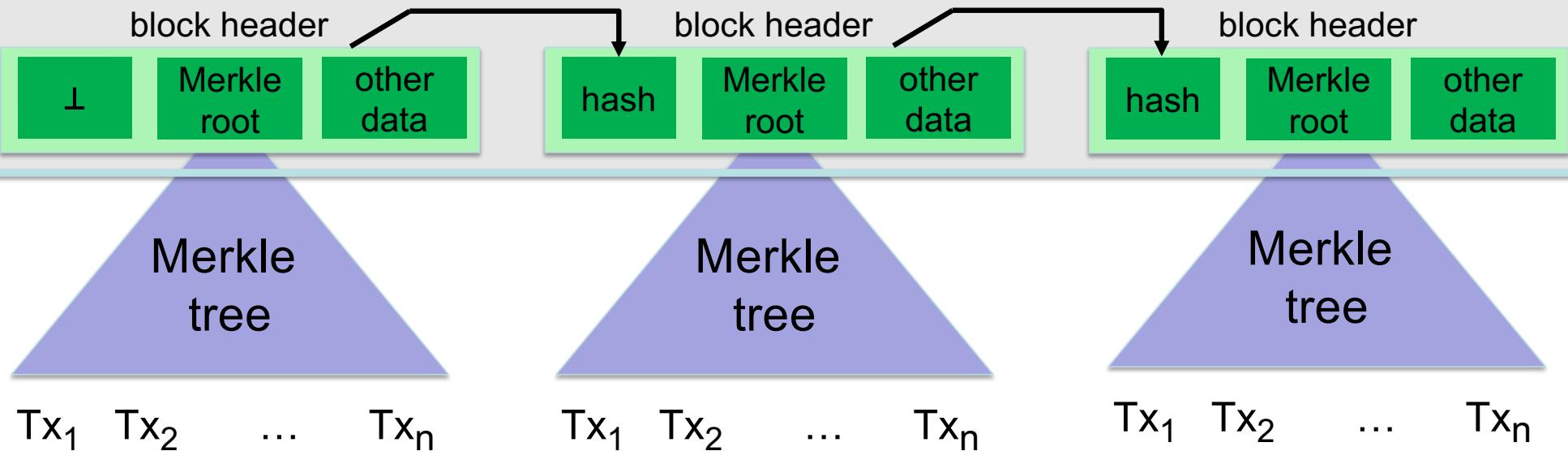
Thm: H CRH  $\Rightarrow$  adv. cannot find  $(S, i, m, \pi)$  s.t.  $m \neq S[i]$  and  
verify( $h, i, m, \pi$ ) = accept where  $h = \text{commit}(S)$   
(to prove, prove the contra-positive)

How is this useful? Super useful. Example:

- When writing a block of transactions  $S$  to the blockchain, suffices to write  $\text{commit}(S)$  to chain. Keep chain small.
- Later, can prove contents of every Tx.

# Abstract block chain

blockchain



Merkle proofs are used to prove that a Tx is “on the block chain”

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Next segment: digital signatures

How to authorize transactions??



# Cryptographic Background: Digital Signatures

<https://defi-learning.org/>

# Digital Signatures

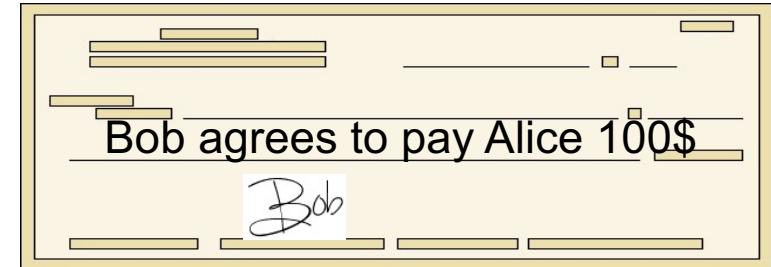
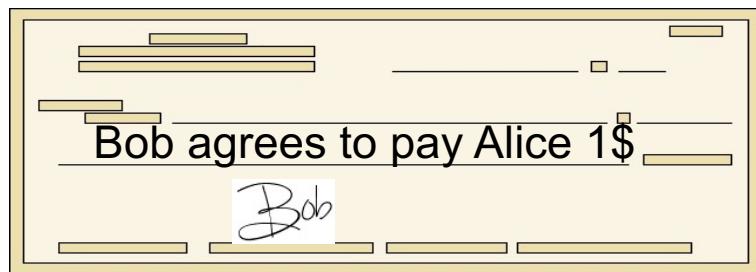
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- In the last segment we looked at cryptographic hash functions.
- In this segment we will look at digital signatures:  
how to approve a transaction?

# Signatures

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Physical signatures: bind transaction to author

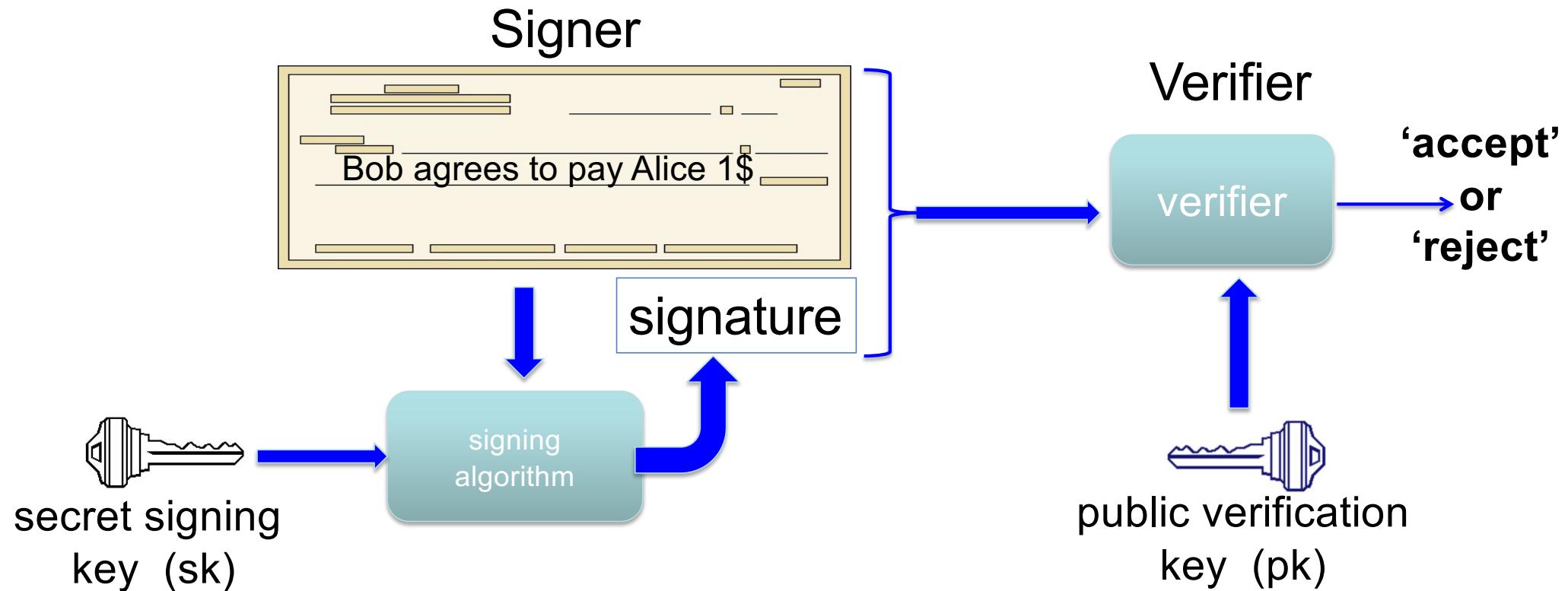


Problem in the digital world:

anyone can copy Bob's signature from one doc to another

# Digital signatures

Solution: make signature depend on document



# Digital signatures: syntax

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Def: a signature scheme is a triple of algorithms:

- **Gen()**: outputs a key pair  $(pk, sk)$
- **Sign**( $sk, msg$ ) outputs sig.  $\sigma$
- **Verify**( $pk, msg, \sigma$ ) outputs ‘accept’ or ‘reject’

Secure signatures: (informal)

Adversary who sees  $pk$  and sigs on many messages of her choice, cannot forge a signature on a new message.

# Families of signature schemes

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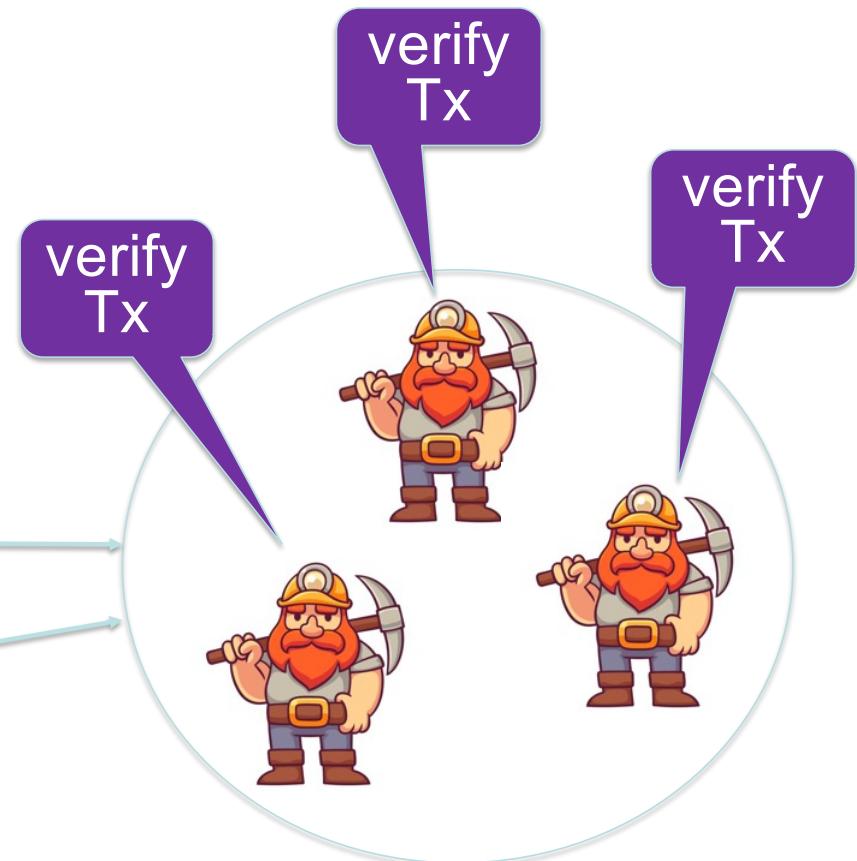
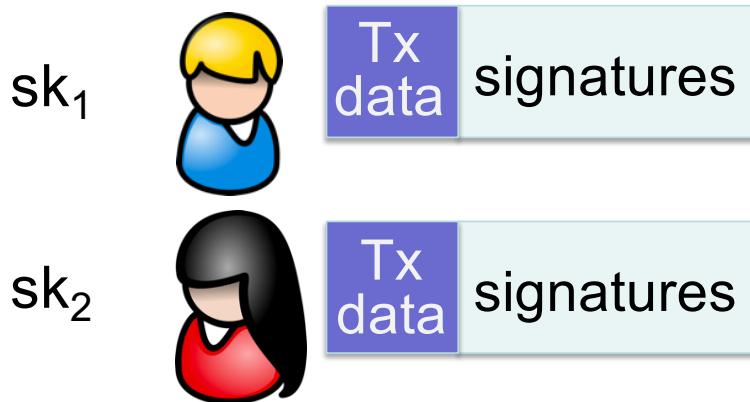
1. RSA signatures (not used in blockchains):
  - long sigs and public keys ( $\geq 256$  bytes), **fast to verify**  
传统的RSA比较快  
现在更多使用的是Schnorr，而不是ECDSA
2. Discrete-log signatures: Schnorr and ECDSA (Bitcoin, Ethereum)
  - short sigs (48 or 64 bytes) and public keys (32 bytes)
3. BLS signatures: 48 bytes, **可聚**, aggregatable, easy threshold  
可以通过聚合的方式大大降低对于空间需求，这里以太坊已经在使用了  
(Ethereum 2.0, Chia, Dfinity)
  -
4. Post-quantum signatures: long ( $\geq 768$  bytes)

# Signatures on the blockchain

签名在区块链中的一些用处

Signatures are used everywhere:

- ensure Tx authorization,
- governance votes,
- consensus protocol votes.



# SNARK proofs

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We covered two important cryptographic primitives:

1. Collision resistant hash functions and Merkle trees,
2. Digital signatures.

Another important cryptographic primitive is a **SNARK proof**:

- Used for scaling and privacy
- We will discuss SNARKs in detail in the lecture on privacy

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Next segment: scaling the blockchains

Can we make it fast??



# Scaling Blockchains

区块链可扩展性

<https://defi-learning.org/>

# Scaling

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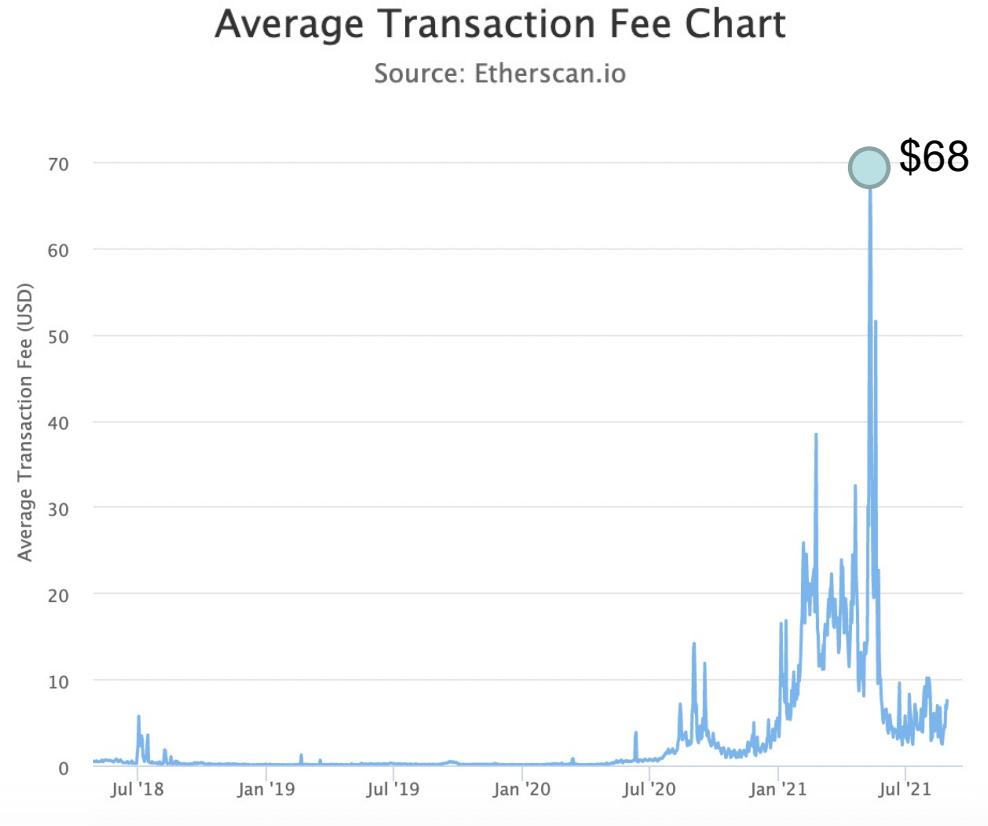
Transaction rates (Tx/sec):

- Bitcoin: can process about **5 (Tx/sec)**
  - Ethereum: can process about **20 (Tx/sec)**
- 
- Tx Fees fluctuate:  
2\$ to 60\$ for simple Tx

吞吐量低下，交易费用波动大、普遍高昂

# Ethereum Tx fees (gas prices)

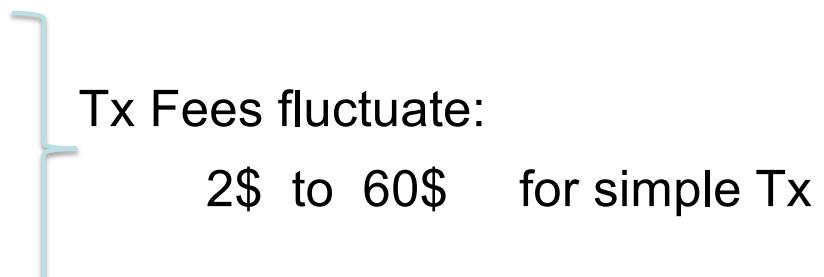
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# Scaling

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Transaction rates (Tx/sec):

- Bitcoin: can process about **5 (Tx/sec)**
  - Ethereum: can process about **20 (Tx/sec)**
  - The visa network: can process up to **24,000 (Tx/sec)**
- 
- Tx Fees fluctuate:  
2\$ to 60\$ for simple Tx

visa相对的交易速度快，手续费低

**Can we scale blockchains to visa speeds? ... with low Tx fees**

# Scaling approaches

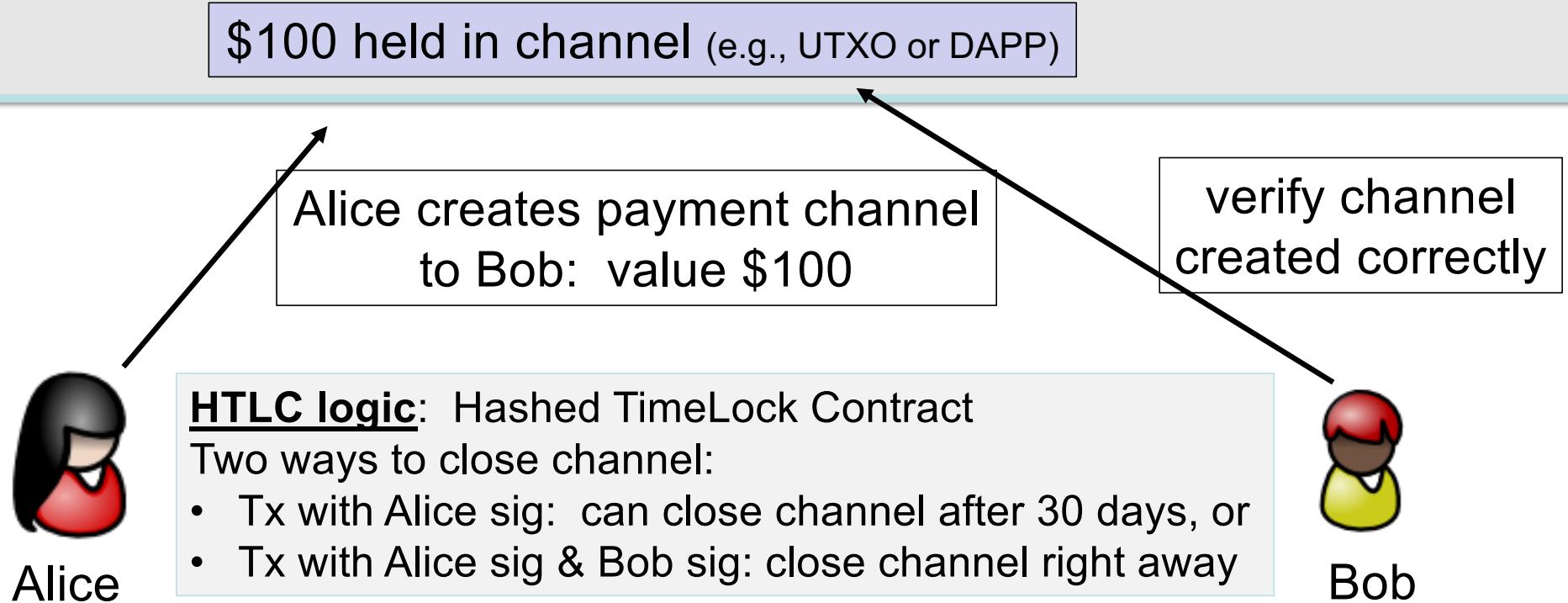
Many approaches to scaling blockchains:

- **Faster consensus**: 更快的共识 modern blockchains (e.g., Solana, Polkadot, Avalanche, ...)
- **Payment channels**: most Tx are off chain Peer-to-Peer (e.g., Lightening)  
离线进行点对点交易
- **Layer 2 approaches**:  
**zkRollup, optimistic Rollup**: batch many Tx into a single Tx
- Sidechains: Polygon and others  
从这个角度来说的话，layer2 中并不完全包括 lightening network。并且ZK和OProllup也不全算是 Layer2的解决方案
- many other ideas ...

闪电网络

# (1) Payment channels (high level idea)

blockchain



# (1) Payment channels (high level idea)

blockchain

\$100 held in channel (e.g., UTXO)

Bob can sign Tx and close channel  
... but he would rather wait (up to 30 days)



Alice

pay Bob: 5\$

Tx: distribute funds: Alice: 95; Bob: 5

$\text{sig}_{\text{Alice}}$



Bob

( off chain message! )

# (1) Payment channels (high level idea)

blockchain

\$100 held in channel (e.g., UTXO)



Alice

another payment: pay Bob: 15\$

Tx: distribute funds: Alice: 80; Bob: 20

$\text{sig}_{\text{Alice}}$

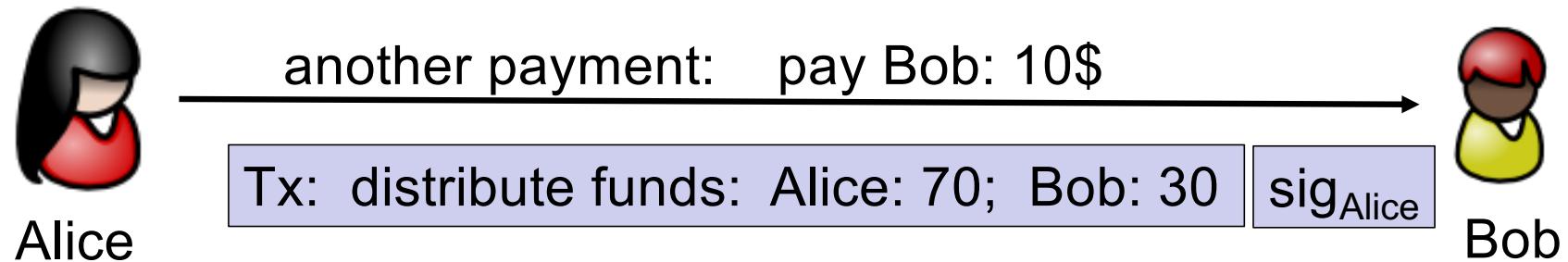


Bob

# (1) Payment channels (high level idea)

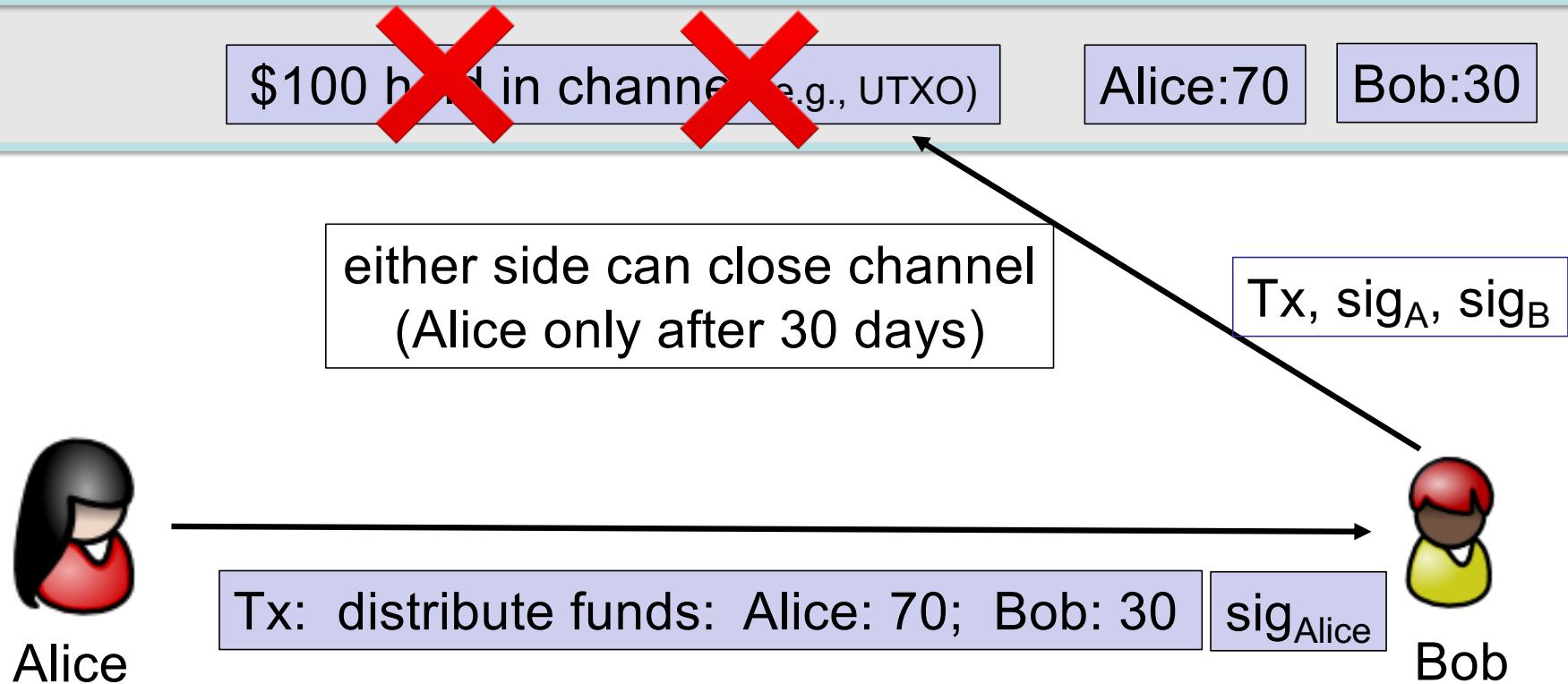
blockchain

\$100 held in channel (e.g., UTXO)



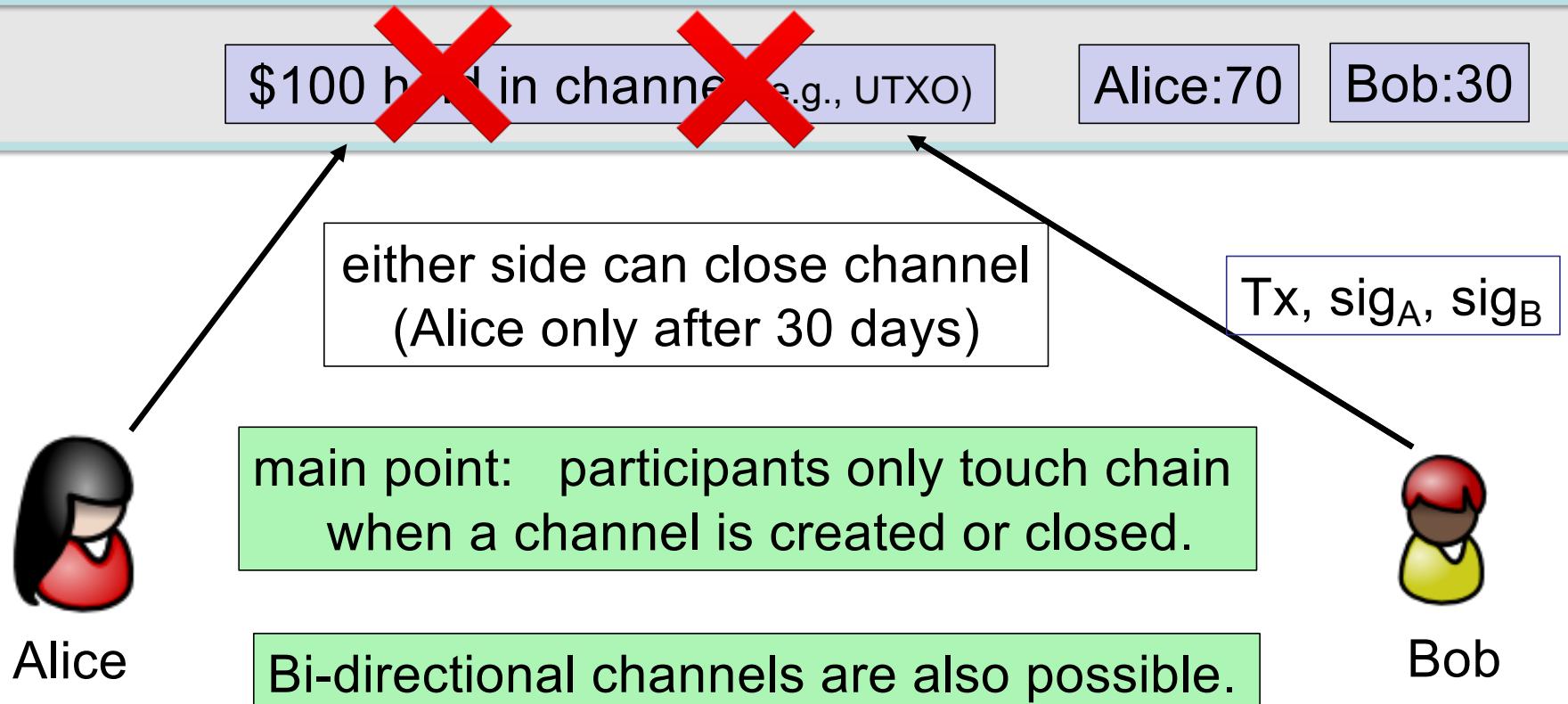
# (1) Payment channels (high level idea)

blockchain



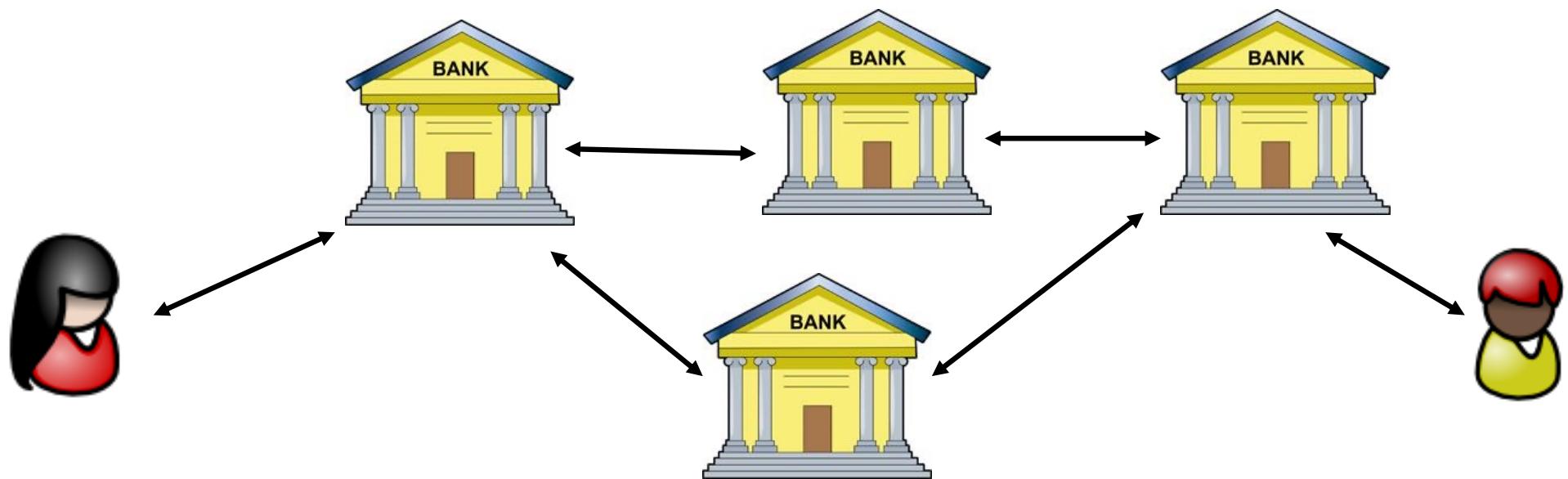
# (1) Payment channels (high level idea)

blockchain



# Payment networks

Lots of bi-directional payment channels



Alice pays Bob by finding the cheapest route through the network  
⇒ while channels are open, nothing touches the blockchain

# The case of El Salvador

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June 24, 2021

8:18 PM PDT

Last Updated 2 months  
ago

Technology

## Bitcoin to become legal tender in El Salvador on Sept 7, 2021

3 minute read

Reuters

Payment channels are necessary to enable state-wide adoption

- Strike wallet: connects to the Bitcoin Lightning network

## (2) Scaling Ethereum Using Rollup

## (2) Scaling Ethereum Using Rollup

Main tool: SNARK (much more on SNARKs later)

C: a program that always terminates in  $\leq B$  steps

x: public input to C,      **w**: private input to C

(C, x, **w**)



prover

short proof  $\pi$

(C, x)



verifier

## (2) Scaling Ethereum Using Rollup

### Main point:

Verifier's run time is  
\*much\* less than running C

x: public input to C,

more on SNARKs

finishes in  $\leq B$  steps

I am convinced  
prover knows w  
s.t.  $C(x, w) = 1$

private input to C

(C, x, **w**)



prover

short proof  $\pi$

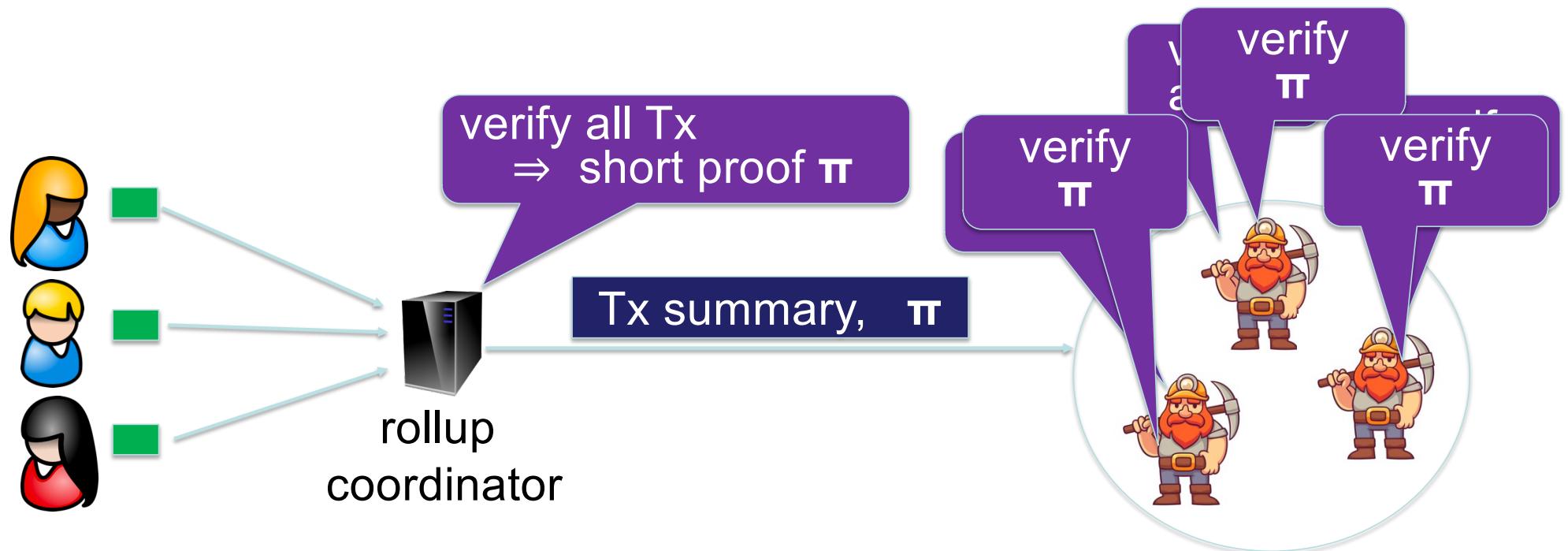
(C, x)



verifier

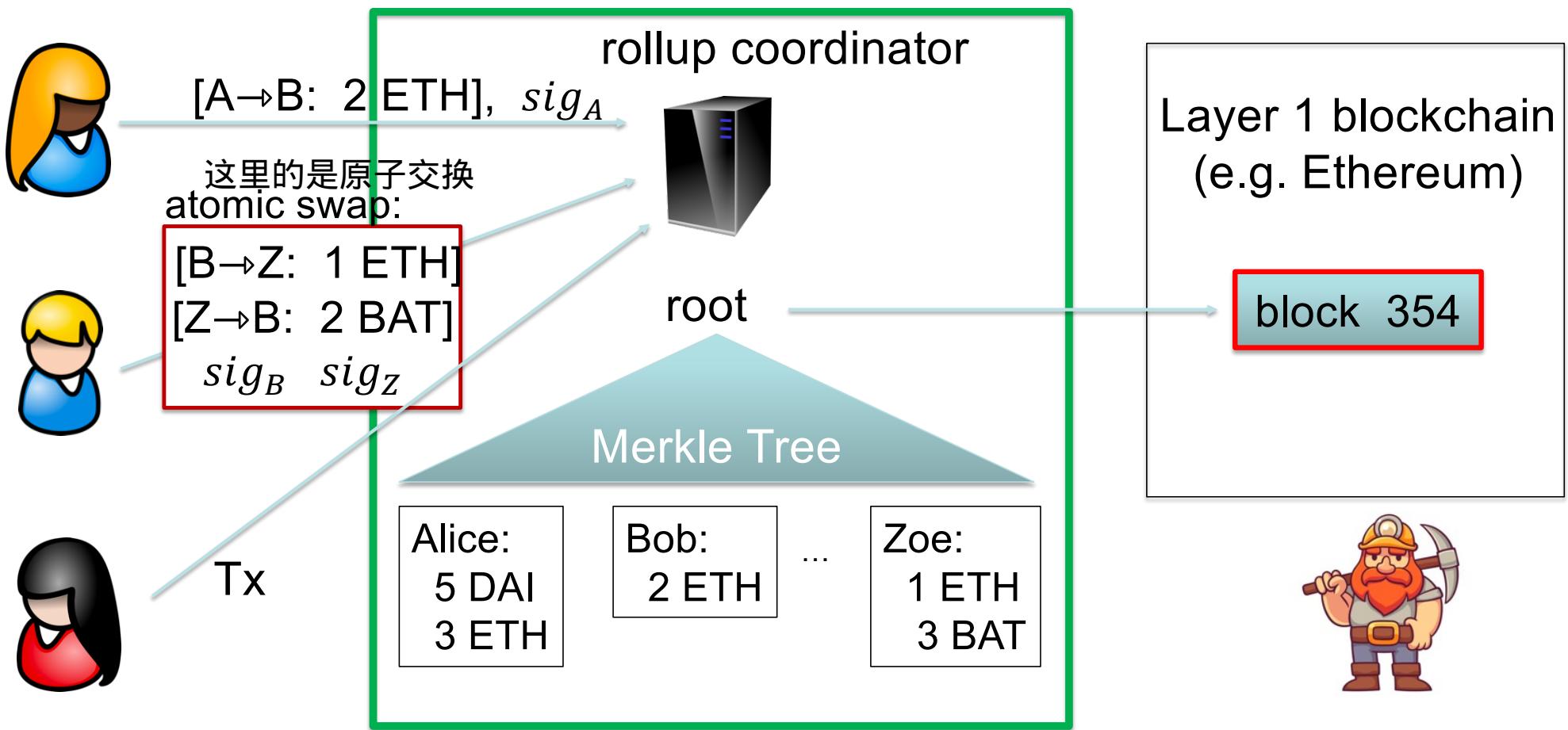
## (2) Rollup: zk and optimistic

Standard L1 chains: every miner must verify every posted Tx

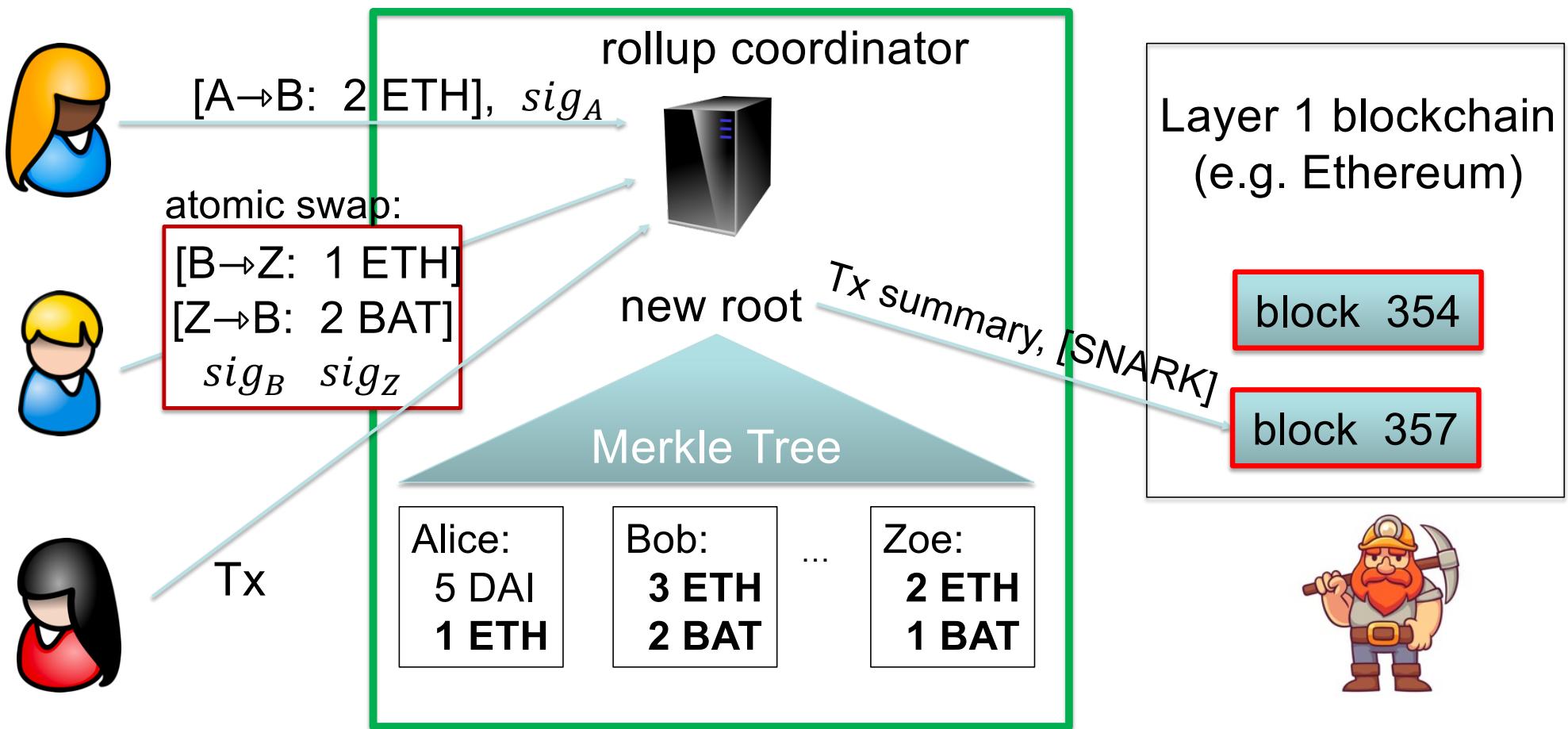


Rollup coordinator: compresses a thousand Tx into one on-chain proof (SNARK)

# zkRollup (simplified)



# zkRollup (simplified)



# Transferring assets to and from L2

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- Transactions within a Rollup system are easy:
  - Batch settlement on L1 network (e.g., Ethereum) 因为L2的工作本身就是以L1为基础的
- Moving funds in to or out of Rollup system ( $L1 \leftrightarrow L2$ )
  - Requires posting more data on L1 network  $\Rightarrow$  higher Tx fees.

L2和L1上的互转手续费比L1到L1的互转还要更高
- Moving funds from one Rollup system to another ( $L2 \leftrightarrow L2$ )
  - Either via L1 network (expensive), or via a direct  $L2 \leftrightarrow L2$  bridge (cheap)

经过一层网络很贵

只经过二层网络则相对低一点点

# Migrating a project from L1 Ethereum to L2 zkRollup

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Upcoming development: **zkEVM** (e.g., MatterLabs and others).

## Solidity compatibility:

- Coordinator can produce a SNARK proof for the execution of a short Solidity program:
  - ⇒ easy to migrate a DAPP from L1 Ethereum to L2 zkRollup.
  - ⇒ reduced Tx fees and increased Tx rate compared to L1

简化的rollup方案

## Optimistic Rollup (simplified) [e.g., Optimism, Arbitrum]

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Same principle as zkRollup, but no SNARK proof

Instead: coordinator posts Tx data on chain without a proof

- Then give a few days for validators to complain:  
if a posted Tx is invalid  $\Rightarrow$   
anyone can submit a **fraud proof** and win a reward,  
Rollup server gets slashed.

Benefit: simple full EVM compatibility, less work for server.

# Data availability: zkSync vs. zkPorter

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Is the coordinator a central point of failure? (centralization fears??)

Answer: No!

coordinator fails  $\Rightarrow$  users find another coordinator to produce proofs

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- Complication: new coordinator needs all current account information
  - How to get the data if the old coordinator is dead?
- Two solutions: zkSync and zkPorter. They work concurrently.

# Data availability: zkSync vs. zkPorter

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- **zkSync**: store all Tx summaries on the L1 blockchain (Ethereum)
  - L1 chain accepts Tx batch only if it includes summary of all Tx
  - Other coordinators can reconstruct L2 state from L1 blockchain
  - Downside: higher Ethereum Tx fees. Good for high value assets  
缺点
- **zkPorter**: store Tx data on a new blockchain
  - maintained by a set of staked coordinators
  - Cheap off-chain storage, but lower guarantee than zkSync  
便宜的链外存储
- Customer can choose how coordinator will store its account.

That's it on this topic ...

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Next segment: interoperability

How to move assets from one chain to another



# Interchain Interoperability

<https://defi-learning.org/>



Solana

**Serum  
DEX**



Bitcoin

Ethereum



Flow

Can I use  
Serum??



Polkadot

20 DOT

交互性  
Interoperability

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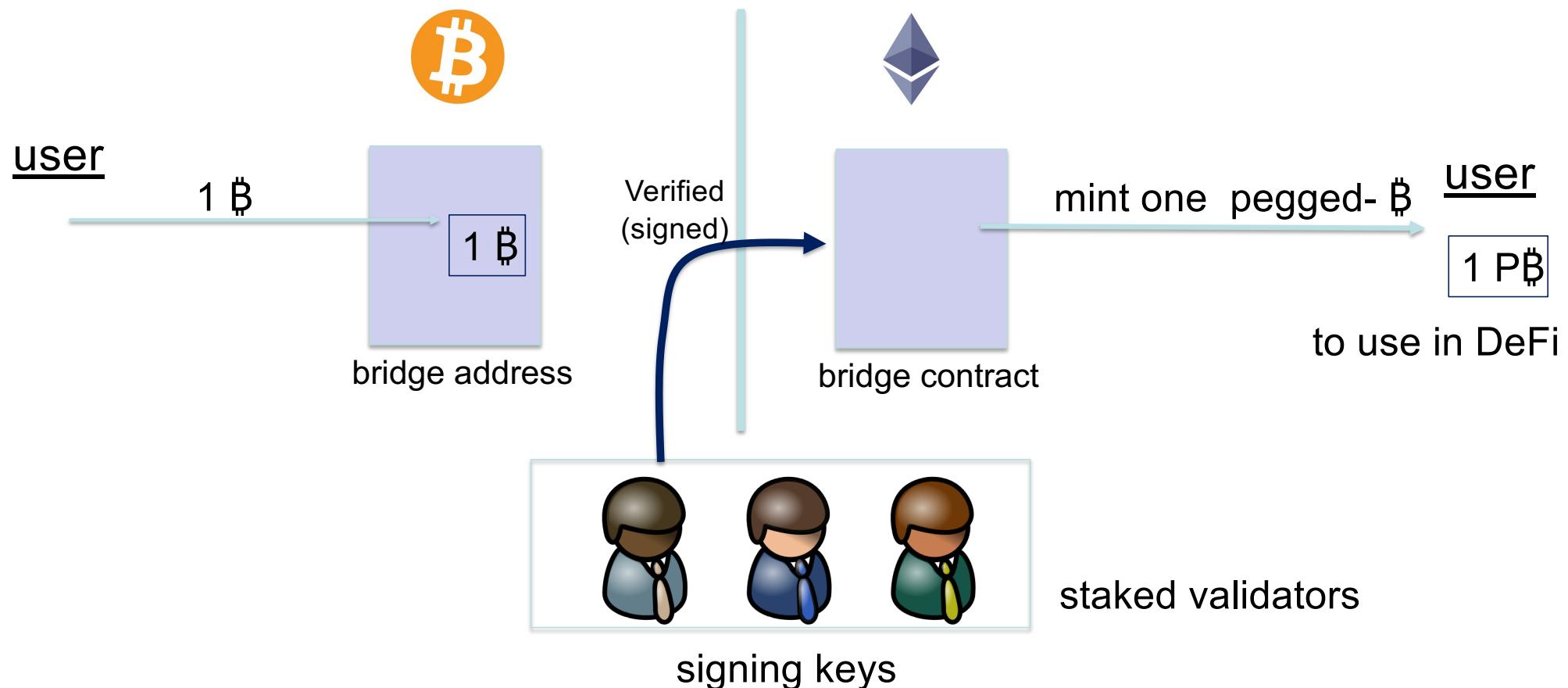
- **Interoperability:**
  - a user owns funds or assets on one blockchain system.  
Goal: enable the user to move funds and/or assets to another system.
- **Composability:**
  - enable a DAPP on one blockchain to call a DAPP on another

Both are easy if the entire world used Ethereum

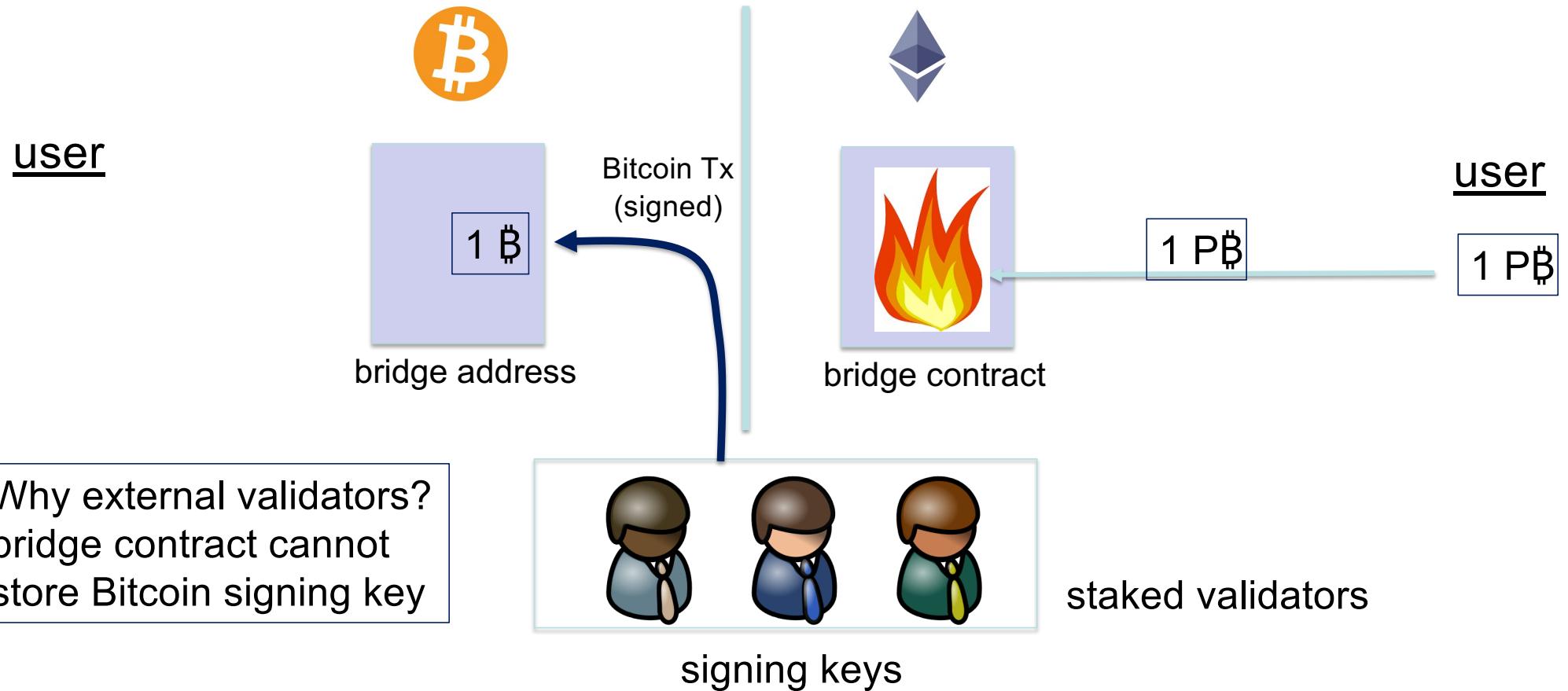
- In reality: many blockchain systems that need to interoperate
- Several cross-chain protocols: ~~XCMP, IBC, ...~~

联邦链桥

## How to move assets? Building a federated bridge (simplified)



# How to move assets? Building a federated bridge (simplified)



# End of lecture: quick review

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Cryptographic primitives:

- Hash functions: committing to large amounts of data
- Digital signatures: authorizing actions

Scaling the blockchain

- Payment channels and Rollups

Interoperability: via bridges and pegged coins.

A dark, grainy image of Earth at night, viewed from space. The planet's curvature is visible against the black void of space. City lights are scattered across continents as glowing yellow and white dots, while the aurora borealis appears as a faint, greenish glow in the upper atmosphere.

END OF TOPIC

<https://defi-learning.org/>