# Exercise 6

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In this text, a class means the collection of all sets that satisfy a given property (first-order formula); V is the universe  $\{x: x=x\}$ ;  $n^+$  means  $n \cup \{n\}$ .

# Problem 1.

*Proof.* (a)  $\forall y (y \notin \emptyset)$ . Thus  $\forall x \forall y ((x \in y \land y \in \emptyset) \rightarrow x \in \emptyset)$  holds.

(b) Suppose n is  $\in$ -transitive. Let x, y be such that  $x \in y$  and  $y \in n \cup \{n\}$ . If  $y \in n$ , then  $x \in n \subseteq n \cup \{n\}$  by the  $\in$ -transitivity of n. If  $y \in \{n\}$ , then  $x \in y$  implies x = n. Thus  $x \in n \cup \{n\}$  trivially. By proof by cases,  $n \cup \{n\}$  is  $\in$ -transitive.

#### Problem 2.

*Proof.* (a)  $0 = \emptyset$  has no nonempty subsets. Thus 0 is  $\in$ -well-ordered by definition.

(b) Suppose n is  $\in$ -well-ordered. Let  $x \subseteq n \cup \{n\}$ . Let  $z = x \setminus \{n\} \subseteq n$ . Let s be an  $\in$ -least element of z. For each  $a \in x$ , if  $a \in z$  then s = a or  $s \in a$  by definition; if a = n then  $s \in z \subseteq n = a$ . Hence s is an  $\in$ -least element of x, completing the proof.

#### Problem 3.

*Proof.* Suppose x, y are such that x, y are inductive;  $x \subseteq z$  for every inductive set z;  $y \subseteq z$  for every inductive set z. Then taking z = y leads to  $x \subseteq y$ , and taking z = x leads to  $y \subseteq x$ . Thus x = y, as desired.

### Problem 4.

*Proof.*  $\emptyset \in u$  and  $\emptyset \in v$  leads to  $\emptyset \in u \cap v$ . For every  $x \in u \cap v$ , we have  $x \in u$  and  $x \in v$ , which implies that  $x \cup \{x\} \in u$  and  $x \cup \{x\} \in v$ , which imply that  $x \cup \{x\} \in u \cap v$ . Thus  $u \cap v$  is inductive.

## Problem 5.

- *Proof.* (a) For each inductive set  $w \subseteq u$ ,  $\emptyset \in w$ . Thus  $\emptyset \in v$ . For each  $x \in v$ ,  $x \in w$  for all inductive set  $w \subseteq u$ , which implies that  $x \cup \{x\} \in w$  for all inductive set  $w \subseteq u$ , which implies that  $x \cup \{x\} \in v$ . Hence v is inductive.
  - (b) For any inductive set  $w \subseteq u$ ,  $v \subseteq w$  by the definition of v. The desired result follows from that v is inductive.

#### Problem 6.

<sup>&</sup>lt;sup>1</sup>Note that we do not need to distinguish whether  $n \in x$  or not.

*Proof.* Suppose there exists an inductive set u. Let<sup>2</sup>

$$N = \bigcap_{X \text{ is inductive}} X.$$

N is indeed a set because  $N \subseteq u$ . We prove that N is the smallest inductive set.

For each inductive set X,  $\emptyset \in X$ . Thus  $\emptyset \in N$ . For each  $x \in N$ ,  $x \in X$  for all inductive set X, which implies that  $x \cup \{x\} \in X$  for all inductive set X, which implies that  $x \cup \{x\} \in N$ . Hence N is inductive. For any inductive set X,  $X \subseteq N$  by definition. Hence N is the smallest inductive set.

#### Problem 7.

*Proof.* (a)  $\mathbb N$  and X are both inductive by definition. Then by Problem 4,  $\mathbb N\cap X$  is inductive.

(b) By definition of  $\mathbb{N}$  (the smallest inductive set),  $\mathbb{N} \subseteq \mathbb{N} \cap X \subseteq \mathbb{N}$ , which implies that  $\mathbb{N} \subseteq X$ , completing the proof.

# Problem 8. <sup>3</sup>

**Lemma 1.** We define that a set T is transitive if  $\forall x (x \in T \to x \subseteq T)$ . Every  $n \in \mathbb{N}$  is transitive.

*Proof.* Let  $E = \{n \in \mathbb{N} : n \text{ is transitive}\}$ . We prove by induction that  $E = \mathbb{N}$ .  $\forall x (x \in 0 \to x \subseteq 0)$  follows from  $0 = \emptyset$ . Assume that  $n \in \mathbb{N}$  is transitive. For each  $x \in n^+ = n \cup \{n\}$ , either  $x \in n$  or x = n. If  $x \in n$  then  $x \subseteq n \subseteq n^+$  by induction hypothesis; if x = n then  $x \subseteq n^+$  trivially. That closes the induction and completes the proof.

**Lemma 2.** Suppose  $n, m \in \mathbb{N}$ . If  $n^+ = m^+$  then n = m.

*Proof.* Suppose  $n \cup \{n\} = m \cup \{m\}$ . Then  $n \in m \cup \{m\}$ . Thus n = m or  $n \in m$ . Similarly m = n or  $m \in n$ . If  $n \neq m$  then  $n \in m$  and  $m \in n$ . By Lemma 1,  $n \subseteq m$  and  $m \subseteq n$ , which implies that n = m.

Proof of the Problem. Fix x. Let R be the property defined by

$$R(A) = ((0, x) \in A \land \forall n \forall y ((n, y) \in A \to (n^+, y \times x) \in A)).$$

Let class  $F = \{(u, v) : P(u, v)\}$ , which is equivalent to

$$F = (\mathbb{N} \times V) \cap \bigcap_{R(A)} A.$$

First we prove that R(F). For every A such that R(A),  $(0, x) \in A$ . Thus  $(0, x) \in F$ . For any  $(n, y) \in F$ ,  $(n, y) \in A$  for each A with property R, which implies that  $(n^+, y \times x) \in A$  for each A with property R, which implies that  $(n^+, y \times x) \in F$ . Hence R(F).

Let

$$X = \{(n, y) \in F : ((n, y) = (0, x) \lor \exists n' \exists y' ((n', y') \in F \land n = n'^+ \land y = y' \times x))\}.$$

<sup>&</sup>lt;sup>2</sup>Note that we have not claimed that  $N = \mathbb{N}$ .

 $<sup>^3</sup>$ Erratum: the seemingly fancy proof below is actually a false one (I applied a property (P in this case) to classes (F in this case)). None have found a proper solution to this problem yet. I have proved the problem with the Recursion Theorem, which is not meaningful.

Now we prove that R(X).  $(0,x) \in X$  trivially. For every  $(n,y) \in X$ ,  $(n,y) \in F$ . Thus  $(n^+, y \times x)$  satisfies the second condition in the definition of X, which implies that  $(n^+, y \times x) \in X$ . Hence R(X).

By definition of  $F, F \subseteq X$ . Let  $E = \{n \in \mathbb{N} : \forall y \forall z (((n,y) \in F \land (n,z) \in F) \to y = z)\}$ . We prove by induction that  $E = \mathbb{N}$ .

First we prove that  $0 \in E$ . Suppose  $(0, y), (0, z) \in F$ . Because  $\forall n \in \mathbb{N}, n^+ \neq \emptyset$ , (0, y), (0, z) must satisfy the first condition in the definition of X, which implies that y = z = x.

Assume  $n \in E$  for some  $n \in \mathbb{N}$ . Suppose  $(n^+, y), (n^+, z) \in F$ . Because  $n^+ \neq \emptyset$ , then  $(n^+, y), (n^+, z)$  must satisfy the second condition in the definition of X. Thus  $\exists (m_1, y'), (m_2, z') \in F$  such that  $n^+ = m_1^+ = m_2^+$  and  $y = y' \times x$  and  $z = z' \times x$ . By Lemma 1,  $m_1 = m_2 = n$ . Furthermore, by induction hypothesis, y' = z', which implies that y = z. That closes the induction, proving that  $E = \mathbb{N}$ , completing the proof.