

I am sorry about my poor English ability.

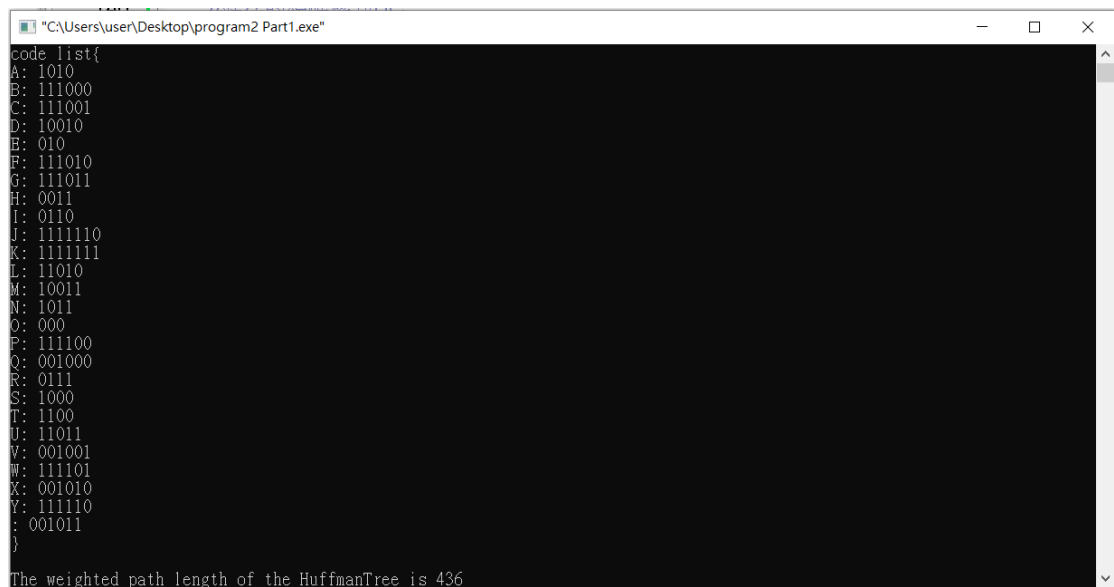
1. 簡單描述何為 huffman encoding 及其用途

Huffman coding is used to code and compress datas, and the followings are the processes of Huffman coding:

- (1) According to the frequency that characters appear, we construct Huffman tree.
- (2) We code the data by the Huffman tree.

Because of using Huffman coding, we can greatly decrease the complexity of the data and increase the efficiency of compressing. Huffman coding is a character coding way which can waste smallest data space.

2. Part 1 的結果



```
"C:\Users\user\Desktop\program2 Part1.exe"
code list{
A: 1010
B: 111000
C: 111001
D: 10010
E: 010
F: 111010
G: 111011
H: 0011
I: 0110
J: 1111110
K: 1111111
L: 11010
M: 10011
N: 1011
O: 000
P: 111100
Q: 001000
R: 0111
S: 1000
T: 1100
U: 11011
V: 001001
W: 111101
X: 001010
Y: 111110
Z: 001011
}
The weighted path length of the HuffmanTree is 436
```

The code list from A to Z is:

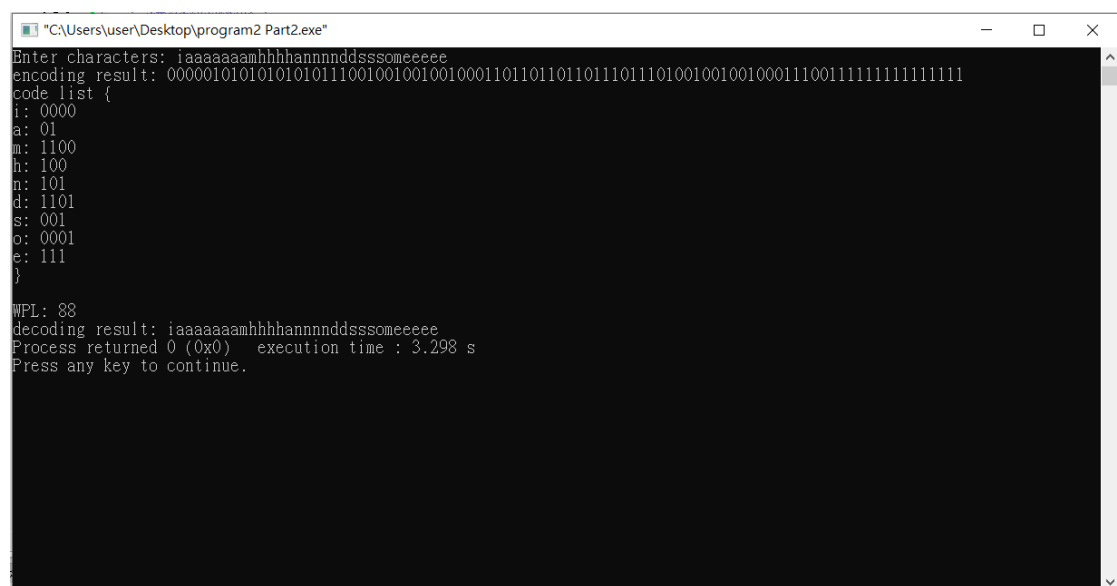
A: 1010, B: 111000, C: 111001, D: 10010, E: 010, F: 111010,

G: 111011, H: 0011, I: 0110, J: 1111110, K: 1111111, L: 11010

M: 10011, N: 1011, O: 000, P: 111100, Q: 001000, R: 0111, S: 1000
T: 1100, U: 11011, V: 001001, W: 111101, X: 001010, Y: 111110
Z: 001011

The WPL is 436.

1. Part 2 的結果



```
"C:\Users\user\Desktop\program2 Part2.exe"
Enter characters: iaaaaaaamhhhhannnnddsssomeeeee
encoding result: 000001010101010101110010010010010001101101101101101100100100100011100111111111111111
code list {
i: 0000
a: 01
m: 1100
h: 100
n: 101
d: 1101
s: 001
o: 0001
e: 111
}
WPL: 88
decoding result: iaaaaaaamhhhhannnnddsssomeeeee
Process returned 0 (0x0)   execution time : 3.298 s
Press any key to continue.
```

The encoding result:

000001010101010101011100100100100100011011011011011101110100100100
10001110011111111111111111

Code list: i: 0000, a: 01, m: 1100, h: 100, n: 101, d: 1101, s: 001

o: 0001, e: 111

WPL:88

Decording result: iaaaaaaamhhhhannnnddsssomeeeee

```
"C:\Users\user\Desktop\program2 Part1.exe"
code list{
A: 1010
B: 111000
C: 111001
D: 10010
E: 010
F: 111010
G: 111011
H: 0011
I: 0110
J: 1111110
K: 1111111
L: 11010
M: 10011
N: 1011
O: 000
P: 111100
Q: 001000
R: 0111
S: 1000
T: 1100
U: 11011
V: 001001
W: 111101
X: 001010
Y: 111110
Z: 001011
}
The weighted path length of the HuffmanTree is 436
```

The encoding result: 011010111111111011100000011001101101010010

Code list:h: 0110,o: 10,w: 111,a: 0111,r: 00,e: 1100,y: 1101

u: 010

WPL:42

Decording result: howwwarrreyoooou

4. 如何使用程式碼實作 huffman encoding 並得到 Part 1 和 Part 2 的結果

Part1: (1)define a structure including weight, left and right

child node

```
typedef struct huffNode {
    double weight; /*權重*/
    int lchild, rchild, parent; /*左右子節點和父節點*/
} HTNode, * HuffTree;
```

(2)Input two tables:one for characters A to Z,one for the

frequwncies

```
N = 26; //共26字母
//第0個保留不用
ElemType data[N] = {"", "A", "B", "C", "D", "E", "F", "G", "H", "I", "J", "K", "L", "M", "N", "O", "P", "Q", "R", "S", "T", "U", "V", "W", "X", "Y", "Z"};
//第0個保留不用
double w[N] = {0.7, 2.2, 3.11, 2.2, 6.6, 1.1, 4.3, 7.9, 2.1, 6.6, 8.4, 1.2, 1.2, 1};
```

(3) Create the huffmantree:

```
void createHT(HuffTree& HT, HuffCode& HC, double* w, int n) {
    int s1, s2, m = 2 * n - 1; //m: 由n個節點組成的赫夫曼會有2n-1個節點
    char* code; //暫存
    HT = new HTNode[m + 1]; //第0個不使用

    for (int i = 1; i <= n; i++) {
        //處理初始化前的第0個節點
        HT[i] = { w[i], 0, 0, 0 };
    }

    for (int i = n + 1; i <= m; i++) {
        //處理初始化後n-1個節點 (找出最小兩節點的父節點)
        HT[i] = { 0, 0, 0, 0 };
    }

    //赫夫曼樹建構
    for (int i = n + 1; i <= m; i++) {
        //找出前i-1個節點中權值最小的節點
        select(HT, i - 1, s1, s2);
        HT[s1].parent = i;
        HT[s2].parent = i;
        HT[i].lchild = s1;
        HT[i].rchild = s2;
        HT[i].weight = HT[s1].weight + HT[s2].weight;
    }

    HC = new char* [n];
    /*這裡以下我真的不知道自己在幹嘛*/
    code = new char[n];

    for (int i = 1; i <= n; i++) {
        //k: 現在的節點, 用0和1表示, f: k的父節點, j: 記錄編碼的位置
        int k = i, f = HT[k].parent, j = 0;
        //從葉子到根走一遍
        while (f != 0) {
            if (HT[f].lchild == k) {
                code[j] = '0';
            }
            else if (HT[f].rchild == k) {
                code[j] = '1';
            }
            k = HT[k].parent;
            f = HT[k].parent;
            j++;
        }
        //標記尾巴位置
        code[j] = '\0';
        reverseChars(code, j);
        //站存的編碼移到HC
        HC[i] = new char[n];
        strcpy(HC[i], code);
    }
}
```

(4) Show the Huffman code

(5) Use DFS to calculate WPL

```

int getWPL(HuffTree& HT, int idx, int depth) {
    //執行dfs直到遇到葉子
    if (HT[idx].lchild == 0 && HT[idx].rchild == 0) {
        return HT[idx].weight * depth;
    }
    return getWPL(HT, HT[idx].lchild, depth + 1) + getWPL(HT, HT[idx].rchild, depth + 1);
}

```

(6)Show WPL

Part2: (1)Initialize the frequency=0

(2)Find the frequency of every characters and array them
from low to high

```

void frequent(string str)
{
    int length = str.length(); /*長度*/
    minnode* node = new minnode[length]; /*宣告最0節點*/
    int i, j;
    for (i = 0; i < length; i++) /*初始化頻度*/
    {
        node[i].ch_num = 0;
    }
    int char_type_num = 0; /*初始化為0種字元*/
    for (i = 0; i < length; i++)
    {
        for (j = 0; j < char_type_num; j++)
        {
            if (str[i] == node[j].ch || (node[j].ch >= 'a' && node[j].ch <= 'z' && str[i] + 32 == node[j].ch))
            {
                break;
            }
        }
        if (j < char_type_num)
        {
            node[j].ch_num++;
        }
        else
        {
            if (str[i] >= 'A' && str[i] <= 'Z')
            {
                node[j].ch = str[i] + 32;
            }
        }
    }

    /*按照頻度從小到大排列*/
    for (i = 0; i < char_type_num; i++)
    {
        for (j = i; j < char_type_num; j++)
        {
            if (node[j].ch_num < node[j + 1].ch_num) /*如果前一個小於後面一個 兩者交換*/
            {
                int temp;
                char ch_temp;
                temp = node[j].ch_num;
                ch_temp = node[j].ch;
                node[j].ch_num = node[j + 1].ch;
                node[j].ch = node[j + 1].ch;
                node[j].ch_num = temp;
                node[j].ch = ch_temp;
            }
        }
    }
}

```

(3) Initialize the nodes

```
huffmanTree* huff = new huffmanTree[2 * char_type_num - 1]; /*位於確定char_type_num*/
huffmanTree temp;
string* code = new string[2 * char_type_num - 1];
for (i = 0; i < 2 * char_type_num - 1; i++) /*節點初始化*/
{
    huff[i].parent = -1;
    huff[i].lchild = -1;
    huff[i].rchild = -1;
    huff[i].flag = -1;
}
for (j = 0; j < char_type_num; j++) /*將排序後的第0個節點權重值賦予樹節點*/
{
    huff[j].weight = node[j].ch_num;
}
int min1, min2;
for (int k = char_type_num; k < 2 * char_type_num - 1; k++) /*賦予0級以上的節點值*/
{
    coding(length, huff, k, min1, min2);
    huff[min1].parent = k;
    huff[min2].parent = k;
    huff[min1].flag = "0";
    huff[min2].flag = "1";
    huff[k].lchild = min1;
    huff[k].rchild = min2;
    huff[k].weight = huff[min1].weight + huff[min2].weight;
}
```

(4) Give every nodes values

```
for (i = 0; i < char_type_num; i++) {
    temp = huff[i];
    while (1) {
        code[i] = temp.flag + code[i];
        temp = huff[temp.parent];
        if (temp.parent == -1) break;
    }
}
```

(5) Use DFS to calculate WPL

(6) Show all results

```
int main(void)
{
    int length = 0; /*字串長度*/
    string str; /*目標字串*/
    cout << "Enter characters: ";
    cin >> str;
    frequent(str); /*求各個字串頻度*/
    cout << endl;
    cout << "decoding result: "<< str;
    return 0;
}
```