# CH-231-A Algorithms and Data Structures ADS

Lecture 15

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#### Heap as a Data Structure

- ► Heaps are a data structure that can be used for other purposes, as well.
- ► In particular, a max-heap is often used to build a max-priority queue.

#### Max-Priority Queues

#### Definition (priority queue):

► A priority queue is a data structure for maintaining a set *S* of elements, each with an associated value called a key.

#### Definition (max-priority queue):

- A max-priority queue is a priority queue that supports the following operations:
  - ightharpoonup Maximum(S): return element from S with largest key.
  - Extract-Max(S): remove and return element from S with largest key.
  - ▶ Increase-Key(S, x, k): increase the value of the key of element x to k, where k is assumed to be larger or equal than the current key.
  - ▶ Insert(S, x): add element x to set S.

### Maximum(S)

HEAP-MAXIMUM(A)
1 return A[1]

Time complexity: O(1)

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#### Extract-Max(S)

```
HEAP-EXTRACT-MAX(A)

1 if A.heap-size < 1

2 error "heap underflow"

3 max = A[1]

4 A[1] = A[A.heap-size]

5 A.heap-size = A.heap-size - 1

6 MAX-HEAPIFY(A, 1)

7 return max
```

Time complexity:  $O(\lg n)$ 

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#### Increase-Key(S, x, k)

```
HEAP-INCREASE-KEY (A, i, key)

1 if key < A[i]

2 error "new key is smaller than current key"

3 A[i] = key

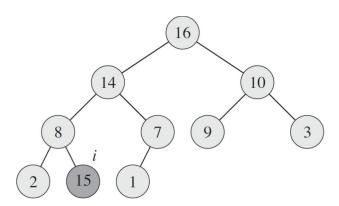
4 while i > 1 and A[PARENT(i)] < A[i]

5 exchange A[i] with A[PARENT(i)]

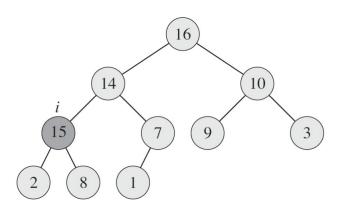
6 i = PARENT(i)
```

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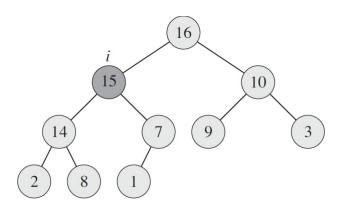
## Increase-Key(S, x, k): Example (1)



# Increase-Key(S, x, k): Example (2)



## Increase-Key(S, x, k): Example (3)



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#### Increase-Key(S, x, k)

```
HEAP-INCREASE-KEY (A, i, key)

1 if key < A[i]

2 error "new key is smaller than current key"

3 A[i] = key

4 while i > 1 and A[PARENT(i)] < A[i]

5 exchange A[i] with A[PARENT(i)]

6 i = PARENT(i)
```

Time complexity:  $O(\lg n)$ 

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### Insert(S, x)

```
Max-Heap-Insert (A, key)
```

- 1 A.heap-size = A.heap-size + 1
- $2 \quad A[A.heap-size] = -\infty$
- 3 HEAP-INCREASE-KEY (A, A. heap-size, key)

Time complexity:  $O(\lg n)$ 

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