CH-231-A

Algorithms and Data Structures

ADS

Lecture 24

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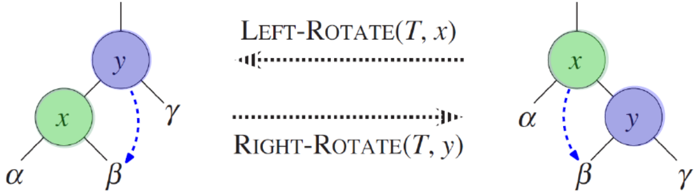
Rotations (1)

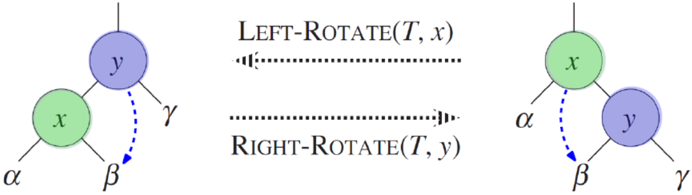
- *Right-Rotate*(T, y):
  - node y becomes right child of its left child x.
  - new left child of y is former right child of x.
- *Left-Rotate*(T, x):
  - node x becomes left child of its right child y.
  - new right child of x is former left child of y.



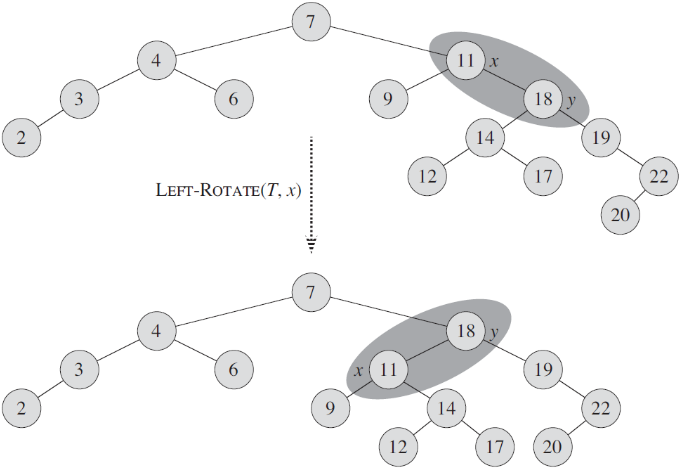
Rotations (2)

BST property is preserved:

- (left): $\mathit{key}(\alpha) \leq x.\mathit{key} \leq \mathit{key}(\beta) \leq y.\mathit{key} \leq \mathit{key}(\gamma)$
- (right): $\mathit{key}(\alpha) \leq x.\mathit{key} \leq \mathit{key}(\beta) \leq y.\mathit{key} \leq \mathit{key}(\gamma)$



Rotation: Example



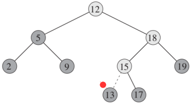
Rotation Pseudocode

```
LEFT-ROTATE(T, x)
  y = x.right                 // set y
  x.right = y.left            // turn y's left subtree into x's right subtree
  if y.left ≠ T.nil
      y.left.p = x
  y.p = x.p                   // link x's parent to y
  if x.p == T.nil
      T.root = y
  elseif x == x.p.left
      x.p.left = y
  else x.p.right = y
  y.left = x                  // put x on y's left
  x.p = y
```

Time complexity: $O(1)$

Insertion

```
TREE-INSERT(T, z)
y = NIL
x = T.root
while x ≠ NIL
    y = x
    if z.key < x.key
        x = x.left
    else x = x.right
z.p = y
if y == NIL
    T.root = z
elseif z.key < y.key
    y.left = z
else y.right = z
```




```
RB-INSERT(T, z)
y = T.nil
x = T.root
while x ≠ T.nil
    y = x
    if z.key < x.key
        x = x.left
    else x = x.right
z.p = y
if y == T.nil
    T.root = z
elseif z.key < y.key
    y.left = z
else y.right = z
z.left = T.nil
z.right = T.nil
z.color = RED
RB-INSERT-FIXUP(T, z)
```

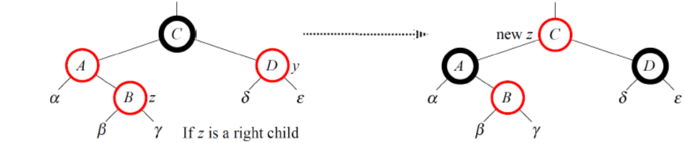
Fixing Red-Black Tree Properties

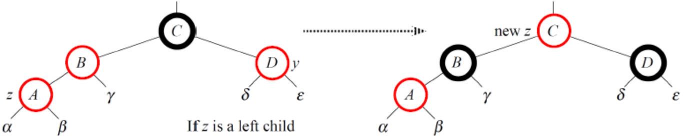
- We are inserting a **red** node to a valid red-black tree.
- Which properties may be violated?
  1. Duh: Cannot be violated. ✓
  2. RooB: Violated if inserted node is root. ✗
  3. LeaB: Inserted node is not a leaf, i.e., no violation. ✓
  4. BredB: Violated if parent of inserted node is red. ✗
  5. BH: Not affected by red nodes, i.e., no violation. ✓

Fixing BredB

- BredB for node z is violated, if $z.p$ is red.
- Then, $z.p.p$ is black. (BredB property)
- We need to consider different cases depending on the uncle y of z, i.e., the child of $z.p.p$ that is not $z.p$.
- There are 6 cases:
  - $z.p$ is left child of $z.p.p$
    - y is red (Case 1)
    - y is black
      - z is right child of $z.p$ (Case 2)
      - z is left child of $z.p$ (Case 3)
  - $z.p$ is right child of $z.p.p$
    - y is red (symmetric to Case 1)
    - y is black
      - z is right child of $z.p$ (symmetric to Case 3)
      - z is left child of $z.p$ (symmetric to Case 2)

Case 1 (Red Uncle)

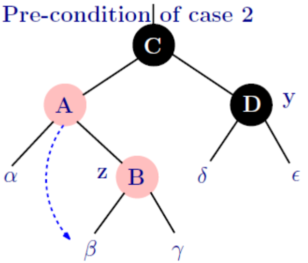


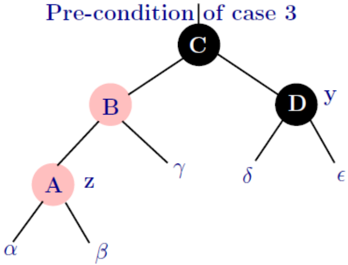



```
if z.p == z.p.p.left
    y = z.p.p.right
    if y.color == RED
        z.p.color = BLACK          Case 1
        y.color = BLACK            Case 1
        z.p.p.color = RED          Case 1
        z = z.p.p                  Case 1
```

There is a new problem, if *z.p.p.p* is red.
Algorithm needs to continue with *z.p.p*.

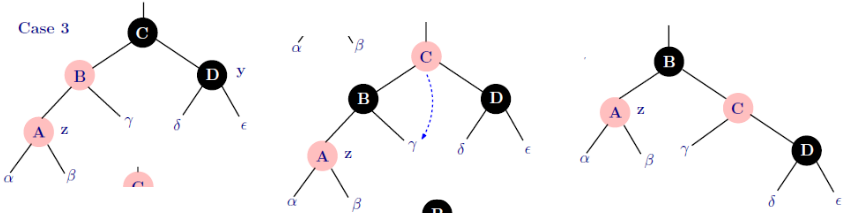
Case 2 (Black Uncle, z Right Child)






```
else if z == z.p.right
    z = z.p                  Case 2
    LEFT-ROTATE(T, z)
```

Case 3 (Black Uncle, z Left Child)



$z.p.color = \text{BLACK}$
$z.p.p.color = \text{RED}$
$\text{RIGHT-ROTATE}(T, z.p.p)$

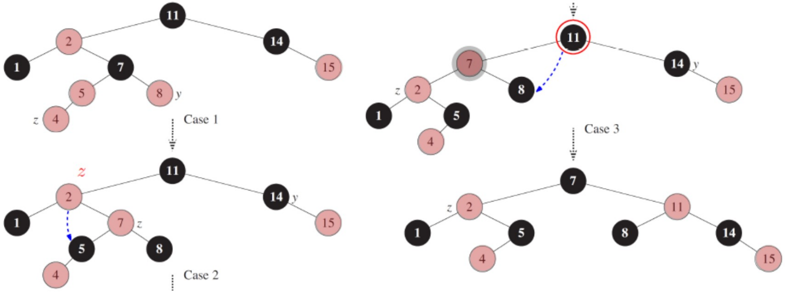
Case 3

Putting It All Together

- We need to put the 3 cases (and the 3 symmetric cases) together.
- Moreover, we need to propagate the considerations upwards (see Case 1).
- Finally, we have to fix RooB.

```
RB-INSERT-FIXUP(T, z)
while z.p.color == RED
    if z.p == z.p.p.left
        y = z.p.p.right
        if y.color == RED
            z.p.color = BLACK
            y.color = BLACK
            z.p.p.color = RED               Case 1
            z = z.p.p
        else if z == z.p.right
                z = z.p                     Case 2
                LEFT-ROTATE(T, z)
            z.p.color = BLACK
            z.p.p.color = RED               Case 3
            RIGHT-ROTATE(T, z.p.p)
    else (same as then clause
            with "right" and "left" exchanged)
T.root.color = BLACK
```

Insert Example



Time Complexity

- In worst case, we have to go all the way from the leaf to the root along the longest path within the tree.
- Hence, running time is $O(h) = O(\lg n)$ for the fixing of the red-black tree properties.
- Overall, running time for insertion is $O(h) = O(\lg n)$.
- Example for building up a red-black tree by iterated node insertion:
  `http://www.youtube.com/watch?v=vDHFF4wjWYU`