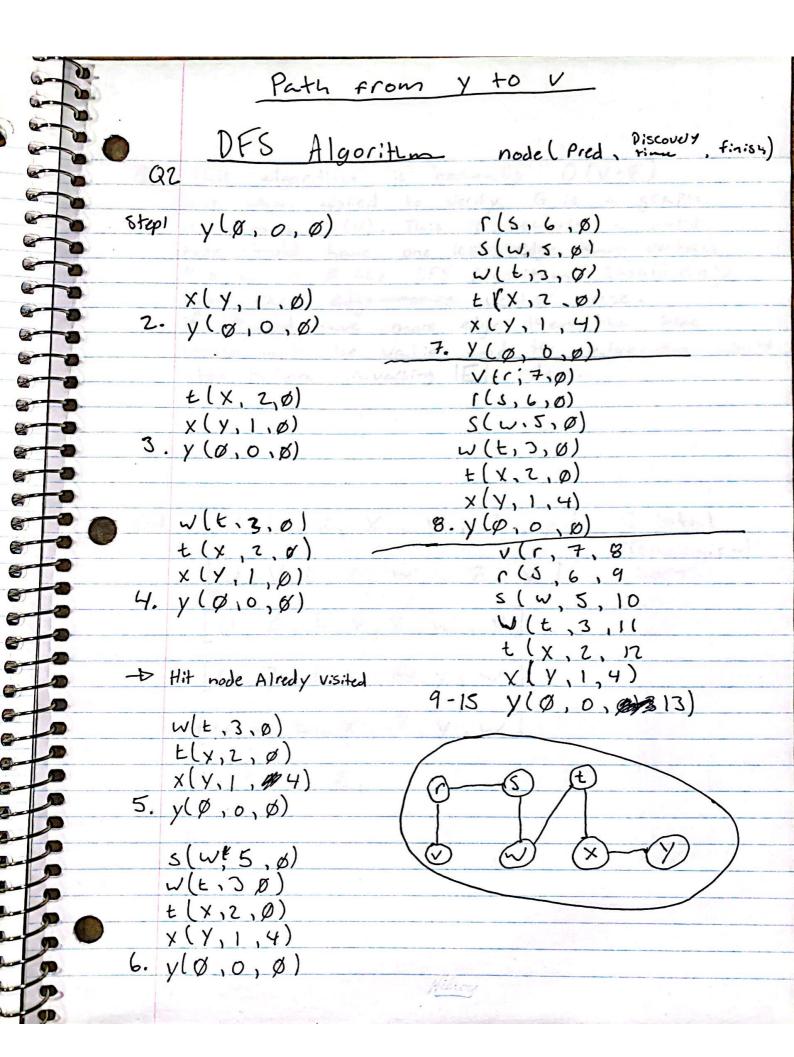
	CP312	Assignment # 5
		Mannet Busilousle
Q1	BFS Algorithm	130951550
	711901710017	(18-6-6)
	example: y (ø,	o) node (Pred, Distance)
		0 (6,5,0%
Chal	Queue y (ø, o)	s (w, 3)
Step 1	y (ø, o)	Step 6 W(x, 2)
		t(x,z)
		U(y,')
	X (Y, 1)	X(y,1)
2.	y (ø, o)	(y(Ø,0)
		$\omega(\mathbb{C}, \mathfrak{I}, \mathfrak{I})$
		1(5,4)
	U(y,1)	S(w,3)
	X (Y , 1)	ω(x, z)
3.	$y(\emptyset, 0)$	t(x, z)
		u(y,i)
	(() (() () () () () () () ()	$\chi(y,1)$
	t(x,2)	\mathcal{F}_{\cdot} $y(\emptyset,0)$
	u (y , 1)	t (x, z, n
	$\chi(\gamma, 1)$	VLY, 1, 4)
4.	$y(\emptyset, 0)$	V(r, 5)
	1.0	((5,4)
		S(W,3)
les .	$\omega(x,\overline{z})$	$\omega(x,z)$
	t (x,2)	t (x, 2)
¥	U (Y, 1)	4(411)
	x(y,1)	x(y,1)
5.	y (\$, 0)	8. y (ø, o)
		Done
		History

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(25. assume For contridiction that all vertices $V \in V$ indegree = 71.

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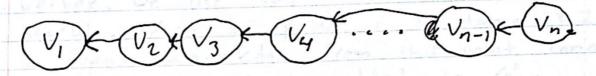
- C

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if all vertices V, to Un mare connected by out degree there must be an edge connecting every vertex by in degree as well



However the only way to connect all points in these two ways of creates a cycle for a finite graph. Since a DAG must connect all points with a degree and connecting all points with attindegree creates a cycle this contradicts the definition of a DAG.

As well & Dag's must be topologically Searchable and without a Starting point Violate the definition of a Dag