

ICT1008 Data Structures and Algorithms

Lecture 5: Sorting Algorithms

Agenda



- Bubble Sort
- Selection Sort
- Insertion Sort
- Merge Sort
- Quick Sort

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Recommended Readings



- 1. Runestone Interactive book: "Problem Solving with Algorithms and Data Structures Using Python"
 - Section: "Sorting & Searching"
- 2. Algorithms by Robert Sedgewick and Kevin Wayne. Addison-Wesley Professional. 4th edition, 2011.
 - Chapter 2: "Sorting"

Bubble Sort



- 1. Step through the list to be sorted.
- 2. Compare two adjacent items at a time until you reach the end of the list; swap the two adjacent items if they are in the wrong order.
- 3. Repeat from the beginning of the list until no swaps are needed.

Bubble Sort Algorithm



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Bubble Sort Analysis



- Regardless of how the items are arranged in the initial list,
 (n−1) passes will be made to sort a list of size n.
- Total number of comparisons

= 1 + 2 + 3 + ... + (n-1) =
$$\frac{1}{2}n^2 - \frac{1}{2}n$$
.

- Complexity is $O(n^2)$ comparisons.
- In the best case, if the list is already ordered, no exchanges will be made.
- In the worst case, every comparison will cause an exchange.
- On average, bubble sort exchanges half of the time.

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Selection Sort



The array is divided into two parts: sorted & unsorted.

Initially, the sorted part is empty.

- 1. Find the minimum value in the list.
- 2. Swap it with the value in the first position of the array. The sorted part grows from here.
- Find the next minimum value in the list.
- 4. Swap it with the value in the (second) next position of the array.
- Repeat steps 3 & 4 till the end of the list.

https://youtu.be/Ns4TPTC8whw

Selection Sort Algorithm



```
1 def selectionSort(alist):
     for fillslot in range(len(alist)-1,0,-1):
         positionOfMax=0
         for location in range (1, fillslot+1):
              if alist[location] > alist[positionOfMax]:
                  positionOfMax = location
         temp = alist[fillslot]
          alist[fillslot] = alist[positionOfMax]
         alist[positionOfMax] = temp
12 \text{ alist} = [54, 26, 93, 17, 77, 31, 44, 55, 20]
13 selectionSort(alist)
14 print (alist)
```

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Selection Sort Analysis

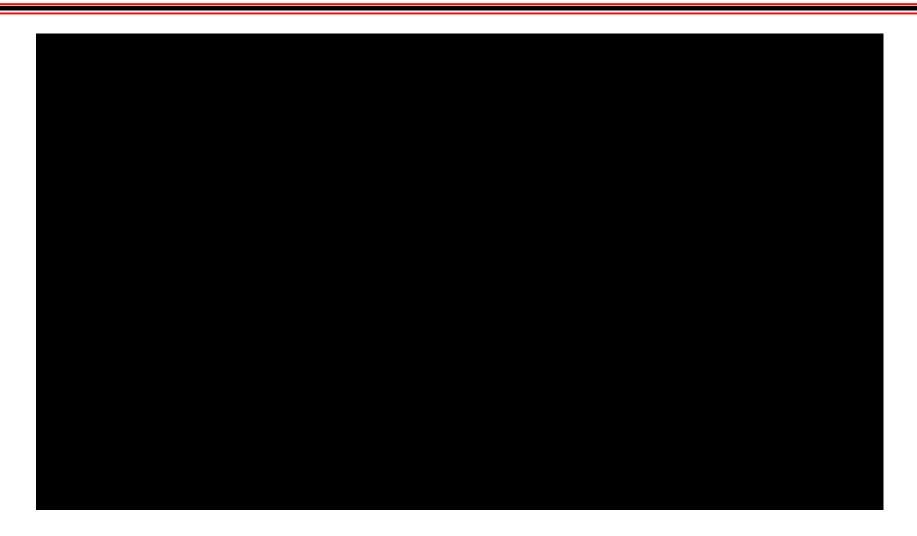


- Selection sort makes the same number of comparisons as bubble sort.
- Complexity is therefore also $O(n^2)$.
- However, due to the reduction in the number of exchanges, selection sort typically executes faster than bubble sort in benchmark studies.

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Selection Sort Demo





https://youtu.be/Ns4TPTC8whw

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Insertion Sort



The array is divided into two parts: sorted & unsorted.

Initially, the sorted part is empty.

- 1. Every iteration of insertion sort removes an element (normally the first one) from the input data, and inserts it into the correct position in the already-sorted list.
- 2. Repeat step 1 until no input element remains.

https://youtu.be/ROalU379l3U

Insertion Sort Algorithm



```
1 def insertionSort(alist):
     for index in range(1,len(alist)):
       currentvalue = alist[index]
       position = index
       while position>0 and alist[position-1]>currentvalue:
           alist[position] = alist[position-1]
           position = position-1
       alist[position]=currentvalue
12
13 alist = [54,26,93,17,77,31,44,55,20]
14 insertionSort (alist)
15 print (alist)
```

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Insertion Sort Analysis



- Insertion Sort uses (n-1) passes to sort n items.
- Maximum number of comparisons is O(n²) complexity:

= 1 + 2 + 3 + ... + (n-1) =
$$\frac{1}{2}n^2 - \frac{1}{2}n$$
.

 Best case: only one comparison needs to be done on each pass for an already sorted list.

Insertion Sort Demo





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Merge Sort

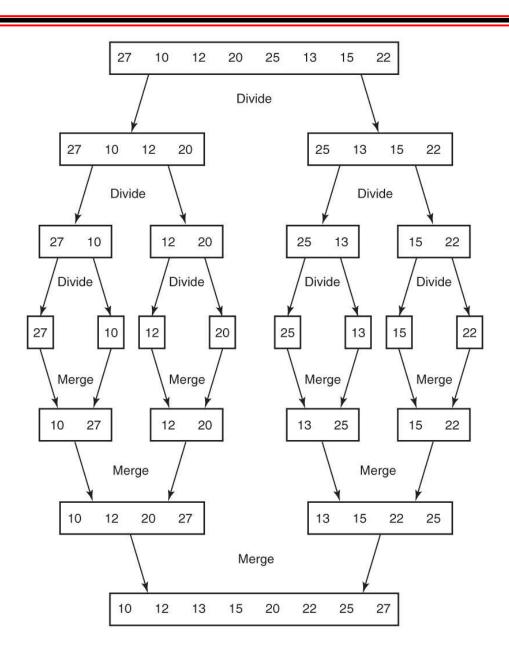


Method: Divide & conquer

- 1. Divide the unsorted list into two nearly equal size sub-lists.
- 2. Sort each sub-list recursively by applying merge sort.
- 3. Merge the two sub-lists back into one sorted list.

Merge Sort: Example Lecture 5 slide 9





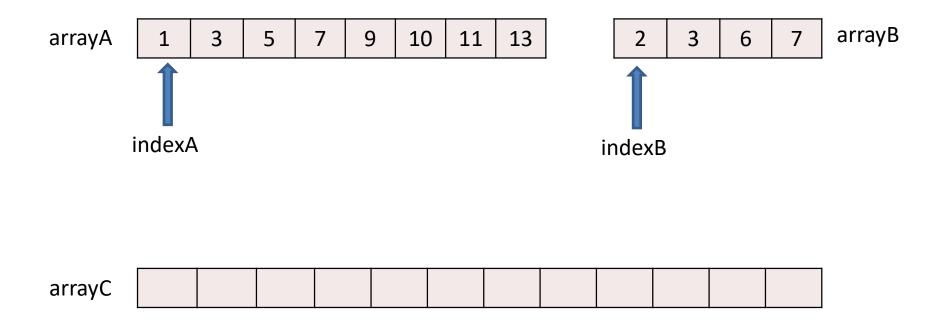
Merge algorithm



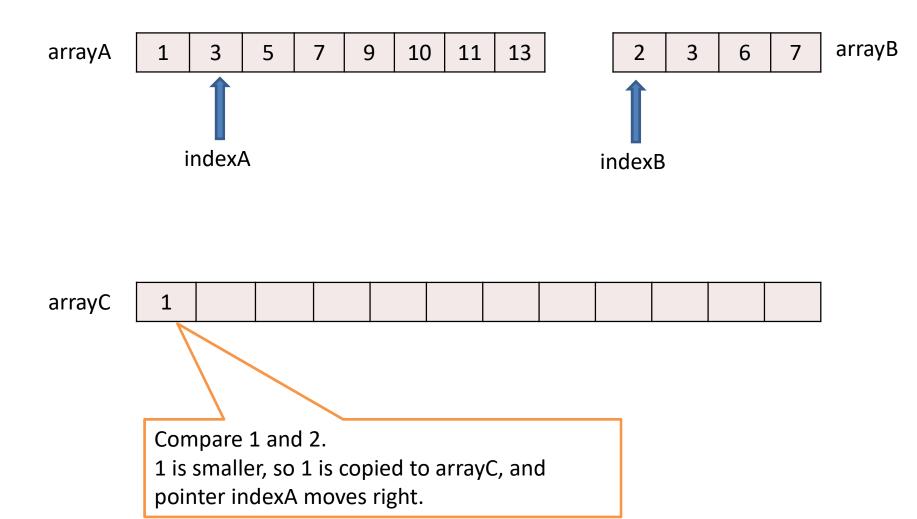
```
def merge(arrayA, arrayB):
     arrayC = []
     sizeA = len(arrayA)
     sizeB = len(arrayB)
     indexA = indexB = 0
     while indexA < sizeA and indexB < sizeB:
          if arrayA[indexA] < arrayB[indexB]:</pre>
              arrayC.append(arrayA[indexA])
              indexA = indexA+1
          else:
              arrayC.append(arrayB[indexB])
              indexB = indexB+1
     for i in range (indexA, sizeA):
         arrayC.append(arrayA[i])
     for i in range (indexB, sizeB):
          arrayC.append(arrayB[i])
     return arrayC
```

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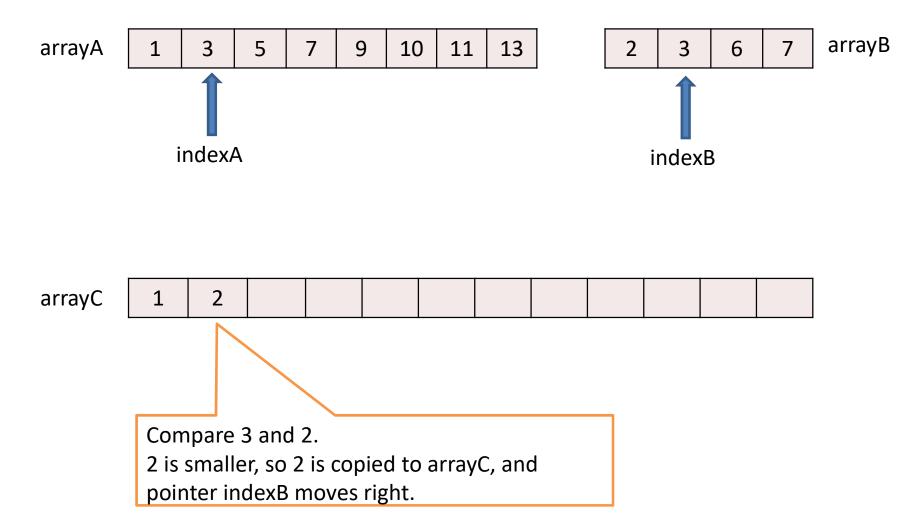




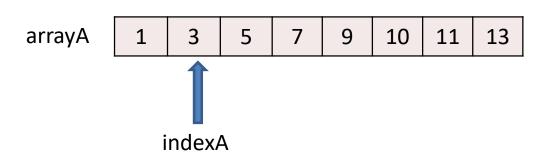


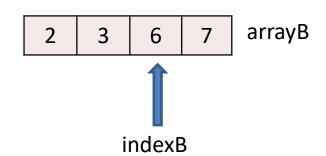


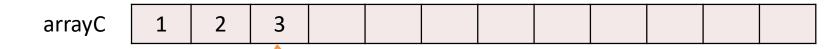








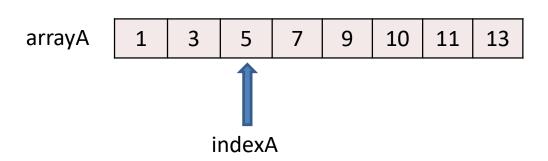


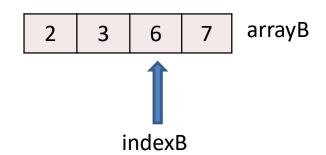


Compare 3 and 3.

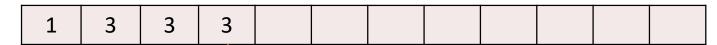
3 in arrayA is not smaller than 3 in arrayB, so the 3 in arrayB is copied to arrayC, and pointer indexB moves right.







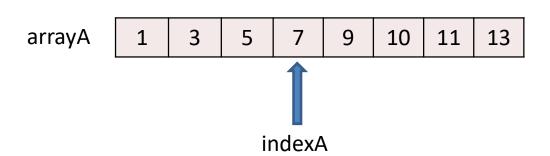


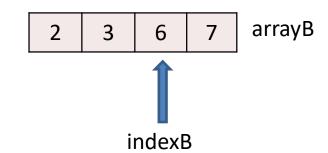


Compare 3 and 6.

3 is smaller than 6, so the 3 is copied to arrayC, and pointer indexA moves right.







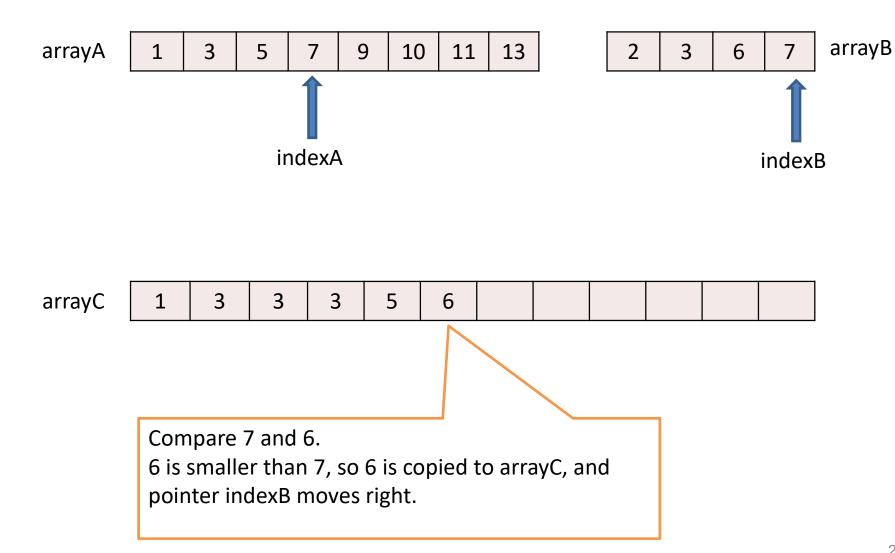




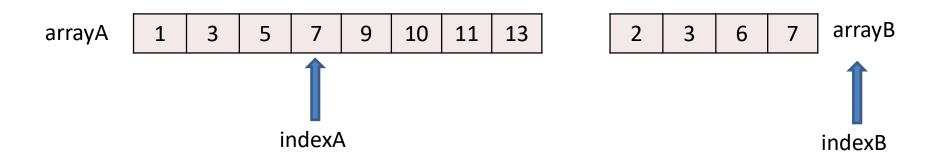
Compare 5 and 6.

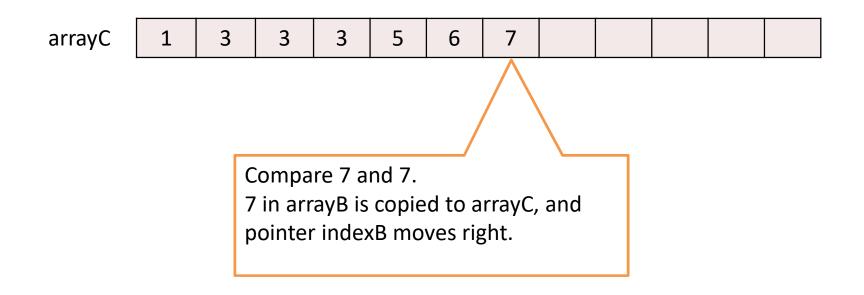
5 is smaller than 6, so 5 is copied to arrayC, and pointer indexA moves right.



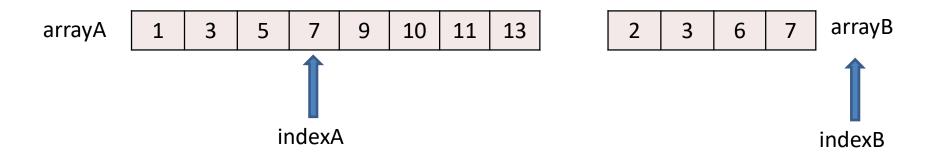














The remaining elements of arrayA are copied into arrayC.

Merge Sort algorithm



```
Base case:
def mergeSort(array):
                                 When the array has one element,
     size = len(array)
                                 it is already sorted.
                                 So return the array for merging.
      if size is 1:
          return array
                                              Divide the array into two
                                              almost equal halves:
     midIndex = size/2
                                              firstHalf and
     firstHalf = array[0:midIndex]
                                              secondHalf.
     secondHalf = array[midIndex:size]
                                               Recursively MergeSort
     firstHalf = mergeSort(firstHalf)
                                               divide firstHalf and
     secondHalf = mergeSort(secondHalf)
                                               divide secondHalf.
     array = merge(firstHalf, secondHalf)
      return array
                                                 Merge the sorted
                                                 firstHalf and
 print mergeSort([27,10, 20,25,13,15,22])
                                                 secondHalf.
```

Return the sorted array.

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Merge Sort – Complexity



- Time complexity
 - -merge is O(n).
 - -merge is called $O(\log n)$ times recursively.
 - -mergeSort is $O(n \log n)$.
- Space complexity
 - merge uses an additional arrayC.
 - If arrayC was local inside merge, much more storage would be used because of recursive calls.
 - Consider using a global arrayC
 in the implementation.

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Quick Sort



Method: Divide-and-conquer.

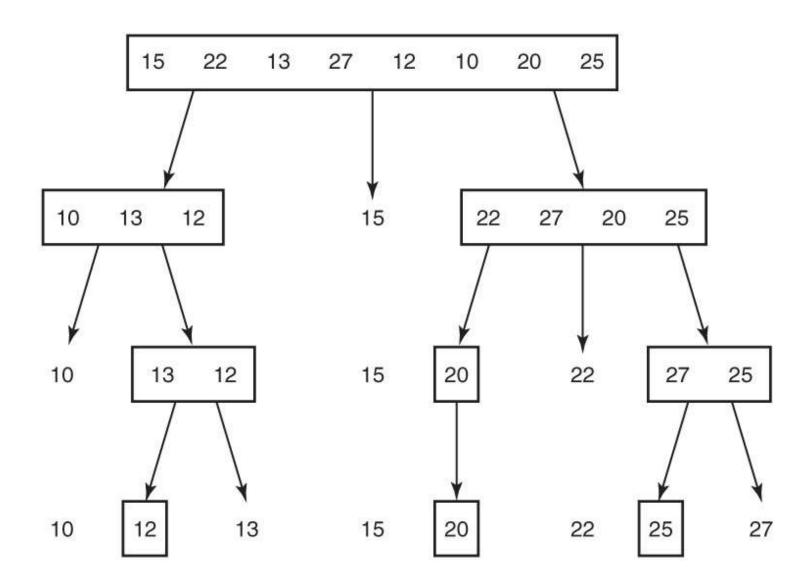
- Pick an element (pivot) from the list.
 - *pivot* is arbitrarily chosen.
 - Normally, the first element is selected.
- Partition the list into two halves such that:
 - All the elements in the first half are smaller than the pivot.
 - All the elements in the second half are greater than or equal to the pivot.

1st half pivot 2nd half

- Quick-sort the 1st half.
- Quick-sort the 2nd half.

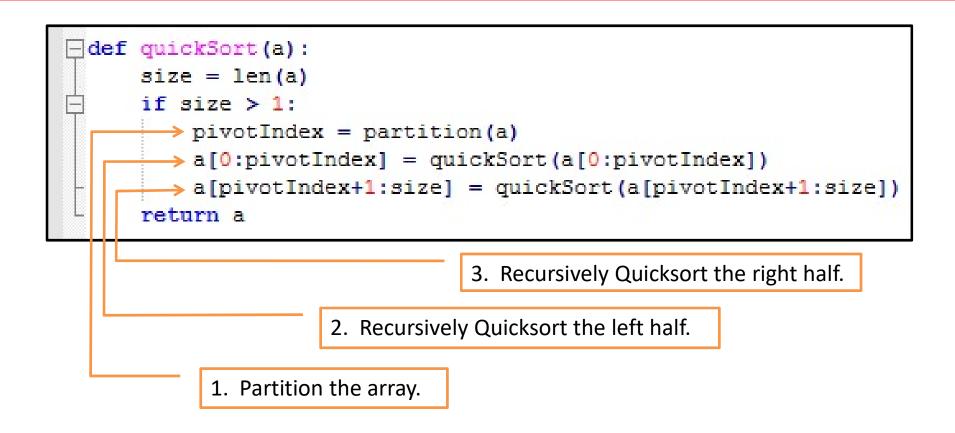
Quick Sort: Example Lecture 5 slide 10





Quick Sort algorithm





Partition Algorithm



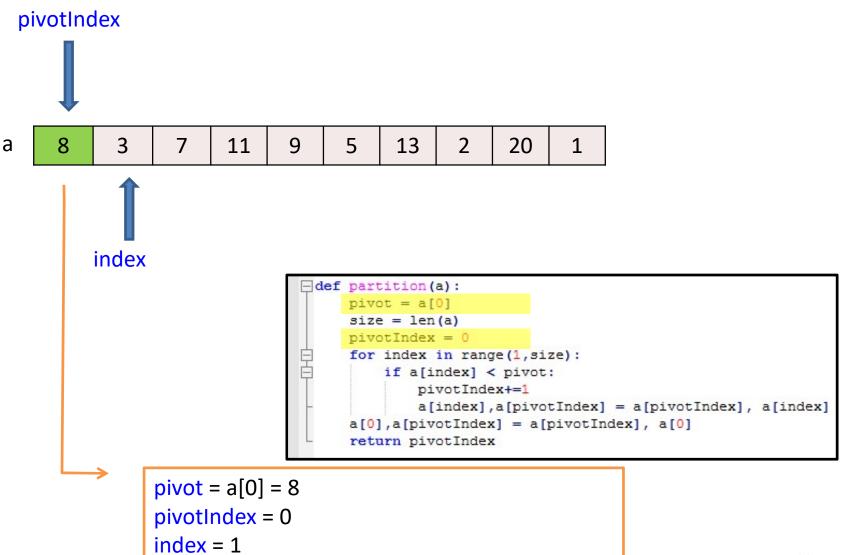
- 3. Move the pivot into place.
- 2. Move element into left half.
- 1. Element at index is smaller than pivot
- => Belongs to left half. Increment pivotIndex to increase size of left half by 1.

Partition Algorithm

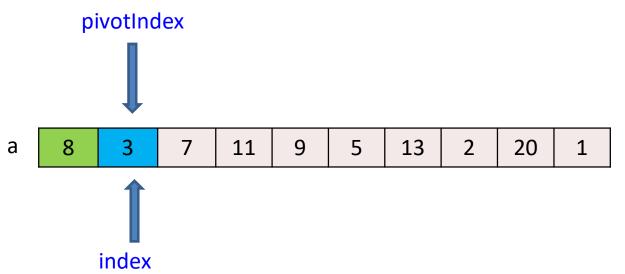


```
def quickSort(a):
    size = len(a)
    if size > 1:
        pivotIndex = partition(a)
        a[0:pivotIndex] = quickSort(a[0:pivotIndex])
        a[pivotIndex+1:size] = quickSort(a[pivotIndex+1:size])
    return a
```



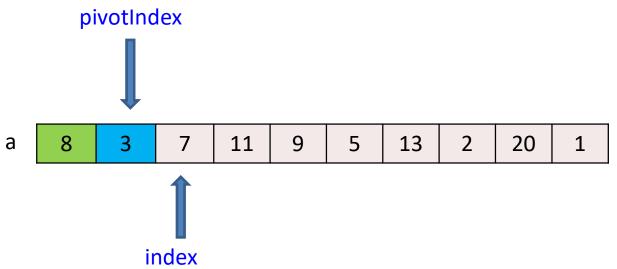






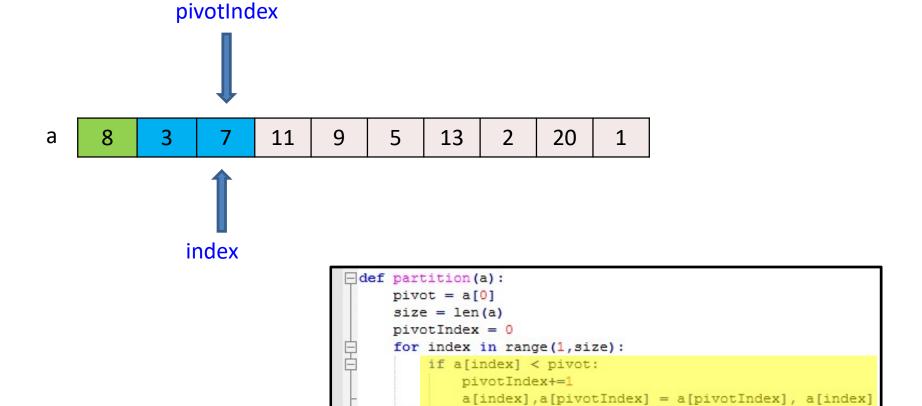
```
3 < 8.</li>pivotIndex shift right.Swap pivotIndex with index => no change to array.
```





index shift right.





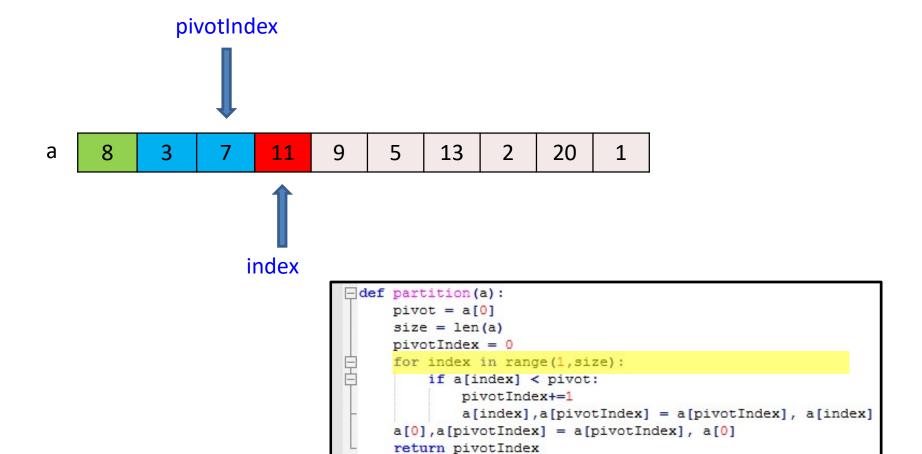
7 < 8.
pivotIndex shift right.

return pivotIndex

a[0],a[pivotIndex] = a[pivotIndex], a[0]

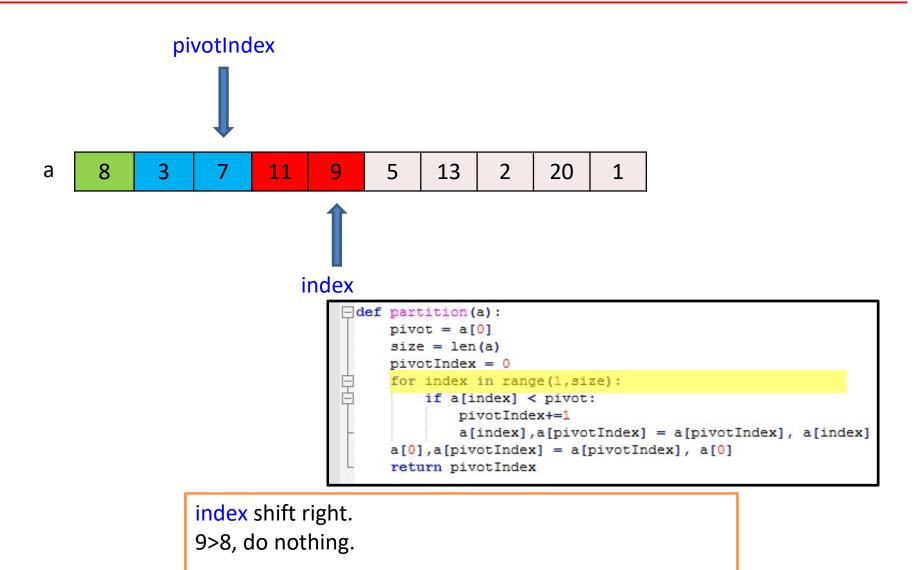
Swap pivotIndex with index => no change to array.



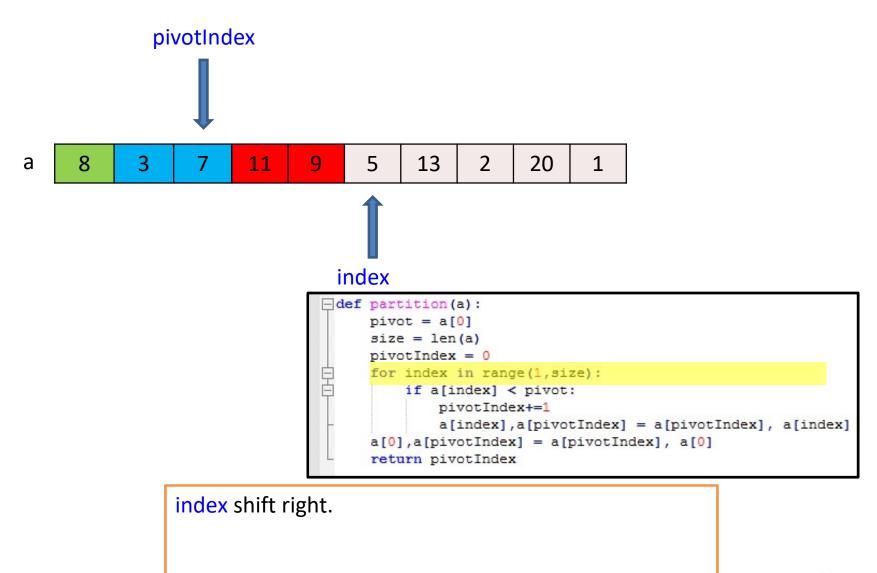


index shift right.11>8, do nothing.

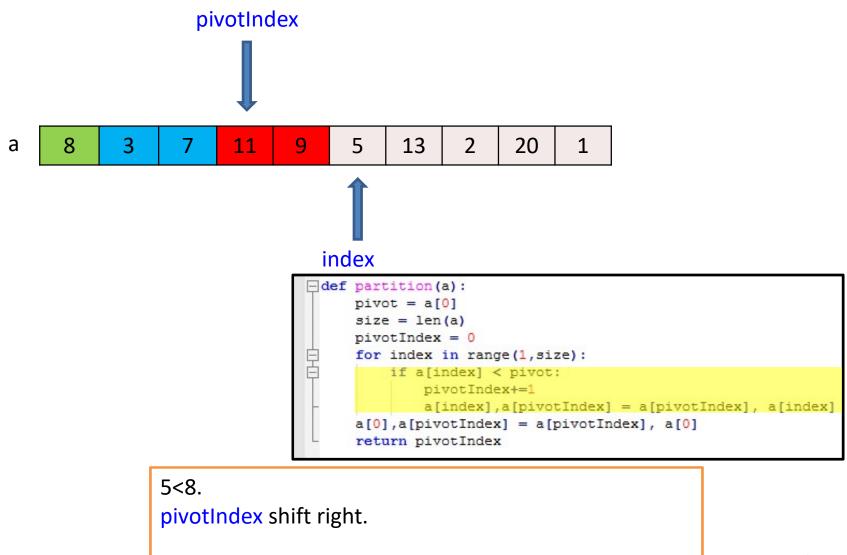




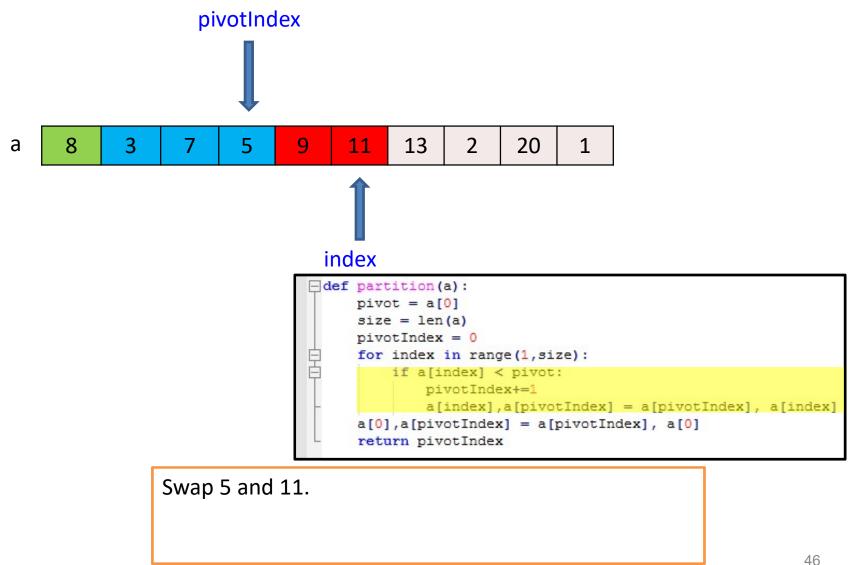




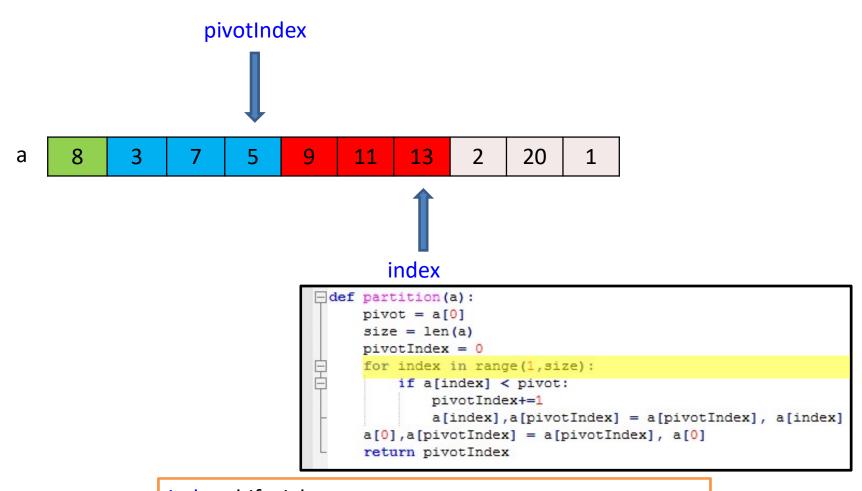






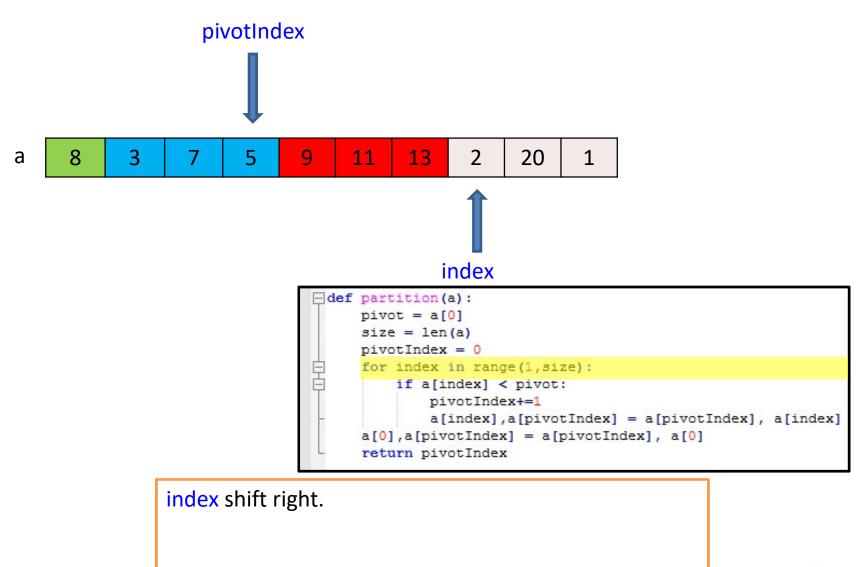




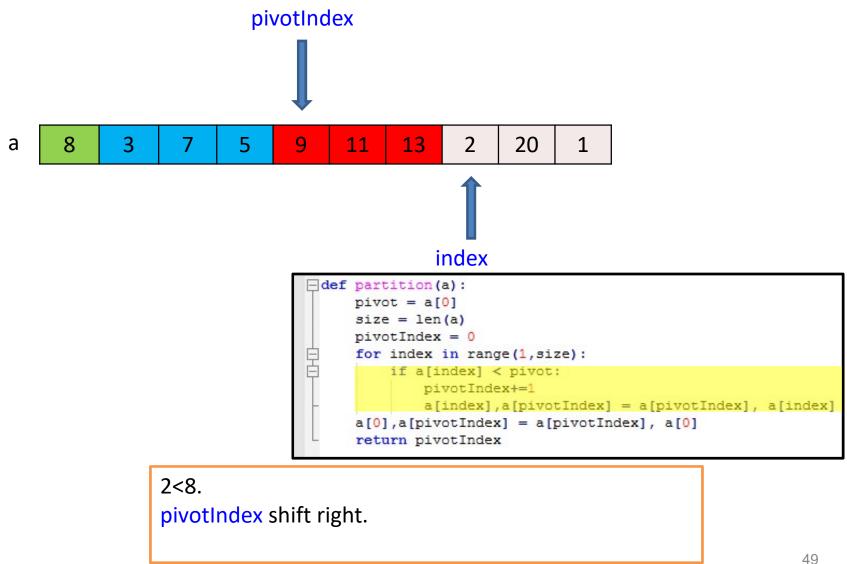


index shift right.13>8, do nothing.

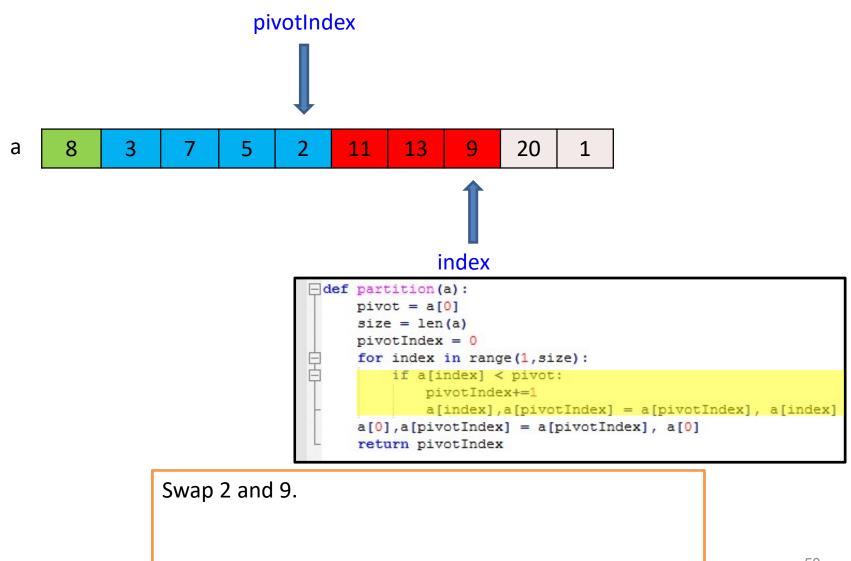




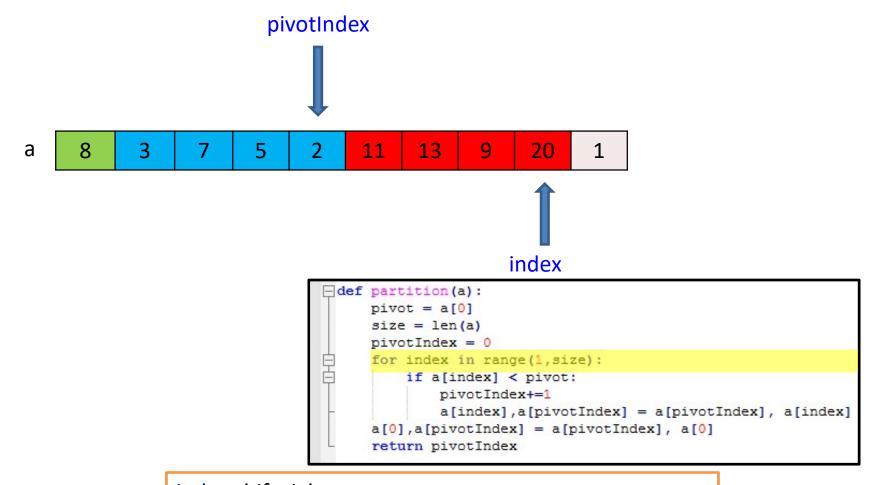






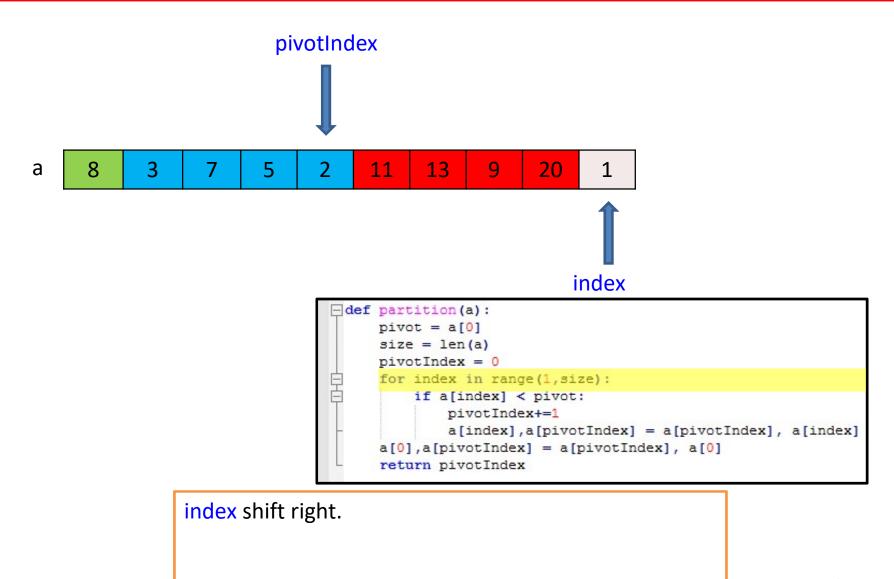




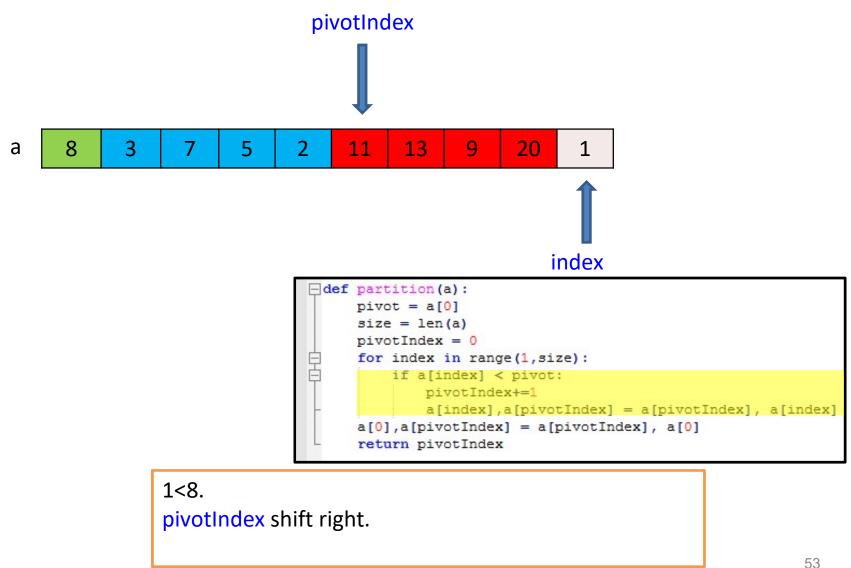


index shift right.20>8, do nothing.

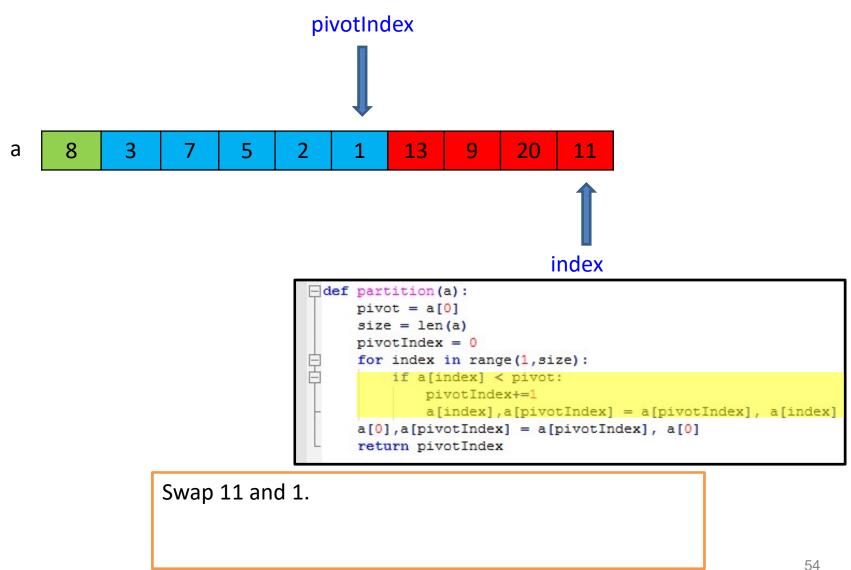




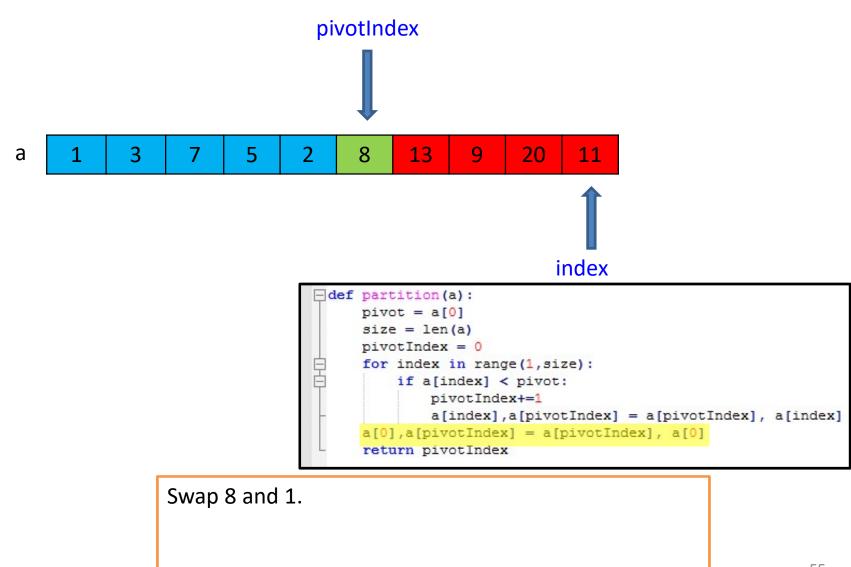












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Quick Sort - Complexity



Time complexity

- On average, each partition halves the size of the array to be sorted.
- On average, each partition swaps half the elements.
- On average, algorithm is $O(n \log n)$.
- Worst case, algorithm is $O(n^2)$.

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Quick Sort – Choice of Pivot



- In this version of quicksort, the leftmost element of the partition is used as the pivot element.
- Unfortunately, this causes worst-case behavior on already sorted arrays because the size of sub-array is only reduced by 1.
- This problem is easily solved by choosing:
 - 1. a random index for the pivot, or
 - 2. the middle index of the partition for the pivot, or
 - 3. the <u>median</u> of the first, middle and last elements of the partition for the pivot.