

Analyzing the Constraints

9:05pm

Constraint	T.C
$n \leq 10^6$	$O(N)$, $O(N \log N)$
$n \leq 20$	$O(N!)$, $O(2^N)$ →
$n \leq 10^{10}$	$O(\log N)$, $O(\sqrt{N})$

→ 10^9 → 1s
 10^{10}

→ Generic Guidelines.

Amazon, Adobe, Microsoft.

Q Given an array of 1's and 0's, you are allowed to replace only one 0 with 1. Find the maximum number of consecutive 1's that can be obtained after making the replacement.

Input = [1, 1, 0, 1, 1, 0, 1, 1]

Output = 5

= 10^9

Q

[1 1 0 1 1 0 1 1 1]

5 6

ans = 6.

Q.

[0, 1, 1, 1, 0, 1, 1, 0, 1, 1, 0]

4 5

ans = 6.

Qbvy

1, 1, 0, 1, 1, 0, 1, 1

L R

2 2

= L + R + 1

= 2 + 2 + 1

= 5

```
int arr[];
int n;
int max_one = 0;
```

→ [1, 1, 1, 1]

4

```
for (i = 0; i < n; i++)
{
    Co(N)
```

[1, 1, 0, 1, 1, 0]

+1

Count the no. of one if == length array return the count.

```
if (arr[i] == 0)
{
    int left = 0, right = 0;
```

$[\downarrow 0, \rightarrow 0, 0]$

$\underline{\underline{1}}$
ans = 1

$k = i - 1$
 $j = i + 1$

iteration
toward right
and counting. 1

while ($j \leq n$ & $arr[j] == 1$) {
 right++;
 j++;
}

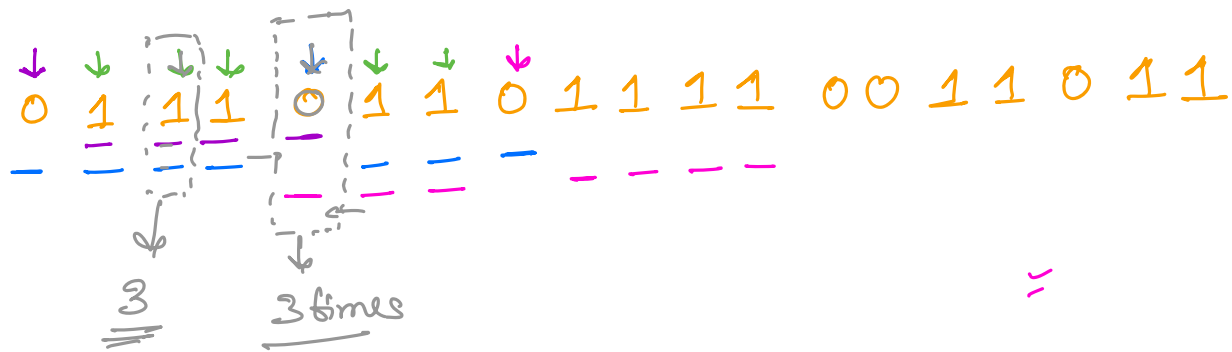
iterating
toward left
and count

while ($k \geq 0$ & $arr[k] == 1$) {
 left++;
 k--;
}

max

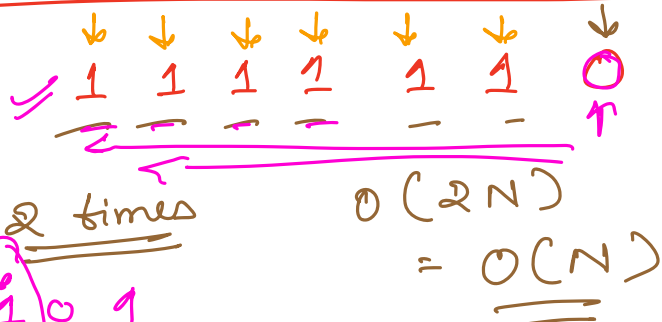
max-one = max (max-one,
 left + right + 1)

return max-one;

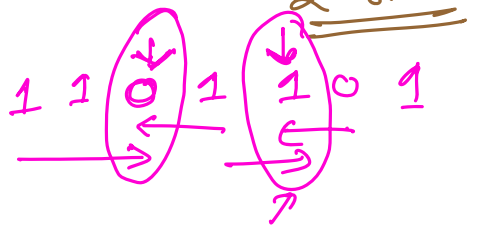


→ iterating 3 times. $O(3N)$
= $O(N)$

S.C. $O(1)$



$O(2N)$
= $O(N)$



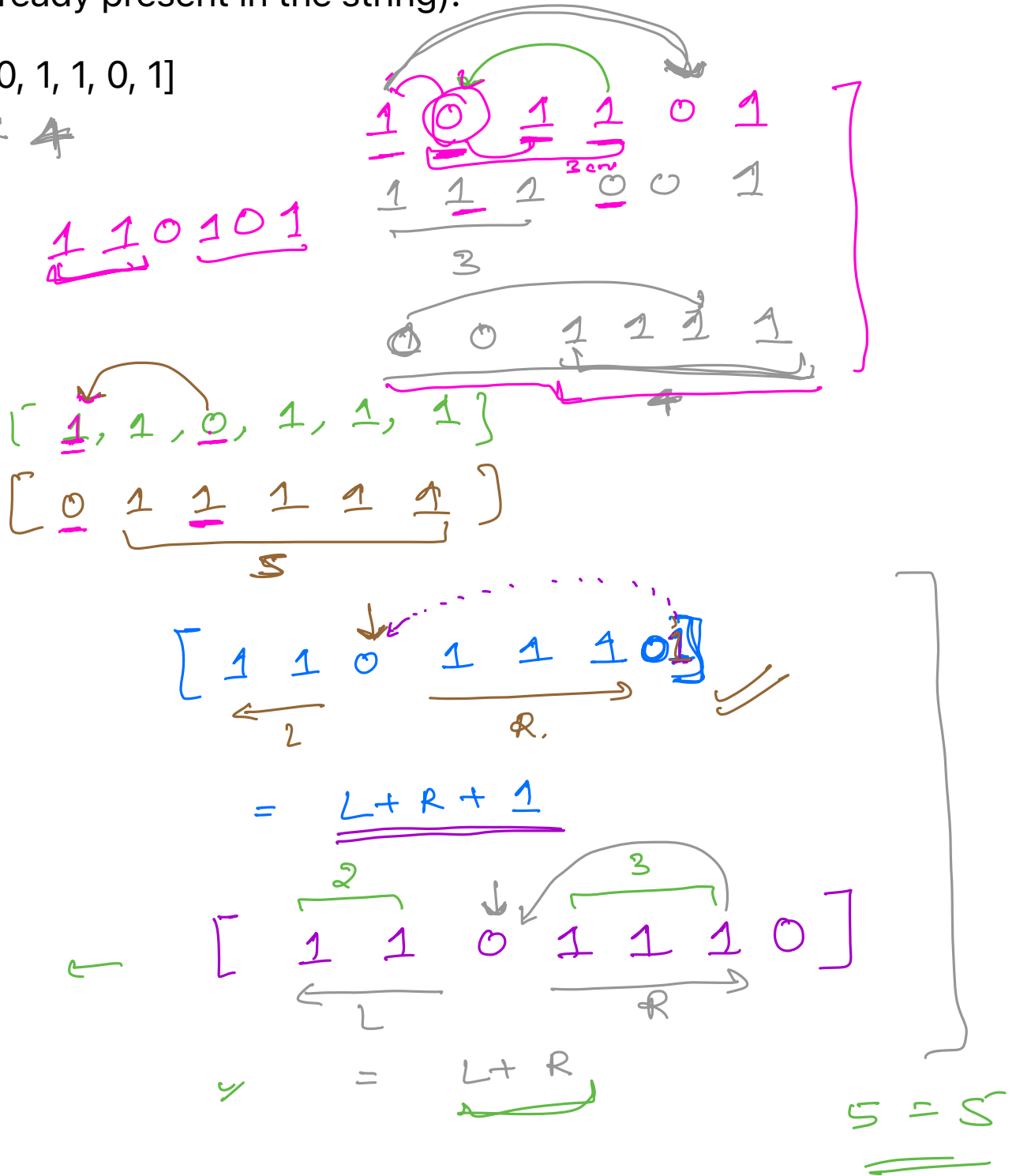
Variation

Amazon, Direct-i, Microsoft, Adobe.

Given an array of 1's and 0's, find the maximum number of consecutive 1's that can be obtained by SWAPPING at most one 0 with 1 (already present in the string).

Input: [1, 0, 1, 1, 0, 1]

Output: ~~5~~ 4



int max_one = 0

```

int arr[]
int n
int totalOne = 0
Count of one. ← [ for ( i = 0; i < n; i++ )
                    if ( arr[i] == 1 :
                        totalOne++; ] [1, 1, 1, 1]

```

```

if ( totalOne == n ) ←
    return n;

```

```

for ( i = 0; i < n; i++ )
{
    if ( arr[i] == 0 )
    {
        int left = 0
        int right = 0
        k = i - 1
        j = i + 1

```

```

while ( j < n && arr[j] == 1 ) {
    right++;
    j++;
}

```

```

while ( k >= 0 && arr[k] == 1 ) {
    left++;
    k--;
}

```

```

if ( left + right == n ) {
    maxOne = max(maxOne, l+r)
}
else {
    maxOne = max(maxOne, l+r+1)
}

```

T.C	$O(N)$
S.C	$O(1)$

Q

Given an array of N integers, find the majority element.

The majority element is the element that occurs more than $n/2$ times where n is size of the array.

$A[] = \{2, 1, 4\}$

$2 \rightarrow 1$ $1 \rightarrow 1$ $4 \rightarrow 1$

Ans = No Majority element ✓ - 1

$A[] = \{3, 4, 3, 2, 4, 4, 4, 4\}$

$3 \rightarrow 2$ $4 \rightarrow 5$ $2 \rightarrow 1$

Ans = 4

$8/2 = 4$

3, 4, 3, 6, 1, 3, 2, 5, 3, 3, 3

3 → 6 ✓

$4 \rightarrow 1$

$6 \rightarrow 1$

$2 \rightarrow 1$

$5 \rightarrow 1$

$11/2 = 5$

4, 6, 5, 3, 4, 5, 6, 4, 4, 4

$4 \rightarrow 5$

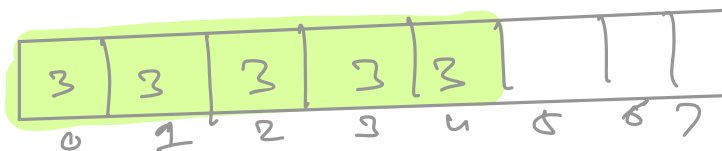
$6 \rightarrow 2$

$5 \rightarrow 2$

$3 \rightarrow 1$

⇒ No Majority element

$10/2 = 5$



$C_1 > n/2$ $C_2 > n/2$

only 1 majority element.

$C_1 + C_2$

$C_1 + C_2 = n$

$n/2 + n/2 \geq n$

~~2~~ ~~1~~ ~~1~~ ~~1~~ ~~1~~ ~~2~~ ~~2~~

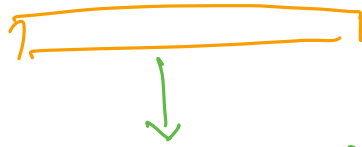
ma = 1
count = 1

Brute Force

①

Sorting -
calculate count

T.C. = $O(n \log n)$
S.C. = $O(1)$ ✓



two loop
and count }

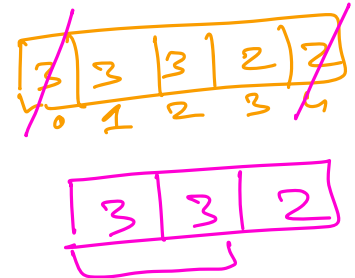
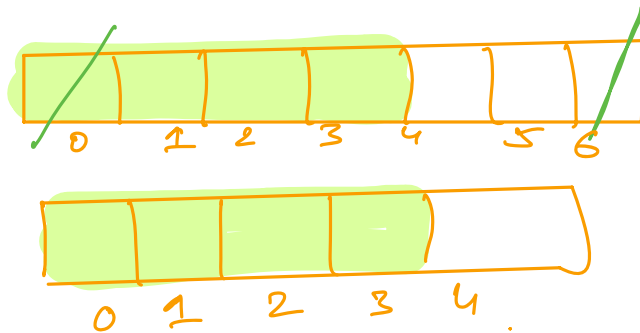
T.C. = $O(N^2)$
S.C. = $O(1)$

③

Hash Map freq
check for $> n/2$

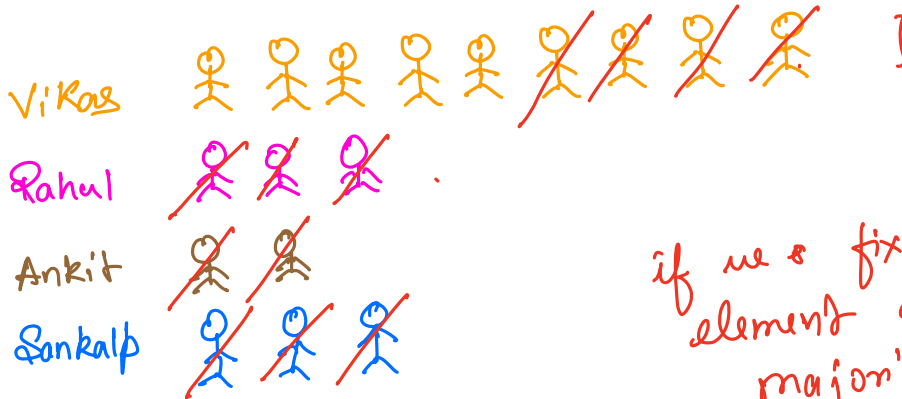
T.C. = $O(n)$ ✓
S.C. = $O(n)$

Optimize it



Obs 2

Elections



if we fix two distinct
element and remove
majority doesn't.

Remove	Orange	Pink	Brown	Blue	Winner
1 Vikas & 1 Rahul	8	2	2	3	orange
1 Vikas & 1 Sankalp	7	2	2	2	orange
1 Vikas & 1 Ankit	6	2	1	2	orange
1 Rahul & 1 Ankit	6	1	0	2	orange
1 Rahul & 1 Sankalp	6	0	0	1	orange
1 Vikas & 1 Sankalp	5	0	0	0	orange

Moore's Voting algorithm

0	1	2	3	4	5	6	7	8	9	10
3	4	3	6	1	3	2	5	3	3	3

majority element = ~~3~~ ~~4~~ ~~3~~ ~~6~~ ~~1~~ ~~3~~ ~~2~~ ~~5~~ ~~3~~ ~~3~~ 3

count = 0 ~~1~~ ~~1~~ ~~1~~ ~~1~~ ~~1~~ ~~1~~ ~~1~~ ~~1~~ ~~1~~ ~~1~~ 3

ans = 3

~~1~~ ~~2~~ ~~1~~ ~~2~~ 1
 $n = 5$
 $n/2 = 2$
3

1	6	1	6	5
---	---	---	---	---

majority element: ~~1~~ ~~6~~ 5

count:

~~1~~ ~~0~~ ~~1~~ ~~0~~ 1

→ No major

$1 > n/2 = 1$
 frequency of the majority
 if it is $> n/2$
 then ans =
 else -1

2	1	<u>1</u>	2	2
--------------	--------------	----------	--------------	--------------

majority 1
 count = 1

majority 1

count = 3
 $n/2$

```
int arr[]
```

```
int n
```

```
int major_index = 0
```

```
int count = 1
```

```
for (i = 1 ; i < n ; i++)
```

```
{
```

```
if (count == 0)
```

```
{
    maj_index = i;
}
```

```
count = 1
```

```
else
```

```
if (arr[maj_index] == arr[i])
```

```
count++;
```

T.C. $O(N)$

S.C. $O(1)$

✓

⇒

⇒

$$\begin{bmatrix} 1 & 2 & -3 & -4 \\ -5 & -6 & 7 & 0 \\ -9 & -2 & 0 & -4 \end{bmatrix}$$

T.C. $O(N^2)$
S.C. $O(1)$

H.W

2nd it ⇒

$$\begin{bmatrix} 1 & 2 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

Algo → (1) find zero and make every element in the row and column of that zero negative.
 $O(N^2)$

+ (2) traverse the matrix and convert all the negative to zero.
 $O(N^2)$