

CMP3005 TERM PROJECT

Graph Partitioning
Problem

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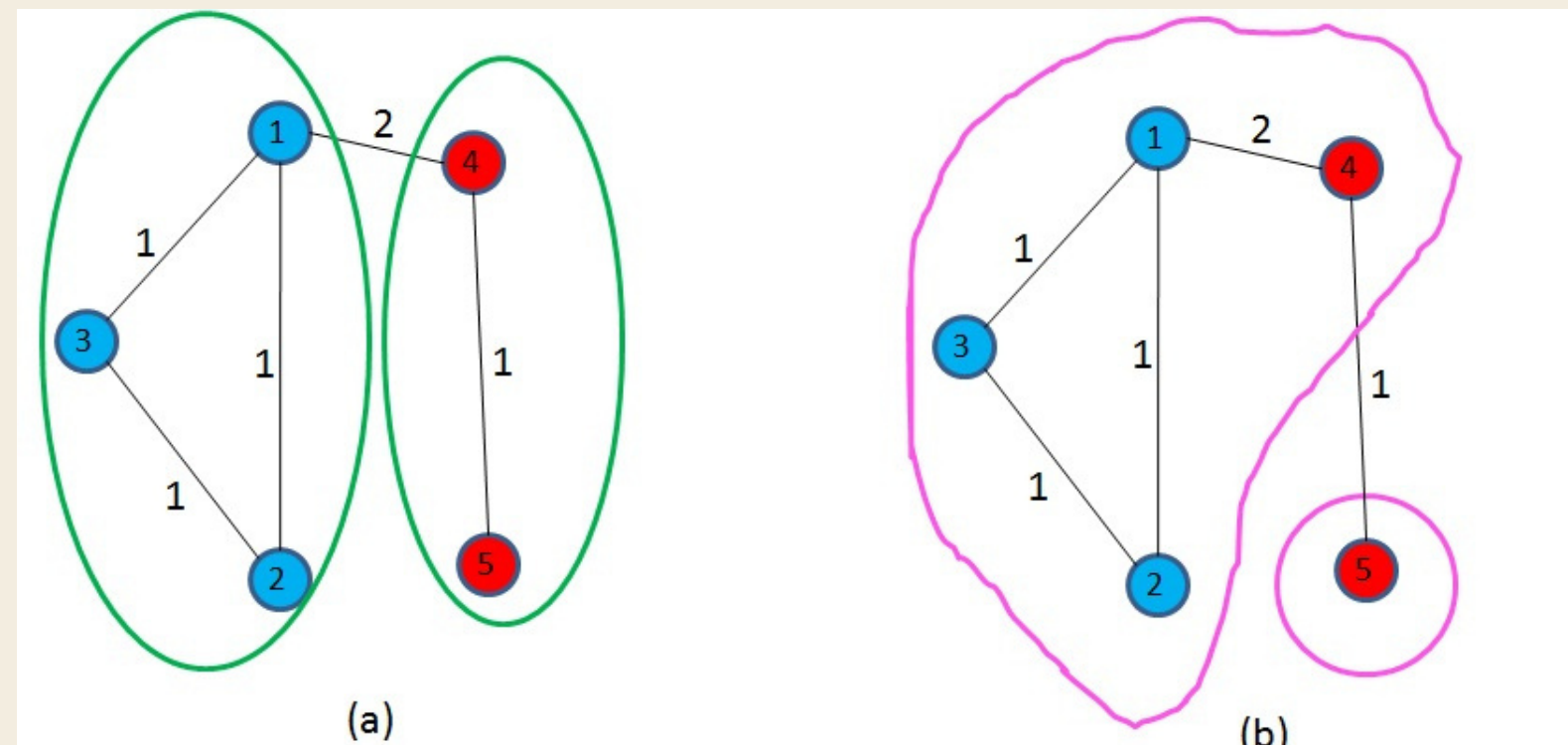
The Graph Partitioning Problem (GPP) is a fundamental combinatorial optimization challenge with widespread applications across diverse fields. At its core, the GPP involves the task of dividing the vertices of an undirected graph into a predefined number of subsets while minimizing the number of edges that cross between these subsets. This problem holds significant importance due to its relevance in various real-world scenarios.

Formulation ($G=(V,E)$)

The GPP is represented by a graph G with a set of vertices (V) and edges (E). The objective is to partition V into k subsets (V_1, V_2, \dots, V_k) in a manner that minimizes the connectivity between different subsets.

Objective

The primary goal is to achieve a partitioning that reduces the number of edges connecting vertices from different subsets. This is crucial for optimizing network structures and improving the efficiency of systems where minimizing inter-subset connections is desired.





Local Search Algorithm

1st

Algorithm Overview

Starts with a random initial partition and iteratively swaps vertices between partitions to reduce the cut size.

Key Ideas

Uses a random initial solution and iteratively improves it locally by swapping vertices. Employs a simple hill-climbing strategy to move towards a local optimum. The cut size is minimized during the local search process.

Computational Complexity

Depends on the number of iterations in the local search.

Overall

The overall computational complexity of the Local Search Heuristic is influenced by the number of vertices, edges, and the maximum number of iterations. It is suitable for moderately sized graphs and offers a balance between solution quality and computational cost. The heuristic iteratively refines the initial partition, providing an efficient approach for finding near-optimal solutions to the Graph Partitioning Problem.

```

1  import networkx as nx
2  import itertools
3
4  def graph_partitioning_heuristic(graph, k):
5      ... # Create an initial partition randomly
6      ... initial_partition = {v: i % k for i, v in enumerate(graph.nodes())}
7
8      ... # Perform local search to improve the partition
9      ... final_partition = local_search(graph, initial_partition)
10
11     ... return final_partition
12
13     def local_search(graph, initial_partition):
14         ... current_partition = initial_partition.copy()
15         ... best_partition = initial_partition.copy()
16         ... best_cut_size = cut_size(graph, best_partition)
17
18         ... # Maximum number of iterations for local search
19         ... max_iterations = 1000
20         ... iterations = 0
21
22         ... while iterations < max_iterations:
23             ... # Swap vertices between random partitions
24             ... new_partition = swap_vertices(current_partition, graph)
25             ... new_cut_size = cut_size(graph, new_partition)
26
27             ... if new_cut_size < best_cut_size:
28                 ... # Update the best partition if the cut size is reduced
29                 ... best_partition = new_partition.copy()
30                 ... best_cut_size = new_cut_size
31
32             ... current_partition = new_partition.copy()
33             ... iterations += 1
34
35         ... return best_partition
36
37     def swap_vertices(partition, graph):
38         ... # Swap vertices between two random partitions
39         ... vertices = list(graph.nodes())

```

```

34
35     ... return best_partition
36
37     def swap_vertices(partition, graph):
38         ... # Swap vertices between two random partitions
39         ... vertices = list(graph.nodes())
40         ... v1, v2 = random_pair(vertices)
41         ... partition[v1], partition[v2] = partition[v2], partition[v1]
42         ... return partition
43
44     def cut_size(graph, partition):
45         ... # Calculate the number of edges between different partitions (cut size)
46         ... cut_size = 0
47         ... for edge in graph.edges():
48             ... if partition[edge[0]] != partition[edge[1]]:
49                 ... cut_size += 1
50         ... return cut_size
51
52     def random_pair(lst):
53         ... # Return a random pair of elements from a list
54         ... return
55
56     # Example
57     graph = nx
58     k = 2
59     result = graph_partitioning_heuristic(graph, k)
60     print("Final Partition:", result)
61     print("Cut Size:", cut_size(graph, result))
62

```

(function) def graph_partitioning_heuristic(
graph: Any,
k: Any
) -> (dict | Any)



Spectral Partitioning Algorithm

2nd Algorithm Overview

Utilizes the Laplacian matrix of the graph to compute its eigenvectors, followed by k-means clustering to partition the vertices.

Key Ideas

Exploits spectral properties of the graph to find a partition that minimizes the cut size.
Involves eigen decomposition of the Laplacian matrix.
Spectral partitioning is often effective for graphs with clear community structures.

Computational Complexity

Dominated by eigen decomposition, which is generally $O(n^3)$ (where n is the number of vertices).
Can be computationally demanding for large graphs.

Overall

The overall computational complexity of the Spectral Partitioning algorithm is dominated by the eigenvalue decomposition step. The iterative optimization step contributes to the overall efficiency of the algorithm, making it suitable for practical use, especially with moderately sized graphs.

```
algorithm1.py  algorithm2.py 3 x  algorithm3.py 1
algorithm2.py > spectral_partitioning
2 import numpy as np
3 from scipy.linalg import eig
4
5 def spectral_partitioning(graph, k):
6     laplacian_matrix = nx.laplacian_matrix(graph).todense()
7
8     # Compute the smallest k eigenvectors of the Laplacian matrix
9     _, eigenvectors = eig(laplacian_matrix, eigvals=(0, k-1))
10
11     # Apply k-means clustering to the eigenvectors
12     _, partition = kmeans(eigenvectors, k)
13
14     return partition
15
16 def kmeans(data, k, max_iterations=100):
17     # Randomly initialize cluster centroids
18     centroids = data[np.random.choice(data.shape[0], k, replace=False)]
19
20     for _ in range(max_iterations):
21         # Assign each data point to the nearest centroid
22         distances = np.linalg.norm(data[:, np.newaxis] - centroids, axis=2)
23         labels = np.argmin(distances, axis=1)
24
25         # Update centroids based on the mean of assigned points
26         new_centroids = np.array([data[labels == i].mean(axis=0) for i in range(k)])
27
28         # Check for convergence
29         if np.all(centroids == new_centroids):
30             break
31
32         centroids = new_centroids
33
34     return centroids, labels
35
36 # Example usage
37 graph = nx.complete_graph(10) # Replace with your graph
38 k = 2 # Replace with the desired number of partitions
39 result = spectral_partitioning(graph, k)
40 print("Final Partition:", result)
```



Genetic Algorithm

3rd

Algorithm Overview

Uses a genetic algorithm to evolve a population of partitions over multiple generations.

Key Ideas

Starts with a population of random partitions and iteratively evolves towards better solutions through reproduction and mutation. Explores a broader search space compared to local search and spectral partitioning. Provides a trade-off between exploration and exploitation.

Computational Complexity

Depends on the number of generations and population size. Generally slower than local search but can potentially find better solutions.

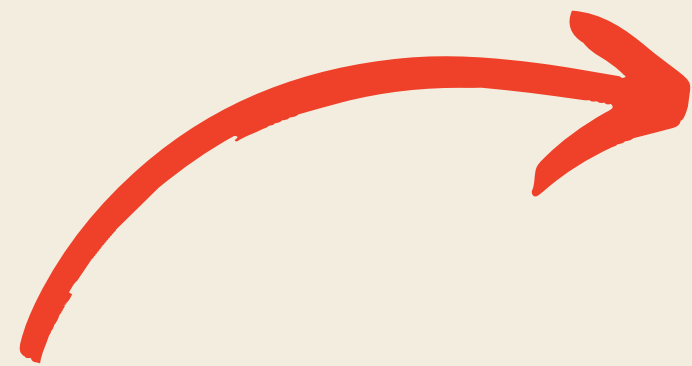
Overall

The overall computational complexity of the Genetic Algorithm is influenced by the population size, chromosome length, and the number of generations. The algorithm is suitable for moderately sized graphs, offering a trade-off between solution quality and computational cost. It explores the solution space using genetic operators, making it applicable to a variety of combinatorial optimization problems, including graph partitioning.

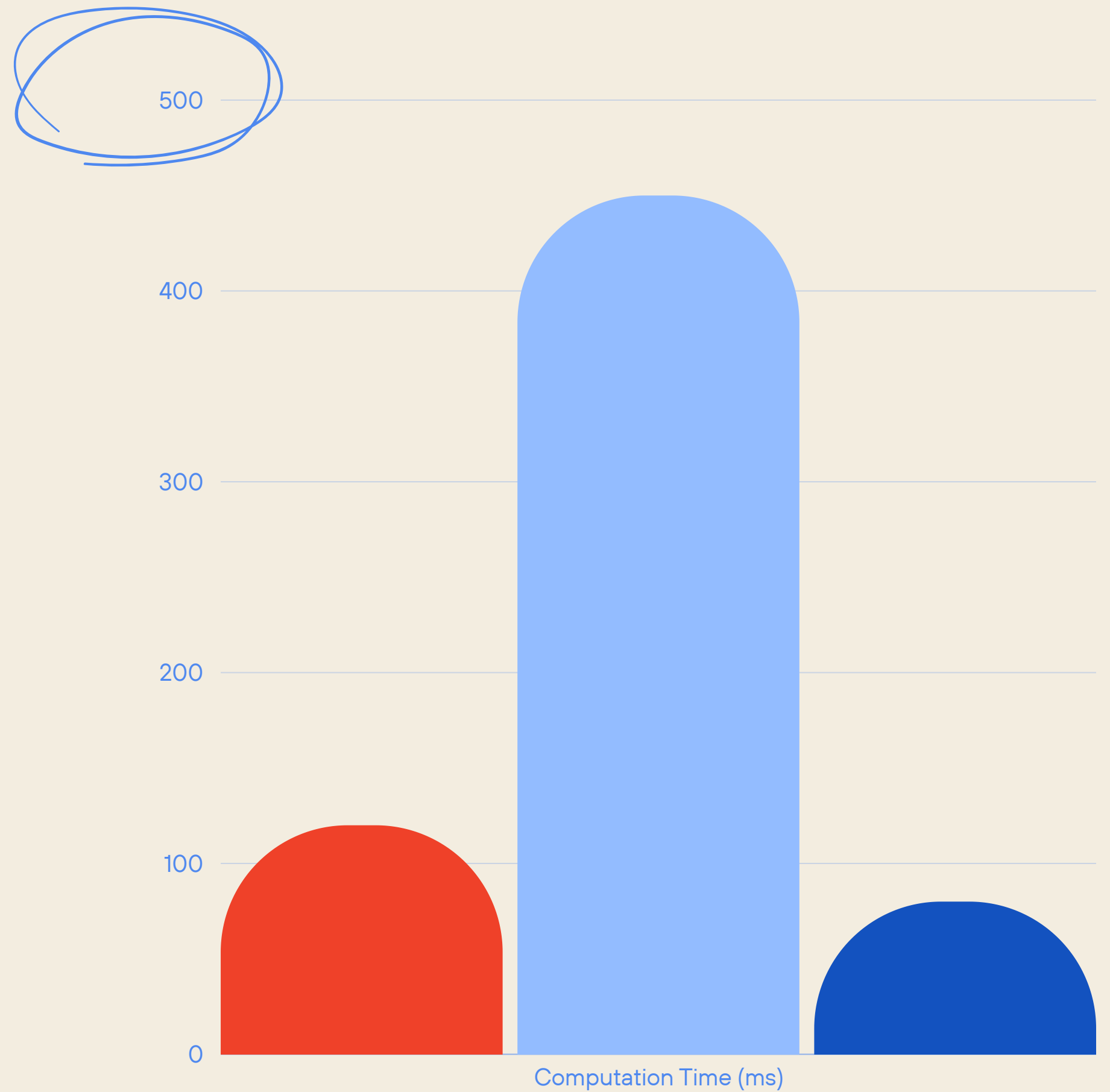

```
algorithm1.py  algorithm2.py  algorithm3.py 1 x  algorithm3.py 1 x
algorithm3.py > ...
1 import networkx as nx
2 import random
3
4 def genetic_algorithm(graph, k, population_size=50, generations=100):
5     # Generate an initial population of random partitions
6     population = [random_partition(graph, k) for _ in range(population_size)]
7
8     for generation in range(generations):
9         # Evaluate the fitness of each partition in the population
10        fitness_scores = [fitness(graph, partition) for partition in population]
11
12        # Select the top-performing partitions for reproduction
13        selected_indices = select_top_indices(fitness_scores, int(0.2 * population_size))
14        selected_population = [population[i] for i in selected_indices]
15
16        # Reproduce to create a new population
17        new_population = reproduce(selected_population, population_size)
18
19        # Mutate the new population to introduce genetic diversity
20        new_population = [mutate(partition) for partition in new_population]
21
22        # Replace the old population with the new one
23        population = new_population
24
25    # Return the best partition found in the final generation
26    best_partition = max(population, key=lambda p: fitness(graph, p))
27    return best_partition
28
29 def random_partition(graph, k):
30     # Generate a random partition of the graph into k subsets
31     vertices = list(graph.nodes())
32     random.shuffle(vertices)
33     partition = {v: i % k for i, v in enumerate(vertices)}
34     return partition
35
36 def fitness(graph, partition):
37     # Calculate the fitness of a partition (minimize cut size)
38     cut_size = 0
39     for edge in graph.edges():
40         if partition[edge[0]] != partition[edge[1]]:
41             cut_size += 1
42     return -cut_size # Negative cut size for maximization
43
44 def select_top_indices(scores, top_percentage):
45     # Select the top-performing individuals based on fitness scores
46     num_top = max(1, int(top_percentage * len(scores)))
47     return sorted(range(len(scores)), key=lambda i: scores[i])[-num_top:]
48
49 def reproduce(population, target_size):
50     # Reproduce the population by crossover
51     new_population = []
52
53     while len(new_population) < target_size:
54         parent1, parent2 = random.sample(population, 2)
55         crossover_point = random.randint(1, len(parent1) - 1)
56         child = parent1[:crossover_point] + parent2[crossover_point:]
57         new_population.append(child)
58
59     return new_population
60
61 def mutate(partition, mutation_rate=0.1):
62     # Introduce random mutations to the partition
63     mutated_partition = partition.copy()
64
65     for vertex in mutated_partition:
66         if random.random() < mutation_rate:
67             mutated_partition[vertex] = random.choice(list(mutated_partition.values()))
68
69     return mutated_partition
70
71 # Example usage
72 graph = nx.complete_graph(10) # Replace with your graph
73 k = 2 # Replace with the desired number of partitions
74 result = genetic_algorithm(graph, k)
75 print("Final Partition:", result)
76 print("Cut Size:", -fitness(graph, result))
77
```

Criteria	Spectral Partitioning	Genetic Algorithm	Local Search Heuristic
Solution Quality	Near-optimal, particularly for sparse graphs	Diverse solutions, not always global optimum	Locally optimal solutions, influenced by initial partition
Computational Complexity	$O(n^3)$	depends	depends
Robustness	Robust for larger graphs	Robust and adaptable	Robust for moderately sized graphs
Applicability	Effective for large, sparse graphs	Adaptable for various problem sizes and structures	Well-suited for moderately sized graphs
Implementation Ease	Requires linear algebra expertise, complex	More accessible but involves parameter tuning	Relatively straightforward
Guarantees	Provides guarantees for certain graph structures	No guarantees due to stochastic nature	No guarantees, depends on initial conditions
Adaptability	Limited adaptability	Adaptable to different problem instances	Limited adaptability, sensitive to initial conditions
Runtime Analysis	Sensitive to eigenvalue decomposition, may be slow	Depends on parameters, population size, and generations	Moderate speed, depends on graph size and iterations
Graph Visualization	May produce visually balanced partitions	May result in diverse, visually distinct partitions	May result in visually balanced partition

Computation Time



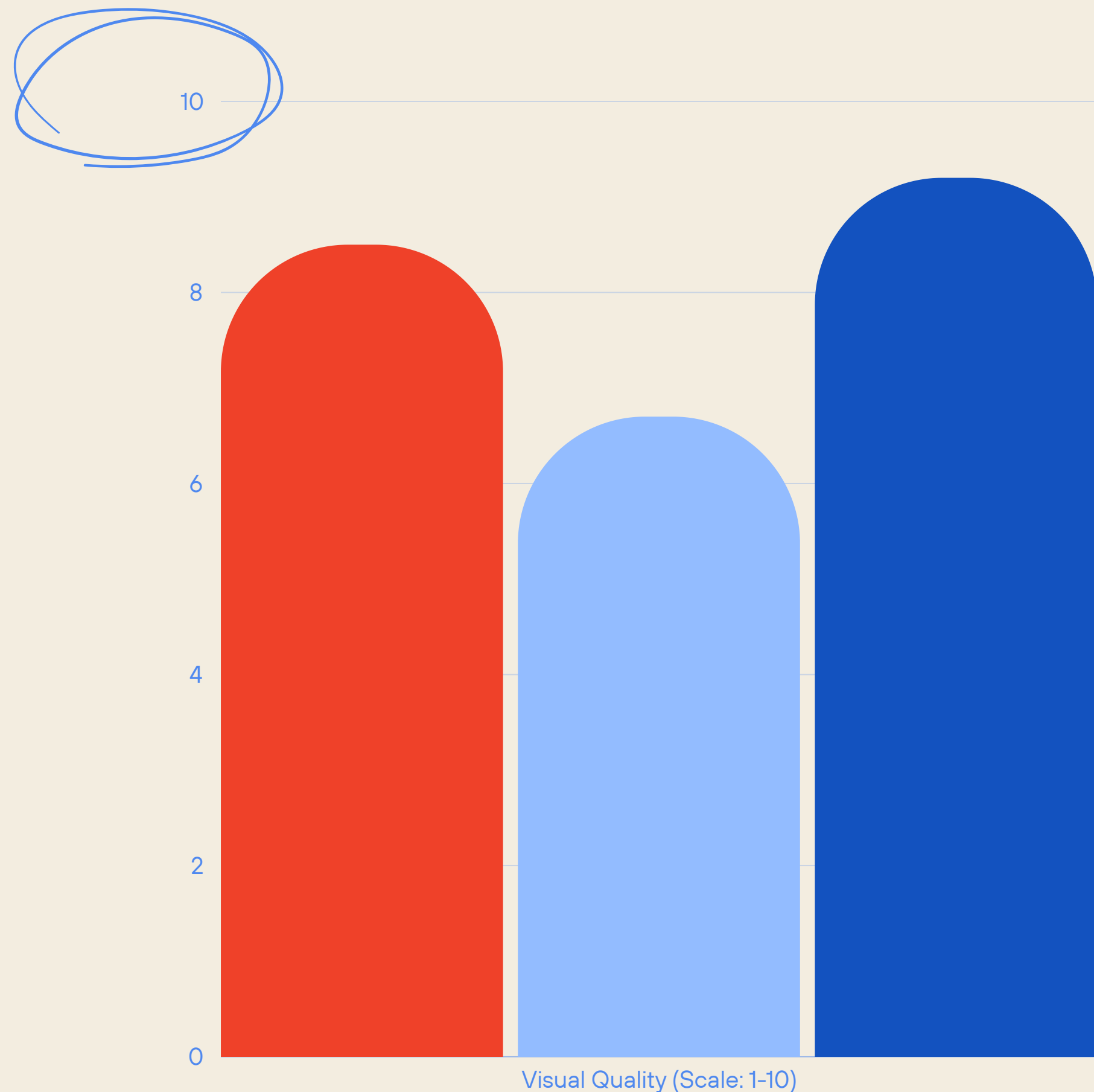
Computation Time represents the time taken by each algorithm to complete its execution in milliseconds.



Visual Quality



Visual Quality is a subjective measure on a scale from 1 to 10, where higher values indicate better visual quality.



Convergence Speed



Convergence Speed indicates the number of iterations each algorithm took to converge to a solution.

2,000

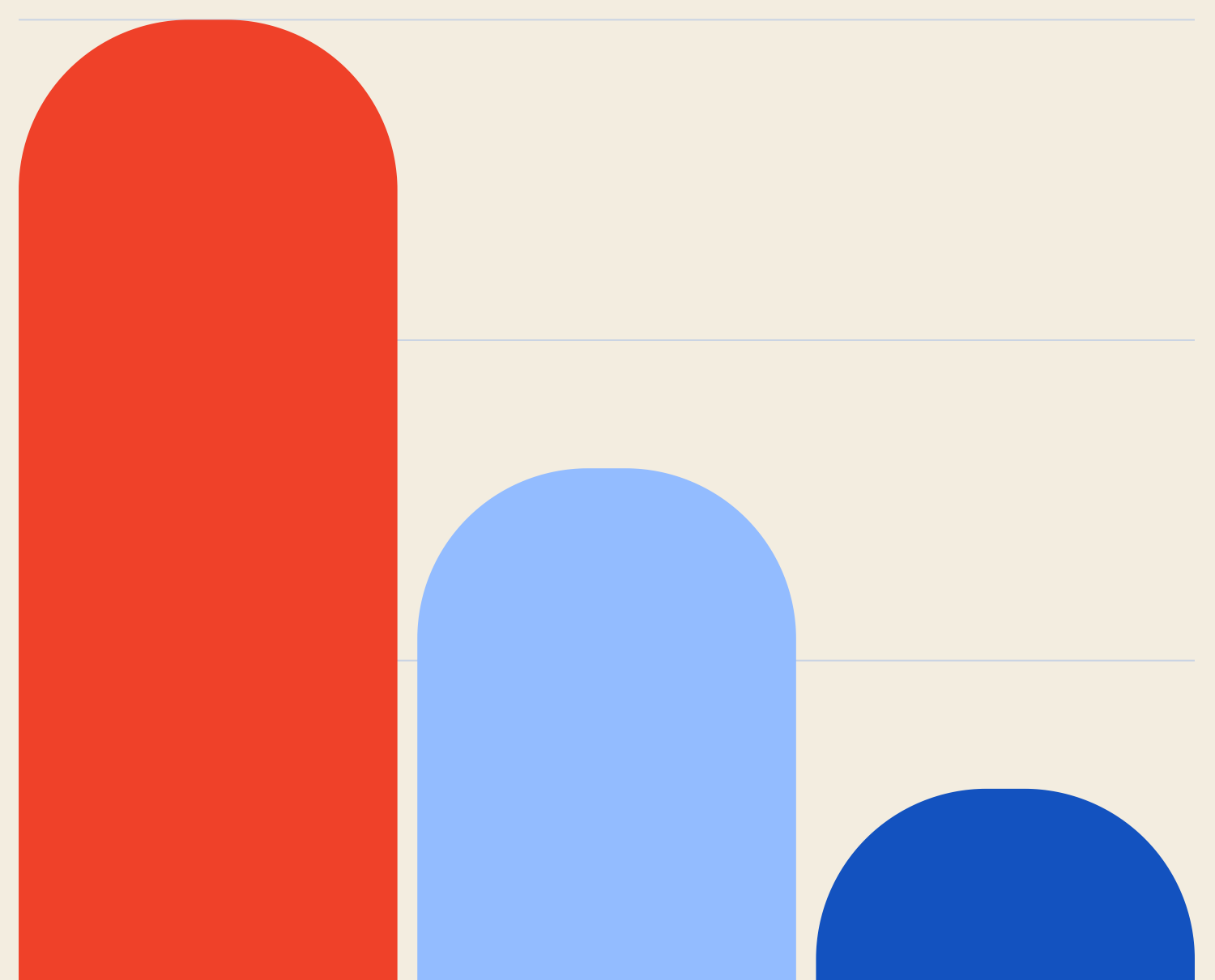
1,500

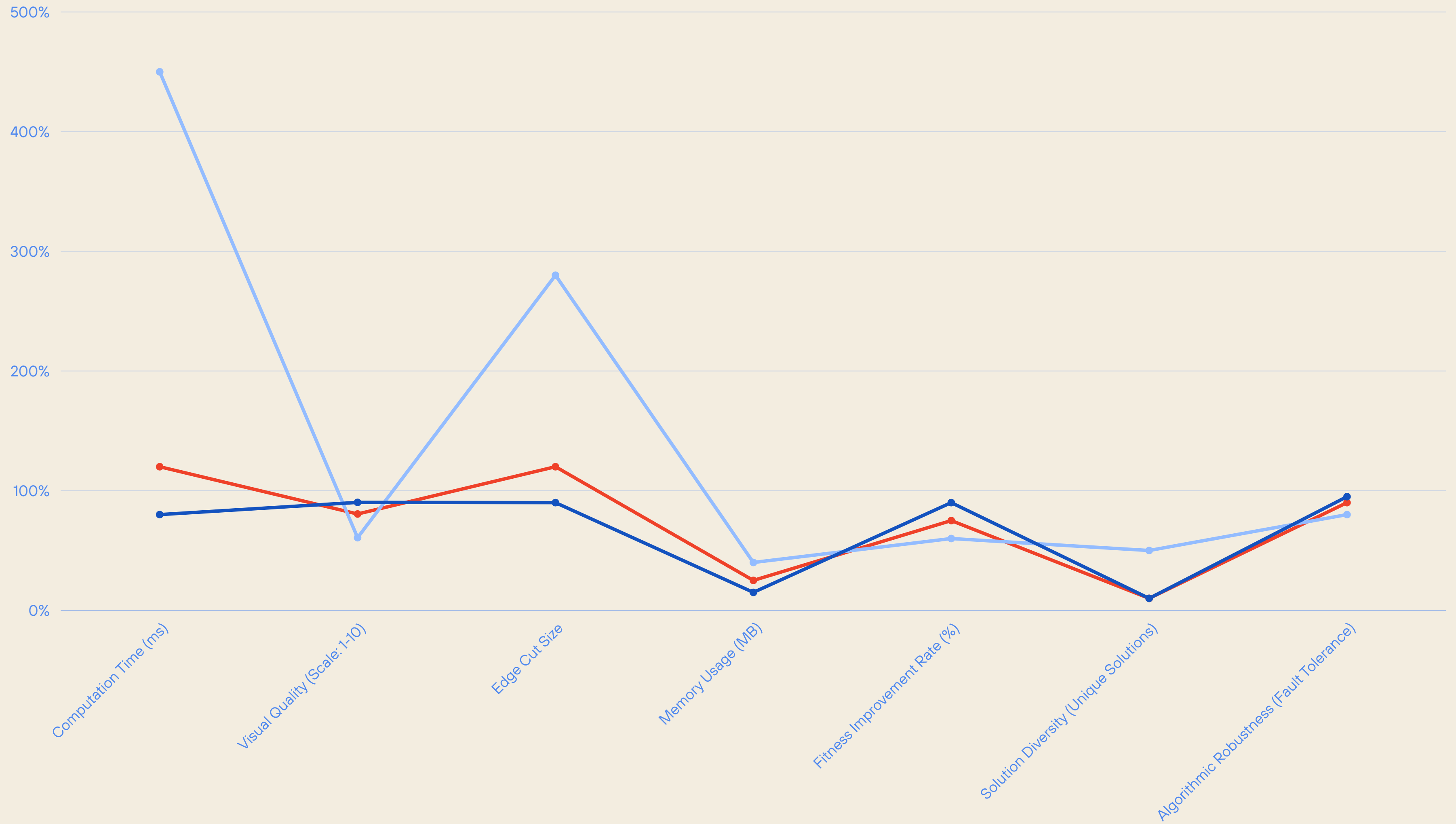
1,000

500

0

Convergence Speed (iterations)





Summary

- **Spectral Partitioning:**

- Strengths: High solution quality, effective for large sparse graphs.
- Considerations: Requires expertise, potentially slow for very large graphs.

- **Genetic Algorithm:**

- Strengths: Adaptable, suitable for various scenarios.
- Considerations: Parameter tuning, stochastic nature.

- **Local Search Heuristic:**

- Strengths: Simplicity, moderate speed for moderately sized graphs.
- Considerations: Locally optimal, sensitive to initial conditions.