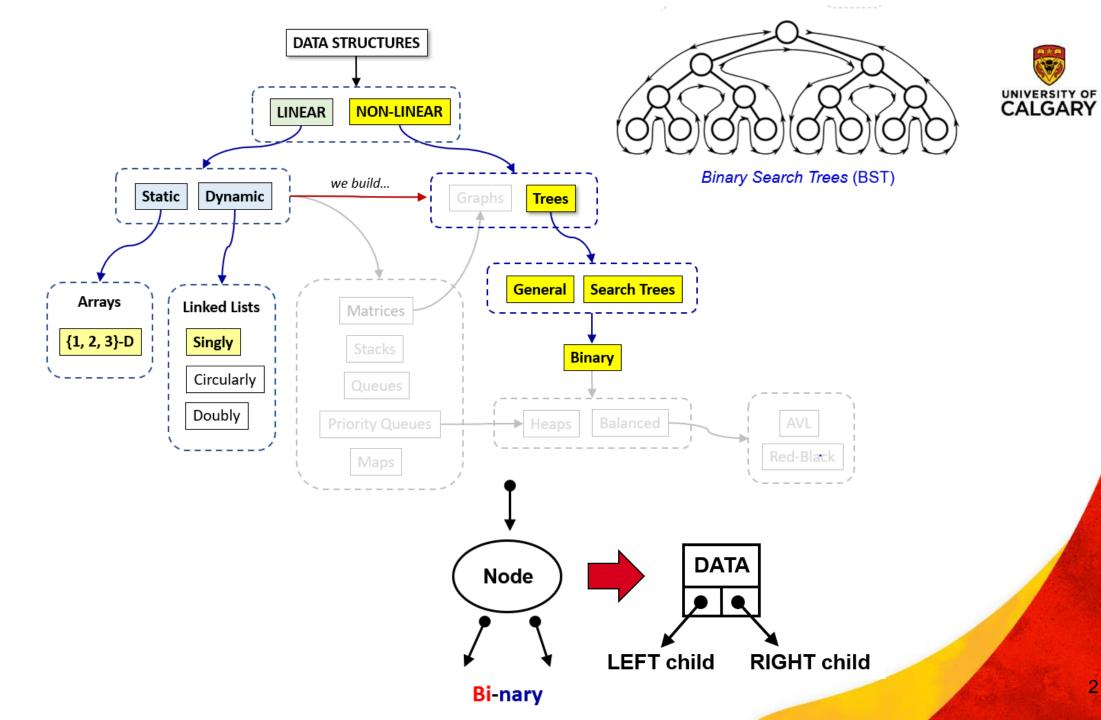


# **Binary Search Trees**

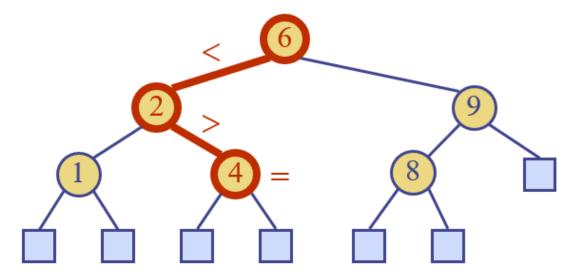




## **Binary Search Tree (BST)**

- A common form of the binary tree found in computer applications is the binary search tree.
- An ordered tree with a search structure where at any position in the tree all nodes less than the current node lie to the left and all nodes greater than the current node lie to the right

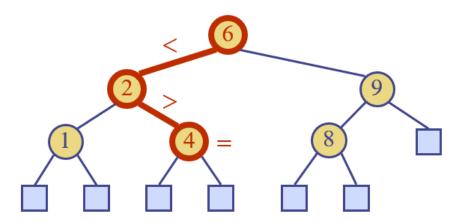
(LEFT child → DATA) < parent < (RIGHT child → DATA)





### **BST Insertions: Scenario**

- We have a Binary Search Tree
  - It can be empty
  - Or have some elements in it already
- We want to add an element to it
  - Inserting/adding involves 2 steps:
    - Find the correct location
    - Do the steps to add a new node
- Must maintain "search" structure:
   (LEFT child → DATA) < parent < (RIGHT child → DATA)</li>

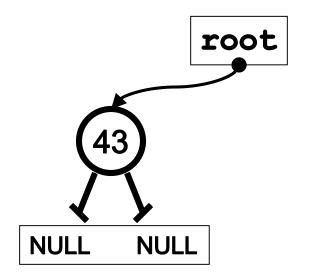






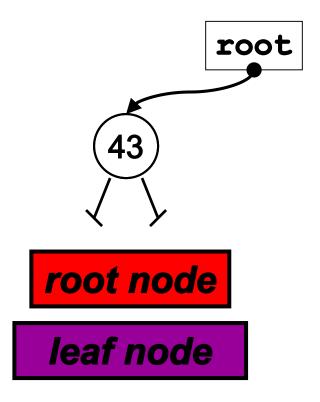
- What follows is a step by step visualization of BST insertion
- An array of integers is the input
- The tree begins empty



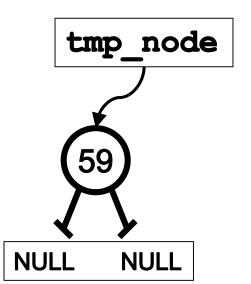


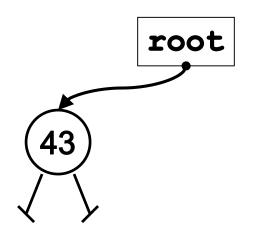


- Allocate Node: root
- Insert Data: root→DATA = 43
- Update Left and Right references (i.e., pointers) to NULL:





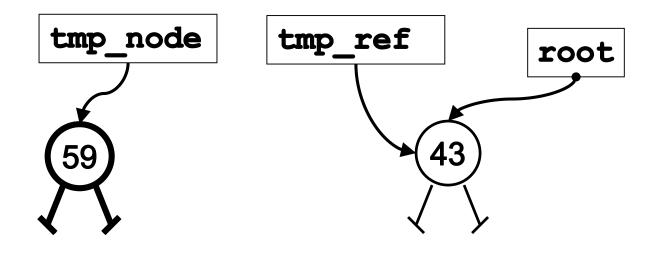






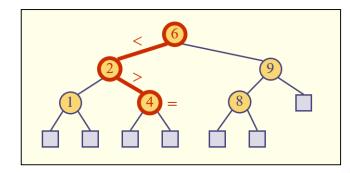
- Allocate New Node: tmp\_node
- Insert Data in New Node: tmp\_node→DATA = 59
- Update Left and Right references (i.e., pointers) to NULL:
  - √ tmp\_node→LEFT = NULL; tmp\_node→RIGHT = NULL

43	59	40	31	64	33	20	56	47	28	89

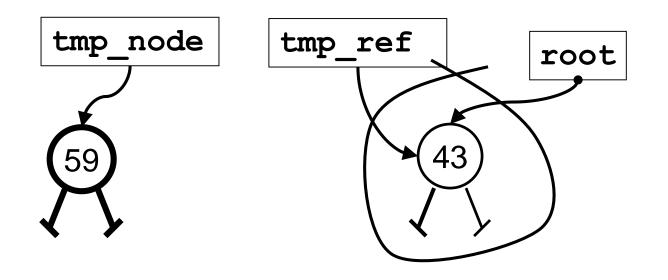




#### **BST** "search" structure

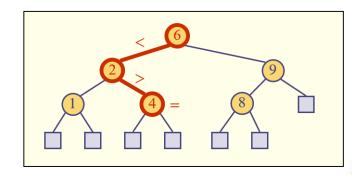


- (...)
- tmp\_ref = root
   i.e., tmp\_ref will reference (i.e., point at)
   the location in the BST to insert the new node

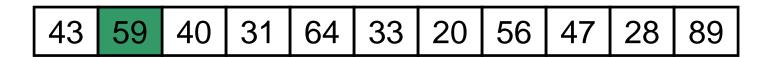


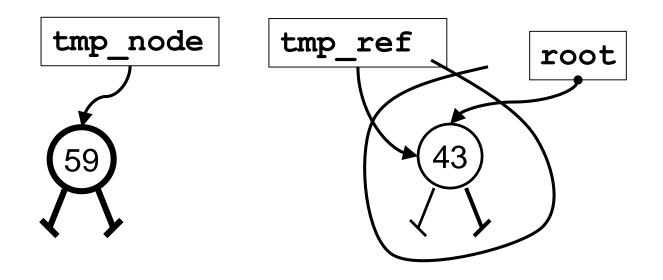


#### BST "search" structure



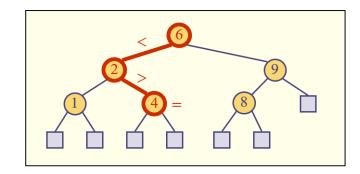
- (...)
- tmp\_node → DATA < tmp\_ref → DATA? Insert @ tmp\_ref → LEFT



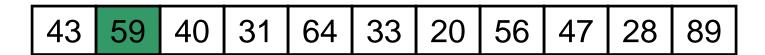


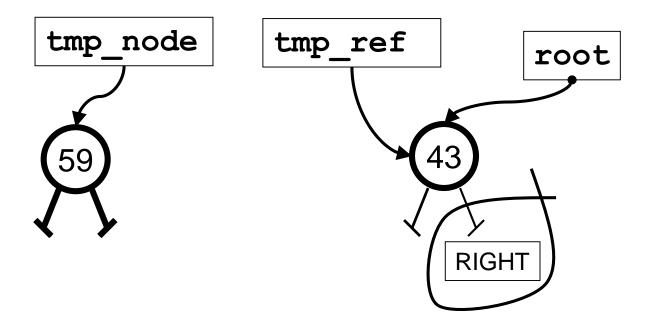


#### **BST** "search" structure



- (...)
- tmp\_node → DATA < tmp\_ref → DATA? Insert @ tmp\_ref → LEFT
- tmp\_node→DATA > tmp\_ref→DATA? Insert @ tmp\_ref→RIGHT

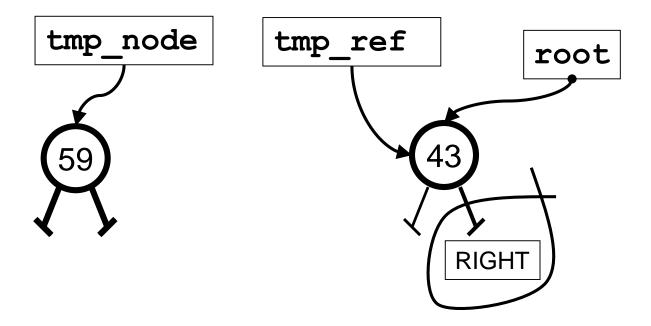






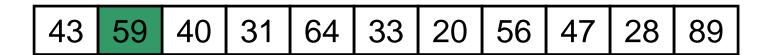
- (...)
- tmp\_node  $\rightarrow$  59 < tmp\_ref  $\rightarrow$  43? NO
- tmp\_node → 59 > tmp\_ref → 43? YES: Insert @ tmp\_ref → RIGHT

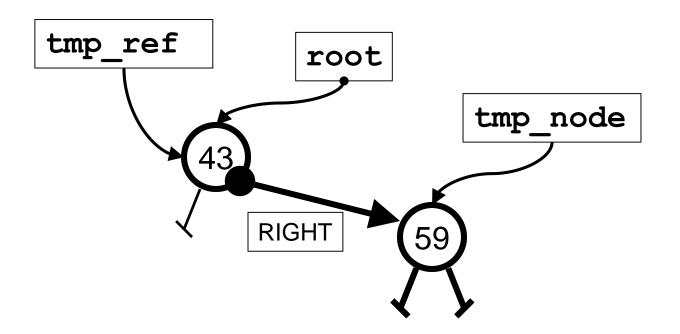






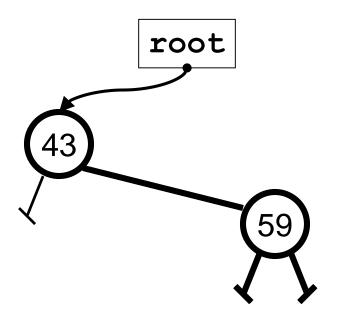
- (...)
- tmp\_node → 59 > tmp\_ref → 43? YES: Insert @ tmp\_ref → RIGHT
   ✓ tmp\_ref → RIGHT = tmp\_node





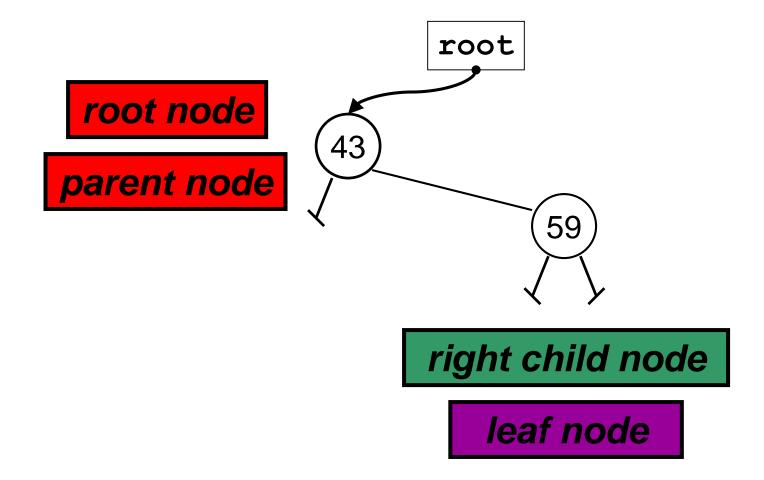


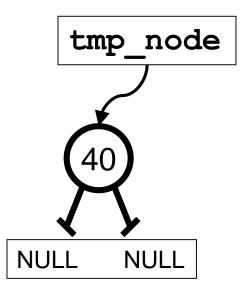
- (...)
- tmp\_node → 59 > tmp\_ref → 43? YES: Insert @ tmp\_ref → RIGHT
   ✓ tmp\_ref → RIGHT = tmp\_node

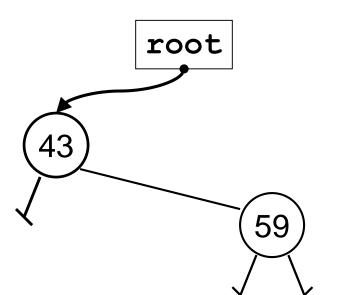








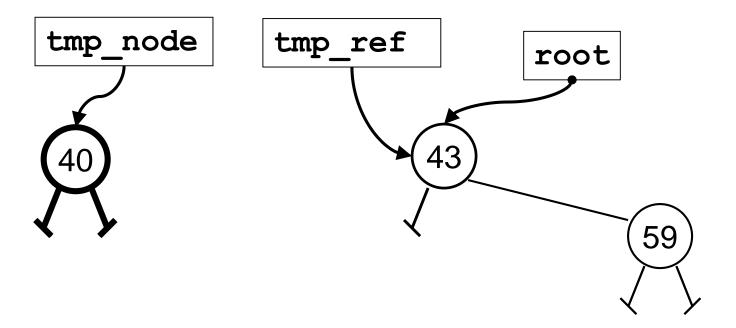






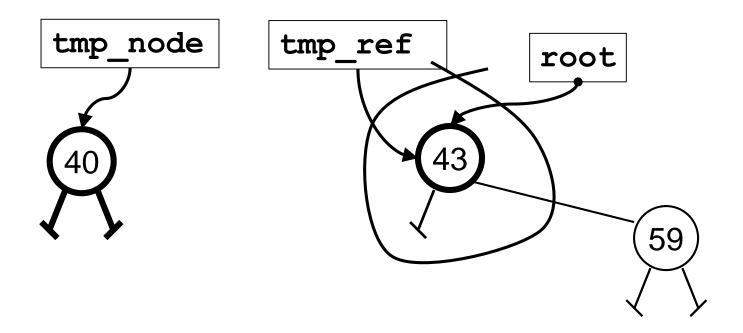
- Allocate New Node: tmp\_node
- Insert Data in New Node: tmp\_node → DATA = 40
- Update Left and Right references (i.e., pointers) to NULL:
  - √ tmp\_node→LEFT = NULL; tmp\_node→RIGHT = NULL

43	59	40	31	64	33	20	56	47	28	89
----	----	----	----	----	----	----	----	----	----	----





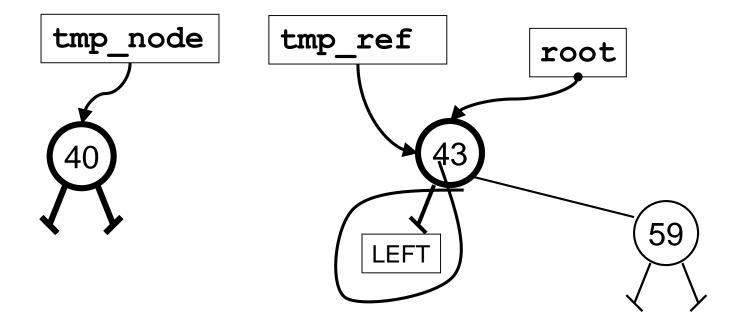
- (...)
- tmp\_ref = root
   i.e., tmp\_ref will reference (i.e., point at)
   the location in the BST to insert the new node





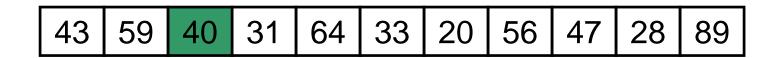
- (...)
- tmp\_node → DATA < tmp\_ref → DATA? Insert @ tmp\_ref → LEFT
- tmp\_node → DATA > tmp\_ref → DATA? Insert @ tmp\_ref → RIGHT

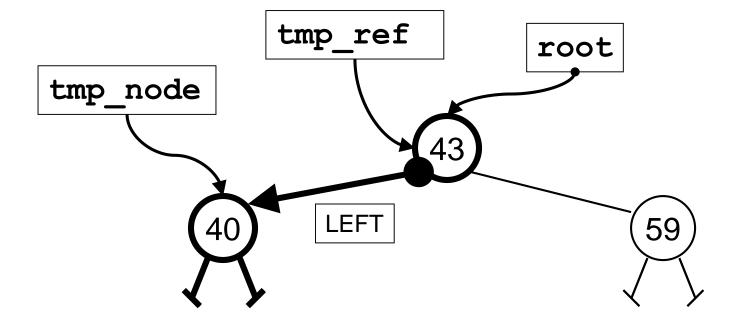
43	59	40	31	64	33	20	56	47	28	89
----	----	----	----	----	----	----	----	----	----	----





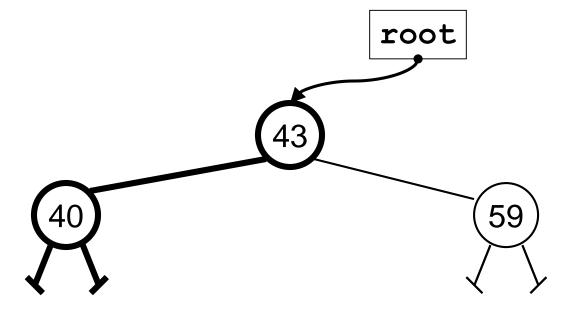
- (...
- tmp\_node → 40 < tmp\_ref → 43? YES: Insert @ tmp\_ref → LEFT





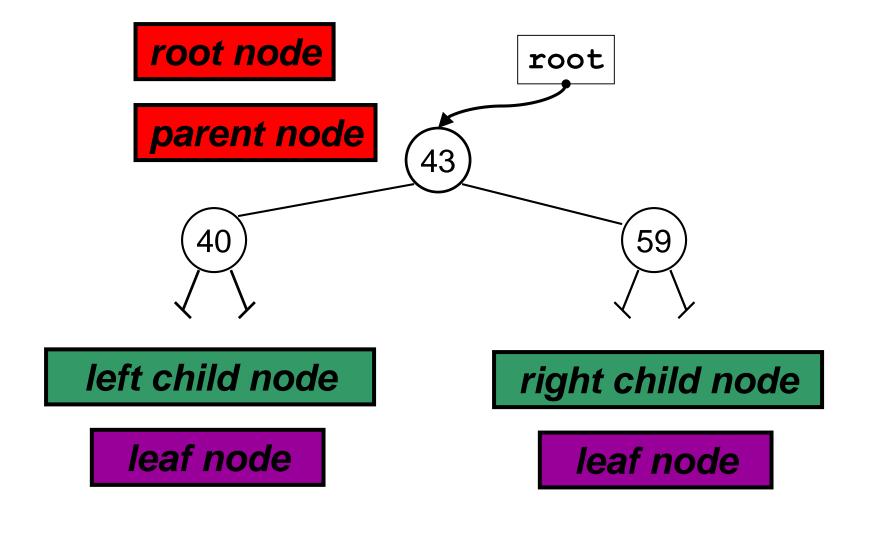


- (...)
- tmp\_node → 40 < tmp\_ref → 43? YES: Insert @ tmp\_ref → LEFT</li>
   ✓ tmp\_ref → LEFT = tmp\_node



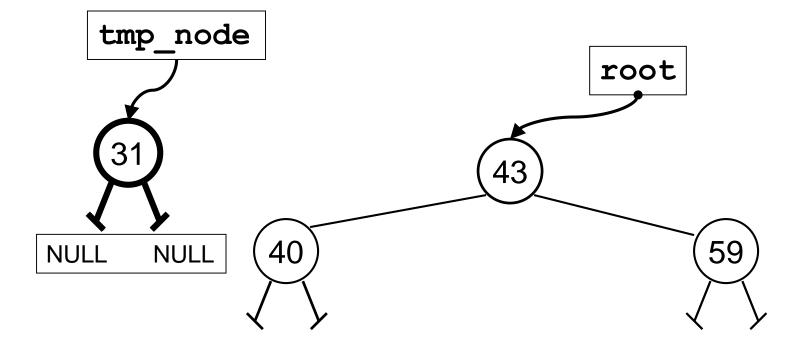






CALGARY

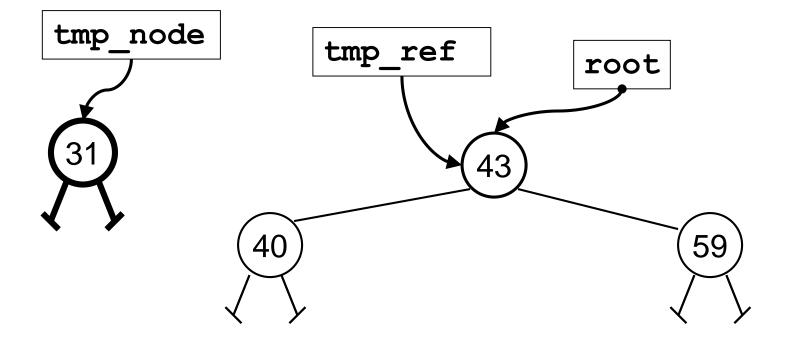






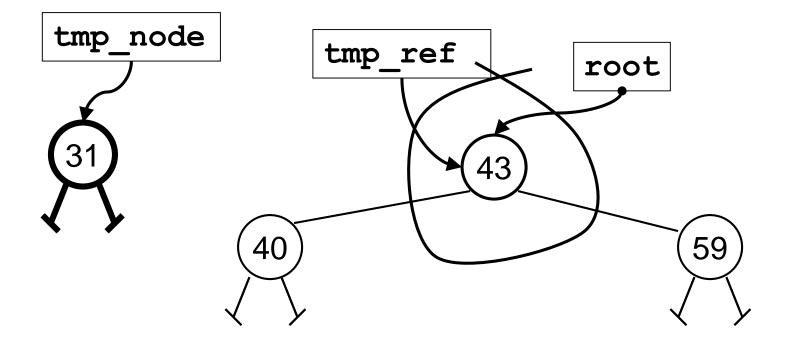
- Allocate New Node: tmp\_node
- Insert Data in New Node: tmp\_node → DATA = 31
- Update Left and Right references (i.e., pointers) to NULL:
  - ✓ tmp\_node→LEFT = NULL; tmp\_node→RIGHT = NULL

43   59   40   <mark>31  </mark> 64   33   20   56   47   28   8
--



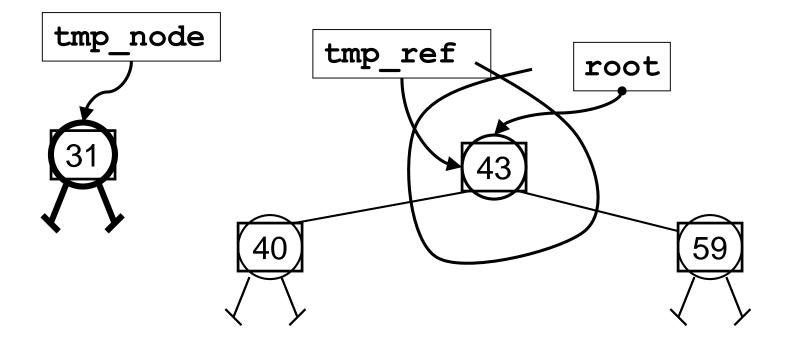


- (...)
- tmp\_ref = root
   i.e., tmp\_ref will reference (i.e., point at)
   the location in the BST to insert the new node



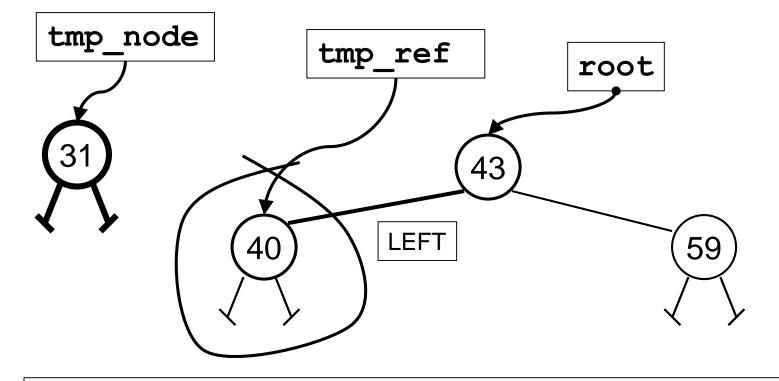


- (...)
- tmp\_node → DATA < tmp\_ref → DATA? Insert @ tmp\_ref → LEFT
- tmp\_node → DATA > tmp\_ref → DATA? Insert @ tmp\_ref → RIGHT



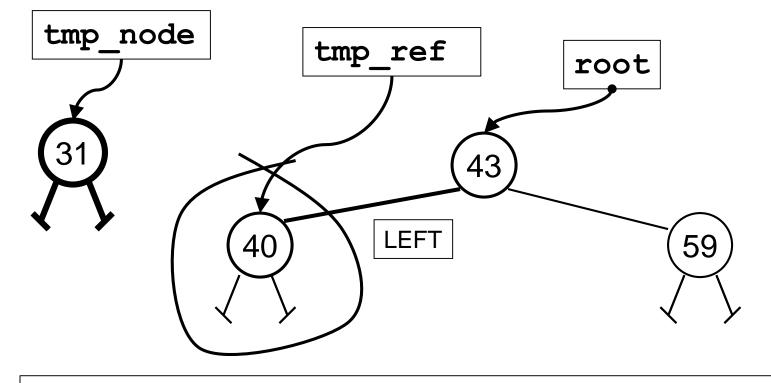


- (...)
- tmp\_node → 31 < tmp\_ref → 43? YES: Insert @ tmp\_ref → LEFT
  - ✓ tmp\_ref→LEFT = NULL? // i.e., Leaf Node?
  - ✓ NO: tmp\_ref = tmp\_ref → LEFT



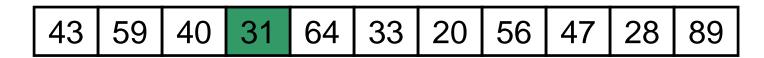


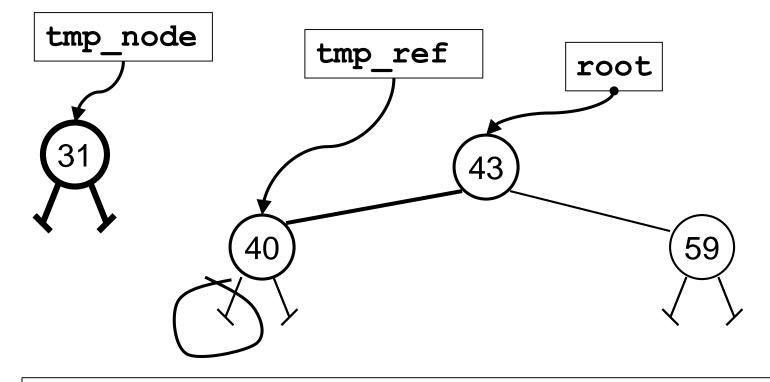
- (...)
- tmp\_node → 31 < tmp\_ref → 43? YES: Insert @ tmp\_ref → LEFT
  - ✓ tmp\_ref→LEFT = NULL? // i.e., Leaf Node?
  - ✓ NO: tmp\_ref = tmp\_ref → LEFT





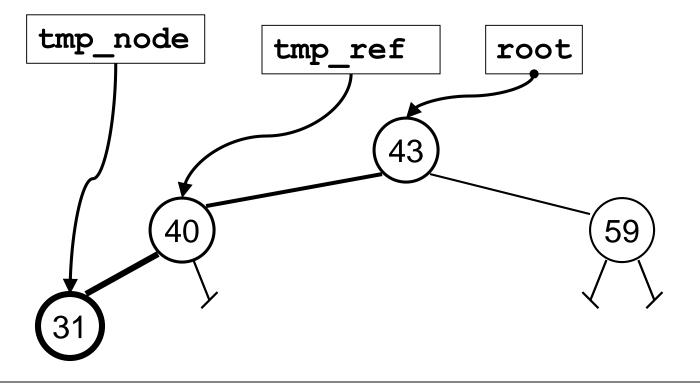
- (...)
- tmp\_node → 31 < tmp\_ref → 40? YES: Insert @ tmp\_ref → LEFT





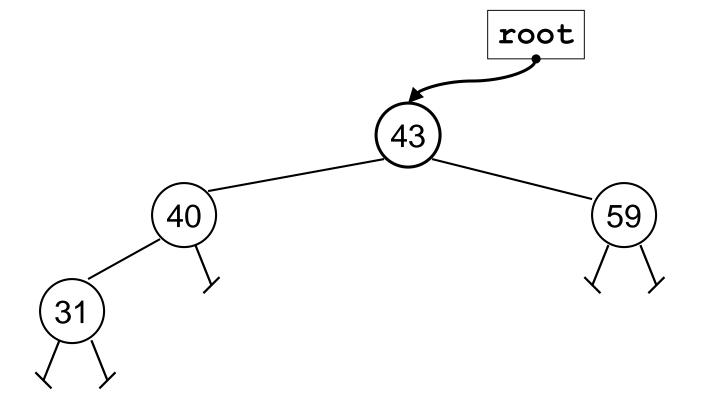


- (...)
- tmp\_node → 31 < tmp\_ref → 40? YES: Insert @ tmp\_ref → LEFT
  - ✓ tmp\_ref→LEFT = NULL? // i.e., Leaf Node?
  - ✓ YES: tmp\_ref = tmp\_ref → LEFT



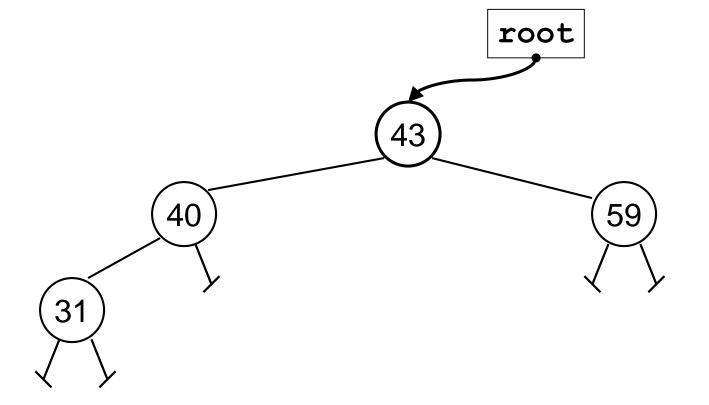


- (...)
- tmp\_node → 31 < tmp\_ref → 40? YES: Insert @ tmp\_ref → LEFT</li>
   ✓ tmp\_ref → LEFT = tmp\_node



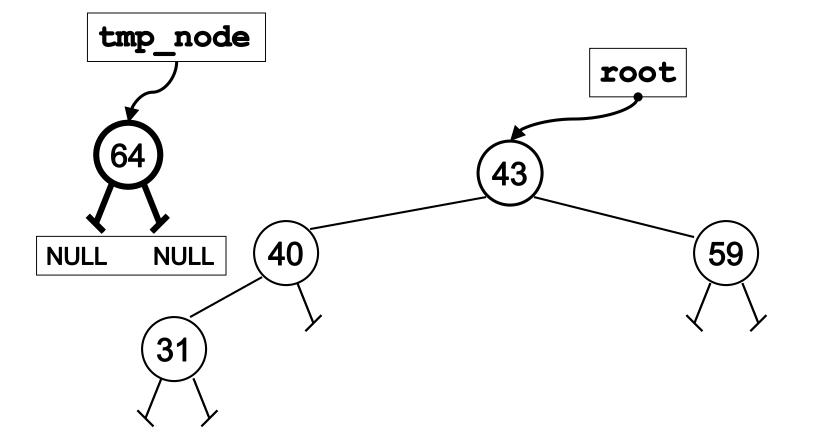






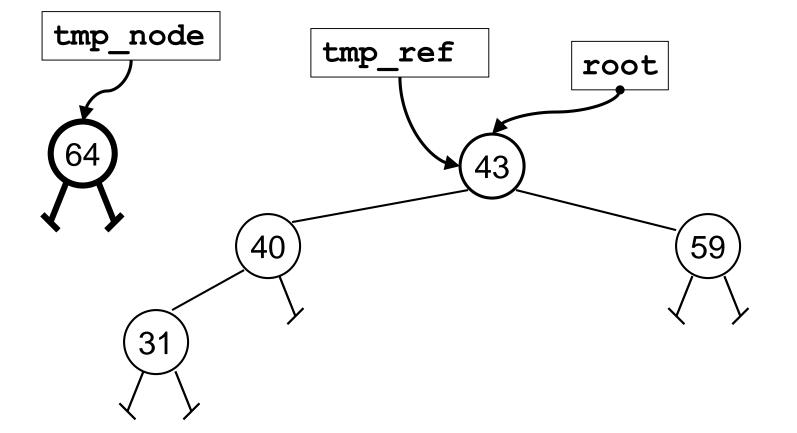






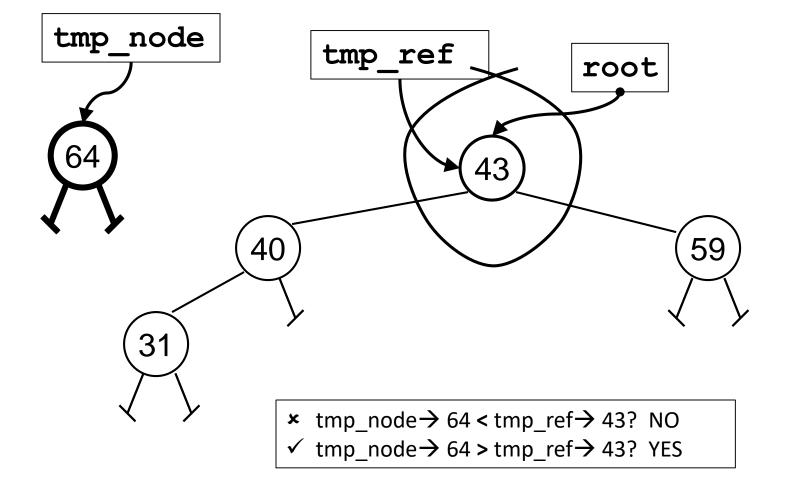






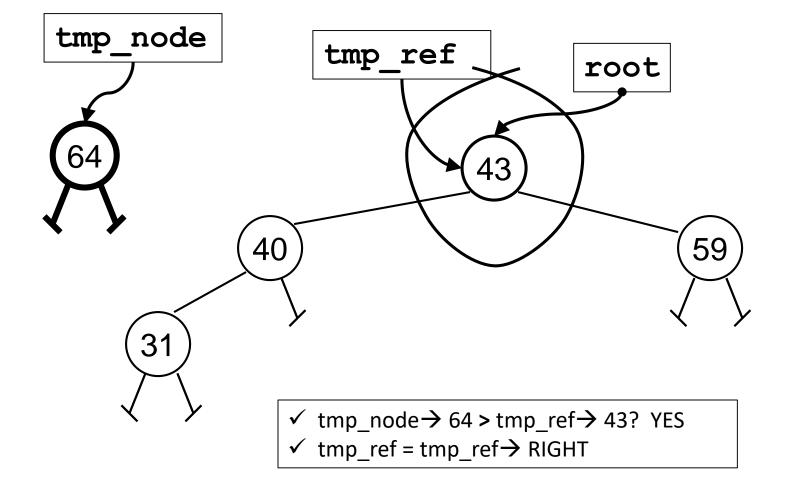




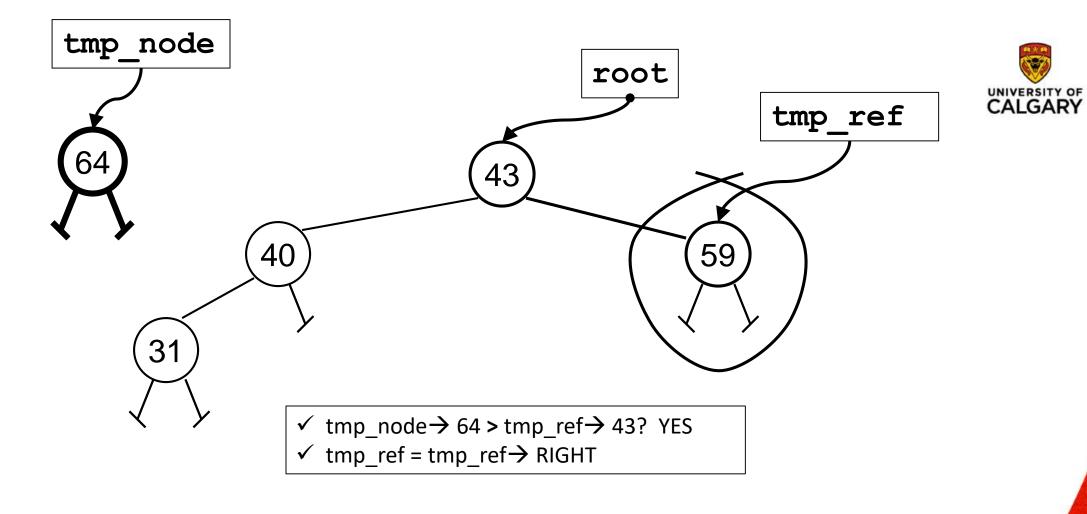


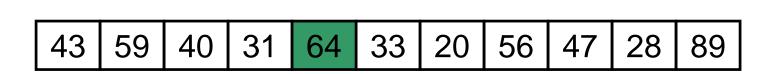


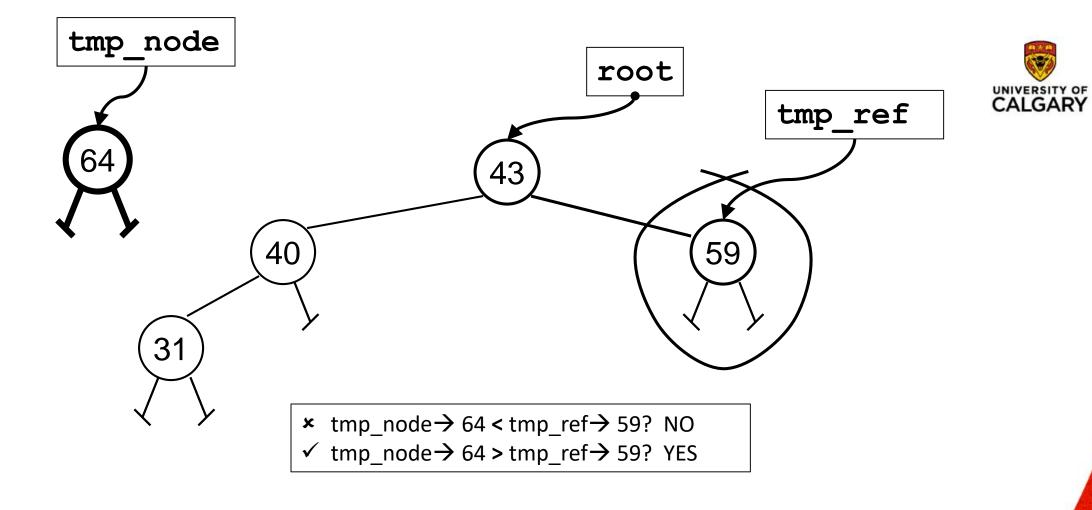
43   59   40   31   64   33   20   56   47   28   89
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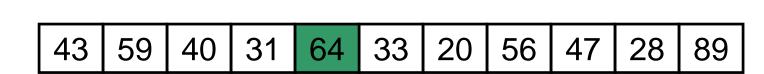


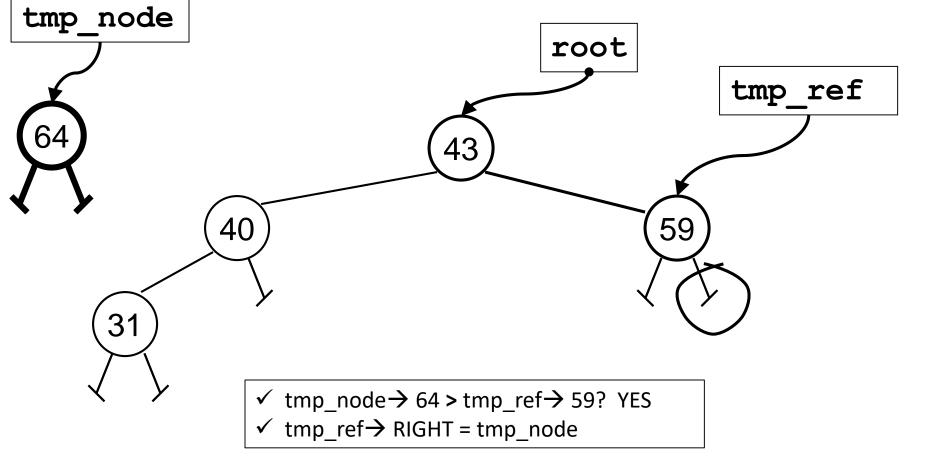






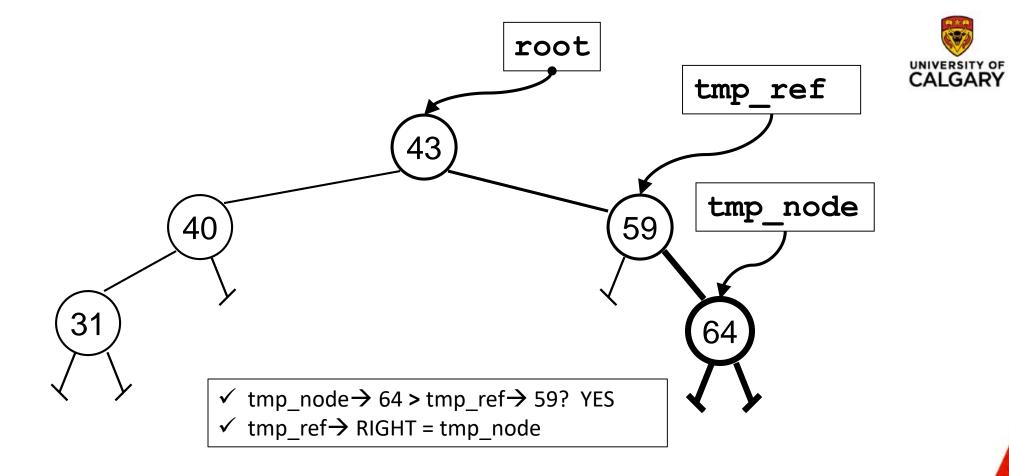


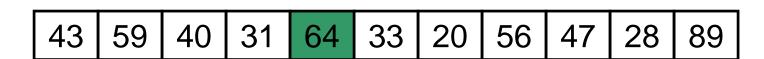


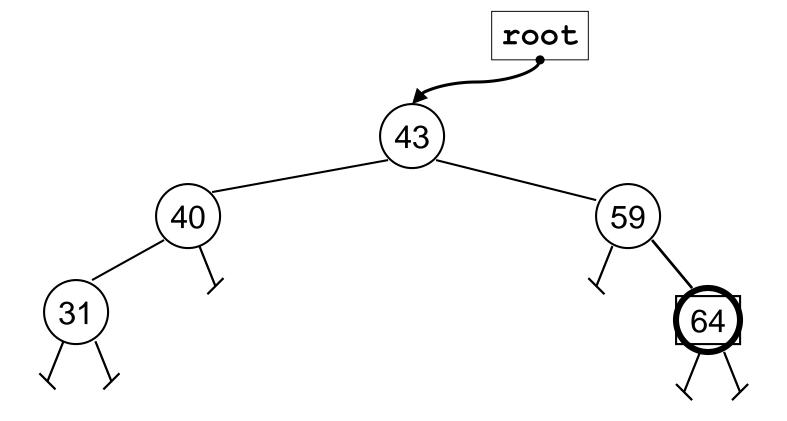




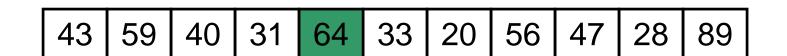
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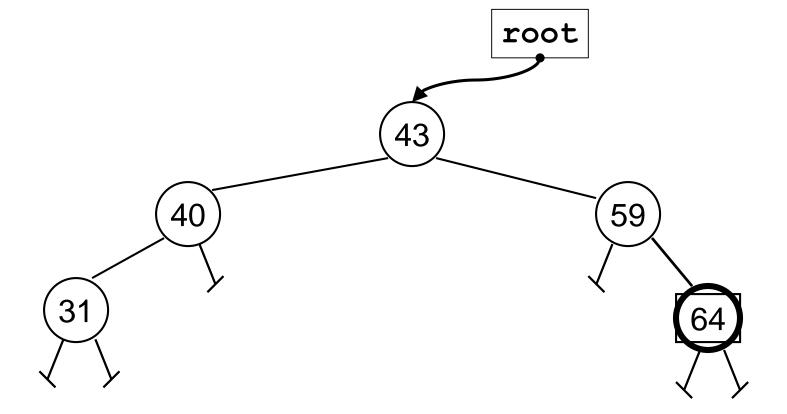






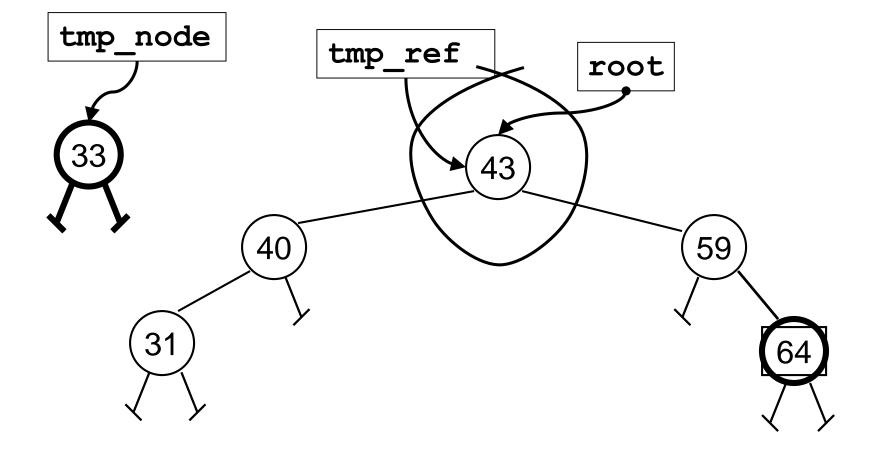






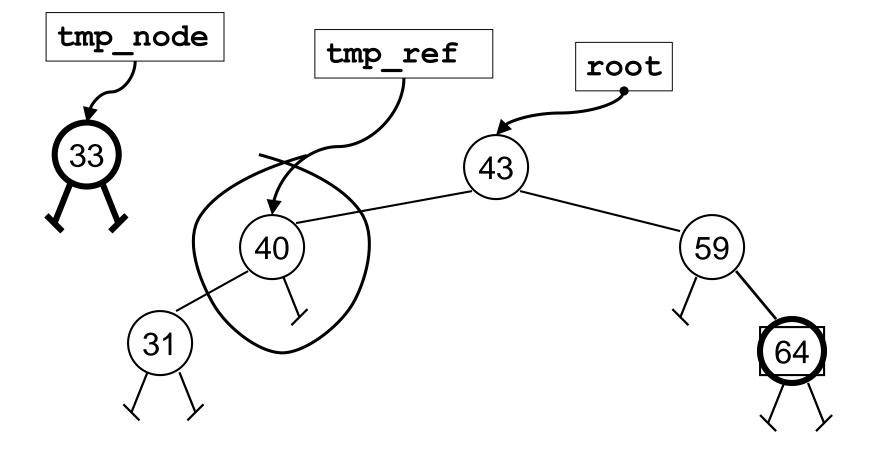






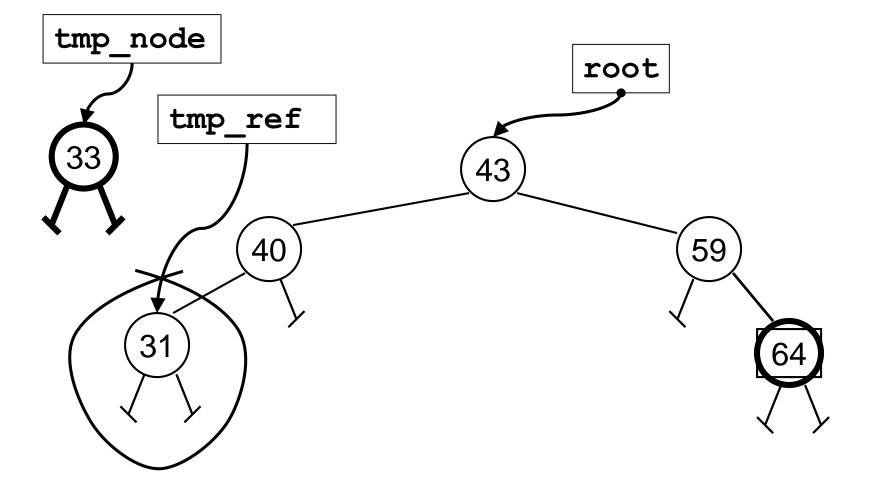






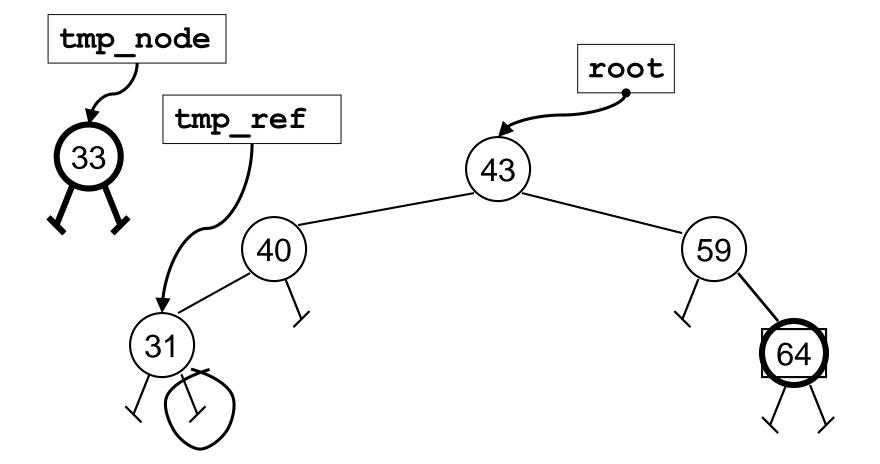






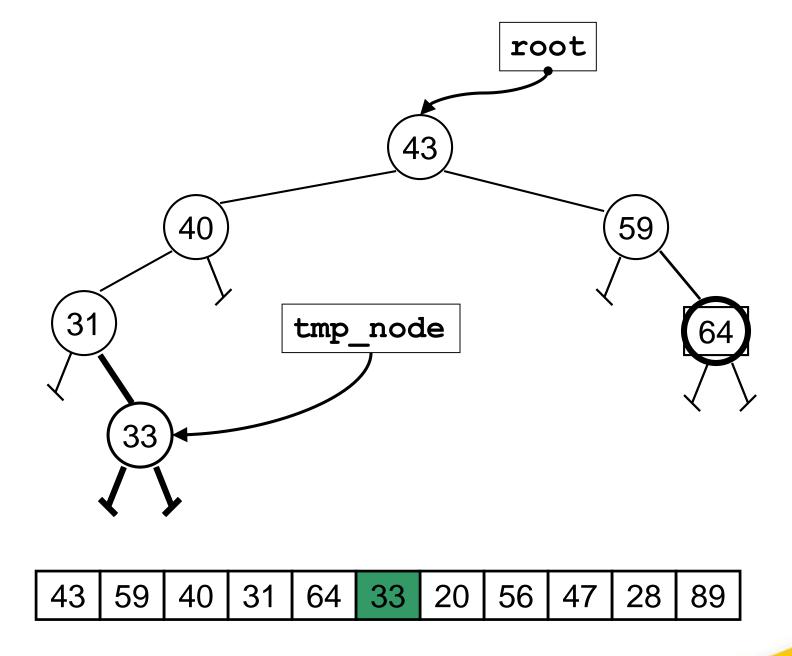




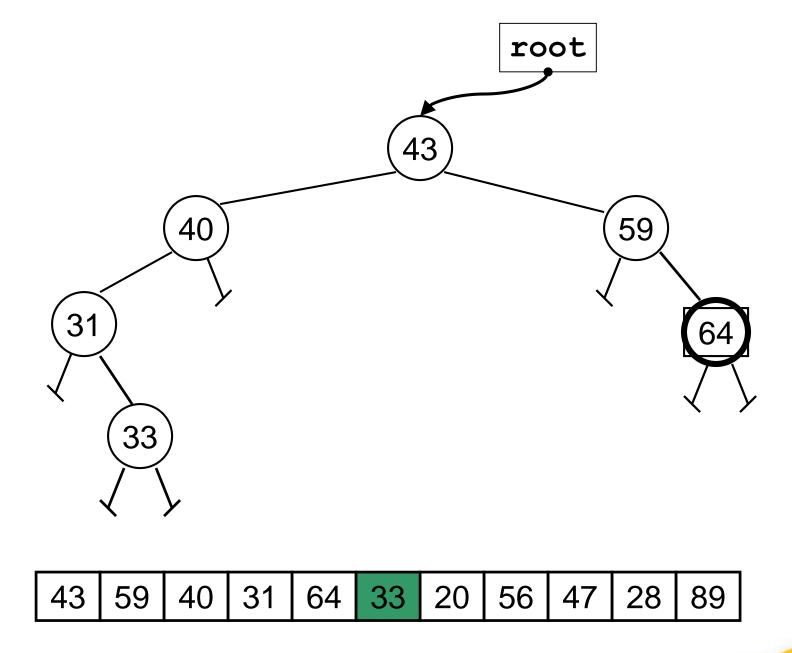




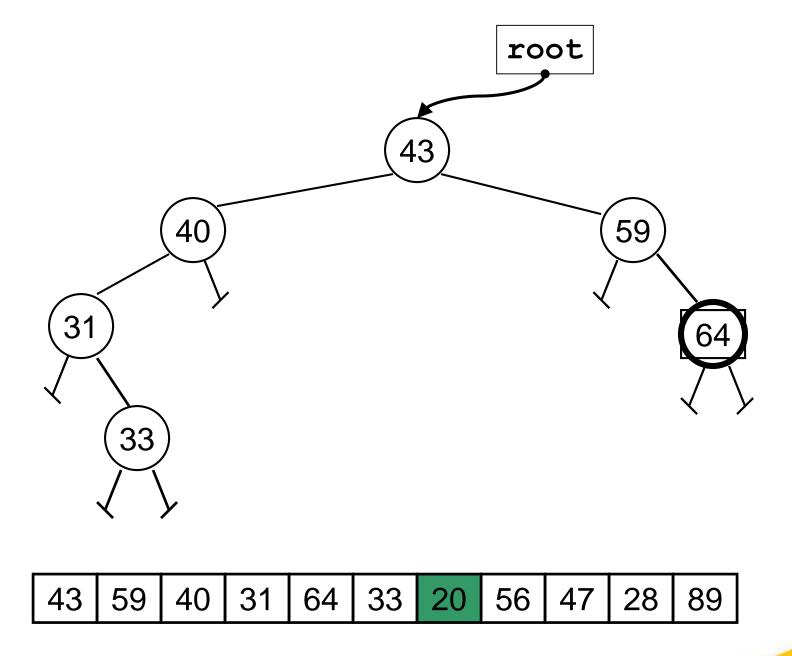




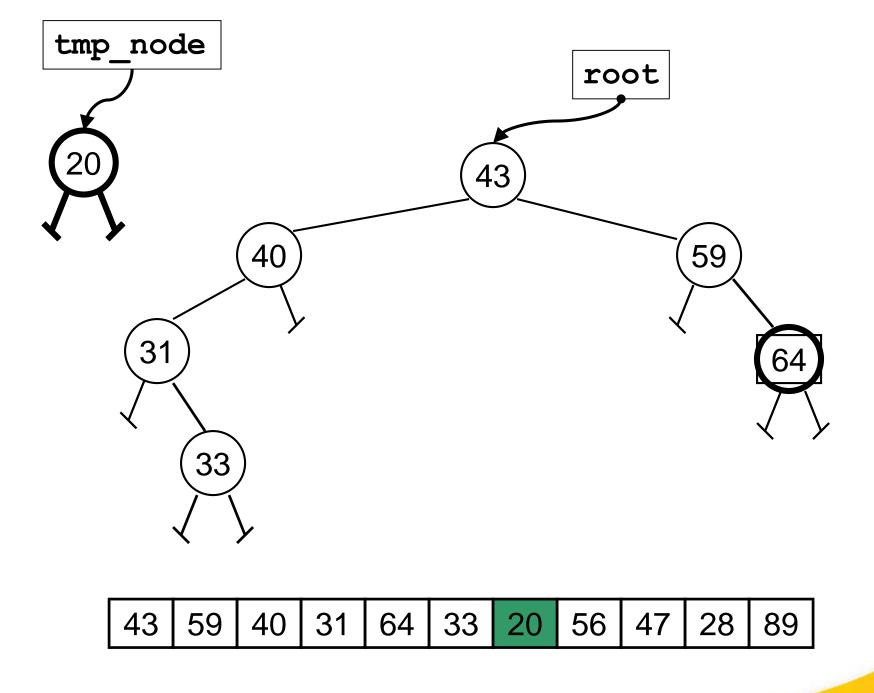




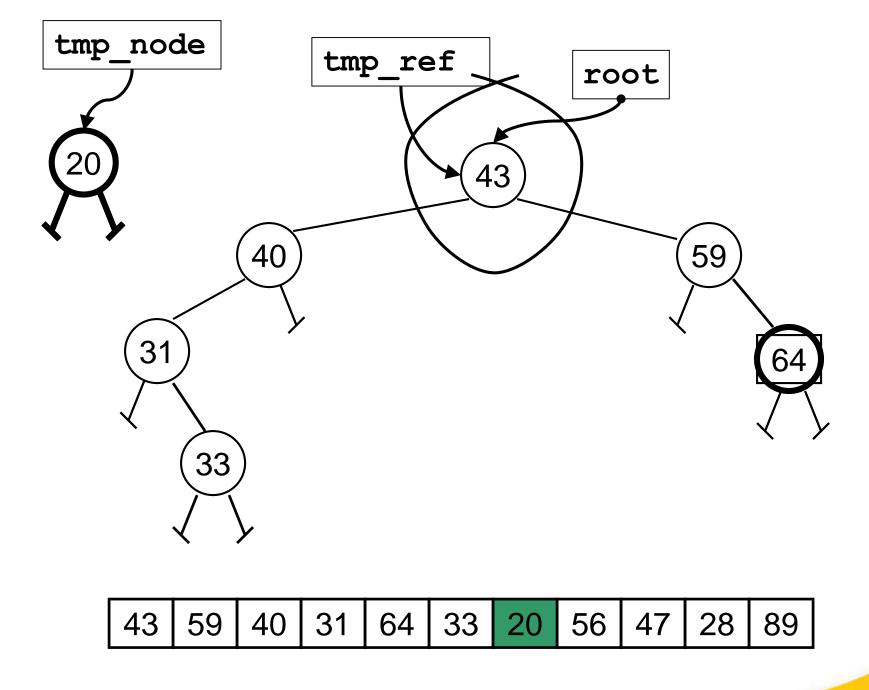




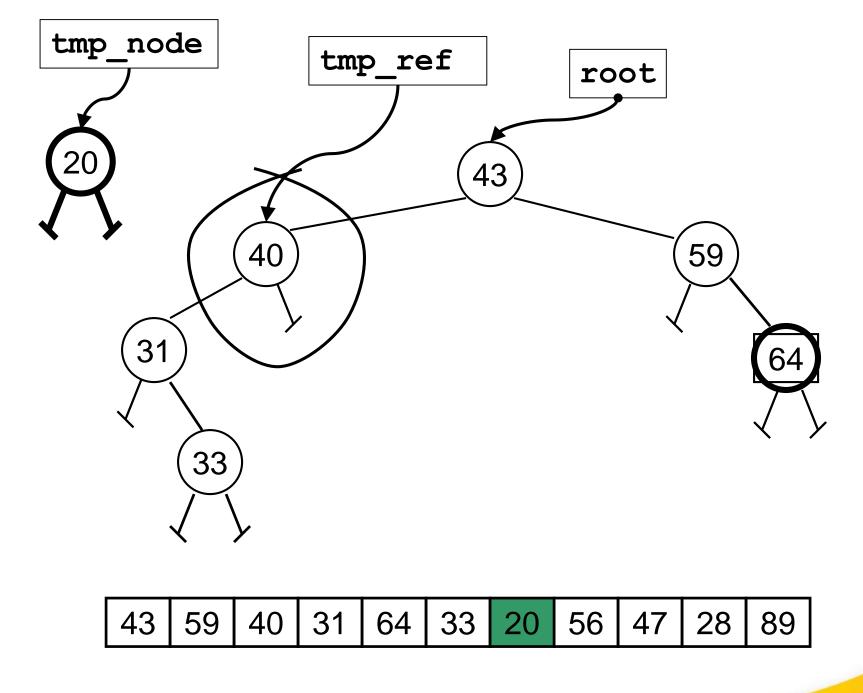




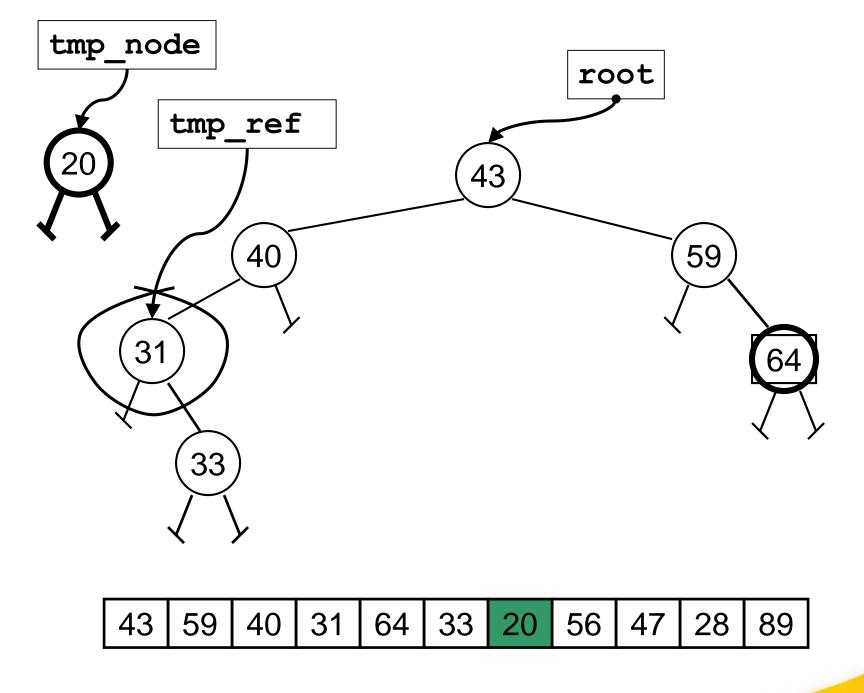




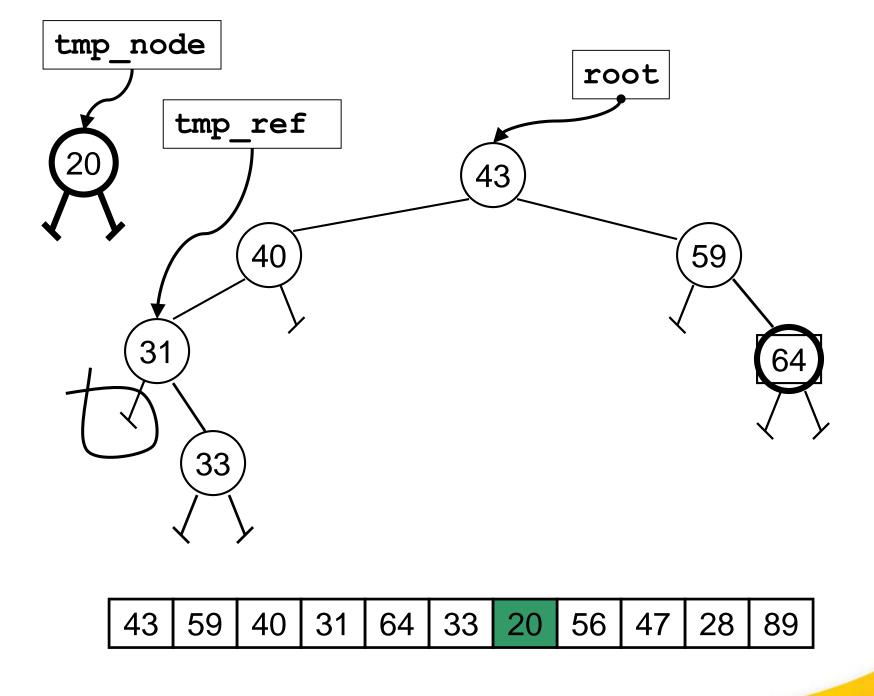




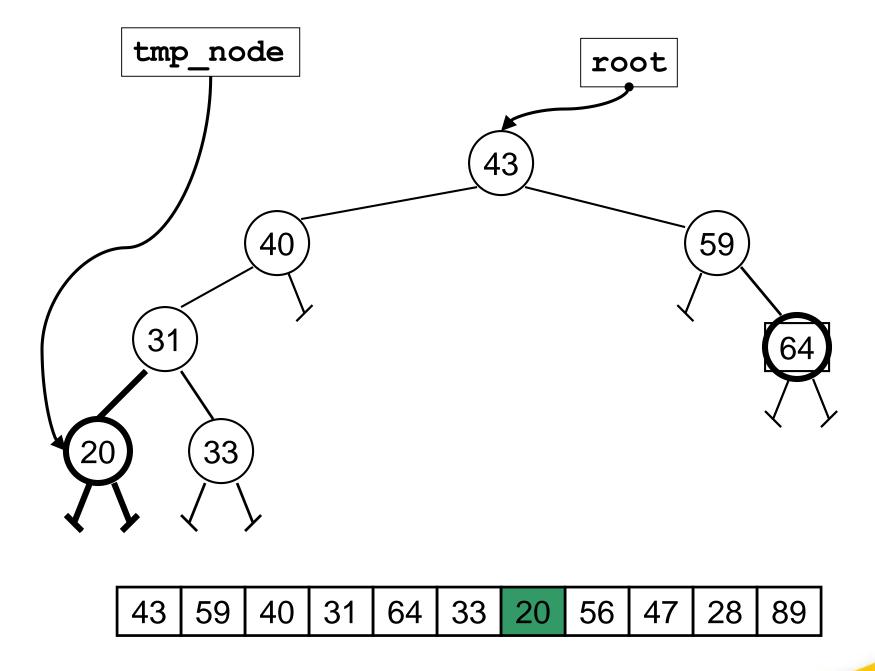




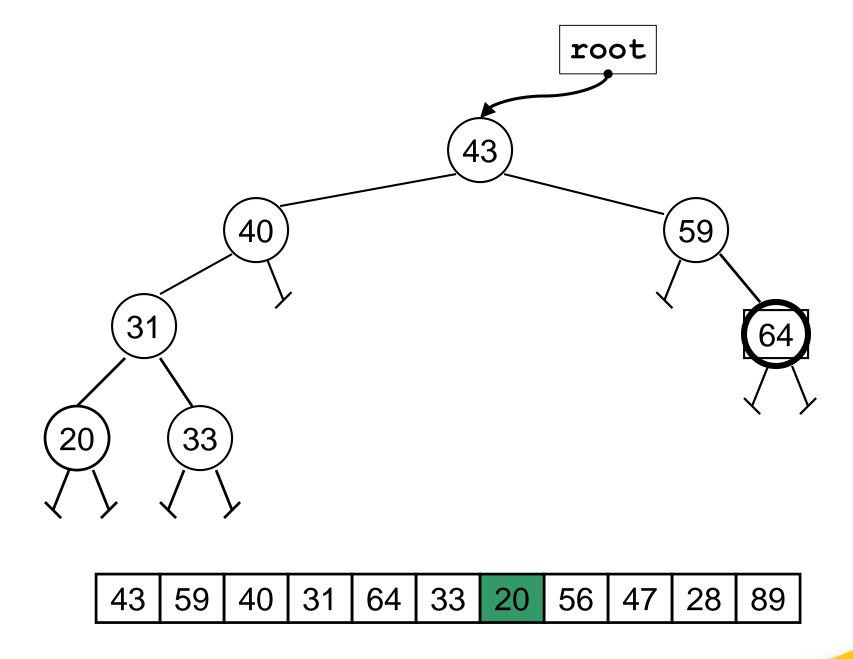




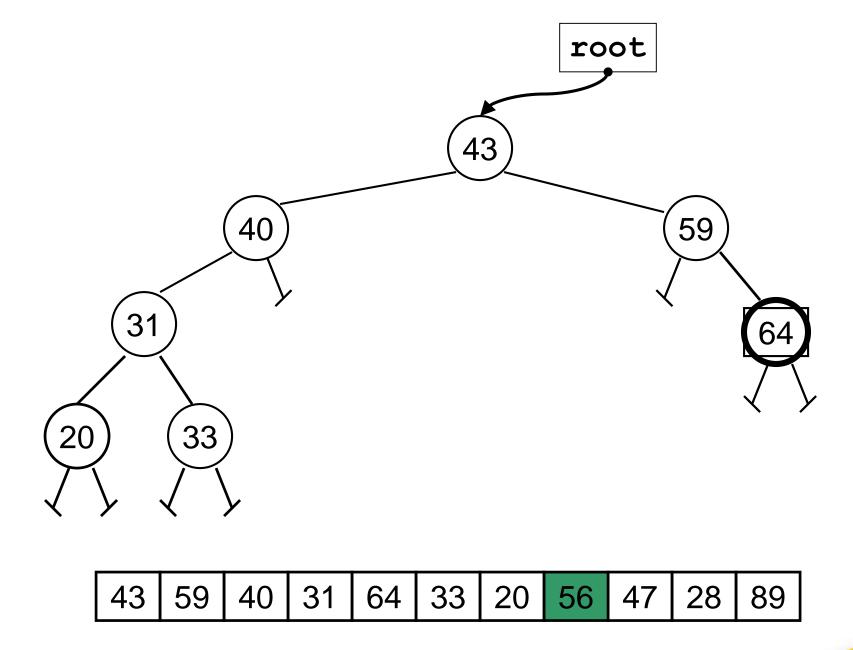




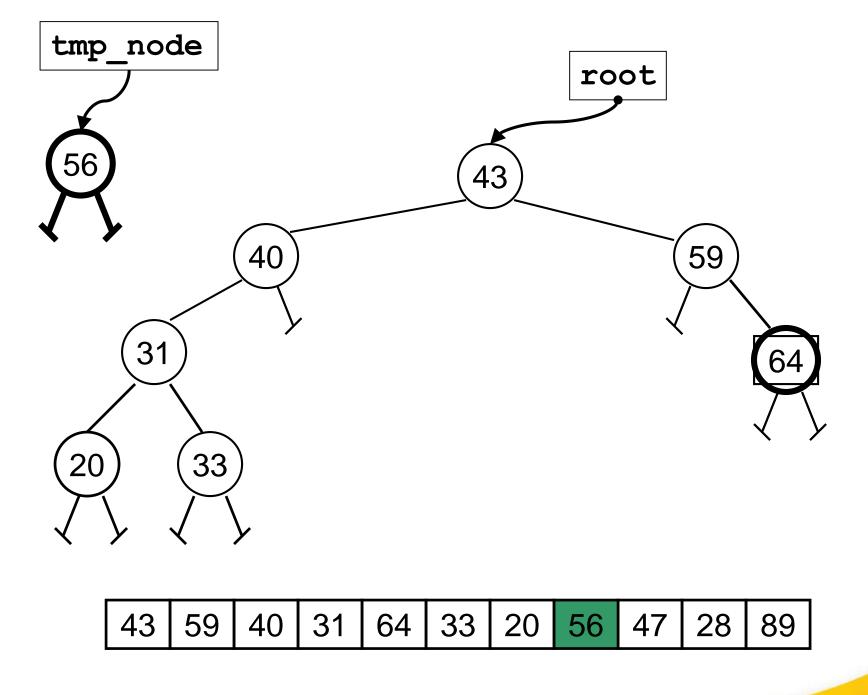




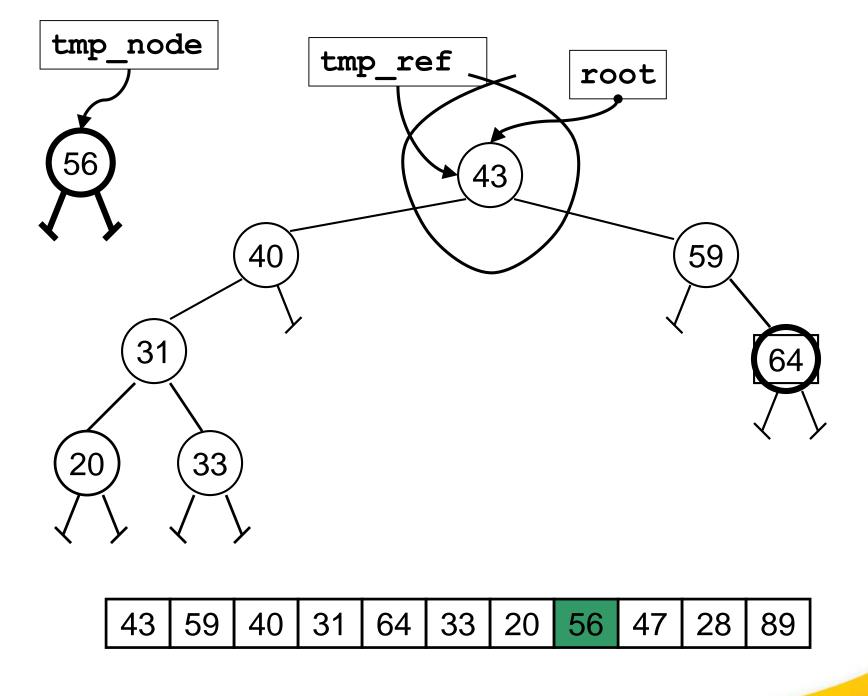




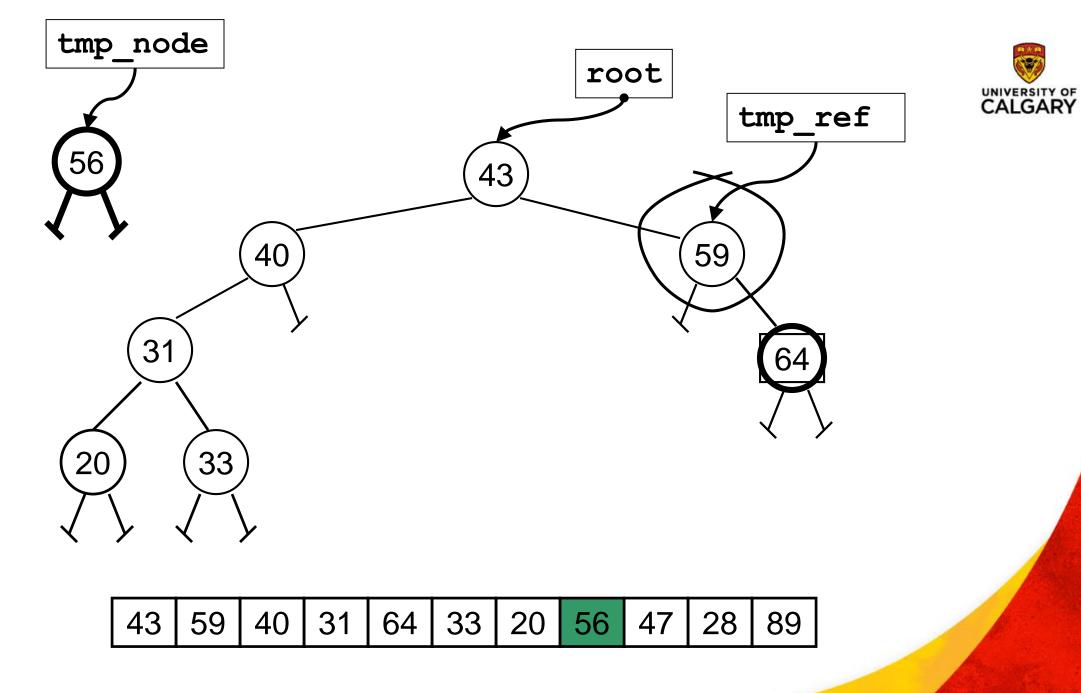


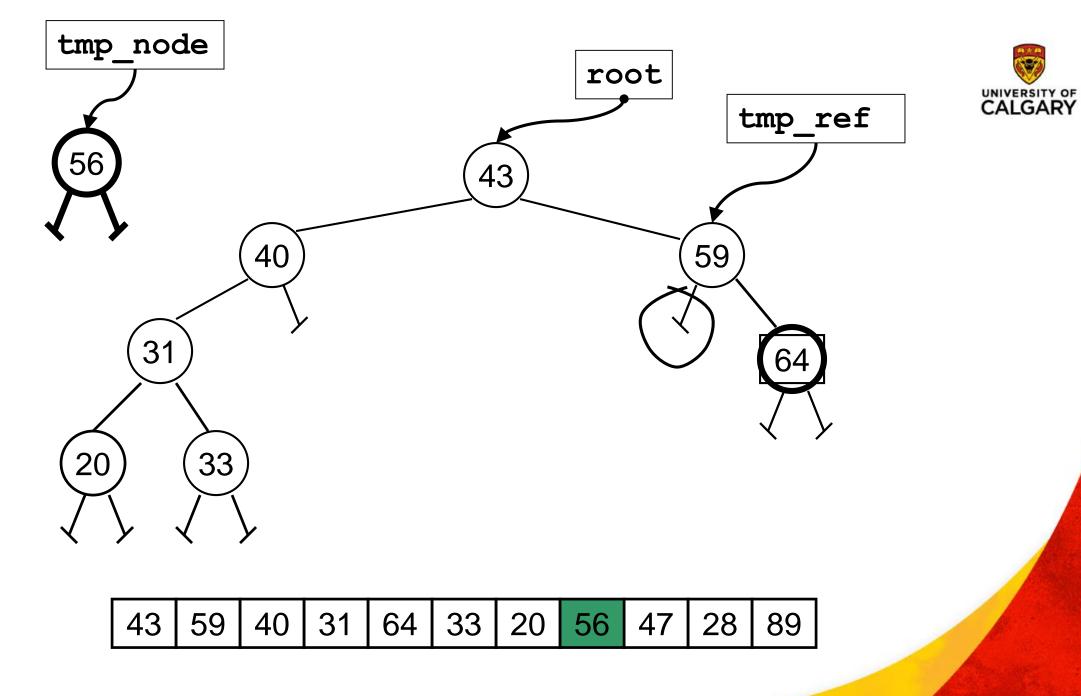


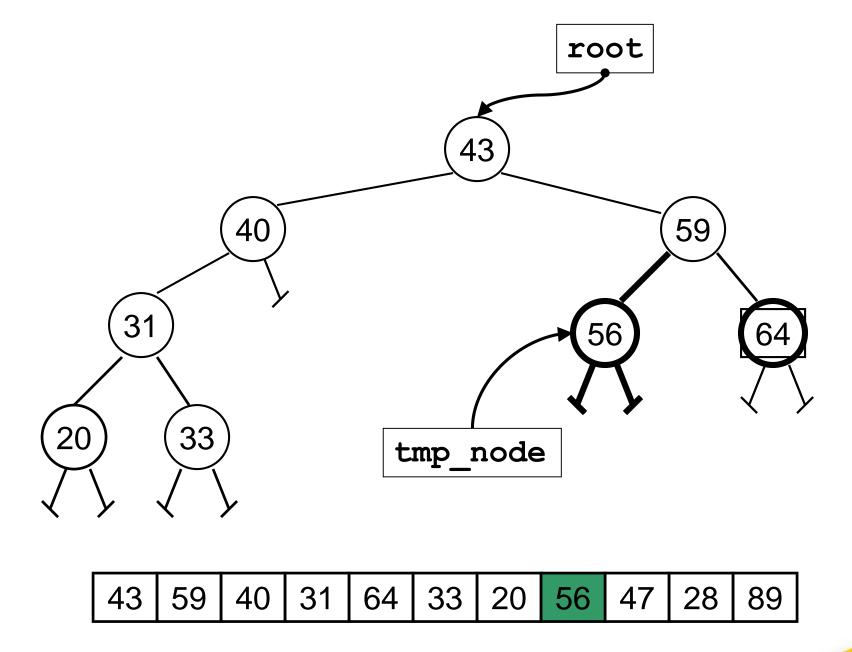




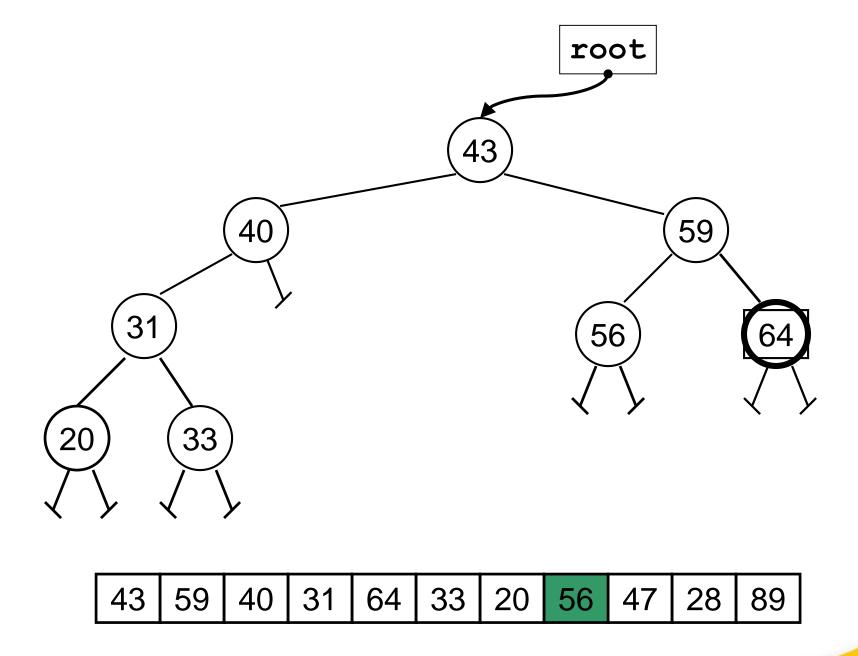




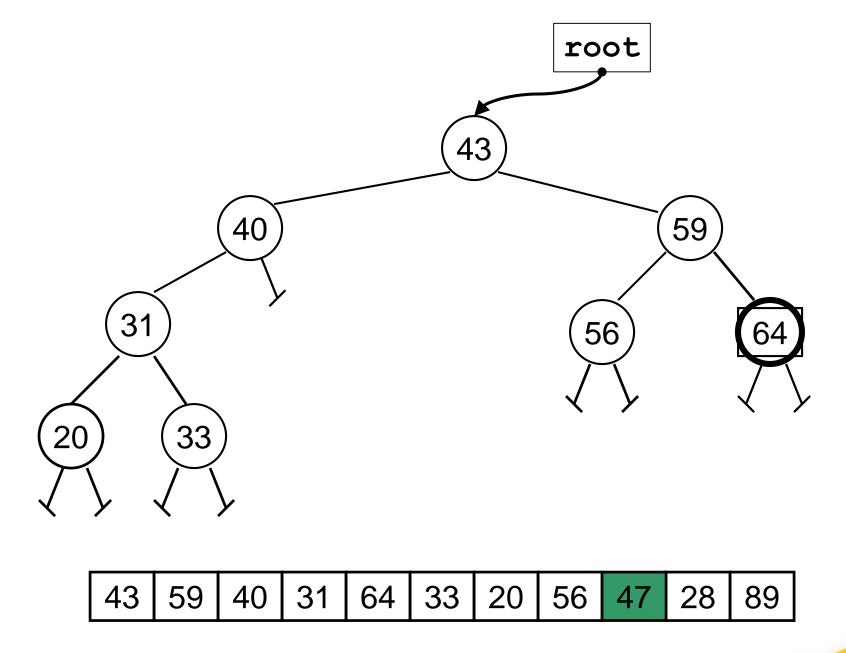




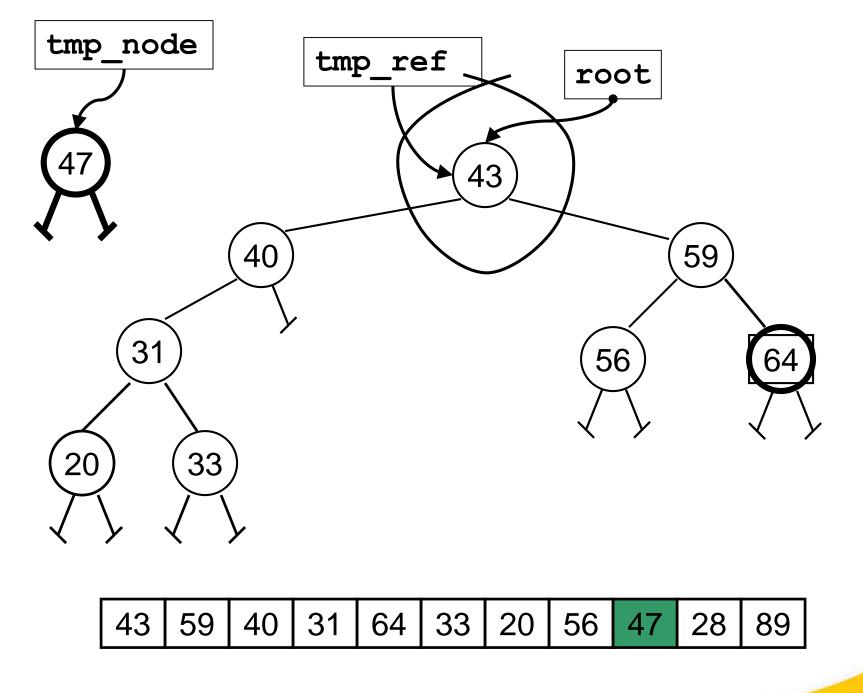




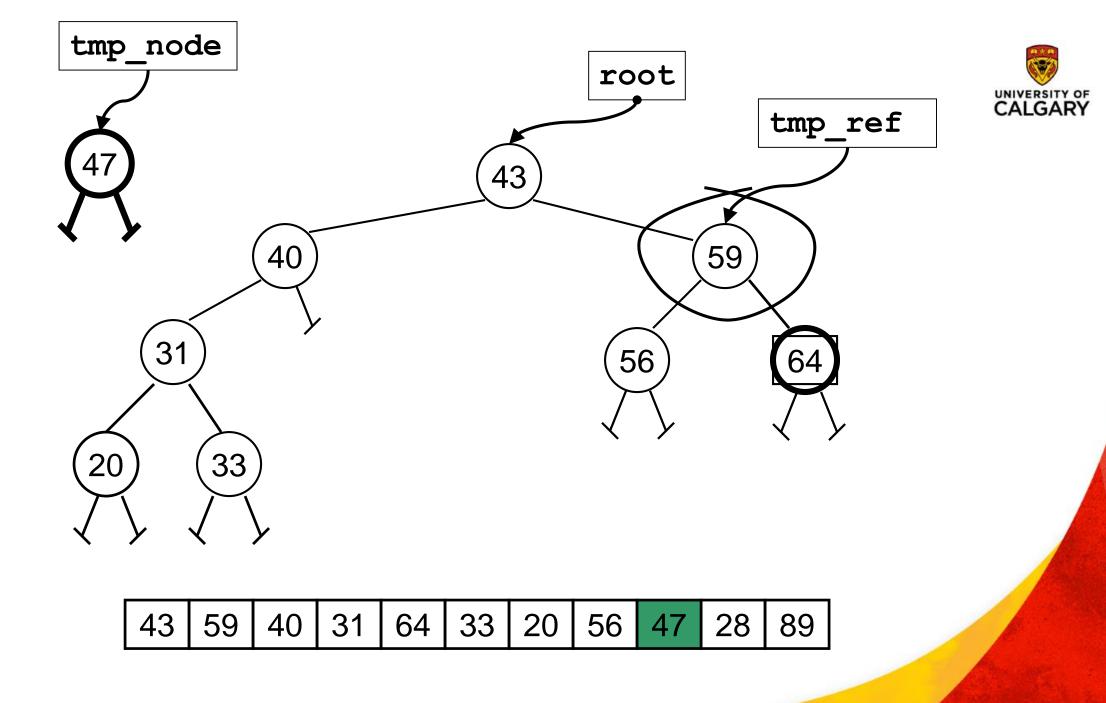


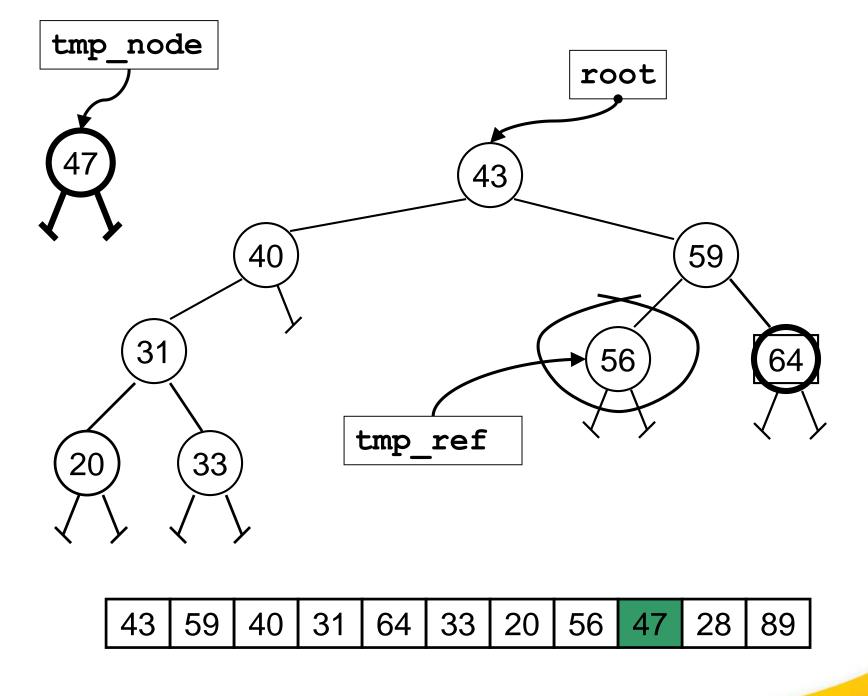




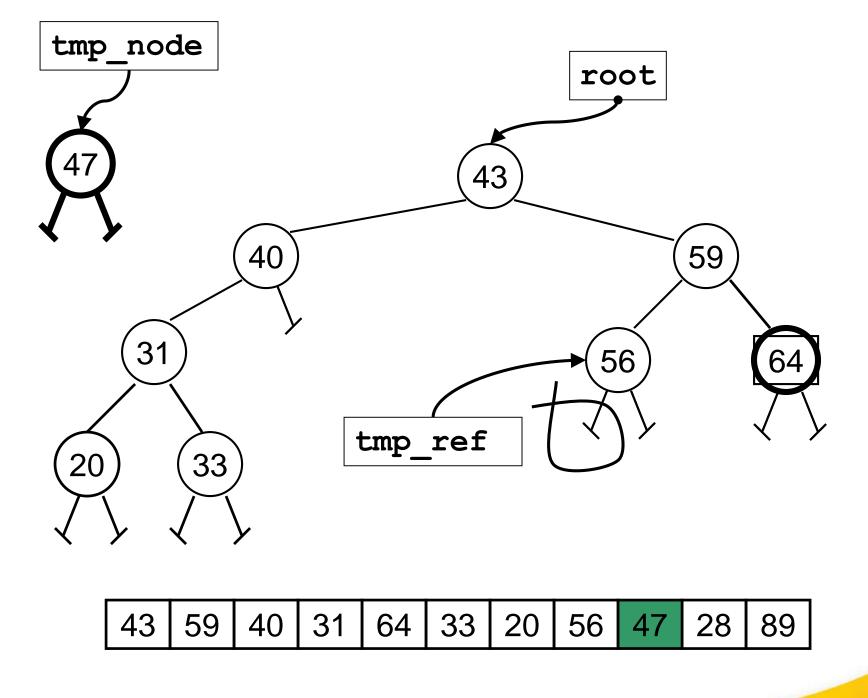




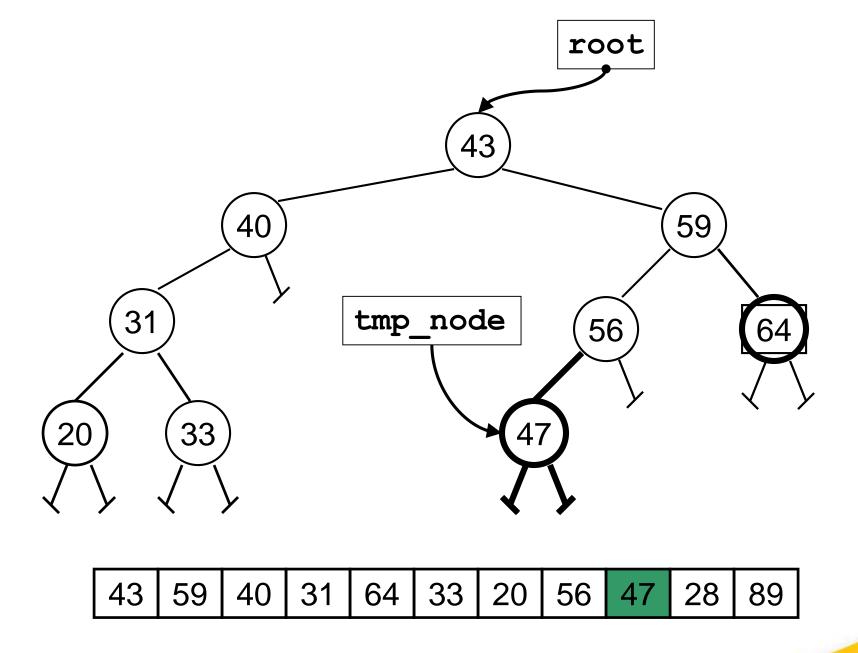




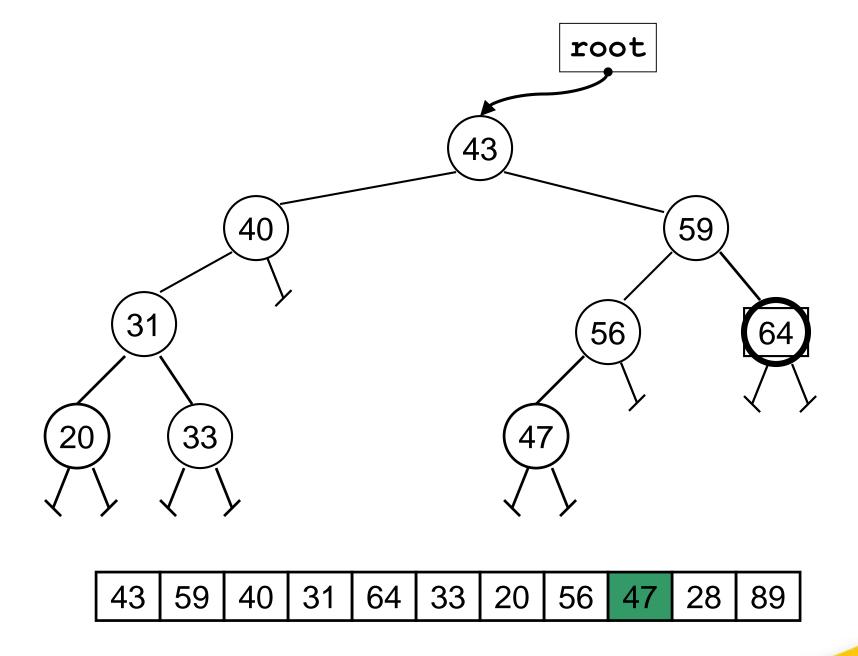




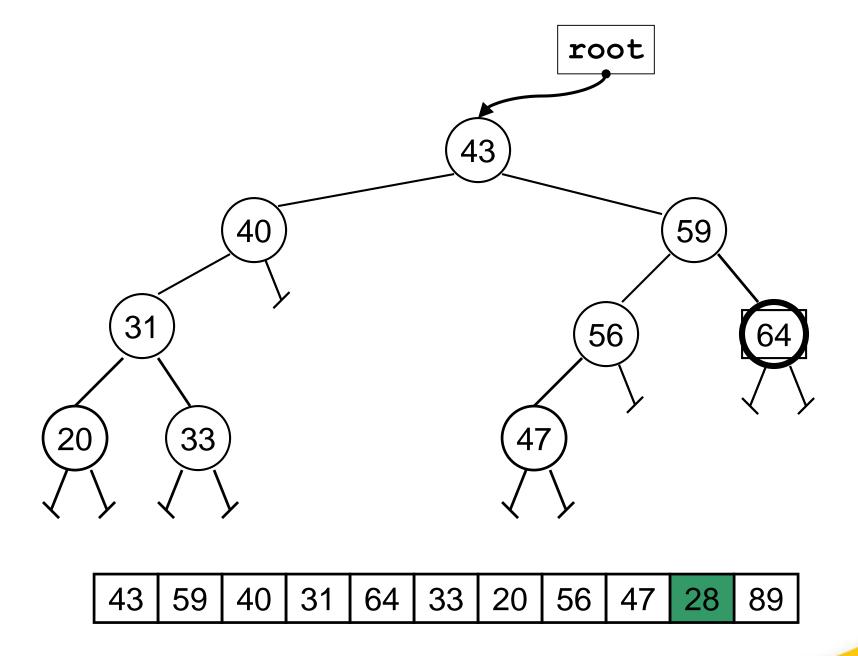




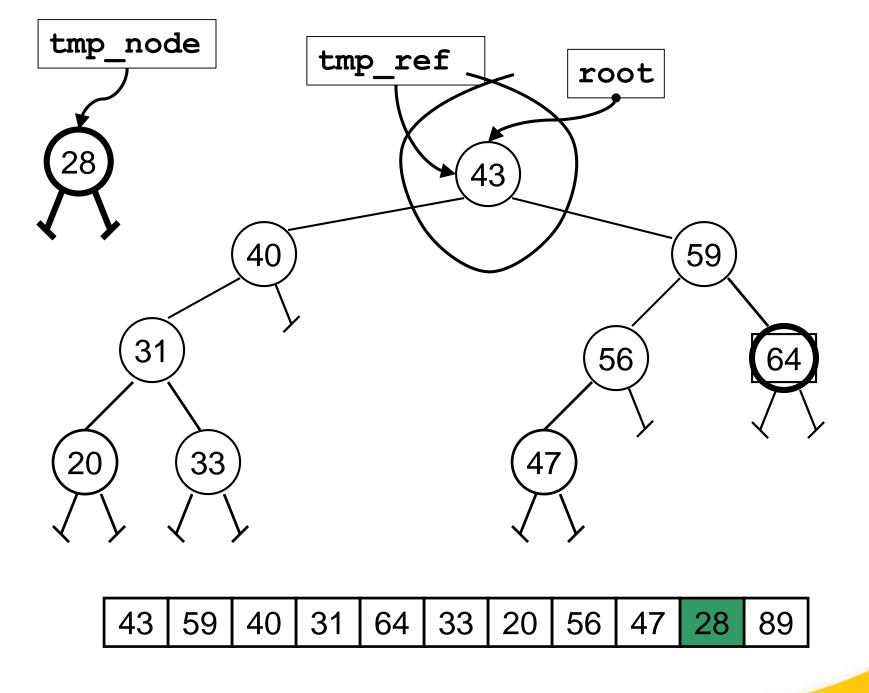




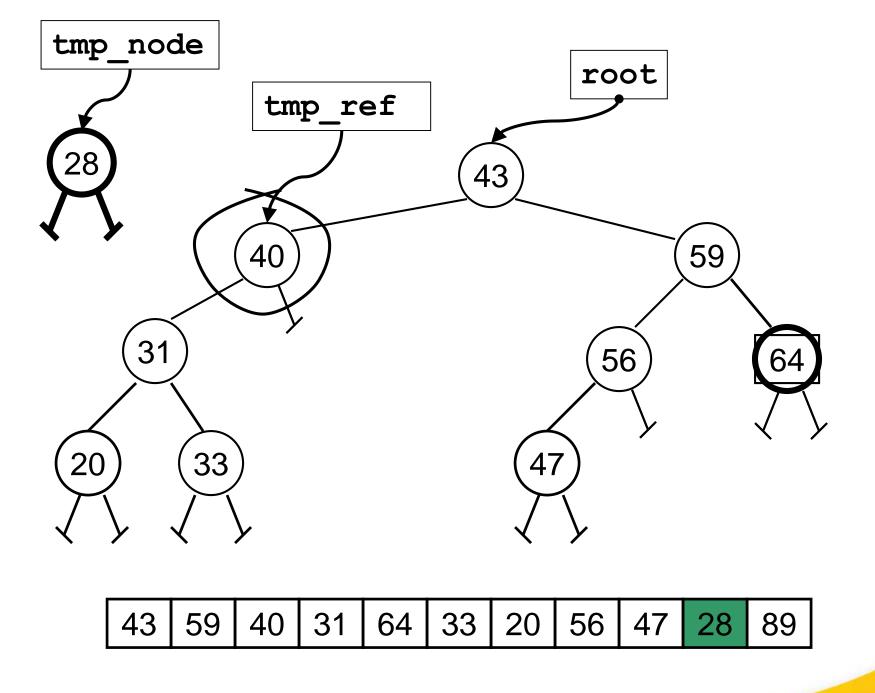




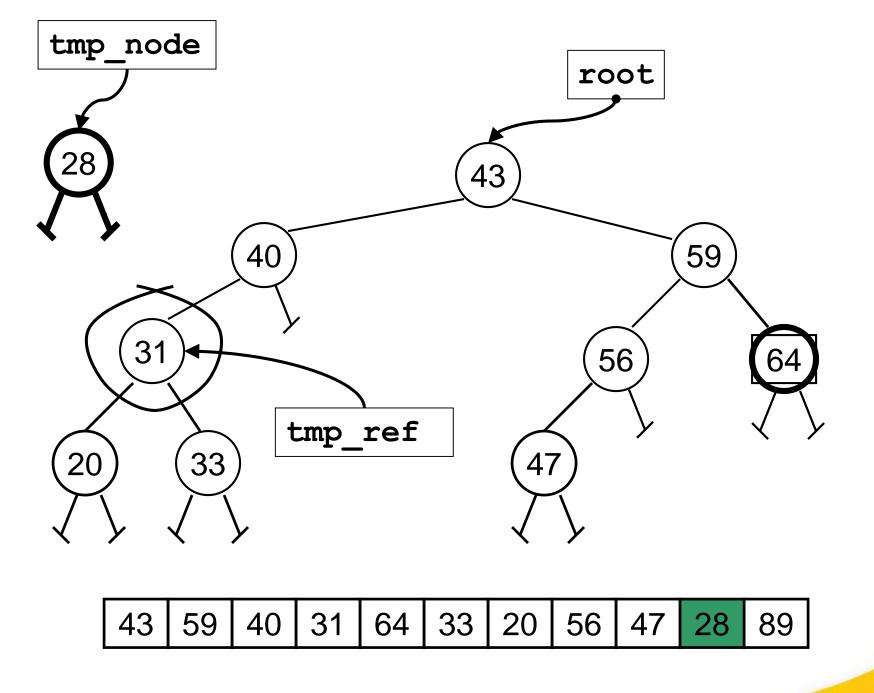




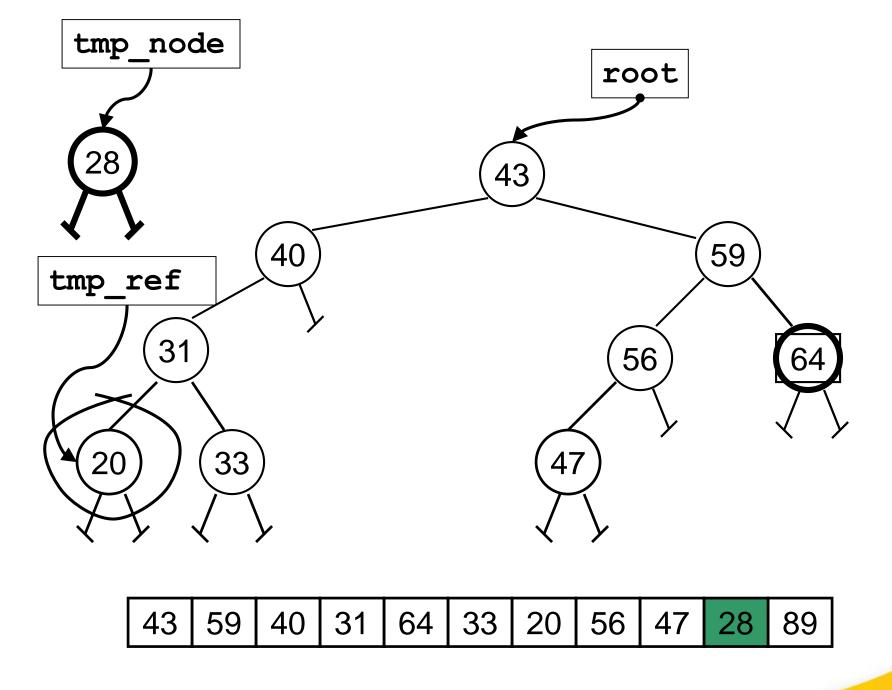




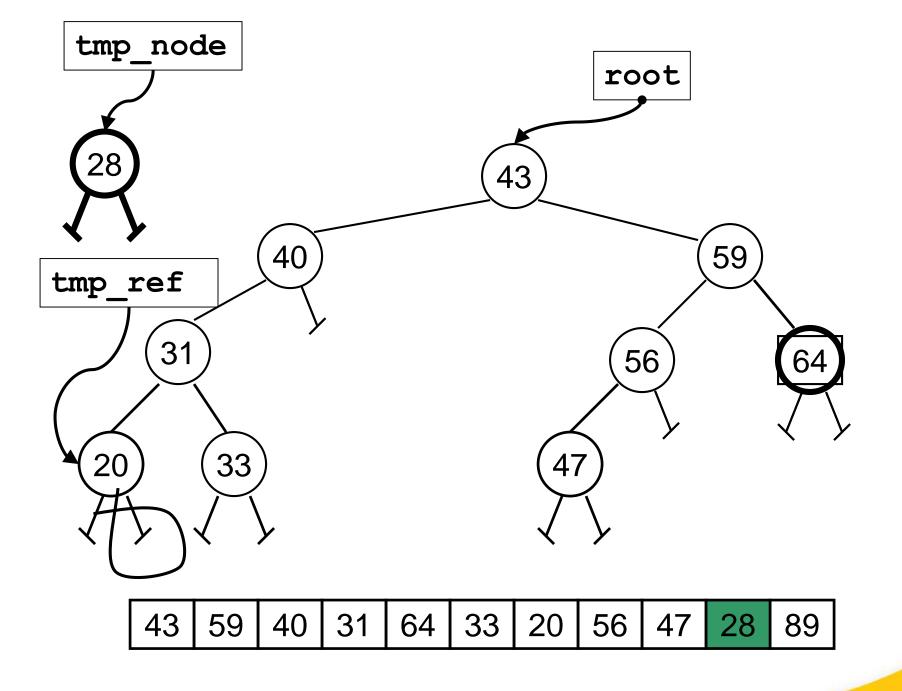




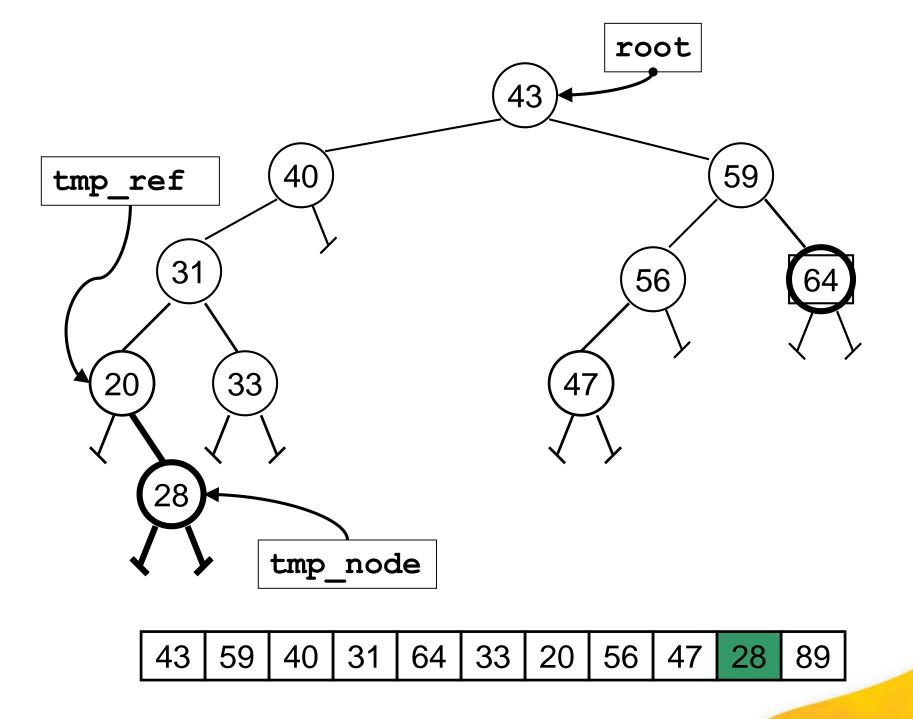




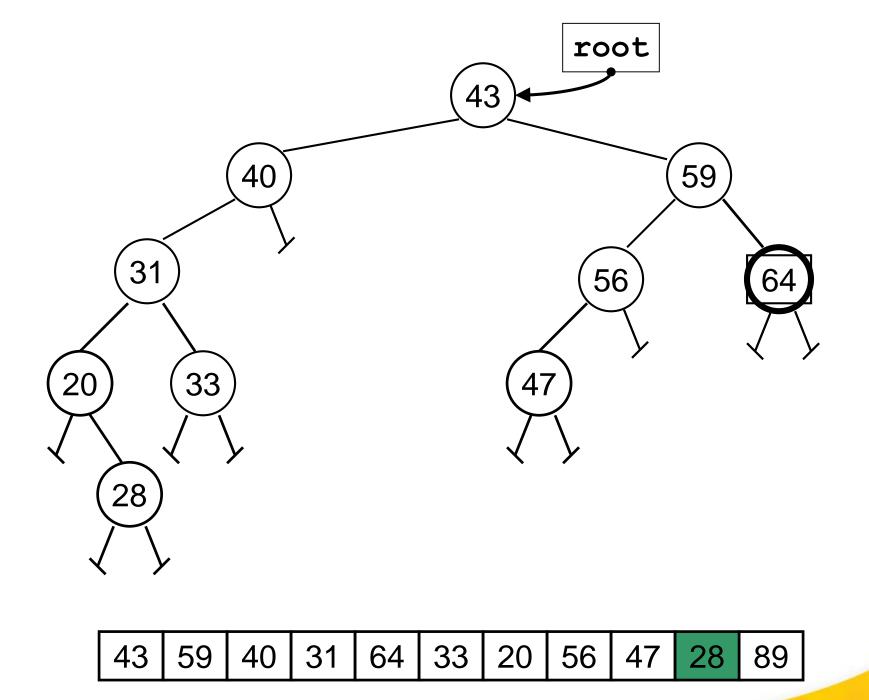




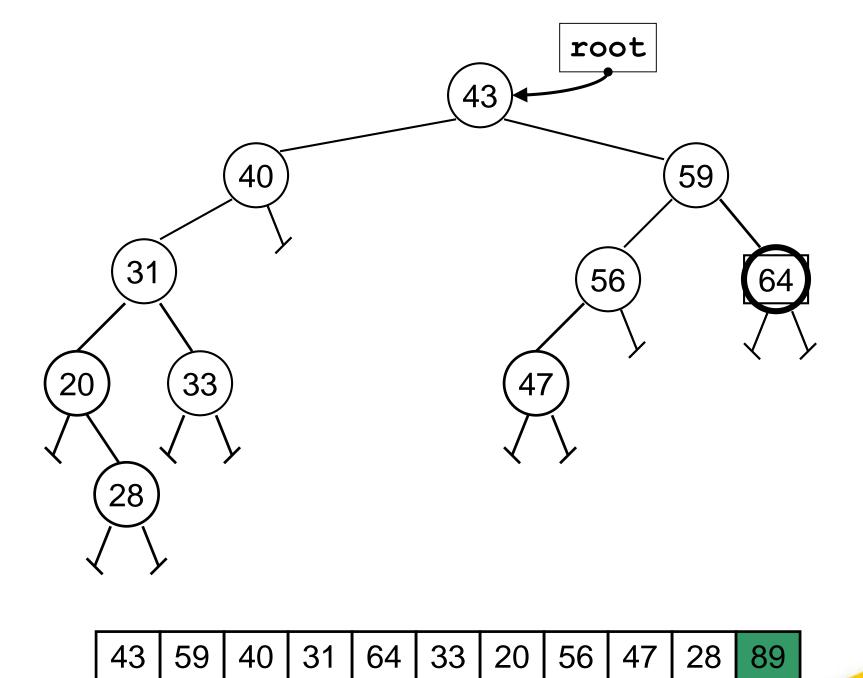




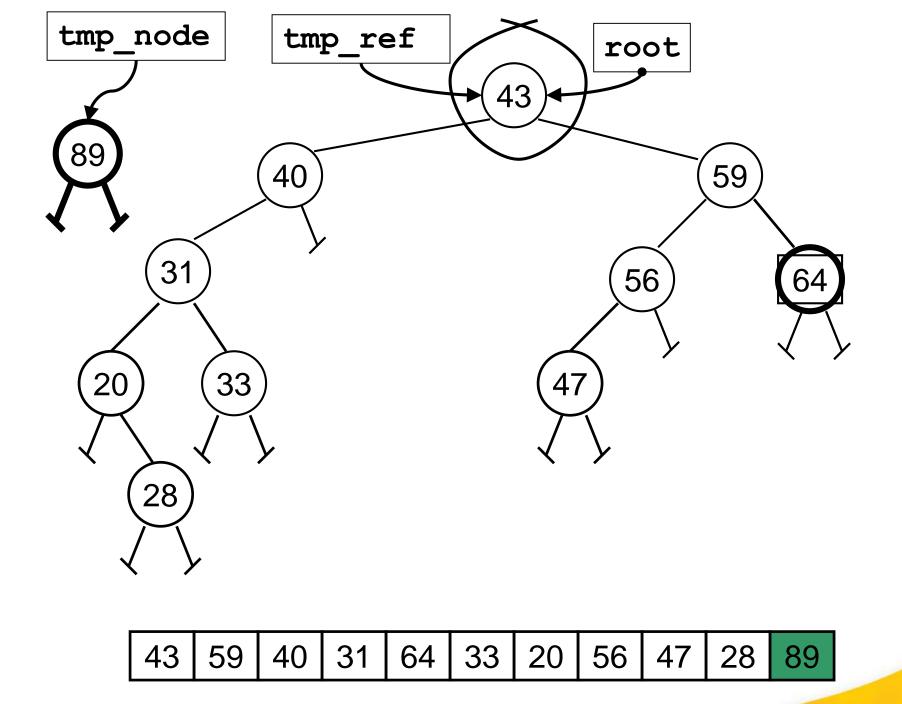




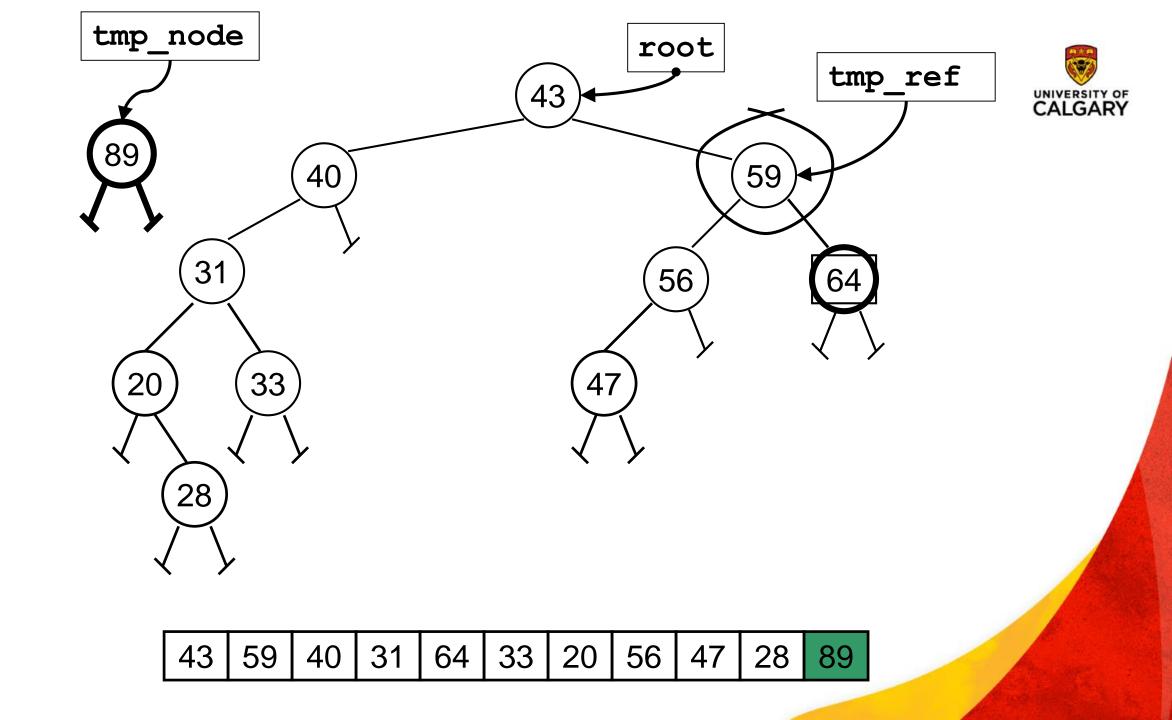


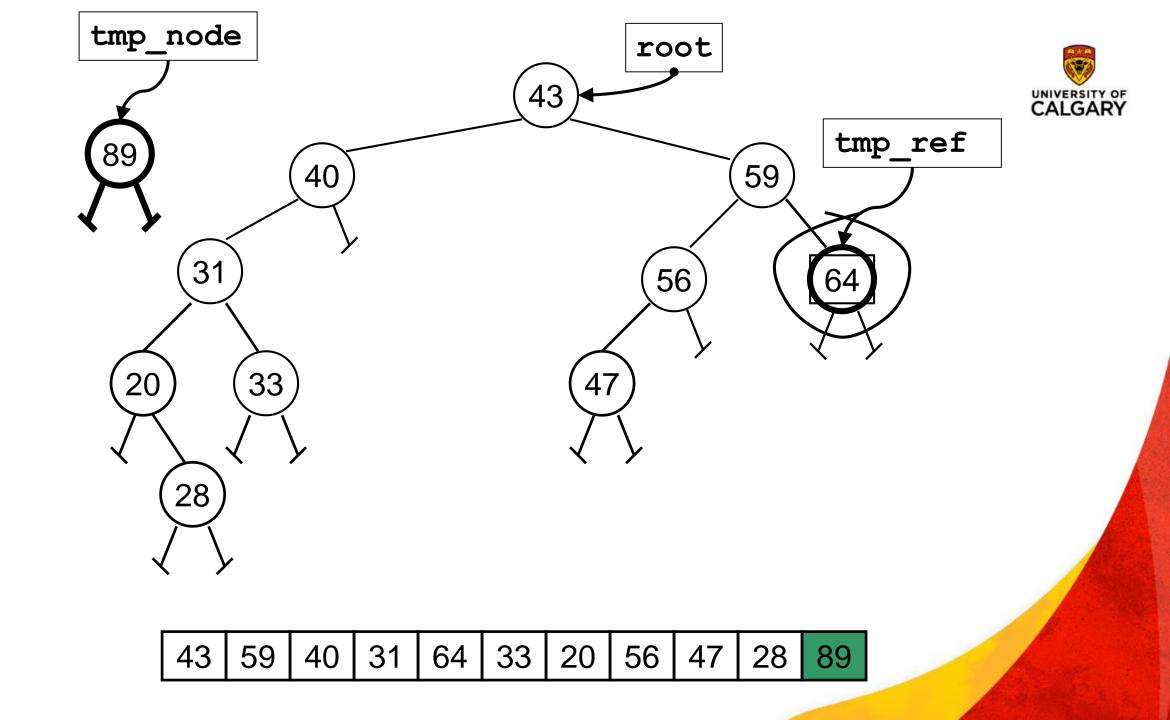


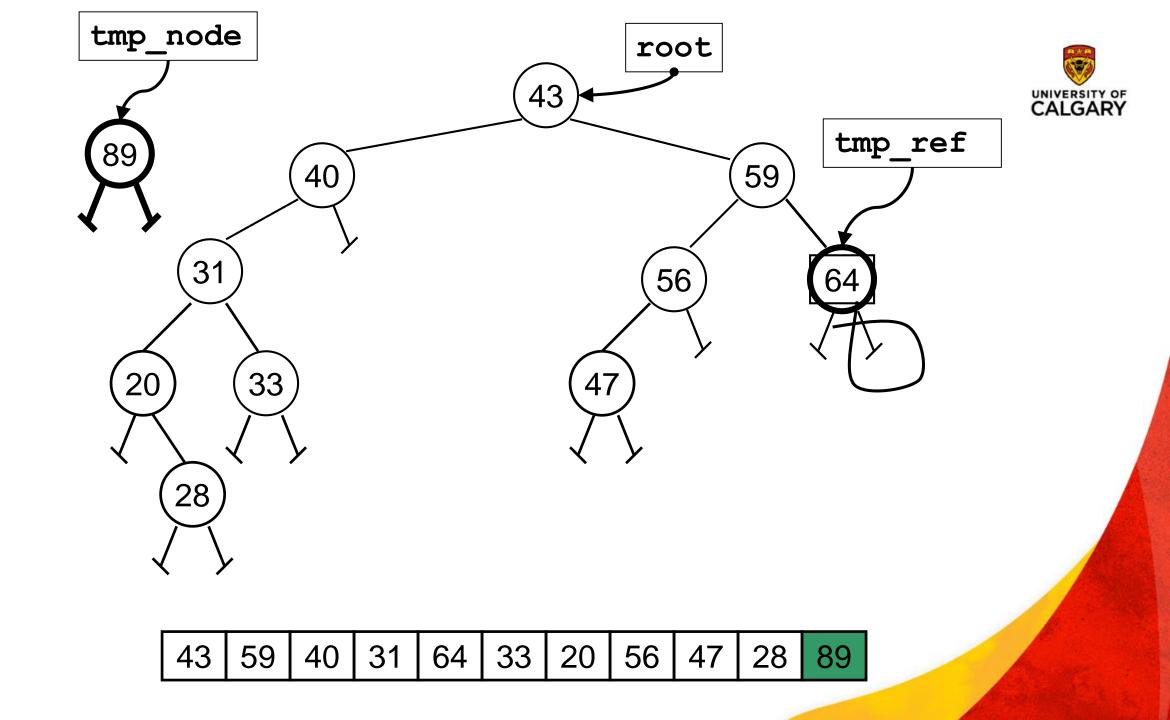


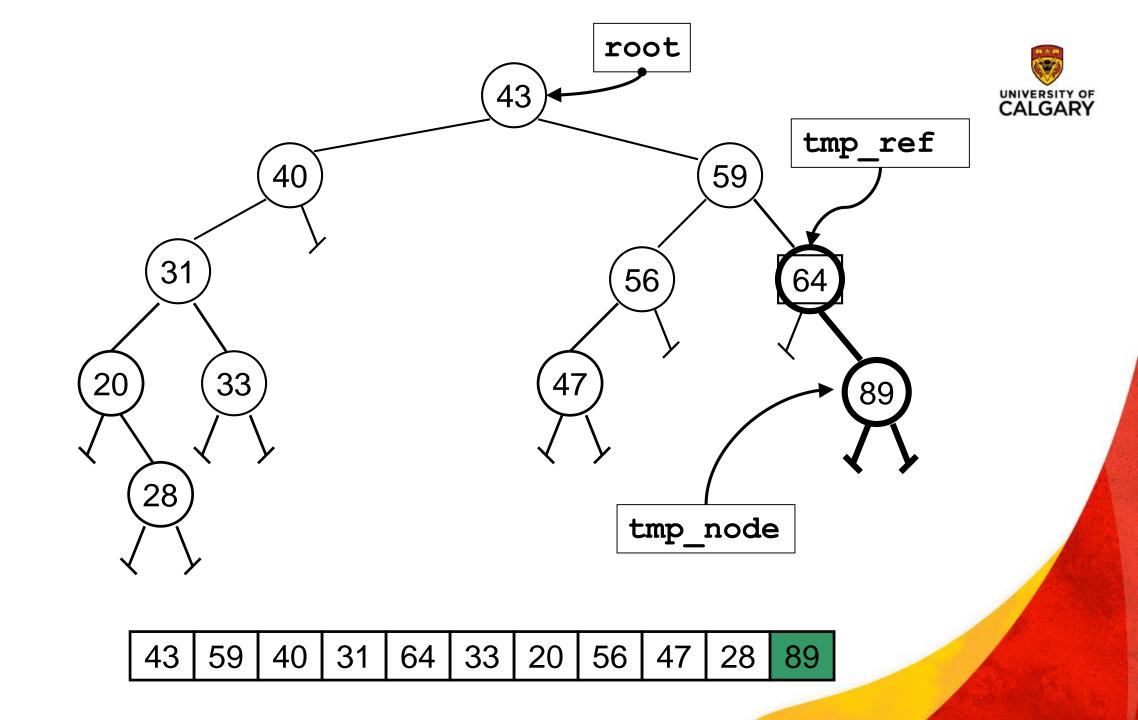




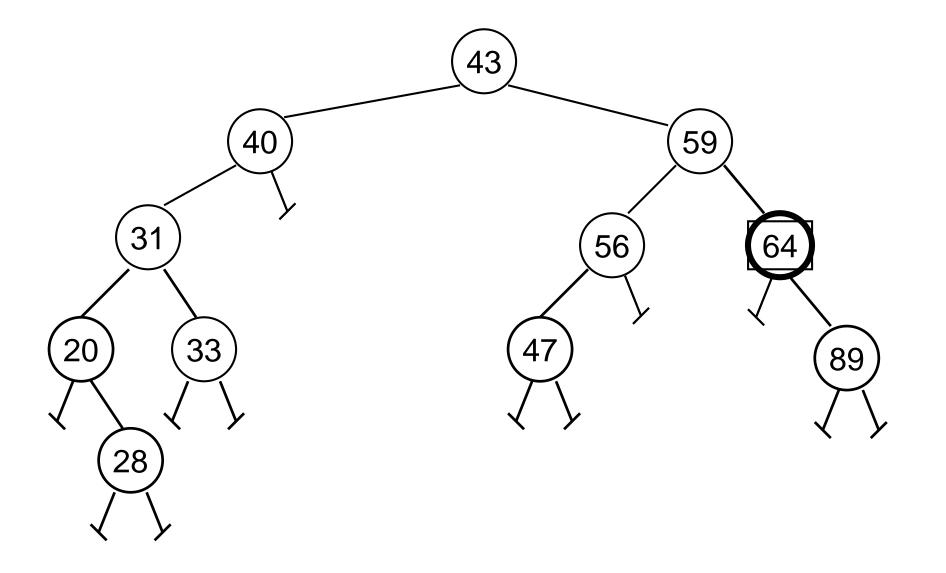












43 59 40 31 64 33 20 56 47 28 89





- Preserve "search" structure!
- Inserting involves 2 steps:
  - 1. Find the correct location
    - For a BST insert, always insert at the "bottom" of the tree (i.e., LEFT or RIGHT child of a leaf node)
  - 2. Do commands to add node
    - Create node
    - Add DATA
    - Make LEFT and RIGHT pointers point to NULL



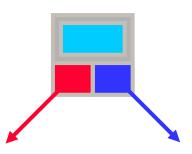


- Now we will explore 'traversal' of a binary tree
- Traversal: How do we move around the tree to interact with (or simply print) the data stored inside of it?

## **In-Order Traversal: The Scenario**



- Imagine we have a binary tree
- We want to traverse the tree
  - It's not linear
  - We need a way to visit all nodes
- Three things must happen:
  - Deal with the entire left sub-tree
  - Deal with the current node
  - Deal with the entire right sub-tree
- Result is IN-ORDER traversal

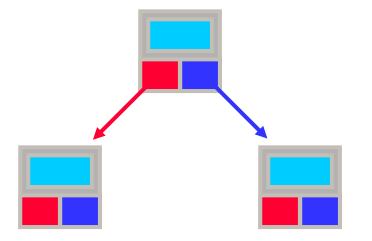






- Three principle steps:
  - Traverse Left
  - Do work (Current)
  - Traverse Right
- Work can be anything
- Separate work from traversal



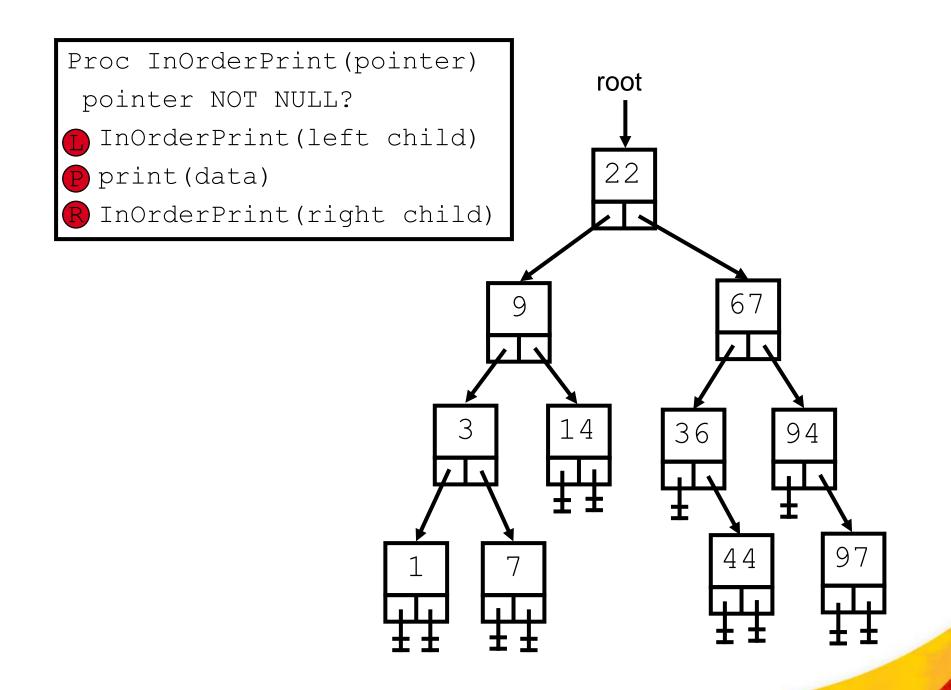


- •Traverse the tree "In order":
  - -Visit the tree's left sub-tree
  - -Visit the current and do work
  - -Visit right sub-tree

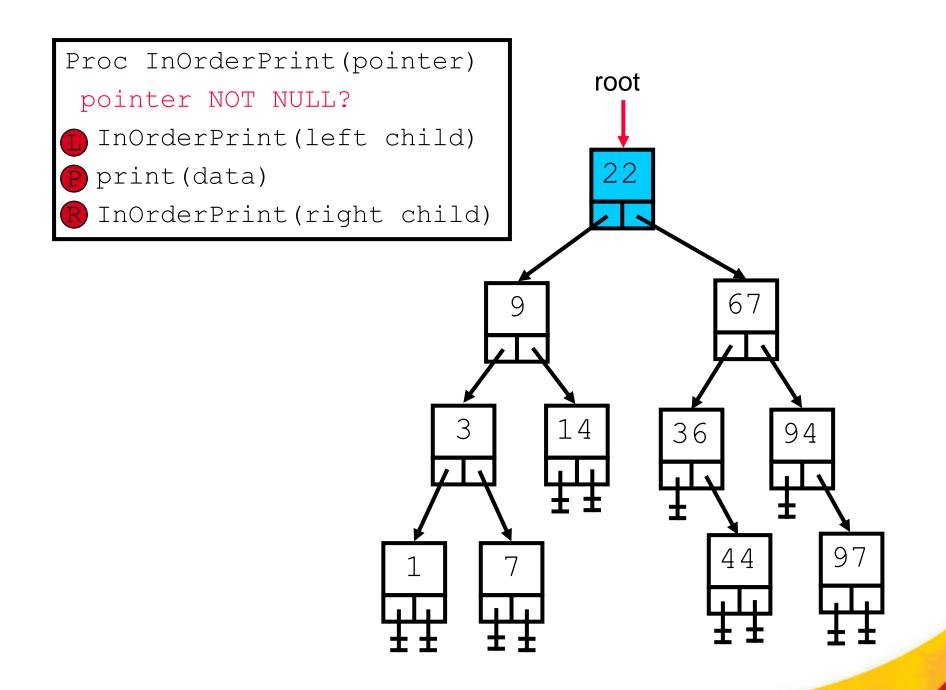




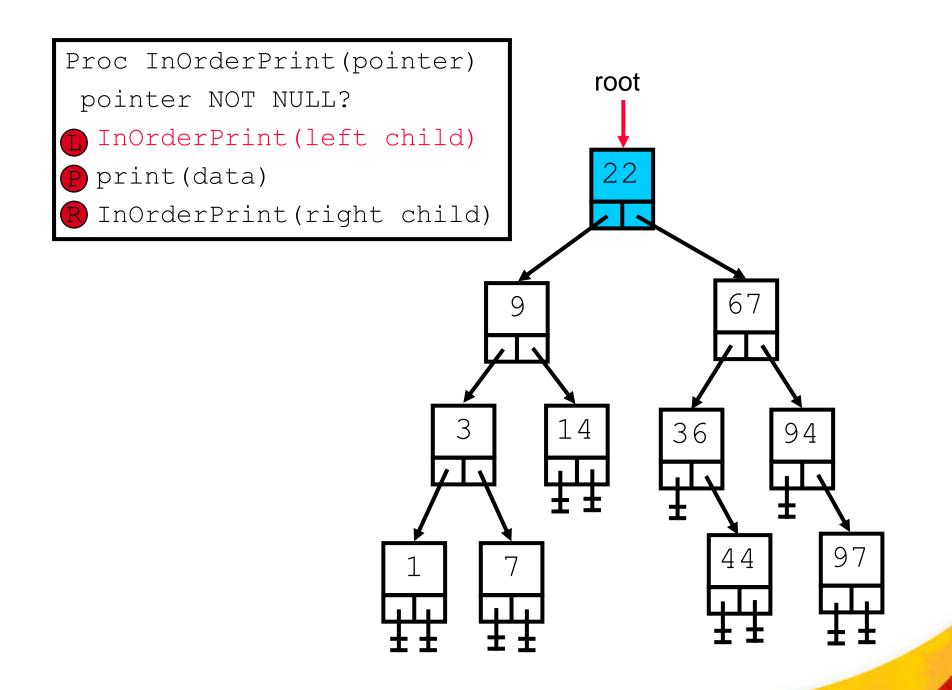
```
procedure In Order (cur iot in Ptr toa Tree Node)
// Purpose: perform in-order traversal, call
            Do Something for each node
// Preconditions: cur points to a binary tree
// Postcondition: Do Something on each tree
//
                  node in "in-order" order
 if( cur <> NULL ) then
   In Order( cur^.left child )
   Do Something (cur^.data)
   In Order( cur^.right child )
 endif
endprocedure //
                  In Order
```



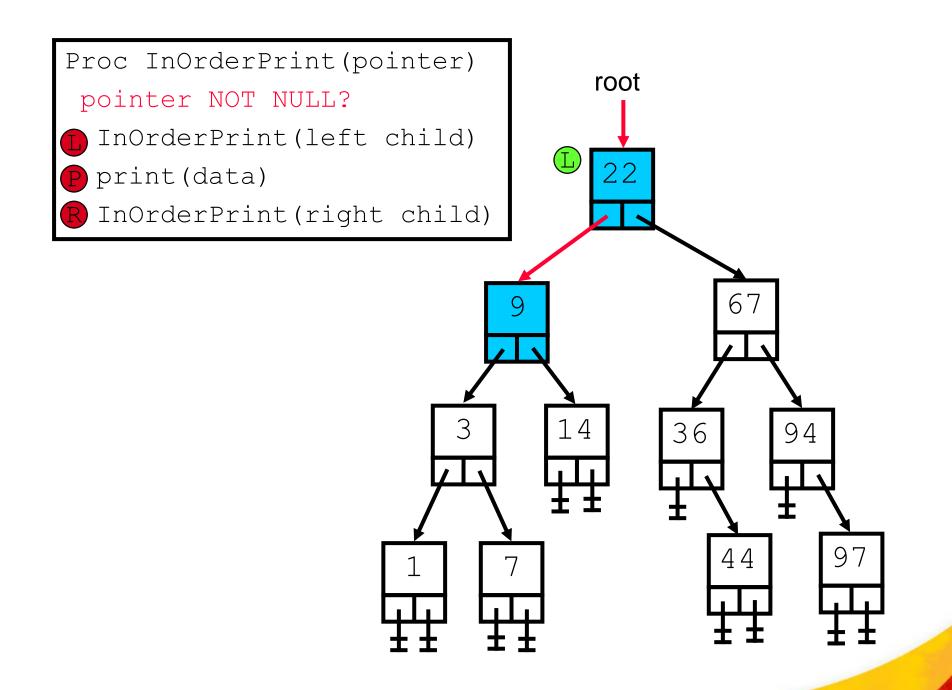




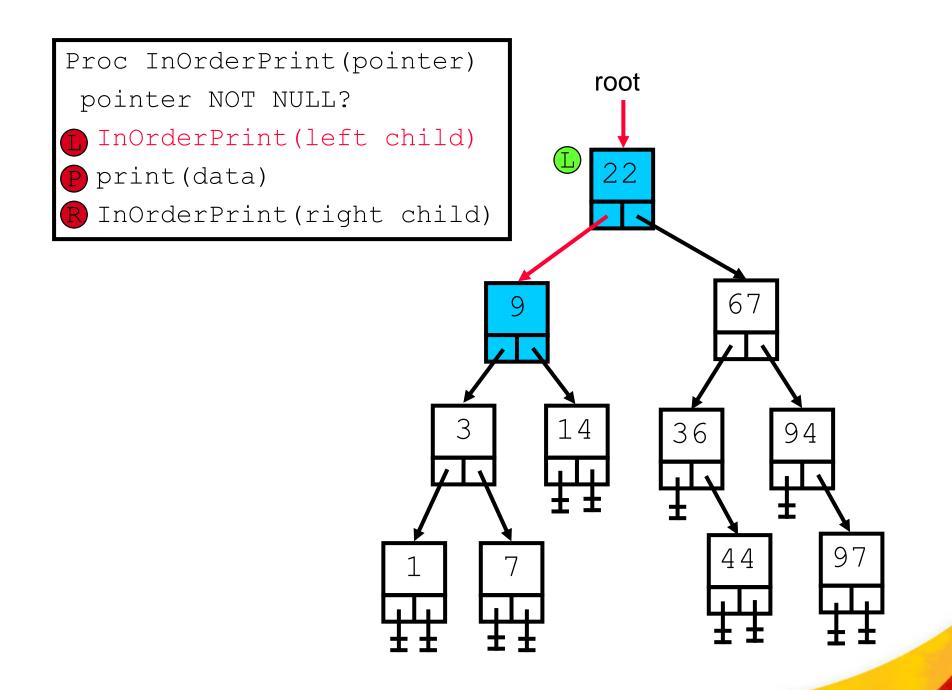




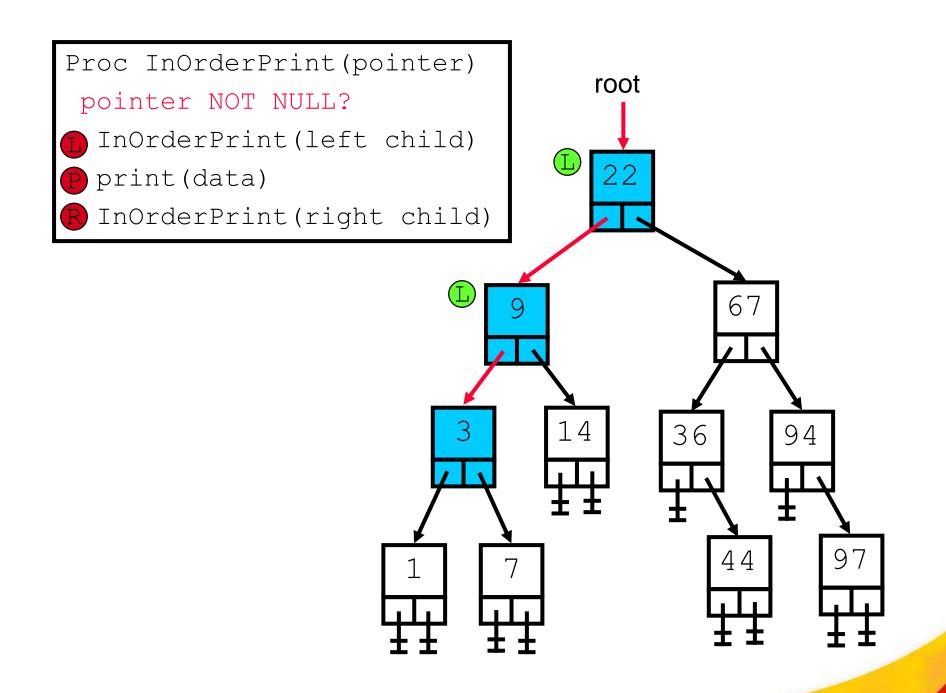




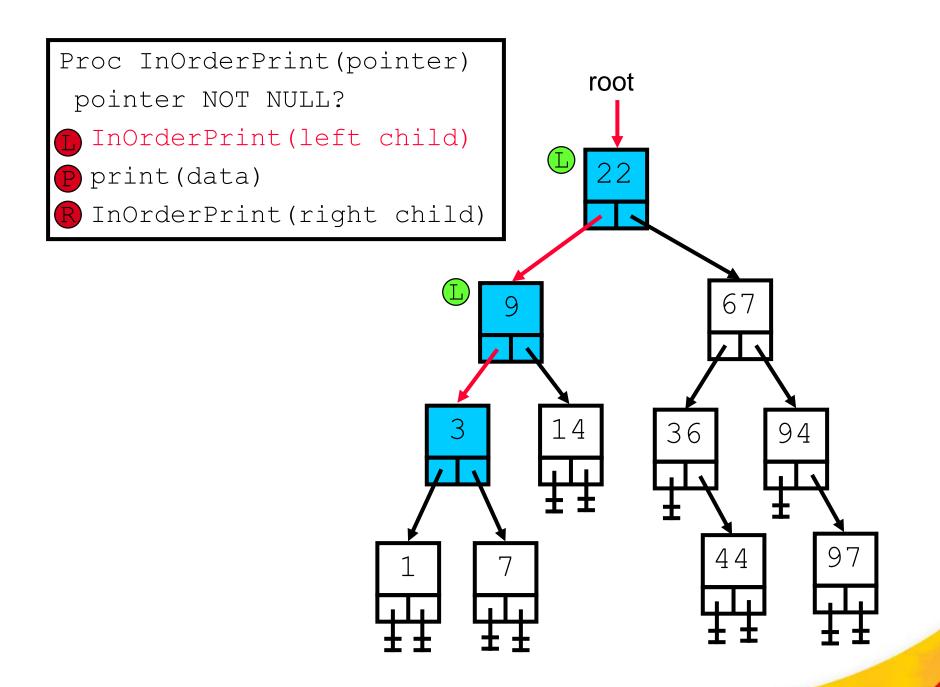




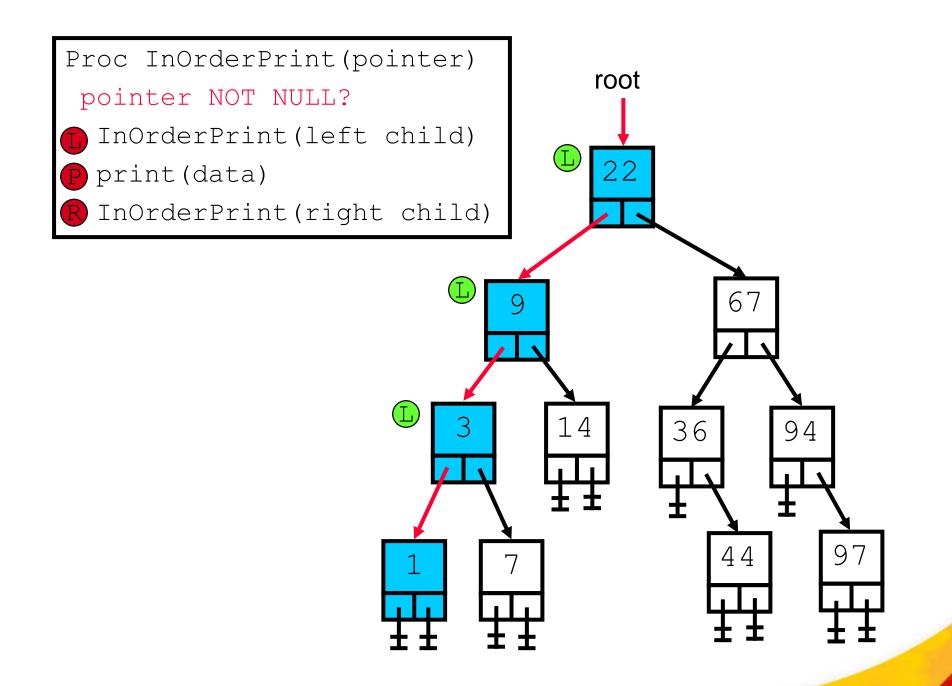




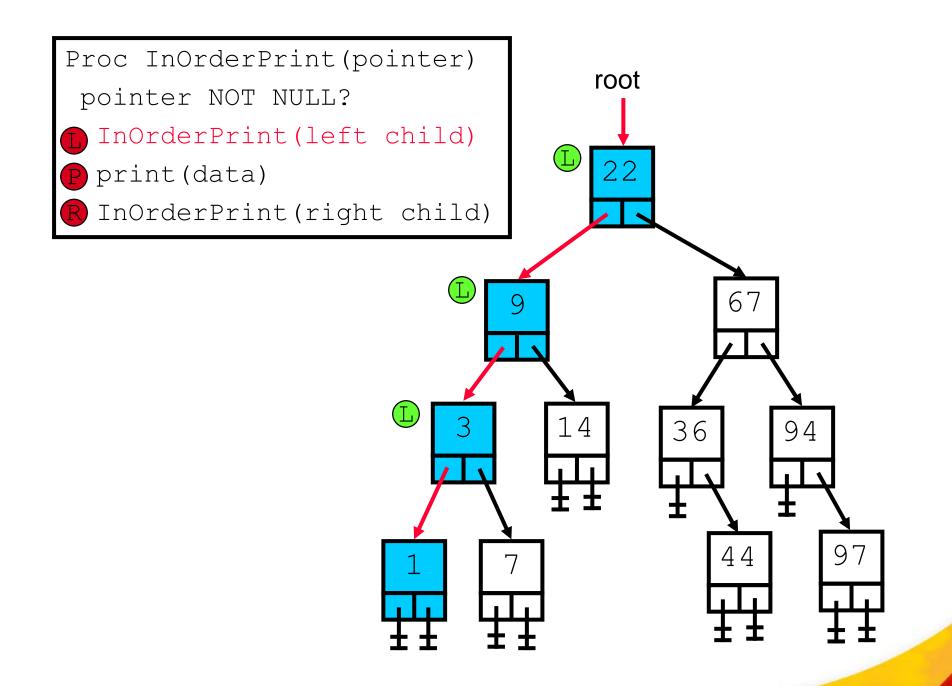




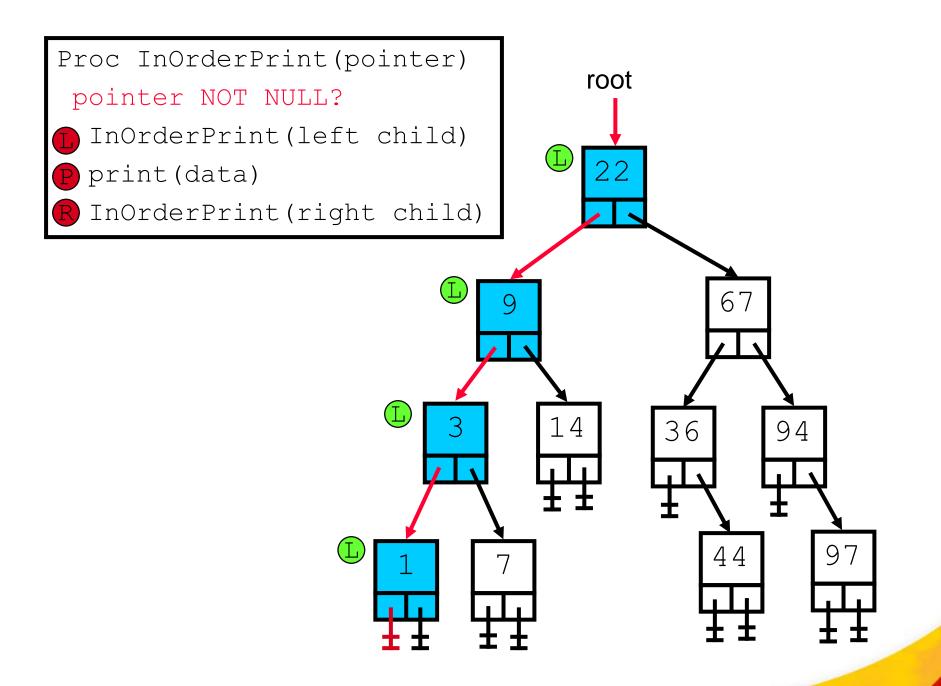




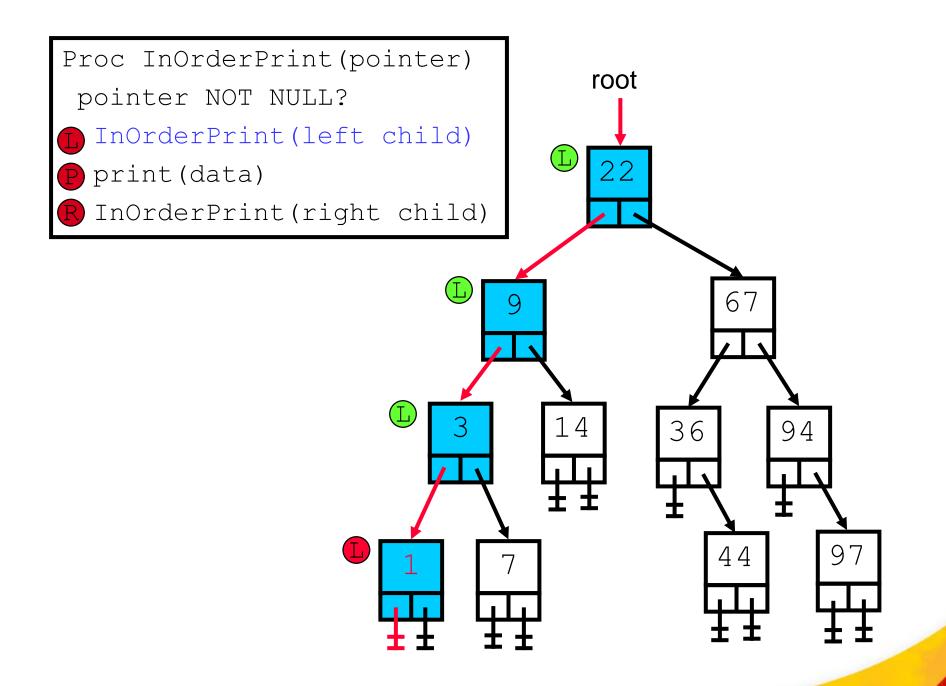




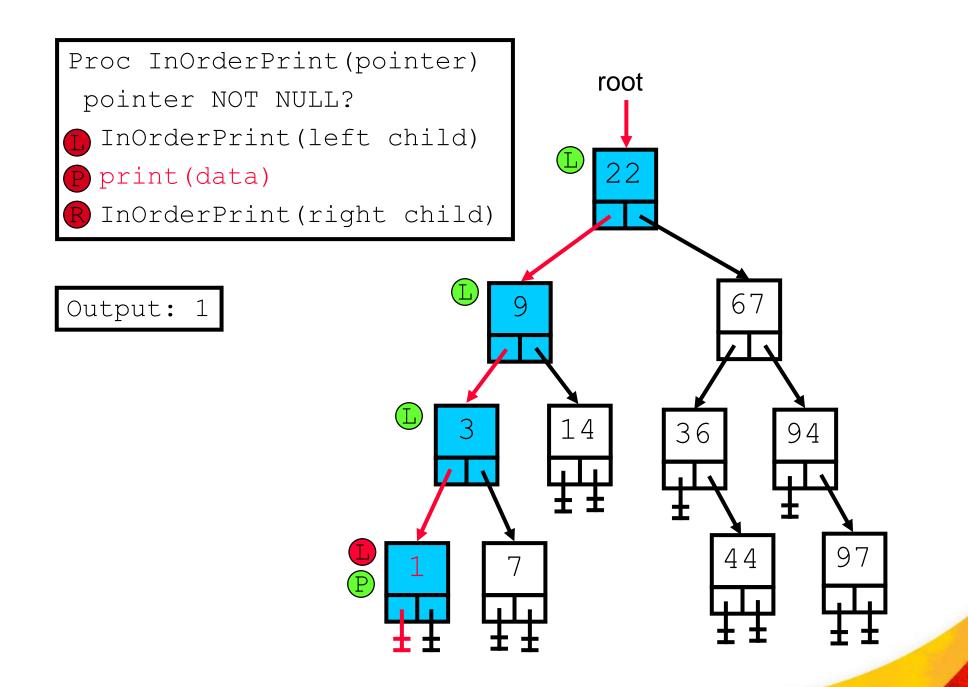




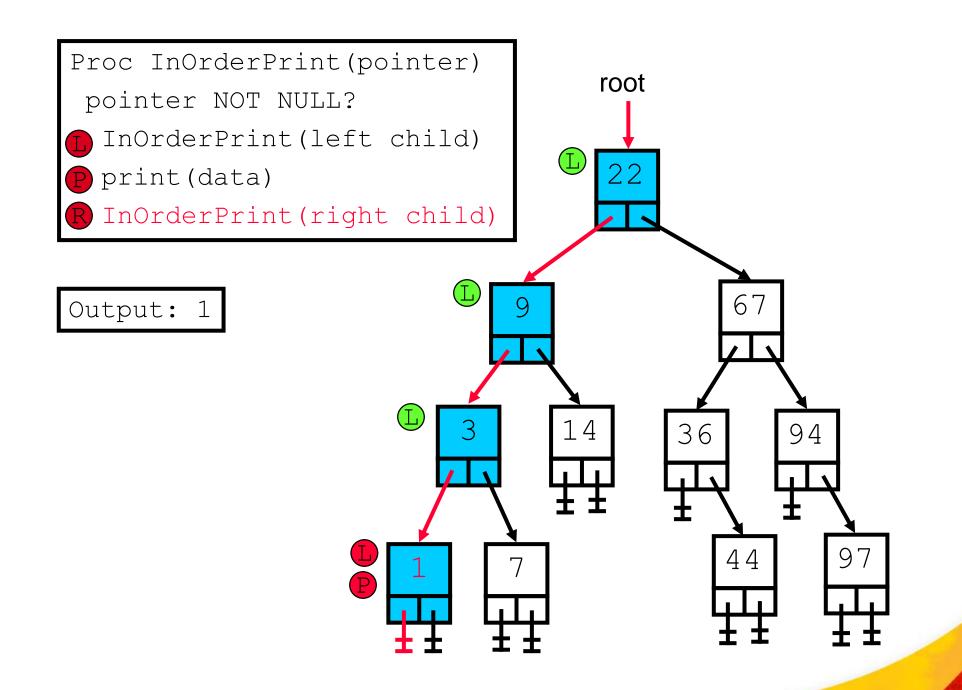




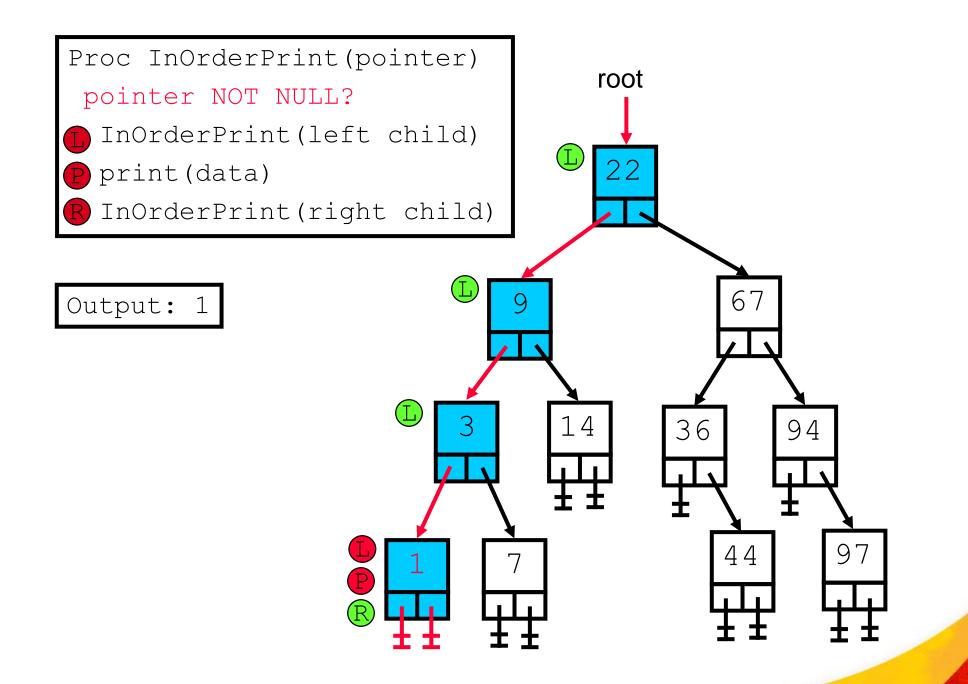




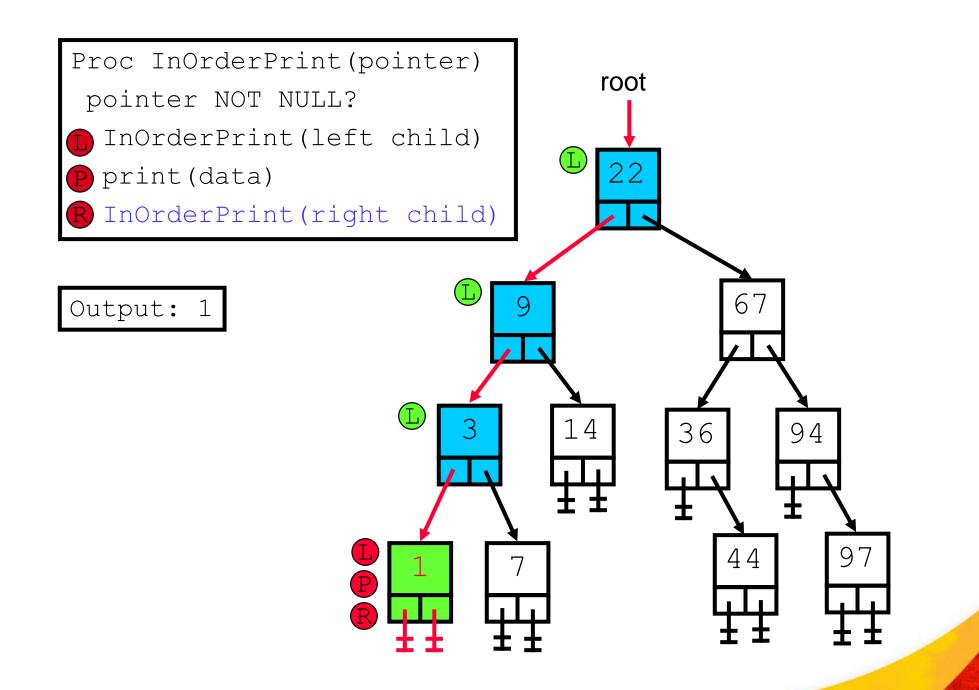




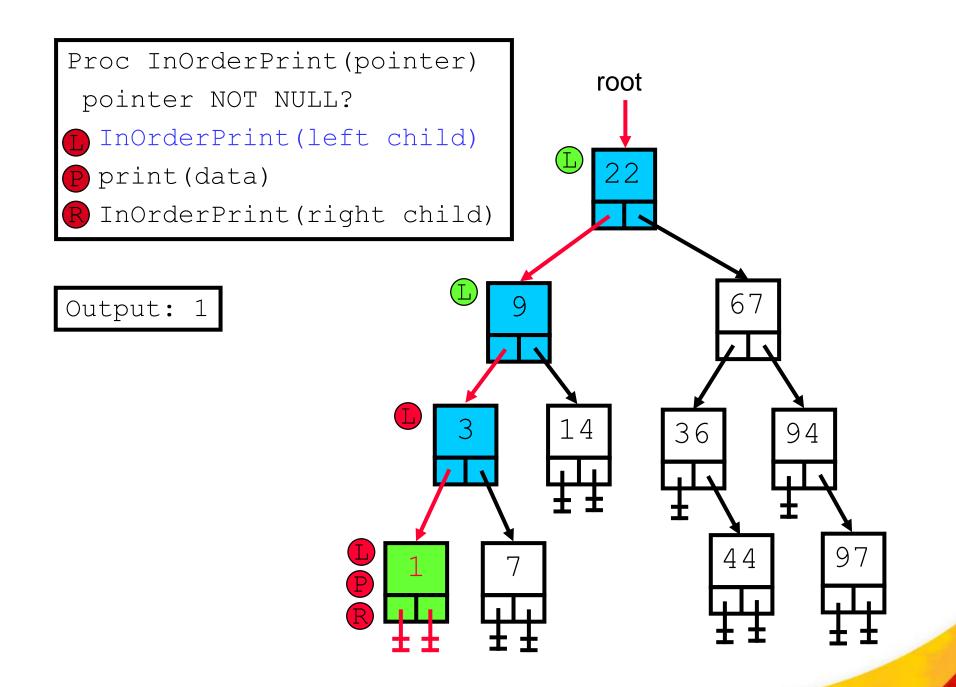




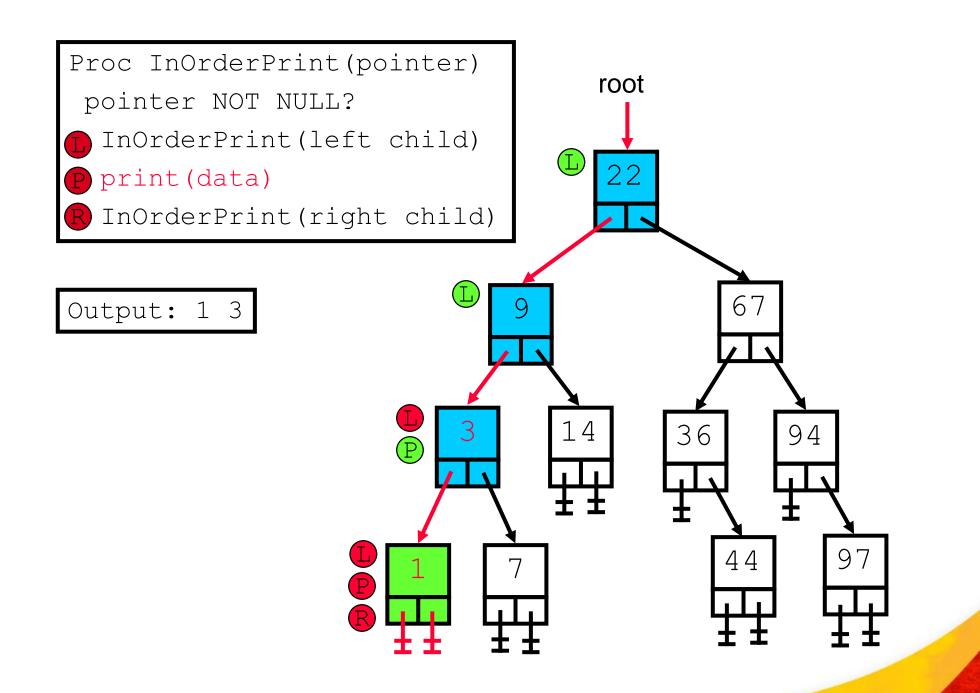




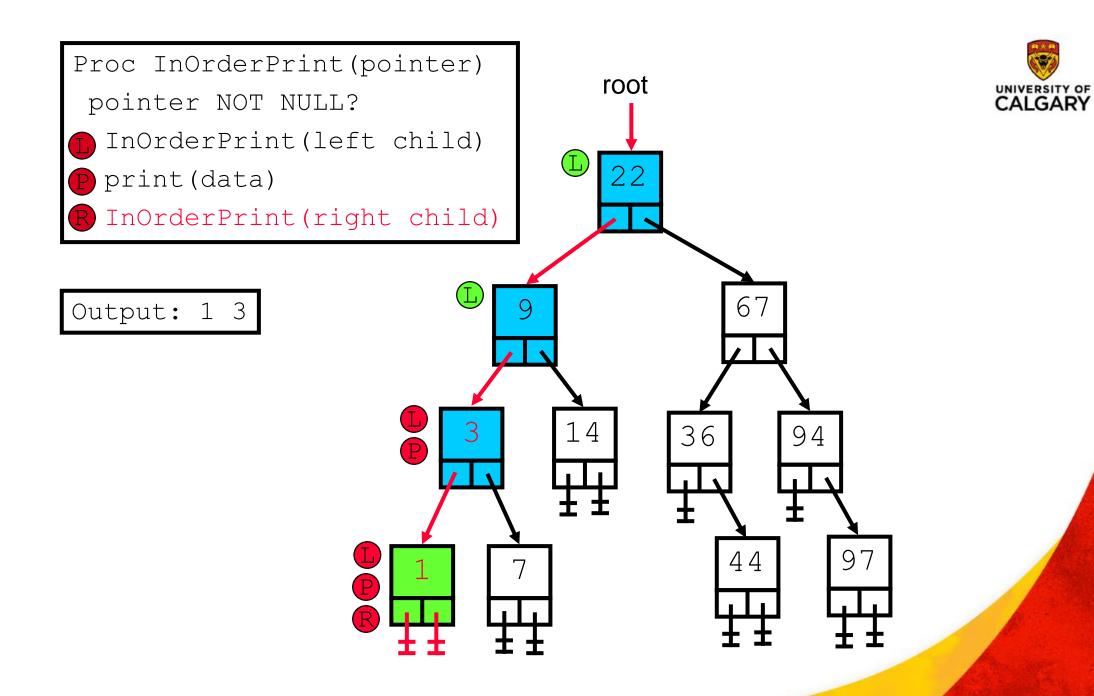


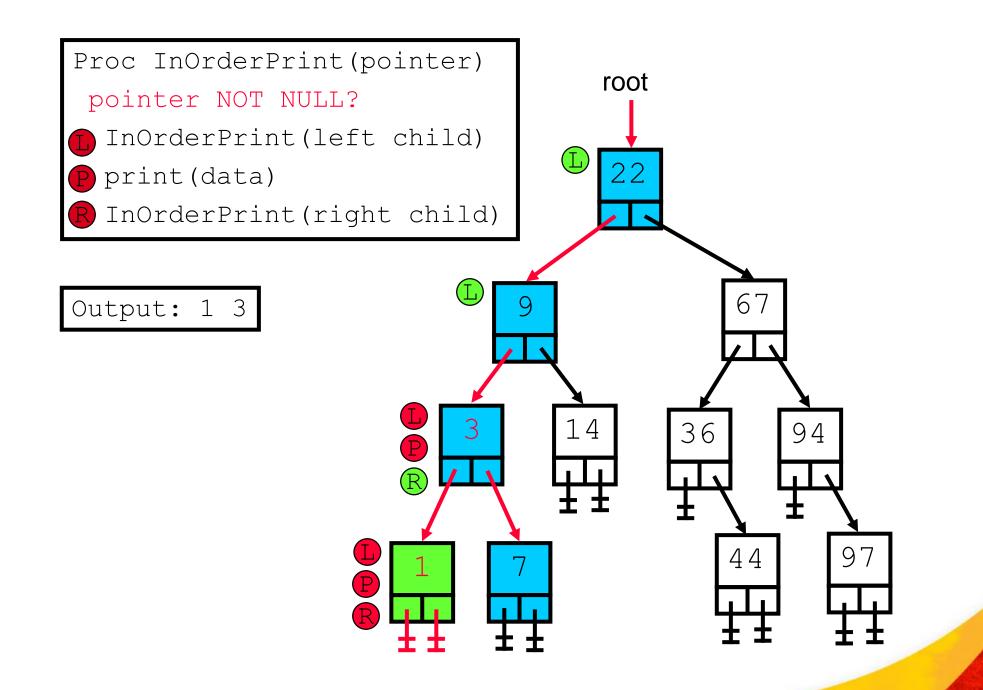




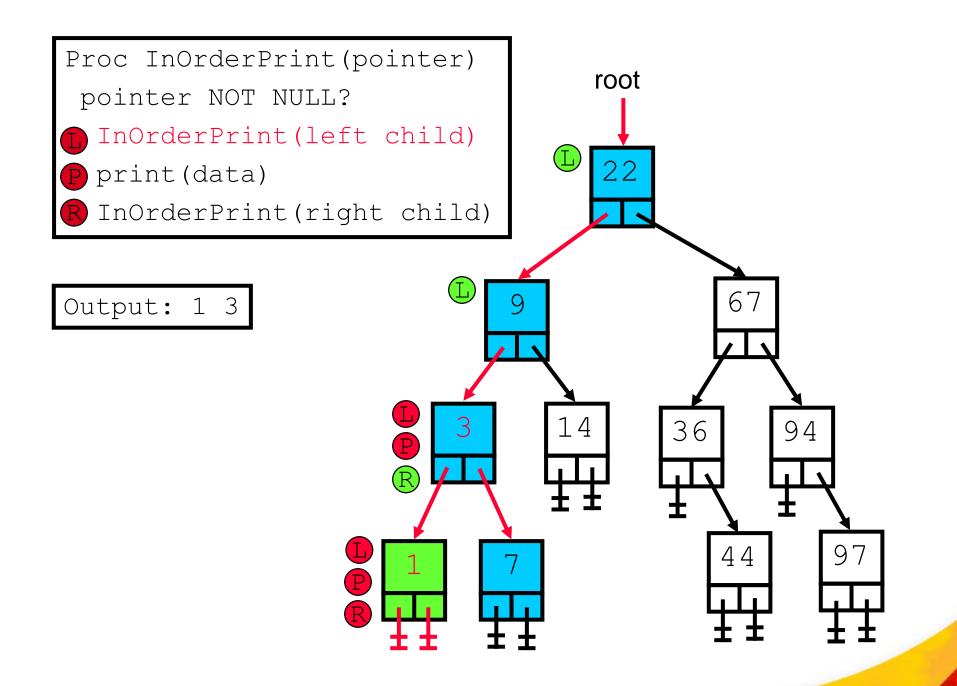




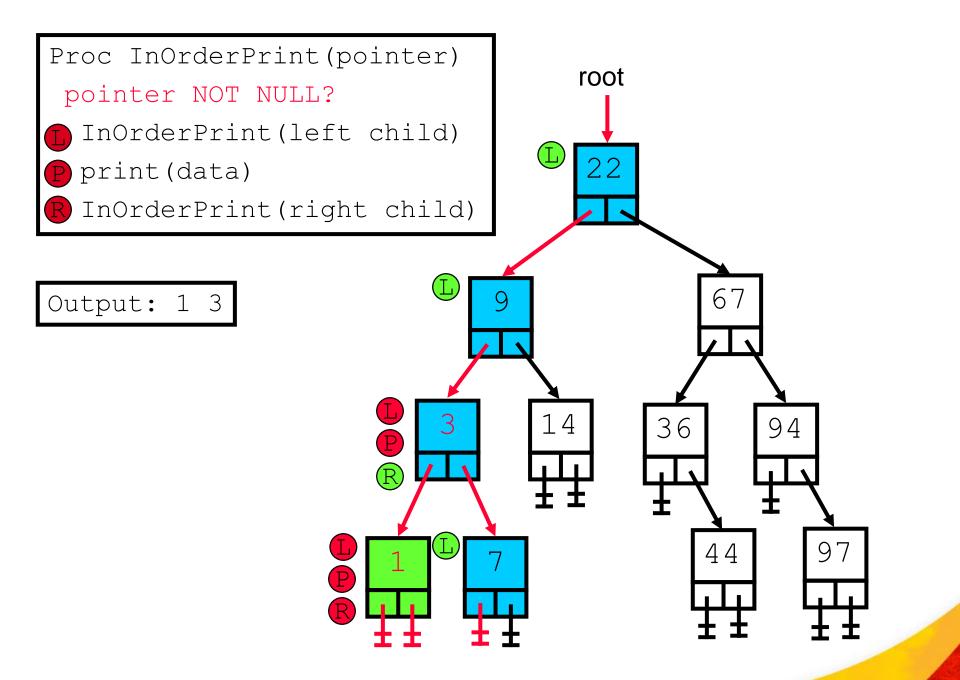




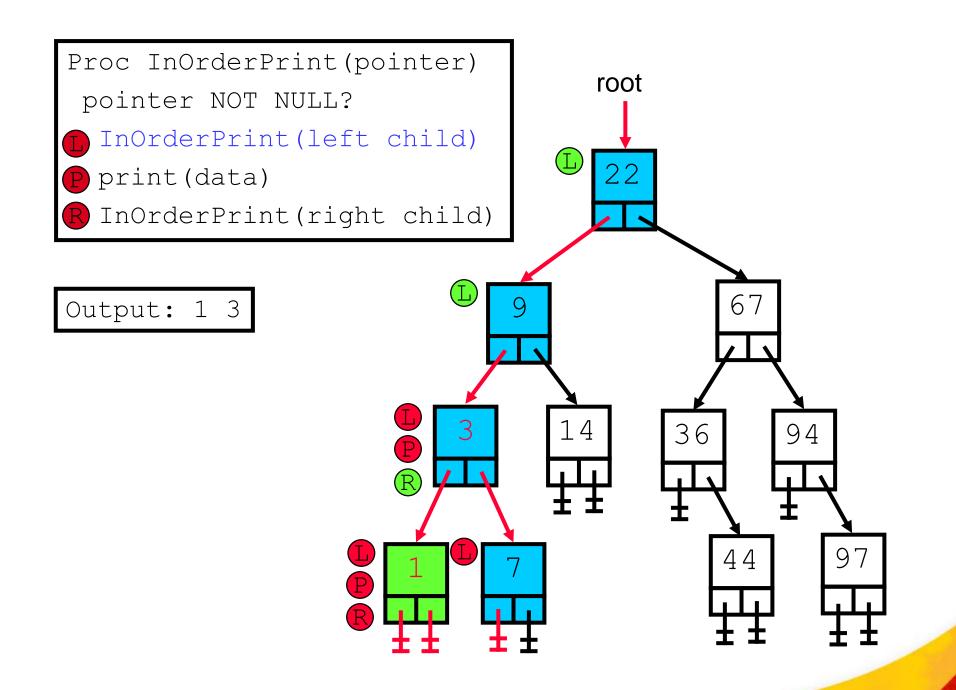




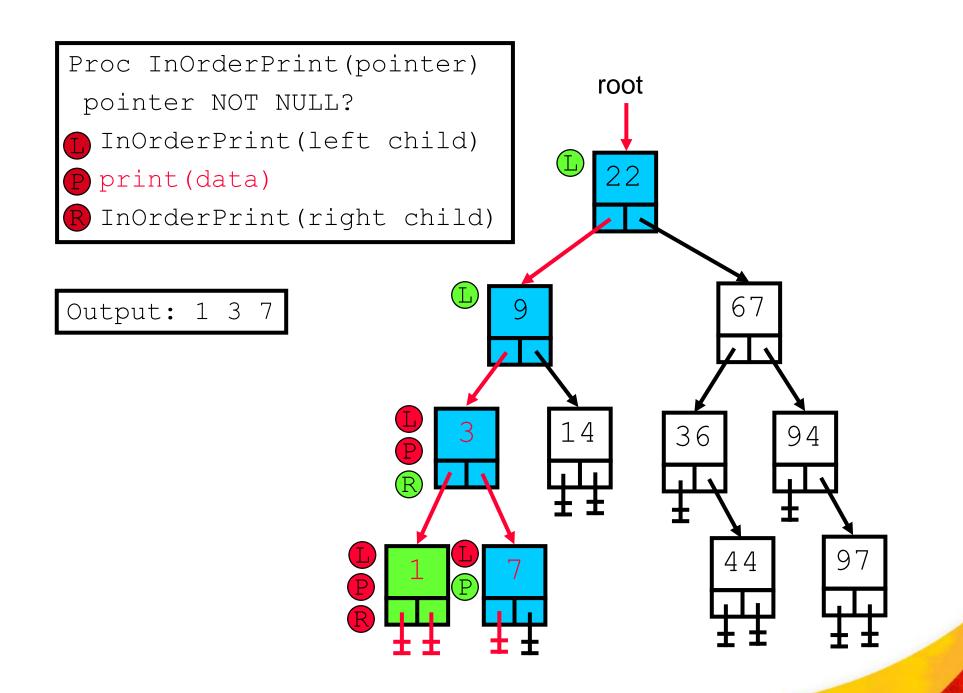




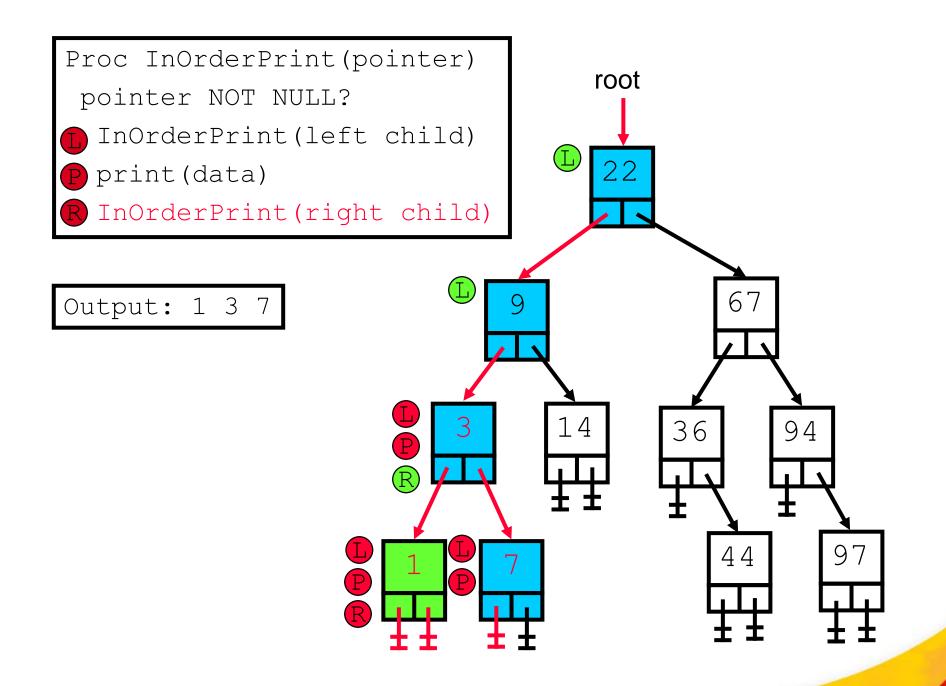




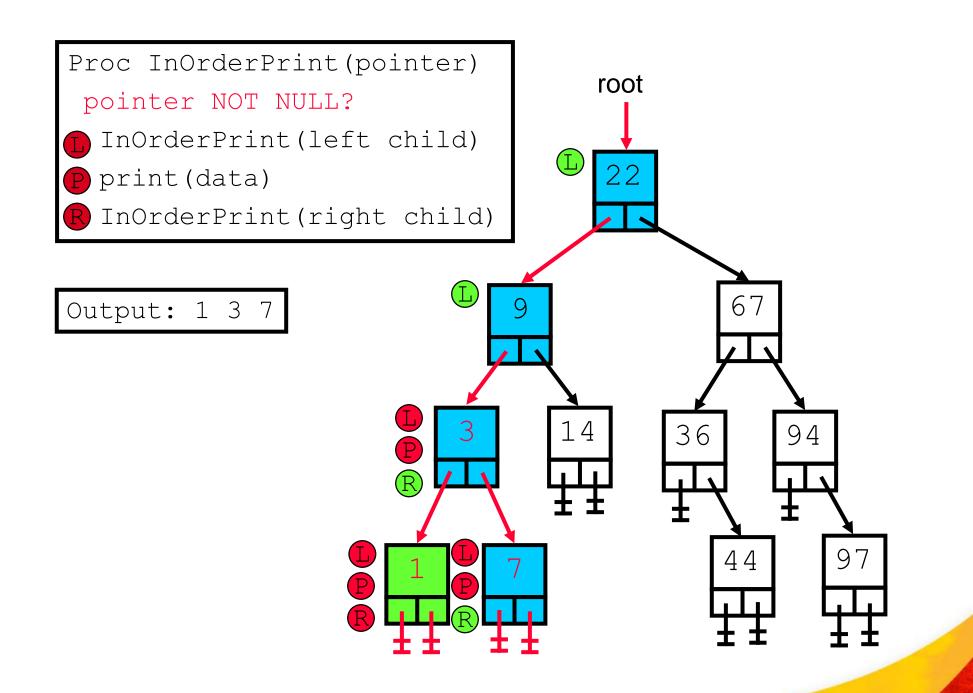




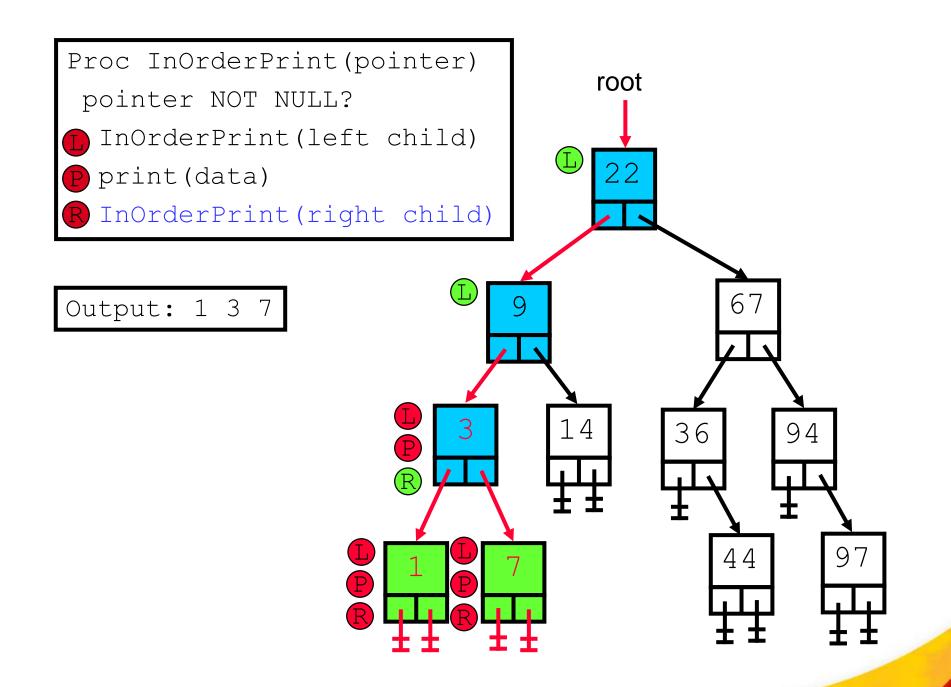
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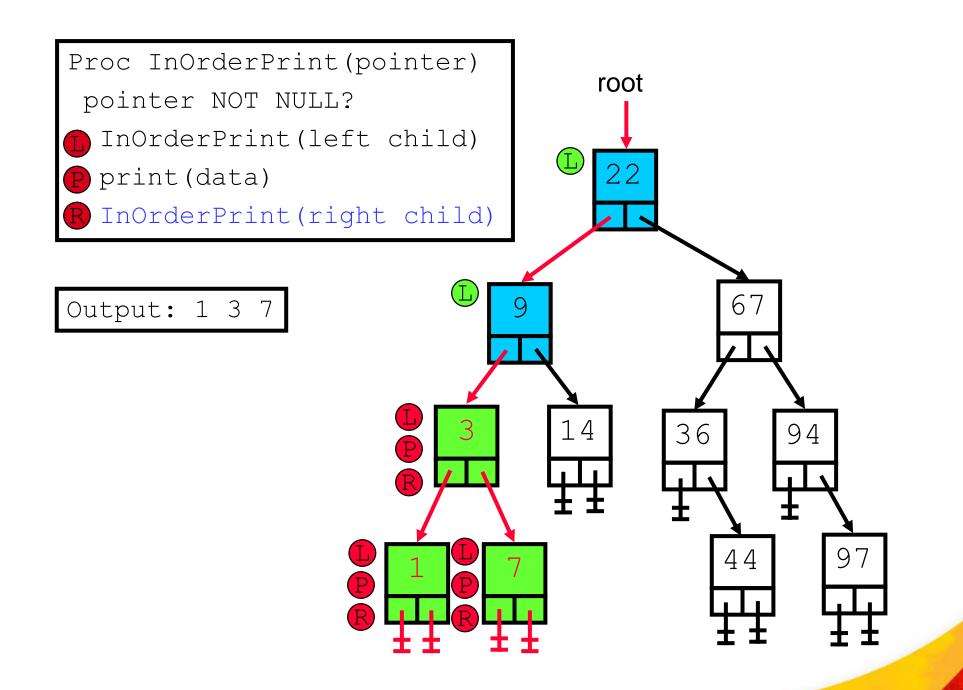




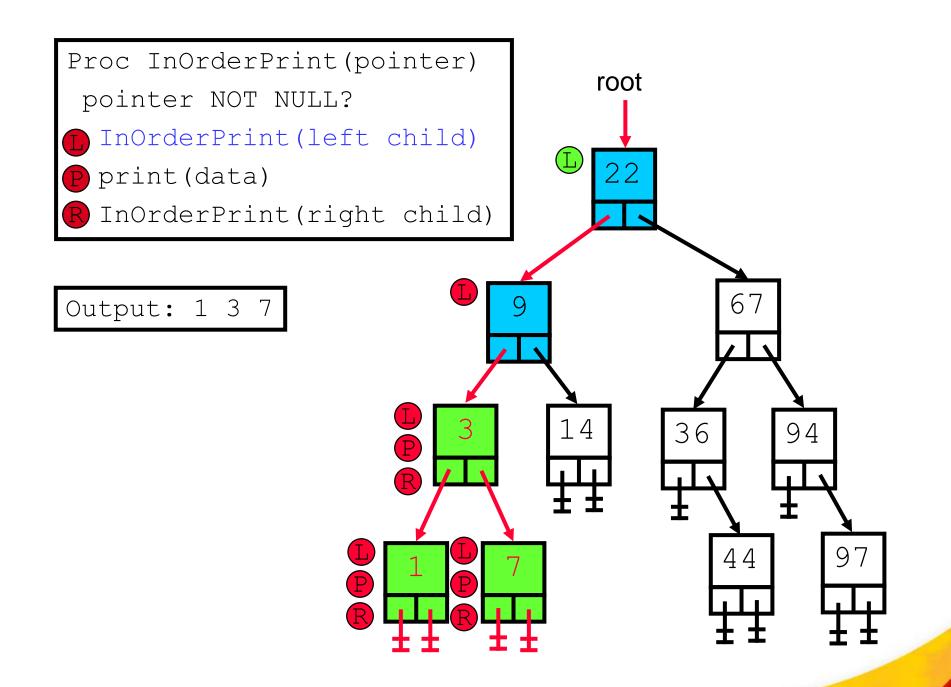




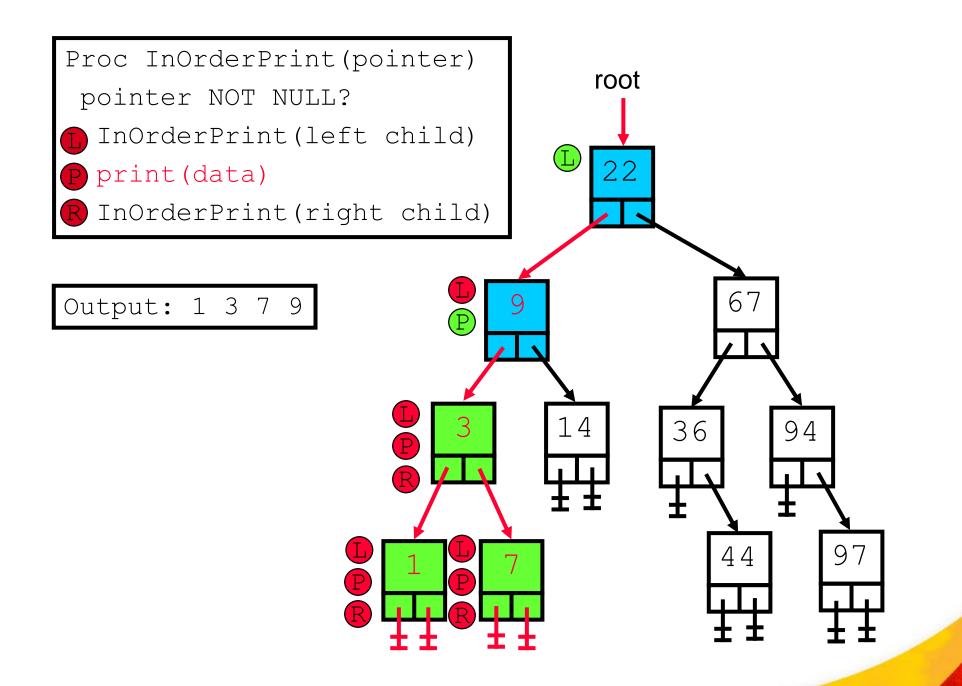




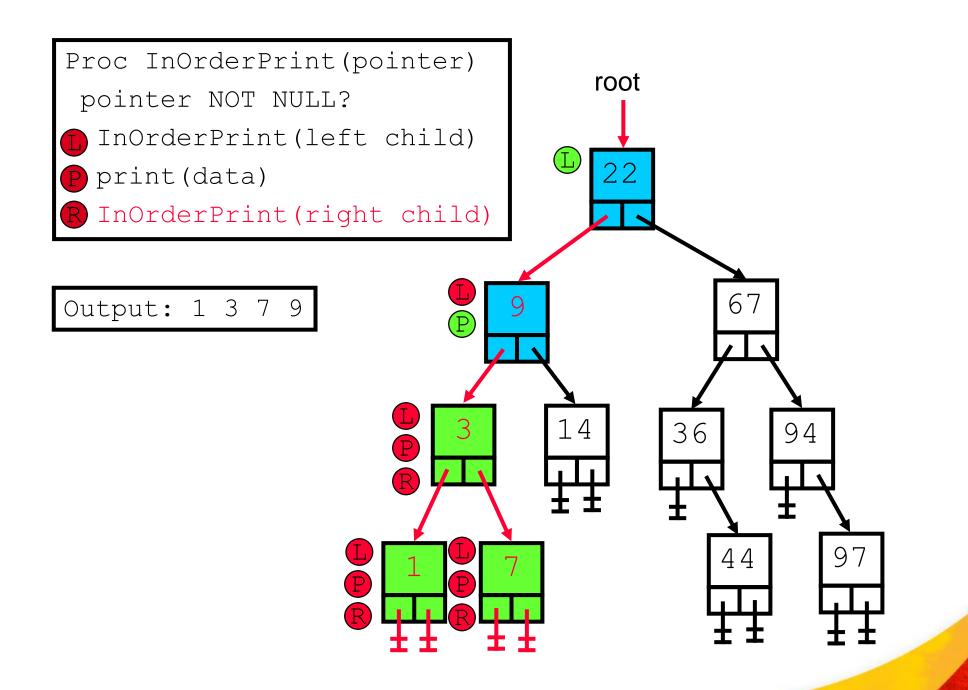




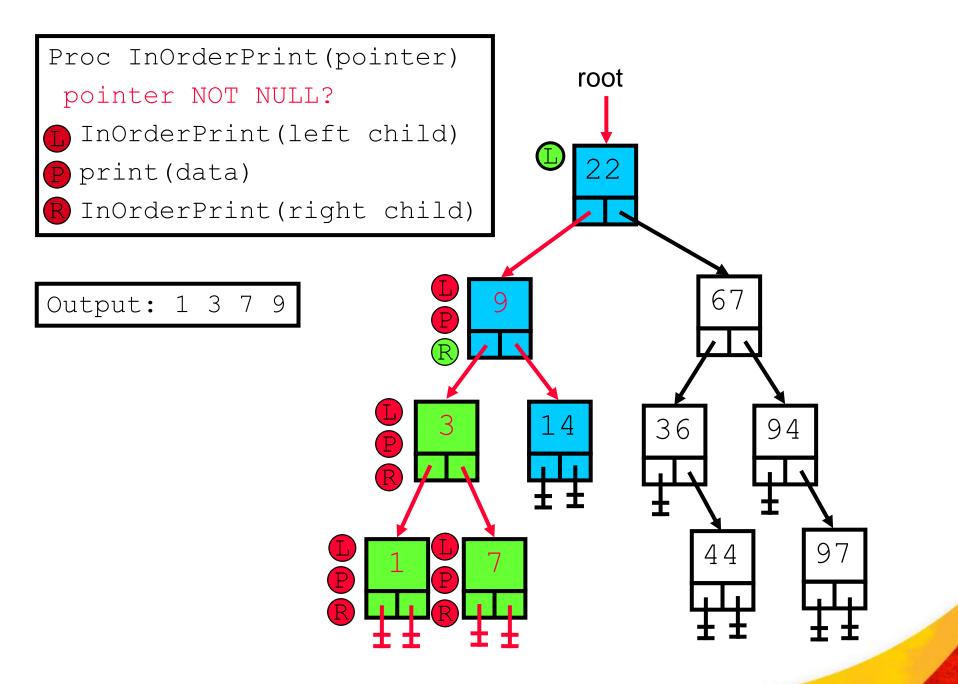




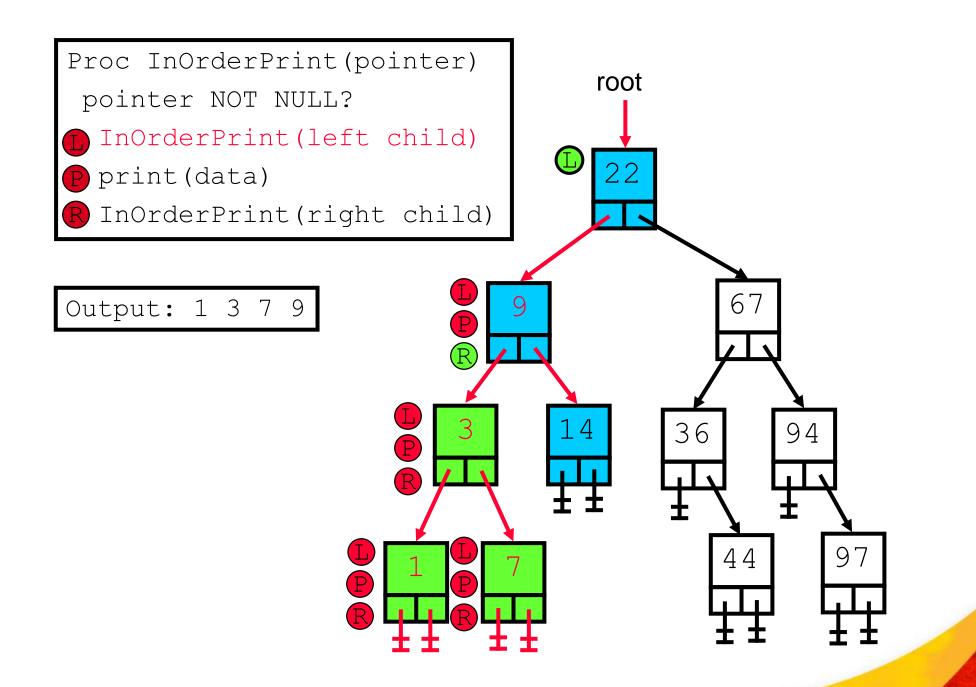




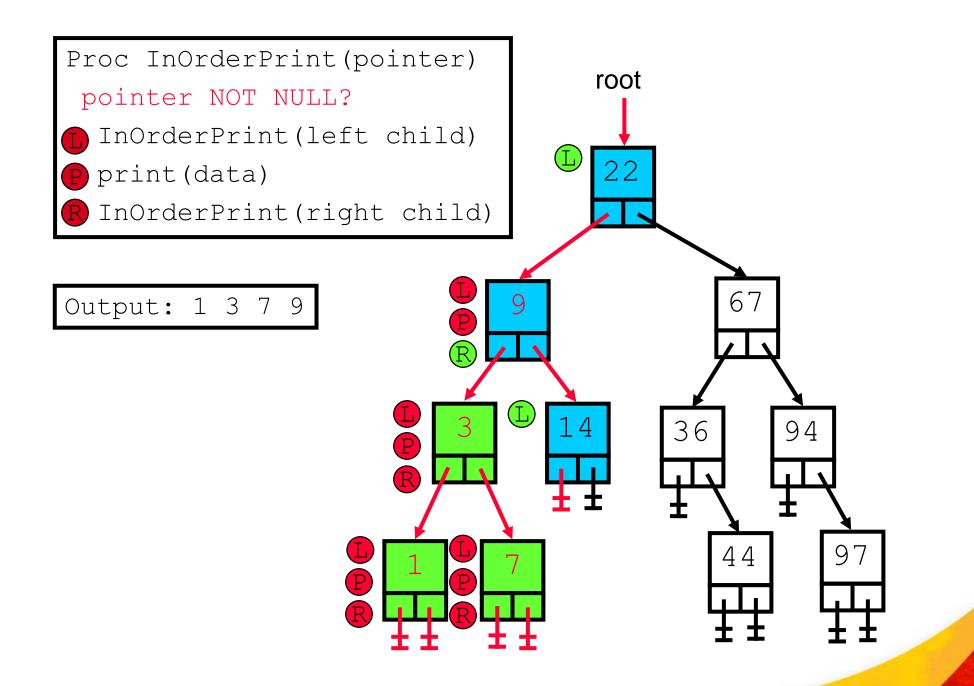




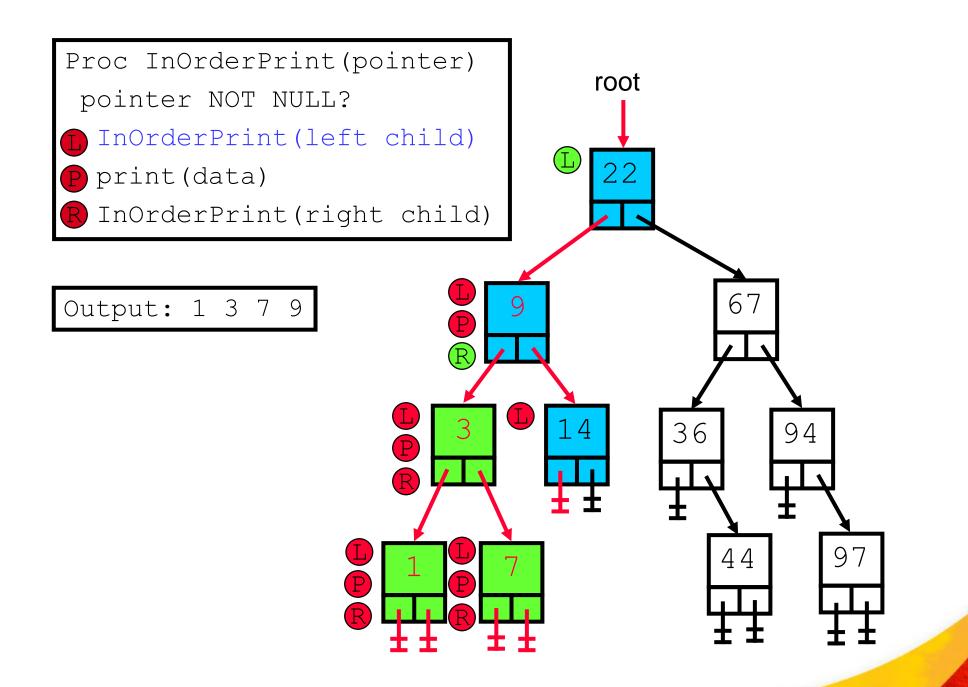




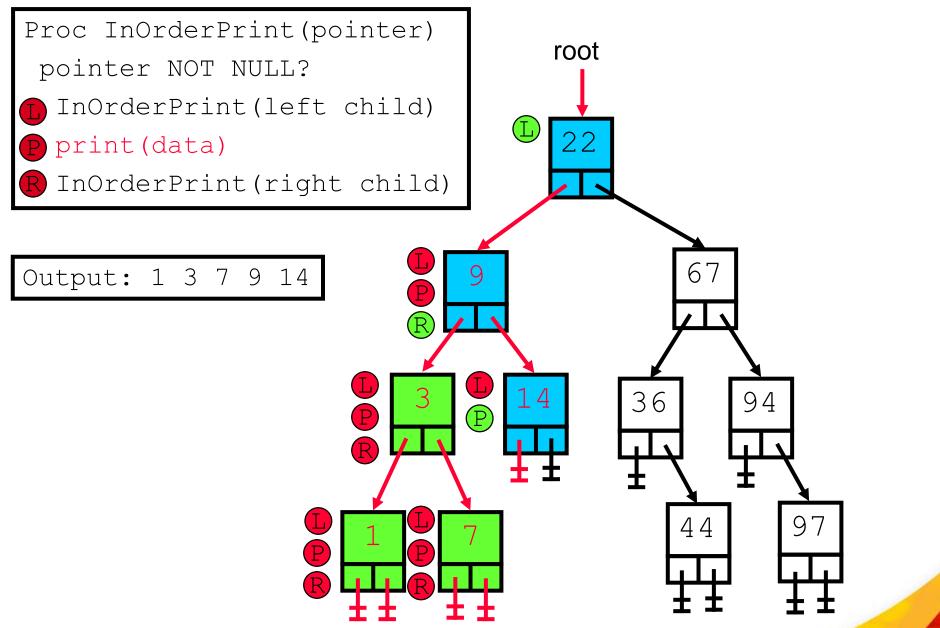




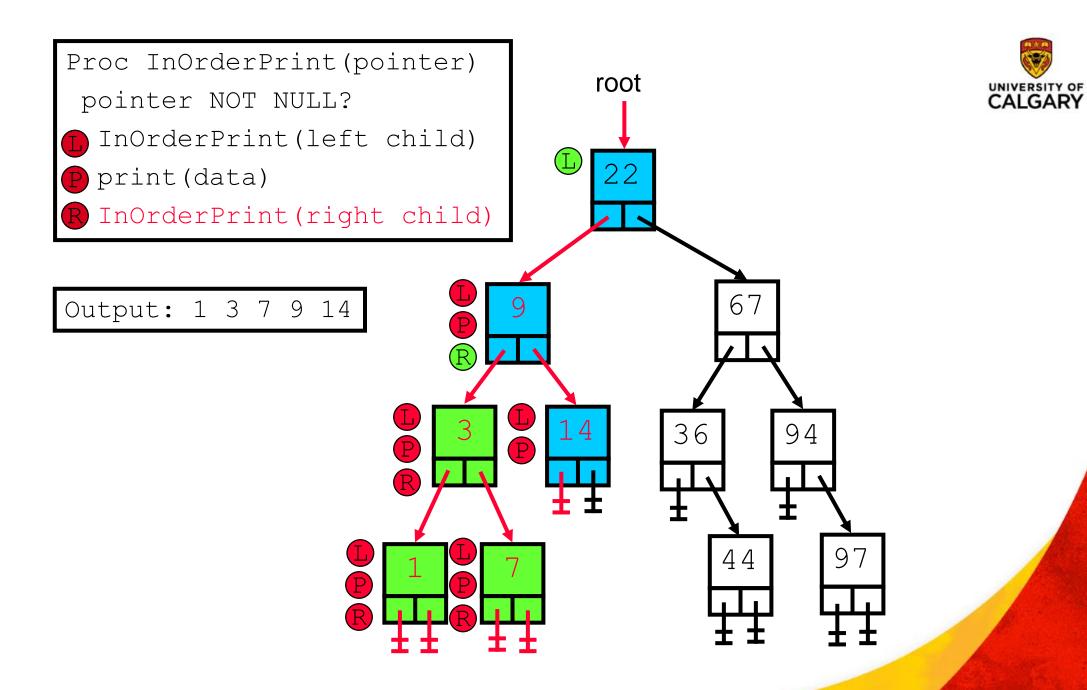


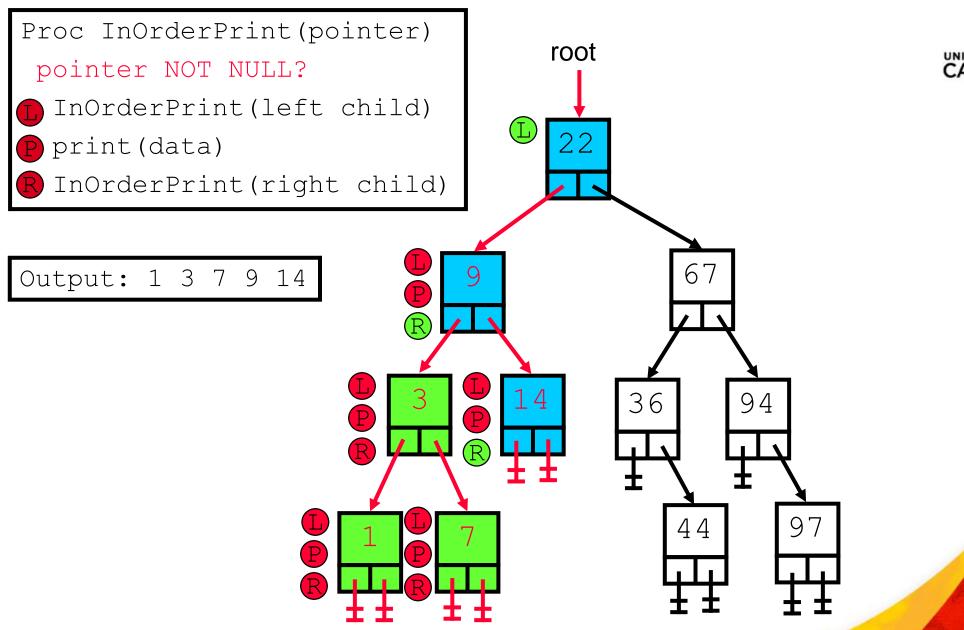




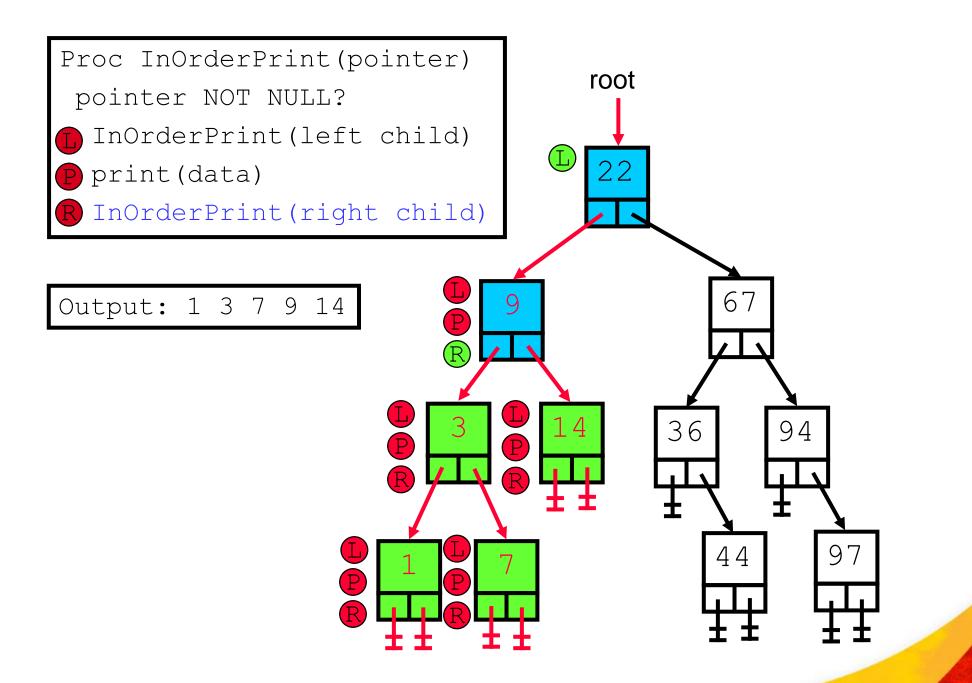




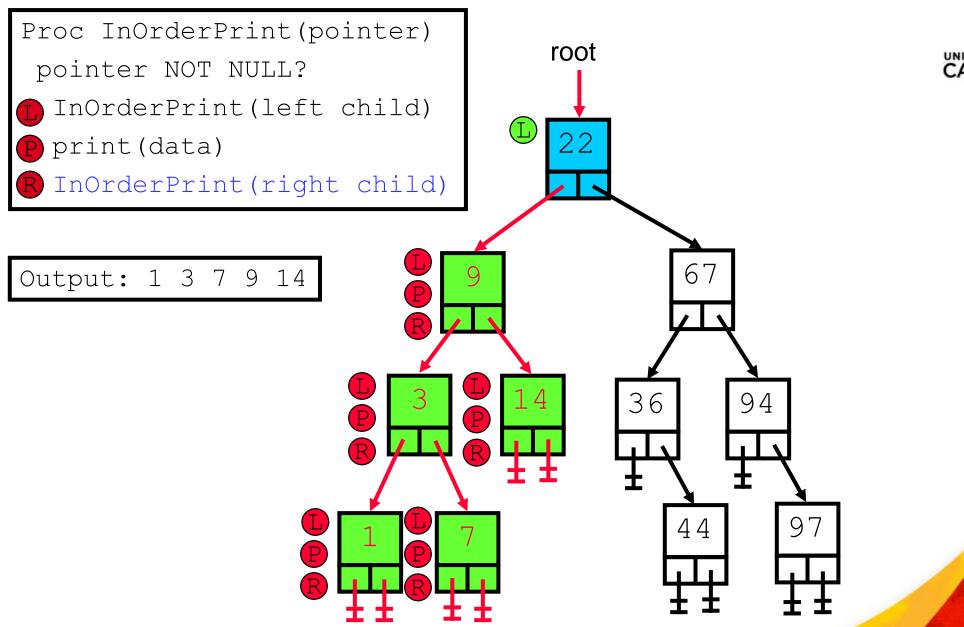




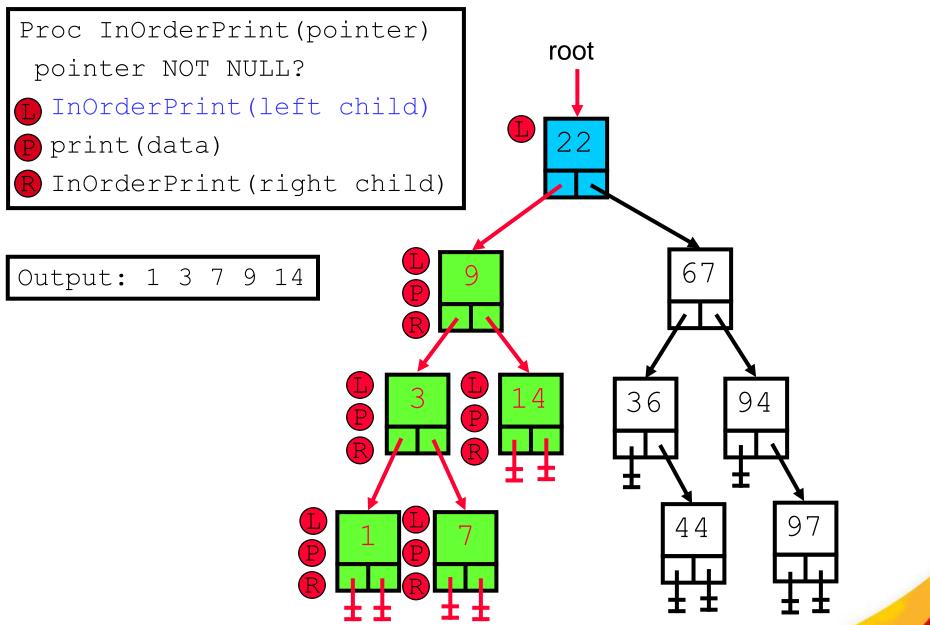




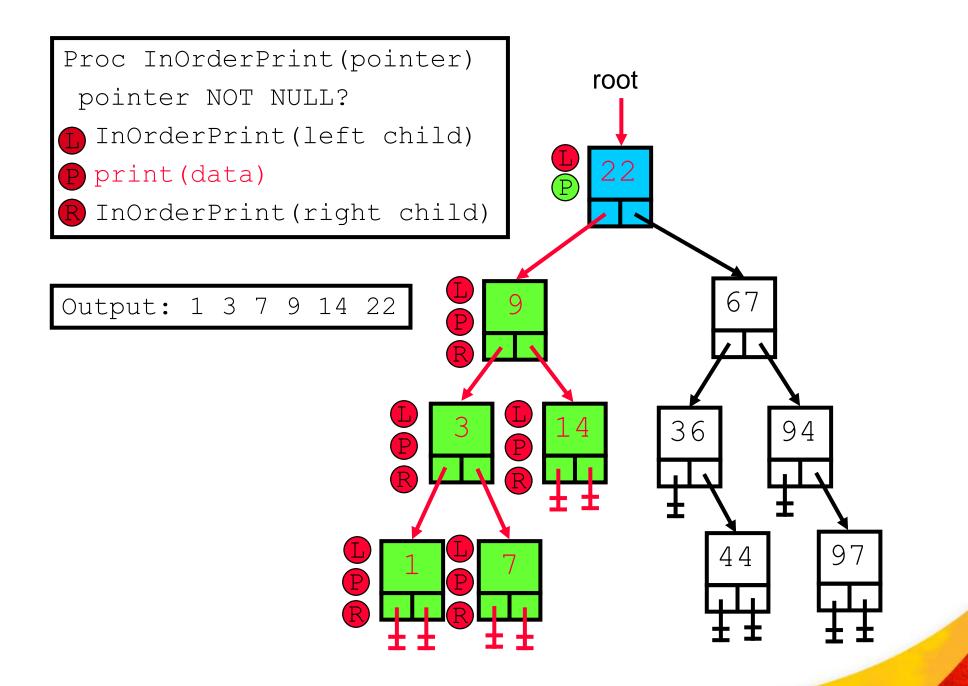




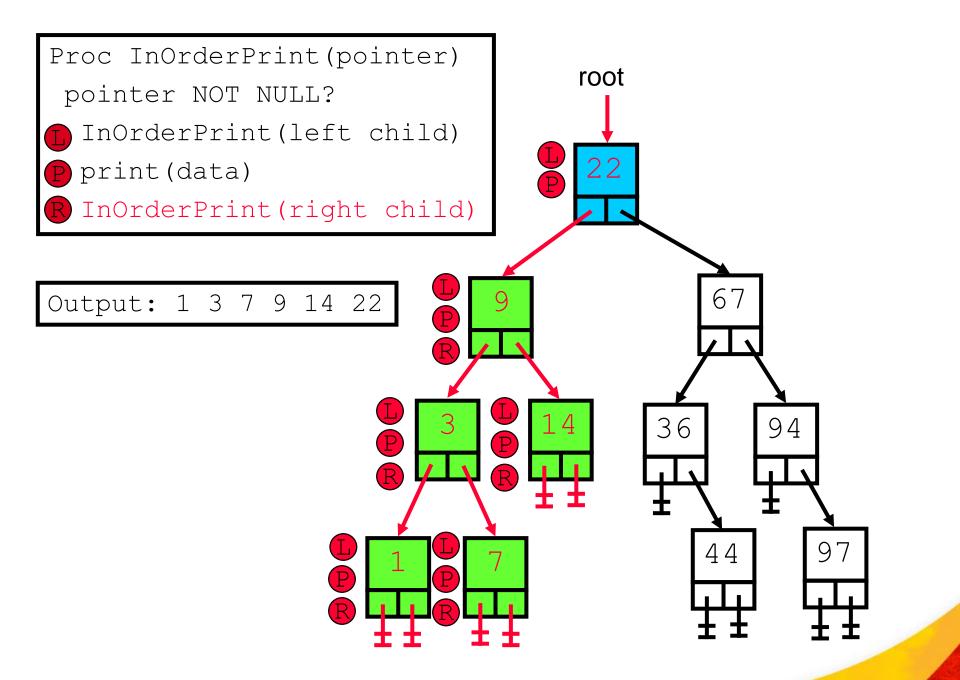




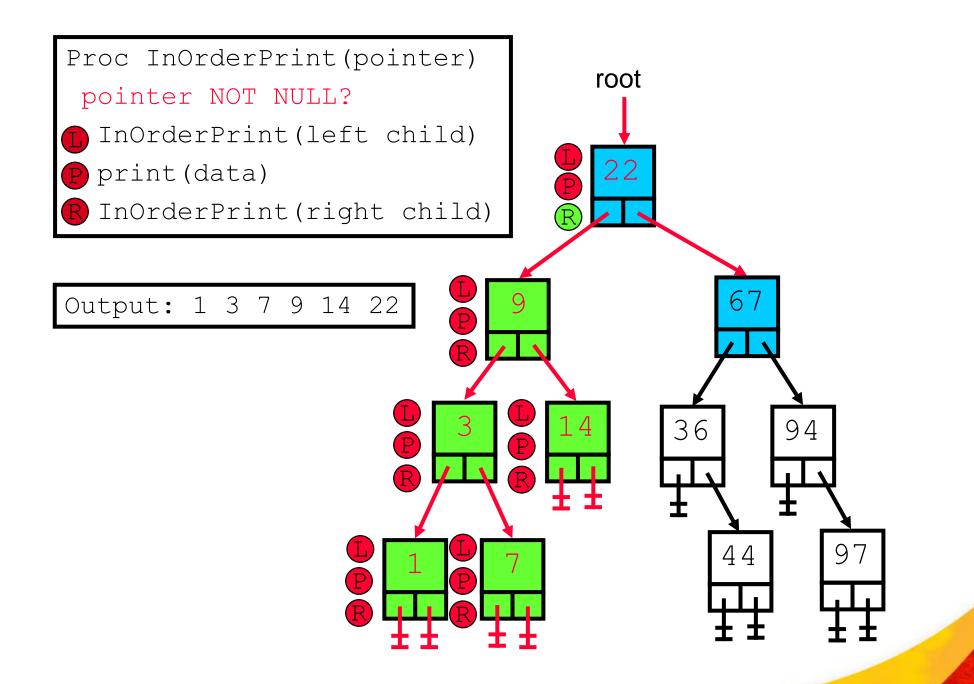




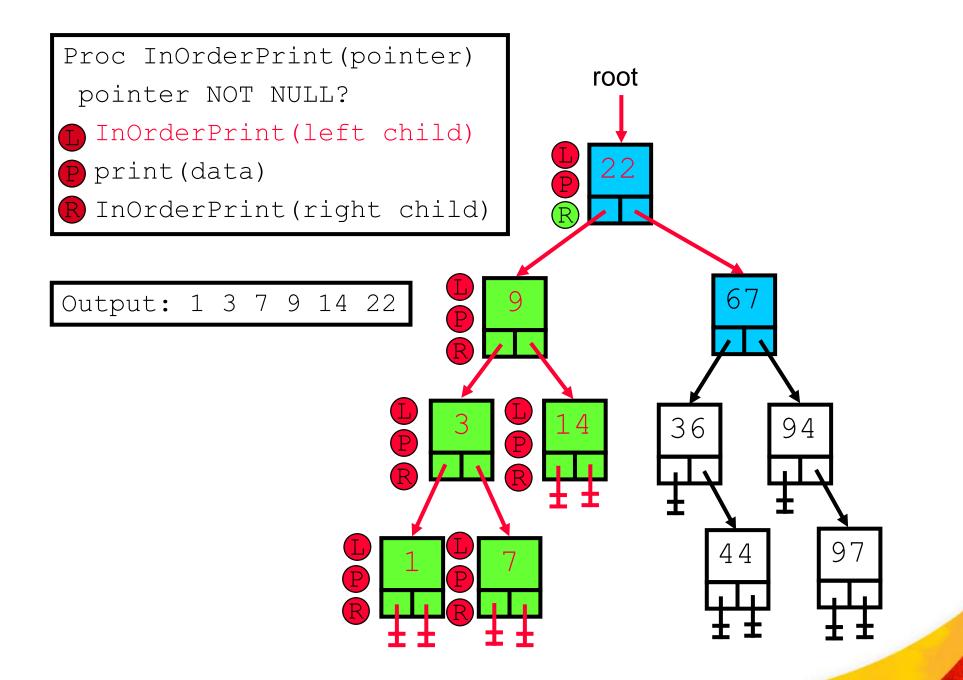




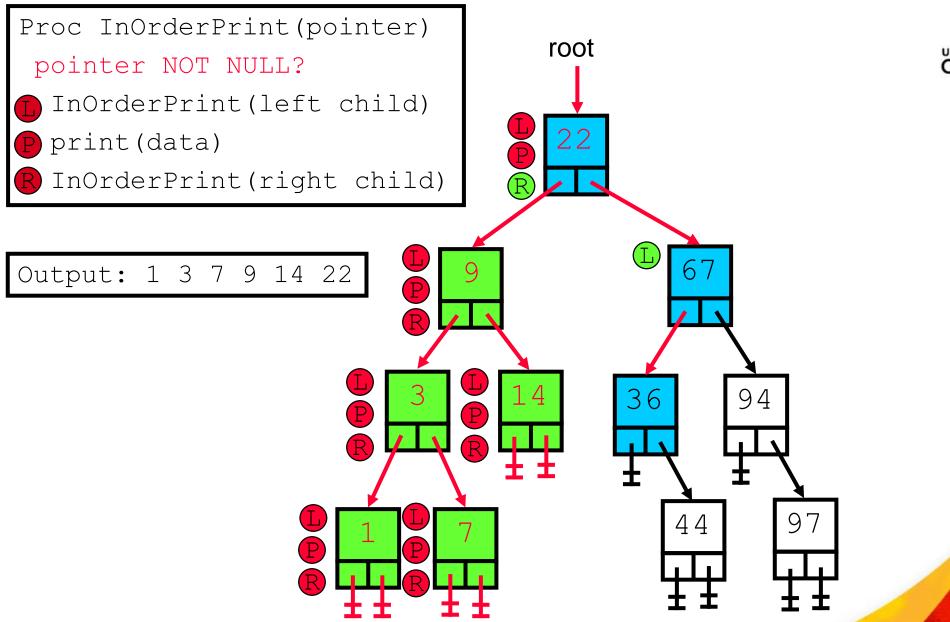




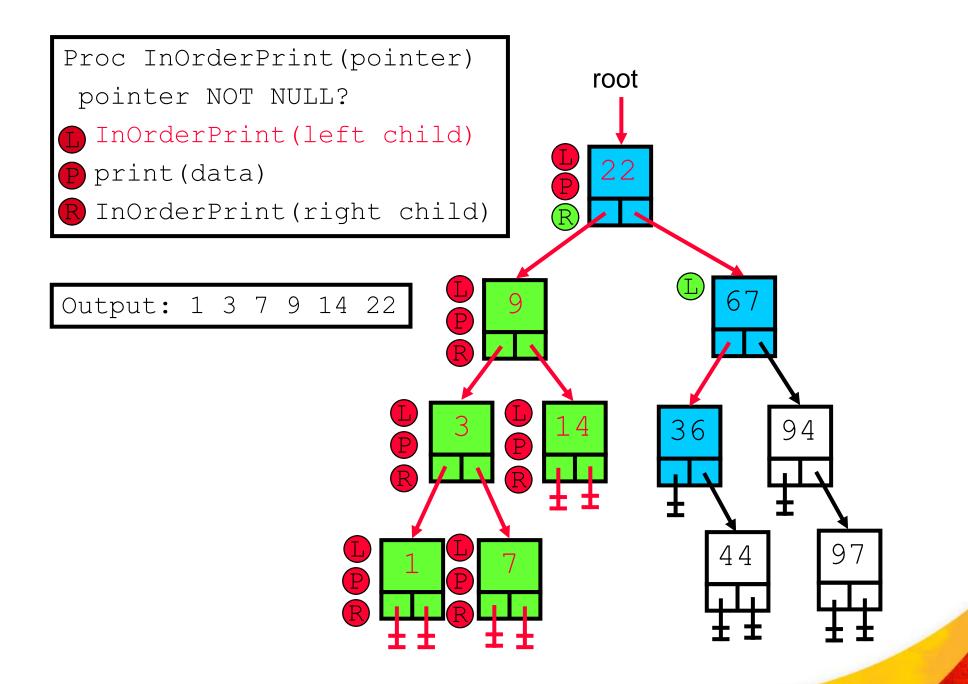




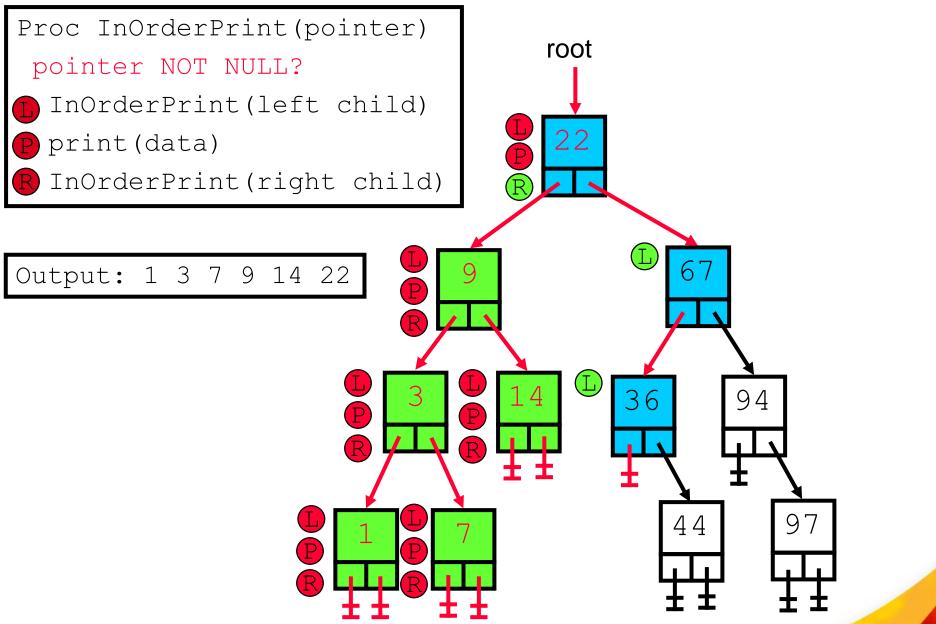




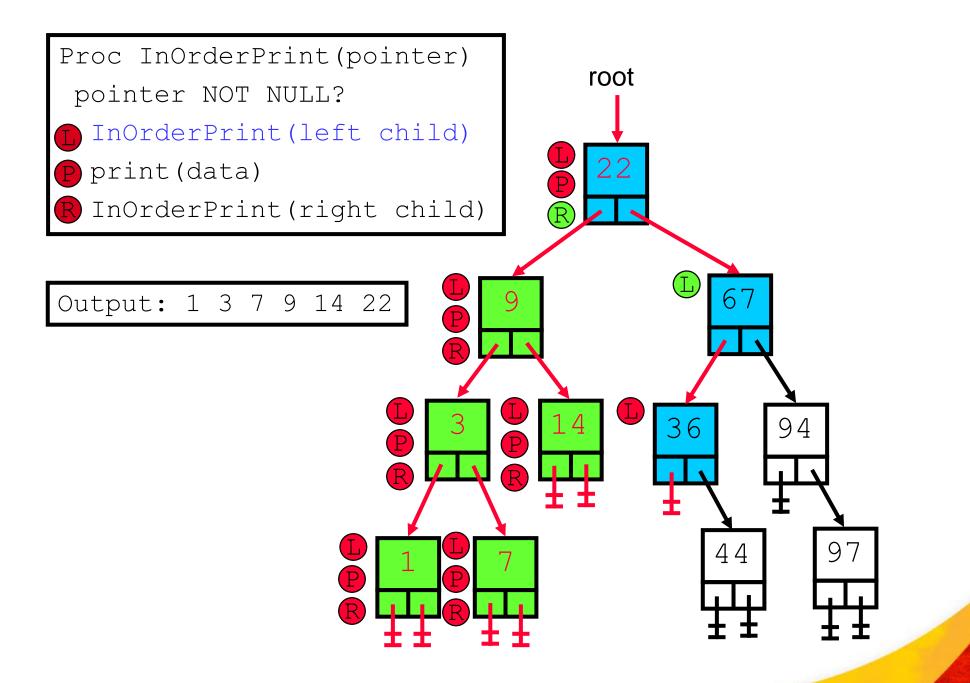




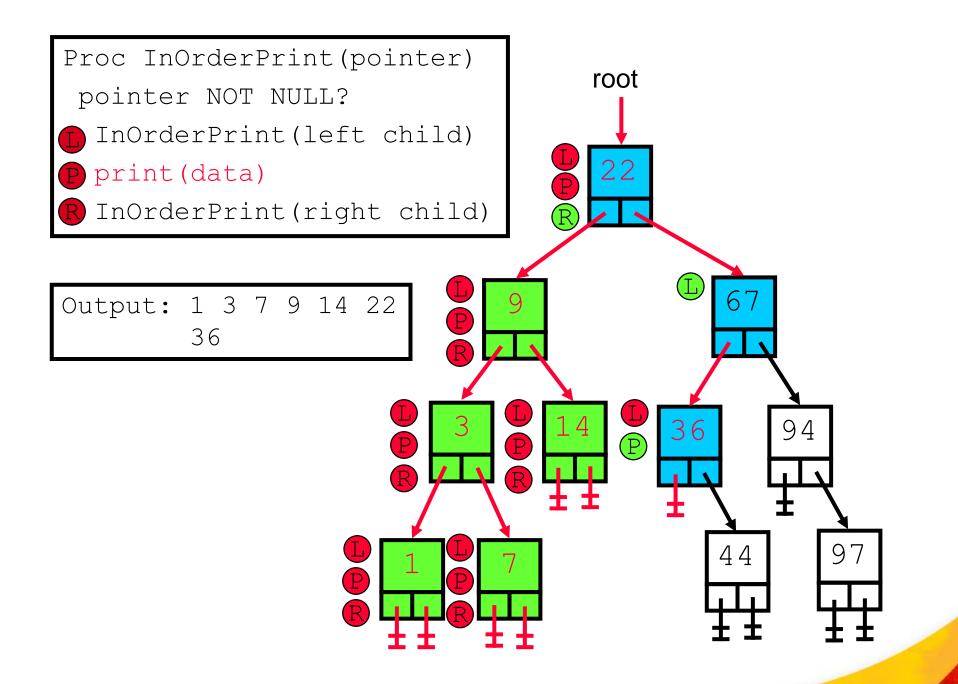




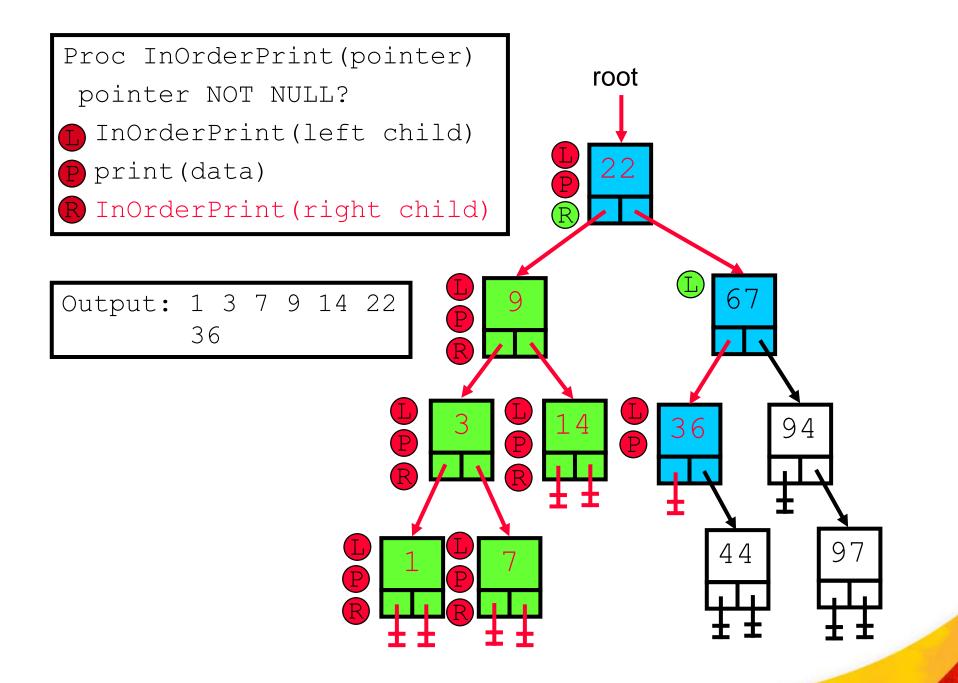




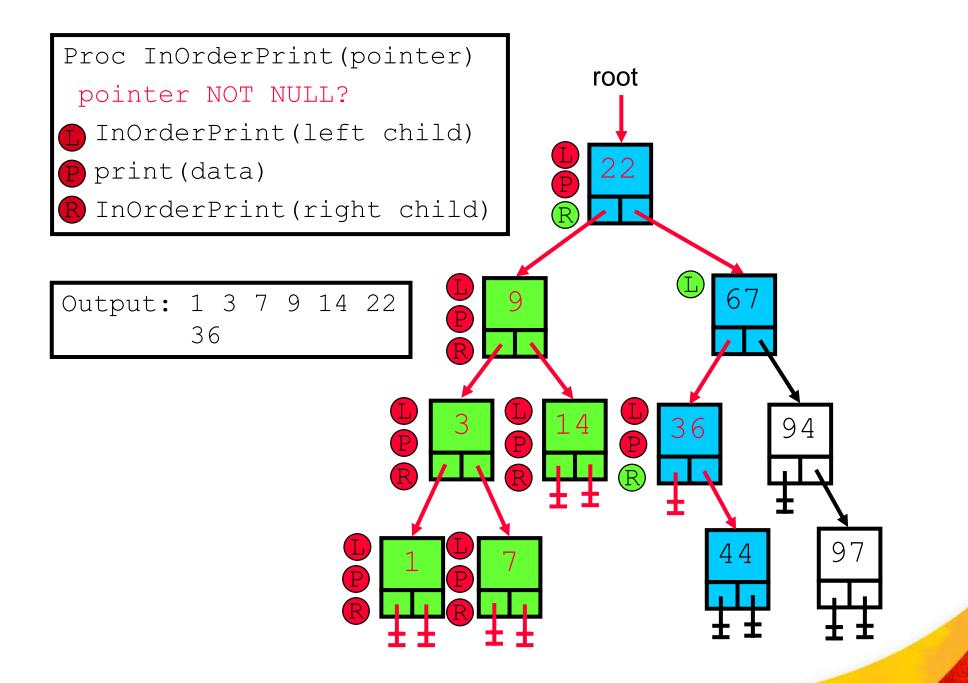




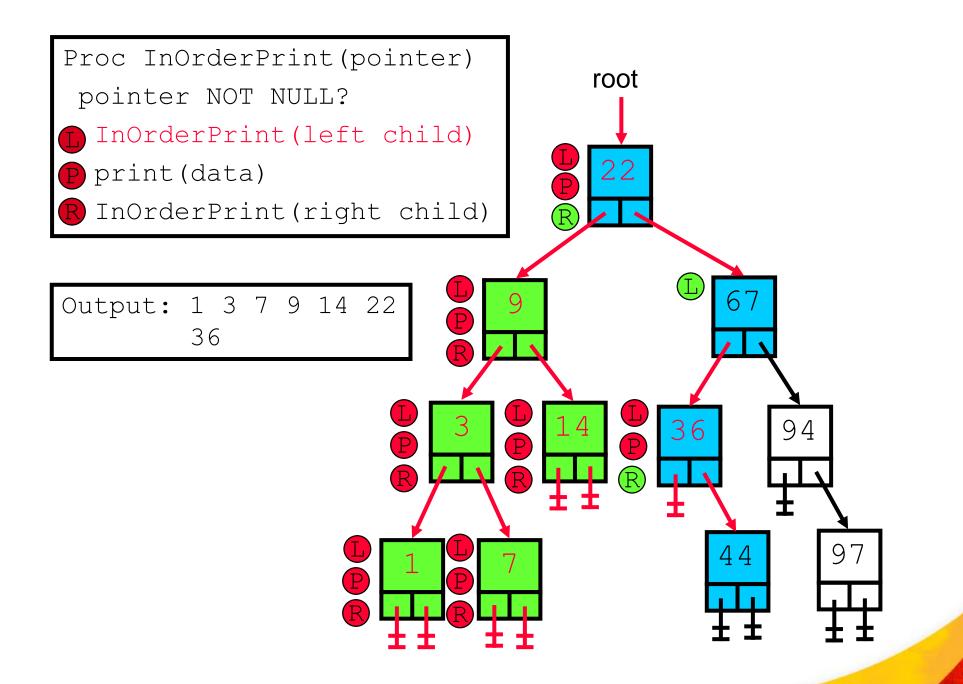




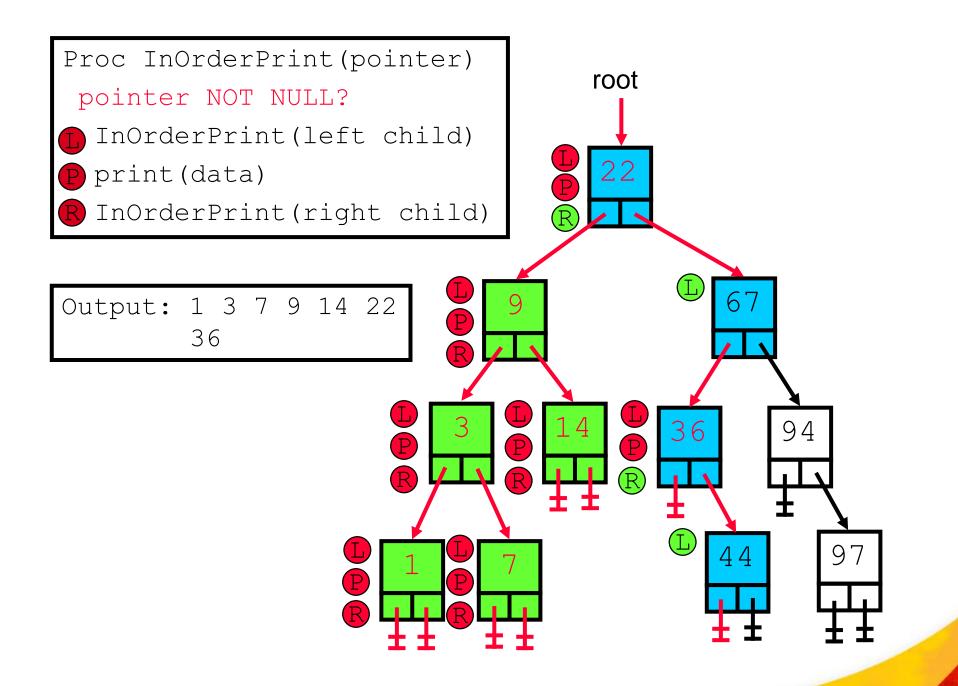




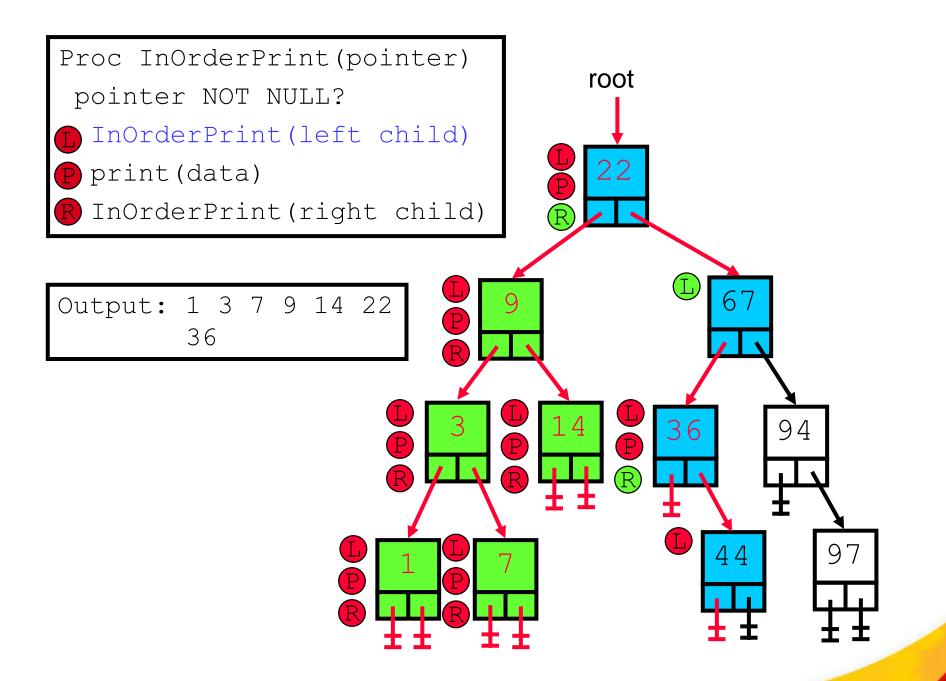




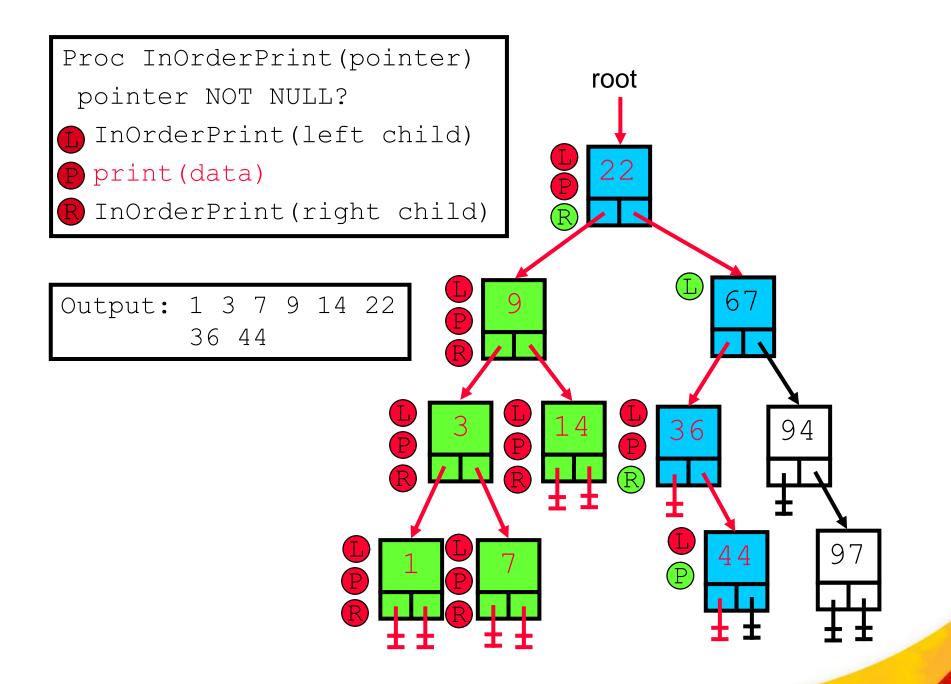




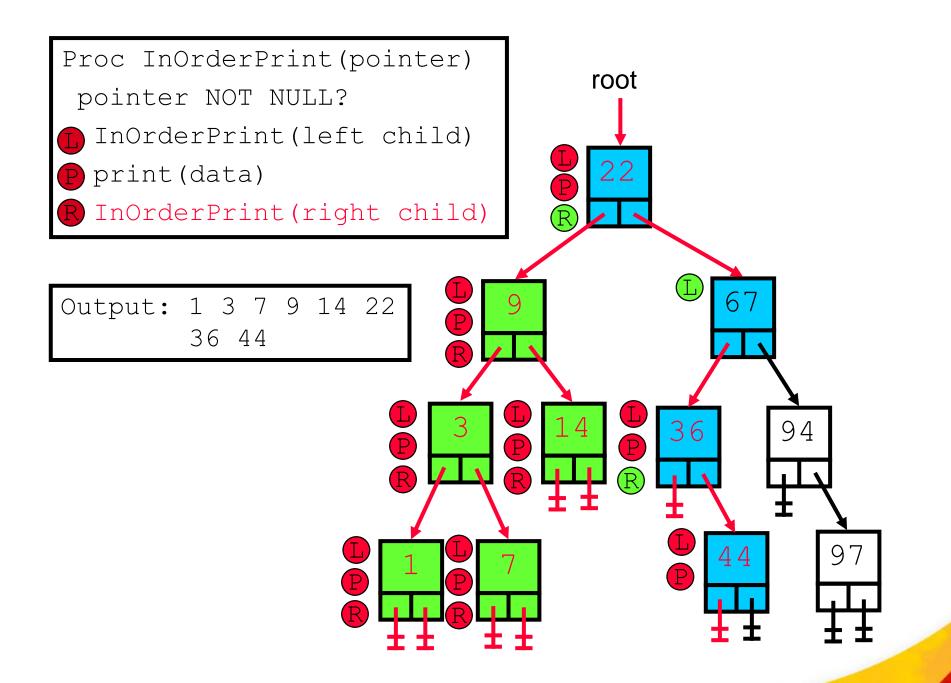




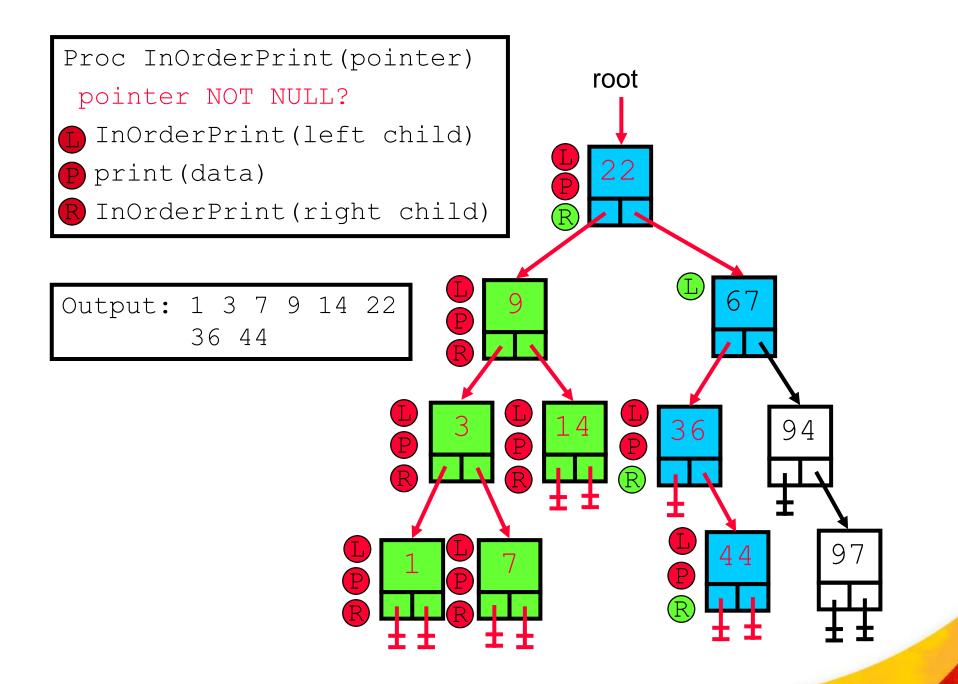




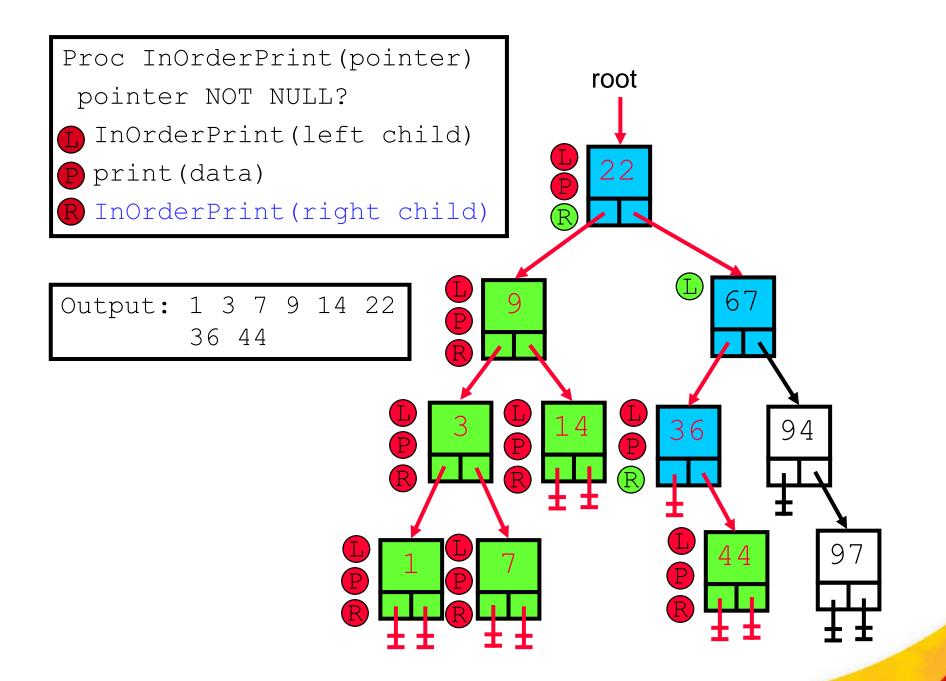




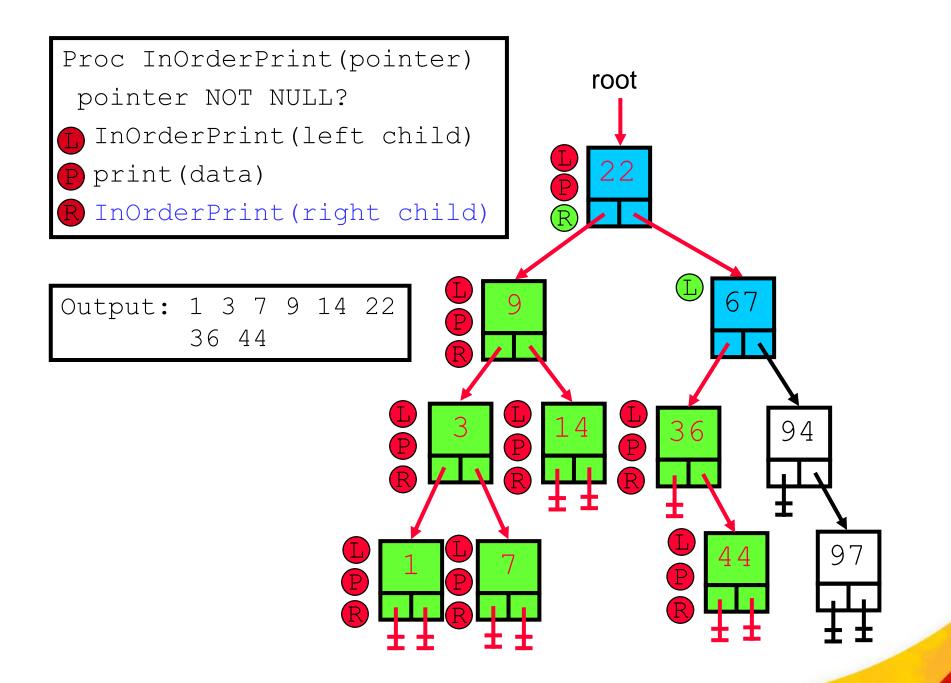




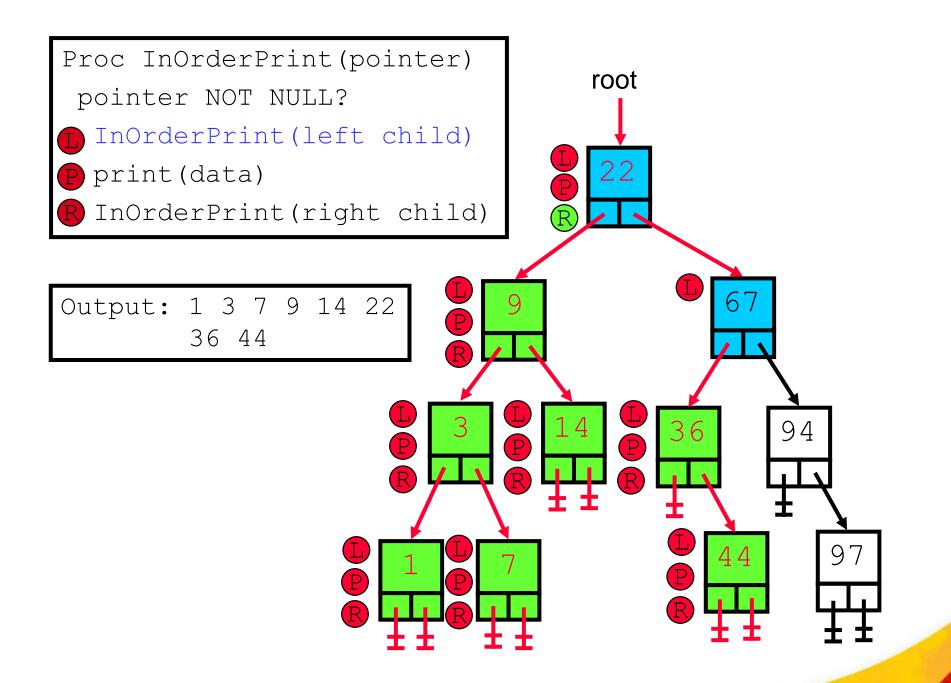




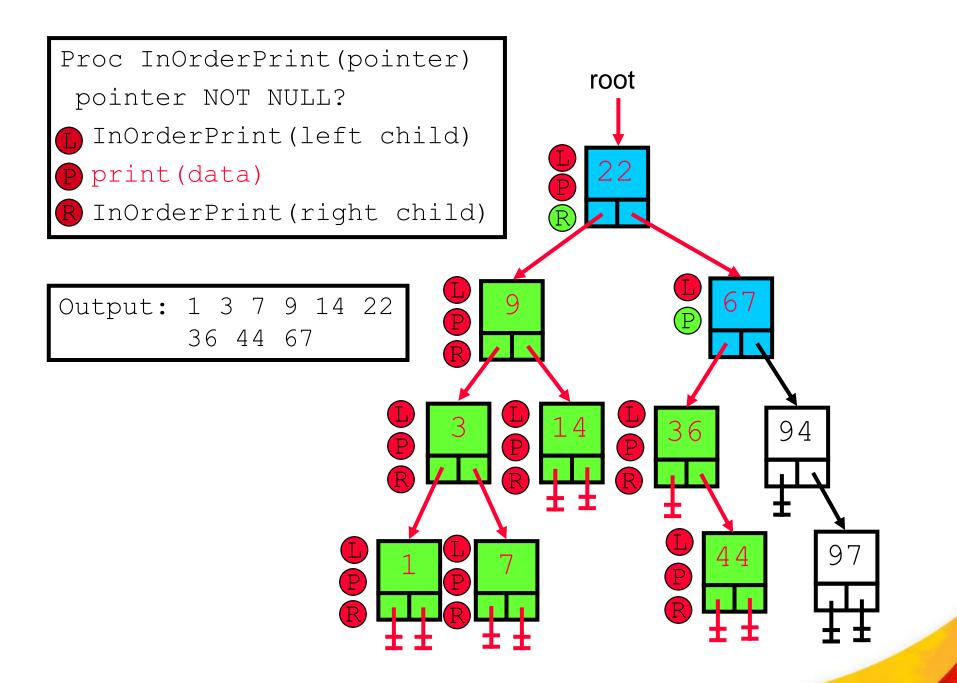




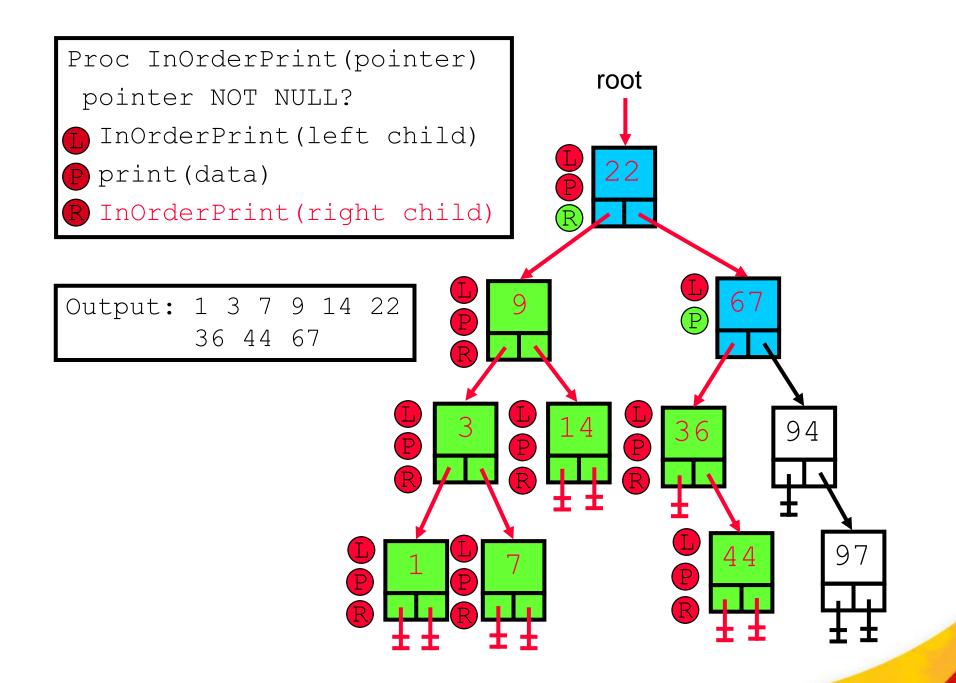




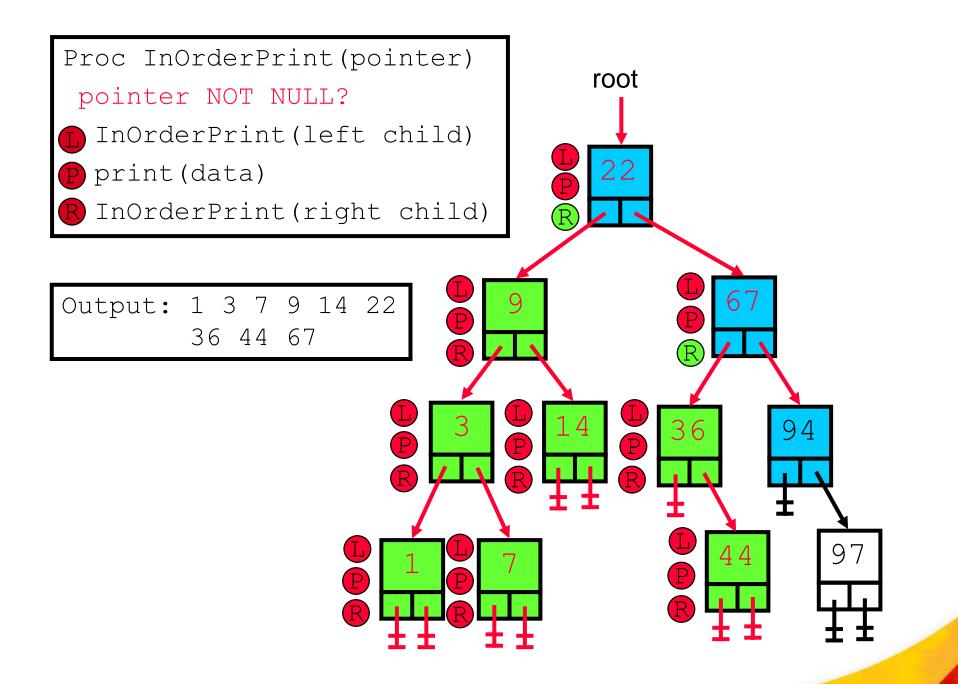




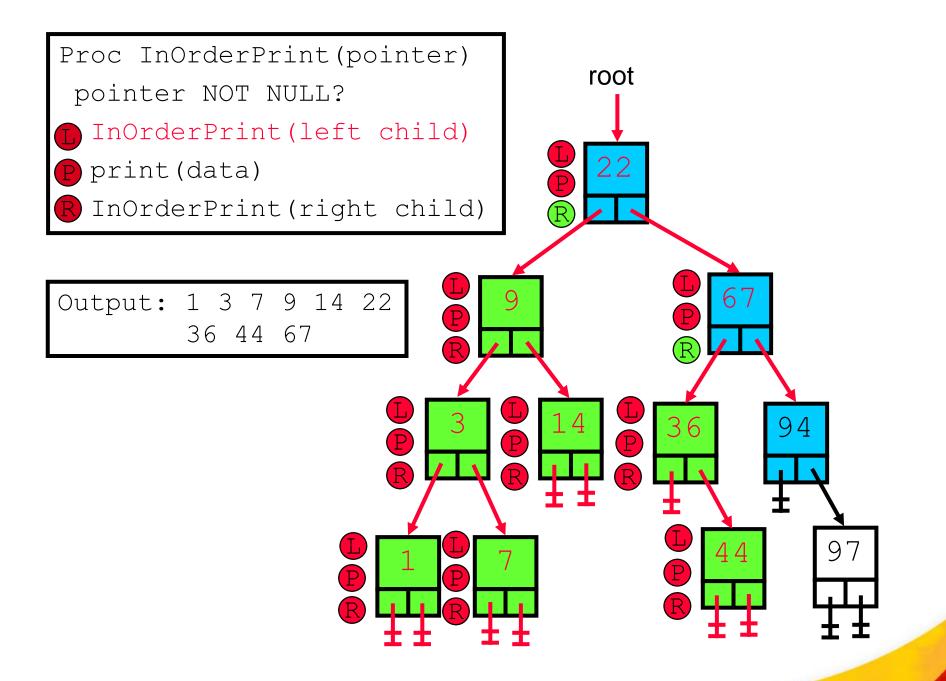




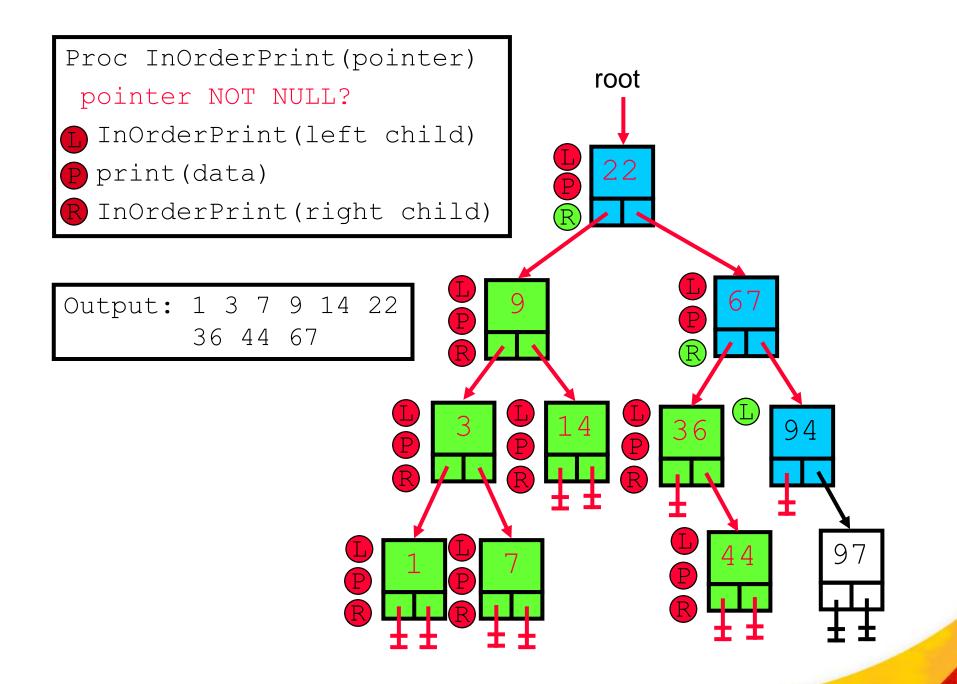




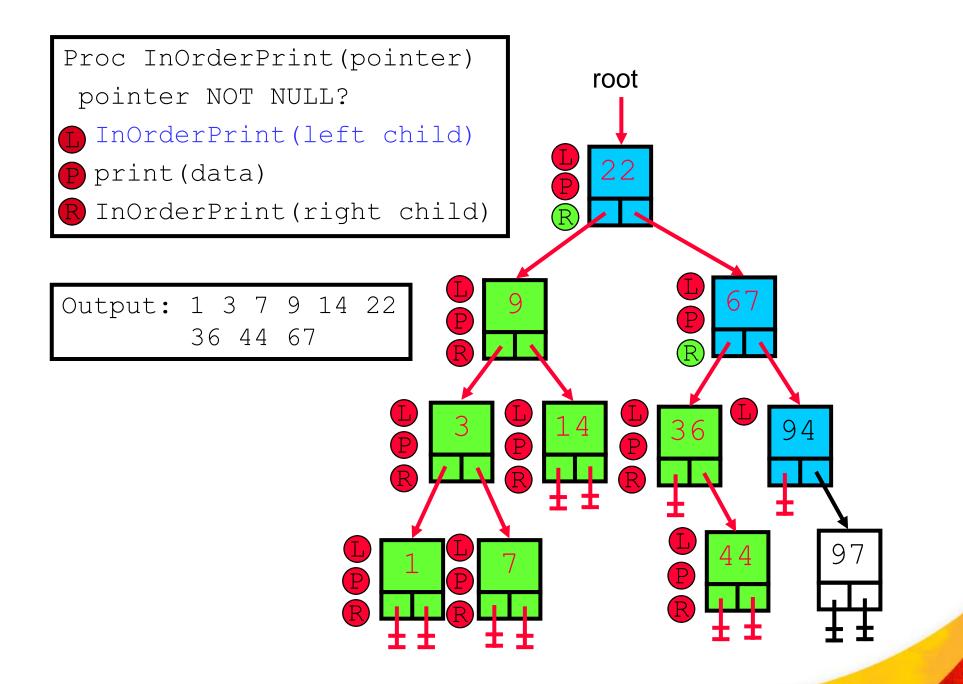




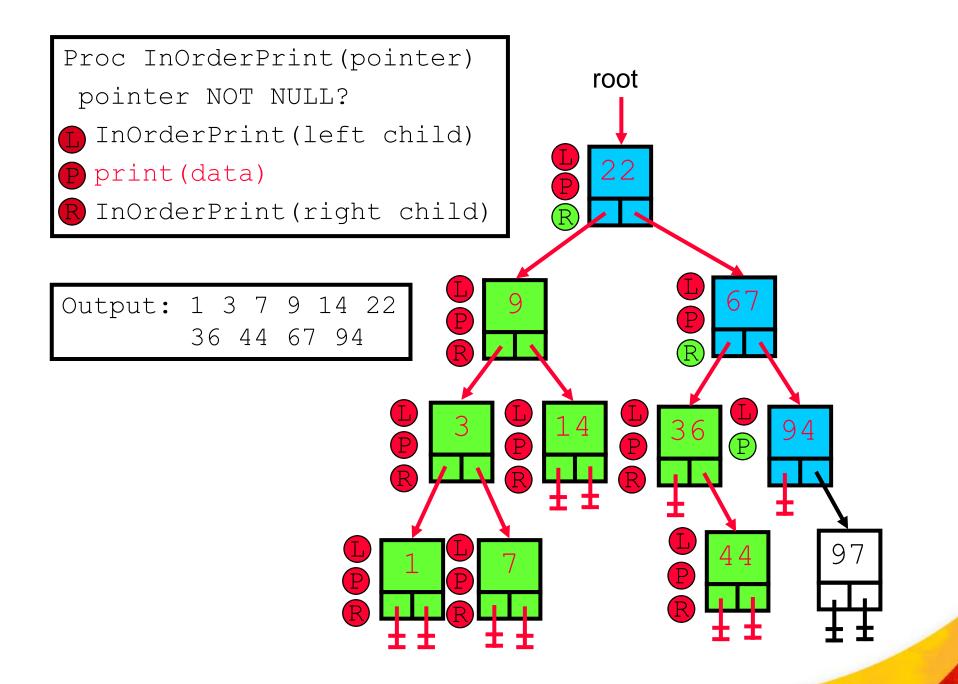




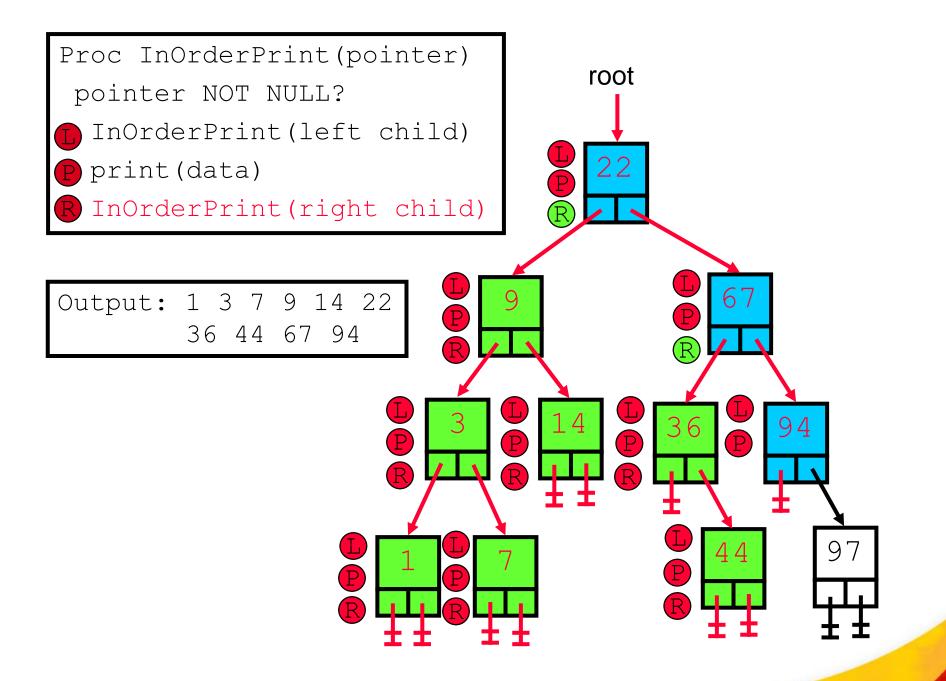




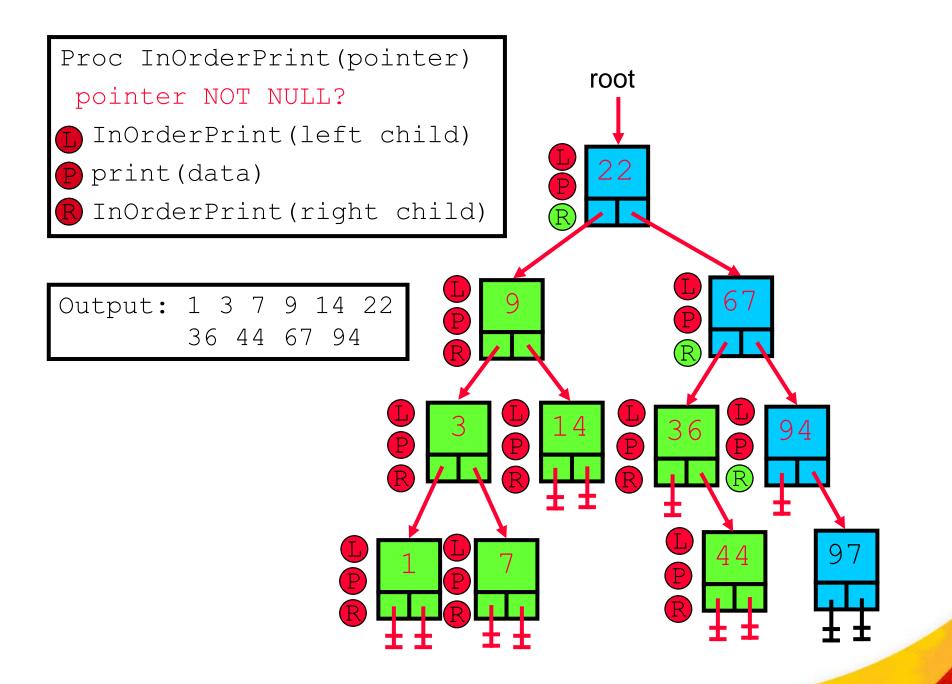




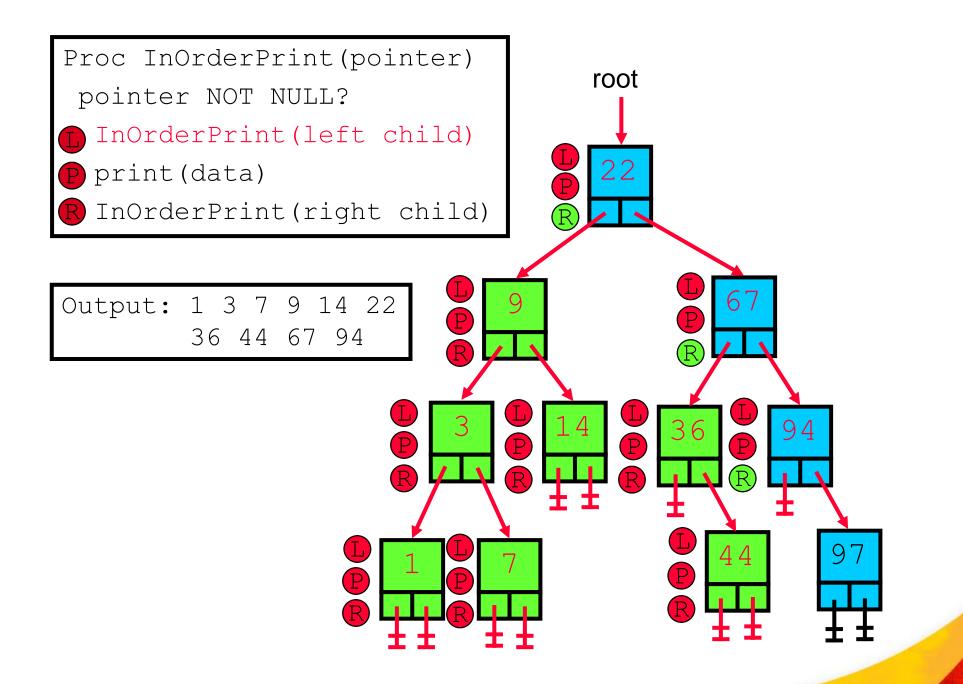




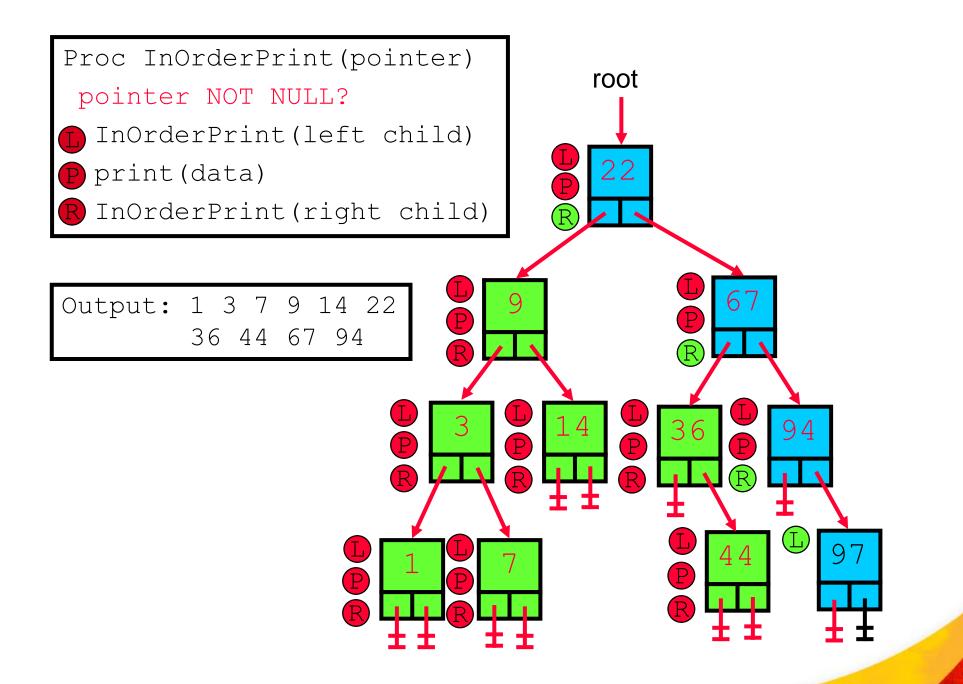




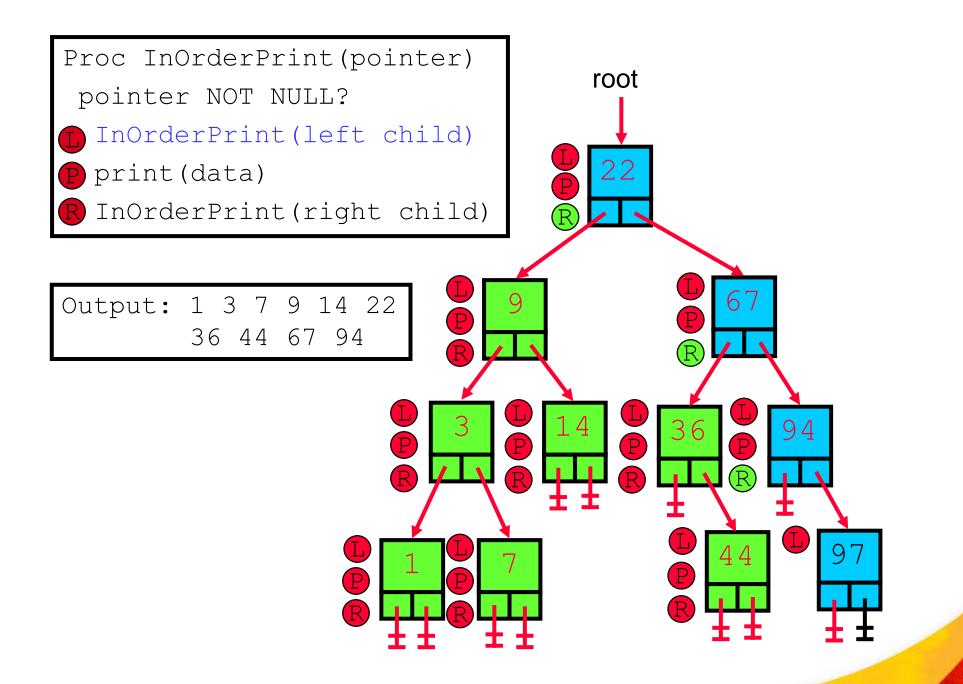




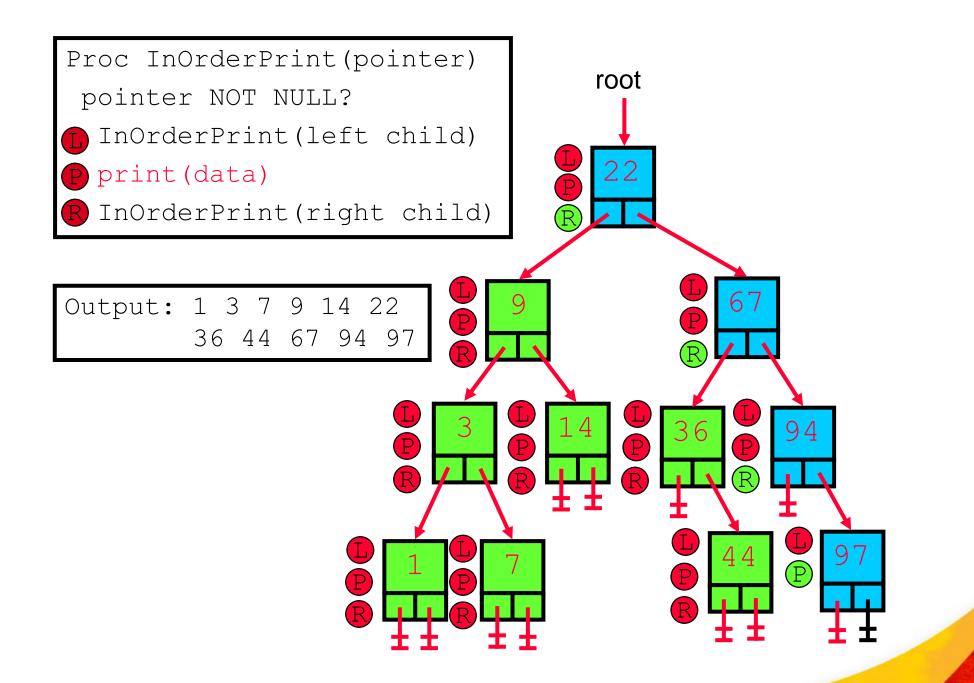




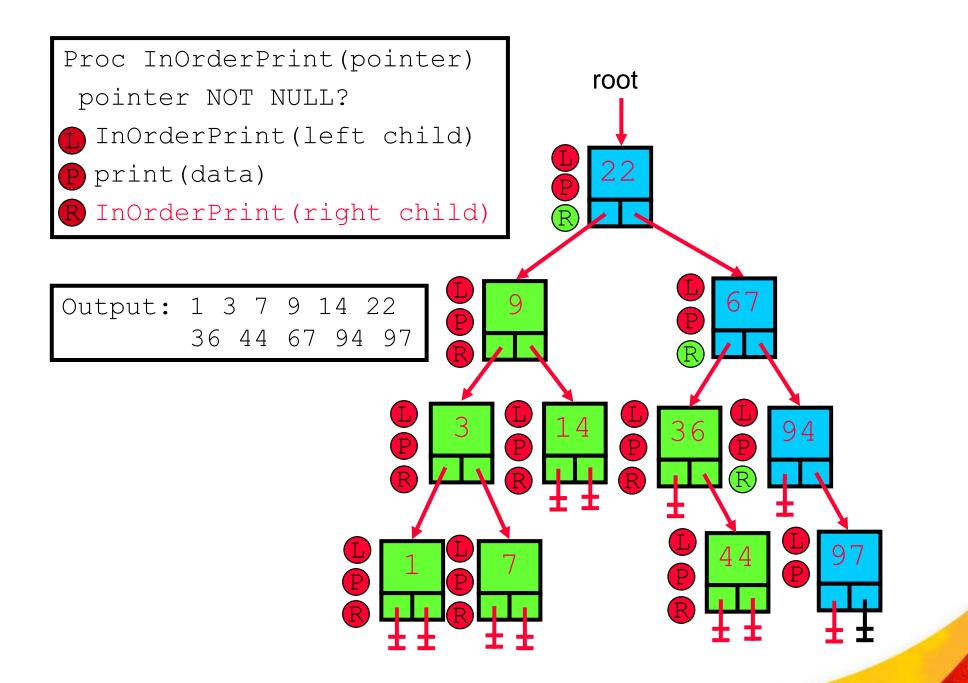




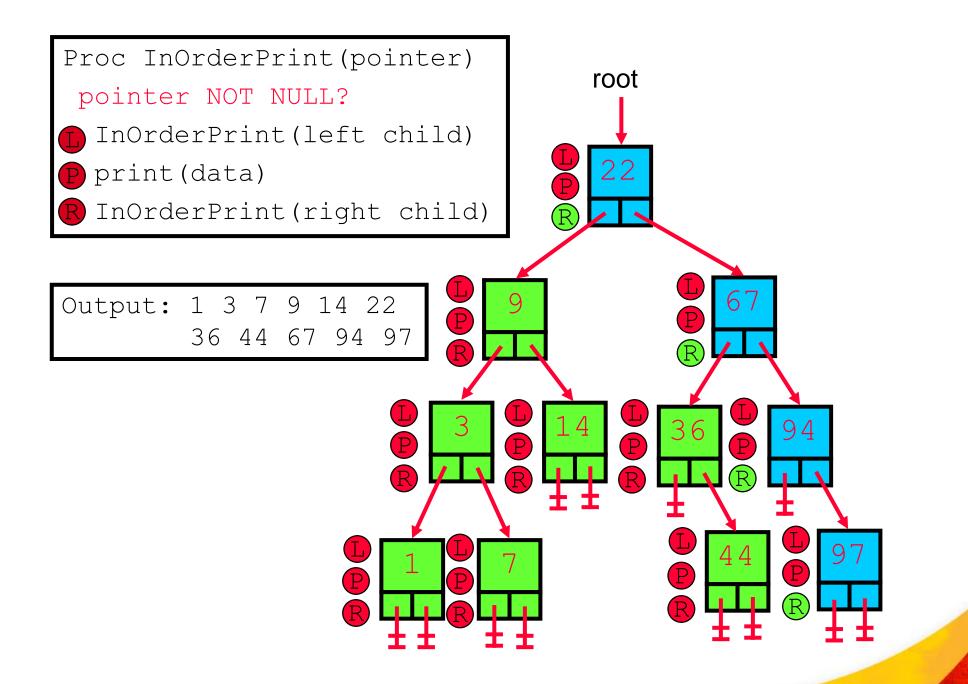




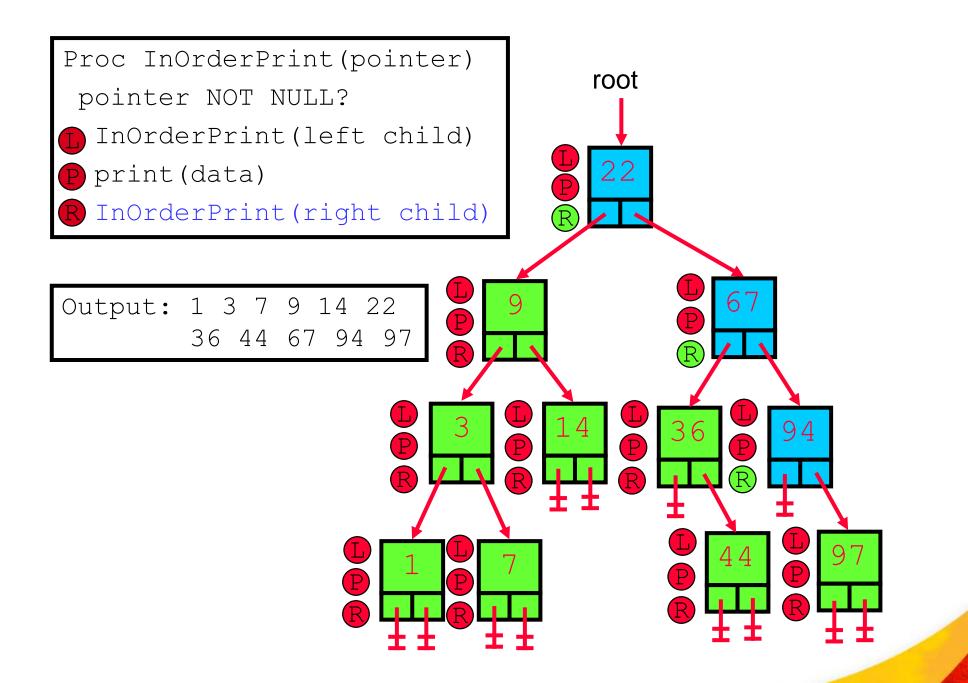




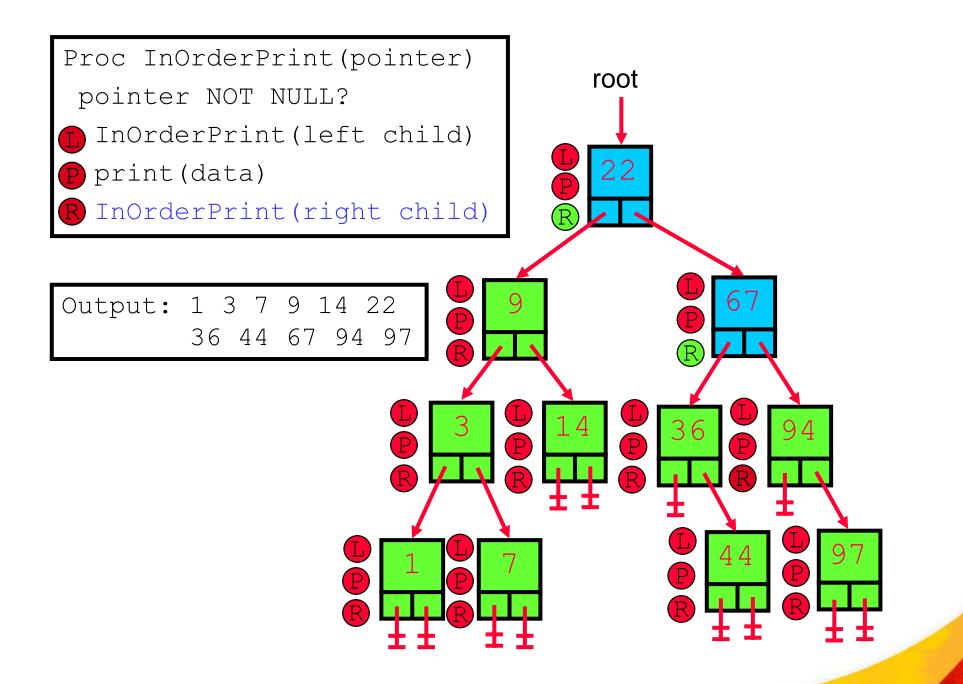




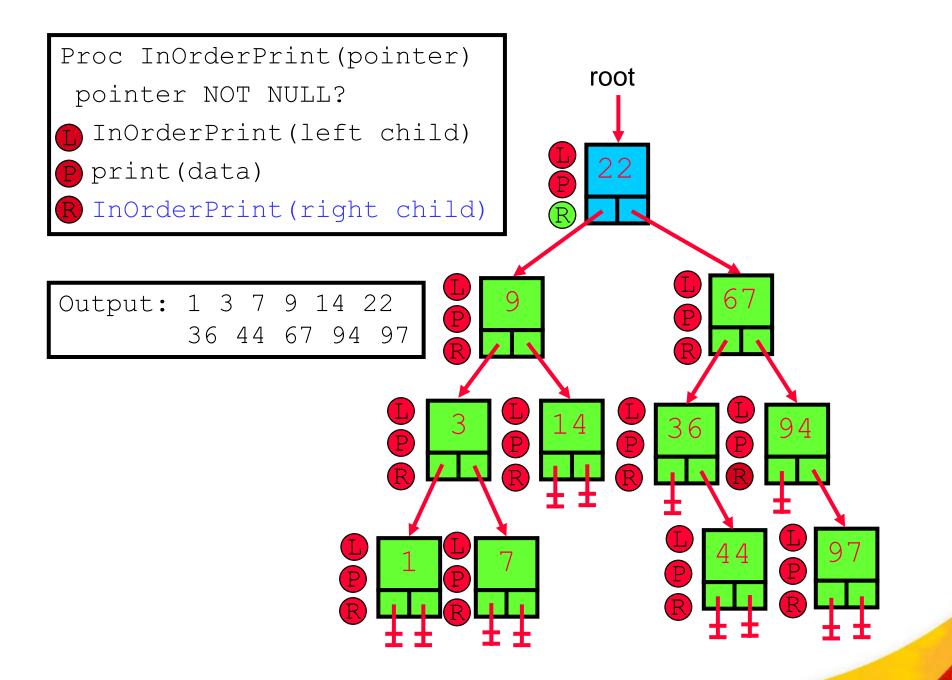




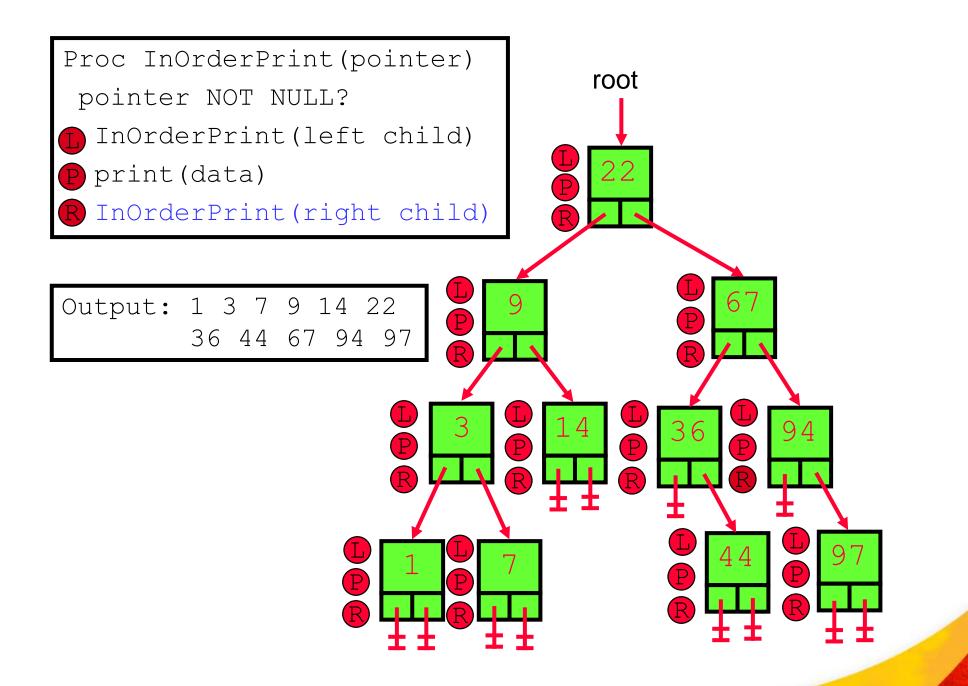




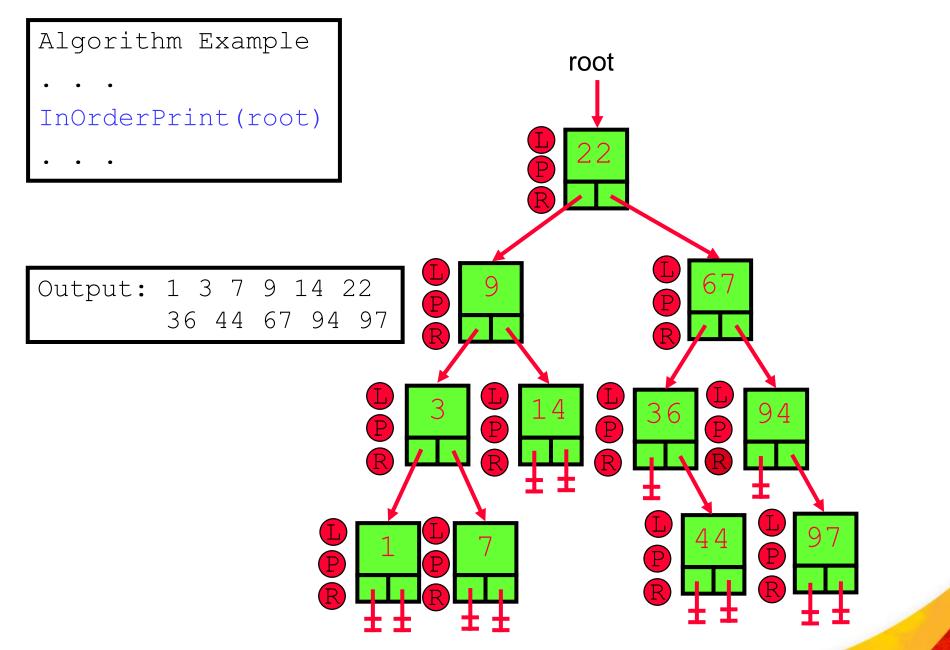


















- An In-Order traversal visits every node
  - Recurse left first
  - Do something with current
  - Recurse right last

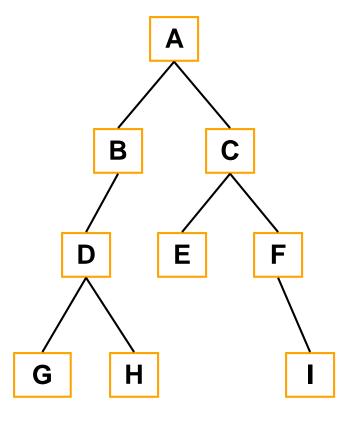
 The "left, current, right" logic is repeated recursively at every node.

 For a BST, an in-order traversal accesses the elements in ascending order.





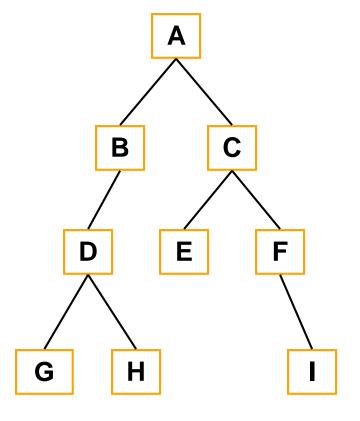
- What if we want to print everything on one level of the tree at a time?
- We can't do simple recursion
- We need to deal with all siblings on a level before we do any of those siblings children





How can we get the following output?

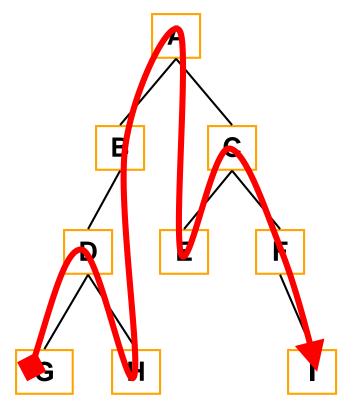
ABCDEFGHI





Inorder (left, node, right)?

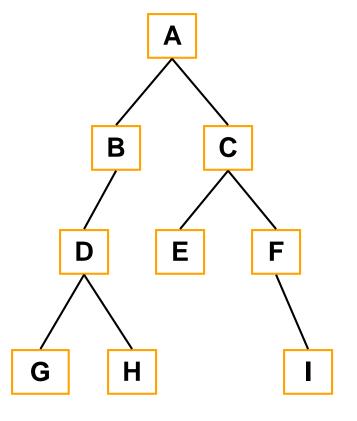
GDHBAECFI





Inorder (left, node, right)?

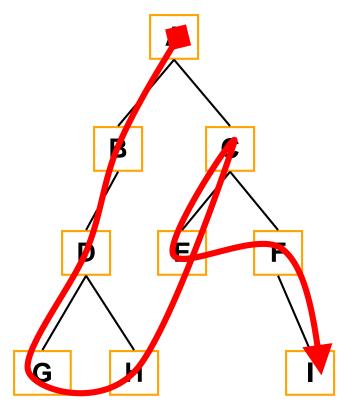
GDHBAECFI





Preorder (node, left, right)?

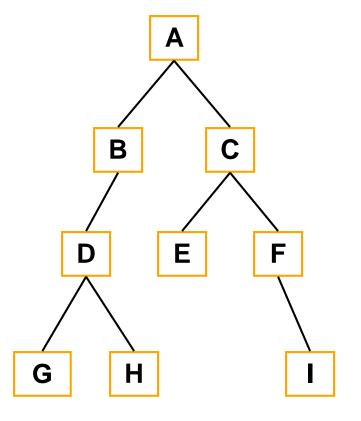
ABDGHCEFI





Preorder (node, left, right)?

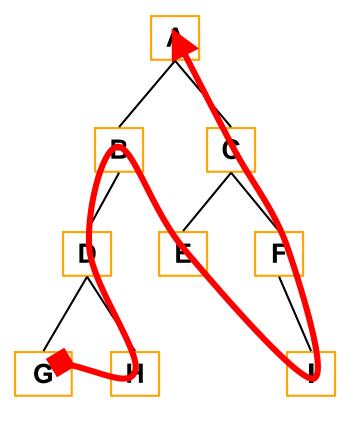
ABDGHCEFI





Postorder (left, right, node)?

GHDBEIFCA



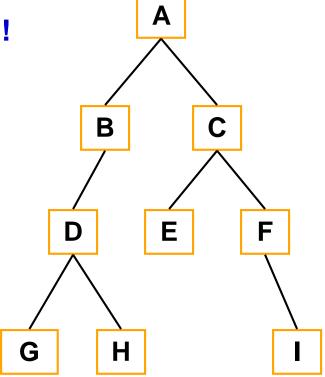


Postorder (left, right, node)?

GHDBEIFCA

We need a new traversal algorithm!



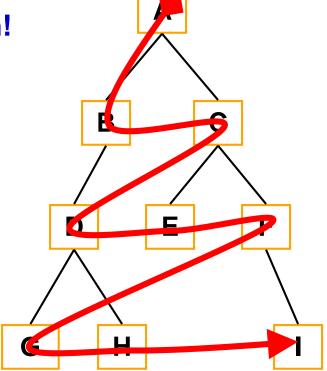


How can we get the following output?

ABCDEFGHI

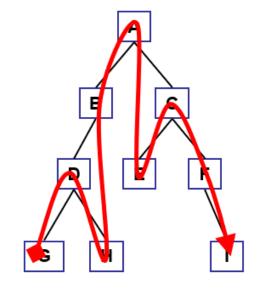
We need a new traversal algorithm!



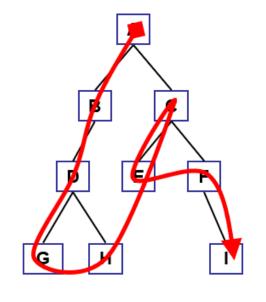


How can we get the following output?

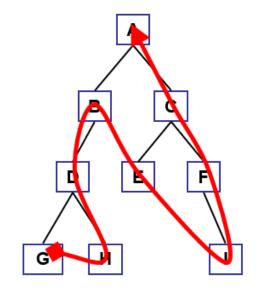
ABCDEFGHI



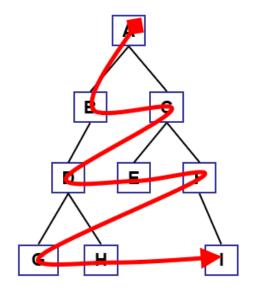




Pre-order



Post-order



Level-order







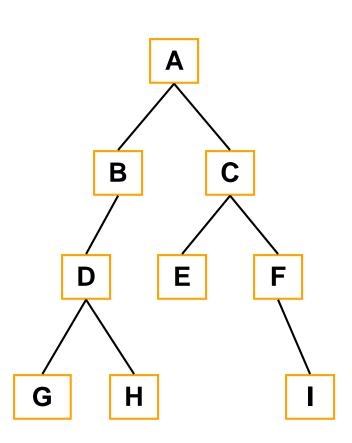
Use a queue to track unvisited nodes

- For each node that is dequeued,
  - enqueue each of its children
  - until queue empty

Also called: Breadth-First traversal

# LevelOrder

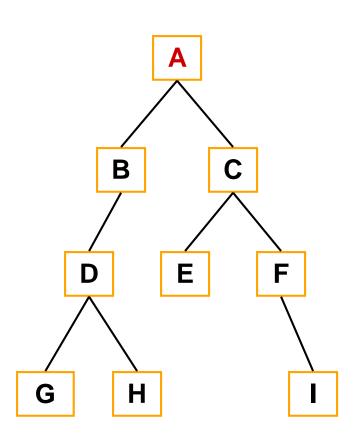




Queue Output

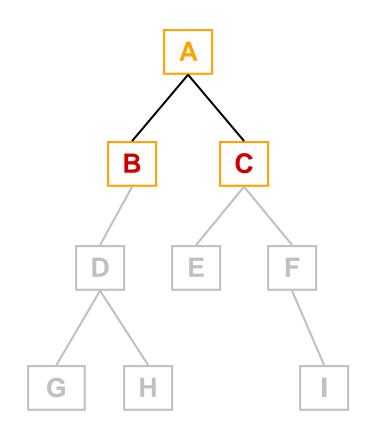
# LevelOrder





Queue Output Init [A] -

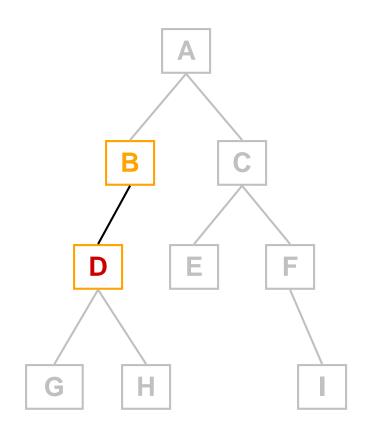




Queue Output Init [A] Step 1 [B,C] A

Dequeue A
Print A
Enqueue children of A





Queue Output Init [A] -

Step 1 [B,C] A

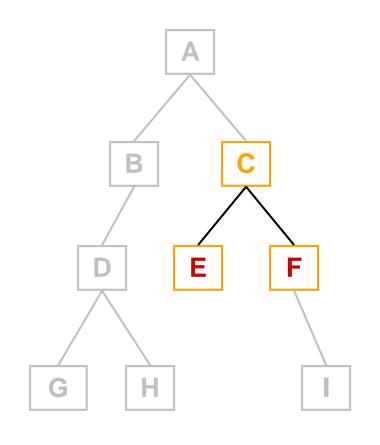
Step 2 [C,**D**] A **B** 

**Dequeue B** 

**Print B** 

**Enqueue children of B** 





Queue Output Init [A] Step 1 [B,C] A

Step 2 [C,D] AB

Step 3 [D,**E**,**F**] A B **C** 

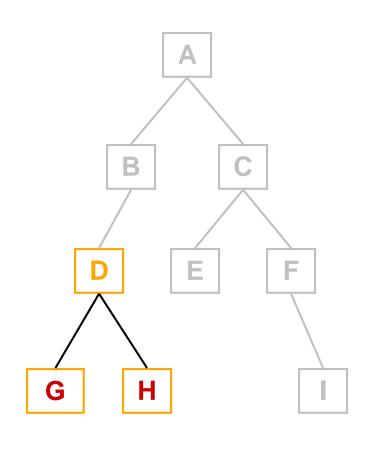
# **Dequeue C**

**Print C** 

**Enqueue children of C** 

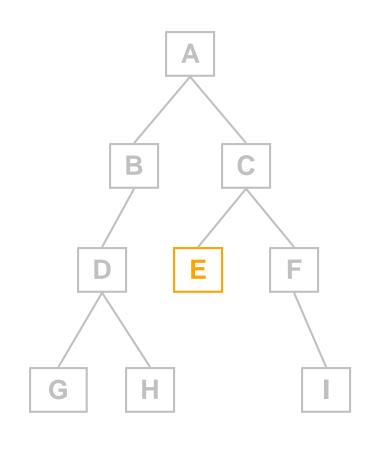
- - -





Init	Queue [A]	Output -
Step 1	[B,C]	Α
Step 2	[C,D]	AB
Step 3	[D,E,F]	ABC
Step 4	[E,F, <b>G</b> , <b>H</b> ]	A B C <b>C</b>

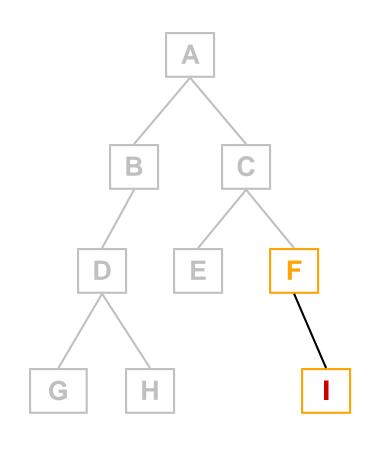




Init	Queue [A]	Output -
Step 1	[B,C]	Α
Step 2	[C,D]	АВ
Step 3	[D,E,F]	ABC
Step 4	[E,F,G,H]	ABCD

Step 5 [F,G,H] ABCD**E** 





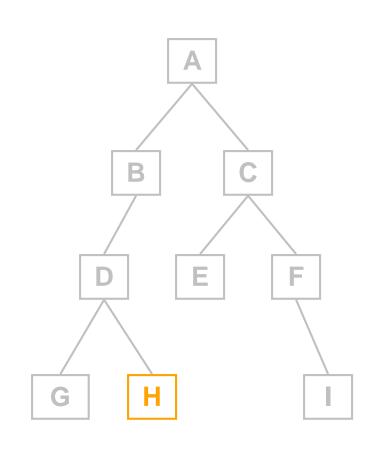
Init	Queue [A]	Output -
Step 1	[B,C]	Α
Step 2	[C,D]	АВ
Step 3	[D,E,F]	ABC
Step 4	[E,F,G,H]	ABCD
Step 5	[F,G,H]	ABCDE
Step 6	[G,H, <b>I</b> ]	ABCDE <b>F</b>



A	
ВС	
D E	F
GH	

Init	Queue [A]	Output -
Step 1	[B,C]	Α
Step 2	[C,D]	АВ
Step 3	[D,E,F]	ABC
Step 4	[E,F,G,H]	ABCD
Step 5	[F,G,H]	ABCDE
Step 6	[ <b>G</b> ,H,I]	ABCDEF
Step 7	[H,I]	ABCDEF <b>G</b>





Init	Queue [A]	Output -
Step 1	[B,C]	Α
Step 2	[C,D]	АВ
Step 3	[D,E,F]	ABC
Step 4	[E,F,G,H]	ABCD
Step 5	[F,G,H]	ABCDE
Step 6	[G,H,I]	ABCDEF
Step 7	[ <b>H</b> ,I]	ABCDEFG
Step 8	[1]	ABCDEFG <b>F</b>

# LevelOrder



	4
В	C
D	E F
G H	

Init	Queue [A]	Output -
Step 1	[B,C]	Α
Step 2	[C,D]	AB
Step 3	[D,E,F]	ABC
Step 4	[E,F,G,H]	ABCD
Step 5	[F,G,H]	ABCDE
Step 6	[G,H,I]	ABCDEF
Step 7	[H,I]	ABCDEFG
Step 8	[1]	ABCDEFGH
Step 9	[]	ABCDEFGHI

#### The Scenario



 We have a Binary Search Tree and want to remove some element based upon a match.

- Must preserve "search" property
- Must not lose any elements (i.e. only remove the one element)

#### **BST Deletion**



Search for desired item.

If not found, then return NULL or print error.

 If found, perform steps necessary to accomplish removal from the tree.





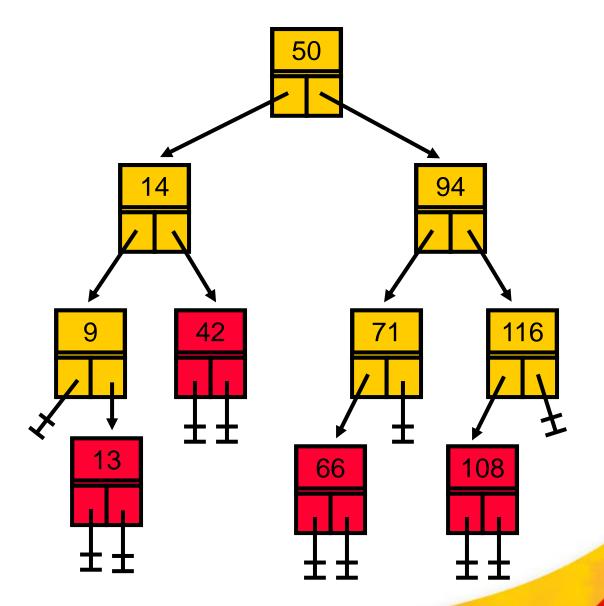
- Delete a leaf node
- Delete a node with only one child (left)
- Delete a node with only one child (right)
- Delete a node with two children
- Cases 2 and 3 are comparable and only need slight changes in the conditional statement used





Set the parent node's child pointer to null

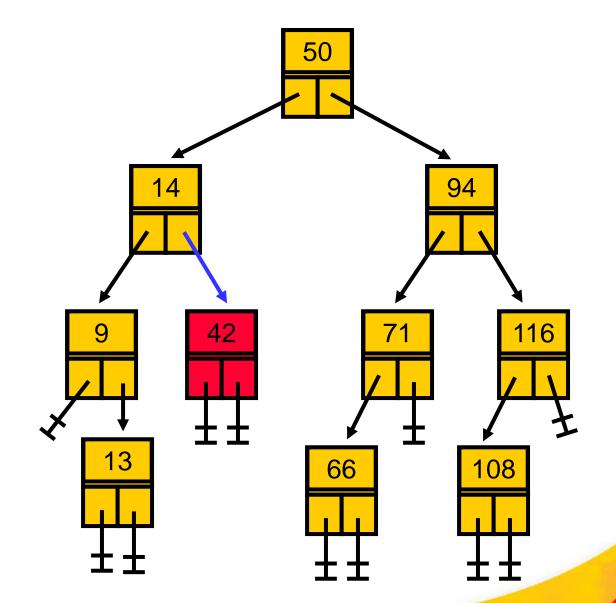
This will remove the node from the tree.



## **Delete a Leaf Node**

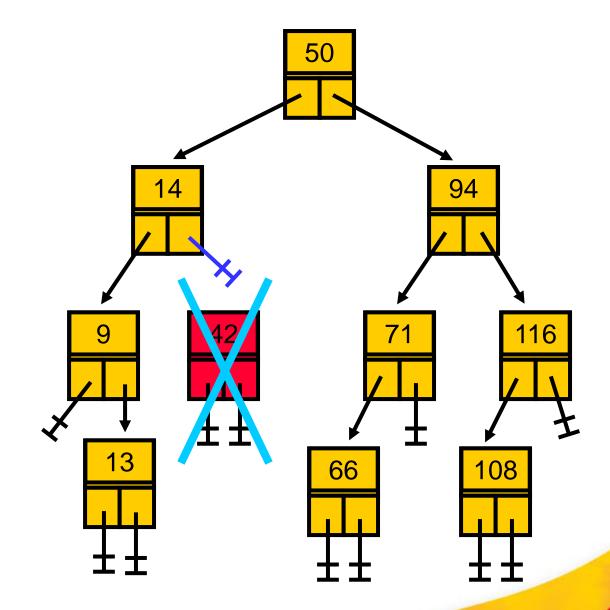


Let's delete 42.



# **Delete a Leaf Node**

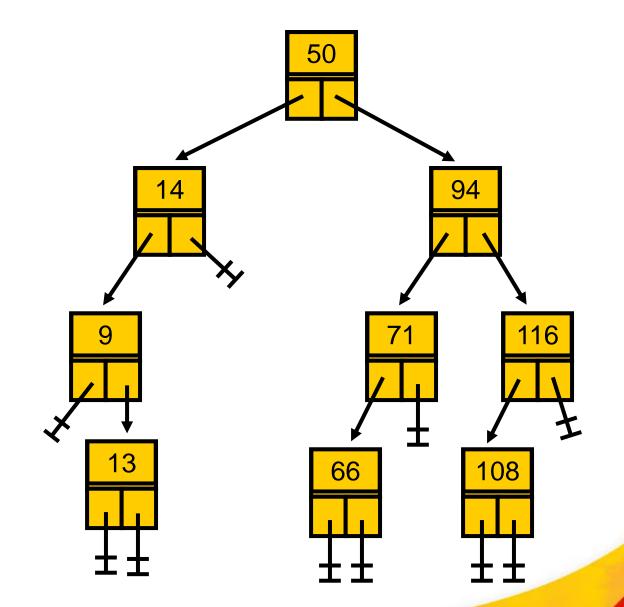




## **Delete a Leaf Node**



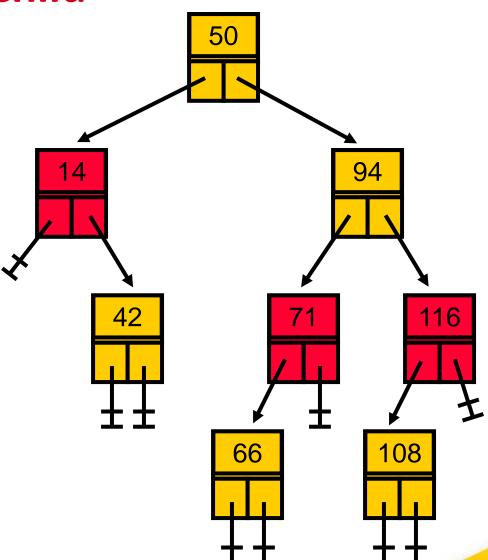
The resulting tree.





Determine if it has a left or a right child.

Point the current pointer to the appropriate child:

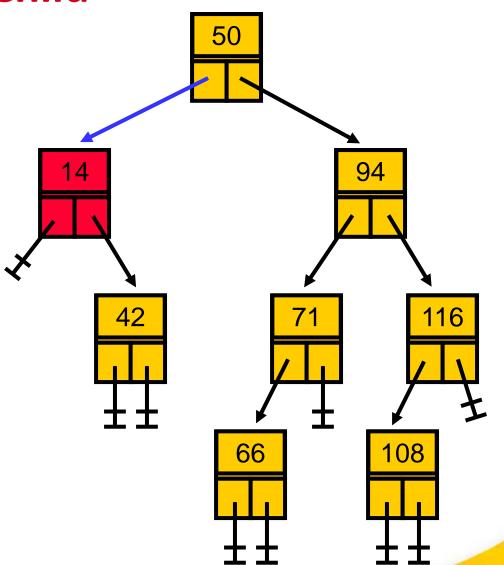




Determine if it has a left or a right child.

Point the current pointer to the appropriate child:

Let's delete 14.



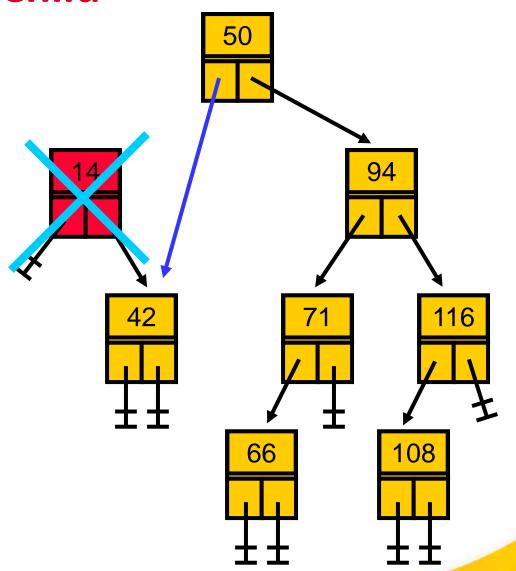


Determine if it has a left or a right child.

Point the current pointer to the appropriate child:

cur <- cur^.right child</pre>

Move the pointer; now nothing points to the node.



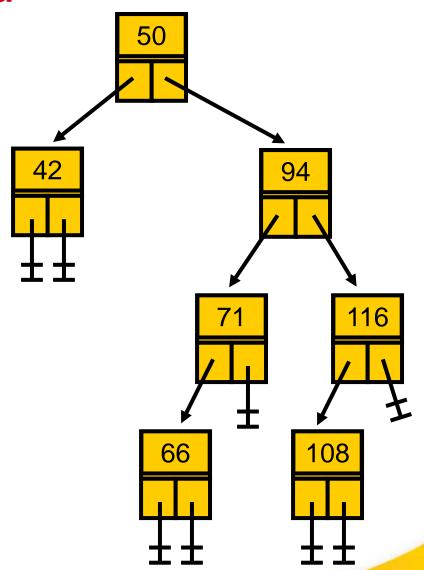


Determine if it has a left or a right child.

Point the current pointer to the appropriate child:

cur <- cur^.right\_child</pre>

The resulting tree.



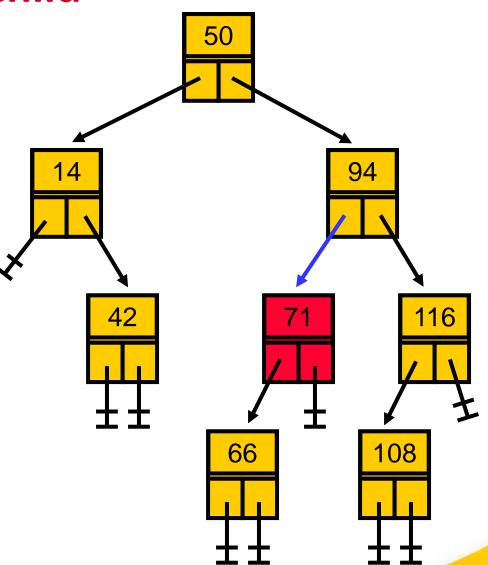


Determine if it has a left or a right child.

Point the current pointer to the appropriate child:

cur <- cur^.left\_child
 Or
cur <- cur^.right\_child</pre>

Let's delete 71.



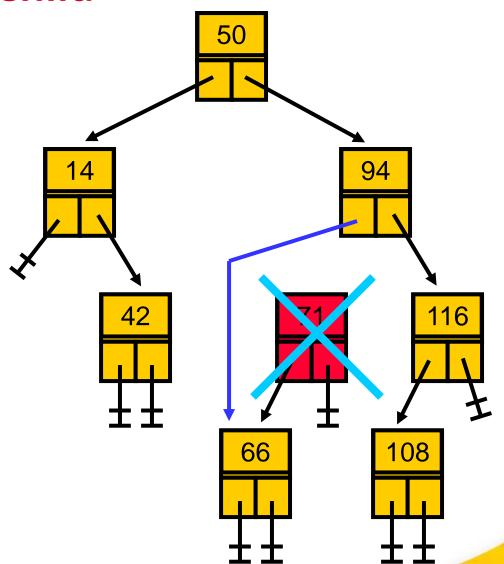


Determine if it has a left or a right child.

Point the current pointer to the appropriate child:

cur <- cur^.left\_child</pre>

Move the pointer; now nothing points to the node.



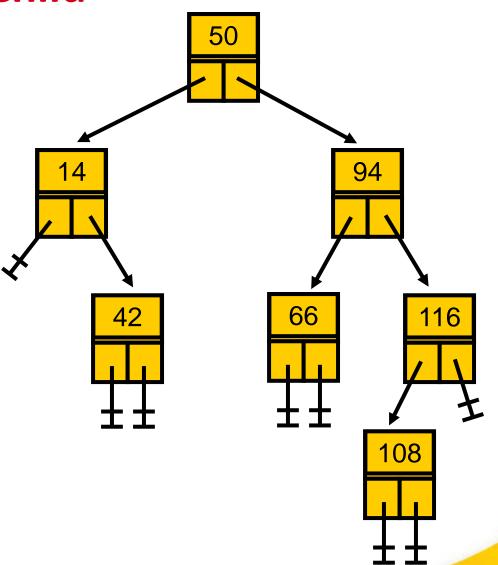


Determine if it has a left or a right child.

Point the current pointer to the appropriate child:

cur <- cur^.left\_child</pre>

The resulting tree.



### **Delete a Node with Two Children**

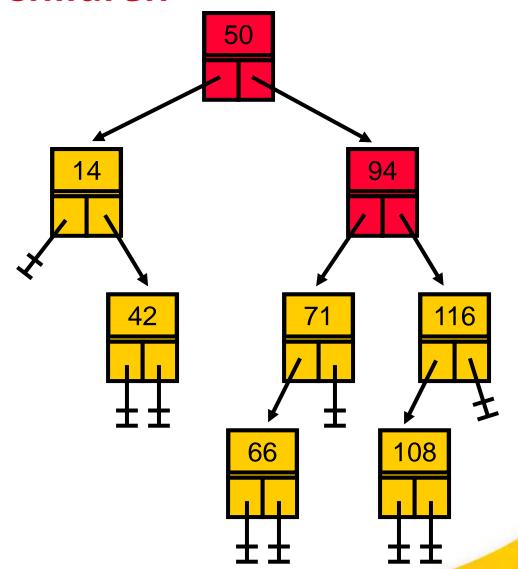


Copy a replacement value from a descendant node.

- Largest from left
- Smallest from right

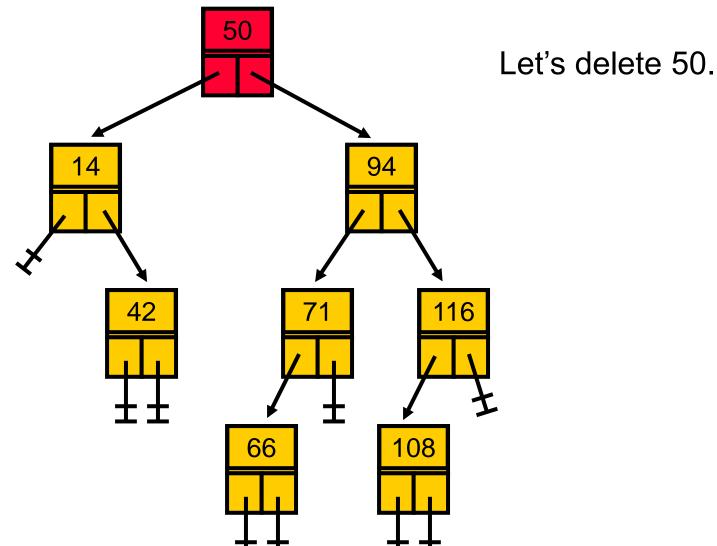
Then delete that descendant node to remove the duplicate value.

- We know this will be an easier case.



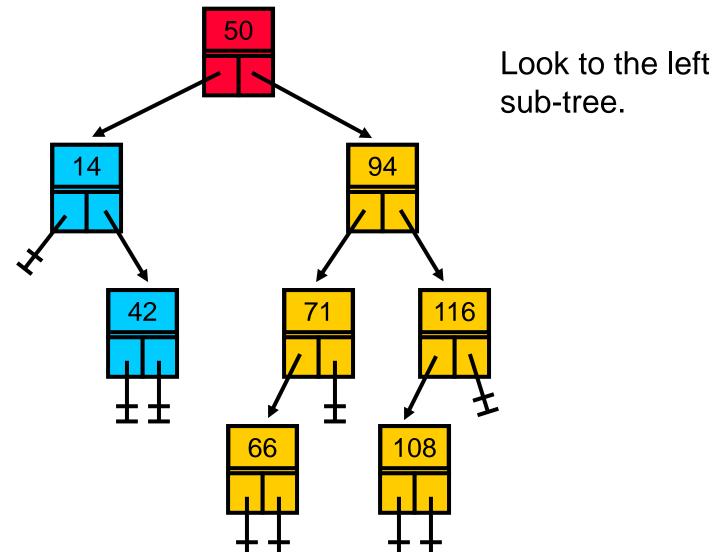






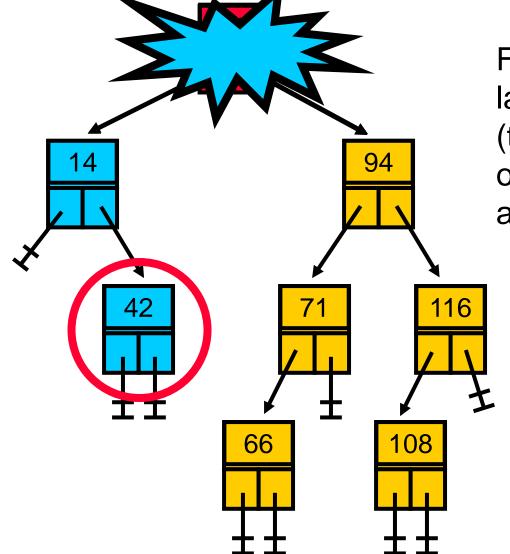








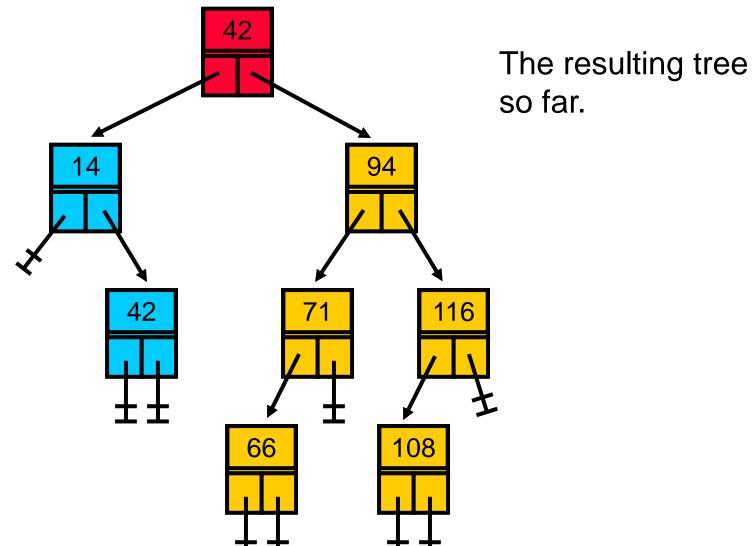
Delete a Node with Two Children



Find and copy the largest value (this will erase the old value but creates a duplicate).

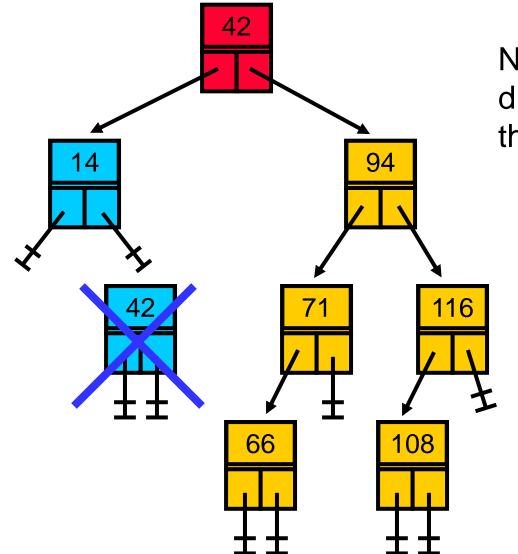








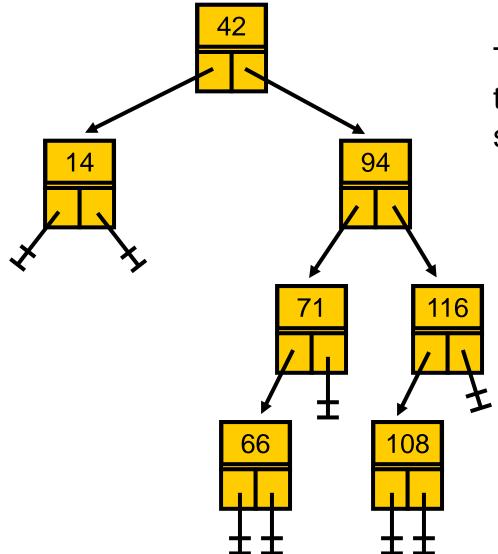




Now delete the duplicate from the left sub-tree.



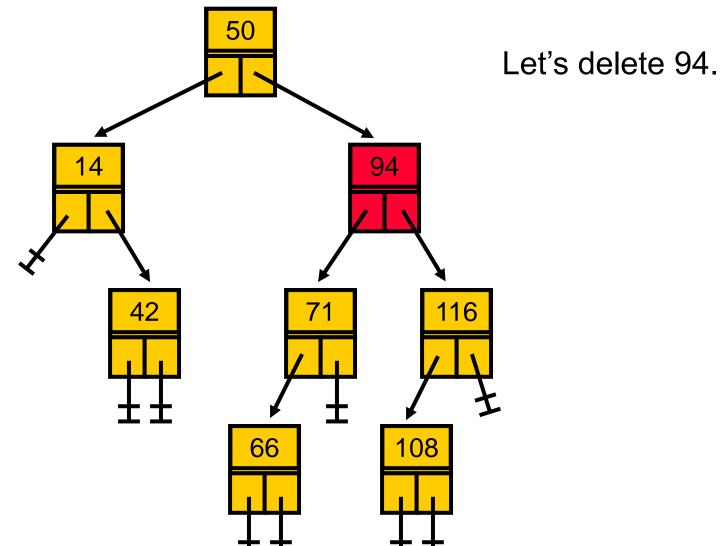




The final resulting tree – still has search structure.

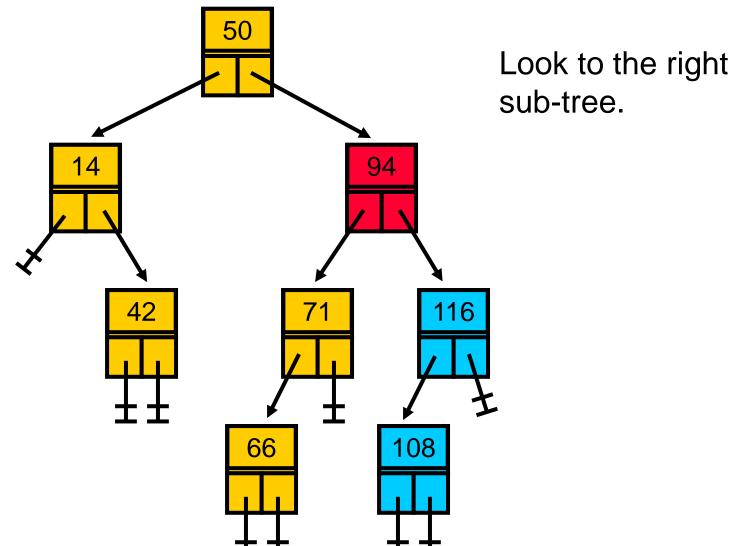






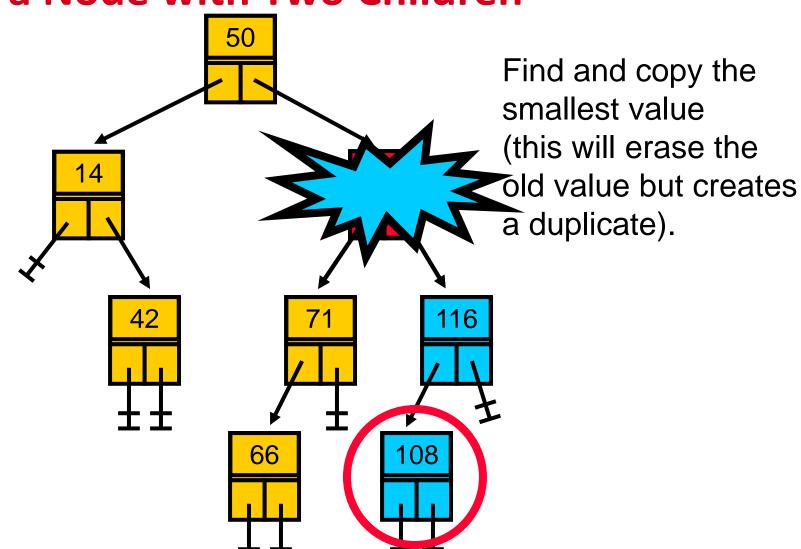








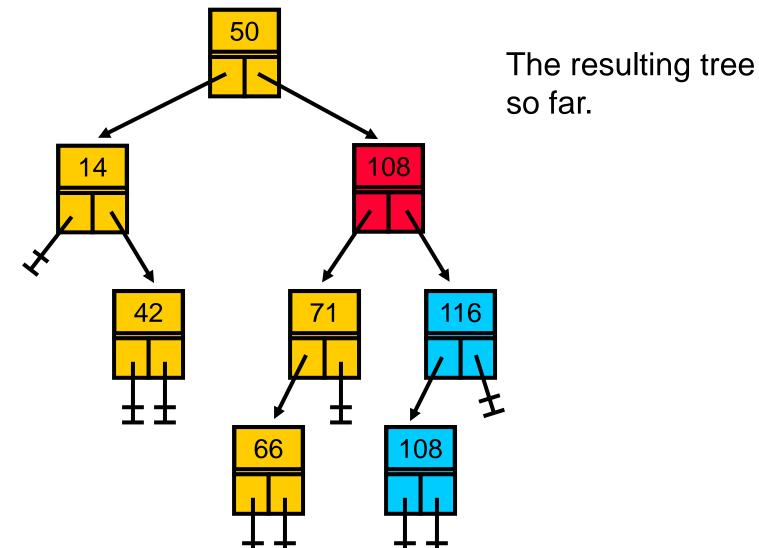
## **Delete a Node with Two Children**





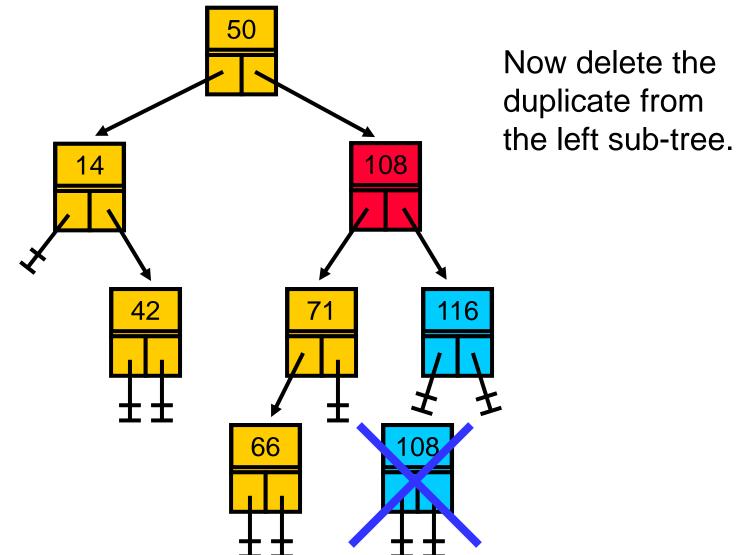


# **Delete a Node with Two Children**





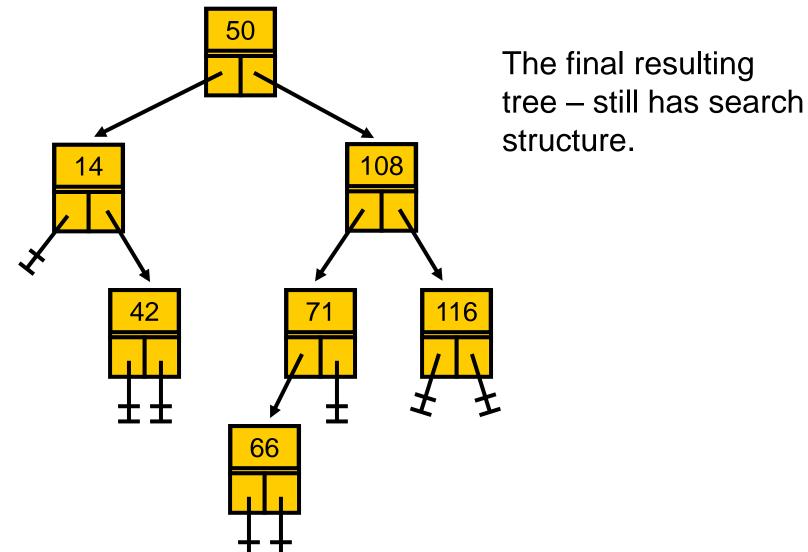




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- Deleting a node from a binary search tree involves two steps:
  - Search for the element
  - Then perform the deletion
- We must preserve the search structure and only delete the element which matches.
- Four cases:
  - Deleting a leaf node
  - Deleting a node with only the left child
  - Deleting a node with only the right child
  - Deleting a node with both children