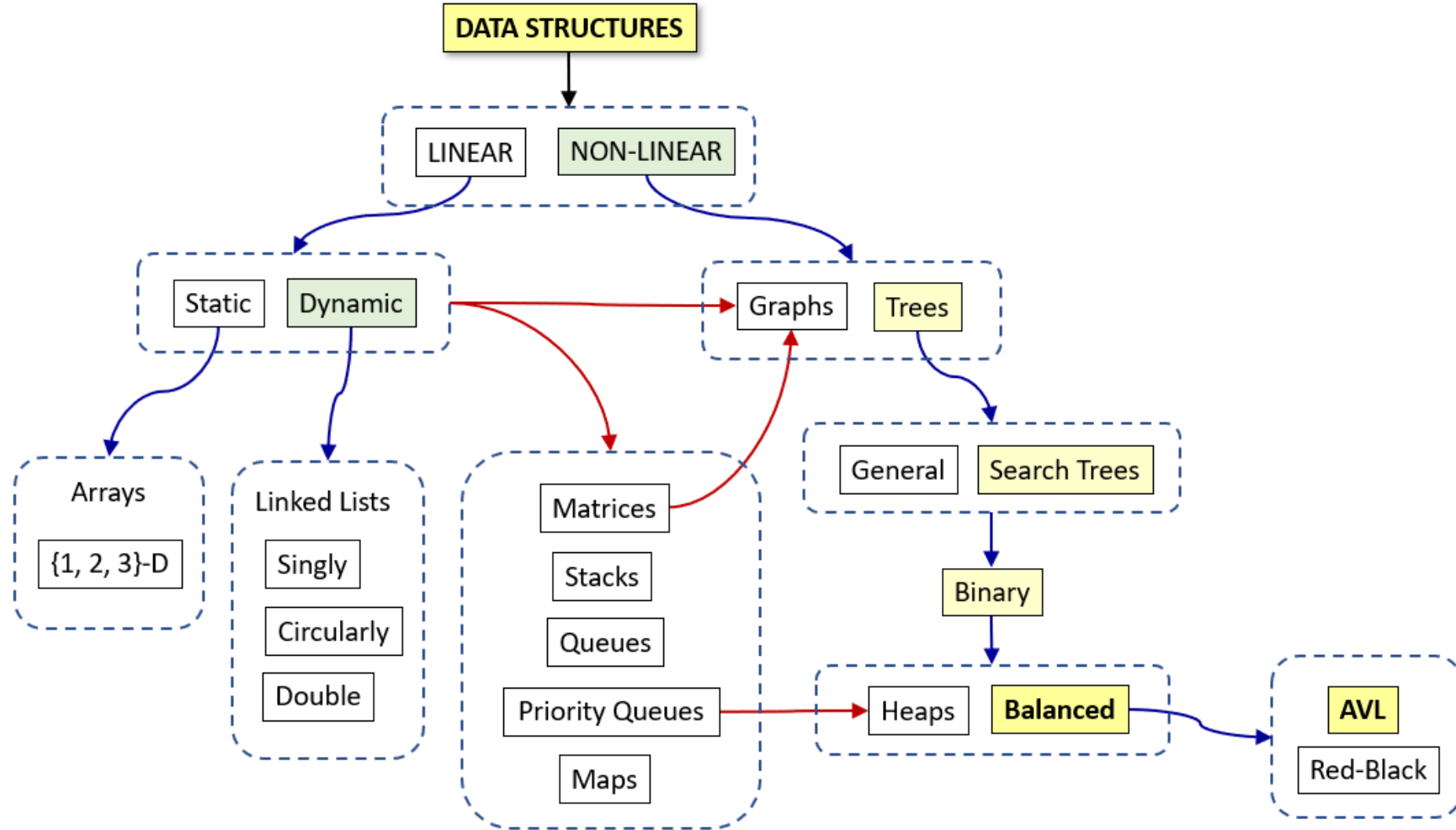


AVL Trees

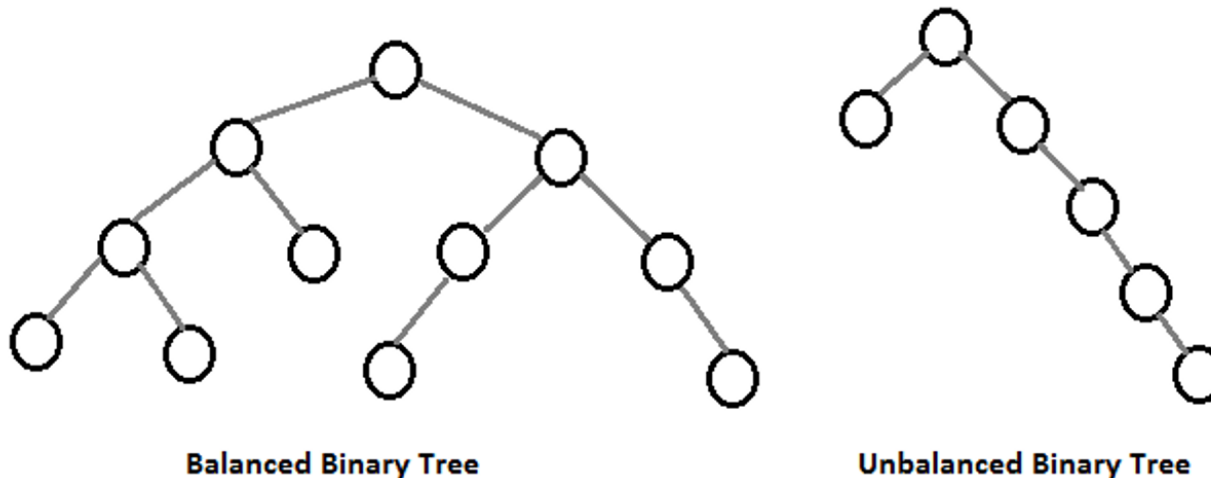


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AVL trees

- Binary search trees
 - Great
 - Search, insert, delete = $O(\log n)$ time complexity
 - **UNLESS they become unbalanced then**
 - **Search, insert, delete = $O(n)$ time complexity**



AVL trees

- Binary search trees
 - Great
 - Search, insert, delete = $O(\log n)$ time complexity
 - UNLESS they become unbalanced then
 - Search, insert, delete = $O(n)$ time complexity
- AVL trees are binary search trees with special insertion and deletion algorithms to keep them well-balanced
- so the time complexity of search insert and delete remain $O(\log n)$

Height of a Tree

- Height of a tree is the length of the longest path from root to some leaf node.
- Height of an empty tree is -1.
- Height of a single node tree is 0.
- Recursive definition:
 - $\text{height}(t) = 0$ if number of nodes = 1
 - $= -1$ if T is empty
 - $= 1 + \max(\text{height}(\text{LT}), \text{height}(\text{RT}))$ otherwise

AVL Property

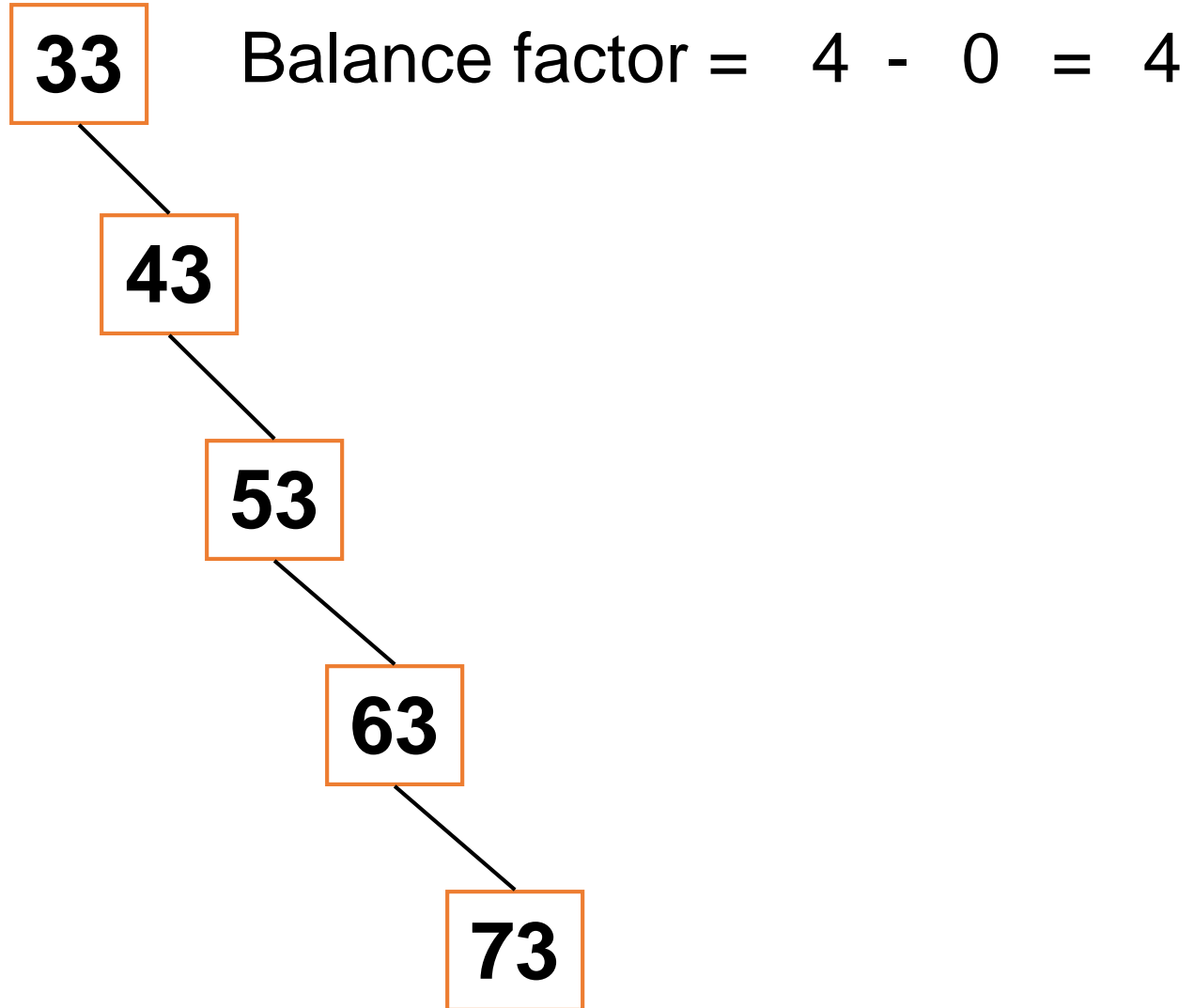
- If N is a node in a binary tree,

node X has AVL property if

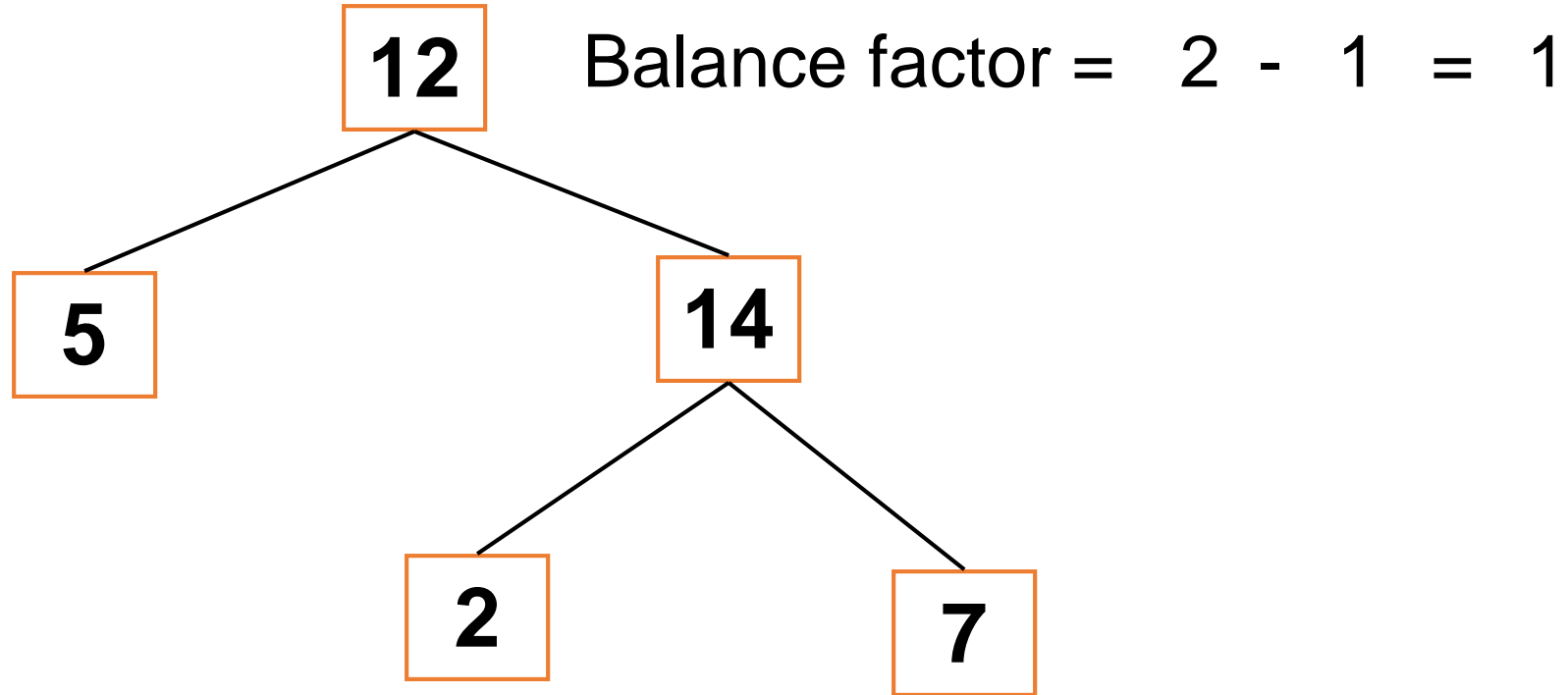
the heights of the left sub-tree and right sub-tree are equal or if they differ by 1.

- Let's look at some examples.

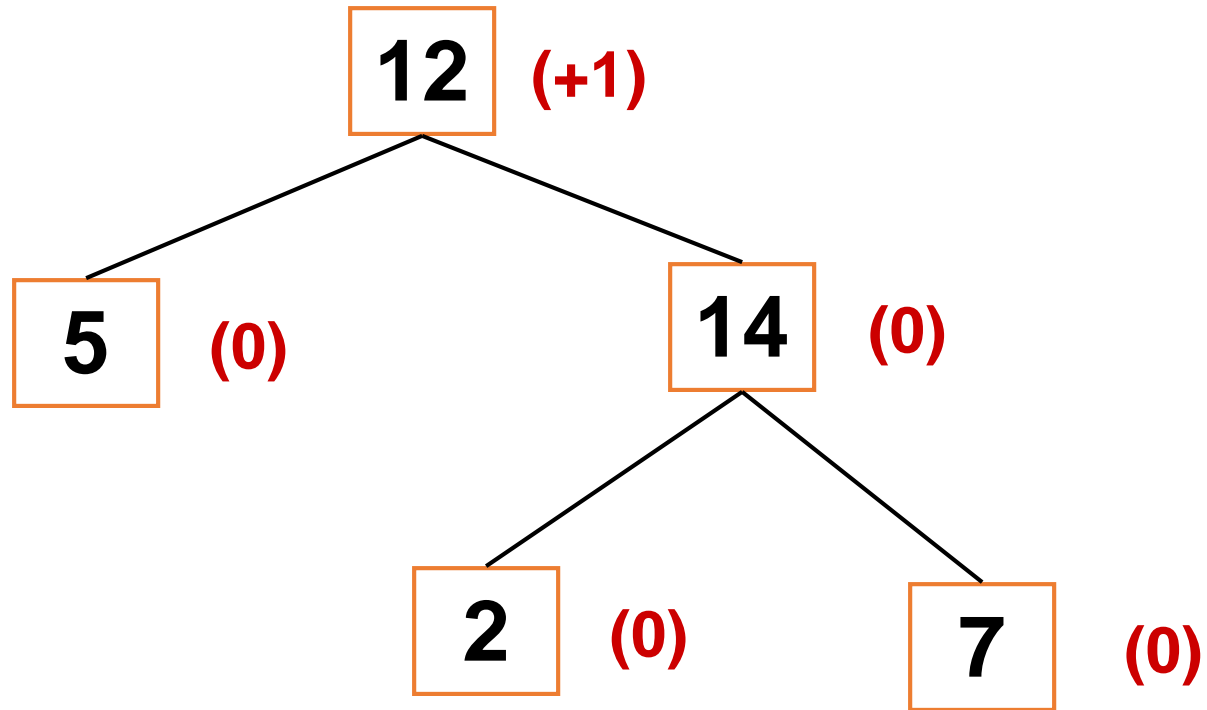
Degenerate BST



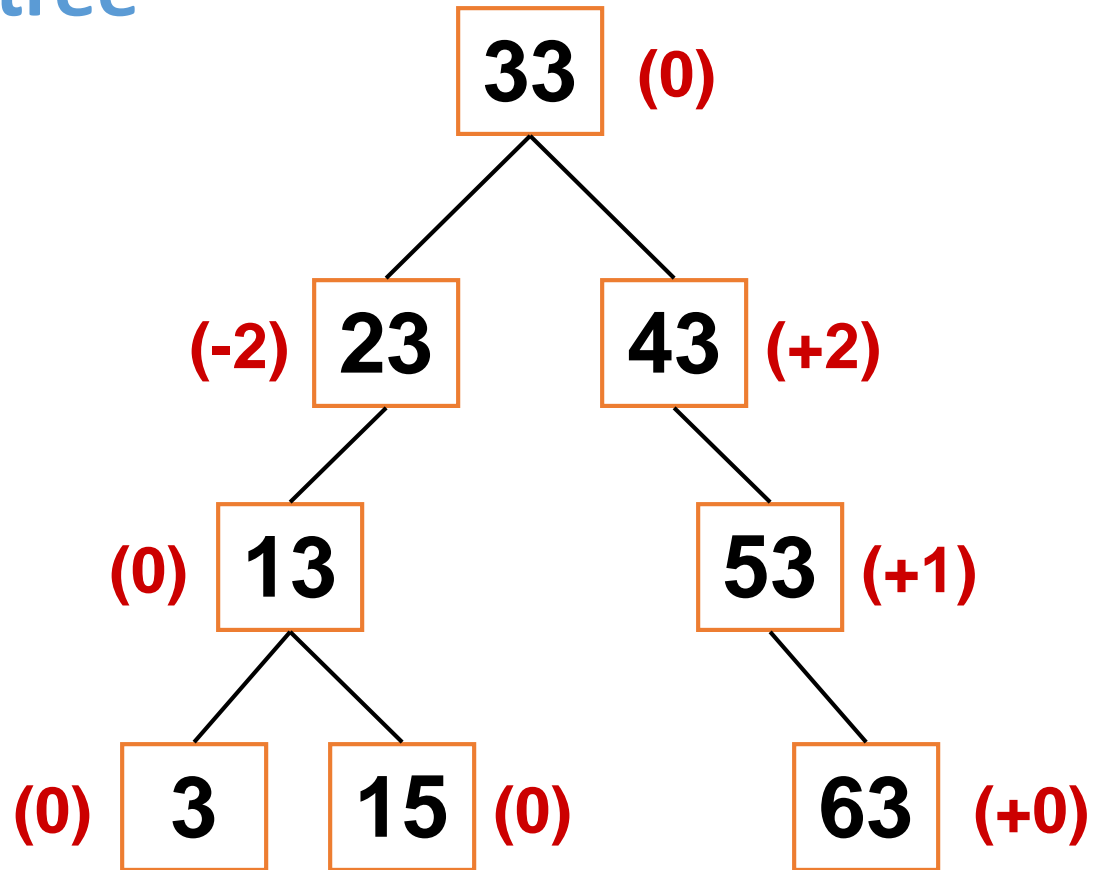
Degenerate BST



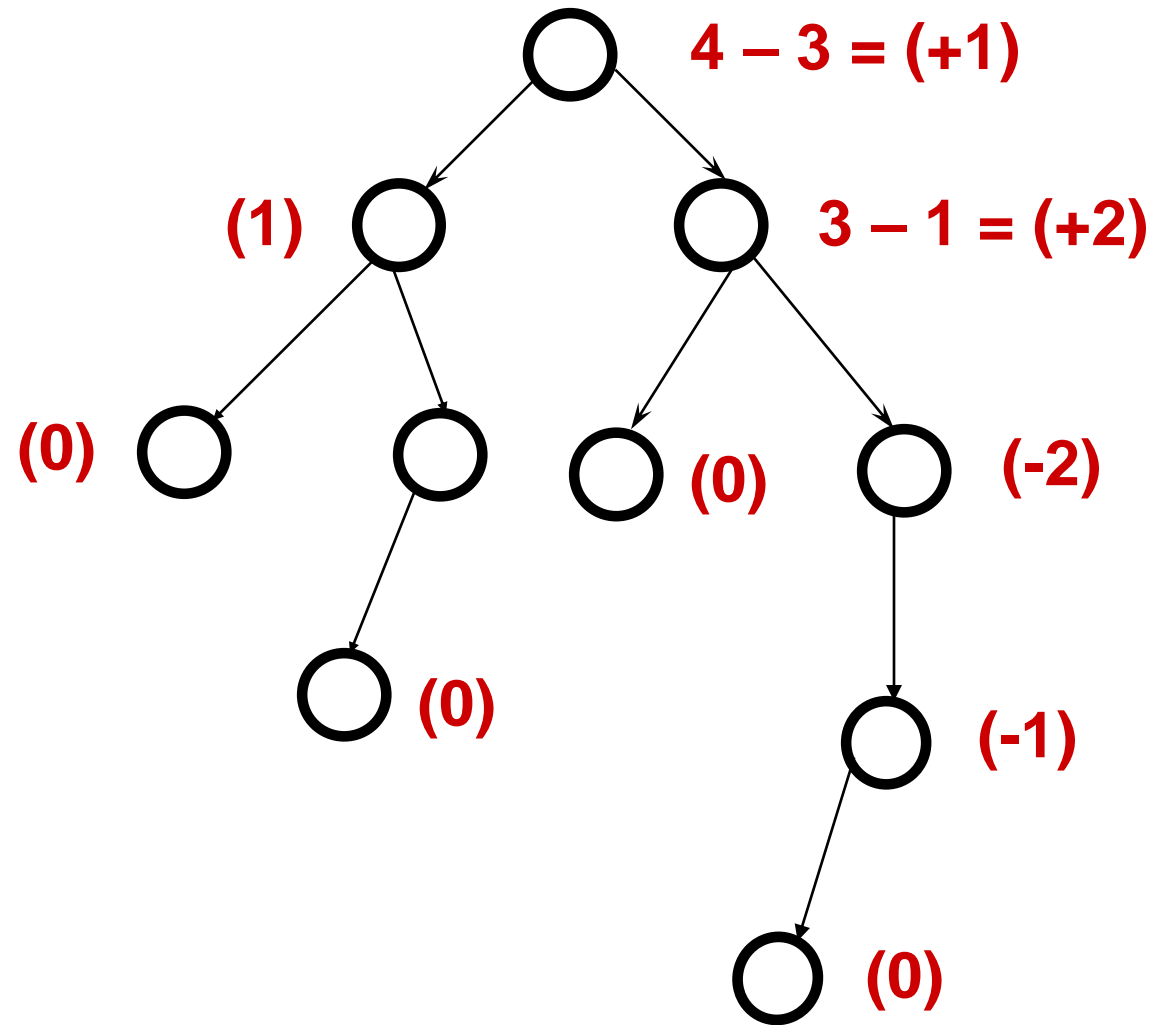
AVL tree



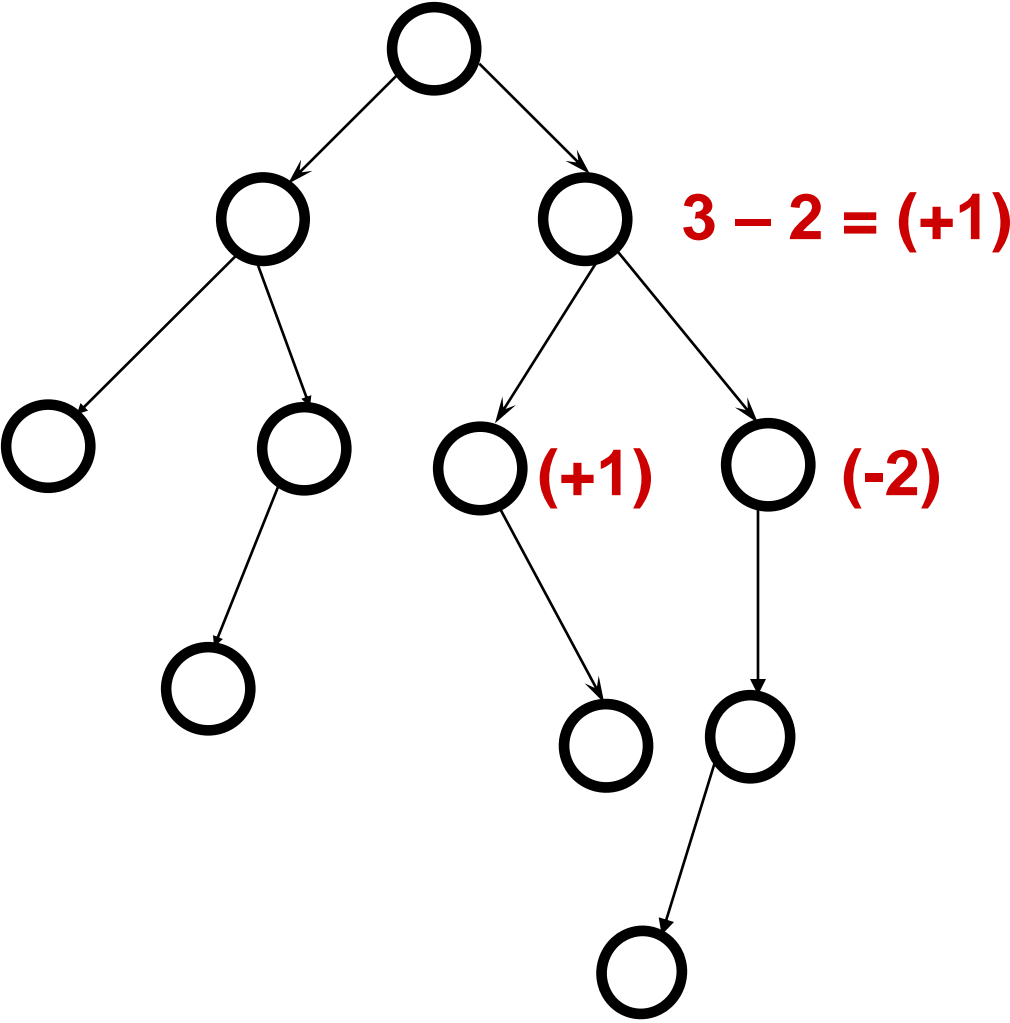
Non-AVL tree



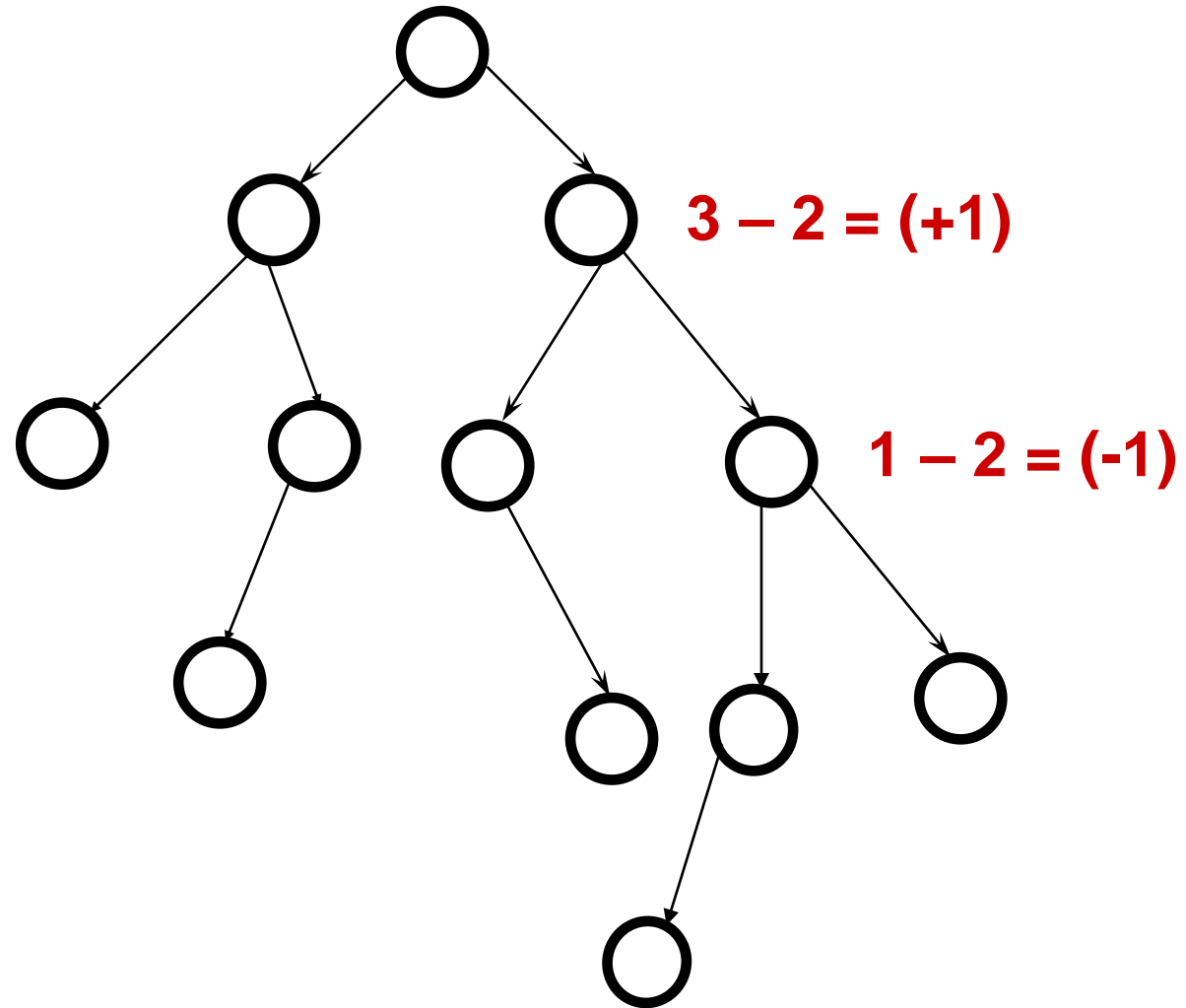
Non-AVL tree



Non-AVL tree



AVL tree

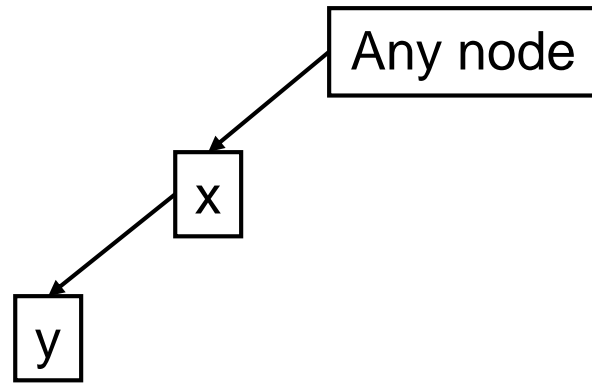


AVL trees

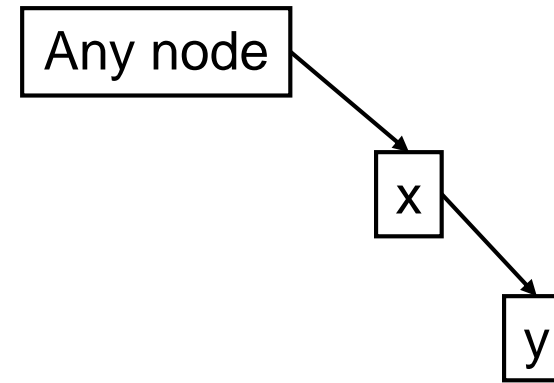
- AVL tree = a binary search tree with the following property: for every node the heights of the left and right subtrees differ at most by one.
- That's all very nice but how do we guarantee it?
- We have to somehow modify the insert and delete functions.
- If, after an insertion or deletion, the property is not satisfied, we “**rotate**” the tree to make it balanced.

AVL trees

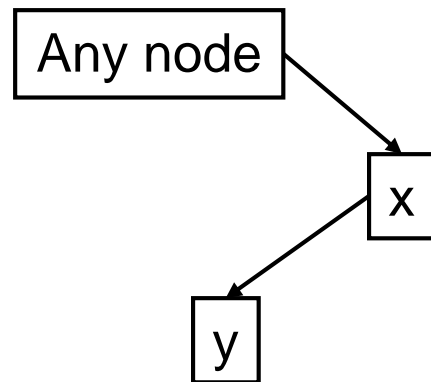
- When can an insertion of a child y at node x cause an imbalance?
 - when both x and y are left children
 - when both x and y are right children
 - when x is a right child and y is a left child
 - when y is a right child and x is a left child



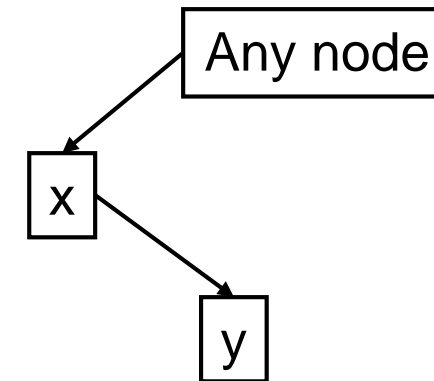
when both x and y are left children



when both x and y are right children



when x is a right child and y is a left child



when y is a right child and x is a left child

AVL trees

- AVL tree = a binary search tree with the following property: for every node the heights of the left and right subtrees differ at most by one.
 - That's all very nice but how do we guarantee it?
- We have to somehow modify the insert and delete functions.
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AVL trees

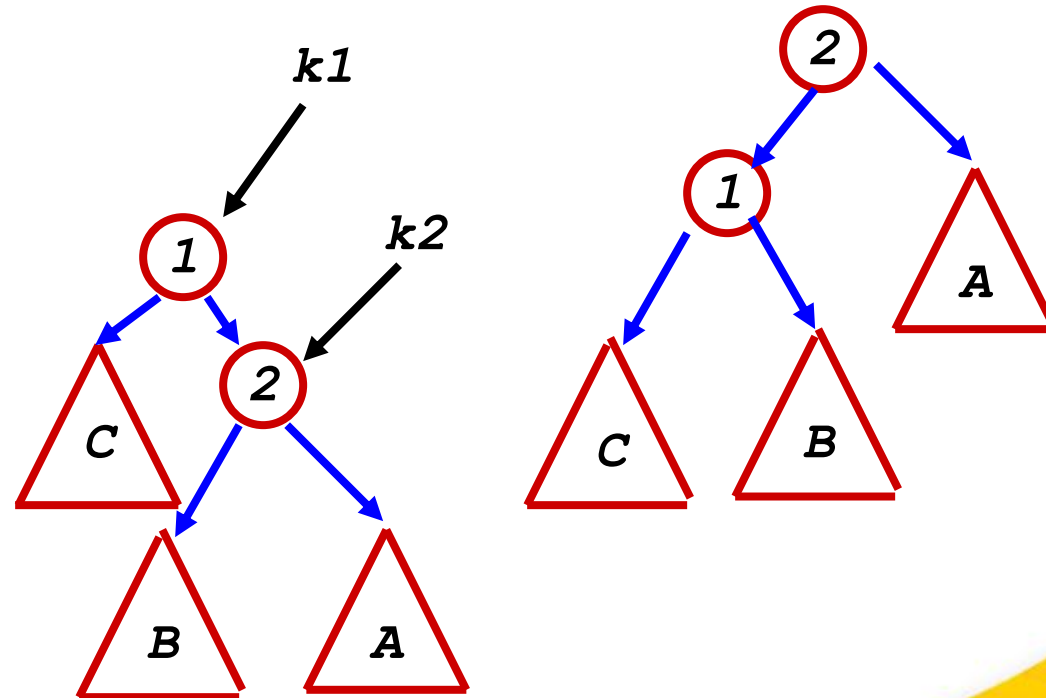
- How/when do we decide whether to rotate?
- Example: insertion
 - step 1: walk down the tree to insert the node in the correct position
 - step 2: walk up the tree checking the property at each node
- we need a helper function to determine the heights of the subtrees of each node.
- we need to be able to determine whether to perform a **single** or a **double** rotation

AVL trees

- Insertion
- may not unbalance the tree
- If it does then two possible responses
 - **Single rotation**
 - **Double rotation**

Single Rotation Left

```
void LeftRotate( BinaryNode * & k1 )  
{  
    BinaryNode *k2 = k1->right;  
    k1->right = k2->left;  
    k2->left = k1;  
    k1 = k2;  
}
```



```
void LeftRotate( BinaryNode * & k1 )
```

```
{
```

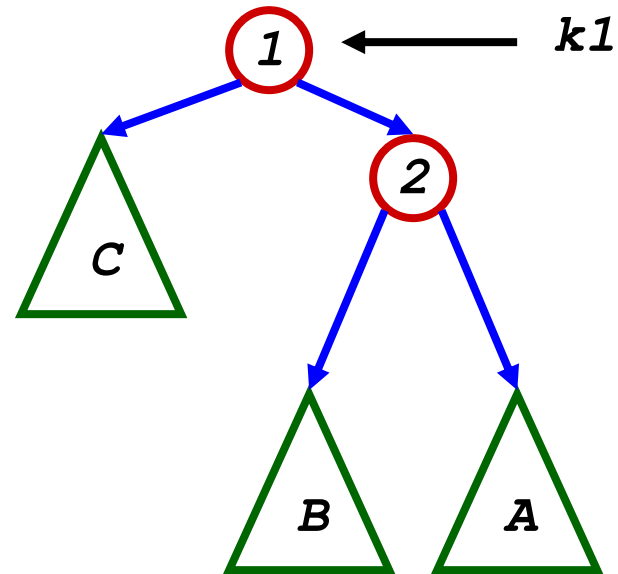
```
    BinaryNode *k2 = k1->right;
```

```
    k1->right = k2->left;
```

```
    k2->left = k1;
```

```
    k1 = k2;
```

```
}
```



```
void LeftRotate( BinaryNode * & k1 )
```

```
{
```

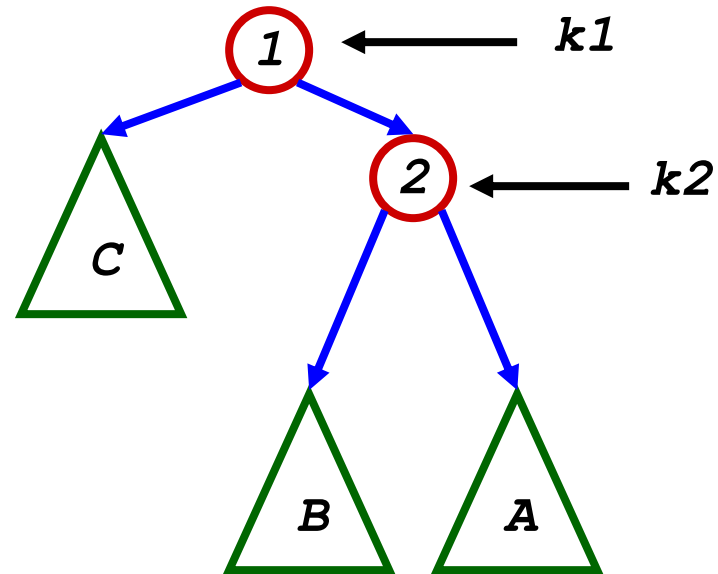
```
    BinaryNode *k2 = k1->right;
```

```
    k1->right = k2->left;
```

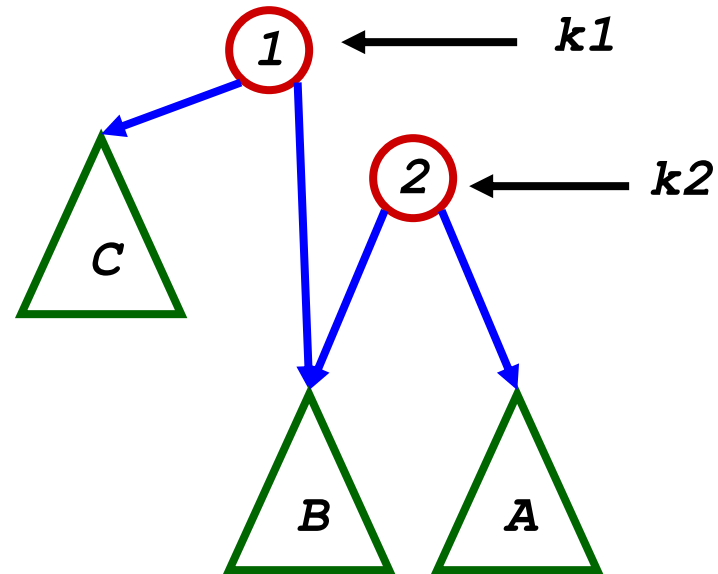
```
    k2->left = k1;
```

```
    k1 = k2;
```

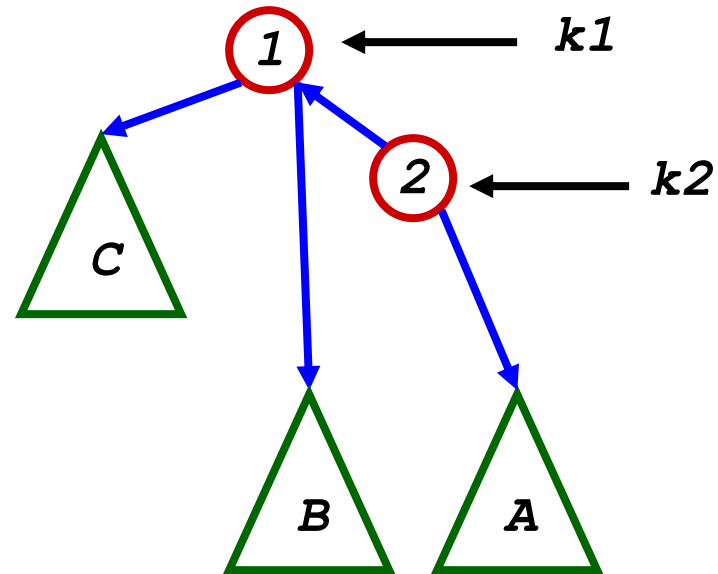
```
}
```



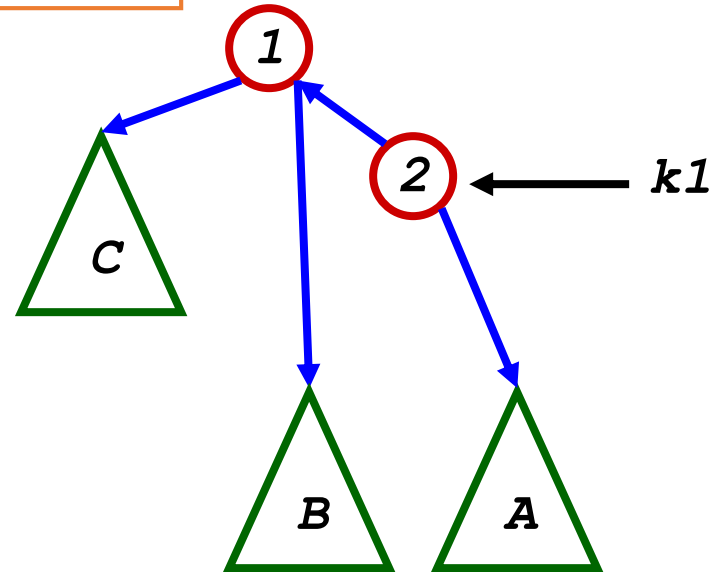
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```



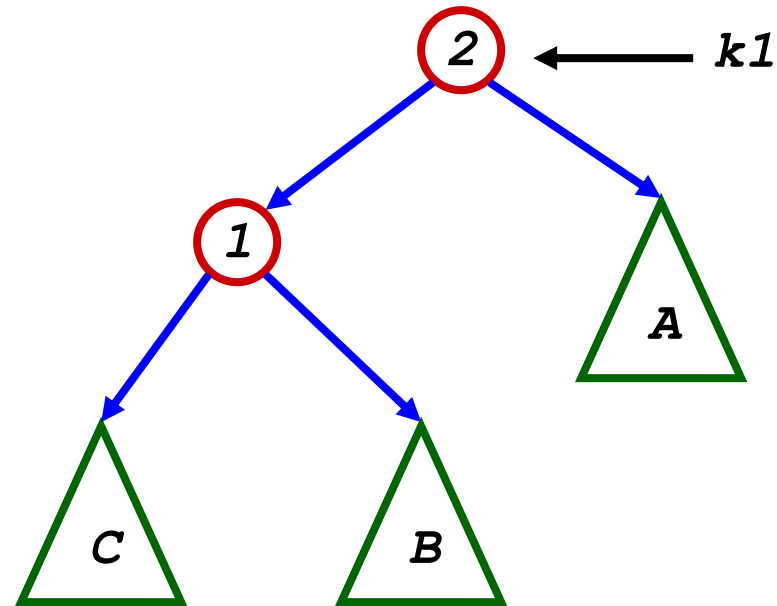
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}
```



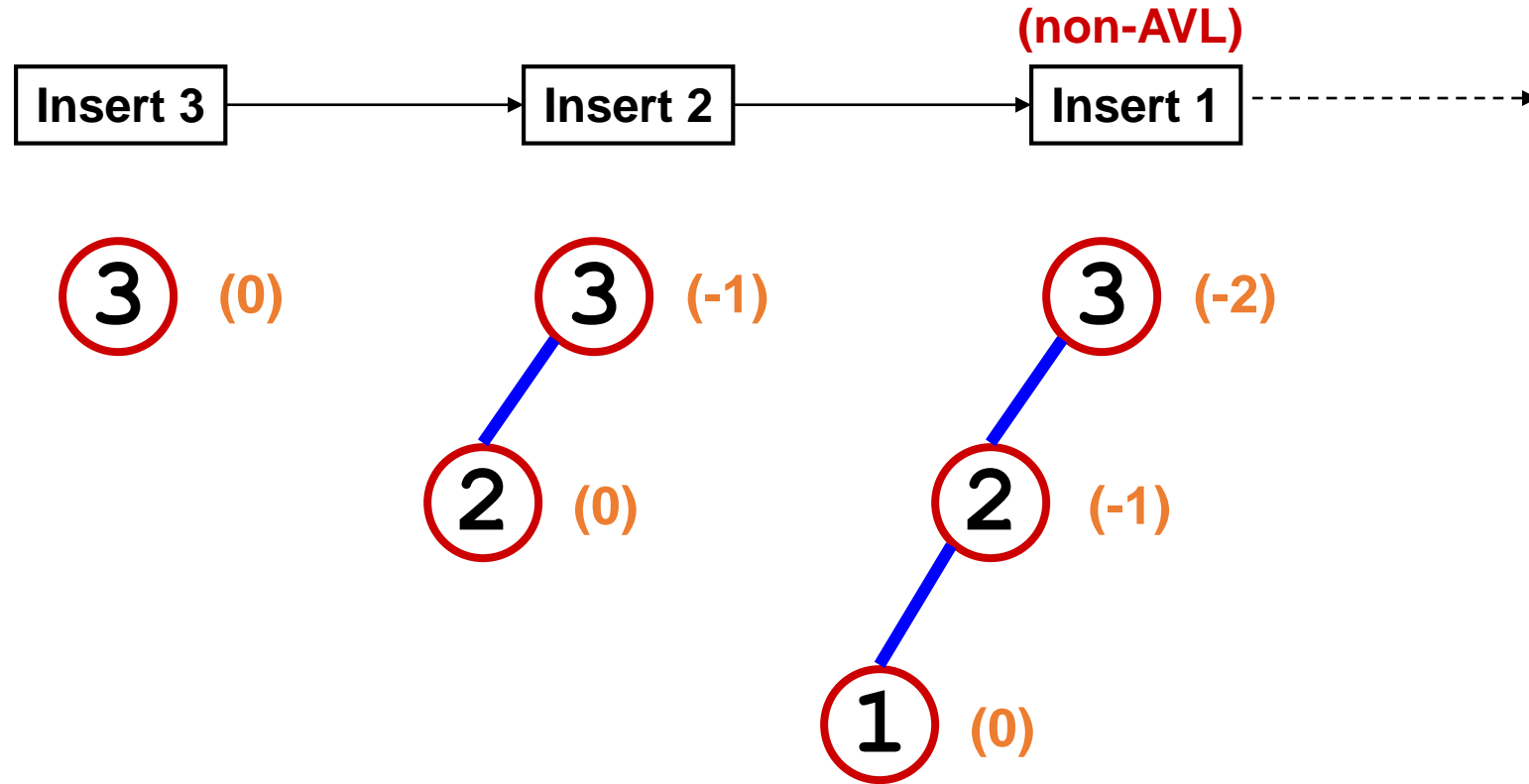

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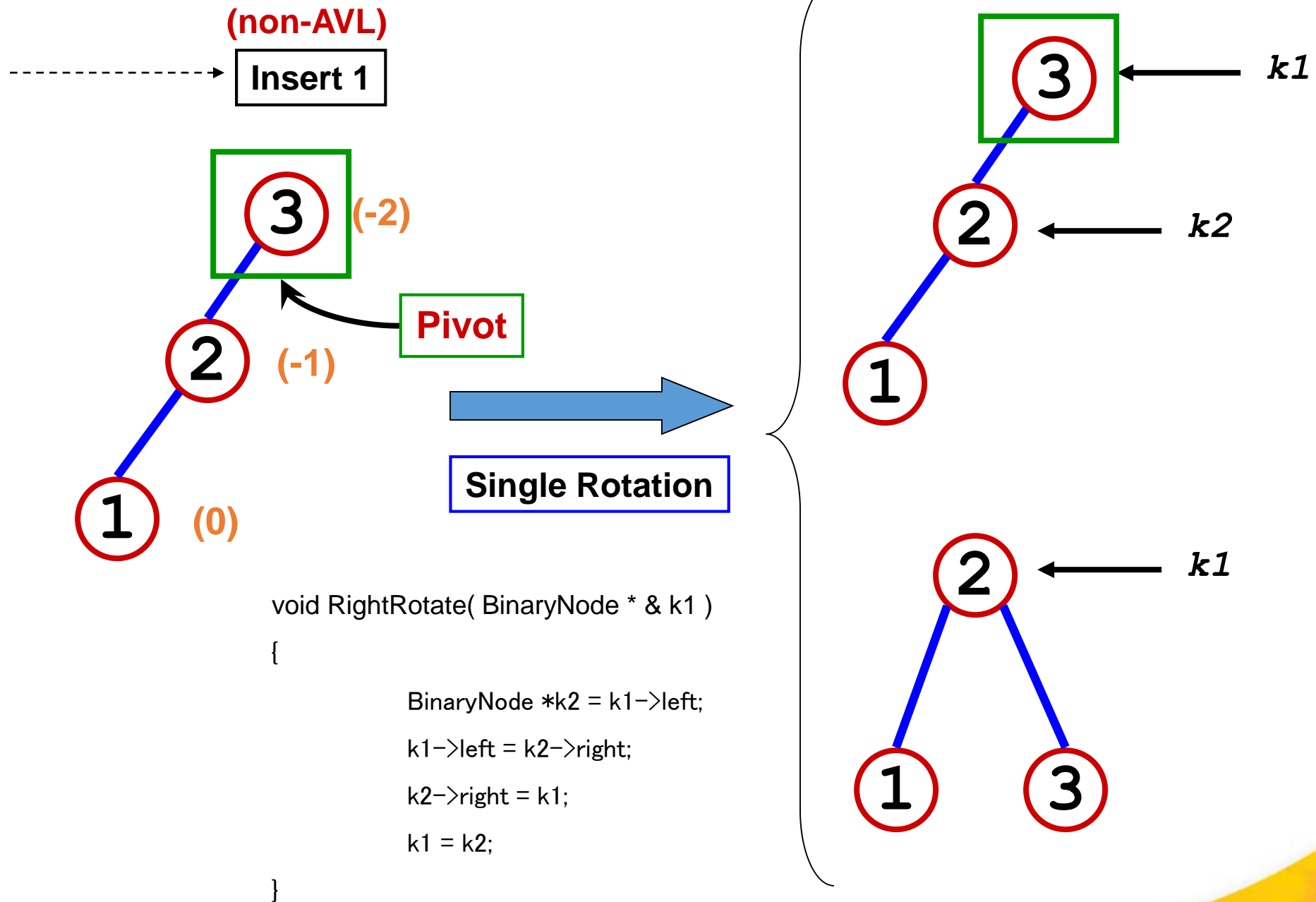


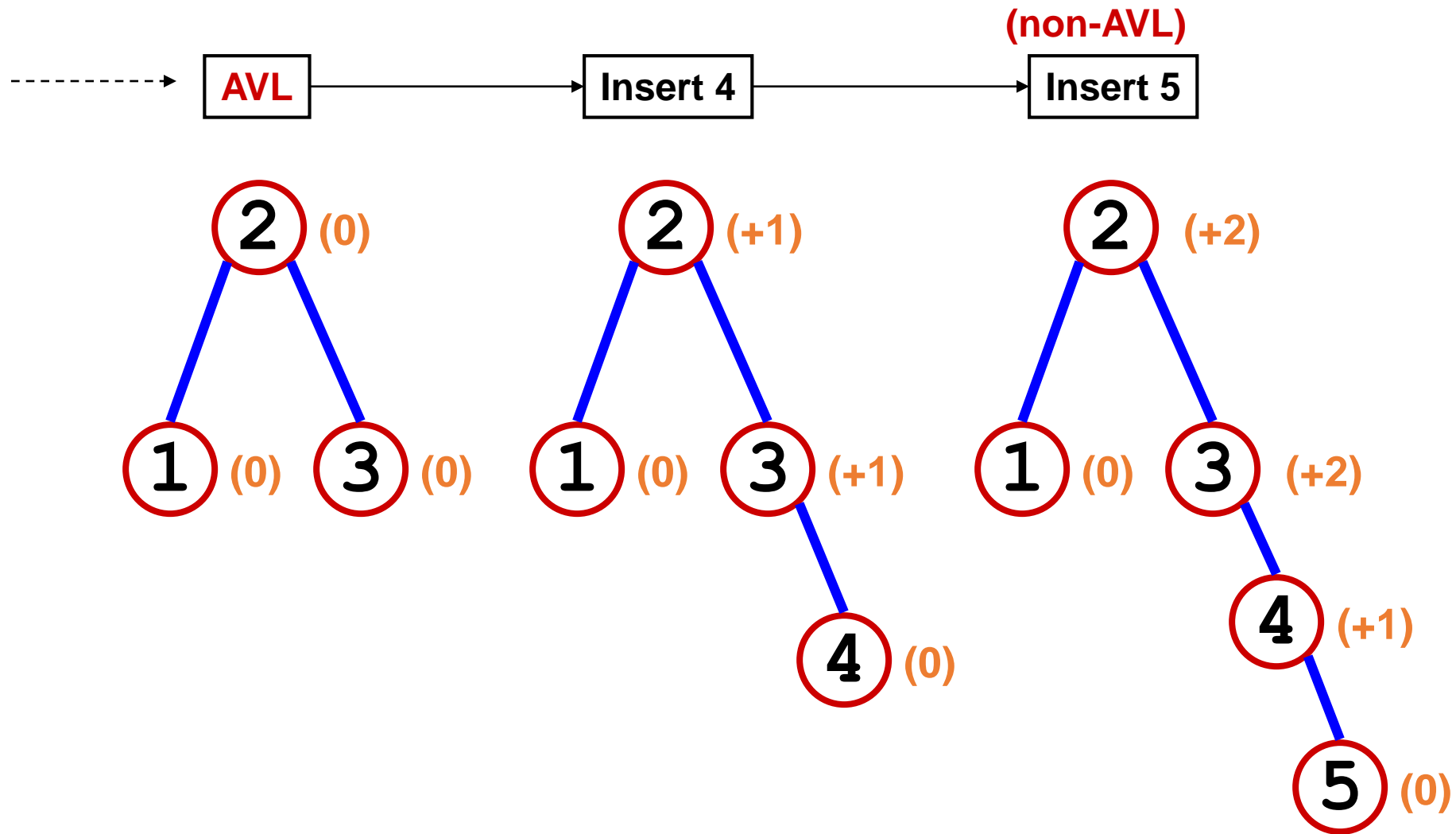
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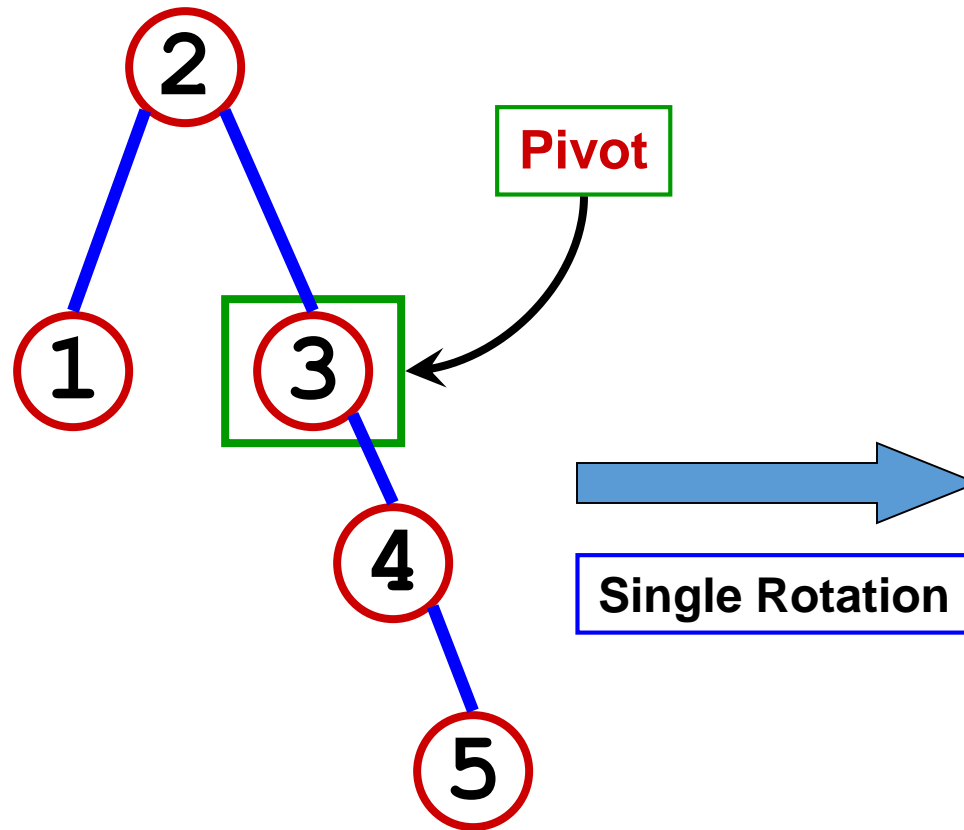


AVL Trees Example

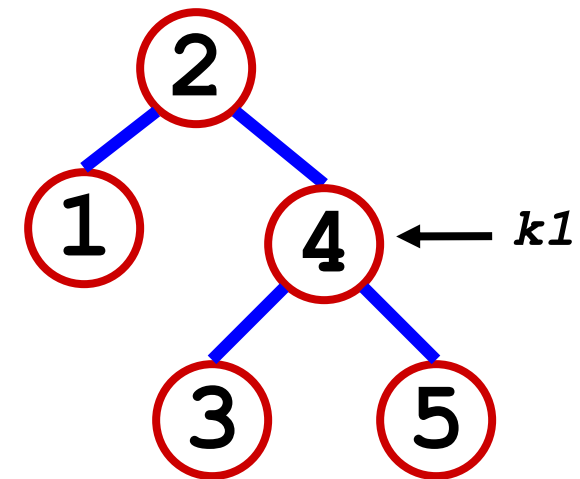
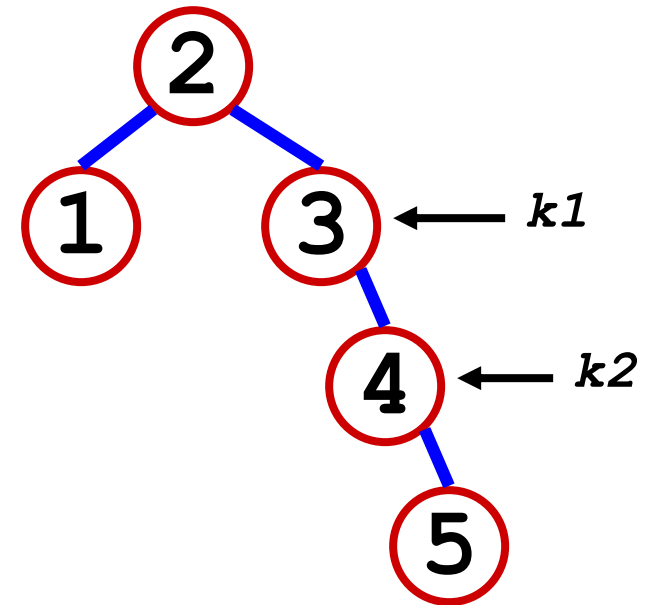


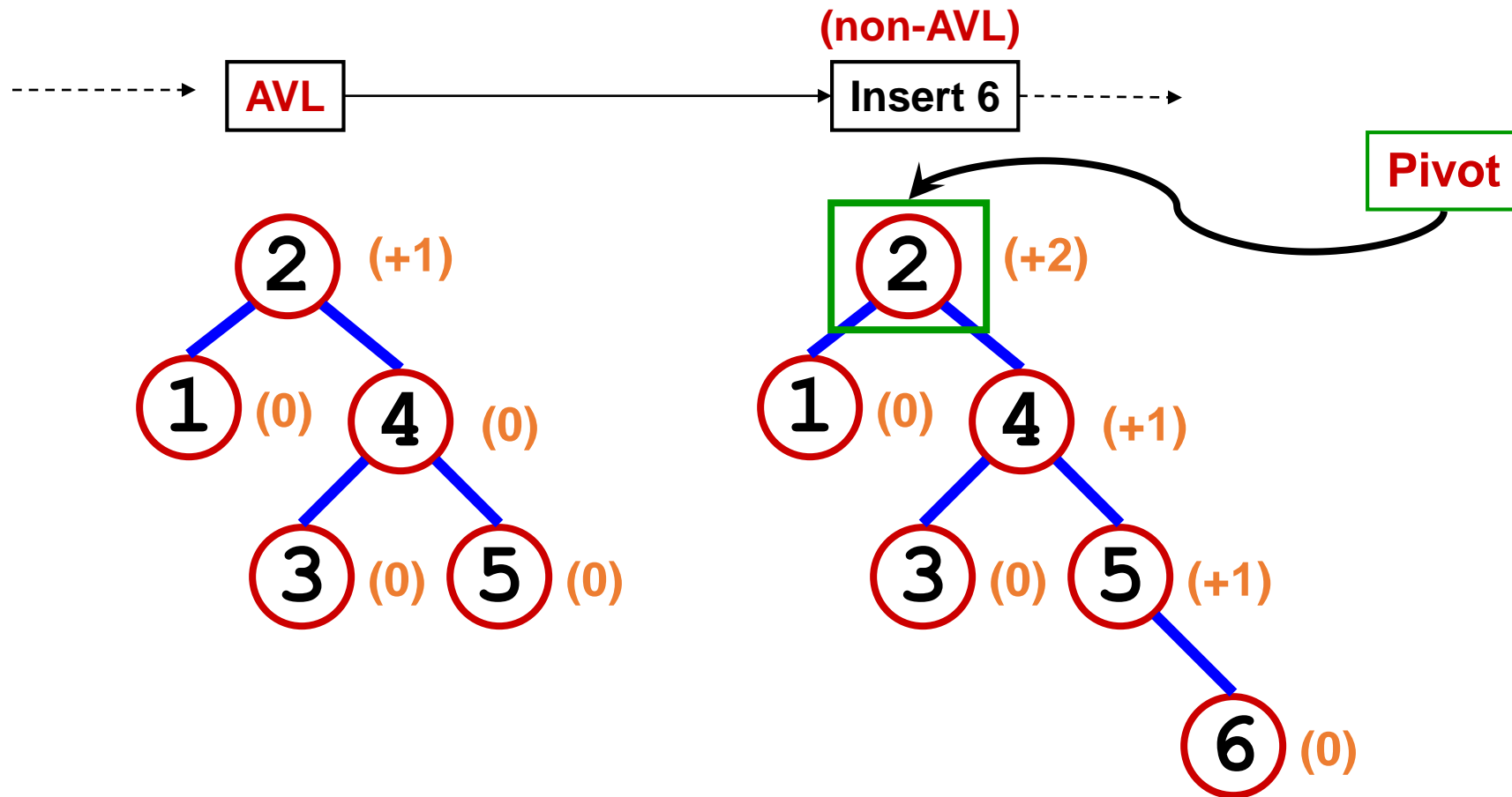


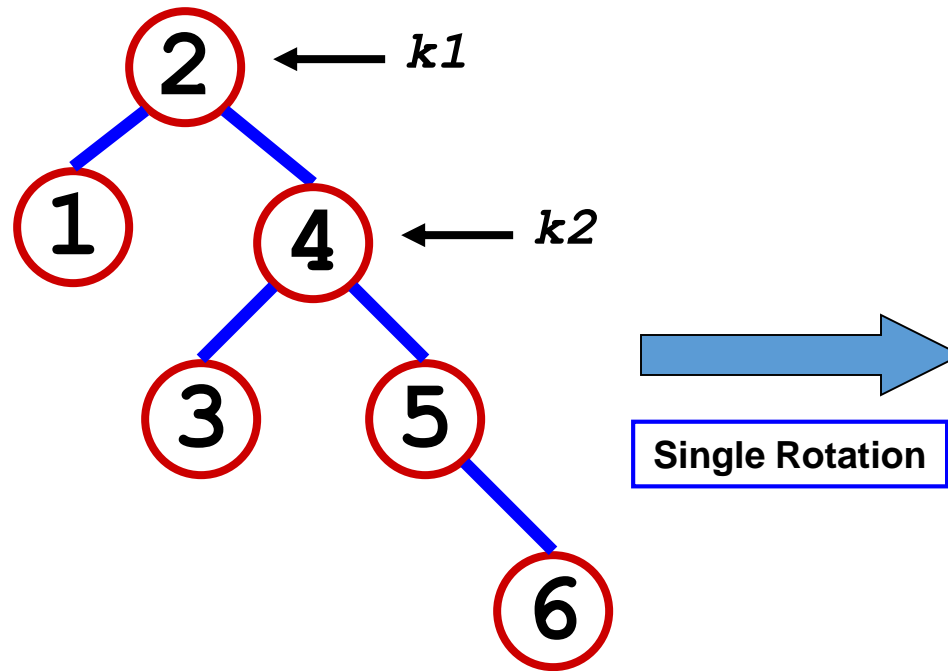




```
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    k1 = k2;
}
```

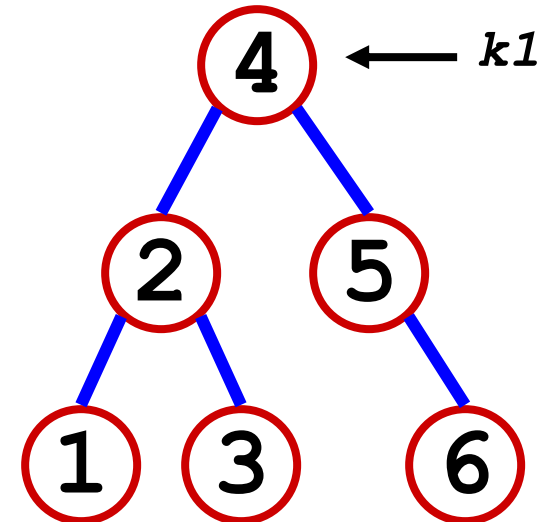
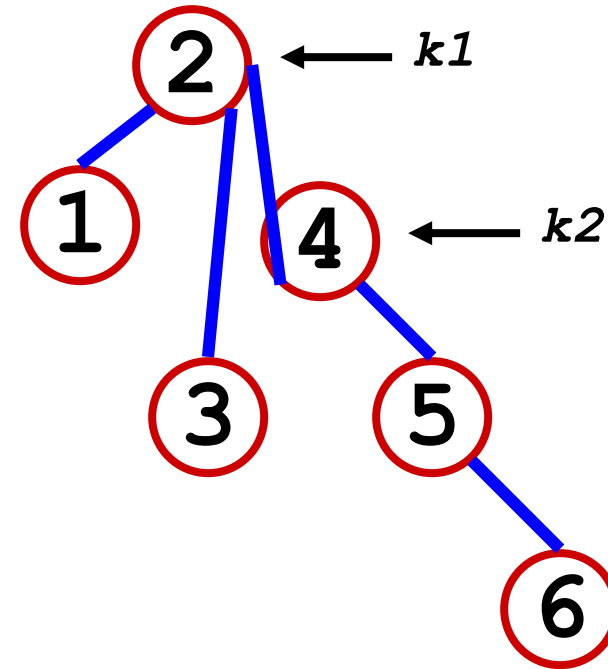


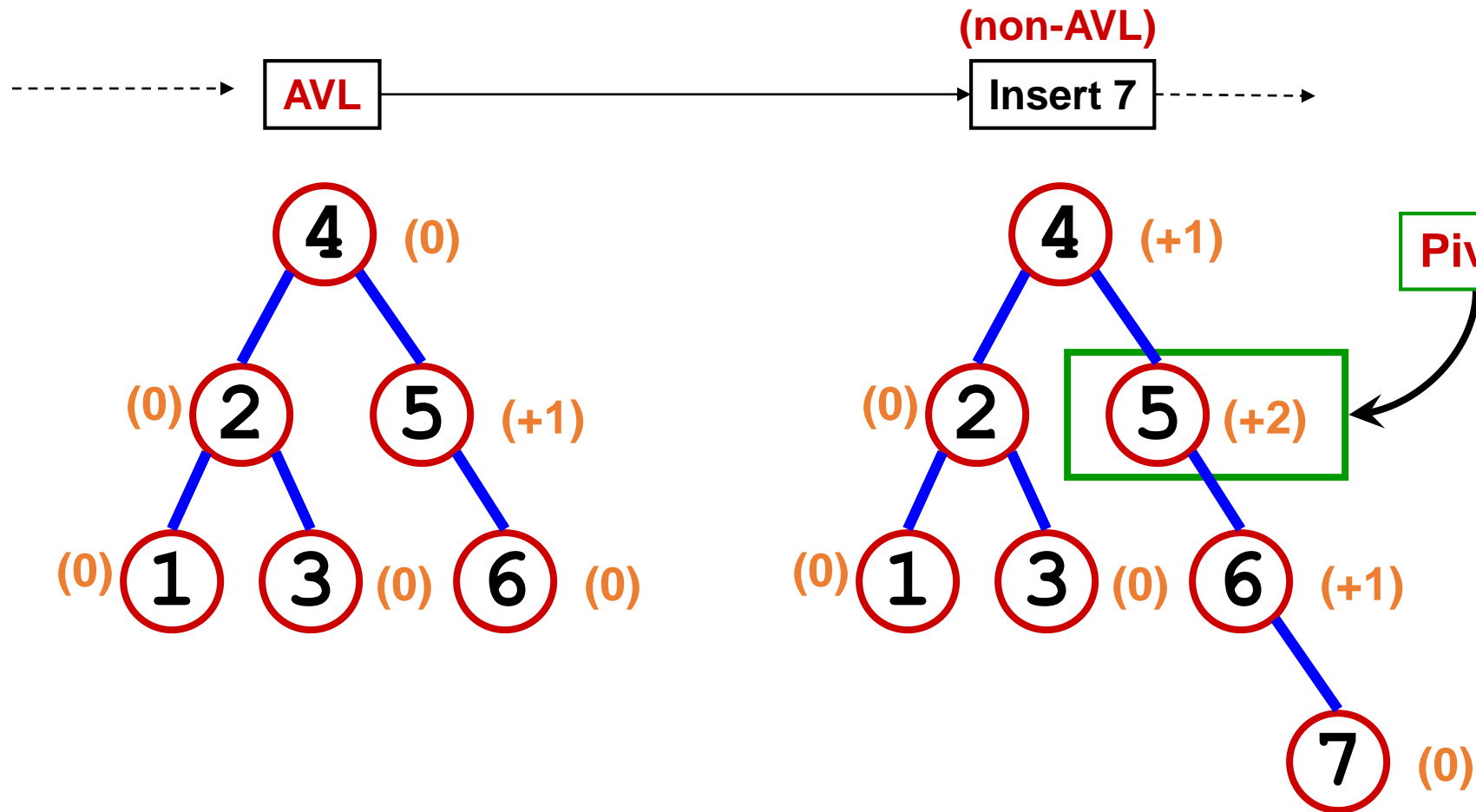


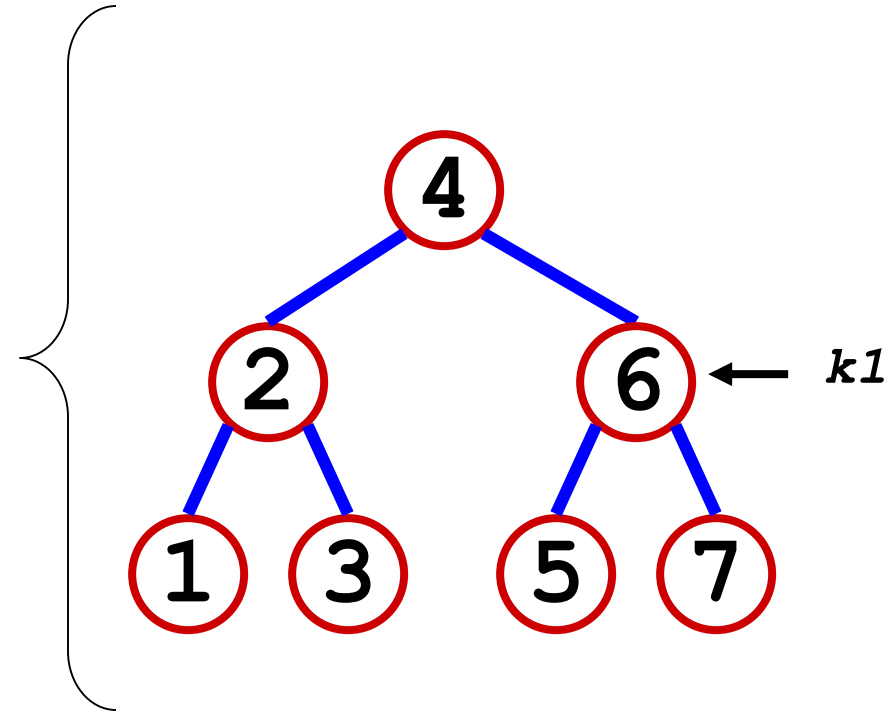
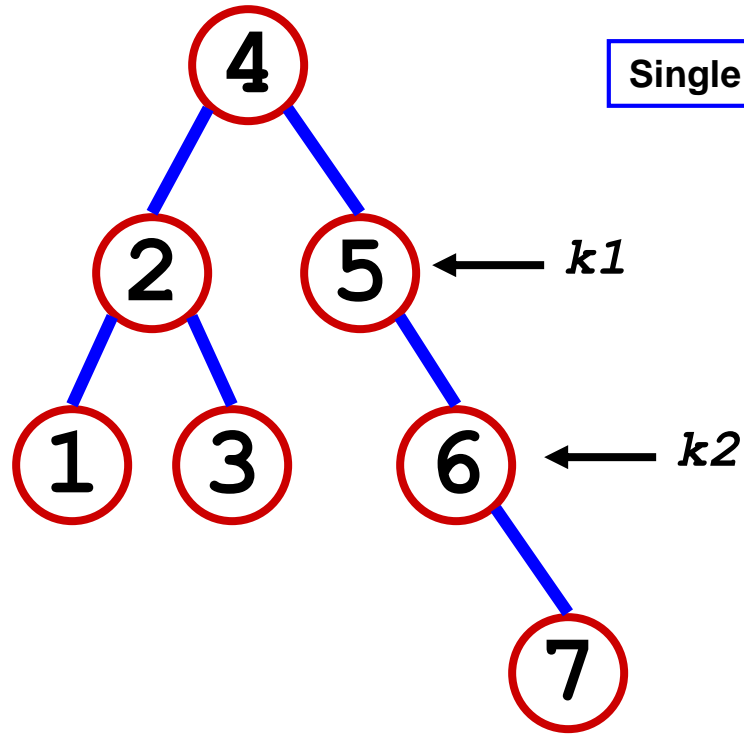


```

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    k1 = k2;
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```

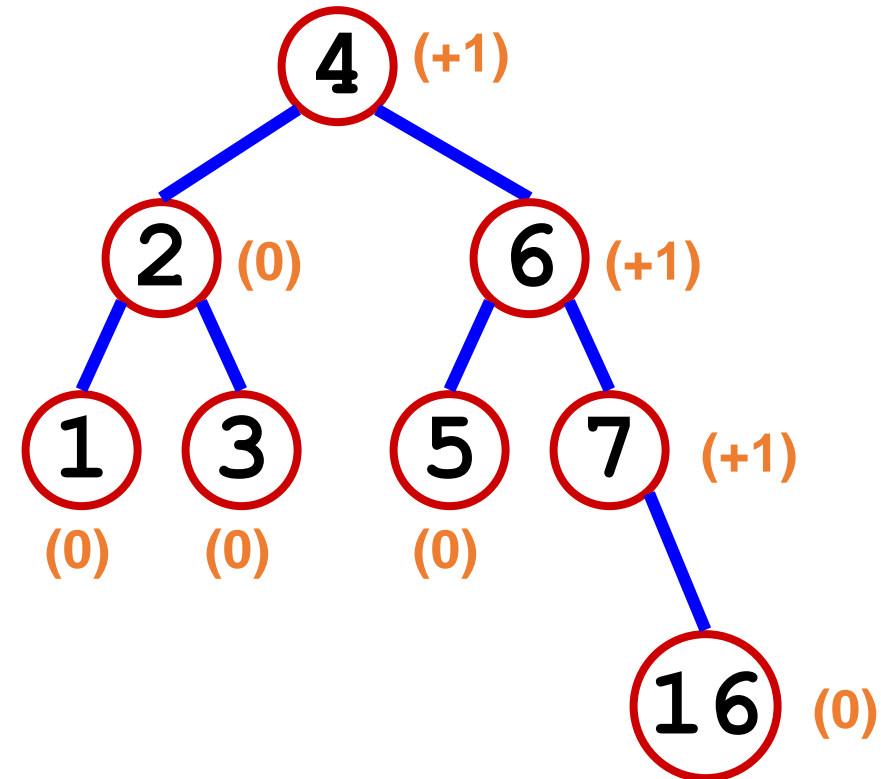
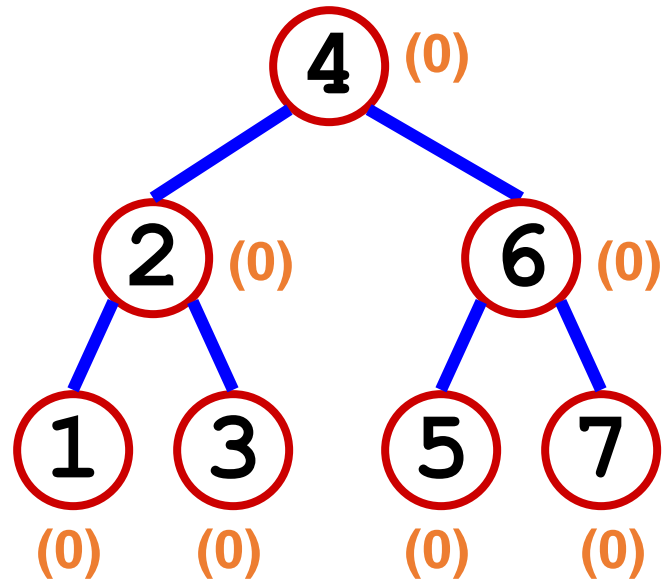
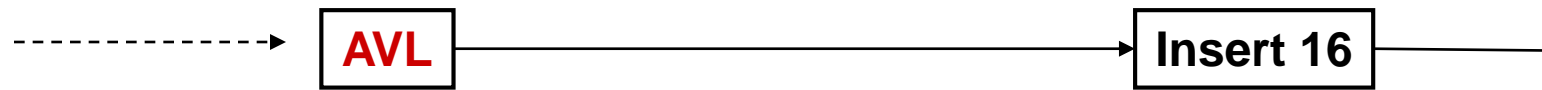




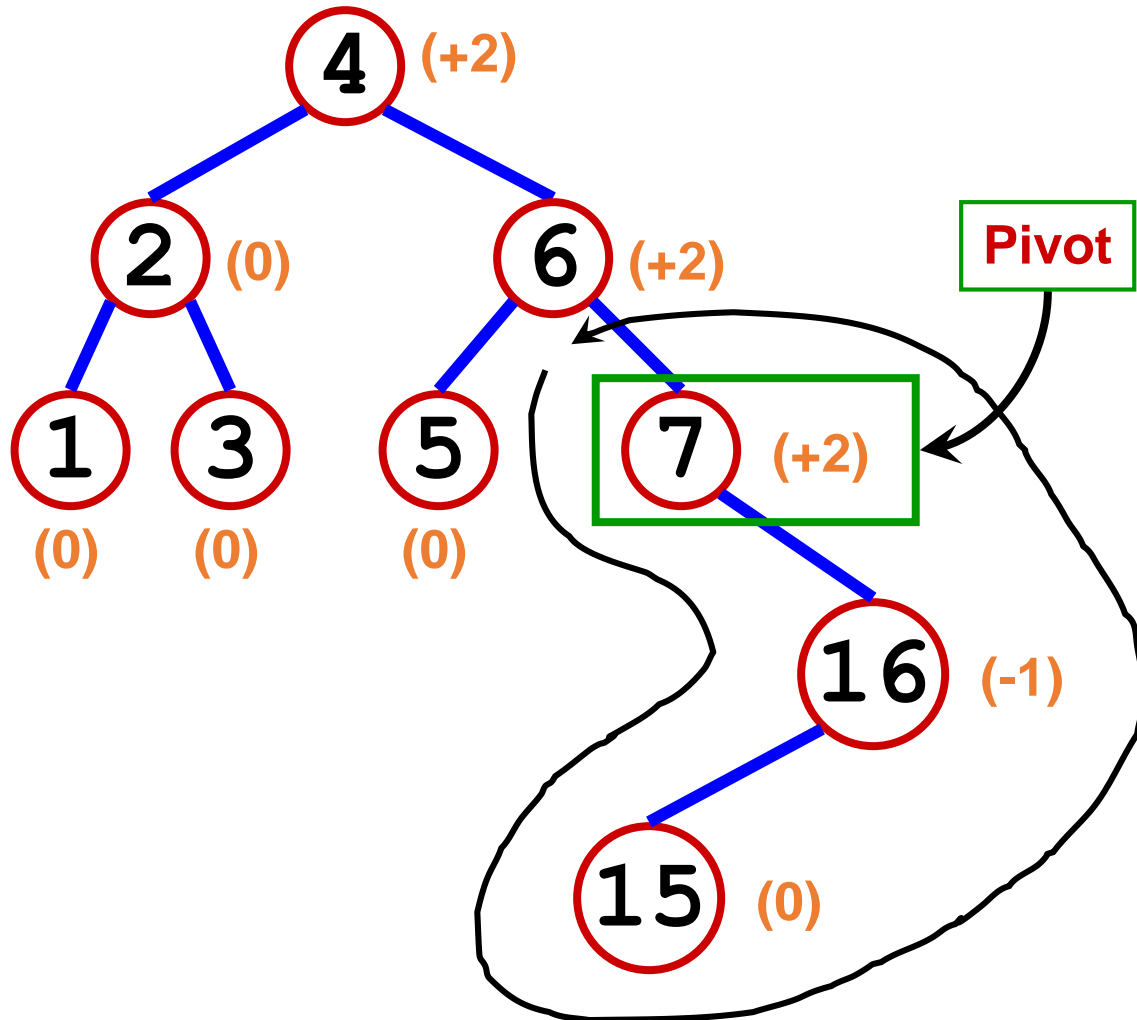


```

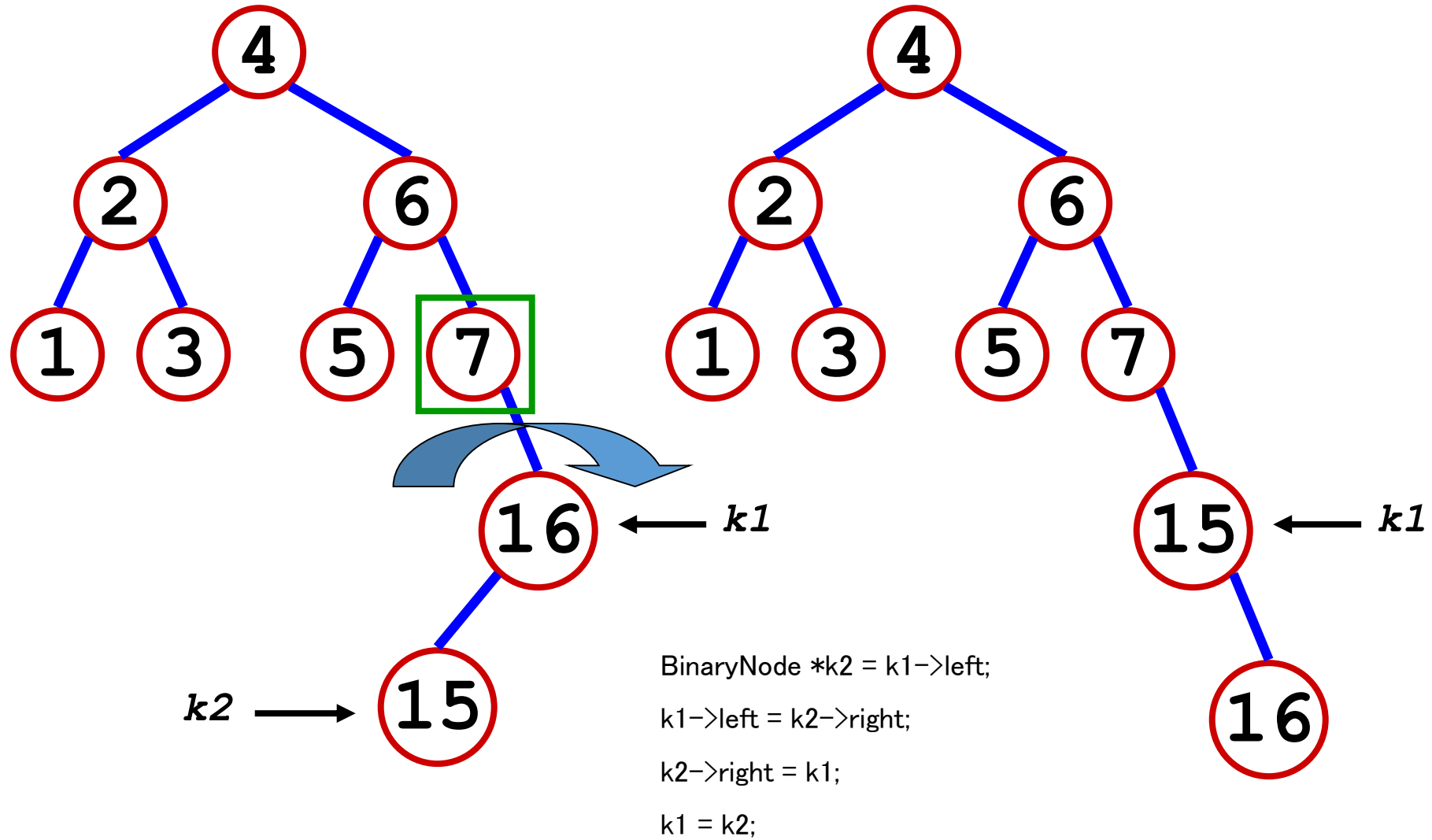
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    k2->left = k1;
    k1 = k2;
}
  
```

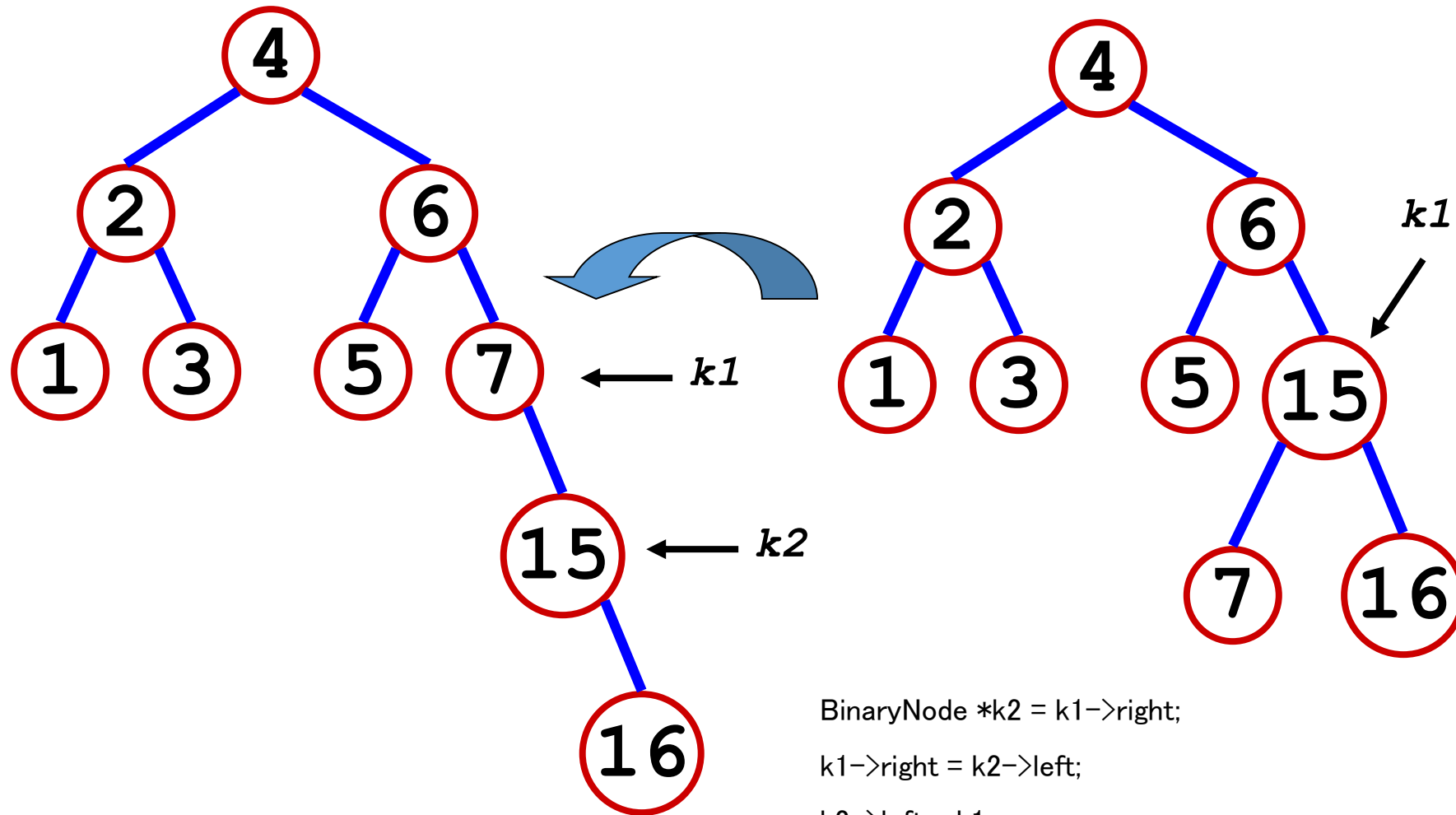


(non-AVL)
Insert 15



Double Rotation



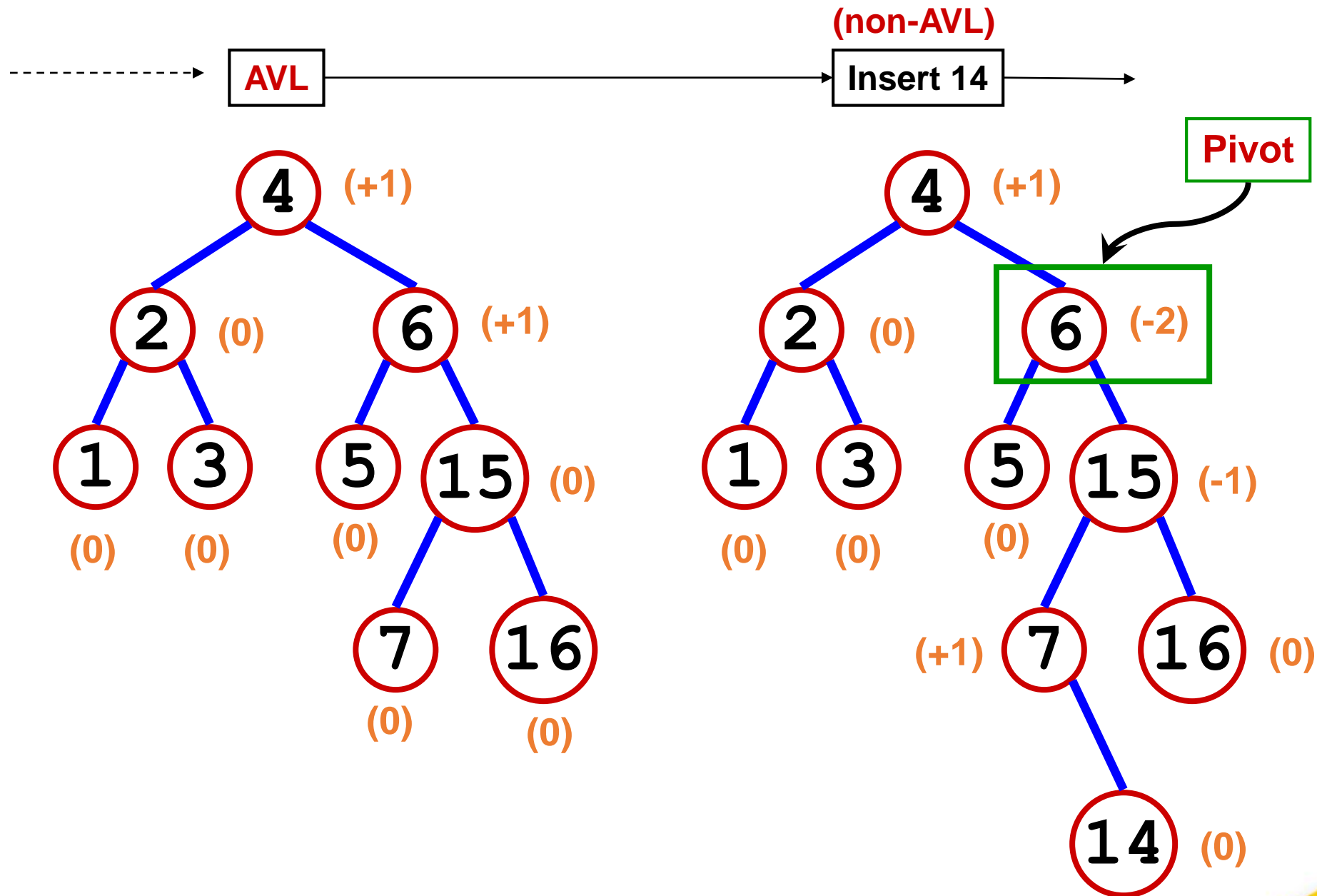


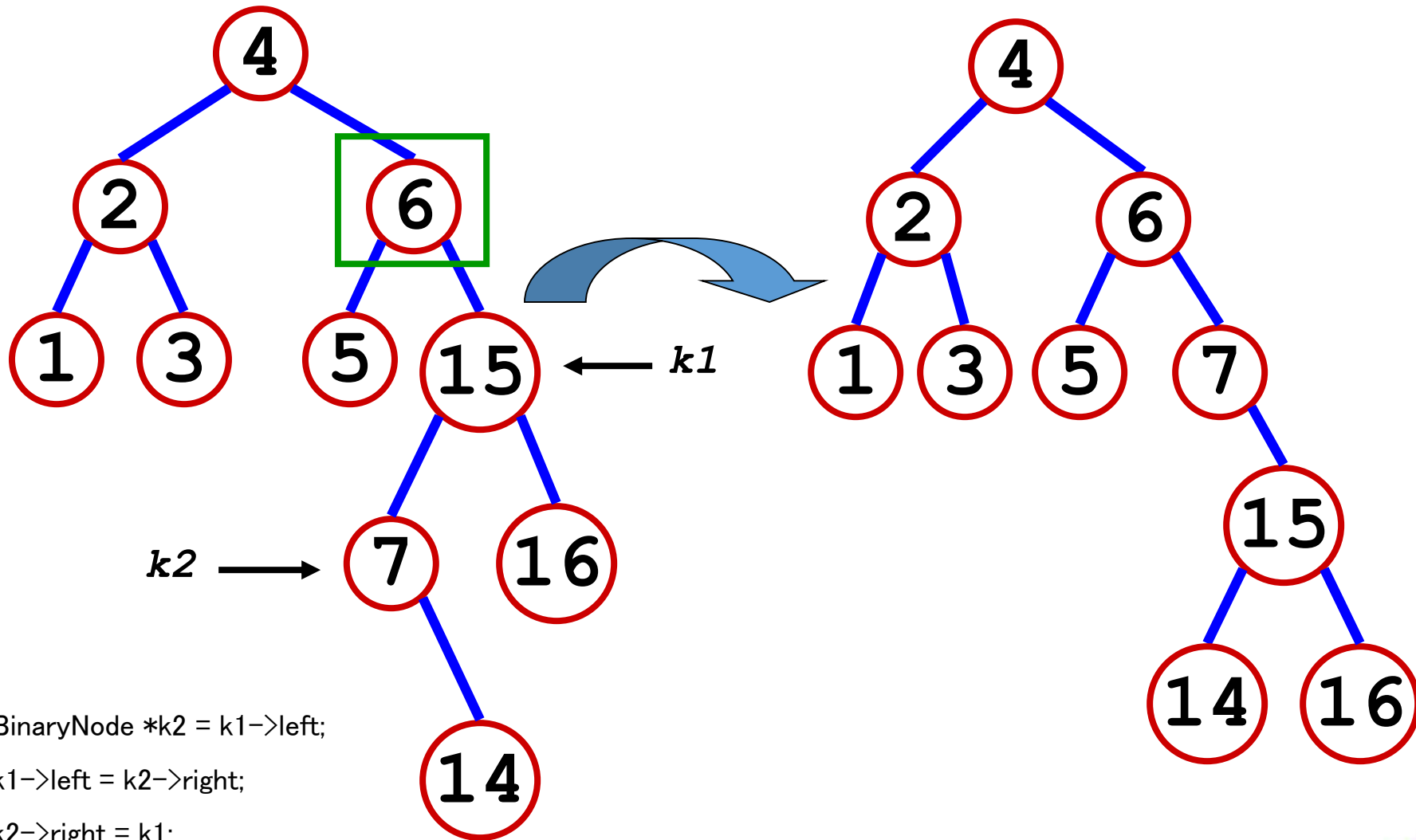
```
BinaryNode *k2 = k1->right;
```

```
k1->right = k2->left;
```

```
k2->left = k1;
```

```
k1 = k2;
```





```
BinaryNode *k2 = k1->left;
```

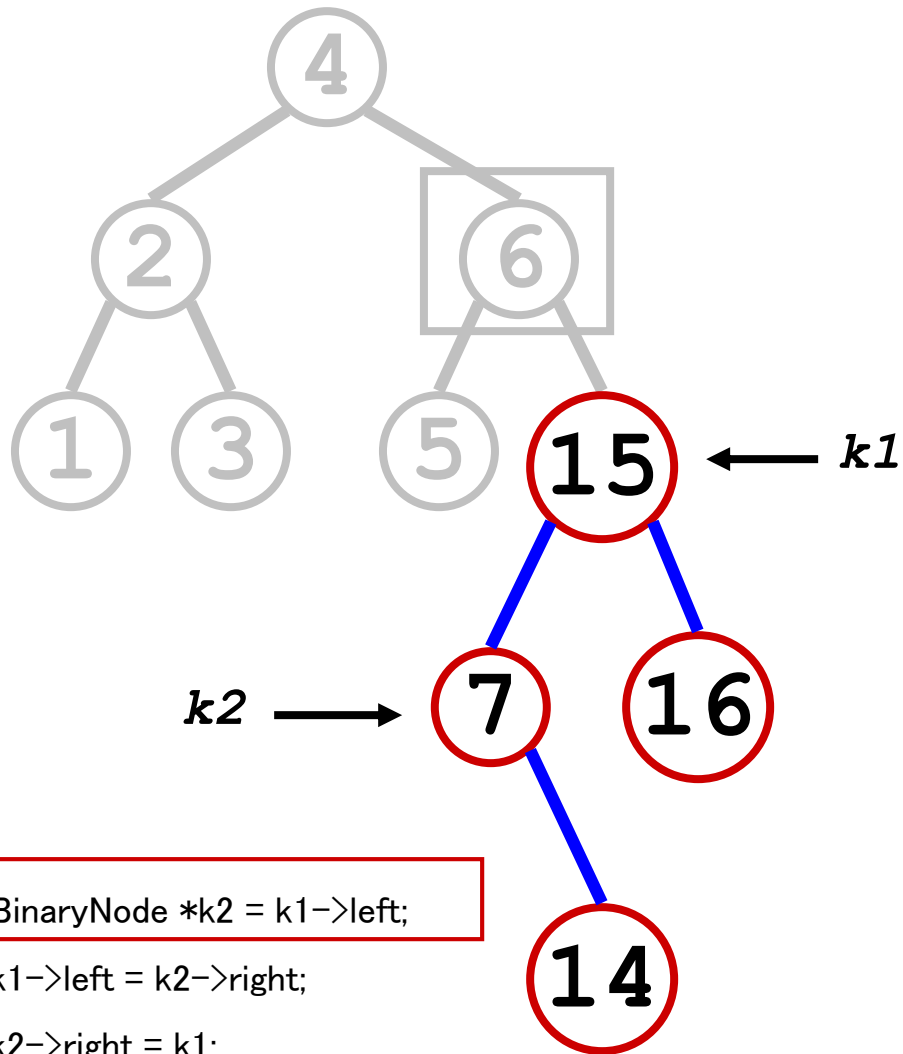
```
k1->left = k2->right;
```

```
k2->right = k1;
```

```
k1 = k2;
```


Double Rotation

Step 1: Rotate child and grandchild



```
BinaryNode *k2 = k1->left;
```

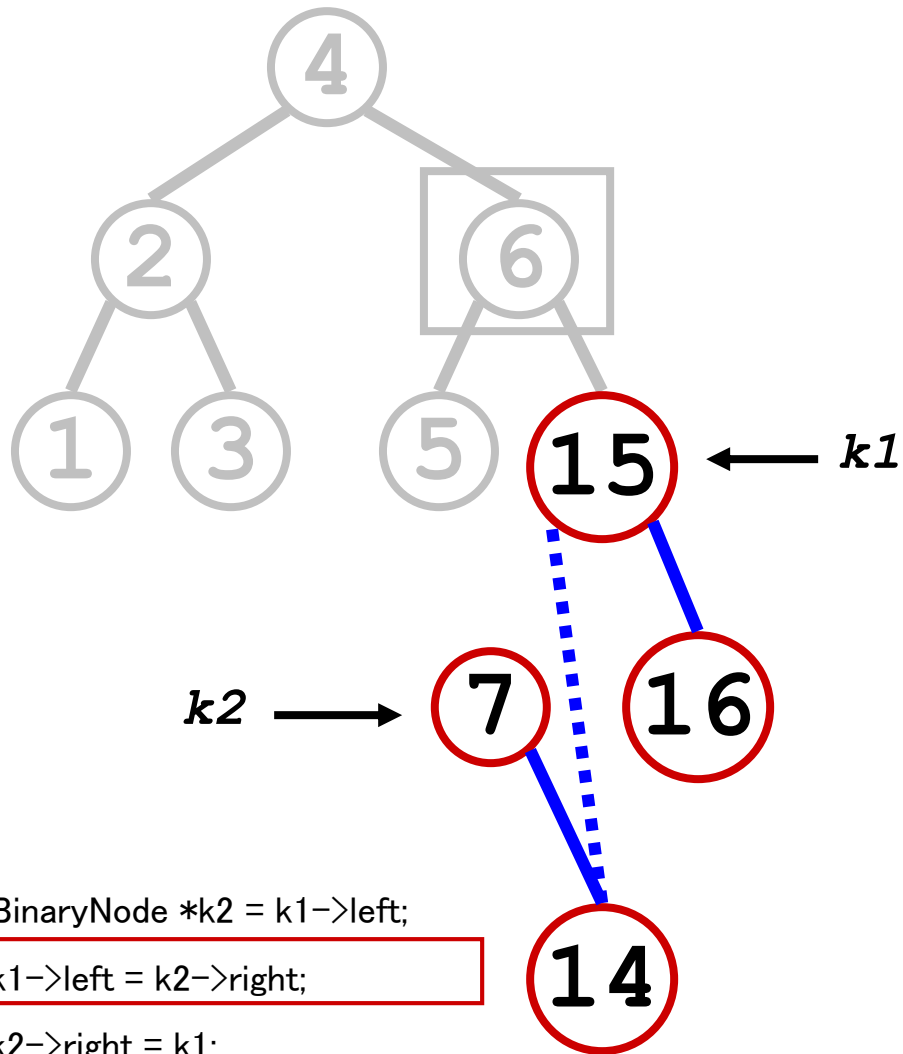
```
k1->left = k2->right;
```

```
k2->right = k1;
```

```
k1 = k2;
```

Double Rotation

Step 1: Rotate child and grandchild



```
BinaryNode *k2 = k1->left;
```

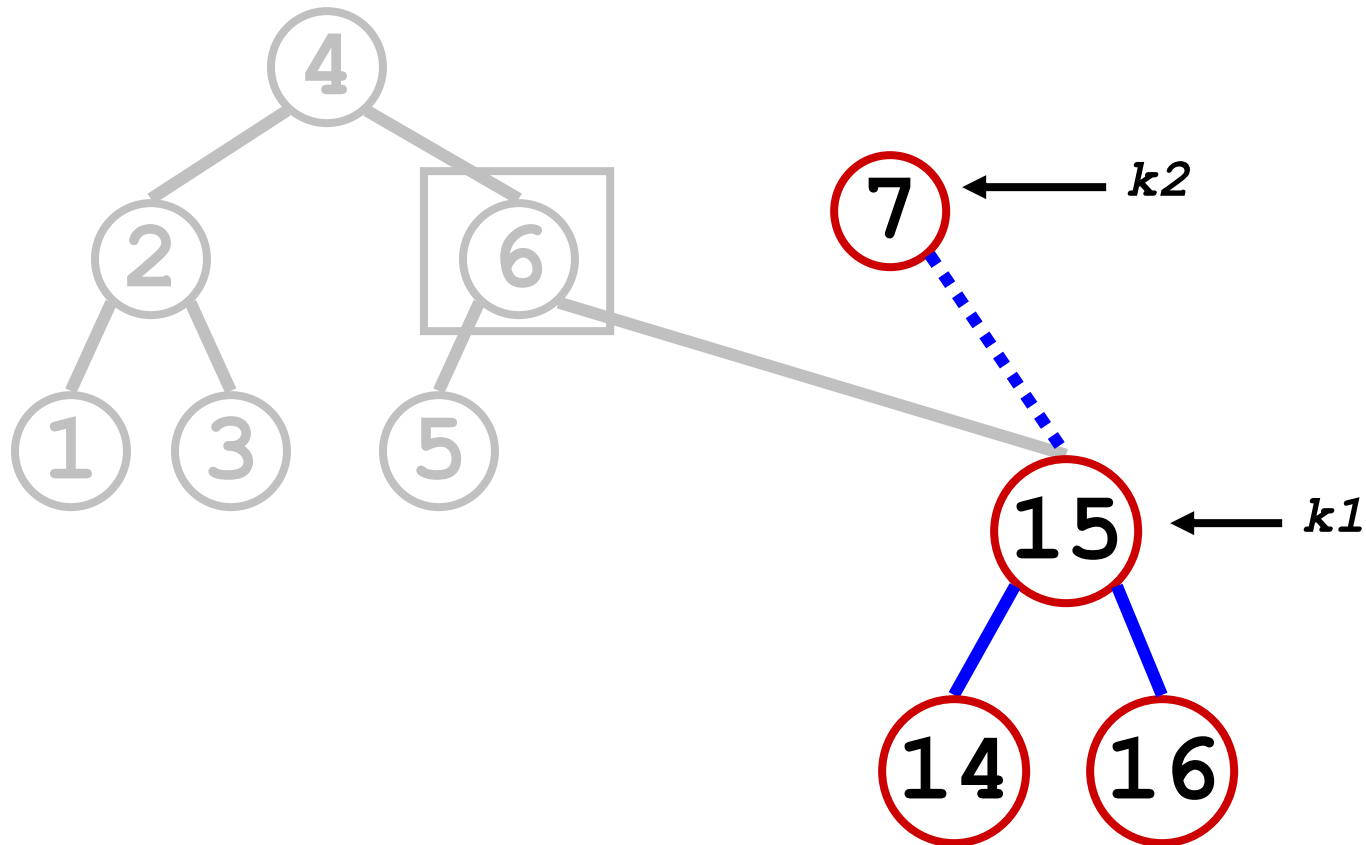
```
k1->left = k2->right;
```

```
k2->right = k1;
```

```
k1 = k2;
```

Double Rotation

Step 1: Rotate child and grandchild



```
BinaryNode *k2 = k1->left;
```

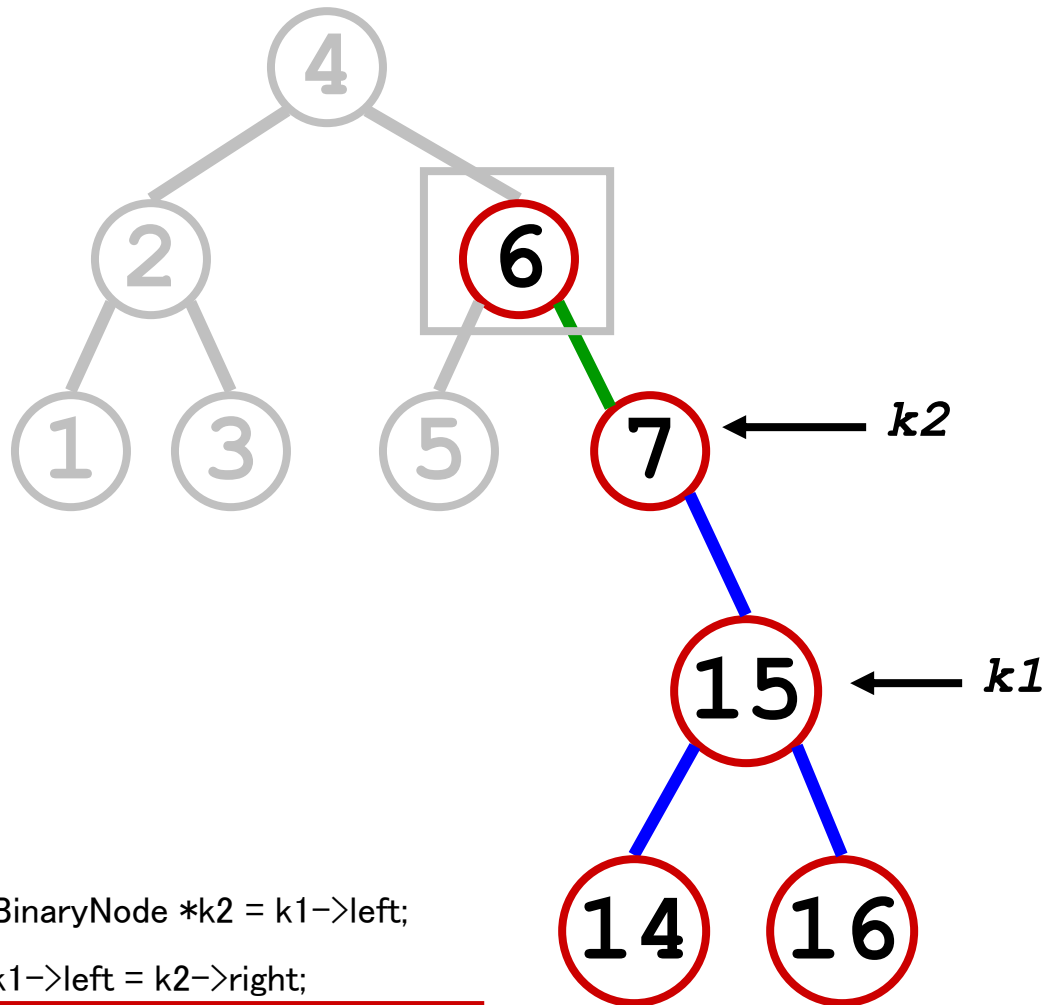
```
k1->left = k2->right;
```

```
k2->right = k1;
```

```
k1 = k2;
```

Double Rotation

Step 1: Rotate child and grandchild



```
BinaryNode *k2 = k1->left;
```

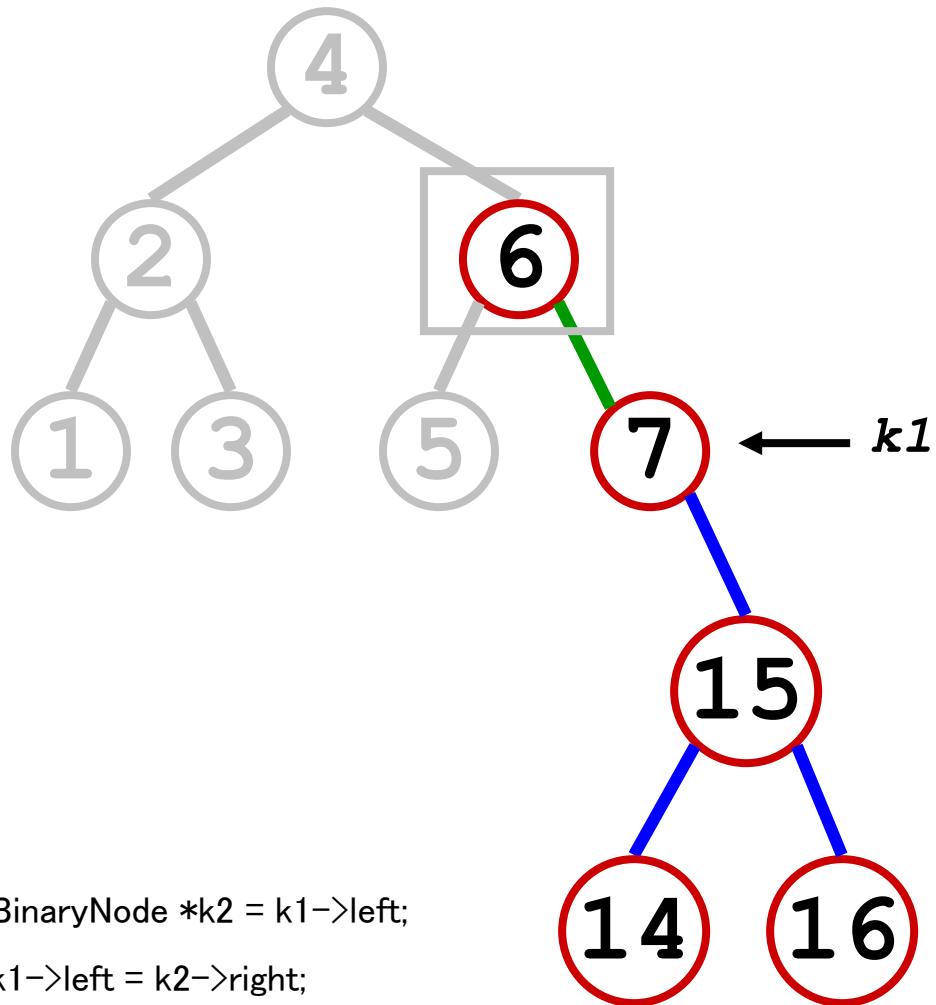
```
k1->left = k2->right;
```

```
k2->right = k1;
```

```
k1 = k2;
```

Double Rotation

Step 1: Rotate child and grandchild

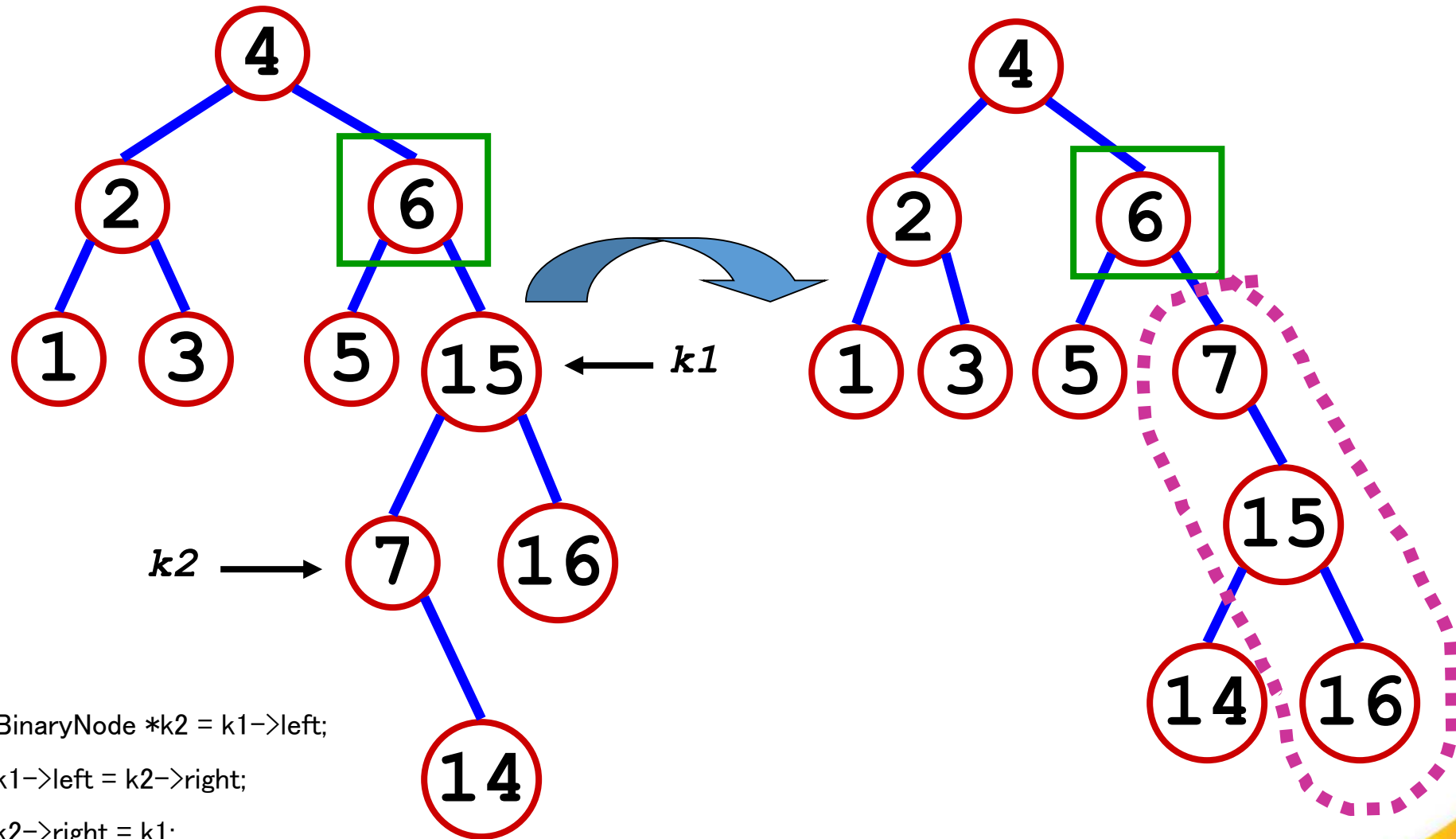


```
BinaryNode *k2 = k1->left;
```

```
k1->left = k2->right;
```

```
k2->right = k1;
```

```
k1 = k2;
```



```
BinaryNode *k2 = k1->left;
```

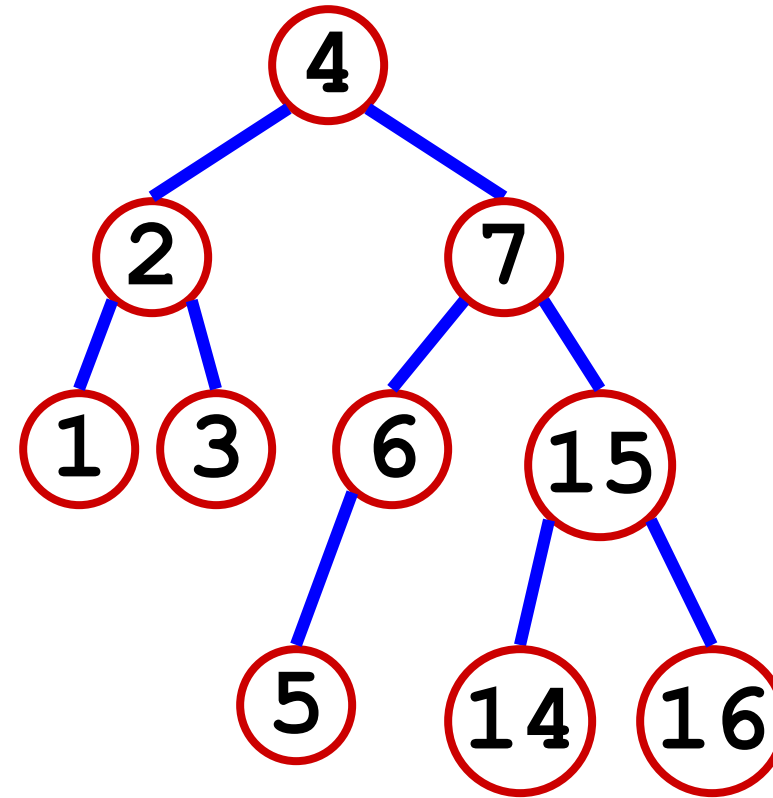
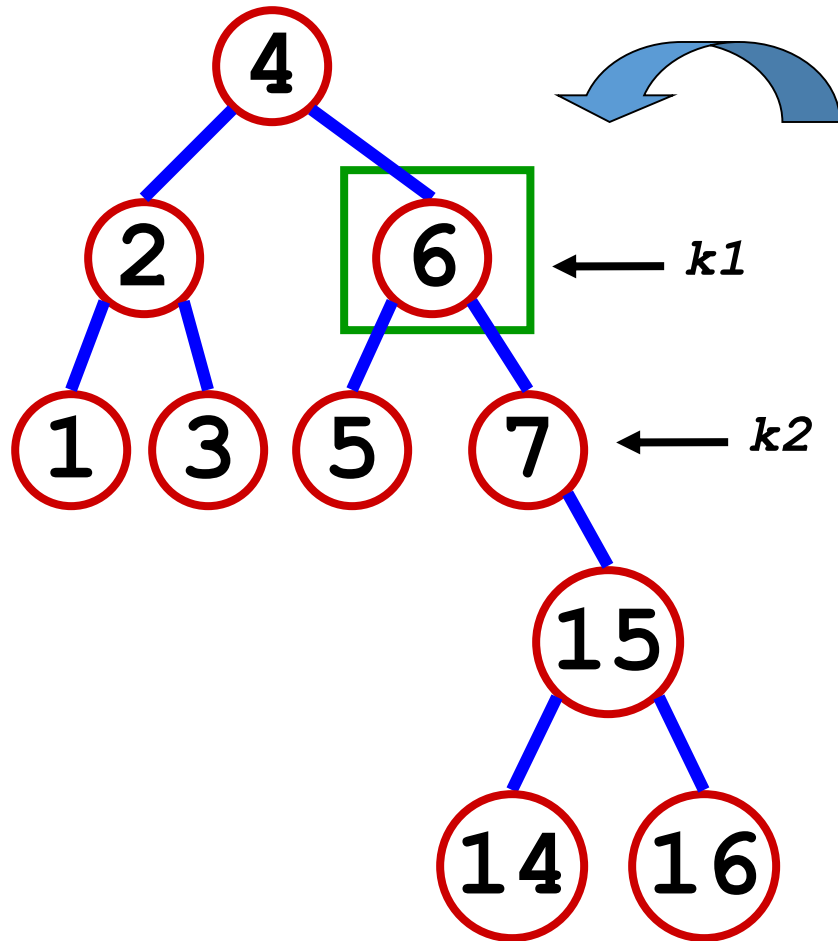
```
k1->left = k2->right;
```

```
k2->right = k1;
```

```
k1 = k2;
```

Double Rotation

Step 2: Rotate node and new child (AVL)



```
BinaryNode *k2 = k1->right;
```

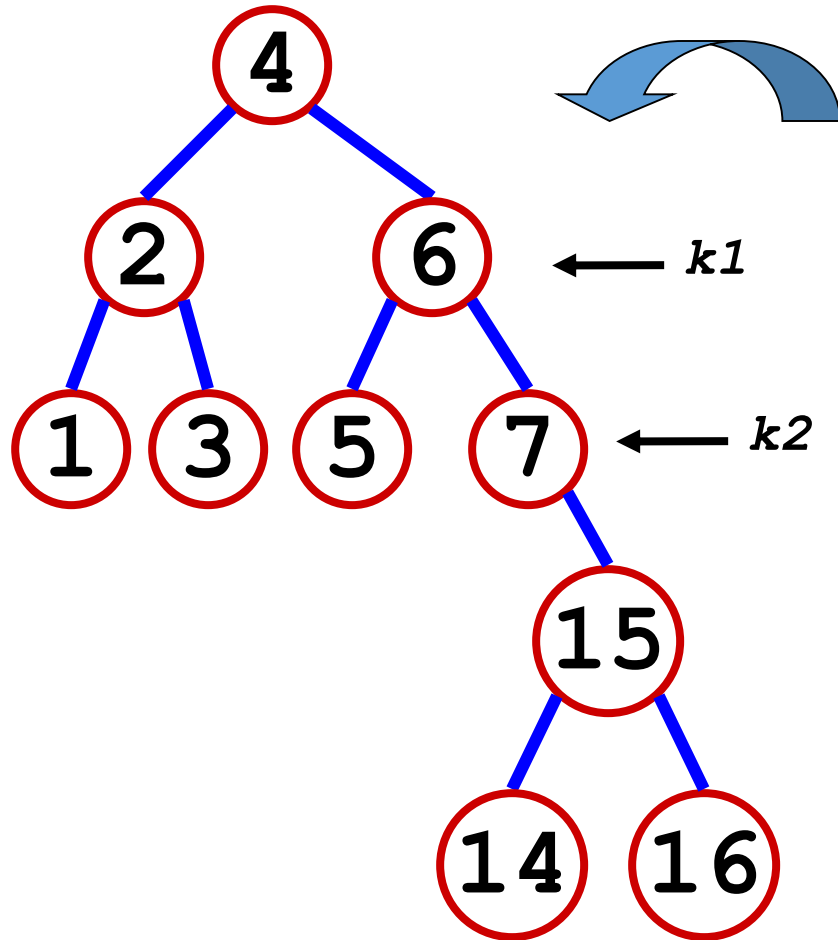
```
k1->right = k2->left;
```

```
k2->left = k1;
```

```
k1 = k2;
```

Double Rotation

Step 2: Rotate node and new child (AVL)

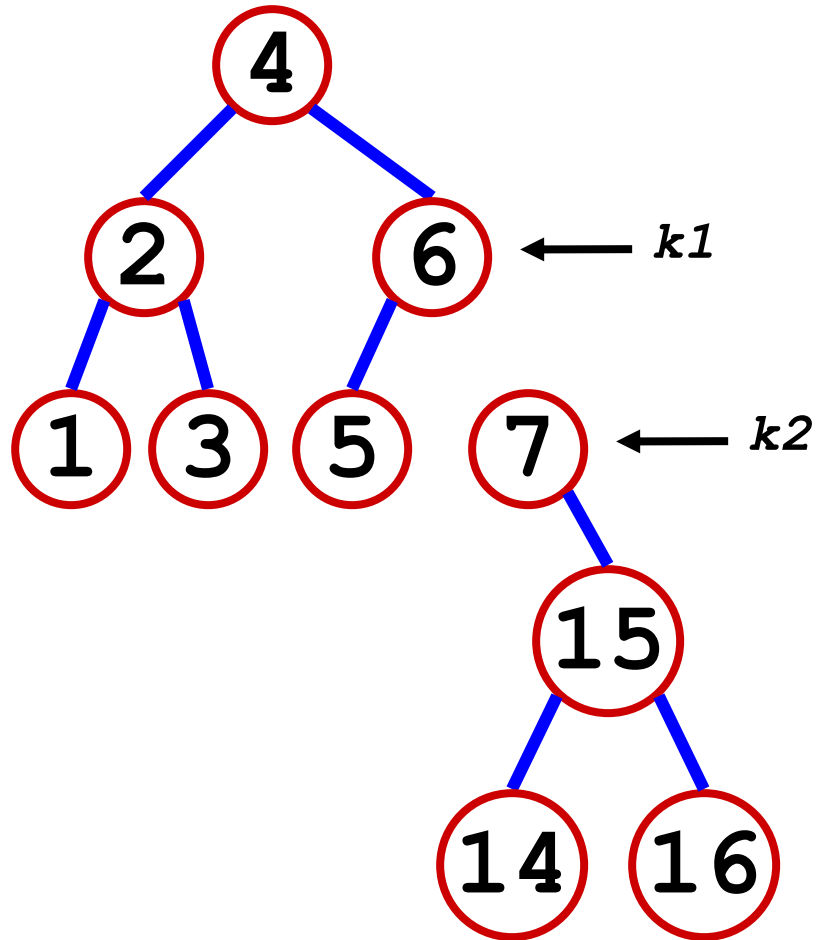


```
BinaryNode *k2 = k1->right;
```

```
k1->right = k2->left;
```

```
k2->left = k1;
```

```
k1 = k2;
```

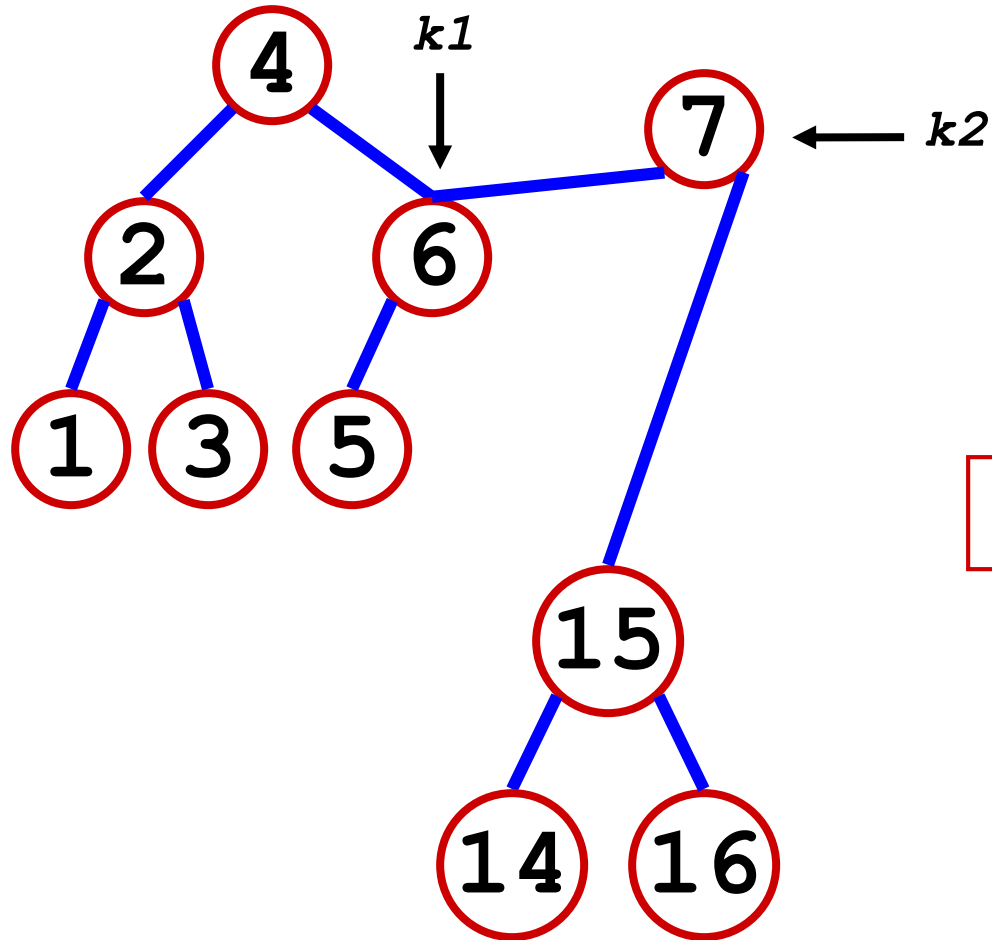



```
BinaryNode *k2 = k1->right;
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k2->left = k1;
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k1 = k2;
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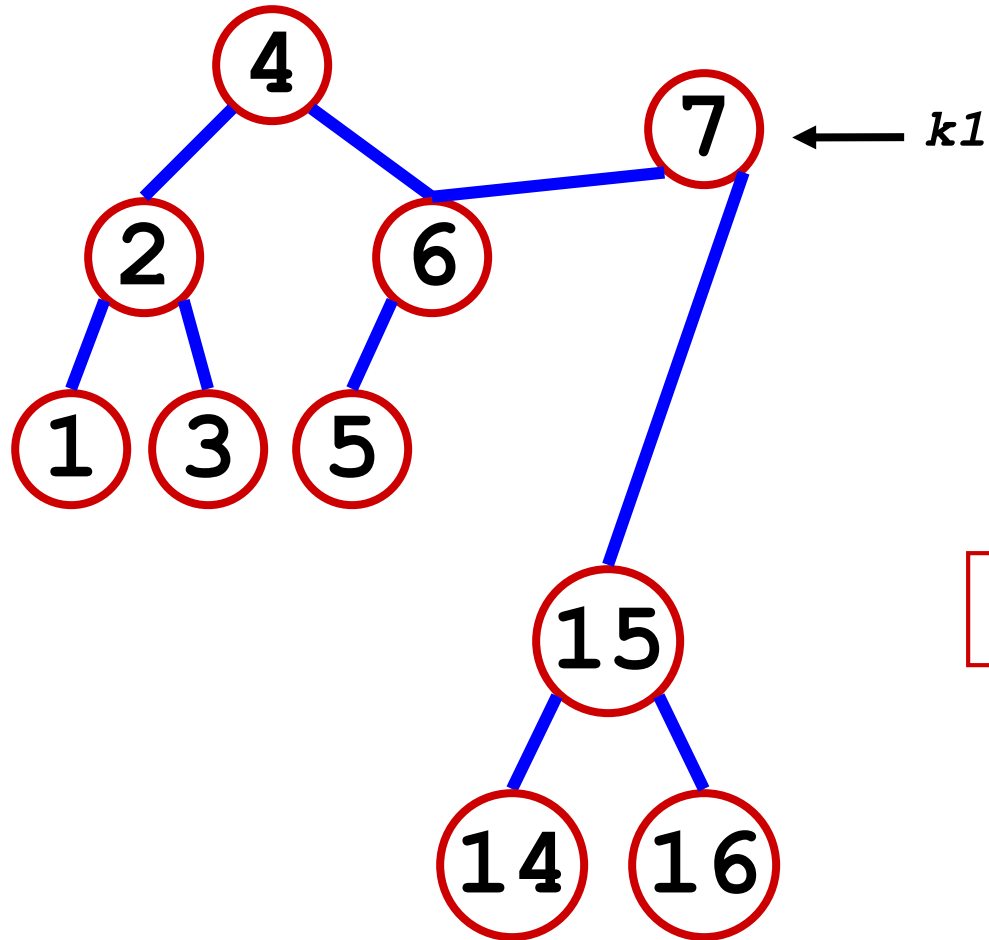


```
BinaryNode *k2 = k1->right;
```

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k1->right = k2->left;
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k2->left = k1;
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```
k1 = k2;
```



```
BinaryNode *k2 = k1->right;
```

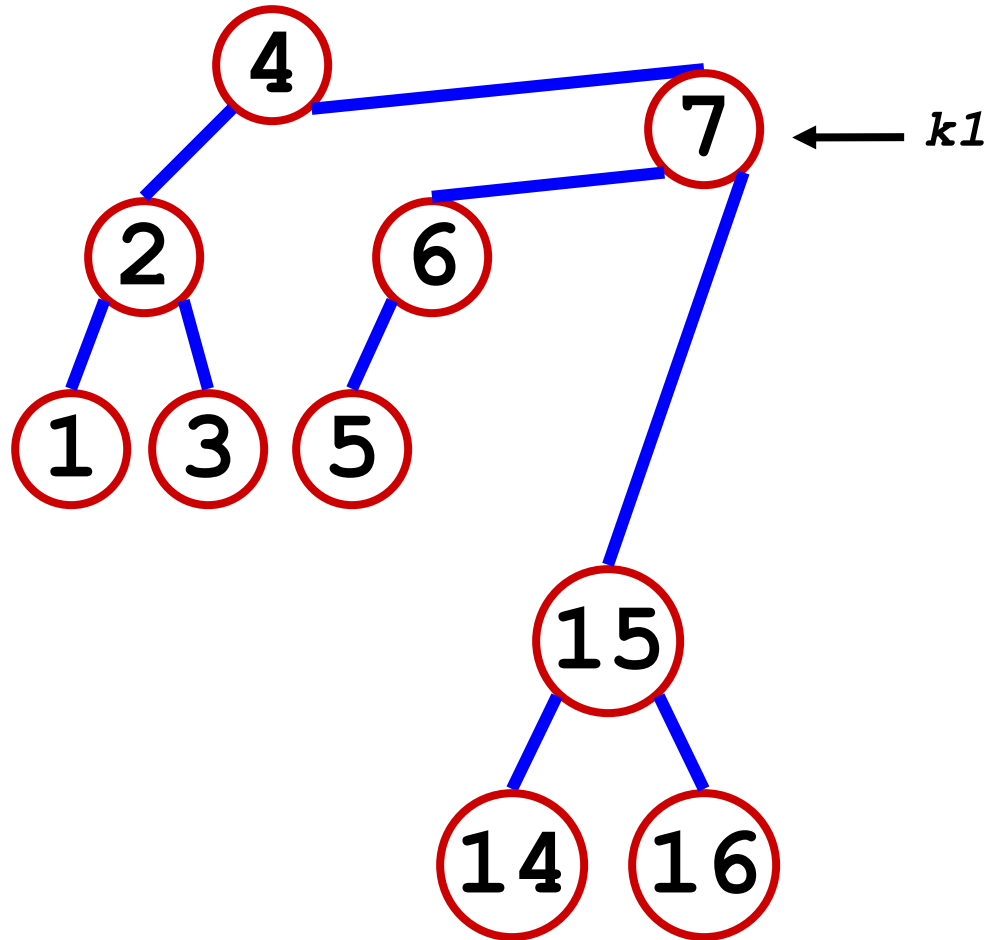
```
k1->right = k2->left;
```

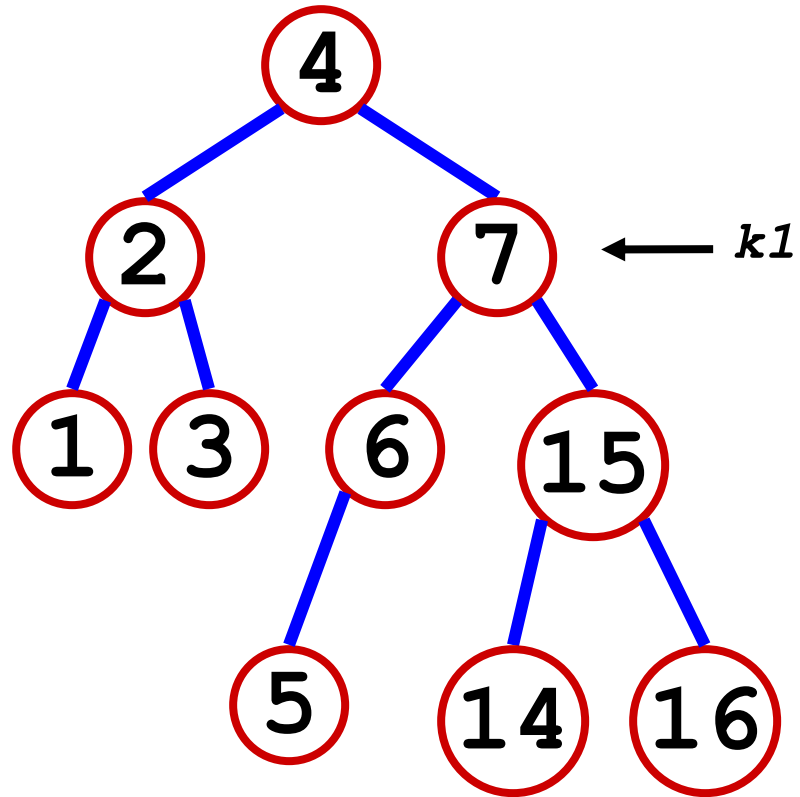
```
k2->left = k1;
```

```
k1 = k2;
```

Double Rotation

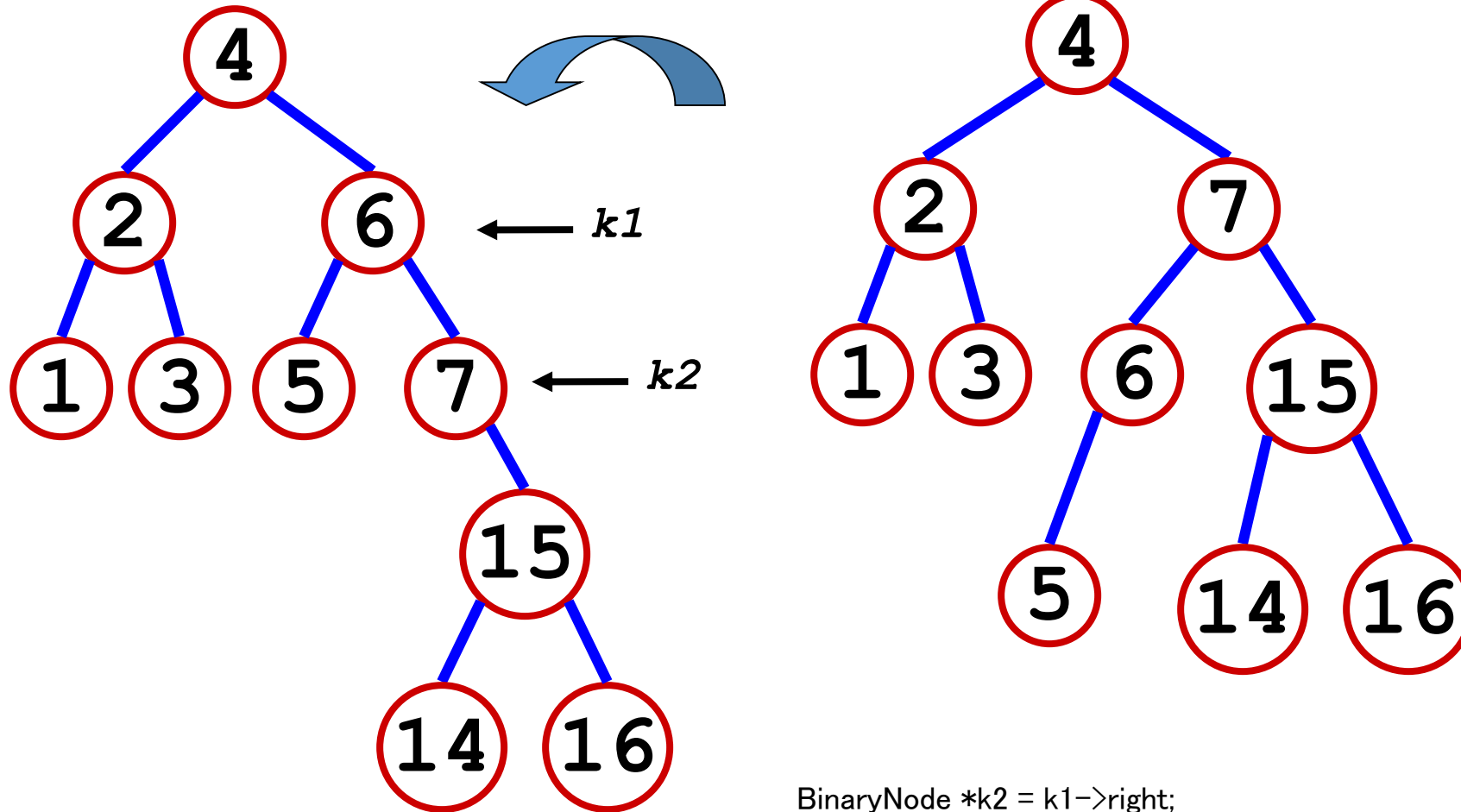
Step 2: Rotate node and new child (AVL)





Double Rotation

Step 2: Rotate node and new child (AVL)



```
BinaryNode *k2 = k1->right;
```

```
k1->right = k2->left;
```

```
k2->left = k1;
```

```
k1 = k2;
```