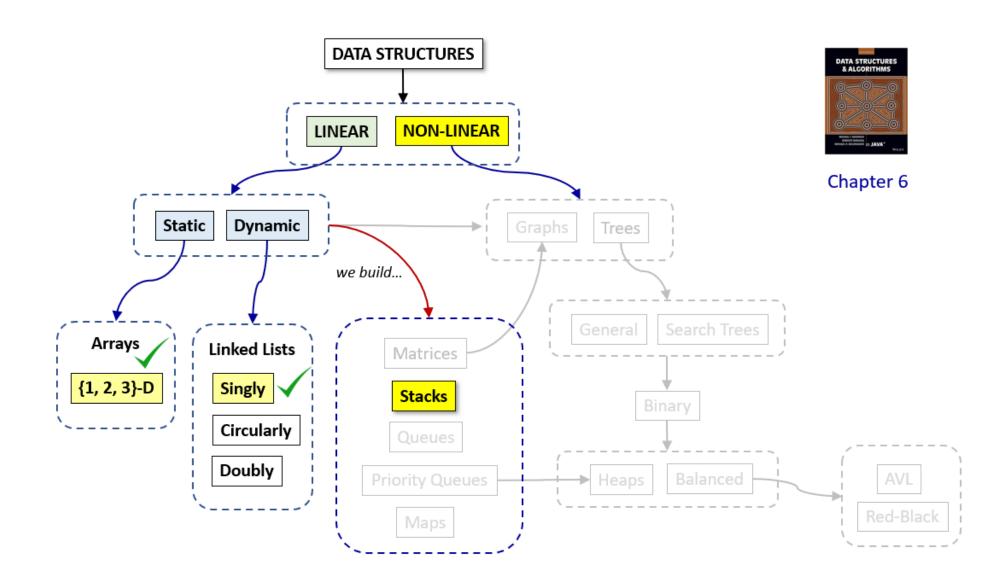
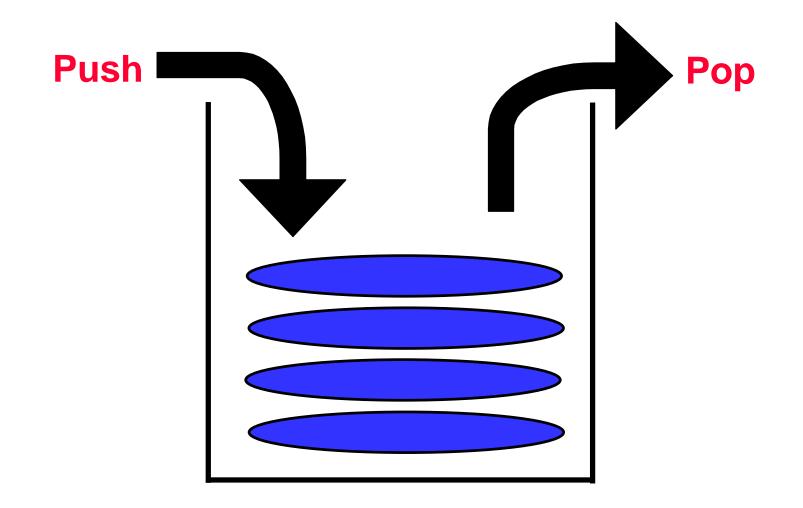


Stacks and Queues



The Stack



The Stack ADT

- The Stack ADT stores arbitrary objects
- Insertions and deletions follow the last-in first-out (LIFO) scheme
- Think of a spring-loaded plate dispenser
- Main stack operations:
 - push(object): inserts an element
 - object pop(): removes and returns the last inserted element

The Stack ADT

- Auxiliary stack operations:
 - object top(): returns the last inserted element without removing it
 - integer size(): returns the number of elements stored
 - boolean isEmpty(): indicates whether no elements are stored

Applications of Stacks

Direct applications

- Page-visited history in a Web browser
- Undo sequence in a text editor
- Activation Stack (recursive calls)
- Chain of method calls in the Java Virtual Machine
- (...)

Properties

```
Idea: a "Last In, First Out" (LIFO) data structure
```

Behaviors:

- Push: Add to top of stack
- Pop: Remove from top of stack (and return that top value)
- Top: Return topmost item (but leave it on the stack)
- Is Full: is it full?
- Is_Empty: is it empty?
- Initialize: empty stack

The Stack as a Logical Data Structure

- The stack is an idea
- It implies a set of logical behaviors
- It can be implemented various ways
 - Using a linked list or a tree or an array
- In this example, we'll focus on dynamic implementations using dynamic data...

Array-based Stack

- A simple way of implementing the Stack ADT uses an array
- We add elements from left to right
- A variable keeps track of the index of the top element

```
Algorithm size()
 return t+1
Algorithm pop()
 if is Empty() then
   throw EmptyStackException
  else
   t \leftarrow t - 1
   return S[t+1]
```



Array-based Stack (cont.)

- The array storing the stack elements may become full
- A push operation will then throw a FullStackException
 - Limitation of the array-based implementation
 - Not intrinsic to the Stack ADT

```
Algorithm push(o)

if t = S.length - 1 then

throw FullStackException

else

t \leftarrow t + 1

S[t] \leftarrow o
```



Performance and Limitations

Performance

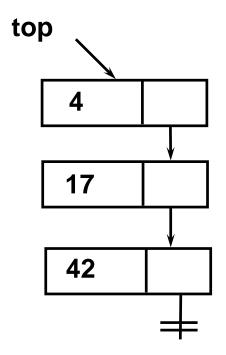
- Let *n* be the number of elements in the stack
- The space used is O(n)
- Each operation runs in time O(1)

Limitations

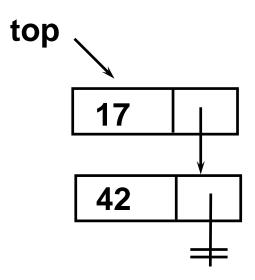
- The maximum size of the stack must be defined a priori and cannot be changed
- Trying to push a new element into a full stack causes an implementationspecific exception

Stacks: Dynamic Implementation

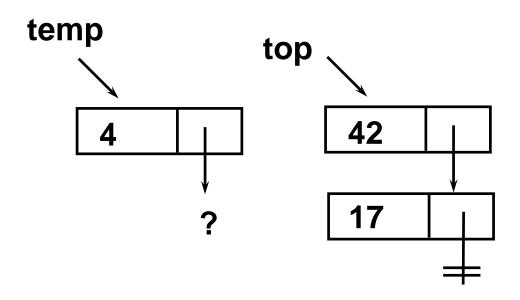
 A singly linked list with restricted set of operations to change its state: only modified from one end



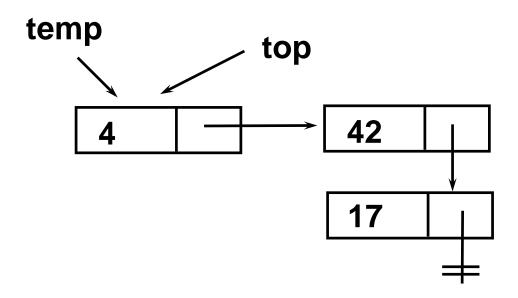
- Create new node
- Add it to the front



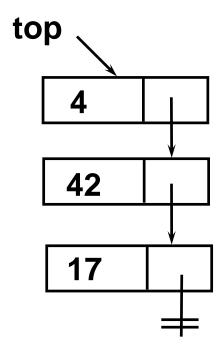
- Create new node
- Add it to the front



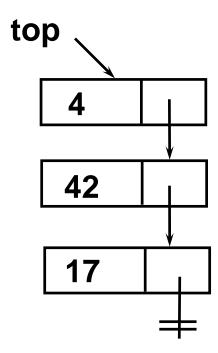
- Create new node
- Add it to the front



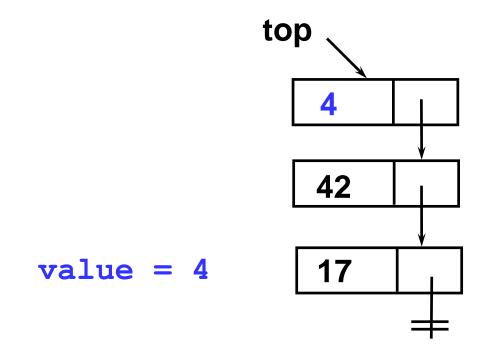
- Create new node
- Add it to the front



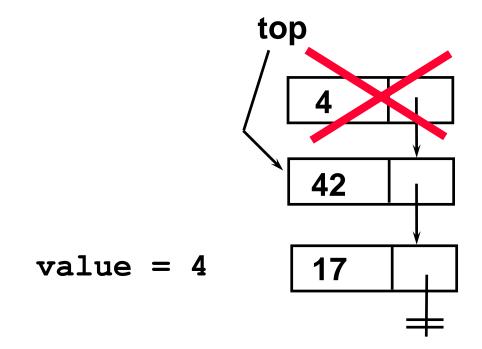
- Capture the first value (to return)
- Remove the first node (move top to next)



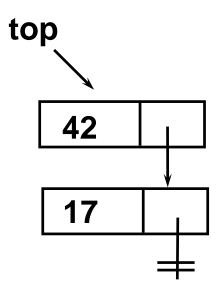
- Capture the first value (to return)
- Remove the first node (move top to next)



- Capture the first value (to return)
- Remove the first node (move top to next)



- Capture the first value (to return)
- Remove the first node (move top to next)



value (4) is returned

Summary: Stack

- Allow us to model "last-in, first-out" (LIFO) behavior
- Can be implemented using different data types
- Behavior is important (and defines a Stack)
 - Push to the front
 - Pop from the front
 - (from the same end)

Why Use an Arrayed Stack?

- No overhead as linked stacks
- Major disadvantage: amount of time to resize → O(n) algorithm

Why Use a Linked Stack?

- Push as many items as you want
- Every operation at the end of the list → pushing and popping are both O(1) algorithms