ENSF 612 Lecture MapReduce Algorithms

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Topics

- Relational operations
- Matrix operations
- Graph algorithms

Relational operations

- Can do SQL-like operations using MapReduce
- Consider select and join operations as examples
- Recap some terminology first
- ◆ Relation (R) a table with names of its columns
- Attributes names of columns
- Tuples rows of a relation
- Schema collection of attributes
 - \bullet denoted by $R(A_1, A_2, ... A_N)$ where the A_S are attributes

Relational operators - cont'd

- How can we implement select using MapReduce?
 - select * from EMPLOYEES where salary>50,000
- Apply condition C on each tuple in a relation
- Can implement using only a map step
- \bullet Input for **Map** is $\langle t, t \rangle$ where t is a tuple in the relation
- \bullet **Map**(< t, t >) outputs < t, t > if C is satisfied; nothing if not.
- Reduce is an identity function simply emits input from Map
- Can either use key or value from Reduce as result

Relational operators- cont'd

- How to implement a join using MapReduce?
- **Compute the natural join** $R(A,B) \bowtie S(B,C)$
 - R and S are each stored in files
 - Tuples are pairs (a,b) or (b,c)

Α	В
a ₁	b ₁
a_2	b ₁
a_3	b_2
a_4	b_3



В	C
b_2	C ₁
b ₂	c_2
b_3	c_3

Α	С
a_3	C ₁
a_3	c_2
a_4	c_3

S

R

Relational operators- cont'd

- **♦** Use a hash function *h* from B-values to 1...k
- **♦** A Map process turns:
 - lacktriangle Each input tuple R(a,b) into key-value pair (b,(a,R))
 - lacktriangle Each input tuple S(b,c) into (b,(c,S))

♦ Map processes

- ◆ send each k-v pair with key b to Reduce process h(b)
- lacktriangle Hadoop does this automatically; just tell it what k is.

♦ Each **Reduce process**

- lacktriangle matches all the pairs (b,(a,R)) with all (b,(c,S))
- lacktriangle outputs (a,b,c).

Graph algorithms

- Graph algorithms have many applications
 - Social network analysis
 - Shortest path in road networks
- Consider the breadth first search (BFS) algorithm
- Traverse a graph in a breadth first manner
- Many applications e.g., calculating shortest path
- How can we implement BFS using MapReduce?

Graph algorithms - cont'd

- BFS terminology
- Graph consists of nodes
- Two nodes may be connected by an edge
- Designate one node as source
- Nodes have colours
 - ◆ White (W) node has not been visited
 - Grey (G) node visited but its kids not visited, i.e., W
 - ◆ Black (B) node visited all its kids are G
- Goal start from source -stop when no G nodes

Wikipedia:

nodes.[2]

https://en.wikipedia.org/wiki/Breadth-first_search

Breadth-first search (BFS) is an algorithm for traversing or searching tree or graph data structures. It starts at the tree root (or some arbitrary node of a graph, sometimes referred to as a 'search key'[1]), and explores all of the neighbor nodes at the present depth prior to moving on to the nodes at the next depth level. It uses the opposite strategy of depth-first search, which instead explores the node branch as far as possible before being forced to backtrack and expand other

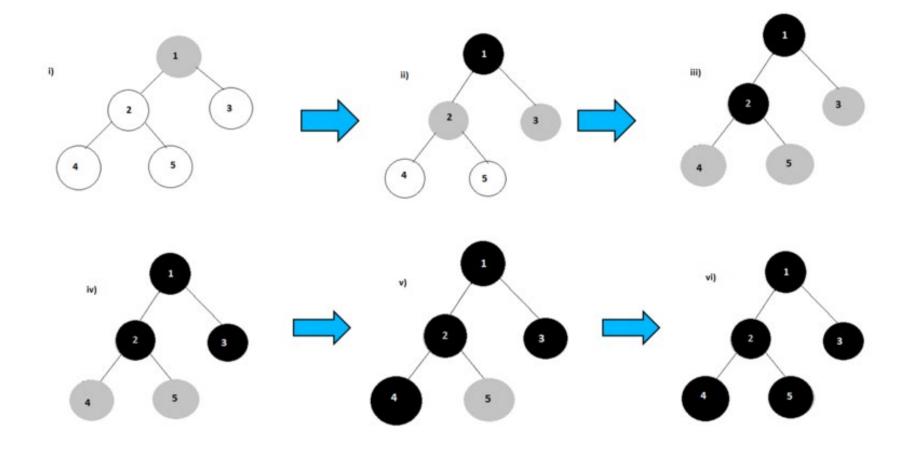
Pseudocode [edit]

Input: A graph G and a starting vertex root of G

Output: Goal state. The *parent* links trace the shortest path back to $root^{[7]}$

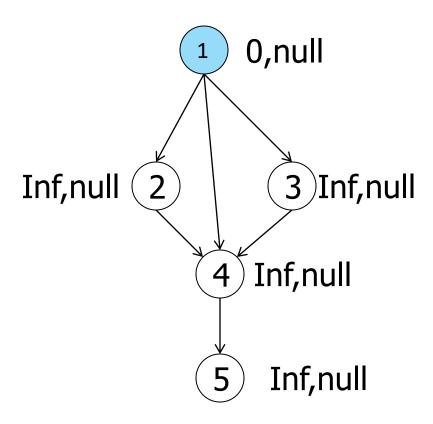
```
procedure BFS(G, root) is
        let Q be a queue
 3
        label root as discovered
        O. engueue (root)
 5
        while Q is not empty do
            v := Q.dequeue()
            if v is the goal then
            for all edges from v to v in
G.adjacentEdges(v) do
                if w is not labeled as discovered
10
then
11
                    label w as discovered
12
                    Q.enqueue(w)
```

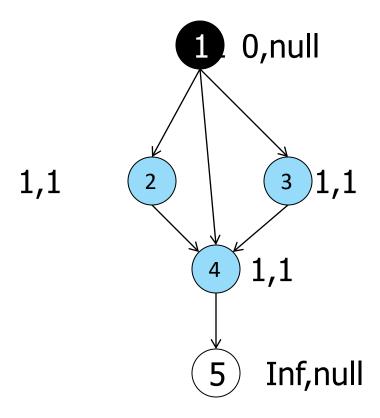
Queue – First in First Out (FIFO) Stack – Last in First Out (LIFO)

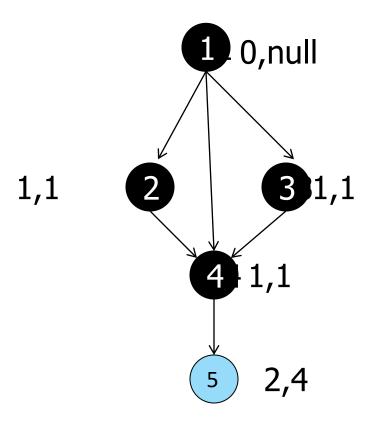


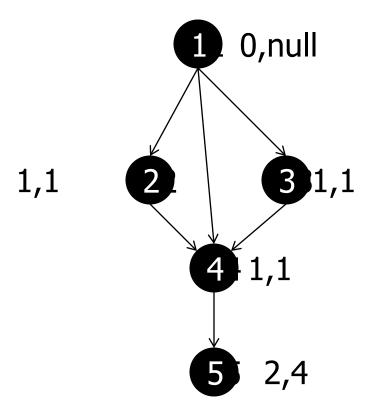
- Using BFS to get shortest path from source to nodes
- Need a couple of extra variables for a node
- Score in how many hops can I get to source node
 - Initially 0 for source and infinity for others
- Parent from which parent did I visit this node
 - Initially null for all nodes
- How does the algorithm work?

- Start from source (process all "G" nodes)
 - Mark from "G" to "B"
 - ♦ Visit unvisited (i.e., "W") kids mark such kids from "W" to "G"
 - Set parent of kids as the node from which visited, i.e., source
 - Set score of kids as score of parent + 1, i.e., 0+1
- Repeat process for all the new "G" nodes
- Keep going till all nodes reachable from source are "B"
- We visit nodes with score K+1 from those with score K
- Score of destination gives shortest hop to source
- Can also trace shortest path by using parent variable









Finding the Shortest Path using BFS in MapReduce

Intuition

- 1. DistanceTo (startNode) = 0
- 2. For all nodes n directly reachable from startNode, DistanceTo (n) = 1
- 3. For all nodes n reachable from some other set of nodes S, DistanceTo (n) = $1 + min(DistanceTo(m), m \in S$

Algorithm

- 1. A map task receives a node n as a key, and (D, points-to) as its value
 - D is the distance to the node from the start
 - "points-to" is a list of nodes reachable from n
 - $\forall p \in points to, emit(p, D + 1)$
- 2. Reduce task gathers possible distances to a given p and selects the minimum one

- How to implement parallel BFS with MapReduce?
- Assume we have large input file describing graph
- Format of input file is as follows:
- node<tab>adjacency_list|score|colour|parentNode
- node is the key, i.e., node id everything else is a value
- * adjacency_list comma separated list of kid nodes
- parentNode "source" for source node, null for others
- Algorithm is iterative
 - Do Map-Reduce pair produce a result
 - Result is input to another Map-Reduce pair
 - Stop when all nodes processed

- Map step
 - Node is "W" − just emit the line
 - \bullet Node is "B'' just emit the line
 - ♠ Node is "*G*"
 - \bullet Make node "B'' emit a line corresponding to this
 - Make kids of node "G"
 - For each kid node, set parent as the node
 - For each kid node, set score as parent node score + 1
 - Emit records corresponding to each kid node
 - How do we handle kids' adjacency lists?

- Reduce step
 - Aggregate records with same key, i.e., node id
 - Will get multiple records for same node
 - Need to aggregate them and emit one valid record
 - A valid record should contain
 - a valid adjacency list for the node
 - the most up to date colour, i.e., state, of the node
 - the correct score and parent

- Reduce step
 - Questions to ponder
 - Why will there be multiple records for same node?
 - Why will the adjacency list of some records be empty?
 - Why will we get a node marked as both "W" and "G"
 - \bullet Why will we get a node marked as both "G' and "B'?
 - Why will we get 2 different scores and parents for same node?
 - Reducer is somewhat more complex than the mapper
- There are libraries which simply do graph processing
 - Pregel, graphX

Topics

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- Matrix operations
- Graph algorithms

Acknowledgements

Portions of these slides were adapted with permission from Leskovec et al.

(http://www.mmds.org)

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Link to material:

https://hadooptutorial.wikispaces.com/Iterative+MapReduce+and+Counters