

**Theorem** in algorithm RecConcave<sup>1</sup> for privacy parameters  $\epsilon, \delta$ , promise  $r$  database  $X$ , domain  $R$ , sensitivity-1 utility function  $u$  and recursion bound  $N = 2$  then: the step 6 will returns a “good interval” in probability at least  $1 - \beta$  if  $r > \frac{16}{\epsilon\alpha} \ln \left( \frac{\ln(T)}{\beta} \right)$

**Proof** notice that when  $N = 2$  the recursion call actually calls the *Exponential-Mechanism* and retrieves its answer

we want to make sure that the *Exponential-Mechanism* outputs with high probability a good interval

meaning that

$$\Pr \left[ q(S, j) < \frac{3\alpha}{\epsilon} r \right] < \beta$$

by The Algorithmic Foundations of Differential Privacy<sup>2</sup> Theorem 3.11

$$\Pr \left[ q(M_E(x, q, R)) \leq OPT_q(x) - \frac{2\Delta q}{\epsilon} \left( \ln \left( \frac{|R|}{|R_{OPT}|} \right) + t \right) \right] \leq e^{-t}$$

and if we switch  $e^{-t} = \beta$  and take the specific parameters used in our case  $\Delta u = 1$ ,  $|R| = \ln(T)$  and  $|R_{OPT}| = 1$

we get

$$\Pr \left[ q(M_E(x, q, R)) \leq OPT_q(x) - \frac{2}{\epsilon} \ln \left( \frac{\ln(T)}{\beta} \right) \right] \leq \beta$$

we also know that  $\frac{\alpha}{2} r \leq OPT_q(x) \leq r$

if we combine all the above we get that we want

$$\frac{3\alpha}{8} r < \frac{\alpha}{2} r - \frac{2}{\epsilon} \ln \left( \frac{T}{\beta} \right) \Rightarrow r > \frac{16}{\epsilon\alpha} \ln \left( \frac{\ln(T)}{\beta} \right)$$

**Remark** we saw that for  $A_{dist}$  to run we must have

$$r > \frac{8 \ln \left( \frac{1}{\beta\delta} \right)}{3\epsilon\alpha}$$

notice that this bound will be less than the one above iff  $\delta > \frac{\beta^5}{\ln^6(T)}$

**Proof**

$$\begin{aligned} \frac{8 \ln \left( \frac{1}{\beta\delta} \right)}{3\epsilon\alpha} < \frac{16}{\epsilon\alpha} \ln \left( \frac{\ln(T)}{\beta} \right) &\Rightarrow \ln \left( \frac{1}{\beta\delta} \right) < 6 \ln \left( \frac{\ln(T)}{\beta} \right) \Rightarrow \frac{1}{\beta\delta} < \left( \frac{\ln(T)}{\beta} \right)^6 \\ &\Rightarrow \delta > \frac{\beta^5}{\ln^6(T)} \end{aligned}$$

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<sup>1</sup>A. Beimel, K. Nissim, and U. Stemmer. Private learning and sanitization- Pure vs. Approximate Differential Privacy

<sup>2</sup>C.Dwork, A.roth