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Prove LPATH is NP-Complete

Definition of Longest Path (LPATH):

LPATH = $\{ \langle G, a, b, k \rangle \mid G \text{ contains a simple path of length at least } k \text{ from } a \text{ to } b \}$

Step 1 (Verify): Given a LPATH of length k between points a and b on a graph G , verify that a simple path of length at least k exists from a to b in polynomial time.

-Check that the LPATH is simple

```
for  $i = 1 \rightarrow n$ 
{
  for  $j = (i + 1) \rightarrow n$ 
  {
    make sure  $vertex_i \neq vertex_j$ 
  }
}
This takes  $O(n^2)$  time.
```

-Check that all edges in path actually exist

```
for  $i = 1 \rightarrow |\text{edges in LPATH}|$ 
{
  for  $j = (i + 1) \rightarrow |\text{edges in } G|$ 
  {
    make sure  $edge_i = edge_j$  for some  $edge_j \in G$ 
  }
}
This takes  $O(n^2)$  time.
```

-Count number of edges in LPATH

This takes $O(n)$ time.

Thus, the entire verify step takes $O(n^2)$ time which is polynomial.

Step 2 (Reduction): Show that LPATH is NP-Hard. We do this by showing reduction from known NP-Complete problem to the LPATH problem.

$P_{known} = \text{UHAMPATH}$

$P_{new} = \text{LPATH}$

$\text{UHAMPATH} = \{ \langle G, s, t \rangle \mid G \text{ is undirected graph that contains a Hamiltonian path from } s \text{ to } t \}$

Reduction:

$\text{UHAMPATH}(G = (V, E), s, t)$

$k = |V| - 1$

Return the answer to $\text{LPATH}(G, s, t, k)$

The reduction step occurs in polynomial time.

Step 3 (Correctness): Graph G has a Longest Path of size $|V| - 1$ between a and b iff G has a Hamiltonian Path between a and b .

-Case 1 Assume G has Longest Path of size $|V| - 1$ between a and b . Show G has a Hamiltonian Path between a and b .

There exists a simple path from a to b with length $|V| - 1$ (from given Hamiltonian Path). Therefore, there exists a Longest Path of at least length $|V| - 1$ from a to b . This Longest Path is the Hamiltonian Path.

-Case 2 Assume G does not have a Hamiltonian Path between a and b . Show G does not have a Longest Path between a and b of size at least $|V| - 1$.

A longest path with length at least $|V| - 1$ from a to b would be a simple path which visits every vertex v , such that $v \in V$. A simple path that starts at a and ends at b , of length $|V| - 1$, is a Hamiltonian Path. However, no hamiltonian path from a to b exists (given). Therefore, no longest path of length $|V| - 1$ exists.