

ENGINEERING ONLINE

Lecture Notes

Course Number: CSC 513

Instructor: Dr. Singh

Lecture Number: 12

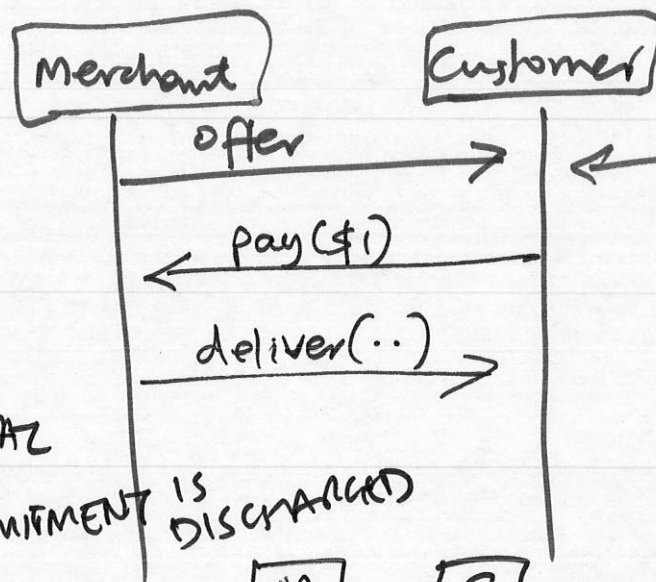


$C_1 = C(\text{MERCHANT, CUSTOMER, } \underline{\text{pay}(\$1)}, \underline{\text{deliver}(\text{marker})})$

Principals
agents
entities that can
enter into
business relationships

propositions: true or false
expressions

BASIS
FOR
JUDGING
CORRECTNESS
OF
ENACTMENTS
(MSCs)



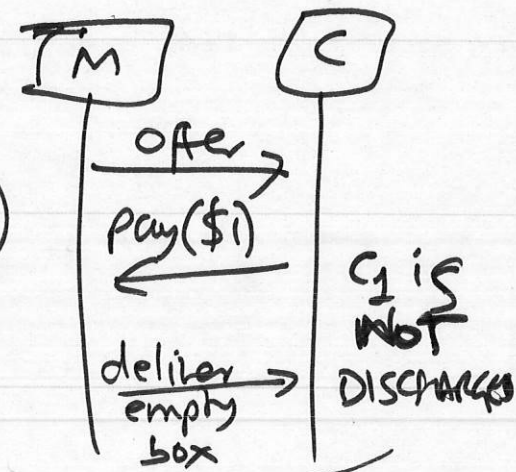
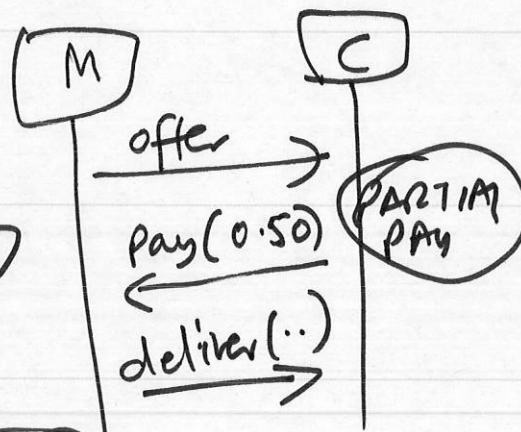
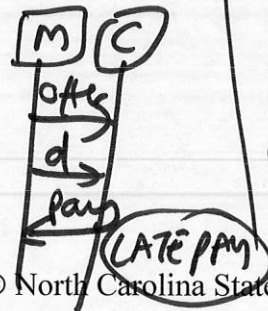
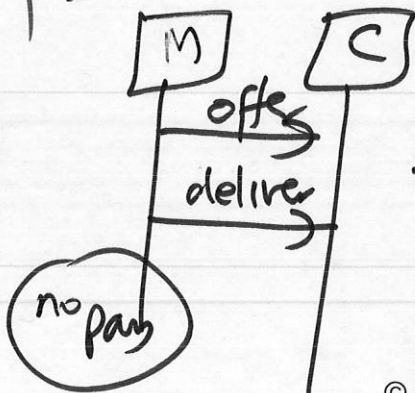
content of the
offer message

What are the next steps
of this MSC?

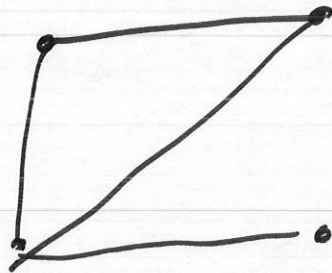
2 IDEAS

THE
COMMITMENT IS
DISCHARGED

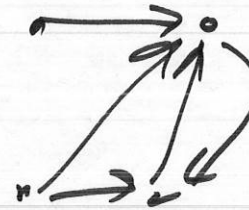
$P \rightarrow Q$
 $Q \rightarrow P$



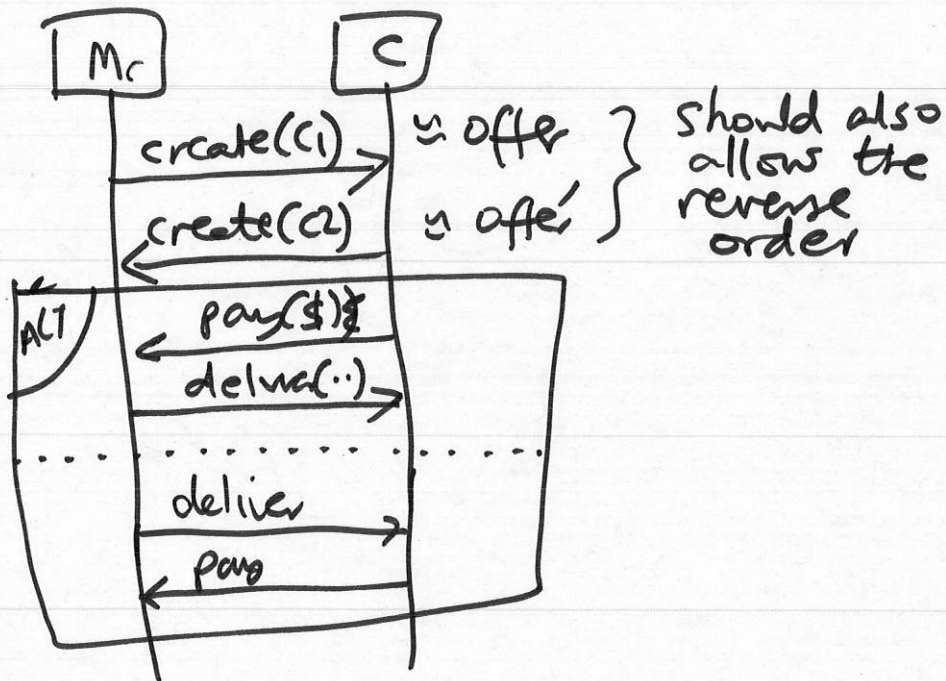
Facebook



Twitter



$C_1 = C(\text{Mer}, \text{Cust}, \text{pay}(\$1), \text{deliver}(\text{marker}))$
 $\wedge C_2 = C(\text{Cust}, \text{Mer}, \text{deliver}(\text{marker}), \text{pay}(\$1))$



Example: Commitment Progression ^{via Logic} _{via Explicit actions}

$C(\text{Buyer, Seller, goods, pay})$ ^{is ACTIVE (AND CONDITIONAL)}

► If goods $\wedge C(\text{Buyer, Seller, goods, pay})$ Then

► $C(\text{Buyer, Seller, T, pay})$

ACTIVE,

UNCONDITIONAL &
= DETACHED
= BASE

► If pay $\wedge C(\text{Buyer, Seller, T, pay})$ Then

► Satisfied

► If pay $\wedge C(\text{Buyer, Seller, goods, pay})$ Then

► Satisfied

NEVER A PAYMENT

(i) Goods (and no pay)
VIOLATION

(ii) no goods (and no pay)

NOT A VIOLATION
NOT SATISFACTION
= EXPIRATION

Can be nested:

$C(\text{Seller, Buyer, pay, } C(\text{Shipper, Buyer, T, deliverGoods}))$

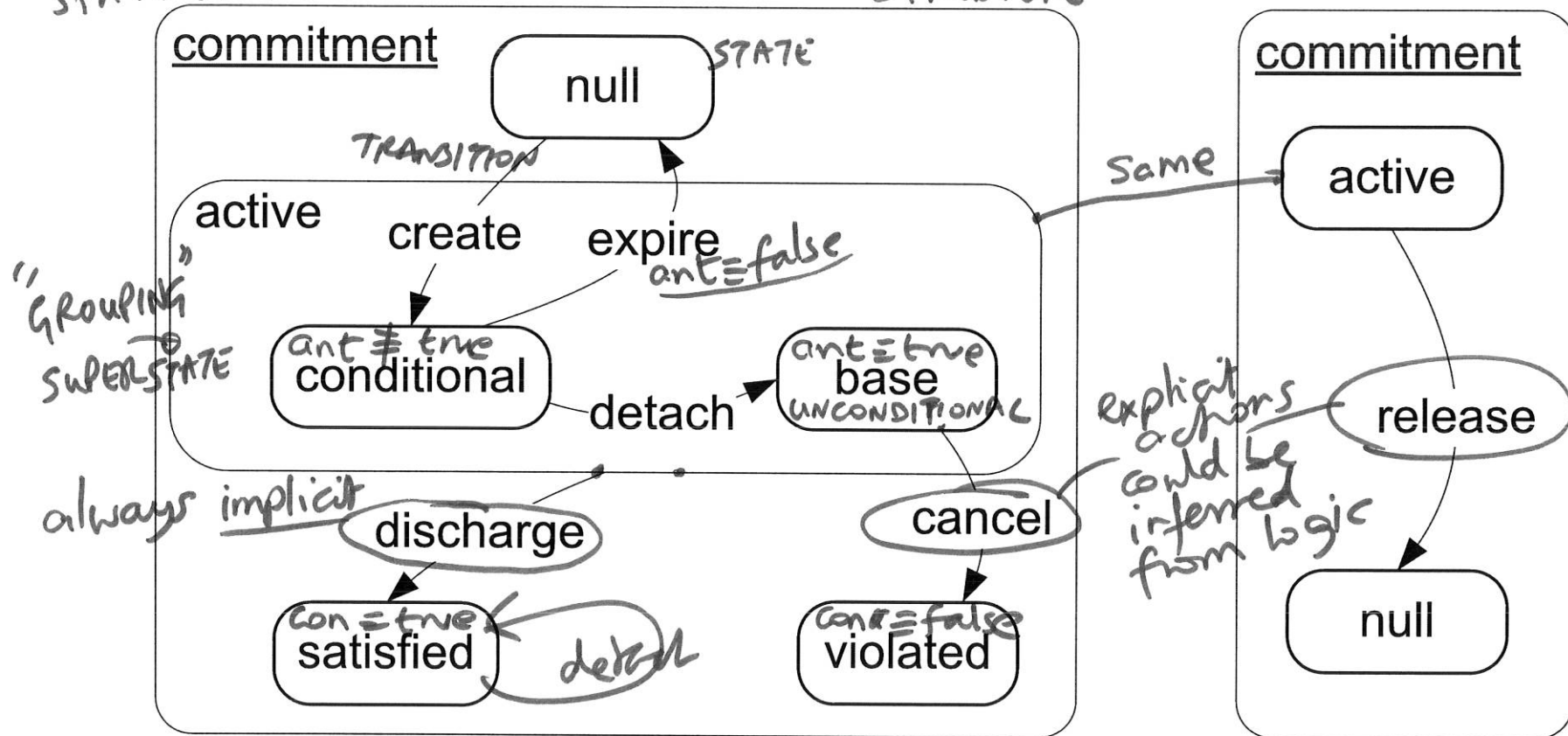
NOT A VIOLATION
NOT SATISFACTION
= EXPIRATION

Commitment Life Cycle (and Patterns)

C(debtor, creditor, ^{IF} antecedent, ^{THEN} consequent)

STATE CHART
STATE DIAGRAM (UML)

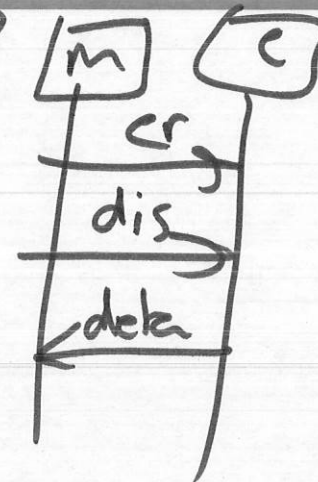
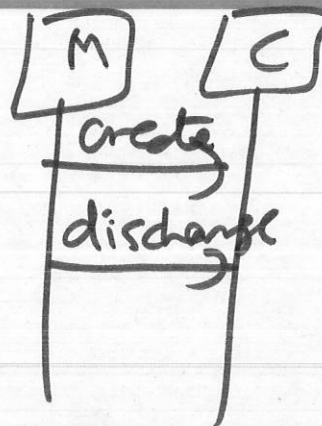
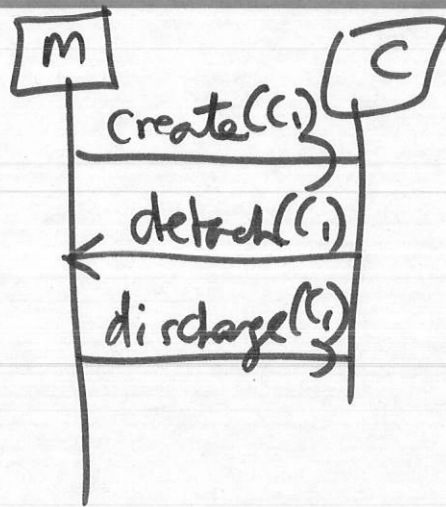
LOGICAL
STRUCTURE



(a) Commit

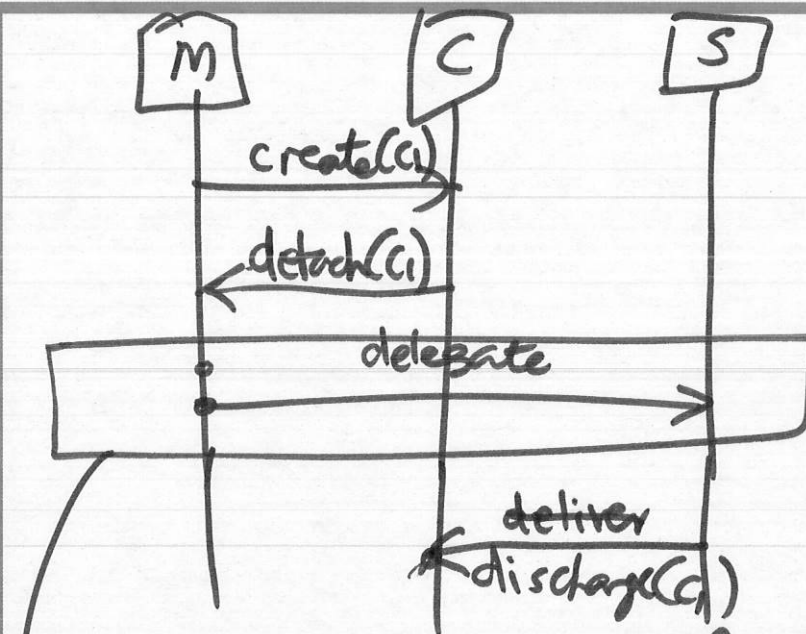
(b) Relieve

Transition = event[condition]/action

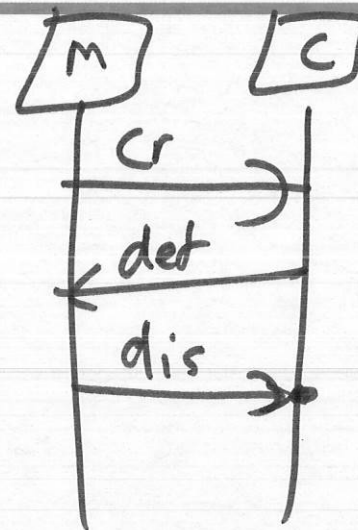


Commitment Operations

- ▶ create($C(x, y, p, q)$) establishes the commitment
- ▶ detach($C(x, y, p, q)$) turns it into a base commitment
- ▶ discharge($C(x, y, p, q)$) satisfies the commitment
- ▶ cancel($C(x, y, p, q)$) cancels the commitment
- ▶ release($C(x, y, p, q)$) releases the debtor from the commitment
- ▶ delegate($z, C(x, y, p, q)$) replaces x by z as the debtor
 - ▶ x remains ultimately responsible (in our work)
- ▶ assign($w, C(x, y, p, q)$) replaces y by w as the creditor



vs



FLEXIBILITY — IN ORDER of EVENTS
 — IN PARTICIPANTS
 SIMPLER REALIZATION

BOTH ARE CORRECT BECAUSE
 THEY END WITH NO VIOLATIONS
 (NO COMMITMENTS SATISFIED or EXPIRED or NULL)

