

Course: CSC707, Automata, Computability and Computational Theory

Reduction Homework: NP-complete problems

Submission: Use Wolfware

File Format: LaTeX and PDF

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Due Date: 2:00 A.M. (EST), Tuesday, February 9, 2010

1. Provide any feedback/questions you may have on this homework (**optional**).
 2. Using LaTeX is required.
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1. Given a NP-complete problem, Vertex Cover, show that the Independent Set is NP-complete.

Independent Set is defined as follows:

INSTANCE: A graph $G = (V, E)$ and a positive integer $k \leq |V|$.

QUESTION: Is there a subset S of k vertices in G such that no pair of vertices in S is connected by an edge in G ?

Solution:

- (a) (Verification): Show that Independent Set is in NP.

Given an independent set $C \subseteq V, |C| = n$ for a graph $G = (V, E)$, we can verify it using the following pseudo-code:

$\forall u \in C,$
 $\forall v \in C,$
check whether $\langle u, v \rangle \in E$

Checking whether $\langle u, v \rangle \in E$ can be done in $T(|E|)$. So the verification algorithm takes $O(n^2|E|)$.

- (b) (Reduction): Show that Independent Set is NP-hard.

Given a G has a VC of size k , we should construct a graph G' has Independent Set of size k' .

Construction Process: Given $\text{VertexCover}(G, k)$ where V_1 is the vertex cover and $k = |V_1|$, we set $G' = G$ and $k' = |V| - k$, then we could return the answer to $\text{IndependentSet}(G', k')$ where $V - V_1$ is the independent set. This takes constant time.

- (c) (Correctness): Show that Independent Set is NP-hard.

We need to show that G has a vertex cover of size k if and only if it has an Independent Set of size $k' = |V| - k$.

Assume $G(V, E)$ has a vertex cover C of size k . Consider two vertices $u \in V - C$ and $v \in V - C$, we can know that $e = \langle u, v \rangle \notin E$ since C is a vertex cover of G . Therefore, no two vertices in $V - C$ are connected by an edge. So $V - C$ is an independent set with size $k' = |V| - k$.

Assume G has an Independent Set S of size $k' = |V| - k$.
 $\forall e \in E, e = \langle u, v \rangle, S$ is independent set $\Rightarrow u \notin S$ or $v \notin S \Rightarrow$
 $u \in V - S$ or $v \in V - S \Rightarrow V - S$ covers $e = \langle u, v \rangle$.