Formal Languages and Automata Theory, SS 2019. Homework 9 (due Week 12)

- 1. Write context-free grammars for the following languages:
 - (a) $L = \{w | w \text{ is a binary palindrome}\}$
 - (b) $L = \{ \text{ the language of balanced paranthesis } \}$ $Examples: ()(((())))((())), (((()))(())); Counterexamples: ((((())))((())), (((())))(())) \}$
 - (c) $L = \{0^m 1^n | m \ge n\}$
 - (d) $L = \{0^n 1^n | n \ge 1\}$ Examples: 01, 00001111
 - (e) $L = \{0^n 1^n | n \ge 0\}$ Examples: λ , 01, 00001111
 - (f) $L = \{$ The set of all strings with an equal number of a's and b's $\}$ $Examples: \lambda$, aabb, bbaa, abbababa, bbabababa.
 - (g) $L = \{ \text{ Binary words with even length} \}$
 - (h) $L = \{0^i 1^j 2^k | i = j \text{ or } j = k, \text{ where } i, j, k \ge 0\}$
 - (i) $L = \{a^i b^j c^k | i + j = k, i, j, k \ge 0\}.$
 - (j) $L = \{a^i b^j c^k | \neg (i = j) \text{ or } \neg (j = k)\}$
 - (k) The language of all binary strings that are not of the form ww, that is, not equal to any repeated string.
 - (l) The language of all strings with twice as many 0's as 1's.
- 2. The following grammar generates the language of regular expression $0^01(0|1)^*$. $S \to A1B, A \to 0A|\varepsilon, B \to 0B|1B|\varepsilon$. Give the leftmost and rightmost derivations of the following strings: (a) 00101, (b) 1001, (c) 00011.
- 3. Consider the CFG G defined by productions: $S \to aS|Sb|a|b$.
 - (a) Prove by induction on the string length that no string in L(G) has ba as a substring.
 - (b) Describe L(G) informally. Justify your answer using point (a).
- 4. Consider the CFG G defined by productions: $S \to aSbS|bSaS|\varepsilon$. Prove that L(G) is the set of all strings with an equal number of a's and b's.
- 5. Consider the grammar: $S \to aS|aSbS|\varepsilon$. This grammar is ambiguous. Show in particular that the string aab has two: (a) parse trees, (b) leftmost derivations, (c) rightmost derivations.
- 6. Find an unambigous grammar for the language generated by the grammar above.