

Automated Theorem Proving (Demonstrarea Automata a Teoremelor) *Resolution*

Mădălina Eraşcu¹ and Tudor Jebelean²

¹West University of Timișoara and Institute e-Austria Timișoara, bvd. V. Parvan 4,
Timișoara, Romania,

`merascu@info.uvt.ro`

²Research Institute for Symbolic Computation, Johannes Kepler University, Linz, Austria,
`tjebelea@risc.jku.at`

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Example 1 Prove by resolution that G is a logical consequence of F_1 and F_2 where

$$\begin{aligned} F_1 : & \quad \forall_x (C[x] \Rightarrow (W[x] \wedge R[x])) \\ F_2 : & \quad \exists_x (C[x] \wedge O[x]) \\ G : & \quad \exists_x (O[x] \wedge R[x]) \end{aligned}$$

Solution. We show that $F_1 \wedge F_2 \wedge \neg G$ is unsatisfiable by resolution. We transform $F_1, F_2, \neg G$ into Skolem standard form. We have

$$\begin{aligned} F_1 : & \quad \forall_x (C[x] \Rightarrow (W[x] \wedge R[x])) \\ \iff & \quad \forall_x (\neg C[x] \vee (W[x] \wedge R[x])) \\ \iff & \quad \forall_x (\neg C[x] \vee W[x]) \wedge (\neg C[x] \vee R[x]) \end{aligned}$$

$$\begin{aligned} F_2 : & \quad \exists_x (C[x] \wedge O[x]) \\ \rightsquigarrow & \quad C[a] \wedge O[a] \end{aligned}$$

$$\begin{aligned} \neg G : & \quad \neg \left(\exists_x (O[x] \wedge R[x]) \right) \\ \iff & \quad \forall_x (\neg O[x] \vee \neg R[x]) \end{aligned}$$

We have the following set of clauses

- (1) $\neg C[x] \vee W[x]$
- (2) $\neg C[x] \vee R[x]$
- (3) $C[a]$
- (4) $O[a]$
- (5) $\neg O[x] \vee \neg R[x]$

By resolution we obtain also the following clauses

$$\begin{array}{ll} (6) & \neg R[a] \quad (4) \wedge (5), \{x \rightarrow a\} \\ (7) & \neg C[a] \quad (6) \wedge (2), \{x \rightarrow a\} \\ (8) & \emptyset \quad (7) \wedge (3) \end{array}$$

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Example 2 Prove by resolution that G is a logical consequence of F_1 and F_2 where

$$\begin{array}{l} F_1 : \quad \exists_x \left(P[x] \wedge \forall_y (D[y] \Rightarrow L[x, y]) \right) \\ F_2 : \quad \forall_x \left(P[x] \Rightarrow \forall_y (Q[y] \Rightarrow \neg L[x, y]) \right) \\ G : \quad \forall_x (D[x] \Rightarrow \neg Q[x]) \end{array}$$

Solution. We show that $F_1 \wedge F_2 \wedge \neg G$ is unsatisfiable by resolution. We transform $F_1, F_2, \neg G$ into Skolem standard form. We have

$$\begin{aligned} F_1 : & \quad \exists_x \left(P[x] \wedge \forall_y (D[y] \Rightarrow L[x, y]) \right) \\ \iff & \quad \exists_x \left(P[x] \wedge \forall_y (\neg D[y] \vee L[x, y]) \right) \\ \iff & \quad \exists_x \forall_y (P[x] \wedge (\neg D[y] \vee L[x, y])) \\ \rightsquigarrow & \quad \forall_y (P[a] \wedge (\neg D[y] \vee L[a, y])) \\ \\ F_2 : & \quad \forall_x \left(P[x] \Rightarrow \forall_y (Q[y] \Rightarrow \neg L[x, y]) \right) \\ \iff & \quad \forall_x \left(P[x] \Rightarrow \forall_y (\neg Q[y] \vee \neg L[x, y]) \right) \\ \iff & \quad \forall_x \left(\neg P[x] \vee \forall_y (\neg Q[y] \vee \neg L[x, y]) \right) \\ \iff & \quad \forall_x \forall_y (\neg P[x] \vee \neg Q[y] \vee \neg L[x, y]) \\ \\ \neg G : & \quad \neg \left(\forall_x (D[x] \Rightarrow \neg Q[x]) \right) \\ \iff & \quad \neg \left(\forall_x (\neg D[x] \vee \neg Q[x]) \right) \\ \iff & \quad \exists_x (D[x] \wedge Q[x]) \\ \rightsquigarrow & \quad D[a] \wedge Q[a] \end{aligned}$$

We have the following set of clauses

$$\begin{array}{l} (1) \quad P[a] \\ (2) \quad \neg D[y] \vee L[a, y] \\ (3) \quad \neg P[x] \vee \neg Q[y] \vee \neg L[x, y] \\ (4) \quad D[a] \\ (5) \quad Q[a] \end{array}$$

By resolution we obtain also the following clauses

$$\begin{array}{ll}
 (6) & L[a, a] \\
 (7) & \neg P[a] \vee \neg Q[a] \\
 (8) & \neg Q[a] \\
 (9) & \emptyset
 \end{array}
 \begin{array}{ll}
 (2) \wedge (4), \{y \rightarrow a\} \\
 (3) \wedge (6), \{x \rightarrow a, y \rightarrow a\} \\
 (1) \wedge (7) \\
 (5) \wedge (8)
 \end{array}$$

Example 3 Prove by resolution that G is a logical consequence of F where

$$\begin{array}{l}
 F : \quad \forall \exists (S[x, y] \wedge M[y]) \Rightarrow \exists (I[y] \wedge E[x, y]) \\
 G : \quad \neg \exists I[x] \Rightarrow \forall (S[x, y] \Rightarrow \neg M[y])
 \end{array}$$

Solution. We show that $F \wedge \neg G$ is unsatisfiable. First we transform the formulas into standard form. We have

$$\begin{aligned}
 F : & \quad \forall_x \left(\exists_y (S[x, y] \wedge M[y]) \right) \Rightarrow \exists_y (I[y] \wedge E[x, y]) \\
 \iff & \quad \forall_x \neg \left(\exists_y (S[x, y] \wedge M[y]) \right) \vee \exists_y (I[y] \wedge E[x, y]) \\
 \iff & \quad \forall_x \left(\forall_y (\neg S[x, y] \vee \neg M[y]) \right) \vee \exists_y (I[y] \wedge E[x, y]) \\
 \iff & \quad \forall_x \left(\forall_y (\neg S[x, y] \vee \neg M[y]) \right) \vee (I[f[x]] \wedge E[x, f[x]]) \\
 \iff & \quad \forall_{xy} (\neg S[x, y] \vee \neg M[y]) \vee (I[f[x]] \wedge E[x, f[x]]) \\
 \iff & \quad \forall_{xy} ((\neg S[x, y] \vee \neg M[y] \vee I[f[x]]) \wedge (\neg S[x, y] \vee \neg M[y] \vee E[x, f[x]])) \\
 \neg G : & \quad \neg \left(\neg \exists_x I[x] \Rightarrow \forall_{x,y} (S[x, y] \Rightarrow \neg M[y]) \right) \\
 \iff & \quad \neg \left(\neg \exists_x I[x] \Rightarrow \forall_{x,y} (\neg S[x, y] \vee \neg M[y]) \right) \\
 \iff & \quad \neg \left(\exists_x I[x] \vee \forall_{x,y} (\neg S[x, y] \vee \neg M[y]) \right) \\
 \iff & \quad \left(\forall_x \neg I[x] \wedge \exists_{x,y} (S[x, y] \wedge M[y]) \right) \\
 \iff & \quad \forall_z \neg I[z] \wedge \exists_{x,y} (S[x, y] \wedge M[y]) \\
 \rightsquigarrow & \quad \forall_z \neg I[z] \wedge S[a, b] \wedge M[b]
 \end{aligned}$$

We have the following set of clauses

$$\begin{array}{l}
 (1) \quad \neg S[x, y] \vee \neg M[y] \vee I[f[x]] \\
 (2) \quad \neg S[x, y] \vee \neg M[y] \vee E[x, f[x]] \\
 (3) \quad \neg I[z] \\
 (4) \quad S[a, b] \\
 (5) \quad M[b]
 \end{array}$$

By resolution we obtain also the following clauses

$$\begin{array}{ll}
 (6) & \neg S[x, y] \vee \neg M[y] \quad (1) \wedge (3), \{z \rightarrow f[x]\} \\
 (7) & \neg M[b] \quad (4) \wedge (6), \{x \rightarrow a, y \rightarrow b\} \\
 (8) & \emptyset \quad (5) \wedge (7)
 \end{array}$$

Example 4 Prove by resolution that G is a logical consequence of F_1, F_2 , and F_3 where

$$\begin{aligned} F_1 : & \quad \forall_x (Q[x] \Rightarrow \neg P[x]) \\ F_2 : & \quad \forall_x \left((R[x] \wedge \neg Q[x]) \Rightarrow \exists_y (T[x, y] \wedge S[y]) \right) \\ F_3 : & \quad \exists_x \left(P[x] \wedge \forall_y (T[x, y] \Rightarrow P[y]) \wedge R[x] \right) \\ G : & \quad \exists_x (S[x] \wedge P[x]) \end{aligned}$$

Solution. We show that $F_1 \wedge F_2 \wedge F_3 \wedge \neg G$ is unsatisfiable. First we transform the formulas into standard form.

$$\begin{aligned} F_1 : & \quad \forall_x (Q[x] \Rightarrow \neg P[x]) \iff \forall_x (\neg Q[x] \vee \neg P[x]) \\ F_2 : & \quad \forall_x \left((R[x] \wedge \neg Q[x]) \Rightarrow \exists_y (T[x, y] \wedge S[y]) \right) \\ & \iff \forall_x \left(\neg (R[x] \wedge \neg Q[x]) \vee \exists_y (T[x, y] \wedge S[y]) \right) \\ & \iff \forall_x \left(\neg R[x] \vee Q[x] \vee \exists_y (T[x, y] \wedge S[y]) \right) \\ & \iff \forall_{xy} (\neg R[x] \vee Q[x] \vee (T[x, y] \wedge S[y])) \\ & \iff \forall_{xy} ((\neg R[x] \vee Q[x] \vee T[x, y]) \wedge (\neg R[x] \vee Q[x] \vee S[y])) \\ & \rightsquigarrow \forall_x ((\neg R[x] \vee Q[x] \vee T[x, f[x]]) \wedge (\neg R[x] \vee Q[x] \vee S[f[x]])) \\ F_3 : & \quad \exists_x \left(P[x] \wedge \forall_y (T[x, y] \Rightarrow P[y]) \wedge R[x] \right) \\ & \iff \exists_x \left(P[x] \wedge \forall_y (\neg T[x, y] \vee P[y]) \wedge R[x] \right) \\ & \iff \exists_{xy} (P[x] \wedge (\neg T[x, y] \vee P[y]) \wedge R[x]) \\ & \rightsquigarrow \forall_y (P[a] \wedge (\neg T[a, y] \vee P[y]) \wedge R[a]) \\ \neg G : & \quad \neg \left(\exists_x (S[x] \wedge P[x]) \right) \\ & \iff \forall_x (\neg S[x] \vee \neg P[x]) \end{aligned}$$

We have the following set of clauses

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|------|---------------------------------------|--|
| (1) | $\neg Q[x] \vee \neg P[x]$ | |
| (2) | $\neg R[x] \vee Q[x] \vee T[x, f[x]]$ | |
| (3) | $\neg R[x] \vee Q[x] \vee S[f[x]]$ | |
| (4) | $P[a]$ | |
| (5) | $\neg T[a, y] \vee P[y]$ | |
| (6) | $R[a]$ | |
| (7) | $\neg S[x] \vee \neg P[x]$ | |
| (8) | $\neg Q[a]$ | $(1) \wedge (4), \{x \rightarrow a\}$ |
| (9) | $\neg R[a] \vee T[a, f[a]]$ | $(8) \wedge (2), \{x \rightarrow a\}$ |
| (10) | $\neg R[a] \vee P[f[a]]$ | $(9) \wedge (5), \{y \rightarrow f[a]\}$ |
| (11) | $P[f[a]]$ | $(10) \wedge (6)$ |
| (12) | $\neg S[f[a]]$ | $(11) \wedge (7)$ |
| (13) | $\neg R[a] \vee Q[a]$ | $(12) \wedge (3)$ |
| (14) | $Q[a]$ | $(13) \wedge (6)$ |
| (15) | \emptyset | $(14) \wedge (8)$ |

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