## Resolution Principle. Examples

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**Example 1** Prove by resolution that G is a logical consequence of  $F_1$  and  $F_2$  where

$$F_1: \quad \forall (C[x] \Rightarrow (W[x] \land R[x]))$$

$$F_2: \quad \exists (C[x] \land O[x])$$

$$G: \quad \exists (O[x] \land R[x])$$

**Solution.** We show that  $F_1 \wedge F_2 \wedge \neg G$  is unsatisfiable by resolution. We transform  $F_1, F_2, \neg G$  into Skolem standard form. We have

$$F_{1}: \ \forall (C[x] \Rightarrow (W[x] \land R[x]))$$

$$\iff \forall (\neg C[x] \lor (W[x] \land R[x]))$$

$$\iff \forall (\neg C[x] \lor W[x]) \land (\neg C[x] \lor R[x])$$

$$F_{2}: \ \exists (C[x] \land O[x])$$

$$\rightsquigarrow C[a] \land O[a]$$

$$\neg G: \ \neg \left(\exists (O[x] \land R[x])\right)$$

$$\iff \forall (\neg O[x] \lor \neg R[x])$$

We have the following set of clauses

$$\begin{array}{llll} (1) & \neg C[x] \lor W[x] \\ (2) & \neg C[x] \lor R[x] \\ (3) & C[a] \\ (4) & O[a] \\ \end{array}$$

 $\neg O[x] \vee \neg R[x]$ 

By resolution we obtain also the following clauses

$$\begin{array}{ll} (6) & \neg R[a] & (4) \wedge (5), \{x \rightarrow a\} \\ (7) & \neg C[a] & (6) \wedge (2), \{x \rightarrow a\} \end{array}$$

 $(7) \wedge (3)$ 

**Example 2** Prove by resolution that G is a logical consequence of  $F_1$  and  $F_2$  where

$$F_{1}: \quad \exists \left(P[x] \land \forall D[y] \Rightarrow L[x,y]\right)$$

$$F_{2}: \quad \forall \left(P[x] \Rightarrow \forall D[y] \Rightarrow \neg L[x,y]\right)$$

$$G: \quad \forall D[x] \Rightarrow \neg D[x]$$

**Solution.** We show that  $F_1 \wedge F_2 \wedge \neg G$  is unsatisfiable by resolution. We transform  $F_1, F_2, \neg G$  into Skolem standard form. We have

We have the following set of clauses

$$\begin{array}{lll} (1) & P[a] \\ (2) & \neg D[y] \ \lor \ L[a,y] \\ (3) & \neg P[x] \ \lor \ \neg Q[y] \ \lor \ \neg L[x,y] \\ (4) & D[a] \\ (5) & Q[a] \end{array}$$

By resolution we obtain also the following clauses

$$\begin{array}{lll} (6) & L[a,a] & (2) \wedge (4), \{y \to a\} \\ (7) & \neg P[a] \vee \neg Q[a] & (3) \wedge (6), \{x \to a, y \to a\} \\ (8) & \neg Q[a] & (1) \wedge (7) \\ (9) & \emptyset & (5) \wedge (8) \end{array}$$

**Example 3** Prove by resolution that G is a logical consequence of F where

$$\begin{array}{lll} F: & \forall \exists \left(S[x,y] \ \land \ M[y]\right) \ \Rightarrow \ \exists \left(I[y] \ \land \ E[x,y]\right) \\ G: & \neg \exists I[x] \ \Rightarrow \ \forall \limits_{x,y} \left(S[x,y] \Rightarrow \neg M[y]\right) \end{array}$$

**Solution.** We show that  $F \wedge \neg G$  is unsatisfiable. First we transform the formulas into standard form. We

We have the following set of clauses

$$\begin{array}{llll} (1) & \neg S[x,y] \ \lor \ \neg M[y] \ \lor \ I[f[x]] \\ (2) & \neg S[x,y] \ \lor \ \neg M[y] \ \lor \ E[x,f[x]] \end{array}$$

$$(2) \quad \neg S[x,y] \ \lor \ \neg M[y] \ \lor \ E[x,f[x]]$$

$$\begin{array}{ccc} (3) & \neg I[z] \\ (4) & S[a,b] \end{array}$$

$$(4)$$
  $S[a,b]$ 

M[b]

By resolution we obtain also the following clauses

$$(6) \quad \neg S|x,y| \lor \neg M|y| \quad (1) \land (3), \{z \to f|x|\}$$

(6) 
$$\neg S[x,y] \lor \neg M[y]$$
 (1)  $\land$  (3),  $\{z \to f[x]\}$   
(7)  $\neg M[b]$  (4)  $\land$  (6),  $\{x \to a, y \to b\}$   
(8)  $\emptyset$  (5)  $\land$  (7)

$$\emptyset$$
  $(5) \wedge (7)$ 

**Example 4** Prove by resolution that G is a logical consequence of  $F_1, F_2$ , and  $F_3$  where

$$F_{1}: \quad \forall (Q[x] \Rightarrow \neg P[x])$$

$$F_{2}: \quad \forall \left( (R[x] \land \neg Q[x]) \Rightarrow \exists (T[x,y] \land S[y]) \right)$$

$$F_{3}: \quad \exists \left( P[x] \land \forall (T[x,y] \Rightarrow P[y]) \land R[x] \right)$$

$$G: \quad \exists (S[x] \land P[x])$$

**Solution.** We show that  $F_1 \wedge F_2 \wedge F_3 \wedge \neg G$  is unsatisfiable. First we transform the formulas into standard form.

$$F_{1}: \begin{tabular}{l} \forall (Q[x] \Rightarrow \neg P[x]) & \Longleftrightarrow & \forall (\neg Q[x] \vee \neg P[x]) \\ F_{2}: \begin{tabular}{l} \forall (R[x] \wedge \neg Q[x]) & \Rightarrow & \exists (T[x,y] \wedge S[y]) \\ & \Longleftrightarrow & \forall \left(\neg (R[x] \wedge \neg Q[x]) \vee & \exists (T[x,y] \wedge S[y]) \right) \\ & \Longleftrightarrow & \forall \left(\neg R[x] \vee Q[x] \vee & \exists (T[x,y] \wedge S[y]) \right) \\ & \Longleftrightarrow & \forall \exists (\neg R[x] \vee Q[x] \vee & (T[x,y] \wedge S[y])) \\ & \Longleftrightarrow & \forall \exists ((\neg R[x] \vee Q[x] \vee & T[x,y]) \wedge & (\neg R[x] \vee Q[x] \vee S[y])) \\ & \longleftrightarrow & \forall ((\neg R[x] \vee Q[x] \vee & T[x,f[x]]) \wedge & (\neg R[x] \vee Q[x] \vee S[f[x]])) \\ & \Longleftrightarrow & \forall \left((\neg R[x] \vee Q[x] \vee & T[x,f[x]]) \wedge & (\neg R[x] \vee Q[x] \vee S[f[x]])) \\ & \longleftrightarrow & \exists \left(P[x] \wedge \forall (T[x,y] \Rightarrow P[y]) \wedge & R[x] \right) \\ & \Longleftrightarrow & \exists \varphi (P[x] \wedge (\neg T[x,y] \vee P[y]) \wedge & R[x]) \\ & \Longleftrightarrow & \exists \varphi (P[x] \wedge (\neg T[x,y] \vee P[y]) \wedge & R[x]) \\ & \hookrightarrow & \forall \varphi (P[x] \wedge (\neg T[x,y] \vee P[y]) \wedge & R[x]) \\ & \hookrightarrow & \forall \varphi (P[x] \wedge (\neg T[x,y] \vee P[y]) \wedge & R[x]) \\ & \Longleftrightarrow & \forall \varphi (\neg S[x] \vee \neg P[x]) \\ & \Longleftrightarrow & \forall \varphi (\neg S[x] \vee \neg P[x]) \\ \end{tabular}$$

We have the following set of clauses

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 \begin{array}{l} \neg Q[x] \ \lor \ \neg P[x] \\ \neg R[x] \ \lor \ Q[x] \ \lor \ T[x,f[x]] \\ \neg R[x] \ \lor \ Q[x] \ \lor \ S[f[x]] \end{array} 
(2)
(3)
(4)
           P[a]
           \neg T[a,y] \lor P[y]
(5)
(6)
           R[a]
           \neg S[x] \vee \neg P[x]
(7)
           \neg Q[a]
(8)
                                                                  (1) \land (4), \{x \rightarrow a\}
           \neg R[a] \ \lor \ T[a, f[a]]
                                                                  (8) \land (2), \{x \rightarrow a\}
(9)
(10)
           \neg R[a] \lor P[f[a]]
                                                                  (9) \land (5), \{y \to f[a]\}
(11)
           P[f[a]]
                                                                  (10) \wedge (6)
           \neg S[f[a]]
                                                                  (11) \wedge (7)
(12)
                                                                  (12) \wedge (3)
(13)
           \neg R[a] \lor Q[a]
(14)
           Q[a]
                                                                  (13) \wedge (6)
(15)
                                                                  (14) \land (8)
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