Formal Methods in Software Development SAT Solving

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Based on slides of the lecture Satisfiability Checking (Erika Ábrahám), RTWH Aachen

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Given:

• Propositional logic formula φ in CNF.

Question:

■ Is φ satisfiable?

(Is there a model for φ ?)

SAT-solving: Components

- Decision (enumeration)
- Boolean constraint propagation (BCP) discussed in the previous lecture
- Conflict resolution and backtracking discussed in the previous lecture

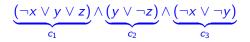
SAT-solving: Components

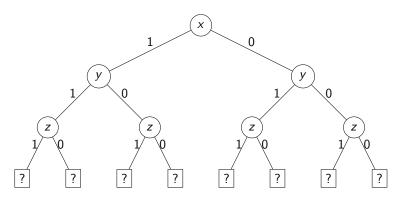
- Decision (enumeration)
- Boolean constraint propagation (BCP)
- Conflict resolution and backtracking

Enumeration algorithm

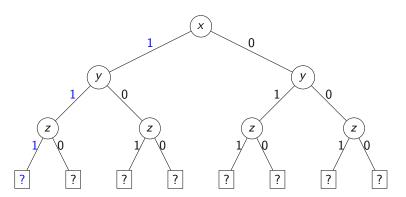
- Naive approach yields 2ⁿ candidate models to check
- Solution: decision heuristics

$$\underbrace{\left(\neg x \lor y \lor z\right)}_{c_1} \land \underbrace{\left(y \lor \neg z\right)}_{c_2} \land \underbrace{\left(\neg x \lor \neg y\right)}_{c_3}$$

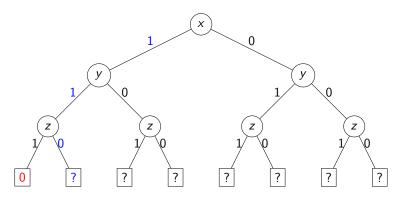




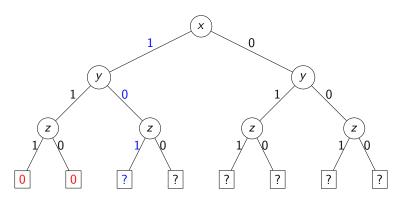




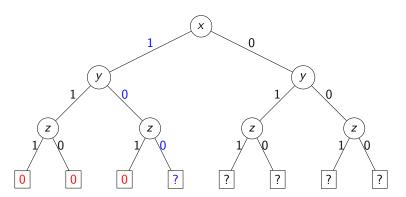




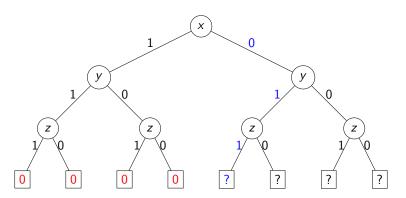


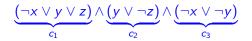


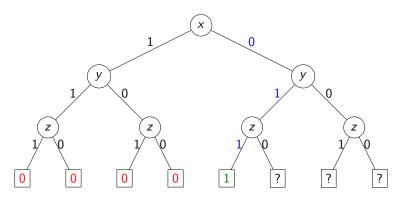




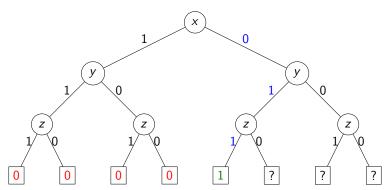




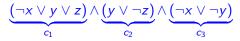




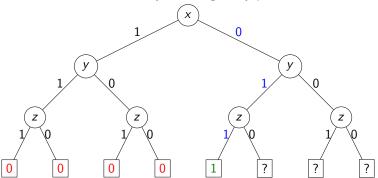




For unsatisfiable problems, all assignments need to be checked. For satisfiable problems, variable and sign ordering might strongly influence the running time.



Static variable order x < y < z, sign: try positive first

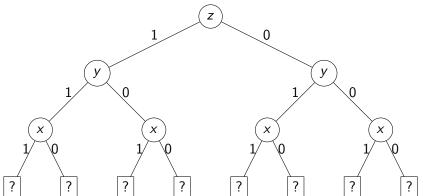


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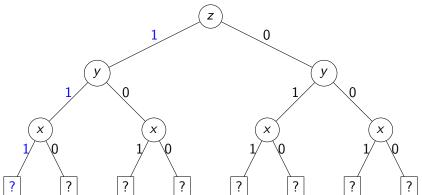
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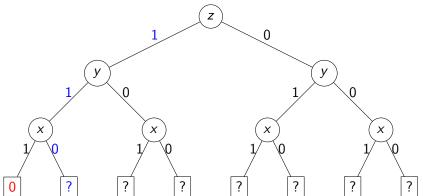




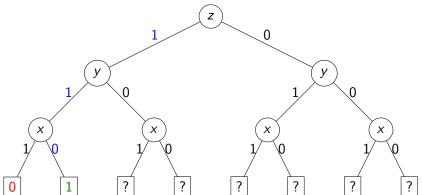
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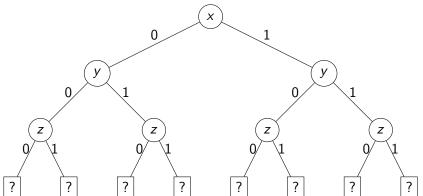
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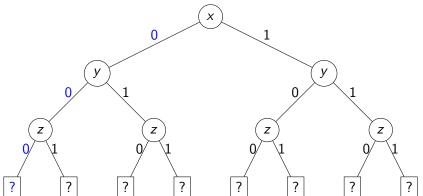
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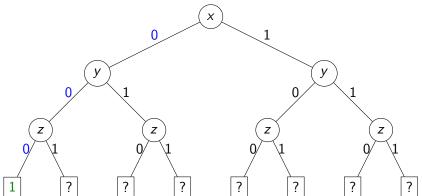
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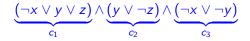


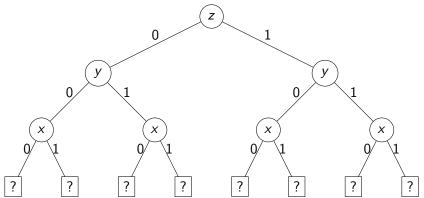
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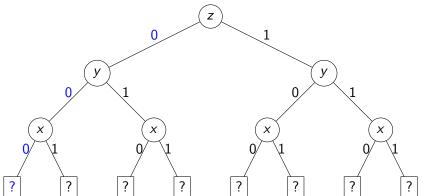
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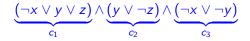
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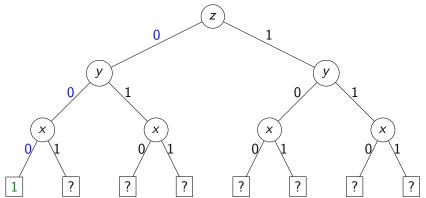












Decision heuristics

Dynamic Largest Individual Sum (DLIS): Choose an assignment that increases the most the number of satisfied clauses

- For each variable x, let C_x be the number of unresolved clauses in which x appears positively.
- For each variable x, let $C_{\neg x}$ be the number unresolved clauses in which x appears negatively.
- Let x be a variable for which C_x is maximal ($C_x \ge C_z$ for all variables z).
- Let y be a variable for which $C_{\neg y}$ is maximal $(C_{\neg y} \ge C_{\neg z})$ for all variables z).
- If $C_x > C_{\neg y}$ choose x and assign it TRUE.
- Otherwise choose *y* and assign it FALSE.
- Requires $\mathcal{O}(\#literals)$ queries for each decision.

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$$\underbrace{\left(\neg x \vee y \vee z\right)}_{c_1} \wedge \underbrace{\left(y \vee \neg z\right)}_{c_2} \wedge \underbrace{\left(\neg x \vee \neg y\right)}_{c_3} \qquad \begin{array}{c} C_x = 0 & C_y = 2 & C_z = 1 \\ C_{\neg x} = 2 & C_{\neg y} = 1 & C_{\neg z} = 1 \end{array}$$

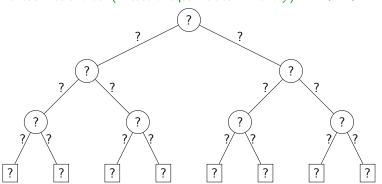
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Dynamic Largest Individual Sum (DLIS) literal order

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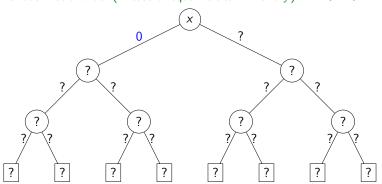
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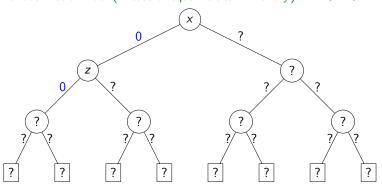
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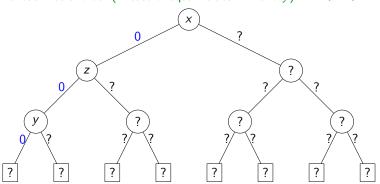
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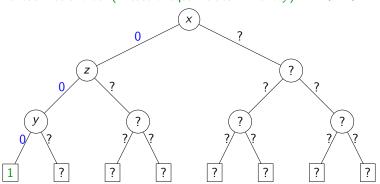
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Dynamic Largest Individual Sum (DLIS) literal order



Decision heuristics

Jersolow-Wang method

Compute for every literal / the following static value:

$$J(I): \sum_{I \in c, c \in \phi} 2^{-|c|}$$

c – clause, ϕ – formula

- Choose a literal I that maximizes J(I).
- This gives an exponentially higher weight to literals in shorter clauses

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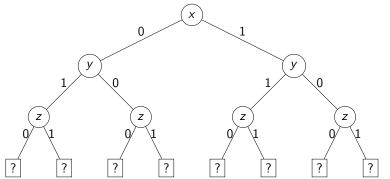
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$$J(x) = 0$$
, $J(\neg x) = \frac{1}{8} + \frac{1}{4}$, $J(y) = \frac{1}{8} + \frac{1}{4}$, $J(\neg y) = \frac{1}{4}$, $J(z) = \frac{1}{8}$, $J(\neg z) = \frac{1}{4}$

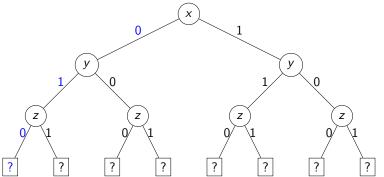
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