# Formal Methods in Software Developement SAT Solving

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Based on slides of the lecture Satisfiability Checking (Erika Ábrahám), RTWH Aachen

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#### Given:

■ Propositional logic formula  $\varphi$  in CNF.

#### Question:

■ Is  $\varphi$  satisfiable?

(Is there a model for  $\varphi$ ?)

#### SAT-solving: Components

- Decision (enumeration)
- Boolean constraint propagation (BCP) discussed in the previous lecture
- Conflict resolution and backtracking discussed in the previous lecture

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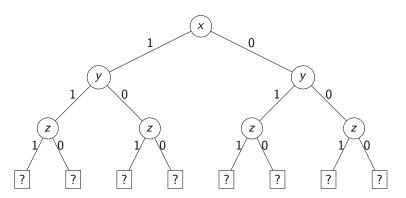
- Decision (enumeration)
- Boolean constraint propagation (BCP)
- Conflict resolution and backtracking

#### Enumeration algorithm

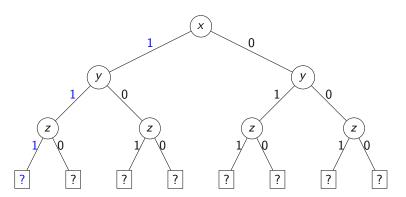
- Naive approach yields 2<sup>n</sup> candidate models to check
- Solution: decision heuristics

$$\underbrace{\left(\neg x \lor y \lor z\right)}_{c_1} \land \underbrace{\left(y \lor \neg z\right)}_{c_2} \land \underbrace{\left(\neg x \lor \neg y\right)}_{c_3}$$

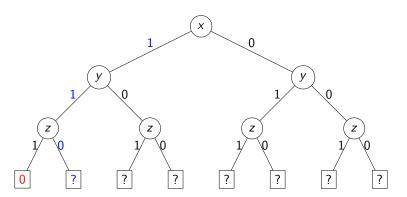




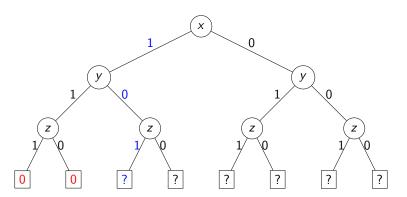




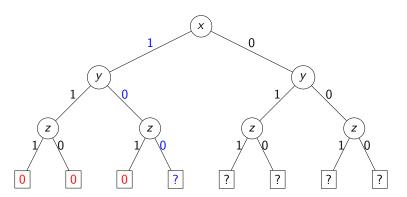




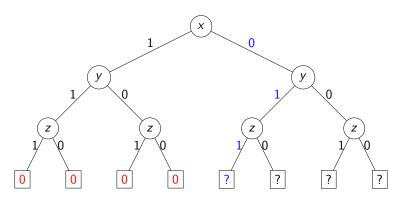




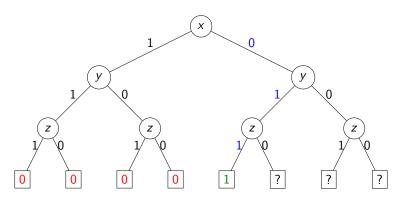




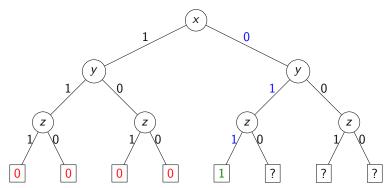




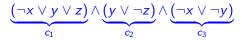




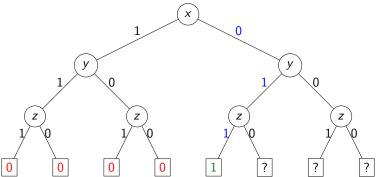




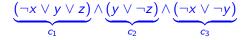
For unsatisfiable problems, all assignments need to be checked. For satisfiable problems, variable and sign ordering might strongly influence the running time.



Static variable order x < y < z, sign: try positive first

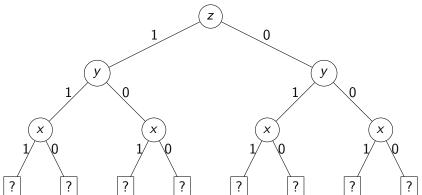


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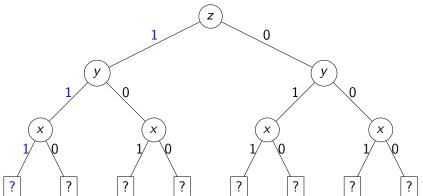


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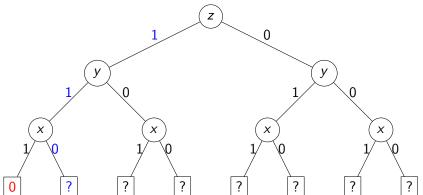




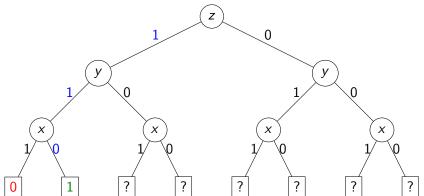
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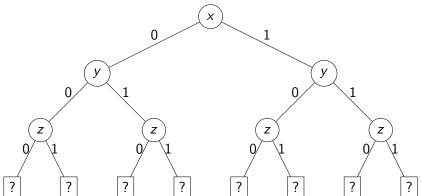




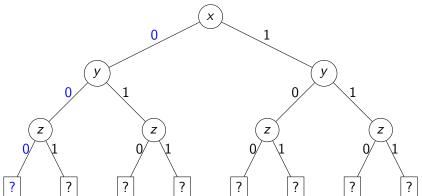
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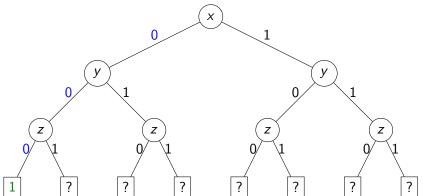
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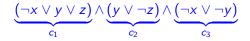


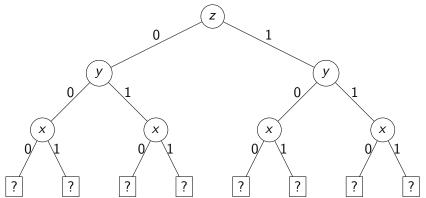
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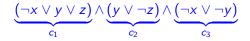


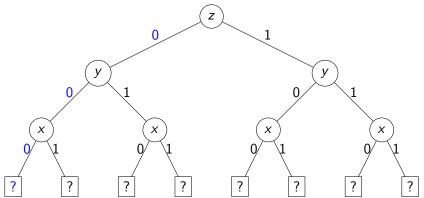
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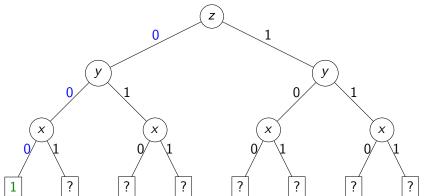












#### Decision heuristics

Dynamic Largest Individual Sum (DLIS): Choose an assignment that increases the most the number of satisfied clauses

- For each variable x, let  $C_x$  be the number of unresolved clauses in which x appears positively.
- For each variable x, let  $C_{\neg x}$  be the number unresolved clauses in which x appears negatively.
- Let x be a variable for which  $C_x$  is maximal ( $C_x \ge C_z$  for all variables z).
- Let y be a variable for which  $C_{\neg y}$  is maximal  $(C_{\neg y} \ge C_{\neg z})$  for all variables z).
- If  $C_x > C_{\neg y}$  choose x and assign it TRUE.
- Otherwise choose y and assign it FALSE.
- Requires  $\mathcal{O}(\#literals)$  queries for each decision.

$$\underbrace{(\neg x \lor y \lor z)}_{c_1} \land \underbrace{(y \lor \neg z)}_{c_2} \land \underbrace{(\neg x \lor \neg y)}_{c_3}$$

$$\underbrace{\left( \neg x \vee y \vee z \right)}_{c_1} \wedge \underbrace{\left( y \vee \neg z \right)}_{c_2} \wedge \underbrace{\left( \neg x \vee \neg y \right)}_{c_3} \qquad \begin{array}{c} C_x = 0 & C_y = 2 & C_z = 1 \\ C_{\neg x} = 2 & C_{\neg y} = 1 & C_{\neg z} = 1 \end{array}$$

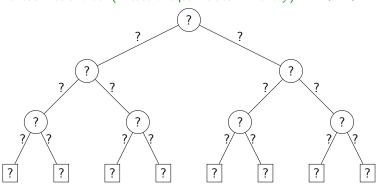
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  $C_y = 2$   $C_z = 1$   
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Dynamic Largest Individual Sum (DLIS) literal order

$$\underbrace{\left(\neg x \vee y \vee z\right)}_{c_1} \wedge \underbrace{\left(y \vee \neg z\right)}_{c_2} \wedge \underbrace{\left(\neg x \vee \neg y\right)}_{c_3} \qquad \begin{aligned} C_x &= 0 & C_y &= 2 & C_z &= 1 \\ C_{\neg x} &= 2 & C_{\neg y} &= 1 & C_{\neg z} &= 1 \end{aligned}$$

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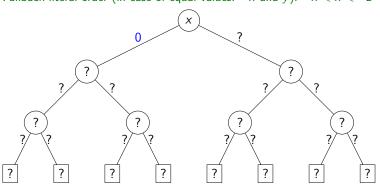
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$$C_x = 0$$
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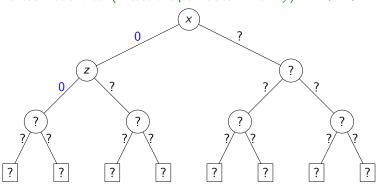
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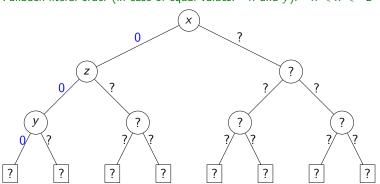
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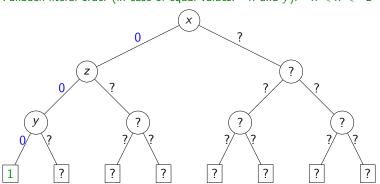
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Dynamic Largest Individual Sum (DLIS) literal order



#### Decision heuristics

#### Jersolow-Wang method

Compute for every literal / the following static value:

$$J(I): \sum_{I \in c, c \in \phi} 2^{-|c|}$$

$$c$$
 – clause,  $\phi$  – formula

- Choose a literal I that maximizes J(I).
- This gives an exponentially higher weight to literals in shorter clauses

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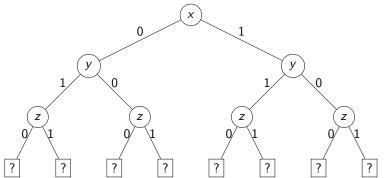
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$$J(x) = 0$$
,  $J(\neg x) = \frac{1}{8} + \frac{1}{4}$ ,  $J(y) = \frac{1}{8} + \frac{1}{4}$ ,  $J(\neg y) = \frac{1}{4}$ ,  $J(z) = \frac{1}{8}$ ,  $J(\neg z) = \frac{1}{4}$ 

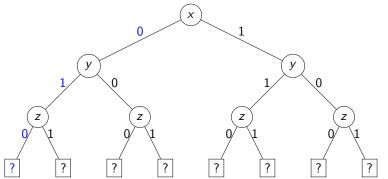
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