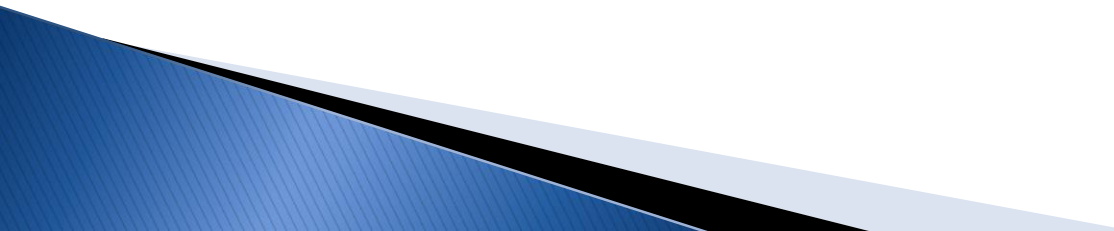


# Internet Routing

Internetworking – Lecture 6

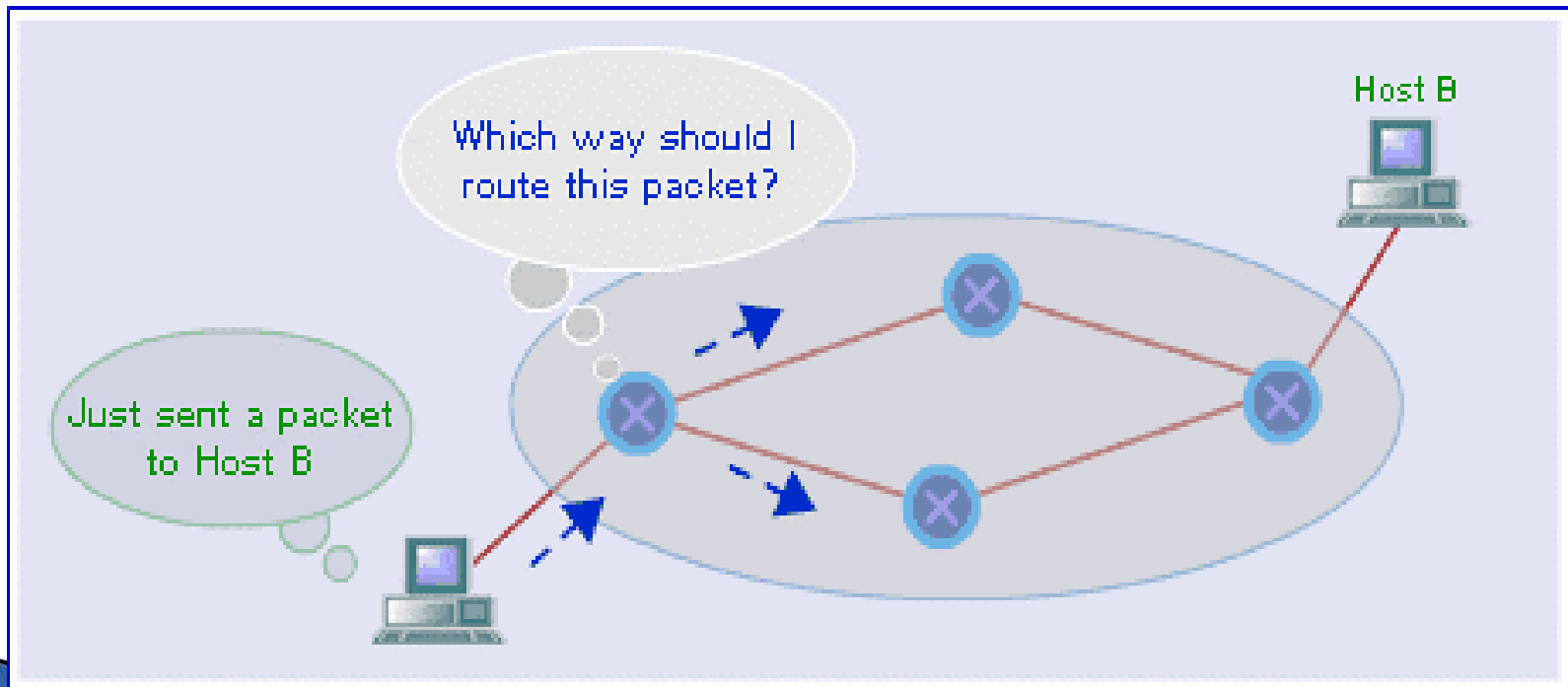
# Internet Routing

## ► Content:

- What is Routing
  - What if's – the Complications of Routing
  - Routing Tables
  - An automated routing solution
  - Defining a routing Protocol
  - Design Considerations
  - Metrics of Routing Protocols
  - Categorising Routing protocols
- 

# What is Routing?

- ▶ The act of forwarding network packets from a source network to a destination network



# What If's', Complications of Routing?

- ▶ When should you route a packet?
- ▶ What is the best route to take?
  - How do you know?
- ▶ Does the topology change?
  - If it does change what then?
- ▶ What if there is a fault in the network?
- ▶ What if the destination does not exist?
- ▶ If a packet has a different network destination network to the host then it is determined that it needs to be forwarded to another network via the IP Address.
- ▶ It is the router that holds the internetwork logic to decide how to forward the packet.

# Routing a Packet to a Destination ?

- ▶ Workstation **A**, sends an email to Workstation **B**  
↓
- ▶ Workstation **A** determines if Workstation **B** is on the same network by checking the local routing table  
↓
- ▶ Determines that Workstation **B** is on a different network  
↓
- ▶ Send packet to the default Gateway

# Static Routing

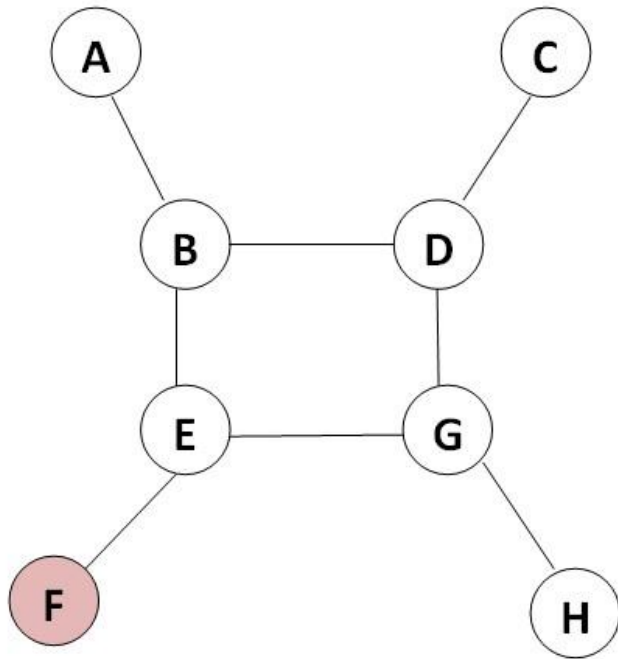
- ▶ Manually populated routing tables
- ▶ An almost impossible task to maintain in modern networks
- ▶ Ideal for small stable networks without redundant network links
- ▶ Dynamically routing protocols use up network resources learning where all the nodes are.
- ▶ Often static routing is coupled with dynamically routing
- ▶ In CISCO static routes can be configured with IP route commands

# Anatomy of a Routing Table

Code	Network, Mask	AD/Metric	Next Hop	Interface
O	10.0.0.0/8	110/20	200.1.1.1	S0
O	172.16.0.0/16	100/15	200.1.1.1	S0
O	192.168.1.0/24	100/20	200.2.2.2	S1
C	210.1.1.4/30	0/0	Directly connected	E0

- ▶ **Code:** what process discovered the route
- ▶ **Network Mask:** address of destination network and its subnet mask
- ▶ **Administrative Distance/Metric:** used to select the best route
- ▶ **Next Hop:** IP address of the next hop router
- ▶ **Interface:** the interface that the packet will be forwarded on

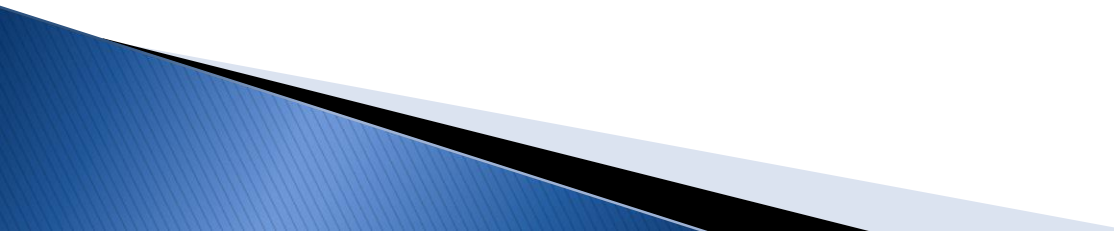
# A simple scenario...



- ▶ Send a packet from node F to node C
- ▶ Which one is the optimal path?
  - What if a link fails?
  - What if a router fails?
- ▶ Can you make an informed decision?



# An Automated Routing Solution

- ▶ We established that static routing is unworkable over the Internet;
  - ▶ The more complex the networks the harder it would be to manually manipulate the routing tables;
  - ▶ Therefore an automated approach to the problem is required
- 

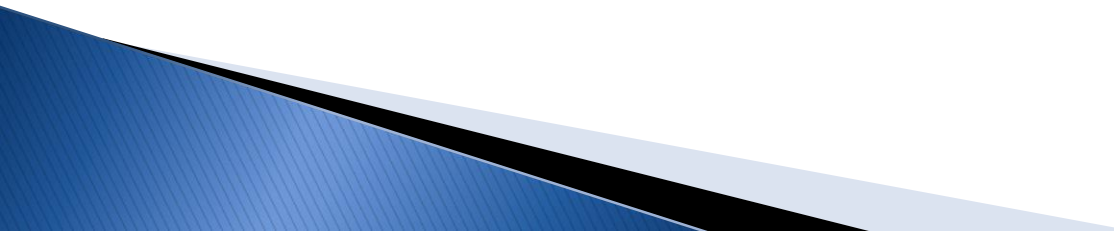
# Key Features of Dynamic Routing

- ▶ Key features of dynamic routing are:
  - They learn about the network
  - Automatically modify the routing tables
  - Dynamic routing should be deployed on any sized network

# What is a Routing Protocol?

- ▶ A set of rules that allow 2 or more routers to exchange information about the networks they are connected to.
- ▶ It is based on an algorithm to solve the communication problem
- ▶ Therefore... it is **a process that runs on the router**
- ▶ Algorithms used are based on graph theory
  - e.g. the router is the dot and they link the networks

# Historical Links

- ▶ Early protocols are based on work by R.Bellman; L.R.Ford; Edsger Dijkstra
  - ▶ Bellman–Ford – Distance Vector algorithms
  - ▶ Dijkstra – Shortest Path first algorithm
  - ▶ **No one protocol has solved all the routing problems to-date!**
- 

# Routing Protocols

## Design Considerations

- ▶ What networking issues need to be taken into consideration?
- ▶ How does the router collate the network data to populate the routing table?
  - The router needs to be able to communicate with others
  - It needs to pass its own knowledge of networks to another router
  - It must be able to receive this data
  - A common language of communication is required
  - Needs to communicate with routers, identify its status and its known routes

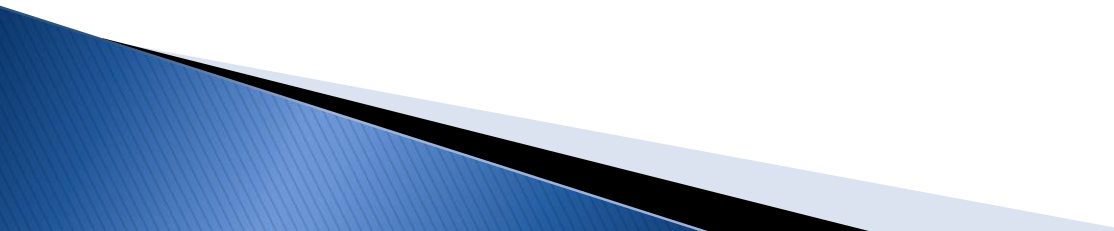
# Routing communication

- ▶ The language and vocabulary used is unique to a particular protocol
- ▶ Communication can only be between routers using the same protocol
- ▶ Routers supporting different protocols can't communicate between each other

# Routing Paths – Convergence

- ▶ If a change to the network occurs, it means the routing table needs updating. The time it takes until this happens is called “CONVERGENCE”
- ▶ What changes triggers an update?
  - If one or more network links fail, every other router needs to be informed
  - If a router crashes it can have a serious impact on the network.

# Characteristics of a Routing Protocol

- ▶ A routing Protocol must incorporate:
    - –Robustness
    - –Optimisation
    - –Flexibility
    - –Speed of convergence
    - –Avoidance of routing loops
    - –Support for classless addressing
    - –Simplicity
- 



# Metric of Routing Protocols

- ▶ How a routing protocol decides which route is best especially if more than 1 route is discovered
  - Each route is assigned a metric value
  - There are numerous factors that the protocols may take into consideration when assigning a metric value

# Metric of Routing Protocols

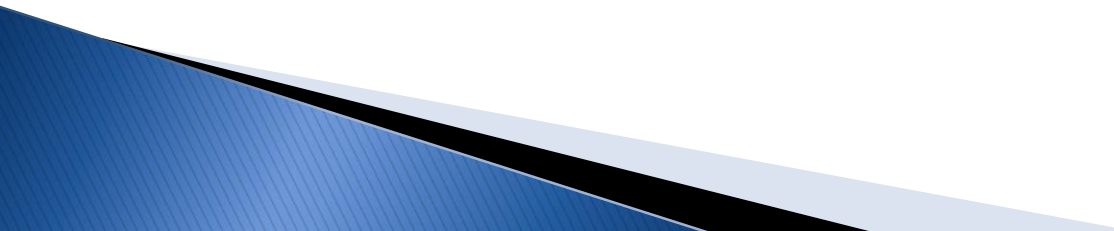
- ▶ **Hop Count:**
  - Number of routers to traverse in order to reach the destination
- ▶ **Path Length:**
  - A refinement of the hop count, Sum of per-link costs
- ▶ **Bandwidth:**
  - Speed of the link between routers
- ▶ **Delay:**
  - Time in milliseconds to cross a link
- ▶ **Load:**
  - Congestion on link due to traffic
- ▶ **Reliability:**
  - Based on bit error rates of path
- ▶ Not all routing protocols use all the variables

# Categorising Dynamic Routing Protocols

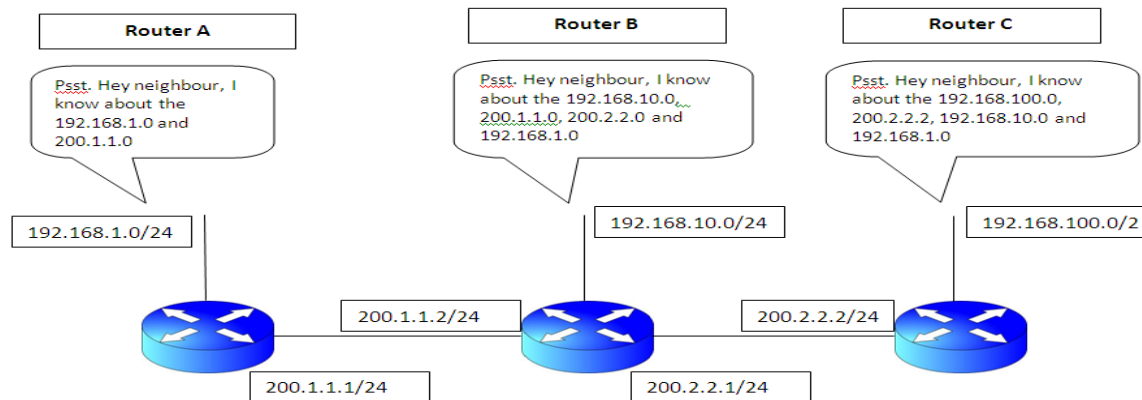
- ▶ Routing protocols are categorised by its designed purpose
- ▶ **Interior Gateway Protocols:**
  - Developed to facilitate routing **within** autonomous systems
- ▶ **Exterior Gateway Protocols:**
  - Developed to facilitate routing **between** autonomous systems
- ▶ Most protocols are interior protocols
  - (Autonomous Systems is a system under a single administration control – e.g the university network)

# Distance Vector

## “Routing by Rumour”

- ▶ Routing information is received from immediate router neighbours only
  - ▶ Sent as routing update packets **via broadcasting**
  - ▶ These updates are then added to the router tables
  - ▶ Then pass this information to their own neighbouring routers
  - ▶ Finally, all routers learn the path to all networks – **network is converged**
- 

# View of Distance Vector Protocol Communication



- ▶ Each router informs its neighbour of its directly connected network
- ▶ Includes networks the router has learned from other neighbours
- ▶ Share the metrics of the routes it knows
- ▶ In distance vector protocols the metrics is distance initially hop count – e.g. how many routers the packet has to cross to reach its destination, 3 hops = the metric 3

# Distance Vector Protocol

- ▶ If two or more paths are discovered for the same destination, the route with the **lowest hop count** would win and be added to the routing table
- ▶ Others have used bandwidth and delay to determine the metric value

# Vectors

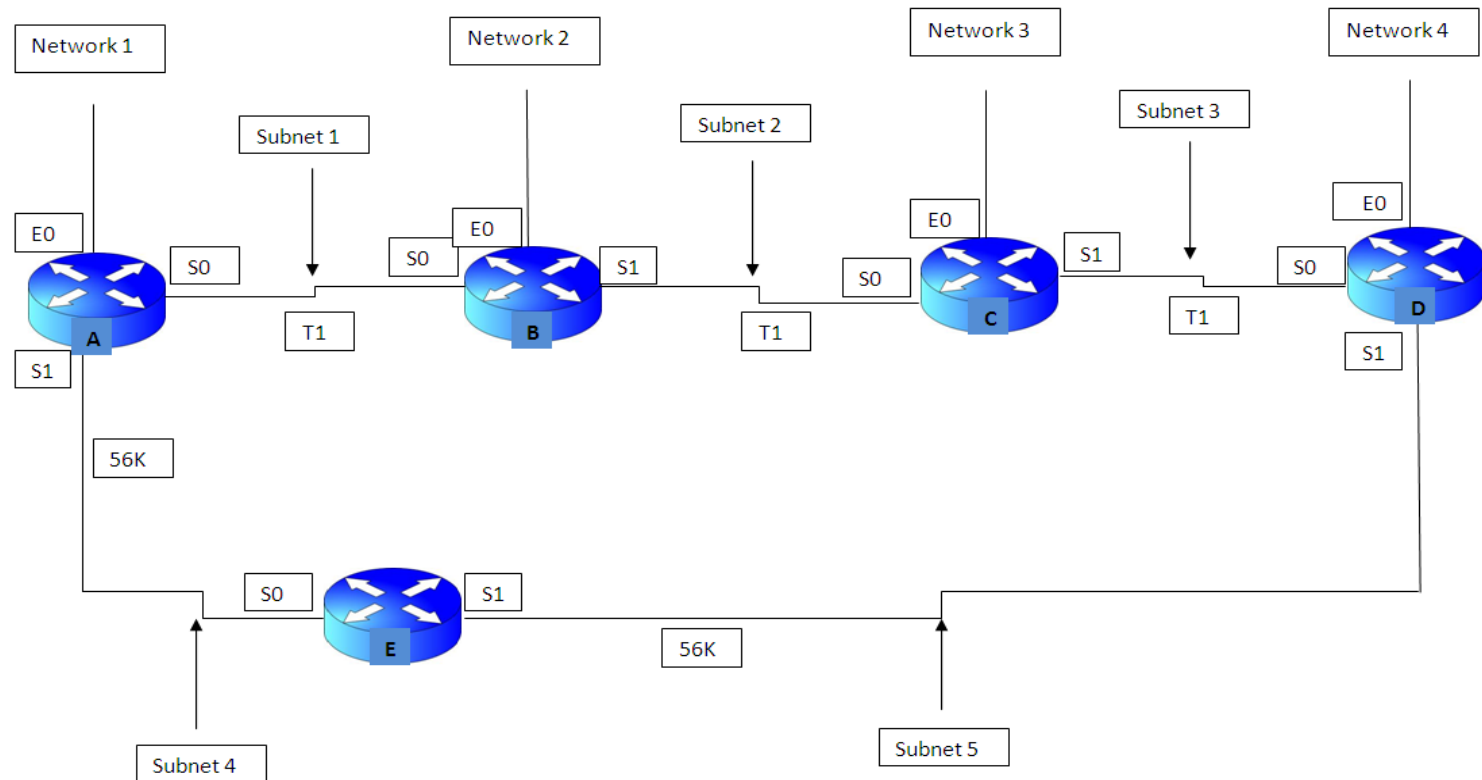
- ▶ In distance vector protocols, vectors are direction of the next hop
  - ▶ Router stores the IP address of the (next hop) router with the lower cost path
  - ▶ Next hop is the next location (router) packets will be forwarded towards the destination
  - ▶ Metric = distance
  - ▶ Direction = vector
- Distance Vector

# Metric

- ▶ Some protocols use hop count as the determinate of the distance to the destination, e.g. the lower the hop count the better the route
- ▶ Works well on networks with a stable transmission speed on network links – smaller controlled systems
- ▶ In larger complex systems with varying bandwidth, hop counts in isolation does not work.



# Determine the best path using the hop count



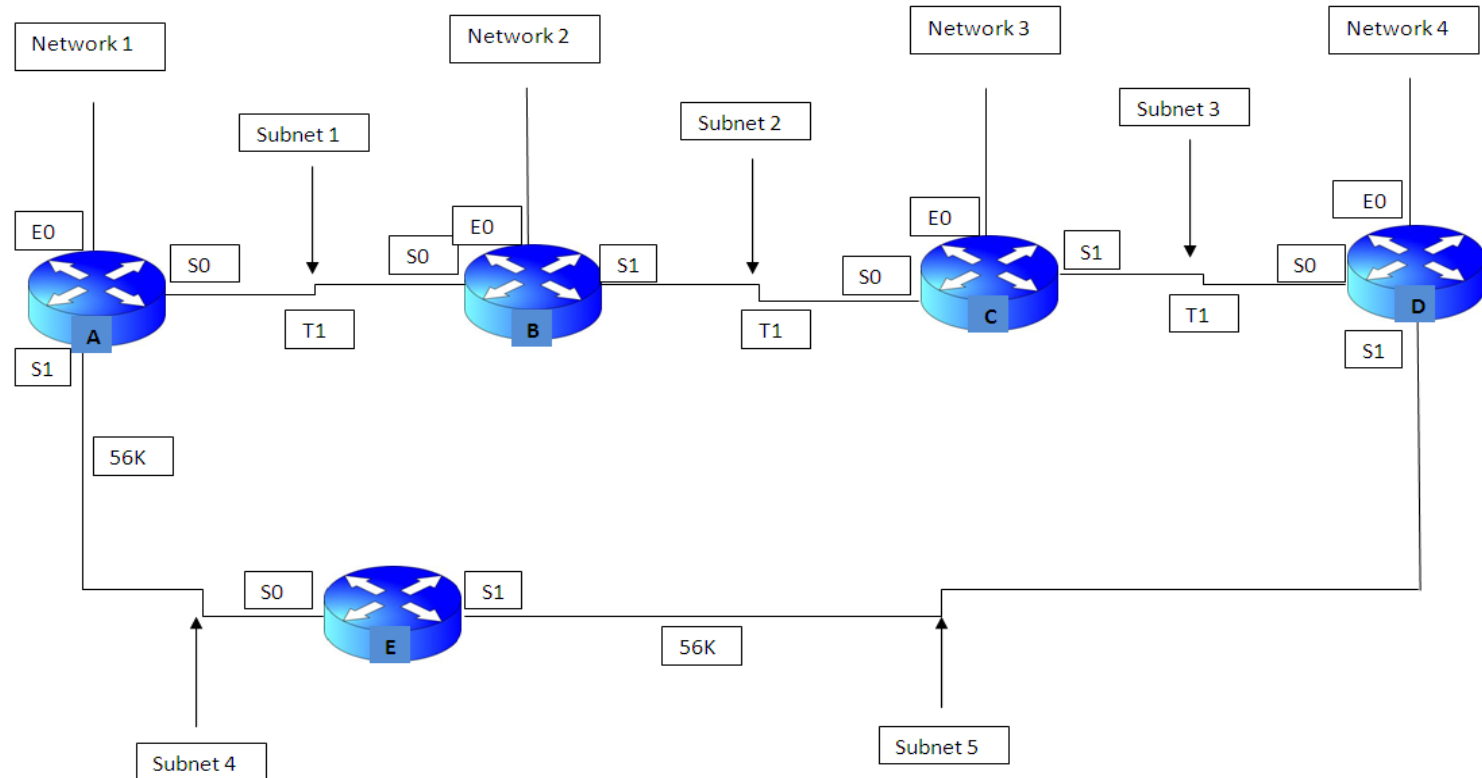
# Router A

- ▶ What does router A know?
- ▶ Its connected to network 1 via Ethernet (port) 0
- ▶ Its connected to subnet1 via serial (port) 0
- ▶ Its connected to subnet 4 via serial 1
- ▶ Network 2 is 1 hop away via serial 0
- ▶ Network 2 is 4 hops away via serial 1
- ▶ Network 3 is 2 hops away via serial 0
- ▶ Network 3 is 3 hops away via serial 1
- ▶ Network 4 is 3 hops away via serial interface 0
- ▶ Network 4 is 2 hops away via serial 1

# Determine the Path

- ▶ If packet X is at network 1 and its destination is network 4
- ▶ Router A identifies 2 paths:
  - Path 1: 3 hops away serial 0
  - Path 2: 2 hops away serial 1
- ▶ With distance vector protocols path 2 will be chosen based on hop count, even though path 1 is faster.
- ▶ The limitations of this was identified therefore bandwidth was used

# Determine the best path now using both hop and bandwidth



# Router A

- ▶ Its connected to network 1 via Ethernet 0
- ▶ Its connected to subnet1 via serial 0
- ▶ Its connected to subnet 4 via serial 1
- ▶ Network 2, 1,544kbps via serial 0
- ▶ Network 2, 2 X 56K links and two 1,544kbps via serial 1
- ▶ Network 3 1,544kbps via serial 0
- ▶ Network 3 2 X 56k, 1 X 1,544kbps via serial 1
- ▶ Network 4 3 X 1,544kbps serial 0
- ▶ Network 4 2 X 56K via serial 1
- ▶ Using bandwidth and hops, the 3 hops on 1,544kbps links would be selected

# Distance Vector problems..

- ▶ Hop count and bandwidth does improve the efficiency of the routing
- ▶ One problem is distance vector consume network resources since the full routing tables can be broadcast every 30 seconds by default
- ▶ Routing tables can be very large
- ▶ This process can also affect convergence due to delay incurred in sending so many update packets
- ▶ Distance vector protocols are prone to loops
- ▶ Routing loops are when two routers point to each other as the path to a network
- ▶ Therefore the packets bounce between the two routers

# Link state Routing Protocols

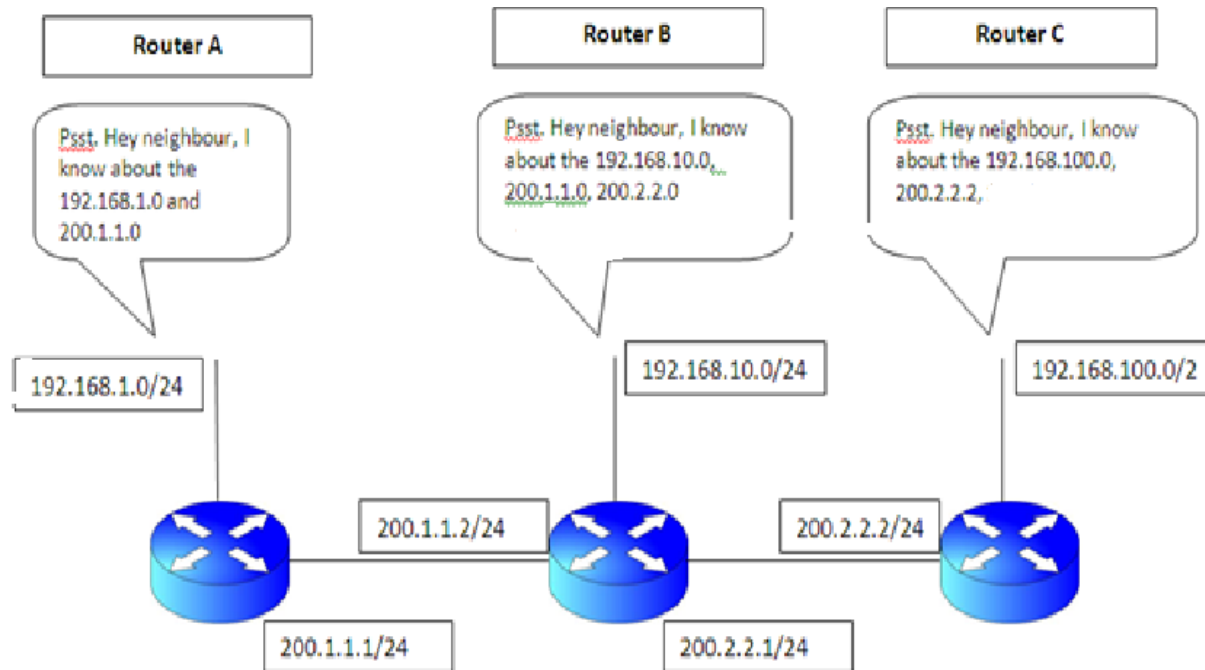
- ▶ Known as shortest path first
- ▶ Based on the Dijkstra algorithm
- ▶ Works on first-hand information not “routing by Rumours”
- ▶ Data is transmitted via Link State Advertisement (LSA)
- ▶ It includes the state of the directly connected routers Links
- ▶ Link state determines how many routers are out there and what networks are connected to them
- ▶ Each router ends up with a topology map of the system

# Link-State Routing Protocols

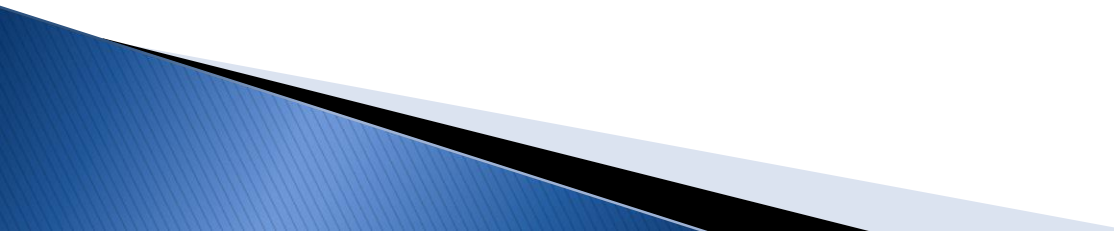
- ▶ As only the link state is communicated and not the whole routing table,
  - Speed of convergence is improved!
- ▶ Table updates are initiated on a change of a link state only
  - Minimises unnecessary use of available bandwidth
- ▶ When the Dijkstra algorithm is run the shortest quickest route is determined to populate the routing table
- ▶ This is less prone to routing loops as each router has a complete map of the system
- ▶ Routers are not tricked into routing packets back to themselves.



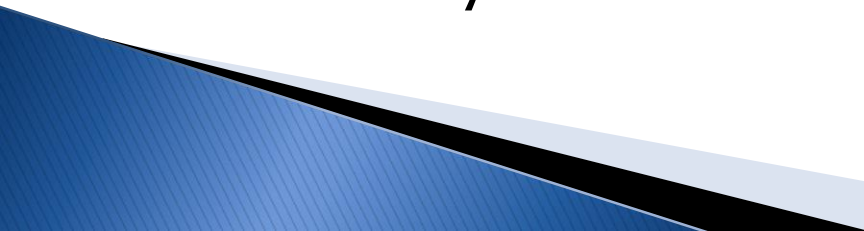
# Link state Routing Protocols



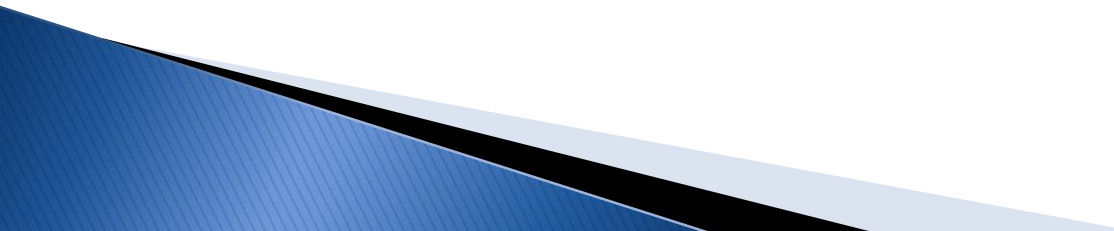
# Benefits of Link State algorithms..

- ▶ Update packets can be sent via multicast rather than broadcast
  - ▶ Reduces processing requirements on routers
  - ▶ Can be configured in a hierarchical fashion
  - ▶ Reduces unnecessary traffic
  - ▶ Eliminates Routing Loops
- 

# Multiple Routing Paths?

- ▶ Multiple paths to a network may exist
  - ▶ Not all routing protocols can actually install multiple paths
  - ▶ If only one path can be installed into the routing table it should be the best path. If this failed then the next best would be installed
  - ▶ If a multipath routing protocol is used a primary path can be identified also packets can be routed via multipaths to reduce throughput and load balancing – multiplexing
  - ▶ This improves network performance and reliability
- 

# Hierarchical Routing

- ▶ To reduce routing update on network bandwidth, routers can be configured in a hierarchical topology
  - ▶ Thereby **routers are grouped into AREAS** and some of the updates are confined to those areas
  - ▶ Areas will communicate as well but the updates are segregated on a need to know basis
  - ▶ Helps with the management of the network resources
- 

# Route Summarisation

- ▶ The concept of reducing the number of entries in route tables while still facilitating paths to all known networks
- ▶ Using subnetting means the route table increases
- ▶ Collating the routing data takes up network resources e.g. bandwidth
- ▶ Large route tables means the lookup process takes longer
- ▶ It also requires larger memory and CPU resources
- ▶ **Route summarisation defines a single path to multiple subnets**
  - Reduces the table size

# Route Summarisation

- ▶ Summarisation can be employed at the address assignment level and organisation level
- ▶ Auto-summarisation is available – the routing protocol summarises routes by default
  - Auto-summarisation can be disabled

# Routing issue

- ▶ A key problem in routing is a routing-loop
  - A packet travelling endlessly around the network without reaching its destination
  - The routing table does not hold the most up to date information
  - Routing decisions are based on incomplete/incorrect information
  - delay in network convergence is often the main cause of this.

# Summary

- ▶ What is Routing
  - ▶ What if's – the Complications of Routing
  - ▶ Routing Tables
  - ▶ An automated routing solution
  - ▶ Defining a routing Protocol
  - ▶ Design Considerations
  - ▶ Metrics of Routing Protocols
  - ▶ Categorising Routing protocols
- 