

# Design Documentation

## *Super Awesome Processor A (SAPA) v1.0*

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# Register Descriptions

Table of Register Descriptions

Register	Number	Availability	Description
<b>\$sn (0-2)</b> <(^.^)>	0-2	Read/write	General purpose registers. The intent is to store long term computation results and save values over functions calls
<b>\$tn (0-3)</b> <(-_-,)>	3-6	Read/write	General purpose registers, to be used like \$sn registers. Will NOT be saved over function calls.
<b>\$v</b> <(-_-,)>	7	Read/write	This register is for storing the return value from a function
<b>\$cr</b> <(-_-,)>	8	Read only*	Accumulator - stores most recently computed value
<b>\$ra</b> <(^.^)>	9	Read/write	Stores the return address of a function
<b>\$an (0-1)</b> <(-_-,)>	10-11	Read/write	These registers are used to store arguments for use in a called function
<b>\$st</b> <(^.^)>	12	Read/write	Reference "top" or lowest memory address of stack
<b>\$in (0-1)</b> :~/	13-14	Not available	Used by assembler for pseudo instructions.
<b>\$k</b> :~/	15	Read/write (while handling exceptions)	Exception registers; only accessible with permissions*

Unsafe between procedure calls = <(-\_-,)>

Safe between procedure calls = <(^.^)>

\* Enforced by exception handler

# Instructions - Assembly and Machine Language

## Procedure Call Conventions:

1. Registers  $\$in$ ,  $\$kn$  (where  $n$  is 0-1) are reserved for the assembler and operating system and should not be used by user programs or compilers.
2. Registers  $\$an$  (where  $n$  is 0-1) are used to pass arguments to procedures, any other arguments should go on the stack. Register  $\$v$  is used to return a value from functions.
3. Registers  $\$an$ ,  $\$tn$ ,  $\$d$ ,  $\$cr$ , and  $\$v$  are temporary and volatile. Expect them to contain different data after a procedure call.
4.  $\$sn$  registers must be backed up on the stack at the beginning of a procedure and restored before returning from the procedure. This preserves values in these registers over procedure calls.
5.  $\$st$  is the stack register. It points to the top memory location in the stack. If the stack is grown at any time in a procedure, it must be reduced before returning from that procedure.
  - a. Memory is allocated to the stack by subtracting from the value in  $\$st$ .  
Memory is deallocated from the stack by adding to the value in  $\$st$ .
6.  $\$ra$  is the return address of a procedure.  $Jal$  will overwrite  $\$ra$  to be the next instruction, so  $\$ra$  must ALWAYS be backed up on the stack before a procedure call and restored after returning from the procedure.
7. The instruction  $jr \$ra$  will return the program to the address.

## Syntax and Semantics:

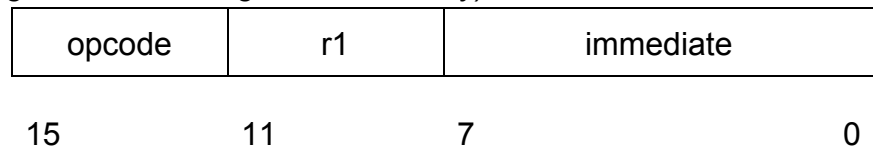
### Basic Instruction Formats

*C-type, Computation types (register to register)*

opcode	r1	r2	func	
15	11	7	3	0

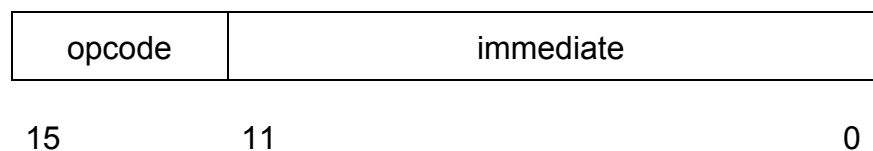
C-type instructions are used for register to register computations. They handle arithmetic and logical computations such as add, sub, and, or, etc. These instructions share a single opcode, and are distinguishable by their func code.

*I-type, Immediate types (register to data, register to memory)*



I-type instructions are used for register to data and register to memory computations. These instructions handle storing data from registers into memory and loading data into registers from memory, immediate values, and the \$scr register. They also handle control flow such as branches and jumps that are necessary for loops and procedure calls.

*L-type, Leap type*



L-type instructions are used for jal and j instructions.

## **Arithmetic and Logical Instructions**

*Arithmetic and logical instructions* are **C-type** instructions. These instructions take two registers as operands to their computations. Their results are always stored in the specialized \$scr register.

Arithmetic and logical instructions are directly translated from their assembly to the machine language. For example, the following assembly would translate accordingly into binary:

*add \$s2, \$t3*

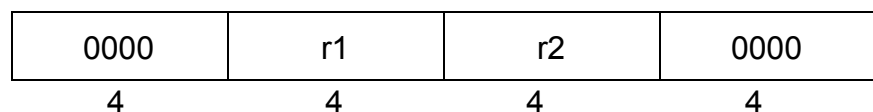


Registers are directly translated from their respective numbers in the registry (see chart in Registers).

### **List of Arithmetic and Logical Instructions:**

#### **Addition:**

*add r1, r2*



Stores the sum of r1 and r2 into register \$cr

**Subtraction:**

*sub r1, r2*

0000	r1	r2	0001
4	4	4	4

Stores the difference between r1 and r2 into register \$cr

**AND:**

*and r1, r2*

0000	r1	r2	0010
4	4	4	4

Stores the logical AND of r1 and r2 into register \$cr

**OR:**

*or r1, r2*

0000	r1	r2	0011
4	4	4	4

Stores the logical OR of r1 and r2 into register \$cr

**NOR:**

*nor r1, r2*

0000	r1	r2	0100
4	4	4	4

Stores the logical NOR of r1 and r2 into register \$cr

**NAND:**

*nand r1, r2*

0000	r1	r2	0101
4	4	4	4

Stores the logical NAND of r1 and r2 into register \$cr

**Exclusive OR:**

*xor r1, r2*

0000	r1	r2	0110
4	4	4	4

Stores the logical Exclusive OR of r1 and r2 into register \$cr

### Set Less Than:

*slt r1, r2*

0000	r1	r2	0111
4	4	4	4

Stores either a 1(True) or a 0(False) in \$cr depending on if r1 is less than r2

## Branch Instructions

Branch instructions are **I-type** instructions, and compare the value in \$cr to the value in r1. The branch instruction will then succeed or fail based on equivalence or inequality.

Branch instructions are PC-relative, meaning they use an 8-bit offset that allows a user to jump  $2^7-1$  instructions forward or  $2^7$  backward. A translation may appear as below (where “loop” is 3 instructions above bieq):

*bieq \$s0, loop*

0001	0000	1101
op	r1	BranchAddr

### List of Branch Instructions:

#### Branch if equal:

*bieq r1, location*

0001	r1	BranchAddr
4	4	8

Conditional branch to address in immediate if r1 is equal to register \$cr

#### Branch not equal:

*bneq r1, location*

0010	r1	BranchAddr
4	4	8

Conditional branch to address in immediate if r1 does not equal register \$cr

## Jump Instructions

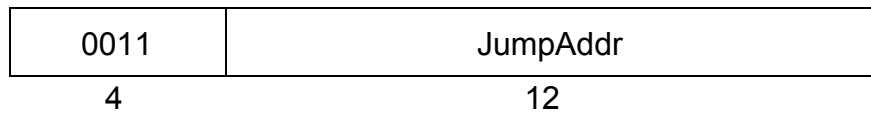
Jump instructions utilize two instructions formats: **L-type** and **I-type**.

Jump instructions use a 12-bit immediate with the top 4 bits of the PC concatenated to create a 16-bit address and allowing for a  $2^{11}$  size block of “jump space”. PC is set to this new address.

### List of Jump Instructions:

#### Jump:

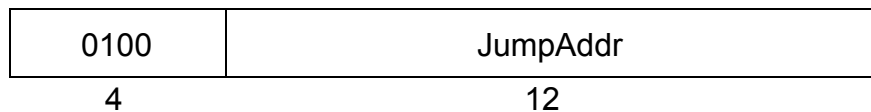
*j*      *location*



Unconditional jump to the address in immediate

#### Jump and link:

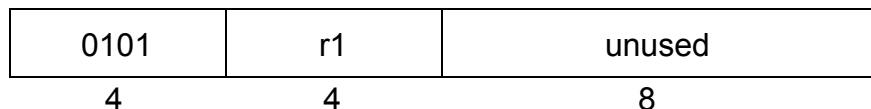
*jal*      *location*



Unconditional jump to the address in immediate, storing the address of subsequent instruction into register \$ra

#### Jump register:

*jr*      *r1*



Unconditional jump to the address in r1

## Load/Store Instructions

Load/Store instructions are **I-type** and often require a value to be computed and stored in \$cr prior to execution. Specific requirements are denoted for each instruction.

Load/store instructions are directly translated similar to arithmetic and logical instructions. The following assembly would translate accordingly:

*lw*      *\$t0*

1010	0101	unused
op	r1	Imm

### List of Load/Store Instructions:

#### Load upper immediate:

*lui r1, upper(big)*

0110	r1	Immediate
4	4	8

Load top half of 16-bit immediate into r1. Bottom bits are set to zero.

#### Load lower immediate:

*lli r1, lower(big)*

0111	r1	Immediate
4	4	8

Load lower half of 16-bit immediate into r1. Top bits are set to sign of immediate.

#### Load to register:

*ltr r1, immediate*

1000	r1	SignExtImm
4	4	8

Take a sign-extended 8-bit immediate value and store in r1. If immediate is greater than 8-bits, utilizes load upper immediate and load lower immediate.

*ltr r1, big*

Translates to:

*lui \$i0, upper(big)*

*lli \$i1, lower(big)*

*or \$i0, \$i1*

*ctr r1*

#### Copy to register:

*ctr r1*

1001	r1	unused
------	----	--------



4

4

8

Take a previously computed value from \$cr and store in r1

### Load word:

*lw r1*

1010	r1	unused
4	4	8

Access memory at the address stored in \$cr and store the value in r1

### Store word:

*sw r1*

1011	r1	unused
4	4	8

Store value in r1 at the previously computed memory address in register \$cr.

## Special Procedures

### Interact with I/O

*syscall*

1110	unused	unused
4	4	8

How to use syscall:

Step 1. Load appropriate code into \$v.

Step 2. Load argument values, if any, into \$a0 or \$a1 as specified.

Step 3. Issue the SYSCALL instruction.

Step 4. Retrieve return values, if any, from result registers as specified.

Example:

```
ltr    $v0, 1      # service 1 is print integer
ltr    $a0, 0      # load desired value into argument register $a0
syscall
```

### Table of Syscall Codes

Function	Integer in \$v	Argument or Return Value
PRINT_INT	1	\$a0 = value
PRINT_STRING	2	\$a0 = address of string
READ_INT	3	Result placed in \$v
READ_STRING	4	\$a0 = address, \$a1 = maximum length
EXIT	5	None
PRINT_CHAR	6	\$a0 low byte = character
READ_CHAR	7	Character returned in low byte of \$v

### **Pseudo Instructions**

These are instructions are not official instructions but sets of instructions that our processor will handle with smaller instructions. All pseudo instructions are **I-type**. It is important to note that the expansion of pseudo instructions often require **C-type** instructions. Therefore, users should generally assume that \$cr will be overwritten when executing pseudo instructions and backup \$cr if they wish to use it afterward.

#### **Set branch comparison:**

*sbc imm*

Sets register \$cr to a sign-extended 8-bit immediate value in preparation for a branch instruction.

Translation:

*ltr \$i0, 0*

*ltr \$i1, imm*

*add \$i0, \$i1*

**Set jump register:**

*sjr     JumpAddr*

Sets register \$ra to a jump address specified by the user.

Translation:

*ltr     \$i1, JumpAddr*  
*ltr     \$i0, 0*  
*add    \$i0, \$i1*  
*ctr    \$ra*

**OR immediate:**

*ori     r1, imm*

Stores the logical OR of r1 and immediate into register \$cr. 8-bit numbers are zero-extended.

8-bit Translation:

*lli     \$i0, imm     # Loads zero-extended lower 8 bits*  
*or     r1, \$i0*

16-bit Translation:

*lui     \$i0, upper(imm)     # Loads upper 8 bits, bottom bits filled with 0*  
*lli     \$i1, lower(imm)     # Loads zero-extended lower 8 bits*  
*or     \$i0, \$i0*  
*or     r1, \$cr*

**ADD immediate:**

*addi   r1, imm*

Stores the arithmetic result of r1 and immediate into register \$cr. 8-bit numbers are zero-extended. For 16-bit numbers, the following translation occurs:

8-bit Translation:

*ltr     \$i0, imm*  
*add    r1, \$i0*

16-bit Translation:

*lui     \$i0, upper(big)*  
*lli     \$i1, lower(big)*  
*or     \$i0, \$i1*  
*add    r1, \$cr*

**Load address:**

*la*      *r1, address*

Stores the sign-extended address in immediate into r1.

Translation:

*lui*      *\$i0, upper(address)*  
*lli*      *\$i1, lower(address)*  
*or*      *\$i0, \$i1*  
*ctr*      *r1*

**Register to Register:**

*rtr*      *r1, r2*

Moves the value in r2 to r1.

Translation:

*ltr*      *\$i0, 0*  
*add*      *r2, \$i0*  
*ctr*      *r1*

**Table of Machine Language Translations for Instructions**

Key: R[8] = \$cr, R[9] = \$ra

Instruction	Type	Code	Description of bits and rules			
Add	C	$R[8] = R[r1] + R[r2]$	0000	r1	r2	0000
Sub	C	$R[8] = R[r1] - R[r2]$	0000	r1	r2	0001
AND	C	$R[8] = R[r1] \& R[r2]$	0000	r1	r2	0010
OR	C	$R[8] = R[r1] \mid R[r2]$	0000	r1	r2	0011
NOR	C	$R[8] = \sim(R[r1] \mid R[r2])$	0000	r1	r2	0100
NAND	C	$R[8] = \sim(R[r1] \& R[r2])$	0000	r1	r2	0101
XOR	C	$R[8] = R[r1] \wedge R[r2]$	0000	r1	r2	0110
Set Less Than	C	$R[8] = (R[r1] < R[r2]) ? 1 : 0$	0000	r1	r2	0111

Branch Equal	I	If ( $R[8] = R[r1]$ ) $PC = PC + 1 +$ BranchAddr	0001	r1	BranchAddr
Branch Not Equal	I	If ( $R[8] \neq R[r1]$ ) $PC = PC + 1 +$ BranchAddr	0010	r1	BranchAddr
Jump	L	$PC = \text{JumpAddr}$	0011	JumpAddr	
Jump and Link	L	$R[9] = PC + 1$ $PC = \text{JumpAddr}$	0100	JumpAddr	
Jump Register	I	$PC = R[r1]$	0101	r1	8'b0
Load Upper Immediate	I	$R[r1] = \{\text{imm}, 8'b0\}$	0110	r1	Immediate
Load Lower Immediate	I	$R[r1] = \{8'b0, \text{imm}\}$	0111	r1	Immediate
Load to Register	I	$R[r1] = \text{SignExtImm}$	1000	r1	Immediate
Copy to Register	I	$R[r1] = R[8]$	1001	r1	8'b0
Load Word	I	$R[r1] = M[R[8]]$	1010	r1	8'b0
Store Word	I	$M[R[8]] = R[r1]$	1011	r1	8'b0
syscall	I	I/O	1110	4'b0	8'b0
Set Branch Comparison	-	$R[8] =$ $\text{SE}(\text{BranchComparison})$	<i>Pseudo instruction</i>		
Set Jump Register	-	$R[9] = \text{JumpAddr}$	<i>Pseudo instruction</i>		
ORi	-	$R[8] = r1 \mid \text{ZeroExtImm}$	<i>Pseudo instruction</i>		
Load Address	-	$R[r1] = \text{SignExtImm}$	<i>Pseudo instruction</i>		

Register to Register	-	R[r1] = R[r2]	<i>Pseudo instruction</i>
Addi	-	R[r1] = SignExtImm	<i>Pseudo instruction</i>

### Table of Common Operations:

	SAPA 1.0	Description
Load Address	<i>la</i> \$s2, 0x4EF6	Loading an address is a pseudo instruction. Refer to <i>Pseudo instruction</i> for full details on how <i>la</i> works.
Arithmetic / Logical	<i>ltr</i> \$t1, 1 <i>ltr</i> \$t2, 1 <i>add</i> \$t1, \$t2	Load values into two registers and call add to store their sum in \$s1. Similar for other arithmetic and logical operations.
Iterations	... <i>ltr</i> \$s1, 15 <i>ltr</i> \$t0, 1 <i>ltr</i> \$t1, 0 <i>ltr</i> \$t2, 3 <i>loop:</i> <i>slt</i> \$s1, \$t1 <i>bieq</i> \$t0, exit <i>add</i> \$t1, \$t2 <i>ctr</i> \$t1 <i>j</i> loop <i>exit:</i> ...	This <i>for-loop</i> keeps adding 3 to register \$s3 (0) until \$t1 becomes greater than \$s1 (15), resulting with a 18 stored in \$t1 after the loop has exited.
Branches	See above	The example above utilizes <i>slt</i> and <i>bieq</i> to create a branch on greater than, which isn't an instruction itself but can be created using multiple instructions.
Jumps	... <i>add2:</i> <i>ltr</i> \$t0, 2 <i>add</i> \$t0, \$a0 <i>ctr</i> \$t1 <i>sbc</i> 20 <i>bneq</i> \$t1, add2	This example of code does recursive addition of adding two until the value of 20 is reached. However, it needs to move to end instead of moving to the next set of recursive code. Unfortunately, it is too far for the regular jump code to reach. The user can instead; however, set the value of

	<i>sjr</i> <i>end</i>	the desired address to the \$ra register and can now jump to it.
	<i>jr</i> <i>\$ra</i>	
	...	
<i>end:</i>	...	

## relPrime and Euclid's Algorithm (Assembly):

```

    ltr    $jr, 65      # preset $jr to 65
    ltr    $a0, 4       # n = 4
relPrimeSetup:
    addi   $st, 0       # $cr gets first address on stack
    sw     $ra          # store $ra on stack
    addi   $st, -4      # $cr gets $st - 4
    ctr    $st          # grow stack by 4
    sw     $s0          # store $s0 on stack
    ltr    $t1, 2       # m = 2
    rtr    $a1, $t1     # $a1 = m
    rtr    $s0, $t1     # $s0 = m
relPrimeLoop:
    jal    gcd
    sbc     1           # $cr = 1
    bieq   $v, cleanup  # if gcd(n, m) == 1, branch to cleanup
    addi   $s0, 1       # m + 1
    ctr    $s0          # m = m + 1
    rtr    $a1, $s0     # $a1 = m
    j      relPrimeLoop
gcd:
    sbc     0           # $cr = 0
    bneq   $a0, gcdLoop # if a != 0, branch to gcdLoop
    rtr    $v, $a1      # $v = b
    jr     $ra          # return to caller
gcdLoop:
    sbc     0           # $cr = 0
    bneq   $a1, subOne  # if b != 0, branch to subOne
    rtr    $v, $a0      # $v = a
    jr     $ra          # return to caller
subOne:
    slt    $a1, $a0     # $cr = b < a

```

```

ltr    $t0, 0
bieq   $t0, subTwo # if !(a > b) go to subTwo
sub    $a0, $a1     # $cr = a - b
ctr    $a0          # a = a - b
j      gcdLoop      # return to gcdLoop
subTwo:
sub    $a1, $a0     # $cr = b - a
ctr    $a1          # b = b - a
j      gcdLoop      # return to gcdLoop
cleanup:
rtr    $v, $s0      # $v = m
addi   $st, 0       # $cr = $st
lw     $s0          # restore $s0
addi   $st, 4       # $cr = $st + 4
ctr    $st          # $st = $st + 4
lw     $ra          # restore $ra
jr     $ra

```

### relPrime and Euclid's Algorithm (Machine Language):

ltr	\$jr, 65	# preset \$jr to 65	1000	1001	01000001
ltr	\$a0, 4	# n = 4	1000	1010	00000100
relPrimeSetup:					
addi	\$st, 0	# \$scr gets first address on stack	1000	1101	00000000
			0000	1100	1101 0000
sw	\$ra	# store \$ra on stack	1011	1001	00000000
addi	\$st,-4	# \$scr gets \$st - 4	1000	1101	11111100



			0000	1100	1101	0000
ctr	\$st	# grow stack by 4				
			1001	1100	00000000	
sw	\$s0	# store \$s0 on stack				
			1011	0000	00000000	
ltr	\$t1, 2	# m = 2				
			1000	0100	00000010	
rtr	\$a1, \$t1	# \$a1 = m				
			1000	1101	00000000	
			0000	0100	1101	0000
			1001	1011	00000000	
rtr	\$s0, \$t1	# \$s0 = m				
			1000	1101	00000000	
			0000	0100	1101	0000
			1001	0000	00000000	
relPrimeLoop:						
jal	gcd		0100	000000011100		
sbc	1	# \$cr = 1				
			1000	1101	00000000	
			1000	1110	00000001	

0000	1101	1110	0000
------	------	------	------

bieq \$v, cleanup # if gcd(n, m) == 1 branch to cleanup

0001	0111	00100000
------	------	----------

addi \$s0, 1 # m + 1

1000	1101	00000001
------	------	----------

0000	0000	1101	0000
------	------	------	------

ctr \$s0 # m = m + 1

1001	0000	00000000
------	------	----------

rtr \$a1, \$s0 # \$a1 = m

1000	1101	00000000
------	------	----------

0000	0000	1101	0000
------	------	------	------

1001	1011	00000000
------	------	----------

j relPrimeLoop

0011	000000010000
------	--------------

gcd:

sbc 0 # \$cr = 0

1000	1101	00000000
------	------	----------

1000	1110	00000000
------	------	----------

0000	1101	1110	0000
------	------	------	------

bneq \$a0, gcdLoop # if a != 0, go to subOne

0010	1010	00000100
------	------	----------

rtr \$v, \$a1 # \$v = b

1000	1101	00000000
------	------	----------

		0000	1011	1101	0000
		1001	0111	00000000	
jr	\$ra	# return to caller			
		0101	1001	00000000	
gcdLoop:					
sbc	0	# \$scr = 0			
		1000	1101	00000000	
		1000	1110	00000000	
		0000	1101	1110	0000
bneq	\$a1, subOne	# if b != 0, branch to subOne			
		0010	1011	00000100	
rtr	\$v, \$a0	# \$v = a			
		1000	1101	00000000	
		0000	1010	1101	0000
		1001	0111	00000000	
jr	\$ra	# return to caller			
		0101	1001	00000000	
subOne:					
slt	\$a1, \$a0	# \$scr = a > b			
		0000	1011	1010	0111
ltr	\$t0, 0				
		1000	0011	00000000	
bieq	\$t0, subTwo	# if !(a > b) go to subTwo			

			0001	0011	00000011	
sub	\$a0, \$a1	# a - b	0000	1010	1011	0001
ctr	\$a0	# a = a - b	1001	1010	00000000	
j	gcdLoop	# return to gcdLoop	0011	000000100100		
subTwo:						
sub	\$a1, \$a0	# b - a	0000	1011	1010	0001
ctr	\$a1	# b = b - a	1001	1011	00000000	
j	gcdLoop	# return to gcdLoop	0011	000000100100		
cleanup:						
rtr \$v, \$s0	# \$v = m	1000	1101	00000000		
		0000	0000	1101	0000	
		1001	0111	00000000		
addi \$st, 0	# \$cr = \$st	1000	1101	00000000		
		0000	1100	1101	0000	
lw \$s0	# restore \$s0					

		1010	0000	00000000
addi	\$t0, 4	# \$cr = \$st + 4		
		1000	1101	00000100
		0000	1100	1101 0000
ctr	\$st	# \$st = \$st + 4		
		1001	1100	00000000
lw	\$ra			
		1010	1001	00000000
jr	\$ra			
		0101	1001	00000000

# Register Transfer Language (RTL)

## Definitions:

CR = Register[8]

RA = Register[9]

I0 = Register[13]

I1 = Register[14]

R1 = Register[IR[11-8]]

R2 = Register[IR[7-4]]

OP = IR[15-12]

IMM8 = IR[7-0]

IMM12 = IR[11-0]

IMM16 = 16-bit immediate

IR = Instruction Register

MDR = Memory Data Register

## Common RTL:

Cycle	Label	RTL
0	FETCH INSTR:	PC = PC + 1 IR = Mem[PC]
1	DECODE:	A = R1 B = R2 C = CR ALUout = PC + SE(IMM8)
Last	DONE:	Goto FETCH INSTR

## Unique RTL:

Instruction -	Type C	Op 0000	func 0000 - 0111
Cycle	Label	RTL	
2	OPERATION:	ALUout = A op B	
3		CR = ALUout	

<b>Instruction</b> Branch if equal / not equal	<b>Type</b> <b>I</b>	<b>Op</b> <b>0001 / 0010</b>	<b>func</b> <b>-</b>
<b>Cycle</b>	<b>Label</b>	<b>RTL</b>	
<b>2</b>	BIEQ:	if (A == C) then PC = ALUout	
	BNEQ:	if (A != C) then PC = ALUout	
<b>3</b>	STALL		

<b>Instruction</b> Load upper / lower Imm.	<b>Type</b> <b>I</b>	<b>Op</b> <b>0110 / 0111</b>	<b>func</b> <b>-</b>
<b>Cycle</b>	<b>Label</b>	<b>RTL</b>	
<b>2</b>	LUI:	R1 = IMM8[8-0] << 8	
	LLI:	R1 = ZE(IMM8[8-0])	

<b>Instruction</b> Load to register	<b>Type</b> <b>I</b>	<b>Op</b> <b>1000</b>	<b>func</b> <b>-</b>
<b>Cycle</b>	<b>Label</b>	<b>RTL</b>	
<b>2</b>	LTR16:	if IMM16 then I0 = IMM16[15-8] << 8 I1 = ZE(IMM16[7-0])	
	LTR8:	if IMM8 then R1 = SE(IMM8)	
<b>3</b>	LTR16:	ALUout = I0   I1	
<b>4</b>	LTR16:	R1 = ALUout	



Instruction Copy to register	Type I	Op 1001	func -
Cycle	Label	RTL	
2	CTR:	R1 = C	

Instruction Load word	Type I	Op 1010	func -
Cycle	Label	RTL	
2	LD FROM MEM:	MDR = Mem[C]	
3	LWSTALL		
4	LOAD TO REG:	R1 = MDR	
5	STALL		

Instruction Store word	Type I	Op 1011	func -
Cycle	Label	RTL	
2	SW:	Mem[C] = A	
3	STALL		

Instruction syscall	Type I	Op 1110	func -
Cycle	Label	RTL	
2	INCR PC:	PC = PC + 1	
?	READ:	if FINISHED then goto RDONE	
	WRITE:	if FINISHED then goto WDONE	

	EXIT:	Close program
?	RLOOP:	if FINISHED then goto RDONE else goto RLOOP
	WLOOP:	if FINISHED then goto WDONE else goto WLOOP
?	RDONE: / WDONE:	

Instruction Jump Register	Type I	Op 0101	func -
Cycle	Label	RTL	
2	JR:	PC = A	
3	STALL		

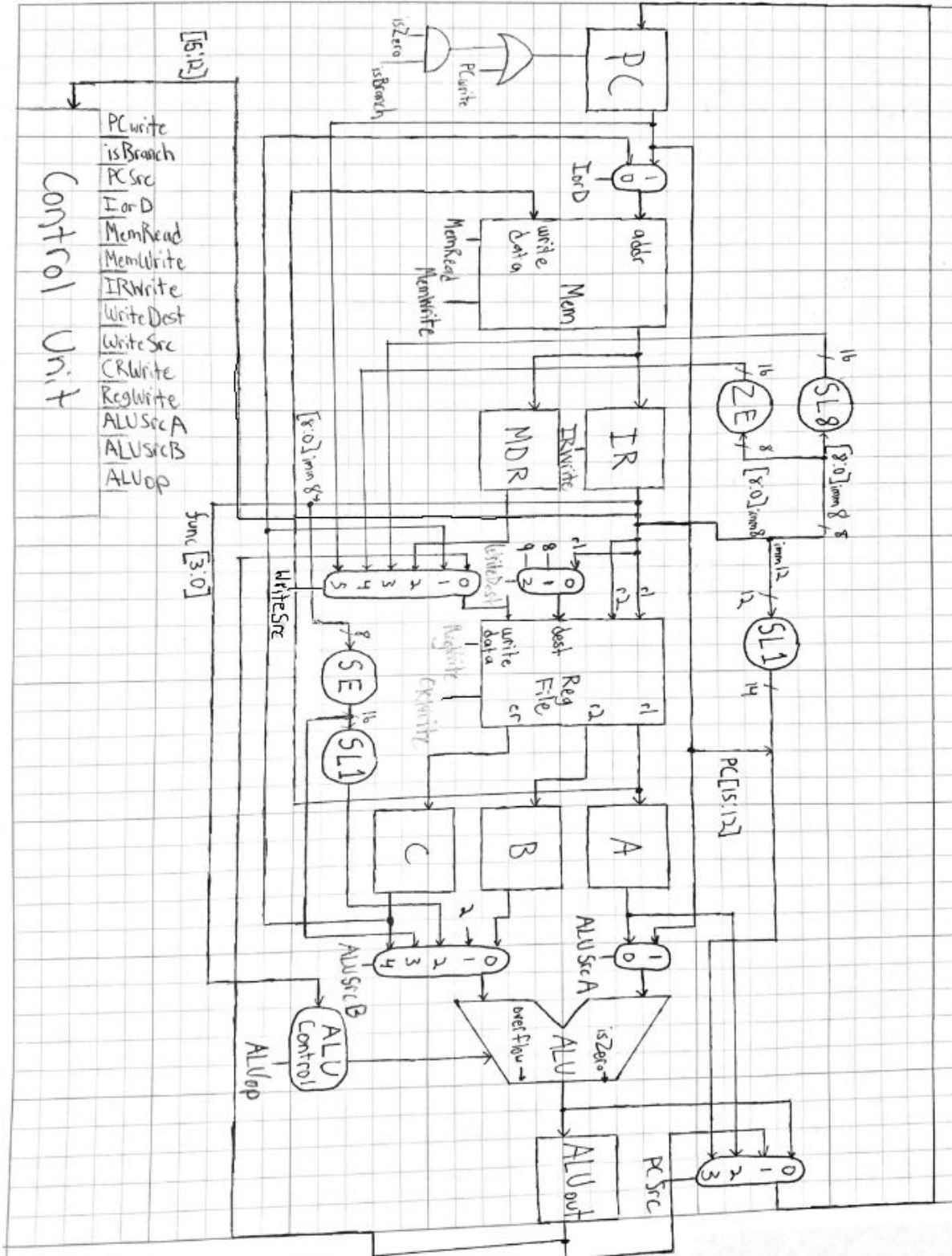
Instruction -	Type L	Op 0011 / 0100	func -
Cycle	Label	RTL	
2	J:	PC = (IMM12)   PC[15-12]	
	JAL:	RA = PC PC = (IMM12)   PC[15-12]	
3	STALL		

# Datapath Components

**Table of Components and Signals**

Signals Components	Inputs (size)	Outputs (size)	Description of Component
PC	- newPC 16'b	- PC 16'b	- PC => Program counter
Memory	- Address 16'b - Write Data 16'b	- MemData 16'b	- Access to memory
IR	- Instruction Data 16'b	- r1 4'b - r2 4'b - imm 8'b or 12'b	- IR => instruction register - Stores instruction data
MDR	- Memory Data 16'b	- Memory Data 16'b	- MDR => Memory data register - Stores memory data
Register File	- r1 4'b - r2 4'b - WriteDest 4'b - Write Data 16'b	- r1 data 16'b - r2 data 16'b	- Registers that store data for reading/writing
A/B Registers	- A or B respectively 16'b	- A or B respectively 16'b	- These registers hold operands for the ALU
ALU	- A 16'b - B 16'b	- isZero 1'b - Overflow 1'b - Operation Result 16'b	- Performs operation on two operands
Shift Left by 1	- Imm 12'b	- Shifted imm size + 1	- Shifts immediate by 1
Shift Left by 8	- Imm 8'b	- Shifted imm size + 8	- Shifts immediate by 8
Sign Extender	- Imm 8'b	- Sign Extended imm 16'b	- Fill in top bits of immediate with sign of immediate
Zero Extender	- Imm 8'b	- Zero extended imm 16'b	- Fill in top bits of immediate with 0s

## Block Diagram of Datapath



# Control Unit

Table of Control Signals

Control Signal	Size	Component	Description
PCSrc	2'b	PC Source (Mux)	Controls the source of PC: incremented PC, jump address, or branch address.
PCWrite	1'b	PC	Controls whether PC should be written to
isZero	1'b	PC	ALU isZero result *Branch condition
isBranch	1'b	PC	Controls whether an instruction is a branch *Branch condition
isBIEQ	1'b	ALU	Determines whether an branch instructions is a bieq
lorD	1'b	MemAddr (Mux)	Decides whether a memory address is coming from PC (instruction) or the C register (data)
MemRead	1'b	Memory	Controls whether memory is being read
MemWrite	1'b	Memory	Controls whether memory is being written to
IRWrite	1'b	IR	Controls whether IR is being written to
WriteDest	2'b	RegDest (Mux)	Controls whether destination of a register write is \$cr or r1
WriteSrc	3'b	WriteData (Mux)	Controls whether source of data to be written to a register (MDR, ALout, C, etc.)
CRWrite	1'b	Register File	Controls whether the \$cr register can be written to (ie disable \$cr being set as r1 in <b>I-type</b> instructions)

<b>RegWrite</b>	1'b	Register File	Controls whether a register is being written to
<b>ALUSrcA</b>	1'b	ALUSrcA (Mux)	Controls whether the source of ALU input A is PC or the A register
<b>ALUSrcB</b>	2'b	ALUSrcB (Mux)	Controls whether the source of ALU input B is one of three immediates or the B register
<b>ALUOp</b>	2'b	ALUControl	Controls the operation sent to the ALUControl component

## Specification for Controls:

### PCSrc

- 00 - PC + 1
- 01 - ALUout
- 10 - Register A
- 11 - (PC[15-12] | IMM12)

### PCWrite

- 0 - turns off; 1 - turns on

### isBranch

- 0 - not branch
- 1 - is branch

### isBIEQ

- 0 - is bieq instruction
- 1 - is not bieq instruction

### lorD

- 0 - C register (data)
- 1 - PC (instruction)

### MemRead

- 0 - turns off; 1 - turns on

### MemWrite

- 0 - turns off; 1 - turns on

### IRWrite

- 0 - turns off; 1 - turns on

### WriteDest

- 00 - r1
- 01 - \$cr
- 10 - \$ra

### WriteSrc

- 000 - ALUout
- 001 - Register C
- 010 - MDR
- 011 - (IMM8 << 8)

100 - ZE(IMM8)

101 - PC

110 - SE(IMM8)

### CRWrite

- 0 - turns off; 1 - turns on

### RegWrite

- 0 - turns off; 1 - turns on

### ALUSrcA

- 0 - Register A
- 1 - PC

### ALUSrcB

- 00 - Register B
- 01 - 1
- 10 - SE(IMM8)
- 11 - Register C

### ALUop

- 00 - Add (PC increment)
- 01 - Sub (Branch)
- 10 - function code (C-type)
- 11 - OR

### func

4'b function determined by C-type

### opcode

Determined by instruction

### current\_state/next\_state

Keeps track of state in multi-cycle

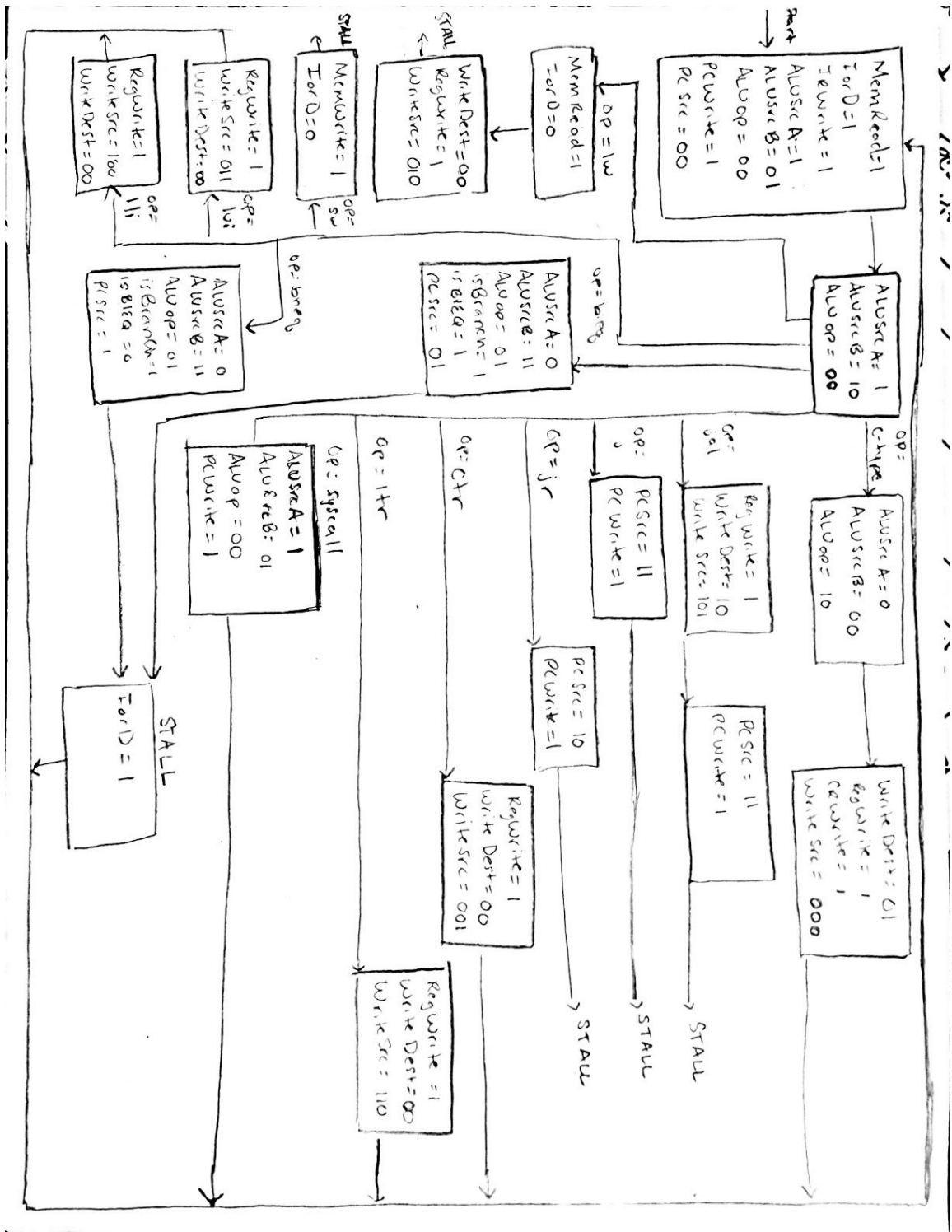
### CLK

Clock speed

### Reset

- 0 - Don't Reset; 1 - Reset

# Finite State Machine Transition Diagram





# Implementation and Testing

## **Hardware Implementation Plan:**

Special Registers (PC, IR, MDR, A, B, C):

- Each register will be able to hold 16-bits and have one output with as much.
- They will have multiple inputs such as the clock signal, a reset signal, and an input for the data.
- These will be made with a Verilog module in Xilinx.

Memory:

- For each bit, use a flip flop.
- Use buses and muxes to combine the data.
- Built as a schematic using these parts in Xilinx.

Register File:

- The register file will be made of registers and parts for the logic required to choose the correct outputs.
- 4'b address input and value output
- The cr register will always be outputted to a special C register.
- The register file will be modified from the resource provided to us to handle 16 registers.

ALU:

- For a 16-bit ALU, we will add a unit for Arithmetic (adds and subs), a unit for Logic (OR, NOR, XOR, AND, NAND), and a unit for less than (slt).
- Each unit will have a 16-bit input and output with their outputs connected to a mux.
- The mux chooses which data to output of the ALU, based on the ALU control signal.
- The ALU can be built in Xilinx using schematics, a modified design from the 4-bit ALU built in Lab 6.

Sign Extender:

- A simple verilog module that takes the most significant bit of the input and extends it to the left, extending the size of the data to 16 bits.

Zero Extender:

- Another verilog module that fills top bits with zeros, extending the size of the data to 16 bits.

Shifters:

- Verilog module that will take input and shift the bits to certain amount.

Muxes (1'b, 2'b, 3'b, 4'b)

- Modify resources given to us

## **Unit Testing:**

### **General Testing:**

- Consider instruction description including inputs and outputs
- Draw out the path of data flow from one component to another
- Confirm desired data source input and output for each component
  - Upon error return to last step and debug
  - Repeat until problem is found
- Confirm desired result/action at the end of calculation

### **Regular and Special Registers (A, B, C, IR, MDR):**

Overall, all of these registers have one thing that they need to accomplish. They need to be able to take in values, hold them and then output that same value on the output for that register.

### **Memory:**

The memory needs to be able to hold the instructions that we will be using for all operations and other values that the user might want to store into it after the processor starts doing operations. The memory should be able to output any instruction we have stored in it as well as any data when called upon.

### **Register File:**

The register file needs to be able to select any registers located within it (Note: The only situation that it should not be able to select the proper register is if the user tries selecting a register that they are not allowed to manipulate themselves). These registers, just like the previous ones, must be able to hold any and all data values presented to them and output them on their respective output. It must be able to write any data to the proper register when called upon (same restrictions as note).

### **ALU:**

The ALU is a combination of different logic gates that can perform different tasks, arithmetic and logical. At the end, it needs to be able to select the correct output among the multiple outputs it will calculate. Below are all the possible operations we will need to test. There are also two flags that we will need to make sure they are set at proper times, they will be tested with the functions that will need to be able to manipulate these flags.

- Addition
  - Needs to be able to handle big additions that cause overflow like numbers out of the range of a 16 bit data type.
  - Needs to be able to handle adding positive and negative numbers, those that will and will not cause overflow. Basically, detect overflow.
- Subtraction

- Be able to set an IsZero flag upon subtracting two values to get zero.
- Needs to be able to handle subtracting positive and negative numbers from each other, those that will and will not cause overflow.
- Needs to be able to set a Negative flag to represent the confirmation that first number in the argument is less than the second number in the argument.
- Bit AND any two numbers of similar bit size
- Bit OR any two numbers of similar bit size
- Bit NOR any two numbers of similar bit size
- Bit XOR any two numbers of similar bit size
- Bit NAND any two numbers of similar bit size

### **Control Unit:**

The control unit must be able to take in an instructions opcode and proceed to handle the flow of cycles. This means that for each instruction, the transition of the datapath's state must be checked at every stage to ensure correct control bits for the respective cycle.

This will likely need to be tested in Verilog with state transitions and simulations for each instruction.

### **ALUControl:**

This will be tested to ensure that it can take in the function bits decoded from the IR register and hold them until the control finishes decides what it needs to do. It should then be able to either send the function code to the ALU if it needs to use it or send a different function code than what was originally sent into the control.

### **Zero/Sign Extenders:**

- Must be able to extend any immediate we give correctly.

### **Shifters:**

- Shift Left 1:
  - Must be able to shift in a zero to the right of any immediate we give it, even a 16 bit immediate (the bit on the far left would be lost after the shift in this situation).
- Shift Left 8:
  - Must be able to shift in a zero to the right, like the Shift Left 1, just seven more times.

## Integration plan and testing:

**General Note:** Assume all steps are independent of each other and only require listed components.

### Step 0:

Goal: Increment PC

Components: PC, ALU (No Mux)

Testing:

General:

- Input A: PC
- Input B: 1
- Op: XXXX (Add)
- Output: PC + 1

Edge Cases:

- N/A

### Step 1:

Goal: Read Memory

Components: PC, MemAddr Mux, Memory, MDR (Read-only)

Testing:

General:

NOTE: PCWrite is set to 1

- MemAddr Mux properly selects 1
- Correct address is returned from MemAddrMux
- With MemRead set to 1, value at PC's memory address is outputted
  - MDR value is now value at PC's memory address
- With MemRead set to 0, no value is outputted
  - MDR value is unchanged

Edge Cases:

- Memory address doesn't exist
- No data in PC
- No data at that memory address
- MemWrite set to 1, behavior is expected

### Step 2:

Goal: Write Memory

Components: MemAddr Mux, Memory (Write only)

Testing:

General:

- MemAddr Mux properly selects 0
- Correct address is returned from MemAddrMux

- With MemWrite set to 1, value at address is set to write data
- With MemWrite set to 0, value at address is unchanged

Edge Cases:

- With MemRead set to 1, behavior is expected
- Memory address doesn't exist

Step 3:

Goal: Read Register File motivated by IR

Components: IR, Register File (Read-only)

Testing:

General:

- When IRWrite is set to 1, IR is overwritten
- When IRWrite is set to 0, IR is unchanged
- Bits 11-8 from the Instruction Register go to r1 in the register file
- Bits 7-4 from the Instruction Register go to r2 in the register file
- The value in r1 is correctly outputted from the register file
- The value in r2 is correctly outputted from the register file
- Value in \$cr is outputted

Edge Cases:

- Bits 11-8 do not refer to an existing register.
- Bits 7-4 do not refer to an existing register.
- The Instruction Register has no data.
- IRWrite is set to 0, and no data is in IR

Step 4:

Goal: Write to Register File

Components: WriteDest Mux, WriteData Mux, Register File (Write only)

Testing:

General:

- WriteDest Mux can select between 2 - \$ra, 1 - \$cr and 0 - another register and output correct value
- WriteData Mux can select between 5 sources and output correct value
- When RegWrite is set to 1 and WriteDest is set to 0, correct register stores correct value after write
- When RegWrite is set to 0 and WriteDest is set to 0, register destination is unchanged
- When CRWrite is set to 1, RegWrite is set to 1, and WriteDest is set to 1, \$cr contains
- When CRWrite or RegWrite is set to 0, and WriteDest is set to 1, \$cr is unchanged

- No other register should be written to

Edge Cases:

- WriteDest does not exist (input of  $\geq 3$ )
- WriteData does not exist (input of  $\geq 5$ )

Step 5:

Goal: ALU ops and outputs

Components: ALUSrcA Mux, ALUSrcB Mux, ALU, ALUControl

Testing:

General:

- ALUSrcA Mux can select between 2 sources and output correct value
- ALUSrcB Mux can select between 4 sources and output correct value
- ALUControl outputs proper value based on ALUop and func
- When computation results in zero, isZero output is 1, otherwise output is 0
- When computation results in overflow, overflow output is 1, otherwise 0

Edge Cases:

- Func code does not exist
- Input for logical operations are different in bitsize
- Overflow occurs
- ALUop requires function code as output

Step 6:

Goal: PCSrc and control signals

Components: PCSrc Mux, PCWrite, isBranch, isZero, PC

Testing:

General:

- PCSrc properly selects between 4 sources and outputs correct value
- PCWrite set to 1, PC is changed to PCSrc output
- PCWrite set to 0, isBranch and isZero set to 1, PC is changed to PCSrc output
- PCWrite set to 0, isBranch or isZero set to 0, PC is unchanged

Edge Cases:

- 

Step 7:

Goal: Shifters and extenders

Components: Shifters, Extenders

Testing:

General:

- Shift Left 1:
  - Ensure any sized input is shifted left one and outputted with original size + 1
- Shift Left 8:
  - Ensure any size input is shifted left eight and outputted with original size + 8
- Zero Extend:
  - 8-bit input is outputted with zero extension
- Sign Extend:
  - 8-bit input is outputted with sign extension

Edge Cases:

- Data of incorrect size is inputted

Final Step:

Combine pieces of processor together and proceed to system testing.

### **System test plan:**

- a. Individual instructions
- b. Small programs
  - i. Code Fragments
    1. Small loop
    2. Small programs
      - a. Simple arithmetic and logic
      - b. Simple jumps
- c. Large programs
  - i. relPrime (see pg. 13-20)