

Pt2. First, when `getpid()` is called, execution starts in the user space at `user.h`, where it is declared as a system call. The call itself is handled in `usys.S` where `SYSCALL(getpid)` is a macro that expands into x86 that initiates a syscall by putting its syscall number (18) into `%eax`. Then the `int $64` instruction is read which triggers an interrupt. Now in the kernel space `trap.c` then catches the interrupt and calls `syscall()`, defined in `syscall.c`. Syscall grabs the syscall number from the `eax` register and assigns it to a variable. This function looks up the syscall number, 18 in this case, from the array of all the syscall numbers and sends the control to the appropriate function. `sys_getpid()` is implemented in `sysproc.c` and returns `myproc()->pid`. `myproc()` returns a pointer to the PCB of the current running process, represented by the struct `proc`, which has PID. execution moves back to `syscall()` in `syscall.c`, which stores the return value into `%eax`, then the execution returns to `trap.c`, resets all the registers back to their original states and gives the return value to `getpid()`, officially going back to user space.