Machine-Level Programming IV: Data

Instructors:

Erol Sahin

Today

Arrays

- One-dimensional
- Multi-dimensional (nested)
- Multi-level

Structures

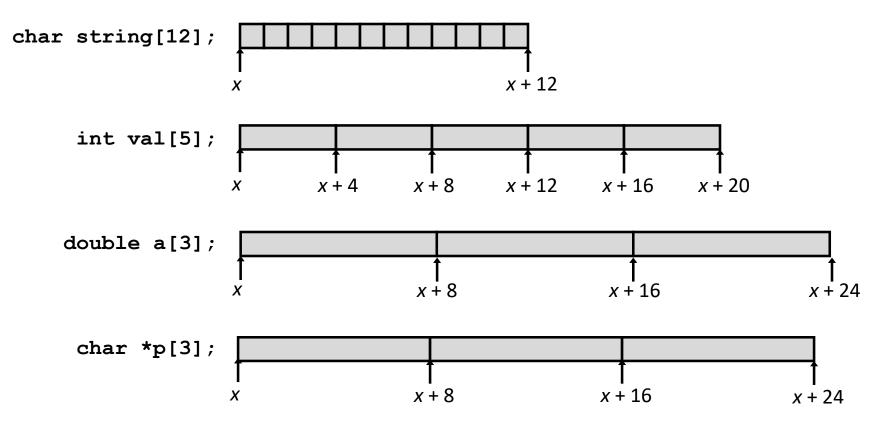
- Allocation
- Access
- Alignment
- Floating Point

Array Allocation

Basic Principle

```
T A[L];
```

- Array of data type T and length L
- Contiguously allocated region of L * sizeof (T) bytes in memory

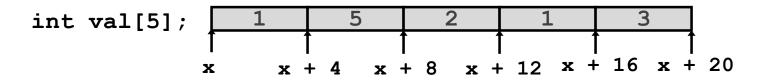


Array Access

■ Basic Principle

```
T \mathbf{A}[L];
```

- Array of data type T and length L
- Identifier A can be used as a pointer to array element 0: Type T*



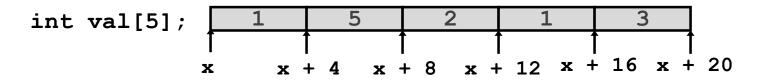
Reference	Type	Value
val[4]	int	3
val	int *	
val+1	int *	
&val[2]	int *	
v al[5]	int	
*(val+1)	int	
val + i	int *	

Array Access

■ Basic Principle

```
T \mathbf{A}[L];
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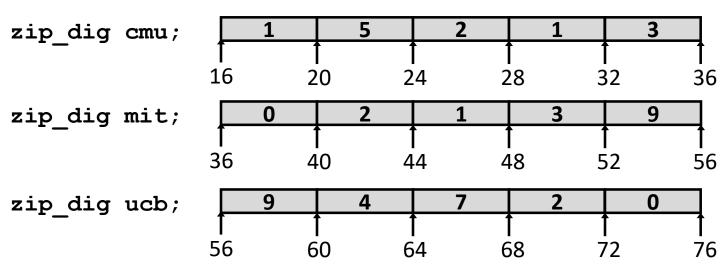


Reference	Type	Value
val[4]	int	3
val	int *	x
val+1	int *	x + 4
&val[2]	int *	x + 8
val [5]	int	??
* (val+1)	int	5 //val[1]
val + <i>i</i>	int *	x + 4 * i //&val[i]

Array Example

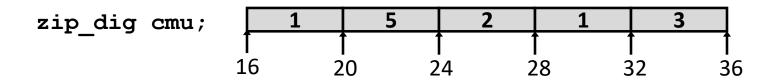
```
#define ZLEN 5
typedef int zip_dig[ZLEN];

zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```



- Declaration "zip dig cmu" equivalent to "int cmu[5]"
- Example arrays were allocated in successive 20 byte blocks
 - Not guaranteed to happen in general

Array Accessing Example



```
int get_digit
  (zip_dig z, int digit)
{
  return z[digit];
}
```

x86-64

```
# %rdi = z
# %rsi = digit
movl (%rdi, %rsi, 4), %eax # z[digit]
```

- Register %rdi contains starting address of array
- Register %rsi contains array index
- Desired digit at %rdi + 4*%rsi
- Use memory reference (%rdi,%rsi,4)

Array Loop Example

```
void zincr(zip_dig z) {
   size_t i;
   for (i = 0; i < ZLEN; i++)
      z[i]++;
}</pre>
```

```
# %rdi = z
                         \# i = 0
 movl $0, %eax
                          # goto middle
 jmp .L3
                         # loop:
.L4:
 addl $1, (%rdi,%rax,4) # z[i]++
 addq $1, %rax
                         # i++
.L3:
                         # middle
 cmpq $4, %rax
                         # i:4
                         # if <=, goto loop</pre>
 jbe .L4
 rep; ret
```

Multidimensional (Nested) Arrays

Declaration

 $T \mathbf{A}[R][C];$

- 2D array of data type T
- R rows, C columns

Array Size

• R * C * sizeof(T) bytes

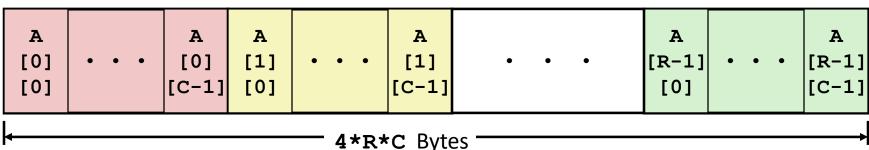
Arrangement

Row-Major Ordering

```
A[0][0] • • • A[0][C-1]

• • • A[R-1][0] • • • A[R-1][C-1]
```

int A[R][C];



Nested Array Example

```
#define PCOUNT 4
 typedef int zip dig[5];
 zip dig pgh[PCOUNT] =
   {{1, 5, 2, 0, 6},
    {1, 5, 2, 1, 3},
    {1, 5, 2, 1, 7},
    {1, 5, 2, 2, 1 }};
zip dig
                                3
             5
                                  1 5
                                         1
                                              1
                                                   2
                  0
                       1
                           2
                                                        1
pgh[4];
         76
                     96
                                116
                                            136
                                                        156
```

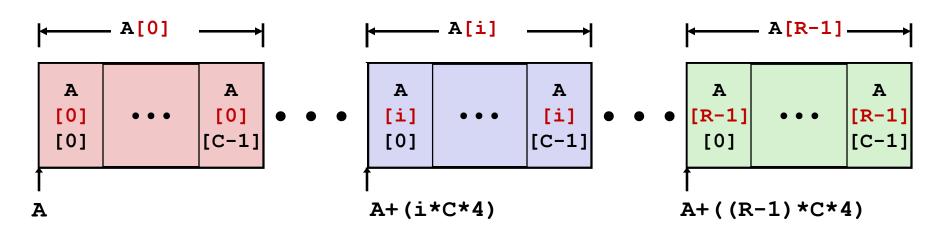
- "zip_dig pgh[4]"equivalent to "int pgh[4][5]"
 - Variable pgh: array of 4 elements, allocated contiguously
 - Each element is an array of 5 int's, allocated contiguously
- "Row-Major" ordering of all elements in memory

Nested Array Row Access

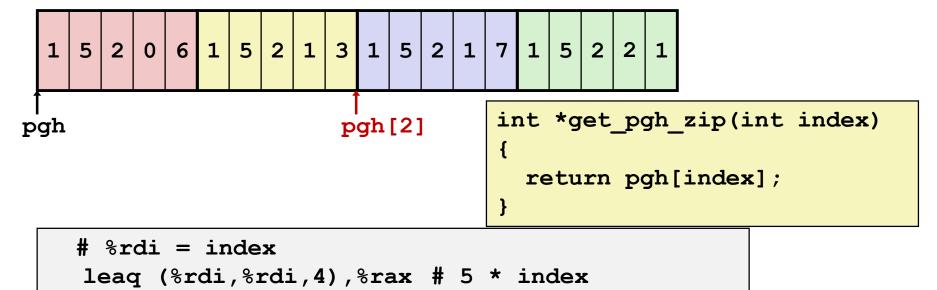
Row Vectors

- **A**[i] is array of *C* elements of type *T*
- Starting address A + i * (C * sizeof(T))

int A[R][C];



Nested Array Row Access Code



Row Vector

- pgh[index] is array of 5 int's
- Starting address pgh+20*index

Machine Code

- Computes and returns address
- Compute as pgh + 4* (index+4*index)

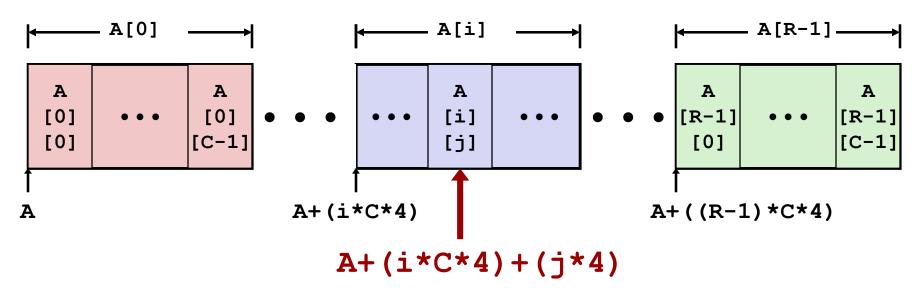
leaq pgh(,%rax,4),%rax # pgh + (20 * index)

Nested Array Element Access

Array Elements

- **A**[i][j] is element of type *T*, which requires *K* bytes
- Address A + i * (C * K) + j * K= A + (i * C + j) * K

int A[R][C];



Nested Array Element Access Code

```
leaq (%rdi,%rdi,4), %rax # 5*index
addl %rax, %rsi # 5*index+dig
movl pgh(,%rsi,4), %eax # M[pgh + 4*(5*index+dig)]
```

Array Elements

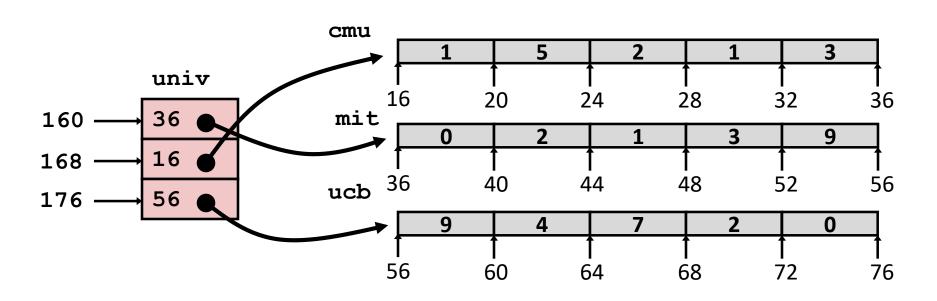
- pgh[index][dig] is int
- Address: pgh + 20*index + 4*dig
 = pgh + 4*(5*index + dig)

Multi-Level Array Example

```
zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```

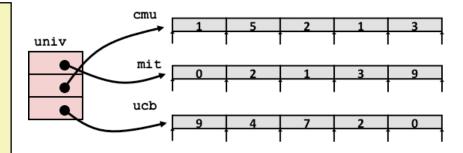
```
#define UCOUNT 3
int *univ[UCOUNT] = {mit, cmu, ucb};
```

- Variable univ denotes array of 3 elements
- Each element is a pointer
 - 8 bytes
- Each pointer points to array of int's



Element Access in Multi-Level Array

```
int get_univ_digit
  (size_t index, size_t digit)
{
  return univ[index][digit];
}
```



```
salq $2, %rsi  # 4*digit
addq univ(,%rdi,8), %rsi # p = univ[index] + 4*digit
movl (%rsi), %eax  # return *p
ret
```

Computation

- Element access Mem [Mem [univ+8*index]+4*digit]
- Must do two memory reads
 - First get pointer to row array
 - Then access element within array

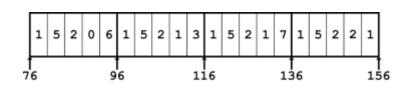
Array Element Accesses

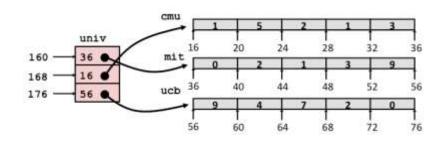
Nested array

```
int get_pgh_digit
   (size_t index, size_t digit)
{
   return pgh[index][digit];
}
```

Multi-level array

```
int get_univ_digit
  (size_t index, size_t digit)
{
  return univ[index][digit];
}
```





Accesses looks similar in C, but address computations very different:

Mem[pgh+20*index+4*digit] Mem[Mem[univ+8*index]+4*digit]

N X N Matrix Code

- Fixed dimensions
 - Know value of N at compile time
- Variable dimensions, explicit indexing
 - Traditional way to implement dynamic arrays
- Variable dimensions, implicit indexing
 - Now supported by gcc

16 X 16 Matrix Access

Array Elements

```
int A[16][16];
Address A + i * (C * K) + j * K

C = 16, K = 4

/* Get element A[i][j] */
int fix_ele(fix_matrix A, size_t i, size_t j) {
  return A[i][j];
}
```

```
# A in %rdi, i in %rsi, j in %rdx
salq $6, %rsi # 64*i
addq %rsi, %rdi # A + 64*i
movl (%rdi,%rdx,4), %eax # M[A + 64*i + 4*j]
ret
```

n X n Matrix Access

Array Elements

```
size_t n;
int A[n][n];
Address A + i * (C * K) + j * K
C = n, K = 4
```

Must perform integer multiplication

```
/* Get element A[i][j] */
int var_ele(size_t n, int A[n][n], size_t i, size_t j)
{
  return A[i][j];
}
```

```
# n in %rdi, A in %rsi, i in %rdx, j in %rcx
imulq %rdx, %rdi  # n*i
leaq (%rsi, %rdi, 4), %rax # A + 4*n*i
movl (%rax, %rcx, 4), %eax # A + 4*n*i + 4*j
ret
```

Example: Array Access

```
#include <stdio.h>
#define ZLEN 5
#define PCOUNT 4
typedef int zip dig[ZLEN];
int main(int argc, char** argv) {
zip dig pgh[PCOUNT] =
    \{\{1, 5, 2, 0, 6\},
    {1, 5, 2, 1, 3},
    {1, 5, 2, 1, 7},
    {1, 5, 2, 2, 1 }};
    int *linear zip = (int *) pgh;
    int *zip2 = (int *) pgh[2];
    int result =
       pgh[0][0] +
       linear zip[7] +
        *(linear zip + 8) +
        zip2[1];
   printf("result: %d\n", result);
    return 0;
```

```
linux> ./array
result: 9
```

Example: Array Access

```
#include <stdio.h>
#define ZLEN 5
#define PCOUNT 4
typedef int zip dig[ZLEN];
int main(int argc, char** argv) {
zip dig pgh[PCOUNT] =
    \{\{1, 5, 2, 0, 6\},
    \{1, 5, 2, 1, 3\},\
    {1, 5, 2, 1, 7},
    {1, 5, 2, 2, 1 }};
    int *linear zip = (int *) pgh;
    int *zip2 = (int *) pgh[2];
    int result =
       pgh[0][0] +
        linear zip[7] +
        *(linear zip + 8) +
        zip2[1];
    printf("result: %d\n", result);
    return 0;
```

```
linux> ./array
result: 9
```

Today

Arrays

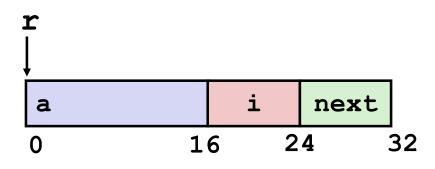
- One-dimensional
- Multi-dimensional (nested)
- Multi-level

Structures

- Allocation
- Access
- Alignment
- Floating Point

Structure Representation

```
struct rec {
   int a[4];
   size_t i;
   struct rec *next;
};
```



- Structure represented as block of memory
 - Big enough to hold all of the fields
- Fields ordered according to declaration
 - Even if another ordering could yield a more compact representation
- Compiler determines overall size + positions of fields
 - Machine-level program has no understanding of the structures in the source code

Generating Pointer to Structure Member

```
struct rec {
   int a[4];
   size_t i;
   struct rec *next;
};
```

```
r r+4*idx
a i next
0 16 24 32
```

Generating Pointer to Array Element

- Offset of each structure member determined at compile time
- Compute as r + 4*idx

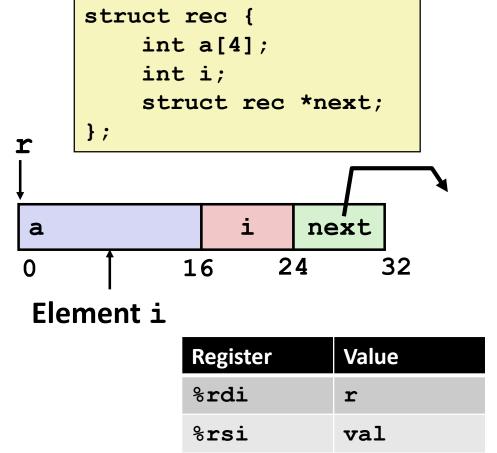
```
int *get_ap
  (struct rec *r, size_t idx)
{
  return &r->a[idx];
}
```

```
# r in %rdi, idx in %rsi
leaq (%rdi,%rsi,4), %rax
ret
```

Following Linked List

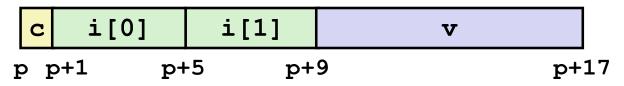
C Code

```
void set_val
  (struct rec *r, int val)
{
  while (r) {
    int i = r->i;
    r->a[i] = val;
    r = r->next;
  }
}
```



Structures & Alignment

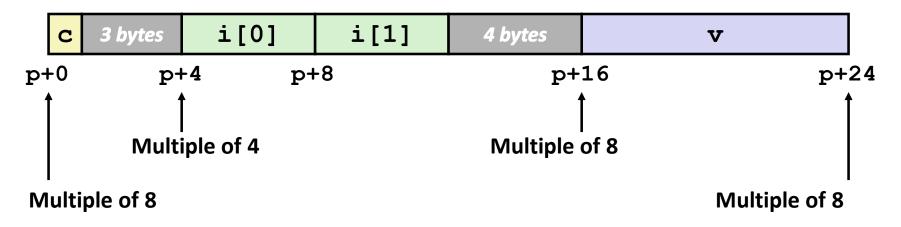
Unaligned Data



```
struct S1 {
  char c;
  int i[2];
  double v;
} *p;
```

Aligned Data

- Primitive data type requires K bytes
- Address must be multiple of K



Alignment Principles

Aligned Data

- Primitive data type requires K bytes
- Address must be multiple of *K*
- Required on some machines; advised on x86-64

Motivation for Aligning Data

- Memory accessed by (aligned) chunks of 4 or 8 bytes (system dependent)
 - Inefficient to load or store datum that spans cache lines (64 bytes). Intel states should avoid crossing 16 byte boundaries.
 - Virtual memory trickier when datum spans 2 pages (4 KB pages)

Compiler

Inserts gaps in structure to ensure correct alignment of fields

Specific Cases of Alignment (x86-64)

- 1 byte: char, ...
 - no restrictions on address
- 2 bytes: short, ...
 - lowest 1 bit of address must be 02
- 4 bytes: int, float, ...
 - lowest 2 bits of address must be 002
- 8 bytes: double, long, char *, ...
 - lowest 3 bits of address must be 0002

Satisfying Alignment with Structures

Within structure:

Must satisfy each element's alignment requirement

Overall structure placement

- Each structure has alignment requirement K
 - **K** = Largest alignment of any element
- Initial address & structure length must be multiples of K

Example:

K = 8, due to double element

```
        c
        3 bytes
        i [0]
        i [1]
        4 bytes
        v

        p+0
        p+4
        p+8
        p+16
        p+24

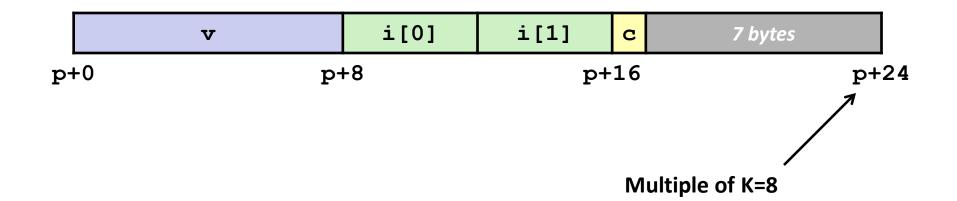
        Multiple of 4
        Multiple of 8
        Multiple of 8

Multiple of 8
```

Meeting Overall Alignment Requirement

- For largest alignment requirement K
- Overall structure must be multiple of K

```
struct S2 {
  double v;
  int i[2];
  char c;
} *p;
```

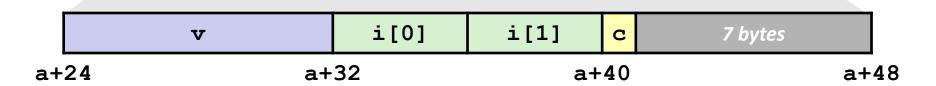


Arrays of Structures

- Overall structure length multiple of K
- Satisfy alignment requirement for every element

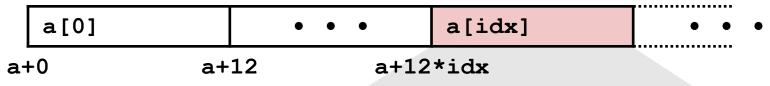
```
struct S2 {
  double v;
  int i[2];
  char c;
} a[10];
```





Accessing Array Elements

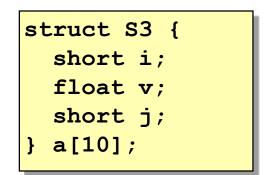
- Compute array offset 12*idx
 - sizeof (S3), including alignment spacers
- Element j is at offset 8 within structure
- Assembler gives offset a+8
 - Resolved during linking





```
short get_j(int idx)
{
  return a[idx].j;
}
```

```
# %rdi = idx
leaq (%rdi,%rdi,2),%rax # 3*idx
movzwl a+8(,%rax,4),%eax
```



Saving Space

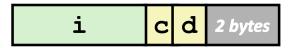
Put large data types first

```
struct S4 {
   char c;
   int i;
   char d;
} *p;
struct S5 {
   int i;
   char c;
   char d;
} *p;

c 3 bytes

i d 3 bytes
```

Effect (largest alignment requirement K=4)



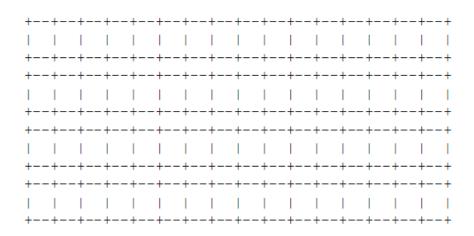
Example Struct Exam Question

Problem 5. (8 points):

Struct alignment. Consider the following C struct declaration:

```
typedef struct {
  char a;
  long b;
  float c;
  char d[3];
  int *e;
  short *f;
} foo;
```

 Show how foo would be allocated in memory on an x86-64 Linux system. Label the bytes with the names of the various fields and clearly mark the end of the struct. Use an X to denote space that is allocated in the struct as padding.



http://www.cs.cmu.edu/~213/oldexams/exam1-f12.pdf

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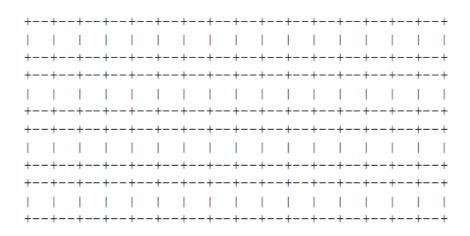
Example Struct Exam Question (Cont'd)

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typedef struct {
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Rearrange the elements of foo to conserve the most space in memory. Label the bytes with the names of the various fields and clearly mark the end of the struct. Use an X to denote space that is allocated in the struct as padding.



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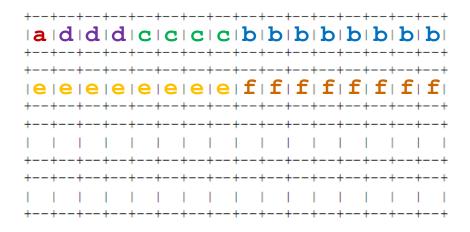
Example Struct Exam Question (Cont'd)

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Structures

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- Access
- Alignment

Floating Point

Background

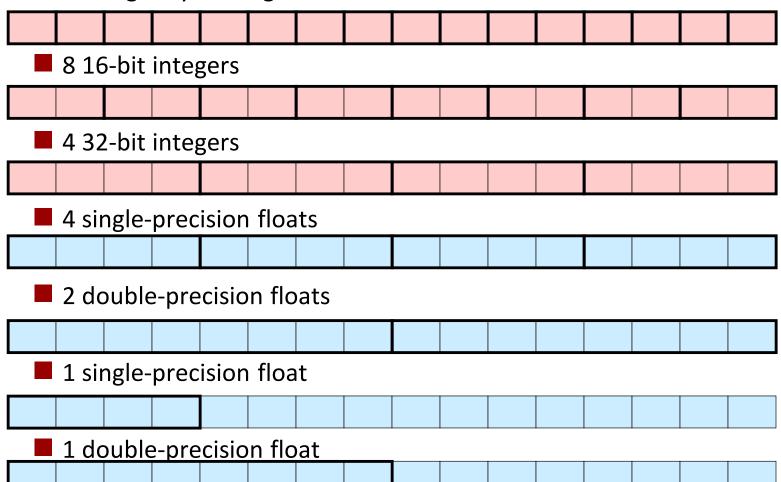
History

- **x87 FP**
 - Legacy, very ugly
- SSE FP
 - Supported by Shark machines
 - Special case use of vector instructions
- AVX FP
 - Newest version
 - Similar to SSE (but registers are 32 bytes instead of 16)
 - Documented in book

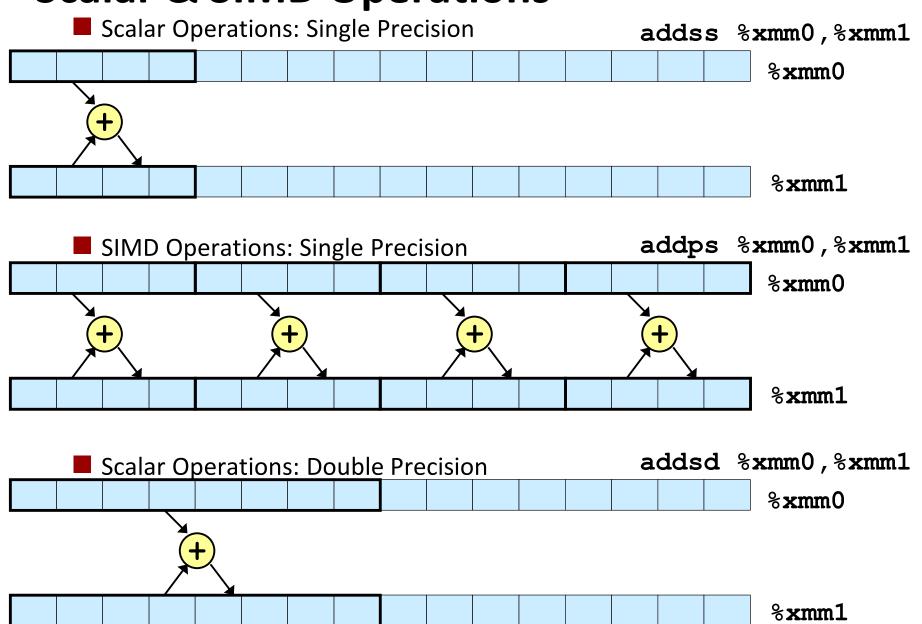
Programming with SSE3

XMM Registers

- 16 total, each 16 bytes
- 16 single-byte integers



Scalar & SIMD Operations



FP Basics

- Arguments passed in %xmm0, %xmm1, ...
- Result returned in %xmm0
- All XMM registers caller-saved

```
float fadd(float x, float y)
{
    return x + y;
}
```

```
double dadd(double x, double y)
{
   return x + y;
}
```

```
# x in %xmm0, y in %xmm1
addss %xmm1, %xmm0
ret
```

```
# x in %xmm0, y in %xmm1
addsd %xmm1, %xmm0
ret
```

FP Memory Referencing

- Integer (and pointer) arguments passed in regular registers
- FP values passed in XMM registers
- Different mov instructions to move between XMM registers,
 and between memory and XMM registers

```
double dincr(double *p, double v)
{
    double x = *p;
    *p = x + v;
    return x;
}
```

```
# p in %rdi, v in %xmm0
movapd %xmm0, %xmm1  # Copy v
movsd (%rdi), %xmm0  # x = *p
addsd %xmm0, %xmm1  # t = x + v
movsd %xmm1, (%rdi) # *p = t
ret
```

Other Aspects of FP Code

Lots of instructions

Different operations, different formats, ...

Floating-point comparisons

- Instructions ucomiss and ucomisd
- Set condition codes CF, ZF, and PF

Using constant values

- Set XMM0 register to 0 with instruction xorpd %xmm0, %xmm0
- Others loaded from memory

Summary

Arrays

- Elements packed into contiguous region of memory
- Use index arithmetic to locate individual elements

Structures

- Elements packed into single region of memory
- Access using offsets determined by compiler
- Possible require internal and external padding to ensure alignment

Combinations

Can nest structure and array code arbitrarily

Floating Point

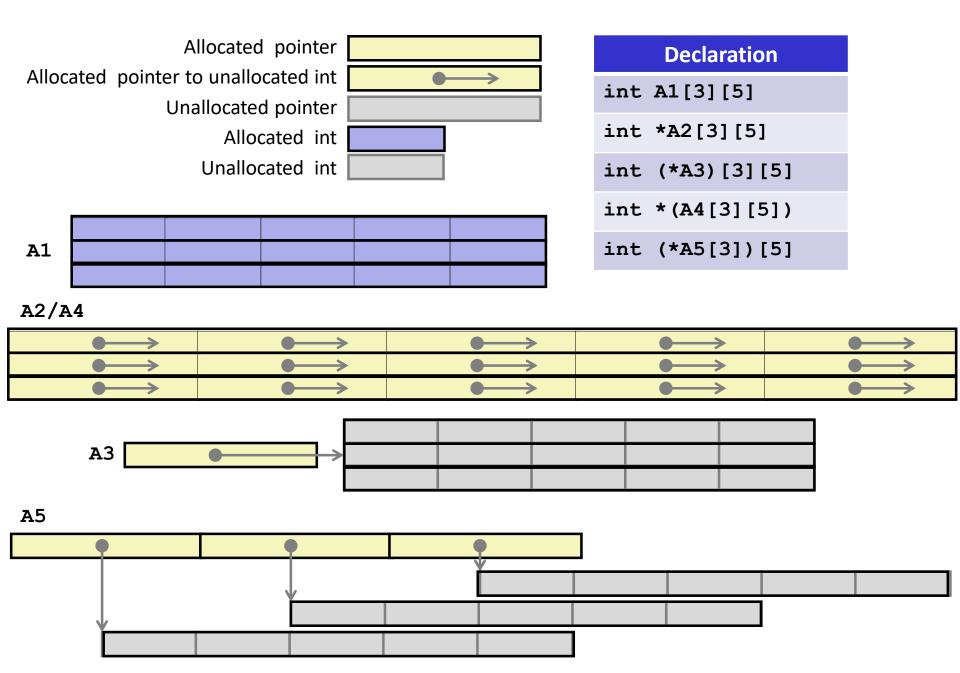
Data held and operated on in XMM registers

Understanding Pointers & Arrays #3

Decl	An			*An			**An		
	Cmp	Bad	Size	Cmp	Bad	Size	Cmp	Bad	Size
int A1[3][5]									
int *A2[3][5]									
int (*A3)[3][5]									
int *(A4[3][5])									
int (*A5[3])[5]									

- Cmp: Compiles (Y/N)
- Bad: Possible bad pointer reference (Y/N)
- Size: Value returned by sizeof

Decl	***An		
	Cmp	Bad	Size
int A1[3][5]			
int *A2[3][5]			
int (*A3)[3][5]			
int *(A4[3][5])			
int (*A5[3])[5]			



Understanding Pointers & Arrays #3

Decl	An			*An			**An		
	Cmp	Bad	Size	Cmp	Bad	Size	Cmp	Bad	Size
int A1[3][5]	Y	N	60	Y	N	20	Y	N	4
int *A2[3][5]	Y	N	120	Y	N	40	Y	N	8
int (*A3)[3][5]	Y	N	8	Y	Y	60	Y	Y	20
int *(A4[3][5])	Y	N	120	Y	N	40	Y	N	8
int (*A5[3])[5]	Y	N	24	Y	N	8	Y	Y	20

- Cmp: Compiles (Y/N)
- Bad: Possible bad pointer reference (Y/N)
- Size: Value returned by sizeof

Decl	***An				
	Cmp	Bad	Size		
int A1[3][5]	N	-	-		
int *A2[3][5]	Y	Y	4		
int (*A3)[3][5]	Y	Y	4		
int *(A4[3][5])	Y	Y	4		
int (*A5[3])[5]	Y	Y	4		