Cours Algorithmique des systèmes parallèles et distribués Exercices

Série :PlusCal pour la programmation répartie ou concurrente (I) par Alessio Coltellacci et Dominique Méry 5 mai 2025

$\textbf{Exercice 1} \ \textit{Etudier le programme Plus Cal suivant}:$

```
----- MODULE pluscaltut1 ------
EXTENDS Integers, Sequences, TLC, FiniteSets
(*
−−wf
--algorithm Tut1 {
variables x = 0;
process (one = 1)
  A: assert x \setminus in \{0,1\};
    x := x - 1;
  B: \ assert \ x \setminus in \{-1,0\};
    x := x * 3;
   BB: assert x \in \{-3, -2, 0, 1\};
};
process (two = 2)
  C: assert x \setminus in \{-3, -2, -1, 0, 1\};
    x := x + 1;
 D:
    assert x \setminus in \{-2, -1, 0, 1, 2\};
};
end algorithm;
\* BEGIN TRANSLATION (chksum(pcal) = "a26cf6c4" /\ chksum(tla) = "72d02274")
VARIABLES x, pc
vars == << x, pc >>
ProcSet == \{1\} \setminus cup \{2\}
Init == (* Global variables *)
```

```
/\setminus x = 0
        [] self = 2 -> "C"]
A == / pc[1] = "A"
     / \setminus x' = x - 1
     / \ pc' = [pc EXCEPT ![1] = "B"]
B == / \setminus pc[1] = "B"
     / \setminus x' = x * 3
     / \ pc' = [pc EXCEPT ![1] = "BB"]
BB == / pc[1] = "BB"
      / \land Assert(x \land in \{-3, -2, 0, 1\},\
                "Failure of assertion at line 15, column 8.")
      / \ pc' = [pc EXCEPT ![1] = "Done"]
      / \setminus x' = x
one == A \setminus / B \setminus / BB
C == / \cdot pc[2] = "C"
     / \land Assert(x \land in \{-3, -2, -1, 0, 1\},\
               "Failure of assertion at line 20, column 6.")
     / \setminus x' = x + 1
     / \ pc' = [pc EXCEPT ![2] = "D"]
D == / \cdot pc/2! = "D"
     / \land Assert(x \land in \{-2, -1, 0, 1, 2\},\
               "Failure of assertion at line 23, column 5.")
     / \ pc' = [pc \ EXCEPT' ! [2] = "Done"]
     / \setminus x' = x
two == C \setminus / D
(* Allow infinite stuttering to prevent deadlock on termination. *)
Terminating == /\ \A self \in ProcSet: pc[self] = "Done"
               /\ UNCHANGED vars
Next == one \setminus / two
           \/ Terminating
Spec == Init / [][Next]_vars
```

```
Termination == <>(\A self \in ProcSet: pc[self] = "Done")
\∗ END TRANSLATION
\textbf{Exercice 2} \ \textit{Etudier le programme Plus Cal suivant}:
  ----- MODULE pluscaltut2 ------
EXTENDS Integers, Sequences, TLC, FiniteSets
--algorithm Tut2 {
variables x = 0;
process (one = 1)
variables temp
A:
       temp := x + 1;
       x := temp;
};
process (two = 2)
variables temp
 CC:
       temp := x + 1;
       x := temp;
};
end algorithm;
\* Process variable temp of process one at line 10 col 11 changed to temp_
CONSTANT\ defaultInitValue
VARIABLES x, pc, temp_, temp
```

```
vars == \langle \langle x, pc, temp_{-}, temp_{-} \rangle
ProcSet == \{1\} \setminus cup \{2\}
Init == (* Global variables *)
          /\setminus x = 0
          (* Process one *)
          / \cdot temp_{-} = defaultInitValue
          (* Process two *)
          / \setminus temp = defaultInitValue
          /\ pc = [self \ in \ ProcSet \ | -> CASE \ self = 1 \ -> "A"]
                                                 [] self = 2 \rightarrow "CC"]
A == / \setminus pc[1] = "A"
      / \cdot temp_{-}' = x + 1
      / \setminus x' = temp_{-}'
      / \ pc' = [pc EXCEPT ![1] = "Done"]
      / \setminus temp' = temp
one == A
CC == / \cdot pc[2] = "CC"
       / \setminus temp' = x + 1
       / \setminus x' = temp'
       /\ pc' = [pc EXCEPT ![2] = "Done"]
       / \setminus temp_{-}' = temp_{-}
two == CC
(* Allow infinite stuttering to prevent deadlock on termination. *)
Terminating == /\ \A self \in ProcSet: pc[self] = "Done"
                  /\ UNCHANGED vars
Next == one \setminus / two
             \/ Terminating
Spec == Init / [][Next]_vars
Termination == <>(\A self \in ProcSet: pc[self] = "Done")
\* END TRANSLATION
test == (\A i \ \n \ ProcSet : pc[i] = "Done") => x \ \n \{2\}
```

Exercice 3 Etudier le programme PlusCal suivant :

```
----- MODULE pluscaltut3 -----
EXTENDS Integers, Sequences, TLC, FiniteSets
--algorithm Tut3 {
variables x = 0;
process (one = 1)
 A:
    x := x + 1;
  B:
    await \quad x = 1;
    print << "x = ", x>>;
};
process (two = 2)
 D:
   await x = 1;
   assert x = 1;
  F:
   x := x -2;
};
end algorithm;
\* BEGIN TRANSLATION
VARIABLES x, pc
vars == \langle \langle x, pc \rangle \rangle
ProcSet == \{1\} \setminus cup \{2\}
Init == (* Global variables *)
        /\ pc = [self \ in \ ProcSet \ | -> CASE \ self = 1 \ -> "A"
```

```
[] self = 2 -> "D"]
A == / \setminus pc[1] = "A"
      /\setminus x' = x + 1
      / \ pc' = [pc EXCEPT ![1] = "B"]
B == / \setminus pc[1] = "B"
      / \setminus x = 1
      / \ pc' = [pc EXCEPT ![1] = "C"]
      / \setminus x' = x
C == / \cdot pc[1] = "C"
      / \land PrintT(<<"x=",x>>)
      / \ pc' = [pc EXCEPT ![1] = "Done"]
      /\setminus x' = x
one == A \setminus / B \setminus / C
D == / \setminus pc[2] = "D"
      / \setminus x = 1
      / \ pc' = [pc EXCEPT ![2] = "E"]
      / \setminus x' = x
E == / \cdot pc[2] = "E"
      \ / \ Assert(x = 1, "Failure of assertion at line 22, column 5.")
      / \ pc' = [pc EXCEPT ![2] = "F"]
      / \setminus x' = x
F == / pc/2! = "F"
      / \setminus x' = x -2
      /\ pc' = [pc EXCEPT ![2] = "Done"]
two == D \setminus / E \setminus / F
(* Allow infinite stuttering to prevent deadlock on termination. *)
Terminating == / \ A \ self \ in \ ProcSet: pc[self] = "Done"
                   /\ UNCHANGED vars
Next == one \setminus / two
              \/ Terminating
Spec == Init / [][Next]_vars
```

 $Termination == <>(\A self \in ProcSet: pc[self] = "Done")$

***** END TRANSLATION

 $test == (\A i \in ProcSet : pc[i]="Done") => x \in \{1, 2\}$

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Exercice 4 pluscalex1.tla

Ecrire un programme Plus Cal qui traduit le protocole suivant : S envoie une valeur à ${\bf R}$

Exercice 5 pluscalex2.tla

 $\label{lem:controlle} \textit{Ecrire un programme PlusCal qui calcule la fonction factorielle de la façon suivante:}$

- *Un processus calcule* $1 \times 2 \times 3.... \times k1$
- Un processus calcule $k2 \times (k2+1) \times ... \times N$
- Les processus stoppent quand la condition k1 < k2 est fausse

Exercice 6 pluscalex3.tla

Ecrire un programme PlusCal qui calcule la fonction L^K la façon suivante :

- Un processus calcule $L \times ... \times L \ k1$ fois.
- Un processus calcule $L \times ... \times L$ k2 fois.
- Les processus stoppent quand la condition k1+K2 < L est fausse