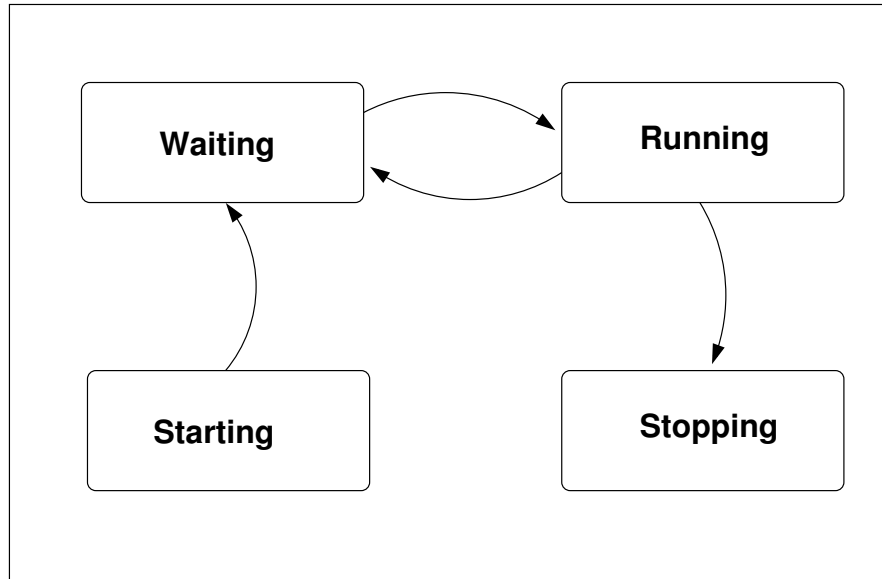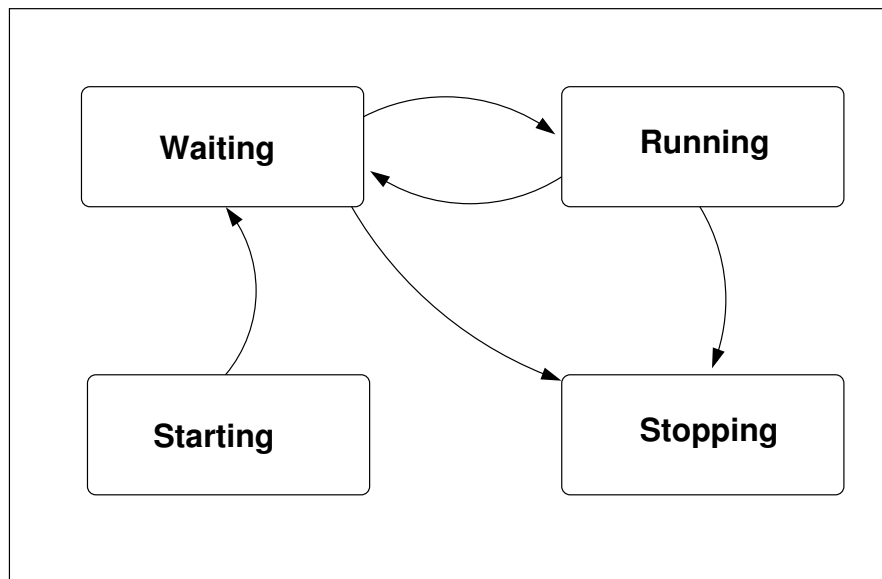**Exercice 1** *ex1-tut1.zip*

Express the following states machine using an Event B machines and check properties on the resulting models.

Exercice 2 *ex2-tut1.zip*

Express the following states machine using an Event B machines and check properties on the resulting models.

Exercice 3 *(ex3-simple-tut.zip)*

Soient trois ensembles A, B et qui sont des parties de U.

    — *Ecrire un modèle Event-B qui utilise deux variables v et w deux sous-ensembles de A et B*

— *Ajouter une fonction partielle de A dans B.*
— *Définir un événement ⊖1 qui enlève un élément de A qui est élément de C.*
— *Définir un événement ⊖2 qui enlève un élément de AB qui est élément de C.*
— *Définir un événement ⊖3 qui crée un lien entre un élément de v et un élément de w à condition que v et w ne contiennet plus d'éléments de C.*

Exercice 4 *ex4-tut1.zip*
We consider a finite sequence of integers $v_1, \ldots, v_n$ where n is the length of the sequence and is supposed to be fixed. Write an Event B specification modelling the computation of the value of the summation of the sequence v. You should define cerafully v, n and the summation of a finite sequence of integers.

Exercice 5 *ex51-tut1.zip and ex52-tut1.zip*
Express the following property in Event B :
— *(ex51-tut1.zip) We assume to have p resources which may be shared by n processes. If a process uses a given resource, the resource can not be used by another process. A process can use only at most one resource.*
— *(ex52-tut1.zip) We assume to have p resources which may be shared by $nmcs$ processes. If a process uses a given resource, the resource can not be used by another process. A process can use possibly more than one ressoyrce. .*

Exercice 6 *ex6-tut1.zip*
A Petri net is a uple R=(S,T,F,K,M,W)
— *S is a finite set of places.*
— *T is a finite set of transitions.*
— *S ∩ T = ∅*
— *F is the flow relation : $F \subseteq S \times T \cup T \times S$*
— *K is expressing the capacity of each place :*
   $K \in S \to Nat \cup \{\omega\}$
— *M is reprenting the initial marking of each place :*
   $M \in S \to Nat \cup \{\omega\}$ and satisfies the following condition $\forall s \in S : M(s) \leq K(s)$.
— *W is the weight of each edge :*
   $W \in F \to Nat \cup \{\omega\}$
THe state of a Petri net R is defined by a set of markings :
— *a marking M for R is a function from S to $Nat \cup \{\omega\}$:*
   $M \in S \to Nat \cup \{\omega\}$ and it satisfies the condition $\forall s \in S : M(s) \leq K(s)$.
— *a transition t of T is ready to fire for a marking M of R, if*

   1. $\forall s \in \{ s' \in S \mid (s',t) \in F \}$:
      M(s) ≥ W(s,t).

   2. $\forall s \in \{ s' \in S \mid (t,s') \in F \}$:
      M(s) ≤ K(s) - W(s,t).

— *$t \in T : Pre(t) = \{s' \in S : (s',t) \in F\}$ and $Post(t) = \{s' \in S : (t,s') \in F\}$*
The simulation of a Petri net is defined by a relation linking three elements : a marking M, a marking M' and a transition t as follows :
— *the new marking M' is defined as follows from M :*

$$M'(s) = \begin{cases} \text{M(S)−W(S,T), SI S} \in \text{PRE(T) - POST(T)} \\ \text{M(S)+W(T,S), SI S} \in \text{POST(T) - PRE(T)} \\ \text{M(S)−W(S,T)+W(T,S), SI S} \in \text{PRE(T)∩POST(T)} \\ \text{M(S), SINON} \end{cases}$$

We consider the following Petri net :

Question 6.1 *Translate this Petri net in Event B.*

Question 6.2 *Express safety properties that you can discover from the diagram.*

Exercice 7 *(ex7-tut1.zip)*
Nous considérons le modèle suivant.

```
MACHINE M1
VARIABLES
      x
INVARIANTS
. . .
EVENTS
EVENT INITIALISATION
  BEGIN
  act1 : x := −10
  END
  EVENT evt1
  WHEN
  grd1 : x ≥ −1
  THEN
  act1 : x := x+1
  END
  EVENT evt2
  WHEN
  grd1 : x ≤ −1
  grd2 : x ≥ −44
  THEN
  act1: x := x−1
  END
END
```

On considère plusieurs cas pour l'invariant.

Question 7.1 *(M1)*

$$inv1 : x \in \mathbb{Z}$$
$$inv3 : x \leq -1$$

Est-ce que toutes les conditions de vérification sont prouvées par le prouveur de l'application Rodin ? Expliquez clairement pourquoi elles sont prouvées ou non.

Question 7.2 *(M2)*

$$inv1 : x \in \mathbb{Z}$$
$$inv3 : x \leq -3$$

Est-ce que toutes les conditions de vérification sont prouvées par le prouveur de l'application Rodin ? Expliquez clairement pourquoi elles sont prouvées ou non.

Question 7.3 *(M3)*

$$inv1 : x \in \mathbb{Z}$$
$$inv4 : -45 \leq x \wedge x \leq -10$$

Est-ce que toutes les conditions de vérification sont prouvées par le prouveur de l'application Rodin ? Expliquez clairement pourquoi elles sont prouvées ou non.

Question 7.4 *(M4)*

$$inv1 : x \in \mathbb{Z}$$
$$inv3 : x \leq -3$$
$$inv4 : -45 \leq x \wedge x \leq -10$$
$$inv2 : x \leq -1$$

Est-ce que toutes les conditions de vérification sont prouvées par le prouveur de l'application Rodin ? Expliquez clairement pourquoi elles sont prouvées ou non.

Exercice 8 *ex8-tut1.zip*

A semaphore s is a shared variable accessible by two operations : $P(s)$ and $V(s$. Informally, we can describe the effect of these two operations as follows :

  — *$P(s)$ is testing if the value of s is greater than 0 and is not equal to 0. Il the value of s is 0, the process which is executing $P(s)$ is inserted in a queue.*
  — *$V(s)$ is increasing the value of s by one, if the queue is non empty. If the queue is non empty, the first waiting process of the queue is awaken and becomes a lively process.*

Using the Event B modelling features, describe a system using the primitives.

Exercice 9 *ex9-tut1.zip*

We assume that two $n{\times}n$ matrices of boolean values are given : A and B.
Write an Event B specification modelling the multiplication of the two matrcices.

Exercice 10 *(ex10-1-tut1.zip and ex10-2-tut1.zip)*

A system is used to sum two numbers x0 and y0 by adding one unit to a variable z. It includes an incx2z event which decreases the value of x by one and increases the value of z by one, and an incy2z event which decreases y by one and increases z by one. The overall process stops when the two variables x and y are zero.
Write a model in Event-B for the syste.