

Cours Algorithmique des systèmes parallèles et distribués
Exercices

Série :PlusCal pour la programmation répartie ou concurrente (I)
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Exercice 1 Etudier le programme PlusCal suivant :

```
----- MODULE pluscaltut1 -----
EXTENDS Integers , Sequences , TLC, FiniteSets
(*
--wf
--algorithm Tut1 {
variables x = 0;

process (one = 1)
{
  A: assert x \in {0,1};
  x := x - 1;
  B: assert x \in {-1,0} ;
  x := x * 3;
  BB: assert x \in {-3,-2,0,1};
};

process (two = 2)
{
  C: assert x \in {-3,-2,-1,0,1};
  x := x + 1;
  D:
    assert x \in {-2,-1,0,1,2};
};

}
end algorithm;

*)
/* BEGIN TRANSLATION (chksum(pcal) = "a26cf6c4" /\ chksum(tla) = "72d02274")
VARIABLES x, pc

vars == << x, pc >>
ProcSet == {1} \cup {2}
Init == (* Global variables *)
```

```

/\ x = 0
/\ pc = [self \in ProcSet |-> CASE self = 1 -> "A"
          [] self = 2 -> "C"]

A == /\ pc[1] = "A"
      /\ Assert(x \in {0,1}, "Failure of assertion at line 11, column 6.")
      /\ x' = x - 1
      /\ pc' = [pc EXCEPT ![1] = "B"]

B == /\ pc[1] = "B"
      /\ Assert(x \in {-1,0}, "Failure of assertion at line 13, column 6.")
      /\ x' = x * 3
      /\ pc' = [pc EXCEPT ![1] = "BB"]

BB == /\ pc[1] = "BB"
      /\ Assert(x \in {-3,-2,0,1},
                 "Failure of assertion at line 15, column 8.")
      /\ pc' = [pc EXCEPT ![1] = "Done"]
      /\ x' = x

one == A \ / B \ / BB

C == /\ pc[2] = "C"
      /\ Assert(x \in {-3,-2,-1,0,1},
                 "Failure of assertion at line 20, column 6.")
      /\ x' = x + 1
      /\ pc' = [pc EXCEPT ![2] = "D"]

D == /\ pc[2] = "D"
      /\ Assert(x \in {-2,-1,0,1,2},
                 "Failure of assertion at line 23, column 5.")
      /\ pc' = [pc EXCEPT ![2] = "Done"]
      /\ x' = x

two == C \ / D

(* Allow infinite stuttering to prevent deadlock on termination. *)
Terminating == /\ \A self \in ProcSet: pc[self] = "Done"
               /\ UNCHANGED vars

Next == one \ / two
       \ / Terminating

Spec == Init /\ [] [Next]_vars

```

```
Termination == <>(&self \in ProcSet : pc[self] = "Done")
```

```
\* END TRANSLATION
```

```
=====
```

Exercice 2 Etudier le programme PlusCal suivant :

```
----- MODULE pluscaltut2 -----
EXTENDS Integers, Sequences, TLC, FiniteSets

(*
--algorithm Tut2 {
variables x = 0;

process (one = 1)

variables temp
{

A:
    temp := x + 1;

    x := temp;

};

process (two = 2)

variables temp
{
    CC:
        temp := x + 1;

        x:= temp;

};

}
end algorithm;

*)
\* BEGIN TRANSLATION (chksum(pcal) = "b54fa406" /\ chksum(tla) = "e84b4125")
\* Process variable temp of process one at line 10 col 11 changed to temp_
CONSTANT defaultInitValue
VARIABLES x, pc, temp_, temp
```

```

vars == << x, pc, temp_, temp >>

ProcSet == {1} \cup {2}

Init == (* Global variables *)
  /\ x = 0
  (* Process one *)
  /\ temp_ = defaultInitValue
  (* Process two *)
  /\ temp = defaultInitValue
  /\ pc = [self \in ProcSet |-> CASE self = 1 -> "A"
           [] self = 2 -> "CC"]

A == /\ pc[1] = "A"
      /\ temp_ = x + 1
      /\ x' = temp_
      /\ pc' = [pc EXCEPT ![1] = "Done"]
      /\ temp' = temp

one == A

CC == /\ pc[2] = "CC"
      /\ temp' = x + 1
      /\ x' = temp'
      /\ pc' = [pc EXCEPT ![2] = "Done"]
      /\ temp_ = temp

two == CC

(* Allow infinite stuttering to prevent deadlock on termination. *)
Terminating == /\ \A self \in ProcSet: pc[self] = "Done"
               /\ UNCHANGED vars

Next == one \ / two
        \ / Terminating

Spec == Init /\ [] [Next]_vars

Termination == <>(\A self \in ProcSet: pc[self] = "Done")

\* END TRANSLATION

test == (\A i \in ProcSet : pc[i] = "Done") => x \in {2}

```

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Exercice 3 Etudier le programme PlusCal suivant :

```
----- MODULE pluscaltut3 -----
EXTENDS Integers, Sequences, TLC, FiniteSets
(*
--algorithm Tut3 {
variables x = 0;

process (one = 1)
{
  A:
    x := x + 1;
  B:
    await x = 1;
  C:
    print <<"x=",x>>;
};

process (two = 2)
{
  D:
    await x = 1;
  E:
    assert x = 1;
  F:
    x := x -2;
};

}
end algorithm;

*)
\* BEGIN TRANSLATION
VARIABLES x, pc

vars == << x, pc >>

ProcSet == {1} \cup {2}

Init == (* Global variables *)
  /\ x = 0
  /\ pc = [self \in ProcSet |-> CASE self = 1 -> "A"
```

```

[] self = 2 -> "D"]

A == /\ pc[1] = "A"
      /\ x' = x + 1
      /\ pc' = [pc EXCEPT !(1) = "B"]

B == /\ pc[1] = "B"
      /\ x = 1
      /\ pc' = [pc EXCEPT !(1) = "C"]
      /\ x' = x

C == /\ pc[1] = "C"
      /\ PrintT(<<"x=",x>>)
      /\ pc' = [pc EXCEPT !(1) = "Done"]
      /\ x' = x

one == A \ / B \ / C

D == /\ pc[2] = "D"
      /\ x = 1
      /\ pc' = [pc EXCEPT !(2) = "E"]
      /\ x' = x

E == /\ pc[2] = "E"
      /\ Assert(x = 1, "Failure of assertion at line 22, column 5.")
      /\ pc' = [pc EXCEPT !(2) = "F"]
      /\ x' = x

F == /\ pc[2] = "F"
      /\ x' = x -2
      /\ pc' = [pc EXCEPT !(2) = "Done"]

two == D \ / E \ / F

(* Allow infinite stuttering to prevent deadlock on termination. *)
Terminating == /\ \A self \in ProcSet: pc[self] = "Done"
                  /\ UNCHANGED vars

Next == one \ / two
                  \ / Terminating

Spec == Init /\ [] [Next]_vars

Termination == <>(\A self \in ProcSet: pc[self] = "Done")
```

* END TRANSLATION

test == ($\forall i \in ProcSet : pc[i] = "Done"$) => x \in \{1, 2\}

=====

Exercice 4 pluscalex1.tla

Ecrire un programme PlusCal qui traduit le protocole suivant : S envoie une valeur à R

Exercice 5 pluscalex2.tla

Ecrire un programme PlusCal qui calcule la fonction factorielle de la façon suivante :

- Un processus calcule $1 \times 2 \times 3 \dots \times k_1$
- Un processus calcule $k_2 \times (k_2 + 1) \times \dots \times N$
- Les processus stoppent quand la condition $k_1 < k_2$ est fausse

Exercice 6 pluscalex3.tla

Ecrire un programme PlusCal qui calcule la fonction L^K la façon suivante :

- Un processus calcule $L \times \dots \times L$ k_1 fois.
- Un processus calcule $L \times \dots \times L$ k_2 fois.
- Les processus stoppent quand la condition $k_1 + K_2 < L$ est fausse