



Cours MALG & MOVEX

Vérification mécanisée de contrats (II) (The ANSI/ISO C Specification Language (ACSL))

Dominique Méry Telecom Nancy, Université de Lorraine

Année universitaire 2024-2025

1 Contracts

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Logic Specification
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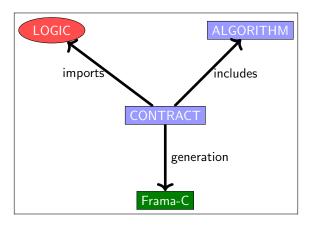
Outline

1 Contracts

Logic Specification Gestion et utilisation des étiquettes pré-définies Validation des annotations (type HOARE)

Sommaire des annotations et autres assertions

- requires
- assigns
- ensures
- decreases
- predicate
- ► logic
- ► lemma



Defining domain properties in logical theory

predicate

```
Listing 1 — project-divers/predicate4.c

//@ predicate pair(integer x) = (x/2)*2==x;
//@ predicate impair(integer x) = (x/2)*2!=x;
//@ lemma ex: \forall integer a, b; a < b \Rightarrow 2*a < 2*b;

/*@ inductive is_gcd(integer a, integer b, integer c) {
    case zero: \forall integer n; is_gcd(n,0,n);
    case deux: \forall integer u, v, w; u > v \Rightarrow is_gcd(u,v-u,w);
    case deux: \forall integer u, v, w; u < v \Rightarrow is_gcd(u, v-u, w);
}

*/
```

(Predicate)

Listing 2 - project-divers/predicate1.c

```
/*0 predicate is_positive(integer x) = x > 0; */

/*0 logic integer get_sign(real x) = 0 x > 0.071:(x < 0.07-1:0);

*/

*/ logic integer max(int x, int y) = x>=y?x:y;

*/
```

(Lemma)

Listing 3 – project-divers/lemma1.c

```
/*@ lemma div.mul.identity: 
@ \forall real x, real y; y != 0.0 \Longrightarrow y*(x/y) = x; @*/
/*@ lemma div.qr: 
@ \forall int a, int b; a >= 0 && b >0 \Longrightarrow 
\exists int q, int r; a = b*q +r && 0<=r && r <b; @*/
```

```
(Definition of fibonacci function)

Listing 4 — project-divers/predicate2.c

/*@ axiomatic mathfibonacci{
    @ logic integer mathfib(integer n);
    @ axiom mathfib0: mathfib(0) = 1;
    @ axiom mathfib1: mathfib(1) = 1;
    @ axiom mathfib1: mathfib(1) = 1;
    @ axiom mathfibrec: \forall integer n; n > 1

mathfib(n) = mathfib(n-1)+mathfib(n-2);
    @ } */
```

```
(Definition of gcd)

Listing 5 — project-divers/predicate3.c

/*@ inductive is_gcd(integer a, integer b, integer d) {
@ case gcd.zero:
@ \forall integer n; is_gcd(n,0,n);
@ case gcd.succ:
@ \forall integer a,b,d; is_gcd(b, a % b, d) \impress is_gcd(a,b,d); @}

@*/
```

- Cette expression est utilisable uniquement dans la postcondition ensures

```
(Valeur initiale x0)
                          Listing 6 – project-divers/old1.c
/#@ requires \valid(a) && \valid(b);
   @ assigns *a, *b;
    @ ensures *b = \langle old(*b) + \langle old(*a) + 2;
    @ ensures *a = \setminus old(*a)+2;
    @ ensures \result == 0;
*/
int old(int *a, int *b) {
  int x,y;
  x = *a;
  y = *b;
  x = x + 1;
  x = x + 1;
   y = y + x;
  *a = x;
  *b = v:
  return 0 ;
```

- ▶ id est une expression parmi Pre, Here, Old, Post, LoopEntry, LoopCurrent, Init

```
(label Pre)
                          Listing 7 – project-divers/at1.c
/*@
  requires \valid(a) && \valid(b);
  assigns *a, *b;
  ensures *a = \setminus old(*a) + 2;
  ensures *b = \langle old(*b)+ \rangle old(*a)+2;
int at1(int *a, int *b) {
//@ assert *a == \at(*a, Pre);
  *a = *a +1:
//@ assert *a == \at(*a, Pre)+1;
  *a = *a +1:
//@ assert *a == \at(*a, Pre)+2;
  *b = *b +*a:
//@ assert *a = \at(*a, Pre)+2 && *b = \at(*b, Pre)+\at(*a, Pre)+2;
  return 0:
```

```
(autre label) Listing 8 — project-divers/at2.c void f (int n) { for (int i = 0; i < n; i++) { /*Ø assert \at(i, LoopEntry) = 0; */ int j=0; while (j++< i) { /*Ø assert \at(j, LoopCurrent) + 1 = j; */ } } }
```

- requires
- assigns
- ensures
- decreases
- predicate
- ► logic
- ► lemma

```
Listing 9 - project-divers/anno0.c

int main(void){
    signed long int x,y,z;
    x = 1;
    /#@ assert x = 1; */
    y = 2;
    /#@ assert x = 1 && y = 2; */
    z = x * y;
    /#@ assert x = 1 && y = 1 && z==2; */
    return 0;
}
```

```
(Contrat valide) Listing 10-\text{project-divers/anno}00.c int main(void){
signed long int x,y,z; // int x,y,z;
x = 1;
/*0 assert x = 1; */
y = 2;
/*0 assert x = 1 && y = 2; */
z = x *y;
/*0 assert x = 1 && y = 2 && z = 2; */
return 0;
}
```

```
. . .
/*@ loop invariant I;
  @ loop assigns L;
*/
. . .
```

```
(Invariant de boucle)
                      Listing 11 – project-divers/anno5.c
/*@ requires a >= 0 && b >= 0;
 ensures 0 \le |result|;
 ensures \result < b;
 ensures \exists integer k; a = k * b + \result;
int rem(int a, int b) {
 int r = a;
 /*0
   loop invariant
   (\exists integer i; a = i * b + r) &&
    r >= 0;
    loop assigns r;
  while (r >= b) \{ r = r - b; \};
  return r:
```

```
(Invariant de boucle)
                      Listing 12 – project-divers/anno6.c
/*@ requires a >= 0 \&\& b >= 0;
 ensures 0 <= \result;
 ensures \result < b;
  ensures \exists integer k; a = k * b + \result;
int rem(int a, int b) {
  int r = a:
 /*@
    loop invariant
   (\exists integer i; a = i * b + r) &&
    r >= 0:
    loop assigns r;
  while (r >= b) \{ r = r - b; \};
  return r:
```

Echec de la preuve

L'invariant est insuffisamment informatif pour être prouvé et il faut ajouter une information sur y.

```
frama-c -wp anno6.c
[kernel] Parsing anno6.c (with preprocessing)
[wp] Warning: Missing RTE guards
[wp] anno6.c:8: Warning: Missing assigns clause (assigns 'everything' i
[wp] 2 goals scheduled
[wp] [Alt-Ergo 2.3.3] Goal typed_f_loop_invariant_preserved : Timeout (
[wp] [Cache] found:1
[wp] Proved goals: 1 / 2
                  1 (0.57ms)
 Qed:
  Alt-Ergo 2.3.3: 0 (interrupted: 1) (cached: 1)
[wp:pedantic-assigns] anno6.c:1: Warning:
  No 'assigns' specification for function 'f'.
  Callers assumptions might be imprecise.
```

Exemple d'invariant de boucle

Analyse avec succès

L'invariant est plus précis et donne des conditions liant x et y.

Résultat de l'analyse

```
frama-c -wp anno7.c
[kernel] Parsing anno7.c (with preprocessing)
[wp] Warning: Missing RTE guards
[wp] anno7.c:8: Warning: Missing assigns clause (assigns 'everything' i
[wp] 2 goals scheduled
[wp] [Cache] found:1
[wp] Proved goals: 2 / 2
Qed: 1 (0.32ms-3ms)
Alt-Ergo 2.3.3: 1 (6ms) (8) (cached: 1)
[wp:pedantic-assigns] anno7.c:1: Warning:
```

No 'assigns' specification for function 'f'. Callers assumptions might be imprecise.

- ▶ Un variant est une quantité qui décroît au cours de la boucle.
- Deux possibilités d'analyse sont possibles :
 - Terminaison d'une boucle (variant)
 - Terminaison de l'appel d'une fonction récursive (decreawse)

```
(Variant)

Listing 14 — project-divers/variant2.c

//@ loop variant e;

//@ decreases e;
```

Terminaison de boucle

- La terminaison est assurée en montrant que chaque boucle termine.
- Une boucle est caractérisée par une expression expvariant(x) appelée variant qui doit décroître à chaque exécution du corps de la boucle S où x_1 et x_2 sont les valeurs de X respectiveuent au début de la boucle S et à la fin de S :

```
\forall x_1, x_2.b(x_1) \land x_1 \xrightarrow{\ \ \ \ } x_2 \Rightarrow \mathsf{expvariant}(x_1) > \mathsf{expvariant}(x_2)
```

```
(Variant)
                    Listing 16 – project-divers/variant3.c
int f() {
int x = 0:
int y = 10;
/*@
    loop invariant
   0 <= x < 11 \&\& x+y == 10;
   loop variant y;
while (y > 0) {
 x++:
  y---:
 return 0;
```

Modèle de mémoire HOARE

- ▶ Pas de gestion de la mémoire comme les pointeurs
- ► Affectation à chaque variable une variable logique
- ► x++ avec x de type int et la C-variable est affectée à deux L-variables x2 = x1 + 1.

Exemples d'annotation

```
(Variant)
                       Listing 18 – project-divers/wp2.c
/*@CONSOLE
#include <LIMITS.h>
int q1() {
 int x=10, y=30, z=20;
//@ assert x== 10 && y == z+x && z==2*x;
y=z+x;
 //@ assert x== 10 && y == x+2*10;
x = x+1:
//@ assert x-1== 10 && y == x-1+2*10;
 return (0);
```

```
(Variant)
                      Listing 19 – project-divers/wp3.c
int q1() {
 int c = 2;
 /*@ assert c == 2; */
  int x;
 /*@ assert c == 2; */
 x = 3 * c;
 /*@ assert x == 6; */
  return (0);
```

```
(Variant)
                      Listing 20 – project-divers/wp4.c
int main()
  int a = 42; int b = 37;
  int c = a+b; // i:1
//@assert b == 37;
 a -= c; // i:2
 b += a; // i:3
//@assert b = 0 \&\& c = 79:
  return(0);
```

```
Listing 21 – project-divers/wp5.c
int main()
 int z; // instruction 8
  int a = 4; // instruction 7
//@assert a == 4 ;
  int b = 3; // instyruction 6
//@assert b = 3 \&\& a = 4:
  int c = a+b; // instruction 4
/*0 assert b = 3 \& c = 7 \& a = 4 ; */
 a += c; // instruction 3
 b += a: // instruction 2
//@ assert a = 11 \&\& b = 14 \&\& c = 7;
//@ assert a +b == 25 ;
 z = a*b; // instruction 1
//@assert a = 11 \&\& b = 14 \&\& c = 7 \&\& z = 154;
  return(0);
```

(Variant)

```
(Variant)

Listing 22 - project-divers/wp6.c

int main()
{
  int a = 4;
  int b = 3;
  int c = a+b; // i:1
  a += c; // i:2
  b += a; // i:3
  //@assert a = 11 && b = 14 && c = 7;
  return(0);
}
```

```
(Variant)
```

Listing 23 – project-divers/wp7.c

```
/*0 ensures x == a;
  ensures y == b;
void swap1(int a, int b) {
  int x = a:
  int y = b;
  //@ assert x == a \&\& y == b;
  int tmp:
  tmp = x;
  x = y;
  y = tmp;
  //@ assert x = a \&\& y = a;
void swap2(int a, int b) {
  int x = a;
  int y = b;
  //@ assert x = a \&\& y = b;
  x = x + y;
  y = x - y:
  x = x - y;
  //@ assert x = b \&\& y = a;
/#@ requires \valid(a);
  requires \valid(b);
  ensures *a = \setminus old(*b);
  ensures *b = \setminus old(*a);
void swap3(int *a, int *b) {
  int tmp:
  tmp = *a;
  *a = *b:
  *b = tmp:
```