
Cours MALG & MOVEX

MALG

Analyse des programmes

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Problem of Software Safety : Ariane 5 – Explosion after launch (1996)

- ▶ The Ariane 5 rocket explodes 37 seconds after launch.
- ▶ Loss of the rocket and its valuable cargo (500 million euros).
- ▶ Software error : Incorrect type conversion : a floating point number (64 bits) was converted to an integer (16 bits) without checking that the value fit in the space.
- ▶ Result : overflow \Rightarrow exception \Rightarrow guidance system shutdown \Rightarrow disorientation \Rightarrow self-destruction.
- ▶ The software came from Ariane 4 and was reused without modification, even though the speeds were much higher in Ariane 5.

- ▶ Several patients received massive doses of radiation (up to 100 times the intended dose).
- ▶ At least six patients were seriously irradiated, with several deaths.
- ▶ Software error : Race conditions between two internal processes : in some cases, a mechanical safety feature was disabled but the software thought it was active.
- ▶ Insufficient safeguards to check the actual status of the device.
- ▶ The software assumed that certain errors were 'impossible'. No alarm was raised.

- ▶ A battery of Patriot missiles failed to intercept an Iraqi Scud missile.
- ▶ Consequences : 28 American soldiers killed in Dhahran, Saudi Arabia.
- ▶ Software error* : Rounding error in floating point. The system used a time counter that accumulated a small error every second.
- ▶ After 100 hours of operation without restarting, the error had accumulated \Rightarrow incorrect calculation of the missile's position \Rightarrow interception failure.
- ▶ This precision bug caused an error of 0.34 seconds ... enough to miss a high-speed missile.

- ▶ Objectives of static program analysis
 - to prove properties about the run-time behaviour of a program
 - in a fully automatic way ie without interaction
 - without actually executing the program
- ▶ Applications
 - code optimisation
 - error detection (array out of bound access, null pointers)
 - proof support for generation of invariant

iOS bug : Emoji crashes Messages (2018)

- ▶ What happened : In 2018, a simple sequence of characters (a Telugu character, an Indian language) sent via iMessage or other apps such as WhatsApp, Facebook Messenger, etc., caused the app to crash or even completely freeze the iPhone.
- ▶ Where ? iPhone (iOS 11.2 to 11.2.5), iPad, Mac (in some apps too)
- ▶ Technical error : The iOS system did not know how to display a Telugu character correctly. The text rendering engine (CoreText) attempted to draw it, but poor memory management caused the system to crash or freeze completely.
- ▶ Consequences : The Messages app could not be opened (it crashed on launch). In some cases, the iPhone would restart in a loop. The conversation had to be deleted from another device or via iCloud.
- ▶ Apple fixed the bug in iOS 11.2.6, which was released as an emergency update shortly afterwards.
- ▶ Everyone uses Messages or Messenger. A single character could crash an iPhone : it was used as a joke on social media. It shows how a small detail (a text character !) can have a disproportionate effect on complex systems.

- ▶ Static Analysis *computes* approximations
- ▶ Abstract Interpretation (AI) provides a mathematical framework for relating *approximations*
- ▶ Properties of programs are generally non computable :
 - the halting problem is undecidable
 - Model checking is computing over finite structures
 - Proof assistant may be useful for proving partial correctness or total correctness by applying induction principles (see Event B)
 - AI provides another solution by transferring results from a *concret* framework to an *abstract* structure

- ▶ $\mathcal{CS}(P)$ is the concrete semantics of a program P : the set of reachable states of P .
- ▶ $\mathcal{AS}(P)$ is the approximation of $\mathcal{CS}(P)$: $\mathcal{CS}(P) \subseteq \mathcal{AS}(P)$.
- ▶ $\mathcal{CS}(P)$ is generally not computable and we will seek for *computable* approximation or abstract semantics $\mathcal{AS}(P)$.
- ▶ Problems : $\mathcal{AS}(P)$ may *loose* the expression of properties.

- ▶ φ is a program property stating the possible bugs or errors which we want to avoid.
- ▶ $\mathcal{CS}(P)$ is the concrete semantics of a program P : the set of reachable states of P .
- ▶ $\mathcal{AS}(P)$ is the approximation of $\mathcal{CS}(P)$: $\mathcal{CS}(P) \subseteq \mathcal{AS}(P)$.
- ▶ Case 1 : $\mathcal{CS}(P) \cap \varphi = \emptyset$ and $\mathcal{AS}(P) \cap \varphi = \emptyset$
- ▶ Case 2 : $\mathcal{CS}(P) \cap \varphi \neq \emptyset$ and $\mathcal{AS}(P) \cap \varphi \neq \emptyset$
- ▶ Case 3 : $\mathcal{CS}(P) \cap \varphi = \emptyset$ and $\mathcal{AS}(P) \cap \varphi \neq \emptyset$

- ▶ Case 1 : $\mathcal{CS}(P) \cap \varphi = \emptyset$ and $\mathcal{AS}(P) \cap \varphi = \emptyset$:
 - P is safe with respect to φ and no error specified by φ is possible for P .
 - Checking is computable on the approximation
- ▶ Case 2 : $\mathcal{CS}(P) \cap \varphi \neq \emptyset$ and $\mathcal{AS}(P) \cap \varphi \neq \emptyset$:
 - An error is detected on the approximation and on the concrete semantics.
 - P is unsafe with respect to φ
 - and an error is detected by the analyser.
- ▶ Case 3 : $\mathcal{CS}(P) \cap \varphi = \emptyset$ and $\mathcal{AS}(P) \cap \varphi \neq \emptyset$:
 - P is safe with respect to φ
 - but an error is detected by the analyser
 - A false alarm is provided by the analyzer
 - Approximation is over-approximating P with respect to φ
 - The analysis should be refined

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- ▶ a web interface to the Interproc analyzer connected to the APRON Abstract Domain Library
- ▶ Analysis of programs using different *abstract domains*
- ▶ <http://pop-art.inrialpes.fr/interproc/>
- ▶ developed by Antoine Miné and his team.

Example 1 : Increment of a value

```
proc incr (x:int) returns (y:int)
begin
  y = x+1;
end
```

```
var i:int;
begin
  i = 0;
  while (i<=10) do
    i = incr(i);
  done;
end
```

Example 1 : Increment of a value

```
proc incr (x : int) returns (y : int) var ;
begin
  /* (L3 C5) [|x>=0; -x+10>=0|] */
  y = x + 1; /* (L4 C10)
                                     [|x>=0; -x+10>=0; y-1>=0; -y+11>=0|] */
end

var i : int;
begin
  /* (L8 C5) top */
  i = 0; /* (L9 C8) [|i>=0; -i+11>=0|] */
  while i <= 10 do
    /* (L10 C18) [|i>=0; -i+10>=0|] */
    i = incr(i); /* (L11 C16)
                                     [|i-1>=0; -i+11>=0|] */
  done; /* (L12 C7) [|i-11=0|] */
end
```

Example 2 : Division

```
var Q : int, R : int, X : int, Y : int;  
begin  
  
    Q = 0;  
    R = Y;  
  
    while R >= Y do  
        Q = Q + 1;  
        R = R - Y;  
    done;  
  
end
```

Example 2 : Division

Annotated program after forward analysis

```
var Q : int, R : int, X : int, Y : int;  
begin  
  /* (L3 C5) top */  
  Q = 0; /* (L5 C8) [|Q=0|] */  
  R = Y; /* (L6 C8) [|Q>=0|] */  
  while R >= Y do  
    /* (L8 C20) [|Q>=0|] */  
    Q = Q + 1; /* (L9 C17) [|Q-1>=0|] */  
    R = R - Y; /* (L10 C17) [|Q-1>=0|] */  
  done; /* (L11 C10) [|Q>=0|] */  
end
```

Example 3 : modified division

```
var Q : int, R : int, X : int, Y : int;  
begin  
  
  Q = 0;  
  R = Y;  
  if Y > 0 then  
    while R >= Y do  
      Q = Q + 1;  
      R = R - Y;  
    done;  
  else  
    skip;  
  endif;  
end
```

Example 3 : modified division

Annotated program after forward analysis

Annotated program after forward analysis

```
var Q : int, R : int, X : int, Y : int;
begin
  /* (L3 C5) top */
  Q = 0; /* (L5 C8) [|Q=0|] */
  R = Y; /* (L6 C8) [|Q=0|] */
  if Y > 0 then
    /* (L7 C15) [|Q>=0; Y-1>=0|] */
    while R >= Y do
      /* (L8 C20) [|Q>=0; R-1>=0; Y-1>=0|] */
      Q = Q + 1; /* (L9 C17)
                  [|Q-1>=0; R-1>=0; Y-1>=0|] */
      R = R - Y; /* (L10 C17)
                  [|Q-1>=0; Y-1>=0|] */
    done; /* (L11 C10) [|Q>=0; Y-1>=0|] */
  else
    /* (L12 C6) [|Q=0; -Y>=0|] */
    skip; /* (L13 C9) [|Q=0; -Y>=0|] */
  endif; /* (L14 C8) [|Q>=0|] */
end
```

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Static analysis

Static program analysis analyses computer software without actually executing programs :

- ▶ absence of run time errors
 - ▶ detecting variables used before initialisation.
- ▶ Data flow analysis
 - ▶ Abstract interpretation
 - ▶ Use of property-preserving abstractions
 - ▶ Programs are interpreted in abstractions

- ▶ Sign analysis is used to determine the sign of variables
- ▶ x is an integer variable and has the following possible *abstract* states :
 - $x > 0$
 - $x \geq 0$
 - $x = 0$
 - $x < 0$
 - $x \leq 0$
 - $x \neq 0$

```
ℓ0 :  
y := -11;  
ℓ0 :  
IF x < y THEN  
    ℓ1 :  
    z := y;  
    ℓ2 :  
ELSE  
    ℓ3 :  
    z := x;  
    ℓ4 :  
FI  
ℓ5 :
```

```
ℓ0 :  
y := -11;  
ℓ0 : y < 0  
IF x < y THEN  
    ℓ1 :  
    z := y;  
    ℓ2 :  
ELSE  
    ℓ3 :  
    z := x;  
    ℓ4 :  
FI  
ℓ5 :
```

```
 $\ell_0 :$   
 $y := -11;$   
 $\ell_0 : y < 0$   
IF  $x < y$  THEN  
   $\ell_1 : y < 0 \quad x < 0$   
   $z := y;$   
   $\ell_2 :$   
ELSE  
   $\ell_3 :$   
   $z := x;$   
   $\ell_4 :$   
FI  
 $\ell_5 :$ 
```

```
 $\ell_0 :$   
 $y := -11;$   
 $\ell_0 : y < 0$   
IF  $x < y$  THEN  
     $\ell_1 : y < 0 \quad x < 0$   
     $z := y;$   
     $\ell_2 : y < 0 \quad x < 0 \quad z < 0$   
ELSE  
     $\ell_3 :$   
     $z := x;$   
     $\ell_4 :$   
FI  
 $\ell_5 :$ 
```

```
 $\ell_0 :$   
 $y := -11;$   
 $\ell_0 : y < 0$   
IF  $x < y$  THEN  
   $\ell_1 : y < 0 \quad x < 0$   
   $z := y;$   
   $\ell_2 : y < 0 \quad x < 0 \quad z < 0$   
ELSE  
   $\ell_3 : y < 0 \quad x \in \mathbb{Z}$   
   $z := x;$   
   $\ell_4 :$   
FI  
 $\ell_5 :$ 
```

```
 $\ell_0 :$   
 $y := -11;$   
 $\ell_0 : y < 0$   
IF  $x < y$  THEN  
   $\ell_1 : y < 0 \quad x < 0$   
   $z := x;$   
   $\ell_2 : y < 0 \quad x < 0 \quad z < 0$   
ELSE  
   $\ell_3 : y < 0 \quad x \in \mathbb{Z}$   
   $z := x;$   
   $\ell_4 : y < 0 \quad x \in \mathbb{Z}$   
FI  
 $\ell_5 :$ 
```



```
 $\ell_0 :$   
 $y := -11;$   
 $\ell_0 : y < 0$   
IF  $x < y$  THEN  
   $\ell_1 : y < 0 \quad x < 0$   
   $z := x;$   
   $\ell_2 : y < 0 \quad x < 0 \quad z < 0$   
ELSE  
   $\ell_3 : y < 0 \quad x \in \mathbb{Z}$   
   $z := x;$   
   $\ell_4 : y < 0 \quad x \in \mathbb{Z}$   
FI  
 $\ell_5 : y < 0 \quad x \in \mathbb{Z} \quad z \in \mathbb{Z}$ 
```

```
 $\ell_0 :$   
 $y := -11;$   
 $\ell_0 : y < 0$   
IF  $x < y$  THEN  
   $\ell_1 : y < 0 \quad x < 0$   
   $z := x;$   
   $\ell_2 : y < 0 \quad x < 0 \quad z < 0$   
ELSE  
   $\ell_3 : y < 0 \quad x \in \mathbb{Z}$   
   $z := x;$   
   $\ell_4 : y < 0 \quad x \in \mathbb{Z}$   
FI  
 $\ell_5 : y < 0 \quad x \in \mathbb{Z} \quad z \in \mathbb{Z}$ 
```

Result

$y < 0 \quad x \in \mathbb{Z} \quad z < 0$ means that $z < 0$ is an information resulting from the analysis over abstract domain of signs.

- ▶ \mathcal{MS} is $(Th(s, c), x, \text{VALS}, \text{INIT}(x), \{r_0, \dots, r_n\})$
- ▶ $\text{NEXT} \stackrel{\text{def}}{=} r_0 \vee \dots \vee r_n.$
- ▶ S is a safety property, when
 $\forall x_0, x \in \text{VALS}. \text{Init}(x_0) \wedge \text{NEXT}^*(x_0, x) \Rightarrow x \in S.$
- ▶ S is a safety property for \mathcal{MS} if, and only if, $\text{REACHABLE}(\mathcal{MS}) \subseteq S$

Characterisation of $\text{REACHABLE}(\mathcal{MS}) \subseteq S$ as a fixed-point

$(\mathcal{P}(\text{VALS}), \subseteq, \emptyset, \cup, \cap)$ is a complete lattice and

$F \in \mathcal{P}(\text{VALS}) \longrightarrow \mathcal{P}(\text{VALS})$ is defined as :

$F(X) = \{x | x \in \text{VALS} \wedge \text{INIT}(x)\} \cup X \cup \longrightarrow [X]$ and satisfies the following properties :

- ▶ F is a monotonic function.
- ▶ $\text{REACHABLE}(\mathcal{MS}) = \mu F$
- ▶ μF is defined as follows :
 - $F^0 = \emptyset$
 - $F^{i+1} = F(F^i), \forall i \in \mathbb{N}$
 - $\mu F = \text{Sup}\{F^i | i \in \mathbb{N}\}$
 - For any safety property S , $\mu F \subseteq S$.

Computing the least fixed-point over a finite lattice

INPUT $F \in T \longrightarrow T$

OUTPUT $result = \mu.F$

VARIABLES $x, y \in T, i \in \mathbb{N}$

$\ell_0 : \{x, y \in T\}$

$x := \perp;$

$y := \perp;$

$i := 0;$

$\ell_{11} : \{x, y \in T \wedge x = F^i \wedge y = \bigcup_{k=0; k=i} F^k \wedge i \leq Card(T) \wedge i = 0\};$

WHILE $i \leq Card(T)$

$\ell_1 : \{x, y \in T \wedge x = F^i \wedge y = \bigcup_{k=0; k=i} F^k \wedge i \leq Card(T)\};$

$x := F(x);$

$\ell_2 : \{x, y \in T \wedge x = F^{i+1} \wedge y = \bigcup_{k=0; k=i} F^k \wedge i \leq Card(T)\};$

$y := x \sqcup y;$

$\ell_3 : \{x, y \in T \wedge x = F^{i+1} \wedge y = \bigcup_{k=0; k=i+1} F^k \wedge i \leq Card(T)\};$

$i := i+1;$

$\ell_4 : \{x, y \in T \wedge x = F^i \wedge y = \bigcup_{k=0; k=i} F^k \wedge i \leq Card(T)+1\};$

OD;

$\ell_5 : \{x, y \in T \wedge x = F^i \wedge y = \bigcup_{k=0; k=i} F^k \wedge i = Card(T)+1\};$

$result := y;$

$\ell_6 : \{x, y \in T \wedge x = F^i \wedge y = \bigcup_{k=0; k=i} F^k \wedge i = Card(T)+1 \wedge result = y\};$

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Examples

- ▶ Computation Traces of Program
- ▶ Transitive Closure of the program transition relation
- ▶ Set of states

The collecting semantics is the semantics which is interesting our analysis and we will consider as collecting semantics the set of states.

► *Collecting semantics*

- Static analysis of a program states a property of program executions defined by a *standard* semantics.
- Defining a so-called *collecting* semantics defining the strongest static property of interest
- Collecting semantics defines the class of static analysis, which approximates it
- State properties are subsets of $\mathbb{I} \times \mathbb{I} \times \mathbb{I} \times \mathbb{I}$ and abstract interpretation executes programs on these properties

► Approximation

- Spaces of values should be restricted to computable entities
- Over-approximation of concrete properties

| | | | |
|--------|-------|---|--|
| $Expr$ | $::=$ | v $?$ x $Expr\ op\ Expr$ | $v \in \mathbb{Z}$ $x \in \mathbb{V}$ $op \in \{+, -, \times, /\}$ |
| $cond$ | $::=$ | $Expr\ relop\ Expr$ not $cond$ $cond$ and $cond$ | $relop \in \{<, \leq, >, \geq, =, \neq\}$ |
| $stmt$ | $::=$ | $\ell[x := Expr]$ $\ell[skip]$ if $\ell[cond]$ then $stmt$ else $stmt$ end if while $\ell[cond]$ do $stmt$ end do $stmt; stmt$ | $\ell \in \mathbb{C}$ |

Two examples of annotated programs

```
 $\ell_0[X := 0];$   
 $\ell_1[Y := Y + X];$   
 $\ell_2[skip]$   
 $\ell_3[X := Y];$ 
```

```
 $\ell_0[Q := 0];$   
 $\ell_1[R := X];$   
IF  $\ell_5[Y > 0]$   
    WHILE  $\ell_2[R \geq Y]$   
         $\ell_3[Q := Q + 1];$   
         $\ell_4[R := R - Y]$   
    ENDWHILE  
ELSE  
     $\ell_6[skip]$   
ENDIF
```

► Semantic Domains

$$Mem \stackrel{def}{=} \mathbb{V} \longrightarrow \mathbb{Z}$$

$$States \stackrel{def}{=} \mathbb{C} \times Mem$$

► Semantics for Expressions

$$\mathcal{E}[[v]](m) \in \mathcal{P}(\mathbb{Z}), e \in Expr, m \in Mem, x \in \mathbb{V}, op \in \{+, -, \times, /\}$$

$$\mathcal{E}[[v]](m) \stackrel{def}{=} \{v\}$$

$$\mathcal{E}[[?]](m) \stackrel{def}{=} \mathbb{Z}$$

$$\mathcal{E}[[x]](m) \stackrel{def}{=} \{m(x)\}$$

$$\mathcal{E}[[e_1 \text{ op } e_2]](m) \stackrel{def}{=} \{v \mid \exists ve_1, ve_2. \left(\begin{array}{l} ve_1 \in \mathcal{E}[[e_1]](m) \\ ve_2 \in \mathcal{E}[[e_2]](m) \\ v = ve_1 \text{ o } ve_2 \end{array} \right) \}$$

► Semantics for conditions

$\mathcal{C}[\![cond]\!](m) \in \mathcal{P}(\mathbb{B})$, $cond \in Cond, m \in Mem, x \in \mathbb{V}$,
 $op \in \{+, -, \times, /\}$

$$\begin{aligned} tt \in \mathcal{C}[\![e_1 \text{ relop } e_2]\!](m) & \stackrel{def}{=} \exists v_1, v_2. \left(\begin{array}{l} v_1 \in \mathcal{E}[\![e_1]\!](m) \\ v_2 \in \mathcal{E}[\![e_2]\!](m) \\ v_1 \text{ relop } v_2 \end{array} \right) \\ ff \in \mathcal{C}[\![e_1 \text{ relop } e_2]\!](m) & \stackrel{def}{=} \exists v_1, v_2. \left(\begin{array}{l} v_1 \in \mathcal{E}[\![e_1]\!](m) \\ v_2 \in \mathcal{E}[\![e_2]\!](m) \\ \neg(v_1 \text{ relop } v_2) \end{array} \right) \\ be_1 \wedge be_2 \in \mathcal{C}[\![be_1 \text{ and } be_2]\!](m) & \stackrel{def}{=} and \left(\begin{array}{l} be_1 \in \mathcal{C}[\![be_1]\!](m) \\ be_2 \in \mathcal{C}[\![be_2]\!](m) \end{array} \right) \end{aligned}$$

- ▶ $(x := e, m) \longrightarrow m[x \mapsto v]$, **where** $v \in \mathcal{E}[[e]](m)$
- ▶ $(skip, m) \longrightarrow m$
- ▶ **If** $(S_1, m) \longrightarrow m'$, **then** $(S_1; S_2, m) \longrightarrow (S_2, m')$.
- ▶ **If** $tt \in \mathcal{C}[[be]]$, **then** $(\text{if } be \text{ then } S_1 \text{ else } S_2 \text{ end if}, m) \longrightarrow (S_1, m)$.
- ▶ **If** $ff \in \mathcal{C}[[be]]$, **then**
 $(\text{if } be \text{ then } S_1 \text{ else } S_2 \text{ end if}, m) \longrightarrow (S_2, m)$.
- ▶ **If** $tt \in \mathcal{C}[[be]]$, **then**
 $(\text{while } be \text{ do } S \text{ end do}, m) \longrightarrow (S; \text{while } be \text{ do } S \text{ end do}, m)$.
- ▶ **If** $ff \in \mathcal{C}[[be]]$, **then** $(\text{while } be \text{ do } S \text{ end do}, m) \longrightarrow m$.

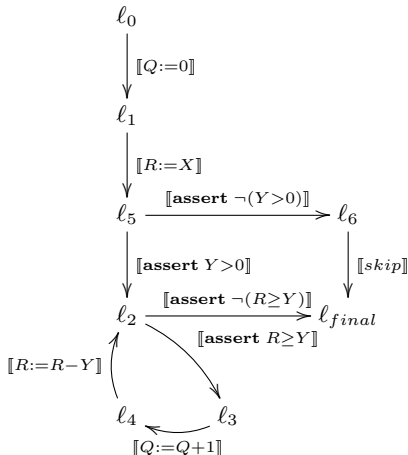
- ▶ A *control flow graph* is generated from the program under consideration namely P .
- ▶ A control flow graph $\mathcal{CFG}\llbracket P \rrbracket$ is defined by nodes ($l \in \mathcal{C}$) which are program control points of P , $\text{Control}\llbracket P \rrbracket$ and by labelled edges with actions ($\text{Actions}\llbracket P \rrbracket$) defined by the following rules :

$$\begin{array}{lcl} \text{actions} & ::= & v := \text{exp} \\ & | & \text{skip} \\ & | & \mathbf{assert} \text{ be} \end{array}$$

- ▶ A *control flow graph* is effectively defined by :
 - $\ell_{init} \in \text{Control}\llbracket P \rrbracket$: the entry point
 - $\ell_{end} \in \text{Control}\llbracket P \rrbracket$: the exit point
 - $\text{Edges}\llbracket P \rrbracket \subseteq \text{Control}\llbracket P \rrbracket \times \text{Actions}\llbracket P \rrbracket \times \text{Control}\llbracket P \rrbracket$
- ▶ $\mathcal{CFG}\llbracket P \rrbracket = (\ell_{init}, \text{Edges}\llbracket P \rrbracket, \ell_{end})$

From program to flowchart

```
ℓ0[Q := 0];  
ℓ1[R := X];  
IF ℓ5[Y > 0]  
    WHILE ℓ2[R ≥ Y]  
        ℓ3[Q := Q + 1];  
        ℓ4[R := R - Y]  
    ENDWHILE  
ELSE  
    ℓ6[skip]  
ENDIF
```



- ▶ $Mem \stackrel{def}{=} \mathbb{V} \longrightarrow \mathbb{Z}$
- ▶ Semantics of actions : $\xrightarrow{a} \subseteq Mem \times Mem$
 - $m \xrightarrow{x:=e} m[x \mapsto v]$ if there is a value $v \in \mathcal{E}[e](m)$
 - $m \xrightarrow{skip} m$
 - $m \xrightarrow{\text{assert } be} m$ if $tt \in \mathcal{C}[be](m)$
- ▶ Semantics for $\mathcal{CFG}[P]$: $\xrightarrow{P} \subseteq States \times States$
 - If $m \xrightarrow{a} m'$ and $(\ell_1, a, \ell_2) \in \mathcal{Edges}[P]$, then $(\ell_1, m) \xrightarrow{P} (\ell_2, m')$
 - The set of initial states is $\{\ell_{init}\} \times Mem$
 - The set of reachable states for P is denoted $\text{REACHABLE}(P)$ and defined by $\llbracket P \rrbracket = \{s \mid \exists s_0 \in \{\ell_{init}\} \times Mem : s_0 \xrightarrow[\star]{P} s\}$.

- ▶ Defining for each control point ℓ of P the set of reachable values :

$$\llbracket P \rrbracket_{\ell}^{coll} = \{s \mid s \in States \wedge s \in \llbracket P \rrbracket \wedge \exists m \in Mem : s = (\ell, m)\}$$

- ▶ Characterizing $\llbracket P \rrbracket_{\ell}^{coll}$: it satisfies the system of equations

$$\forall \ell \in \mathcal{C}(P). X_{\ell} = X_{\ell}^{init} \cup \bigcup_{(\ell_1, a, \ell) \in \mathcal{E}dges[P]} \llbracket a \rrbracket(X_{\ell_1}) \quad (1)$$

- ▶ Let $a \in \mathcal{A}ctions[\llbracket P \rrbracket]$ and $x \subseteq Mem$.

$$\llbracket a \rrbracket(x) = \{e \mid e \in States \wedge \exists f. f \in x \wedge f \xrightarrow{a} e\}$$

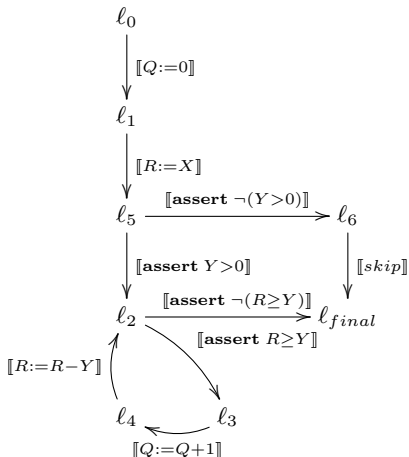
- ▶ $\forall \ell \in \mathcal{C}(P). \left(\begin{array}{l} \ell = \ell_{init} \Rightarrow X_{\ell}^{init} = Mem \\ \ell \neq \ell_{init} \Rightarrow X_{\ell}^{init} = \emptyset \end{array} \right)$

.....
☺ Théorème Let F the function defined as follows :

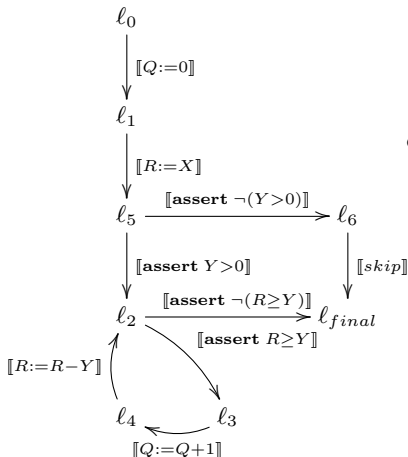
- ▶ n is the cardinality of $\mathcal{C}(P)$.
- ▶ $F \in \mathcal{P}(\text{States})^n \longrightarrow \mathcal{P}(\text{States})^n$
- ▶ If $X \in \mathcal{P}(\text{States})^n$, then $F(X) = (\dots, F_\ell(X), \dots)$
- ▶ $\forall \ell \in \mathcal{C}(P). F_\ell(X) = X_\ell^{init} \cup \bigcup_{(\ell_1, a, \ell) \in \text{Edges}[P]} \llbracket a \rrbracket(X_{\ell_1})$

The function F is monotonic over the complete lattice $(\mathcal{P}(\text{States})^n, \subseteq)$ and has a least fixed-point μF defining the collecting semantics.

From flowchart to equational system

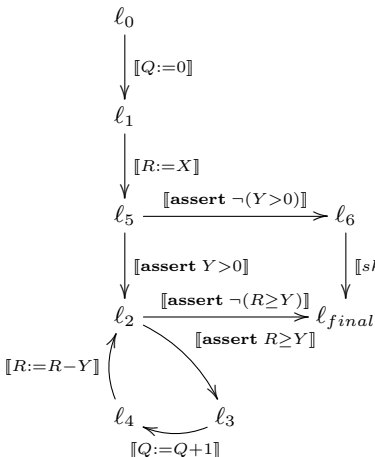


From flowchart to equational system



Defining equational systems for the collecting semantics :

From flowchart to equational system



Defining equational systems for the collecting semantics :

$$\left\{ \begin{array}{l} X_0 = Mem \\ X_1 = \llbracket Q := 0 \rrbracket(X_0) \\ X_5 = \llbracket R := X \rrbracket(X_1) \\ X_2 = \llbracket assert(Y > 0) \rrbracket(X_5) \cup \llbracket R := R - Y \rrbracket(X_4) \\ X_3 = \llbracket assert R \geq Y \rrbracket(X_2) \\ X_4 = \llbracket Q := Q + 1 \rrbracket(X_3) \\ X_6 = \llbracket assert \neg(Y > 0) \rrbracket(X_5) \\ X_7 = X_6 \cup \llbracket assert \neg(R \geq Y) \rrbracket(X_2) \end{array} \right.$$

- ▶ The collecting semantics is the least fixed-point of the system of equations, which exists by fixed-point theorems.
- ▶ Questions :
 - How to compute the solution ?
 - Computing over finite structures, when it is possible....
 - Using an approximation of fixed-points ?
 - What is an approximation ?
 - What is an abstraction ?
 - What is the best abstraction ?

Next step

Defining a framework for computing lfp solution of these equational systems in any case.

A control flow graph with four nodes: l_1 , l_2 , l_3 , and l_4 . The nodes are represented as orange circles. The edges are blue arrows with associated labels:

- From l_1 to l_2 : $v := 1$
- From l_2 to l_4 : $\neg(v < 10)$
- From l_2 to l_3 : $v < 10$
- From l_3 to l_2 : $v := v + 1$

```
graph TD; l1((l1)) -- "v := 1" --> l2((l2)); l2 -- "¬(v < 10)" --> l4((l4)); l2 -- "v < 10" --> l3((l3)); l3 -- "v := v + 1" --> l2;
```

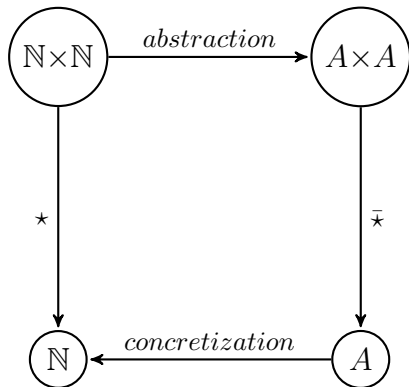
- ## ► Reachability

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- ▶ $n \in \mathbb{N}$: abstraction for n is defined as $modulo(n, 9)$.
- ▶ \mathbb{A} is the set of possible abstract values $\bar{0}, \bar{1}, \bar{2}, \bar{3}, \bar{4}, \bar{5}, \bar{6}, \bar{7}, \bar{8}$
- ▶ $n \star m = \bar{n} \star \bar{m}$
- ▶ $25 \star 25 = 625$, $\bar{25} \star \bar{25} = \bar{625}$, $\bar{7} \star \bar{7} = \bar{4}$, $\overline{7 \star 7} = \bar{4}$, $\overline{49} = \bar{4}$, $\bar{4} = \bar{4}$,
- ▶ But $25 \star 25 \neq 265$ and $\overline{25 \star 25} = \overline{265}$

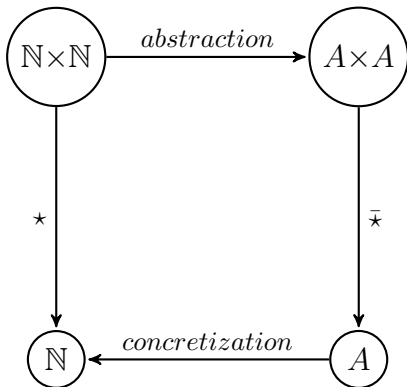
Finding Sound Abstractions for Computing

- ▶ $n \in \mathbb{N}$: abstraction for n is defined as $\text{modulo}(n, 9)$.
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Finding Sound Abstractions for Computing

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.....
☺ Rule

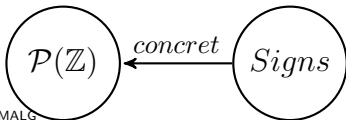
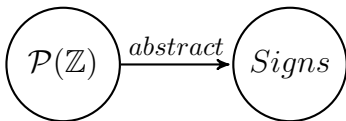
If $\overline{a \star b} \neq \overline{a} \star \overline{b}$,

then the operation is not correct.

- ▶ A number $z \in \mathbb{Z}$ is soundly approximated by an abstract value $abstract(z) \in Signs$.
- ▶ 2 is approximated by $pos : \{2\} \subseteq concrete(pos)$
- ▶ $\{2, 8\}$ is approximated by $pos : \{2, 8\} \subseteq concrete(pos)$
- ▶ -2 is approximated by $neg : \{2\} \subseteq concrete(neg)$
- ▶ 0 is approximated by $zero : \{0\} \subseteq concrete(zero)$
- ▶ $\{-2, -8\}$ is approximated by $pos : \{-2, -8\} \subseteq concrete(neg)$
- ▶ $\{-2, 2, 8\}$ is approximated by $pos : \{-2, 2, 8\} \subseteq concrete(nonzero)$

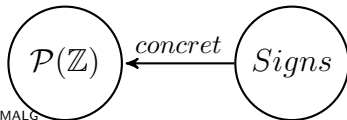
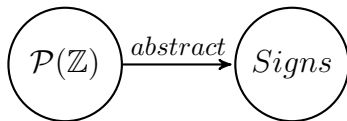
Finding Sound Abstractions for Computing

- ▶ A number $z \in \mathbb{Z}$ is soundly approximated by an abstract value $\text{abstract}(z) \in \text{Signs}$.
- ▶ 2 is approximated by $\text{pos} : \{2\} \subseteq \text{concrete}(\text{pos})$
- ▶ $\{2, 8\}$ is approximated by $\text{pos} : \{2, 8\} \subseteq \text{concrete}(\text{pos})$
- ▶ -2 is approximated by $\text{neg} : \{2\} \subseteq \text{concrete}(\text{neg})$
- ▶ 0 is approximated by $\text{zero} : \{0\} \subseteq \text{concrete}(\text{zero})$
- ▶ $\{-2, -8\}$ is approximated by $\text{pos} : \{-2, -8\} \subseteq \text{concrete}(\text{neg})$
- ▶ $\{-2, 2, 8\}$ is approximated by $\text{pos} : \{-2, 2, 8\} \subseteq \text{concrete}(\text{nonzero})$



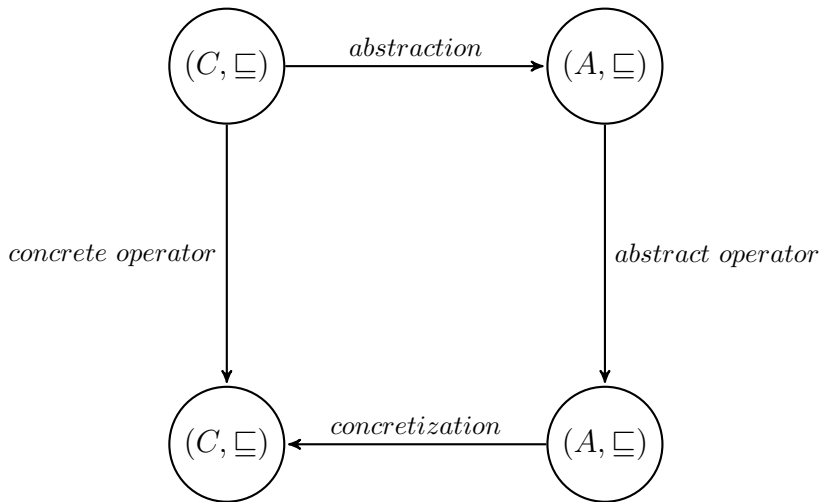
Finding Sound Abstractions for Computing

- ▶ A number $z \in \mathbb{Z}$ is soundly approximated by an abstract value $\text{abstract}(z) \in \text{Signs}$.
- ▶ 2 is approximated by $\text{pos} : \{2\} \subseteq \text{concret}(\text{pos})$
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- ▶ $\{-2, 2, 8\}$ is approximated by $\text{pos} : \{-2, 2, 8\} \subseteq \text{concret}(\text{nonzero})$



- ▶ $\text{concret}(\text{pos}) = \{z \mid z \in \mathbb{Z} \wedge z > 0\}$
- ▶ $\text{concret}(\text{neg}) = \{z \mid z \in \mathbb{Z} \wedge z < 0\}$
- ▶ $\text{concret}(\text{non}) = \emptyset$
- ▶ $\text{concret}(\text{nonzero}) = \{z \mid z \in \mathbb{Z} \wedge z \neq 0\}$

- ▶ A number $z \in \mathbb{Z}$ is soundly approximated by an abstract value $abstract(z) \in Signs$.
- ▶ 2 is approximated by pos :
 - $\{2\} \subseteq concrete(pos)$
 - $abstract(\{2\}) = pos$
- ▶ $\{2, 8\}$ is approximated by pos :
 - $\{2, 8\} \subseteq concrete(pos)$
 - $abstract(\{2, 8\}) = pos$
- ▶ $\{-2, 2, 8\}$ is approximated by pos :
 - $\{-2, 2, 8\} \subseteq concrete(nonzero)$
 - $abstract(\{-2, 2, 8\}) = nonzero$



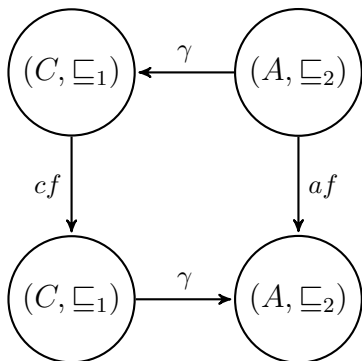
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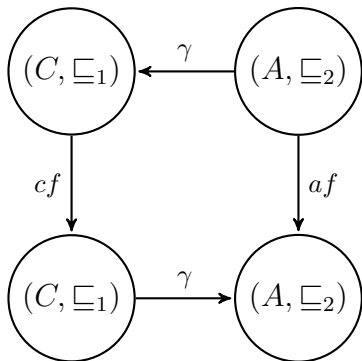
- ▶ Let $(A, \sqsubseteq, \sqcup, \sqcap)$ be a complete lattice. Let Q a subset of A . Q is a Moore family, if for each part Q' of Q , $\sqcap Q' \in Q$.
- ▶ **Property** : Let $(A, \sqsubseteq, \sqcup, \sqcap)$ be a complete lattice and $B \subseteq A$.
 - ① If, for any $p \in A$, $\{q \in B | p \sqsubseteq q\}$ has a least element, then B is Moore family.
 - ② If B is Moore family, then for any $p \in A$, $\{q \in B | p \sqsubseteq q\}$ has a least element.

If B is a Moore family, then it is a good abstraction.

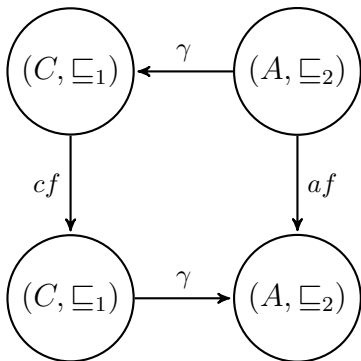
- ▶ Let $(A, \sqsubseteq, \sqcup, \sqcap)$ be a complete lattice.
- ▶ $\rho \in A \longrightarrow A$ is a upper closure operator, if it satisfies the following properties :
 - ρ is monotonic : $\forall x, y \in A. x \sqsubseteq y \Rightarrow \rho(x) \sqsubseteq \rho(y)$.
 - ρ is extensive : $\forall x \in A. x \sqsubseteq \rho(x)$.
 - ρ is idempotent : $\forall x \in A. \rho(x) = \rho(\rho(x))$.
- ▶ **Property** : Let $(A, \sqsubseteq, \sqcup, \sqcap)$ be a complete lattice. and $B \subseteq A$.
 B is a Moore family if, and only if, there exists a closure operator ρ such that $B = \rho(A)$
- ▶ B is a *good* abstraction of C , if it satisfies $B = \rho(C)$ where ρ is a closure operator.

- ▶ Two complete lattices $(C, \sqsubseteq_1, \sqcup_1, \sqcap_1)$ and $(A, \sqsubseteq_2, \sqcup_2, \sqcap_2)$ are supposed to be given.
- ▶ Two functions α and γ are supposed to be defined as follows :
 - $\alpha \in C \longrightarrow A$
 - $\gamma \in A \longrightarrow C$
- ▶ The pair (α, γ) is a Galois connection, if it satisfies the following property : $\forall x_1 \in C, x_2 \in A. \alpha(x_1) \sqsubseteq_2 x_2 \Leftrightarrow x_1 \sqsubseteq_1 \gamma(x_2)$
- ▶ A complete lattice A is a good abstraction of L , when there is a Galois connection between A and L .

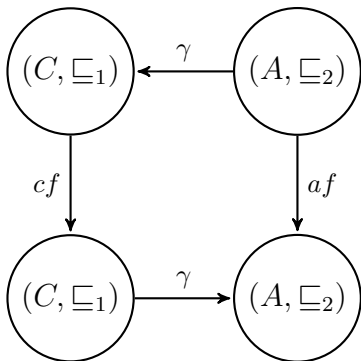




► a is a sound abstraction of c , if $c \sqsubseteq_1 \gamma(a)$.



- ▶ a is a sound abstraction of c , if $c \sqsubseteq_1 \gamma(a)$.
- ▶ functional operator : af is a sound abstraction of cf , if $\forall a \in A. cf(\gamma(a)) \sqsubseteq_1 \gamma(af(a))$



- ▶ a is a sound abstraction of c , if $c \sqsubseteq_1 \gamma(a)$.
- ▶ functional operator : af is a sound abstraction of cf , if $\forall a \in A. cf(\gamma(a)) \sqsubseteq_1 \gamma(af(a))$
- ▶ relational operator : ar is a sound abstraction of cr , if $\forall a \in A. cr(\gamma(a_1), \dots, \gamma(a_n)) \sqsubseteq_1 \gamma(ac(a_1, \dots, a_n))$

Galois Connections

The pair (α, γ) is a Galois connection, if it satisfies the following property : $\forall x_1 \in L, x_2 \in L. \alpha(x_1) \sqsubseteq' x_2 \Leftrightarrow x_1 \sqsubseteq \gamma(x_2)$

Notation : $L \xrightleftharpoons[\alpha]{\gamma} L'$

Properties of a Galois connection $L \xrightleftharpoons[\alpha]{\gamma} L'$

- ▶ α and γ are monotonic over the lattices.
- ▶ $\text{id}(L) \subseteq \gamma \circ \alpha$: $\gamma \circ \alpha$ is extensive.
- ▶ $\alpha \circ \gamma \subseteq \text{id}(L')$: $\alpha \circ \gamma$ is retractive.
- ▶ $\alpha \circ \gamma \circ \alpha = \alpha$ and $\gamma \circ \alpha \circ \gamma = \gamma$
- ▶ $\alpha(x) = \bigcap' \{y \in L' \mid x \sqsubseteq \gamma(y)\}$
- ▶ $\gamma(y) = \bigcup \{x \in L \mid \alpha(x) \sqsubseteq' y\}$

Properties

- ▶ $\gamma \circ \alpha \circ \gamma \circ \alpha = \gamma \circ \alpha$
- ▶ $\alpha \circ \gamma \circ \alpha \circ \gamma = \alpha \circ \gamma$
- ▶ We assume that $\{(\alpha_i, \gamma_i) | i \in \{1 \dots n\}\}$ is a family of Galois connections :

$$L_1 \xleftrightarrow[\alpha_1]{\gamma_1} L_2 \xleftrightarrow[\alpha_2]{\gamma_2} \dots L_{n-1} \xleftrightarrow[\alpha_{n-1}]{\gamma_{n-1}} L_n$$

Then $(\alpha_1; \dots; \alpha_i; \dots; \alpha_{n-1}, \gamma_{n-1}; \dots; \gamma_i; \dots; \gamma_1)$ is a Galois connection. or equivalently

$$L_1 \xleftrightarrow[\alpha_{n-1} \circ \dots \circ \alpha_i \circ \dots \circ \alpha_1]{\gamma_1 \circ \dots \circ \gamma_i \circ \dots \circ \gamma_{n-1}} \text{ is a Galois connection.}$$

- ▶ We assume that $\{(\alpha_i, \gamma_i) | i \in \{1, 2\}\}$ two Galois connections :
 $\alpha_1 = \alpha_2$ if, and only if, $\gamma_1 = \gamma_2$

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- ▶ We consider a transition system (S, I, t) where S is the set of states, I is the set of initial states and t is a binary relation over S .
- ▶ A property P of the transition system is a subset of S : $P \subseteq S$.
- ▶ P holds in $s \in S$, when $s \in P$.
- ▶ Four operators over properties can be defined as follows :

- $\text{pre}[t]P \stackrel{\text{def}}{=} \{s | s \in S \wedge \exists s'. ((s, s') \in t \wedge s' \in P)\}$
- $\tilde{\text{pre}}[t]P \stackrel{\text{def}}{=} \{s | s \in S \wedge \forall s'. ((s, s') \in t \Rightarrow s' \in P)\}$
- $\text{post}[t]P \stackrel{\text{def}}{=} \{s | s \in S \wedge \exists s'. ((s', s) \in t \wedge s' \in P)\}$
- $\tilde{\text{post}}[t]P \stackrel{\text{def}}{=} \{s | s \in S \wedge \forall s'. ((s', s) \in t \Rightarrow s' \in P)\}$

- ▶ Duality of operators :

① $\tilde{\text{pre}}[t]\neg P = \neg \text{pre}[t]P$

② $\tilde{\text{post}}[t]\neg P = \neg \text{post}[t]P$

- ▶ Galois connections over \mathcal{P} , the set of subsets of S :

$$(\mathcal{P}, \subseteq) \xrightleftharpoons[\text{pre}[t]]{\tilde{\text{post}}[t]} (\mathcal{P}, \subseteq)$$

$$(\mathcal{P}, \subseteq) \xrightleftharpoons[\text{post}[t]]{\tilde{\text{pre}}[t]} (\mathcal{P}, \subseteq)$$

- ▶ Let two sets \mathcal{L} standing for labels et \mathcal{M} standing for memories.
- ▶ First step :
 - \sqsubseteq is the partial ordering over functions using the subset relationship over function graphs : $f \sqsubseteq g$ means that $\text{Graph}(f) \subseteq \text{Graph}(g)$.
 - $\alpha_1 = \lambda P. \lambda l. \{m \mid (l, m) \in P\}$
 - $\gamma_1 = \lambda Q. \{(l, m) \mid l \in \mathcal{L} \wedge m \in Q(l)\}$
 - $(\mathcal{P}(\mathcal{L} \times \mathcal{M}), \sqsubseteq) \xleftrightarrow[\alpha_1]{\gamma_1} (\mathcal{L} \longrightarrow \mathcal{P}(\mathcal{M}), \sqsubseteq)$ is a Galois connection
- ▶ Second step :
 - Let two sets $Pred$, set of predicates, and \mathcal{M} , a set of memories.
 - The relationship between both sets is stating as follows : For any given predicate p and any given memory m , p holds in m .
 - We define $B(p) = \{m \mid m \in \mathcal{M} \wedge p(m)\}$, set of memories in which p holds.
 - Next we define :
 - ▶ $\alpha_2 = \lambda Q. \{p \mid p \in Pred \wedge Q \subseteq B(p)\}$
 - ▶ $\gamma_2 = \lambda P. \bigcap \{B(p) \mid p \in P\}$
 - $(\mathcal{P}(\mathcal{M}), \sqsubseteq) \xleftrightarrow[\alpha_2]{\gamma_2} (\mathcal{P}(Pred), \Rightarrow)$ is a Galois connection.

► Third step

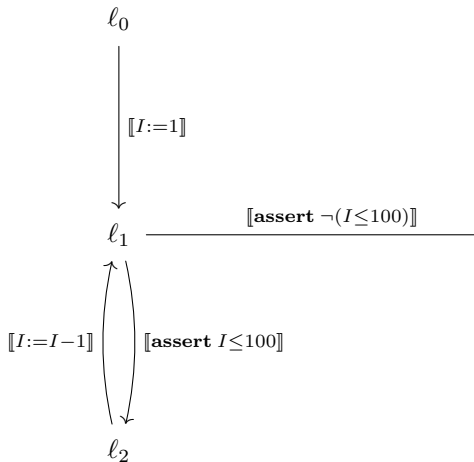
- $\alpha_3 = \lambda\ell.\alpha_2(Q_\ell) : Q \subseteq_1 Q' \stackrel{def}{=} \forall \ell \in \mathcal{L}. Q_\ell \subseteq Q'_\ell.$
- $\gamma_3 = \lambda\ell.\gamma_2(P_\ell) : P \Rightarrow_1 P' \stackrel{def}{=} \forall \ell \in \mathcal{L}. P_\ell \Rightarrow P'_\ell.$
- $(\mathcal{L} \longrightarrow \mathcal{P}(\mathcal{M}), \subseteq_1) \xrightleftharpoons[\alpha_3]{\gamma_3} (\mathcal{L} \longrightarrow \mathcal{P}(Pred), \Rightarrow_1)$ is a Galois connection.

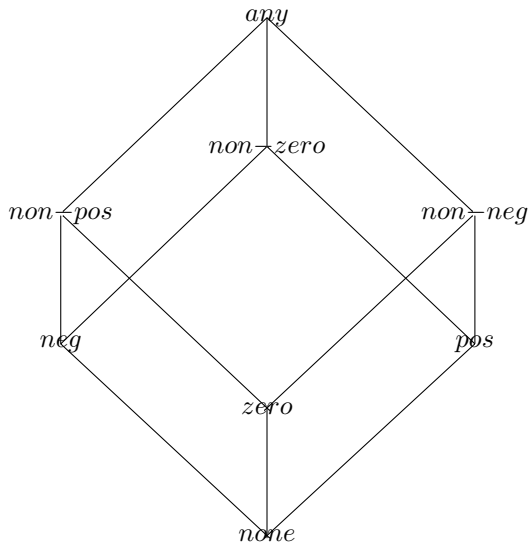
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Examples of Abstractions

```
 $\ell_0[I := 1];$   
while  $\ell_1[I \leq 100]$  do  
   $\ell_2[I := I + 1];$   
end while  
 $\ell_{final}[skip]$ 
```





- ▶ Defining an abstraction for integers

$\alpha \in \mathcal{P}(\mathbb{Z}) \longrightarrow \text{Signs}$

- ▶
$$\begin{cases} z & \alpha(z) \\ z < 0 & \text{neg} \\ z > 0 & \text{pos} \\ z = 0 & \text{zero} \end{cases}$$

- ▶ Abstraction by projection :

$$(\mathcal{P}(Var \longrightarrow \mathbb{Z}), \subseteq) \xleftrightarrow[\alpha_\pi]{\gamma_\pi} (Var \longrightarrow \mathcal{P}(\mathbb{Z}), \subseteq)$$

- ▶ Abstraction of signs

$$(Var \longrightarrow \mathcal{P}(\mathbb{Z}), \subseteq) \xleftrightarrow[\alpha_{sign}]{\gamma_{sign}} (Var \longrightarrow Signs), \subseteq)$$

- ▶ Composition of abstractions :

$$(\mathcal{P}(Var \longrightarrow \mathbb{Z}), \subseteq) \xleftrightarrow[\alpha_{sign} \circ \alpha_\pi]{\gamma_\pi \circ \gamma_{sign}} (Var \longrightarrow Signs), \subseteq)$$

- ▶ $\alpha = \alpha_{sign} \circ \alpha_\pi$ and $\gamma = \gamma_\pi \circ \gamma_{sign}$

- L is the concrete domain and L' is the abstract model :

$$\begin{array}{ccc} L & \xrightarrow{f} & L \\ \gamma \uparrow & & \downarrow \alpha \\ L' & \xrightarrow{f'} & L' \end{array}$$

$$f' = \alpha \circ f \circ \gamma \quad (2)$$

f' is the best approximation of f

- ▶ Concrete states : $cv \in Var \longrightarrow \mathcal{P}(\mathbb{Z})$: if X is in Var , then $cv(X) \in \mathcal{P}(\mathbb{Z})$.
- ▶ Abstract states : $av \in Var \longrightarrow Signs$: if X is in Var , then $av(X) \in Signs$.
- ▶ (α, γ) is extended as :
 (α_1, γ_1) entre $(Var \longrightarrow \mathcal{P}(\mathbb{Z}), \subseteq)$ et $(Var \longrightarrow Signs, \sqsubseteq)$. En particulier, $\alpha_1(cv) = av$ et, pour tout X de Var , $av(X) = \alpha(cv(X))$; $\gamma_1(av) = cv$ et, pour tout X de Var , $cv(X) = \gamma(av(X))$.
- ▶ Any expression e can be evaluated on each domain :
 - concrete domain : $States = Var \longrightarrow \mathcal{P}(\mathbb{Z})$:
 $\llbracket e \rrbracket \in (Var \longrightarrow \mathcal{P}(\mathbb{Z})) \longrightarrow \mathcal{P}(\mathbb{Z})$ and $\llbracket e \rrbracket(cv)$
 - abstract domain : $AStates = Var \longrightarrow Signs$:
 $\llbracket e \rrbracket_a \in (Var \longrightarrow Signs) \longrightarrow Signs$ and $\llbracket e \rrbracket_a(av)$.

- ▶ The best abstraction is simply dedined as follows :

$$\llbracket e \rrbracket_{best}(av) = \alpha \circ \llbracket e \rrbracket \circ \gamma_1(av).$$

- ▶ Applying the best approximation for assignment :

$$\llbracket x := e \rrbracket_{best}(av) = \begin{cases} av(y), y \neq x \\ \llbracket e \rrbracket_{best}(av) \end{cases}$$

- ▶ $(\mathcal{P}(Var \longrightarrow \mathbb{Z}), \subseteq) :$

$$A, B \in \mathcal{P}(\mathbb{Z}) : A+B = \{a+b | a \in A \wedge b \in B\}$$

- ▶ $(Var \longrightarrow Signs), \subseteq) :$

$$x, y \in Signs : x \oplus y = \alpha(\gamma(x) + \gamma(y))$$

- ▶ examples :

- $pos \oplus neg = \alpha(\gamma(pos) + \gamma(neg)) = \alpha((1, +\infty) + (-\infty, -1)) = \alpha((-\infty, +\infty)) = any$
- $pos \oplus zero = \alpha(\gamma(pos) + \gamma(zero)) = \alpha((1, +\infty) + (0)) = \alpha((1, +\infty)) = pos$
- Building a table for the abstract operation \oplus .

- ▶ Applying the analysis on the example

$\ell_0[X := 1];$
 $\ell_1[Y := 5];$
 $\ell_2[X := X+1];$
 $\ell_3[Y := Y-1];$
 $\ell_4[X := Y+X];$
 $\ell_{final}[skip];$

| ℓ | X | Y |
|----------------|----------------|----------------|
| ℓ_0 | <i>any</i> | <i>any</i> |
| ℓ_1 | <i>pos</i> | <i>any</i> |
| ℓ_2 | <i>pos</i> | <i>pos</i> |
| ℓ_3 | <i>pos</i> | <i>pos</i> |
| ℓ_4 | <i>pos</i> | <i>non-neg</i> |
| ℓ_{final} | <i>non-neg</i> | <i>non-neg</i> |

- ▶ ℓ_3 to ℓ_4 : abstract value of Y is *pos* and by γ , we obtain $(1, +\infty)$
and now we can compute in concrete domain \mathbb{Z}
 $(1, +\infty) + (-1) = (0, +\infty)$. By reapplying α we obtain *non-neg*.
- ▶ Computations may be not computable and one should use techniques for accelerating the convergence like widening.
- ▶ Computing is still costly : computing now in the abstraction and defining a sound approximation of f .

- ▶ Evaluation is using the *best* approximation :

$$\llbracket e \rrbracket_{best}(av) = \alpha \circ \llbracket e \rrbracket \circ \gamma_1(av)$$

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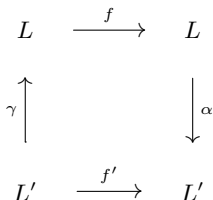
- ▶ Evaluation is using the *best* approximation :
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- ▶ Idea : approximation of the *best* approximation : $\llbracket e \rrbracket_a$ and, for any av abstract state, $\llbracket e \rrbracket_{best}(av) \sqsubseteq \llbracket e \rrbracket_a(av)$.
- ▶ Abstract semantics is defined as follows :
 $av \in Var \longrightarrow Signs :$
 - $\llbracket const \rrbracket_a(av) = \alpha(\{c\})$
 - $\llbracket x \rrbracket_a(av) = av(x)$
 - $\llbracket e_1 + e_2 \rrbracket_a(av) = \llbracket e_1 \rrbracket_a(av) \oplus \llbracket e_2 \rrbracket_a(av)$
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 - $\llbracket e_1 * e_2 \rrbracket_a(av) = \llbracket e_1 \rrbracket_a(av) \otimes \llbracket e_2 \rrbracket_a(av)$
- ▶ $\ell[X := E] : \llbracket E \rrbracket_a$ in av ou encore $\llbracket E \rrbracket_a(av) :$
$$\llbracket Y + X + 6 \rrbracket_a(av) = \llbracket Y \rrbracket_a(av) +_a \llbracket X \rrbracket_a(av) +_a \llbracket 6 \rrbracket_a(av).$$
 - $\llbracket Y - 1 \rrbracket_a(av) = \llbracket Y \rrbracket_a(av) \oplus \llbracket -1 \rrbracket_a(av) = pos \oplus neg = any$
 - $\llbracket Y - 1 \rrbracket_{best}(av) = \alpha \circ \llbracket Y - 1 \rrbracket \circ \gamma_1(av) = \alpha(\llbracket Y - 1 \rrbracket(\gamma_1(av))) = \alpha(\llbracket Y - 1 \rrbracket(\{Y \mapsto (1, +\infty)\})) = \alpha((1 + \infty) + (-1)) = \alpha((0, +\infty)) = non - neg$

Sound approximations of f with respect to a Galois connection

A sound approximation of f with respect to a Galois connection f' satisfies the following property :

$$\forall x \in L, y \in L'. \alpha(x) \sqsubseteq y \Rightarrow \alpha(f(x)) \sqsubseteq f'(y)$$



The four statements are equivalent

- ▶ f' is a sound approximation of f with respect to a Galois connection
- ▶ $\alpha \circ f \sqsubseteq' f' \circ \alpha$
- ▶ $\alpha \circ f \circ \gamma \sqsubseteq' f'$
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- ▶ $\llbracket e \rrbracket_{best}(av) = \alpha \circ \llbracket e \rrbracket \circ \gamma_1(av)$ provide the best abstraction but is costly.
- ▶ Another solution is to define an abstract semantics for expressions : $\llbracket e \rrbracket_a$ such that for any av , $\llbracket e \rrbracket_{best}(av) \sqsubseteq \llbracket e \rrbracket_a(av)$.
- ▶ $av \in Var \longrightarrow Signs :$
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- ▶ $\llbracket Y-1 \rrbracket_a(av) = \llbracket Y \rrbracket_a(av) \oplus \llbracket -1 \rrbracket(av)_a = pos \oplus neg = may$
- ▶ $\begin{aligned} \llbracket Y-1 \rrbracket_{best}(av) &= \alpha_1 \circ \llbracket Y-1 \rrbracket \circ \gamma_1(av) = \alpha_1(\llbracket Y-1 \rrbracket(\gamma_1(av))) = \\ &= \alpha_1(\llbracket Y-1 \rrbracket(\{Y \mapsto (1, +\infty)\})) = \alpha_1((1+\infty)+(-1)) = \\ &= \alpha_1((0, +\infty)) = non-neg \end{aligned}$

- ▶ Applying the analysis on the example

$\ell_0[X := 1];$
 $\ell_1[Y := 5];$
 $\ell_2[X := X + 1];$
 $\ell_3[Y := Y - 1];$
 $\ell_4[X := Y + X];$
 $\ell_{final}[skip];$

| ℓ | X | Y |
|----------------|------------|------------|
| ℓ_0 | <i>any</i> | <i>any</i> |
| ℓ_1 | <i>pos</i> | <i>any</i> |
| ℓ_2 | <i>pos</i> | <i>pos</i> |
| ℓ_3 | <i>pos</i> | <i>pos</i> |
| ℓ_4 | <i>pos</i> | <i>any</i> |
| ℓ_{final} | <i>any</i> | <i>any</i> |

- ▶ The new analysis is less precise but more efficient since we compute in the domain of signs.

- 1 Introduction
- 2 Example of analysis
- 3 Static analysis
- 4 Overview of the methodology
- 5 Standard, Collecting and Abstract Semantics
- 6 Finding Sound Abstractions for Computing
- 7 Galois Connections
- 8 Examples of Galois connections
- 9 Domain of Signs

- ▶ $\mathbb{I}(\mathbb{Z}) = \{\perp\} \cup \{[l, u] \mid l \in \mathbb{Z} \cup \{-\infty\}, u \in \mathbb{Z} \cup \{\infty\}, l \leq u\}$
- ▶ $[l_1, u_1] \sqsubseteq [l_2, u_2]$ si, et seulement si, $l_2 \leq l_1$ et $u_1 \leq u_2$.
- ▶ $(\mathbb{I}(\mathbb{Z}), \sqsubseteq)$ est une structure partiellement ordonnée.
- ▶
 - ① $[l_1, u_1] \sqcup [l_2, u_2] = [\min(l_1, l_2), \max(u_1, u_2)]$
 - ② $[l_1, u_1] \sqcap [l_2, u_2] = \begin{cases} [\max(l_1, l_2), \min(u_1, u_2)] \\ \perp, \text{ si } \max(l_1, l_2) > \min(u_1, u_2) \end{cases}$
- ▶ $(\mathbb{I}(\mathbb{Z}), \sqcup)$ is a complete lattice.
- ▶
 - ① $\alpha(X) = \begin{cases} [\min(X), \max(X)] \\ \perp, \text{ si } X = \emptyset \end{cases}$
 - ② $\gamma([l, u]) = [l..u]$ et $\gamma(\perp) = \emptyset$
- ▶ (α, γ) is a Galois connexion.
- ▶
 - ① $i_1 \oplus i_2 = [l_1 + l_2, u_1 + u_2]$
 - ② $i_1 \ominus i_2 = [l_1 - u_2, u_1 - l_2]$
 - ③ $i_1 \otimes i_2 = [\min(l_1 \cdot l_2, l_1 \cdot u_2, u_1 \cdot l_2, u_1 \cdot u_2), \max(l_1 \cdot l_2, l_1 \cdot u_2, u_1 \cdot l_2, u_1 \cdot u_2)]$
 - ④ $i_1 \oslash i_2 = [\min(l_1 / l_2, l_1 / u_2, u_1 / l_2, u_1 / u_2), \max(l_1 / l_2, l_1 / u_2, u_1 / l_2, u_1 / u_2)]$

- ▶ Computing *collecting semantics* is generally undecidable :
 - S is a safety property for \mathcal{MS} if, and only if, $\text{REACHABLE}(\mathcal{MS}) \subseteq S$.
 - Finding a sound approximation of $\text{REACHABLE}(\mathcal{MS})$, denoted $\alpha(\text{REACHABLE}(\mathcal{MS}))$, and satisfying $\gamma(\alpha(\text{REACHABLE}(\mathcal{MS}))) \subseteq S$.
 - $\text{REACHABLE}(\mathcal{MS}) \subseteq \gamma(\alpha(\text{REACHABLE}(\mathcal{MS})))$ and $\gamma(\alpha(\text{REACHABLE}(\mathcal{MS}))) \subseteq S$.
- ▶ Abstract domains can be finite as the domain of Signs but the domain of intervals is infinite : computing $\text{REACHABLE}(\mathcal{MS})$ remains undecidable but we can approximate its computation.
- ▶ Abstract domains can be infinite : we have to accelerate the computations of fixed-points in the case of loops for instance : widening and narrowing.

- L is the concrete domain and L' is the abstract model :

$$\begin{array}{ccc} L & \xrightarrow{f} & L \\ \gamma \uparrow & & \downarrow \alpha \\ L' & \xrightarrow{f'} & L' \end{array}$$

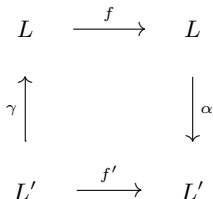
$$f' = \alpha \circ f \circ \gamma \quad (3)$$

f' is the best approximation of f

Sound approximations of f with respect to a Galois connection

A sound approximation of f with respect to a Galois connection f' satisfies the following property :

$$\forall x \in L, y \in L'. \alpha(x) \sqsubseteq y \Rightarrow \alpha(f(x)) \sqsubseteq f'(y)$$



The four statements are equivalent

- ▶ f' is a sound approximation of f with respect to a Galois connection
- ▶ $\alpha \circ f \sqsubseteq' f' \circ \alpha$
- ▶ $\alpha \circ f \circ \gamma \sqsubseteq' f'$
- ▶ $f \circ \gamma \sqsubseteq' \gamma \circ f'$

- ▶ $\llbracket e \rrbracket_{best}(av) = \alpha \circ \llbracket e \rrbracket \circ \gamma_1(av)$ provides the best abstraction but is costly.
- ▶ Another solution is to define an abstract semantics for expressions :
hide $\llbracket e \rrbracket_a$ such that for any av , $\llbracket e \rrbracket_{best}(av) \sqsubseteq \llbracket e \rrbracket_a(av)$.
- ▶ $av \in Var \longrightarrow Signs :$
 - $\llbracket const \rrbracket_a(v) = \alpha(\{c\})$
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 - $\llbracket e_1 + e_2 \rrbracket_a(v) = \llbracket e_1 \rrbracket_a(v) \oplus \llbracket e_2 \rrbracket_a(v)$
 - $\llbracket e_1 * e_2 \rrbracket_a(v) = \llbracket e_1 \rrbracket_a(v) \otimes \llbracket e_2 \rrbracket_a(v)$

- ▶ Suppose that $C \xrightleftharpoons[\alpha]{\gamma} A$ is a Galois connection
- ▶ a function $f \in C \rightarrow C$: to find a function g

$$\begin{array}{ccc} C & \xrightarrow{f} & C \\ \gamma \uparrow & & \downarrow \alpha \\ A & \xrightarrow{g} & A \end{array}$$

- ▶ f is monotone
- ▶ $g = R(\alpha, \gamma, f)$ and $f \sqsubseteq \gamma \circ g \circ \alpha$
- ▶ $f \sqsubseteq \gamma \circ g \circ \alpha$ or equivalently $\alpha \circ f \circ \gamma \sqsubseteq g$
- ▶ $g = \alpha \circ f \circ \gamma$ is the *best* approximation.

Definition of a sound approximation of a function f

A function $g \in A \rightarrow A$ is a sound approximation of a function $f \in C \rightarrow C$, if it satisfies the following condition :

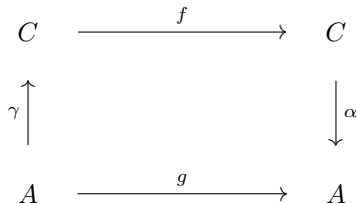
$$\forall c \in C : \forall a \in A : \alpha(c) \sqsubseteq a \Rightarrow \alpha(f(c)) \sqsubseteq g(a)$$

Properties

Suppose that $C \xrightleftharpoons[\alpha]{\gamma} A$ is a Galois connection.

The four statements are equivalent

- ① g is a sound approximation of f with respect to a Galois connection
- ② $\alpha \circ f \sqsubseteq g \circ \alpha$
- ③ $\alpha \circ f \circ \gamma \sqsubseteq g$
- ④ $f \circ \gamma \sqsubseteq \gamma \circ g$
- ⑤ $f \sqsubseteq \gamma \circ g \circ \alpha$



Best abstraction

Suppose that :

- ▶ $C \xrightleftharpoons[\alpha]{\gamma} A$ is a Galois connection.
- ▶ $f \in C \rightarrow C$ is monotonous
- ▶ $g = \alpha \circ f \circ \gamma$

Then $lfp(f) \sqsubseteq \gamma(lfp(g))$ and $\alpha(lfp(f)) \sqsubseteq lfp(g)$
or equivalently rewritten as $\mu f \sqsubseteq \gamma(\mu g)$ and $\alpha(\mu f) \sqsubseteq \mu g$

First theorem

- ▶ Suppose that $C \xleftrightarrow[\alpha]{\gamma} A$ is a Galois connection
- ▶ Two functions $f \in C \rightarrow C$ and $g \in A \rightarrow A$:

$$\begin{array}{ccc} C & \xrightarrow{f} & C \\ \gamma \uparrow & & \downarrow \alpha \\ A & \xrightarrow{g} & A \end{array}$$

- ▶ f and g are monotone
- ▶ $\alpha \circ f = g \circ \alpha$.

Then $\alpha(\mu.f) = \mu.g$.

- ① $\mu g \sqsubseteq \alpha(\mu f)$
 - $f(\mu f) = \mu f$ (fixed-point property)
 - $\alpha(f(\mu f)) = \alpha(\mu f)$ (applying the relation over f and g)
 - $\alpha(f(\mu f)) = g(\alpha(\mu f)) = \alpha(\mu f)$
 - $\alpha(\mu f)$ is a fixed-point of g and $\mu g \sqsubseteq \alpha(\mu f)$
- ② $\alpha(\mu f) \sqsubseteq \mu g$
 - Consider y a fixed-point of $g : g(y) = y$ and $\mu g \sqsubseteq y$.
 - $\gamma(y)$ is a fixed-point of f
 - $\mu f \sqsubseteq \gamma(y)$
 - $\alpha(\mu f) \sqsubseteq y$
 - $\alpha(\mu f) \sqsubseteq \mu g$

Second theorem

- ▶ Suppose that $C \xleftrightarrow[\alpha]{\gamma} A$ is a Galois connection
- ▶ Two functions $f \in C \rightarrow C$ and $g \in A \rightarrow A$:

$$\begin{array}{ccc} C & \xrightarrow{f} & C \\ \gamma \uparrow & & \downarrow \alpha \\ A & \xrightarrow{g} & A \end{array}$$

- ▶ f and g are monotone
- ▶ $\alpha \circ f \sqsubseteq g \circ \alpha$.

Then $\alpha(\mu f) \sqsubseteq \mu g$.

- ▶ $\alpha \in \mathcal{P}(\mathbb{Z}) \rightarrow \text{Signs} : \begin{cases} z & \alpha(z) \\ z < 0 & \text{neg} \\ z > 0 & \text{pos} \\ z = 0 & \text{zero} \end{cases}$
- ▶ $f \in \mathcal{P}(\mathbb{Z}) \rightarrow \mathcal{P}(\mathbb{Z})$ where $f(X) = \{0\} \cup \{x+2 \mid x \in \mathbb{Z} \wedge x \in X\}$
- ▶ $g = \alpha \circ f \circ \gamma$
- ▶ $f^0 = \emptyset, f^1 = \{0\}, f^2 = \{0, 2\}, \dots$
- ▶ $g(\perp) = \perp, g^1 = \alpha \circ f \circ \gamma(\perp) = [0, \infty[, g^2 = [0, \infty[, \dots$ and $\forall i \geq 2 : g^i = [0, \infty[.$
- ▶ $\mu.g = [0, \infty[$

Definition

∇ is a widening operator over (L, \sqsubseteq) ($\nabla \in L \times L \rightarrow L$)

- ▶ For any x and y in L : $x \sqcup y \sqsubseteq x \nabla y$
- ▶ For any sequence $x_0 \sqsubseteq x_1 \sqsubseteq x_2 \sqsubseteq x_3 \dots \sqsubseteq x_i \sqsubseteq x_{i+1} \dots$, the sequence $\{y_i | i \in \mathbb{N}\}$
 - $y_0 = x_0$
 - $y_{i+1} = y_i \nabla x_{i+1}$

stabilizes after a finite amount of time.

Theorem

If ∇ is a widening operator over (L, \sqsubseteq) ($\nabla \in L \times L \rightarrow L$), then the ascending sequence $x_0 \sqsubseteq x_1 \sqsubseteq x_2 \sqsubseteq x_3 \dots \sqsubseteq x_i \sqsubseteq x_{i+1} \dots$ defined by :

- ▶ $x_0 = \perp$
- ▶ $x_{i+1} = x_i \nabla f(x_i)$

is eventually stationary and its limit satisfies $lfp(f) \sqsubseteq\sqsubseteq \{x_i | i \in \mathbb{N}\}$ stabilizes after a finite amount of time.

- ▶ Using ∇ instead of \sqsubseteq for computing approximation of upper bound.

Intervals

- ▶ $\perp \nabla \perp = \perp$
- ▶ $\perp \nabla (l, u) = (l, u) \nabla \perp = (l, u)$
- ▶ $(l1, u1) \nabla (l2, u2) = \left(\begin{pmatrix} -\infty & \text{if } l2 < l1 \\ l1 \end{pmatrix}, \begin{pmatrix} \infty & \text{if } u2 > u1 \\ u1 \end{pmatrix} \right)$

- ▶ $\mathbb{I}(\mathbb{Z}) = \{\perp\} \cup \{[l, u] \mid l \in \mathbb{Z} \cup \{-\infty\}, u \in \mathbb{Z} \cup \{\infty\}, l \leq u\}$
- ▶ $(\mathbb{I}(\mathbb{Z}), \sqsubseteq)$ est une structure partiellement ordonnée.
- ▶ $[l_1, u_1] \nabla [l_2, u_2] = [\text{cond}(l_2 < l_1, -\infty, l_1), \text{cond}(u_1 < u_2, \infty, u_1)]$
- ▶ $[2, 3] \nabla [1, 4] = [-\infty, \infty]$
- ▶ $[0, 1] \sqsubseteq [0, 3]$
- ▶ $[0, 1] \nabla [0, 3] = [0, \infty]$.
- ▶ $[0, 3] \nabla [0, 2] = [0, 3]$.
- ▶ $[0, 2] \nabla ([0, 1] \nabla [0, 2]) = [0, \infty]$
- ▶ $([0, 2] \nabla [0, 1]) \nabla [0, 2] = [0, 2]$

Let us assume that (L, \sqsubseteq) is a complete lattice and f is a monotonic function defined from L to L .

Theorem

If $\nabla \in L \times L \rightarrow L$ is a widening operator, then the sequence $\{x_i | i \in \mathbb{N}\}$ defined by

- ▶ $x_0 = \perp$
- ▶ $x_{i+1} = \begin{cases} x_i & \text{if } f(x_i) \sqsubseteq x_i \\ x_i \nabla f(x_i) & \text{otherwise} \end{cases}$

is eventually stationary and its limit satisfies $lfp(f) \sqsubseteq \bigcup \{x_i | i \in \mathbb{N}\}$

Definition of a sound approximation of a function f

A function $g \in A \rightarrow A$ is a sound approximation of a function $f \in C \rightarrow C$, if it satisfies the following condition :

$$\forall c \in C : \forall a \in A : \alpha(c) \sqsubseteq a \Rightarrow \alpha(f(c)) \sqsubseteq g(a)$$

Properties

Suppose that $C \xrightleftharpoons[\alpha]{\gamma} A$ is a Galois connection.

The four statements are equivalent

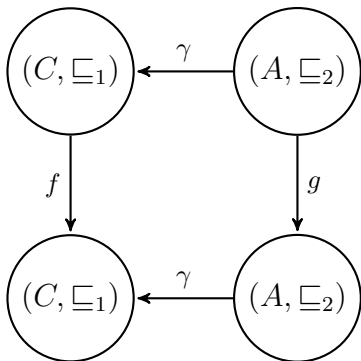
- ① g is a sound approximation of f with respect to a Galois connection
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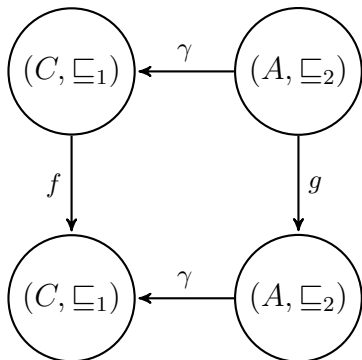
Example of a sound approximation of the invariant of a system

- ▶ C is the set of concrete states : $cv \in Var \longrightarrow \mathcal{P}(\mathbb{Z})$: if X is in Var , then $cv(X) \in \mathcal{P}(\mathbb{Z})$.
- ▶ A is the set of abstract states : $av \in Var \longrightarrow Signs$: if X is in Var , then $av(X) \in Signs$.
- ▶ (α, γ) is extended as :
 (α_1, γ_1) entre $(Var \longrightarrow \mathcal{P}(\mathbb{Z}), \subseteq)$ et $(Var \longrightarrow Signs, \sqsubseteq)$. En particulier, $\alpha_1(cv) = av$ et, pour tout X de Var ,
 $av(X) = \alpha(cv(X))$; $\gamma_1(av) = cv$ et, pour tout X de Var ,
 $cv(X) = \gamma(av(X))$.

Computing the set of computing states of a transition system

- ▶ $Init \subseteq C$ is the set of initial states.
- ▶ **NEXT** defines the transition over concrete states
- ▶ $REACHABLE(TS) = \{u | u \in C \wedge (\exists x_0. x_0 \in C \wedge (x_0 \in Init) \wedge NEXT^*(x_0, x))\}$
- ▶ pour toute partie U de Σ , $U = FP(U)$
- ▶ pour toute partie U de Σ , $FP(U) = Init_S \cup \longrightarrow [U]$



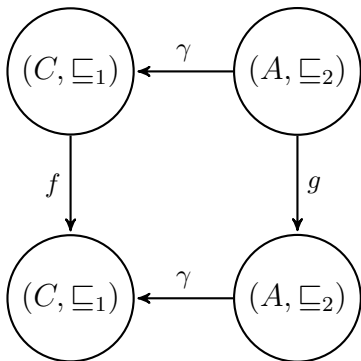


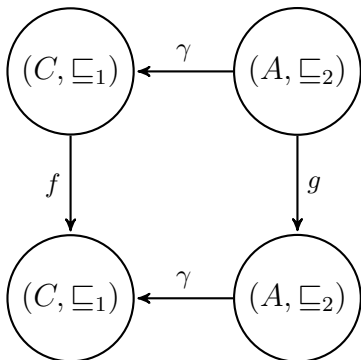
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Second Theorem

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- ▶ $\perp \nabla \perp = \perp$
- ▶ $\perp \nabla (l, u) = (l, u) \nabla \perp = (l, u)$
- ▶ $(l_1, u_1) \nabla (l_2, u_2) = \left(\begin{pmatrix} -\infty \text{ if } l_2 < l_1 \\ l_1 \end{pmatrix}, \begin{pmatrix} \infty \text{ if } u_2 > u_1 \\ u_1 \end{pmatrix} \right)$

- ▶ $\mathbb{I}(\mathbb{Z}) = \{\perp\} \cup \{[l, u] \mid l \in \mathbb{Z} \cup \{-\infty\}, u \in \mathbb{Z} \cup \{\infty\}, l \leq u\}$
- ▶ $(\mathbb{I}(\mathbb{Z}), \sqsubseteq)$ est une structure partiellement ordonnée.
- ▶ $[l_1, u_1] \nabla [l_2, u_2] = [\text{cond}(l_2 < l_1, -\infty, l_1), \text{cond}(u_1 < u_2, \infty, u_1)]$
- ▶ $[2, 3] \nabla [1, 4] = [-\infty, \infty]$
- ▶ $[0, 1] \sqsubseteq [0, 3]$
- ▶ $[0, 1] \nabla [0, 3] = [0, \infty]$.
- ▶ $[0, 3] \nabla [0, 2] = [0, 3]$.
- ▶ $[0, 2] \nabla ([0, 1] \nabla [0, 2]) = [0, \infty]$
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ℓ0 :  
y := -11;  
ℓ0 :  
IF  x < y  THEN  
    ℓ1 :  
    z := y;  
    ℓ2 :  
ELSE  
    ℓ3 :  
    z := x;  
    ℓ4 :  
FI  
ℓ5 :
```

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```

Result : $y < 0 \quad x \in \mathbb{Z} \quad z < 0$ means that $z < 0$ is an information resulting from the analysis over abstract domain of

- MALG & MOVEX 120/135

Computing the least fixed-point over a finite lattice

```
INPUT   $tf \in T \longrightarrow T$ 
OUTPUT  $result = \mu.f$ 
VARIABLES   $x, y \in T, i \in \mathbb{N}$ 
 $\ell_0 : \{x, y \in T\}$ 
 $x := \perp;$ 
 $y := \perp;$ 
 $i := 0;$ 
 $\ell_{11} : \{x, y \in T \wedge x = F^i \wedge y = \bigcup_{k=0; k=i} F^k \wedge i \leq Card(T) \wedge i = 0\};$ 
WHILE   $i \leq Card(T)$ 
   $\ell_1 : \{x, y \in T \wedge x = F^i \wedge y = \bigcup_{k=0; k=i} F^k \wedge i \leq Card(T)\};$ 
   $x := f(x);$ 
   $\ell_2 : \{x, y \in T \wedge x = F^{i+1} \wedge y = \bigcup_{k=0; k=i} F^k \wedge i \leq Card(T)\};$ 
   $y := x \sqcup y;$ 
   $\ell_3 : \{x, y \in T \wedge x = F^{i+1} \wedge y = \bigcup_{k=0; k=i+1} F^k \wedge i \leq Card(T)\};$ 
   $i := i+1;$ 
   $\ell_4 : \{x, y \in T \wedge x = F^i \wedge y = \bigcup_{k=0; k=i} F^k \wedge i \leq Card(T)+1\};$ 
OD;
 $\ell_5 : \{x, y \in T \wedge x = F^i \wedge y = \bigcup_{k=0; k=i} F^k \wedge i = Card(T)+1\};$ 
 $result := y;$ 
 $\ell_6 : \{x, y \in T \wedge x = F^i \wedge y = \bigcup_{k=0; k=i} F^k \wedge i = Card(T)+1 \wedge result = y\};$ 
```

- ▶ Abstract interpretation is a general framework for defining sound approximation of the semantics of computer programs, based on monotonic functions over ordered sets, especially lattices.
- ▶ Main concrete application is formal static analysis, the automatic extraction of information about the possible executions of computer programs.
- ▶ When defining an abstract domain, it can be finite (domain of signs) or infinite (domain of intervals) : it means that we have to manage undecidability questions for computing fixed-points.
- ▶ interproc is a tool that can be used for analysing recursive programs and for playing with abstract interpretation.

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```

Result

$y < 0, x \in \mathbb{Z}, z < 0$ means that $z < 0$ is an information

- MALG & MOVEX 130/135

Computing the least fixed-point over a finite lattice

```
INPUT   $tf \in T \longrightarrow T$ 
OUTPUT  $result = \mu.f$ 
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$$\mathcal{D}, \mathcal{S} \quad \text{SATISFAIT} \quad \mathcal{R}$$

- ▶ \mathcal{R} : pre/post.
- ▶ \mathcal{D} : entiers, réels, ...
- ▶ \mathcal{S} : code, procédure, programme, ...

$$\mathcal{D}, \text{ALG} \quad \text{SATISFAIT} \quad \left\{ \begin{array}{l} \mathbf{pre}(\text{ALG})(v) \\ \mathbf{post}(\text{ALG})(v_0, v) \end{array} \right.$$



Vérification de conditions de vérification

| |
|-------------------------------------|
| \mathcal{D} |
| <hr/> |
| $\mathbf{pre}(\text{ALG})(v)$ |
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| <hr/> |
| ALG |

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- Vérification des conditions de vérification avec un model-checker par exploration de tous les états.

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- ▶ Vérification des conditions de vérification avec un model-checker par exploration de tous les états.
- ▶ Vérification des conditions de vérification avec un outil de preuve formelle.

\mathcal{D}, ALG SATISFAIT $\left\{ \begin{array}{l} \text{requires ALG}(v) \\ \text{ensures ALG}(v_0, v) \end{array} \right.$



Vérification de conditions de vérification

| |
|-------------------------|
| \mathcal{D} |
| requires ALG(v) |
| ensures ALG(v_0, v) |
| ALG |

$$\mathcal{D}, \text{ALG} \text{ SATISFAIT } \begin{cases} \text{requires ALG}(v) \\ \text{ensures ALG}(v_0, v) \end{cases}$$



Vérification de conditions de vérification

| |
|------------------------------|
| \mathcal{D} |
| <hr/> |
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| <hr/> |
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- ▶ Vérification des conditions de vérification avec un outil de preuve formelle QED
- ▶ Vérification des conditions de vérification avec un outil de preuve formelle Alt-Ergo