Introduction to MPI

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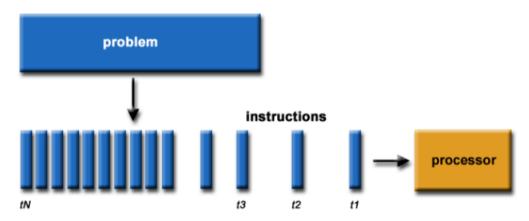
What is Parallel Computing?

Serial Computing:

- Traditionally, software has been written for serial computation:
 - A problem is broken into a discrete series of instructions
 - Instructions are executed sequentially one after another
 - Executed on a single processor
 - Only one instruction may execute at any moment in time

What is Parallel Computing? ...

Serial Computing:



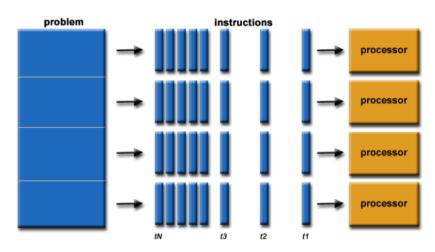
What is Parallel Computing . . .

Parallel Computing:

- In the simplest sense, parallel computing is the simultaneous use of multiple compute resources to solve a computational problem:
 - A problem is broken into discrete parts that can be solved concurrently
 - Each part is further broken down to a series of instructions
 - Instructions from each part execute simultaneously on different processors
 - An overall control/coordination mechanism is employed

What is Parallel Computing . . .

Parallel Computing:



What is Parallel Computing? ...

- The computational problem should be able to:
 - Be broken apart into discrete pieces of work that can be solved simultaneously;
 - Execute multiple program instructions at any moment in time;
 - Be solved in less time with multiple compute resources than with a single compute resource.
- The compute resources are typically:
 - A single computer with multiple processors/cores
 - An arbitrary number of such computers connected by a network

Parallel Computer Memory Architectures

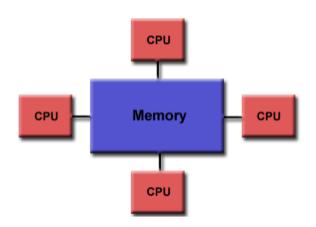
Shared Memory

General Characteristics:

- Shared memory parallel computers vary widely, but generally have in common the ability for all processors to access all memory as global address space.
- Multiple processors can operate independently but share the same memory resources.
- Changes in a memory location effected by one processor are visible to all other processors.
- Historically, shared memory machines have been classified as UMA and NUMA, based upon memory access times.

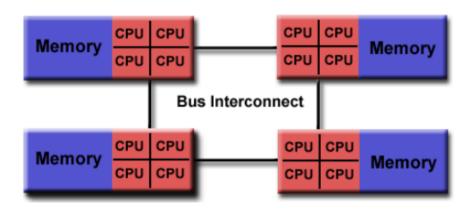
Parallel Computer Memory Architectures . . .

Uniform Memory Access (UMA) Shared Memory



Parallel Computer Memory Architectures . . .

Non-Uniform Memory Access (NUMA) Shared Memory

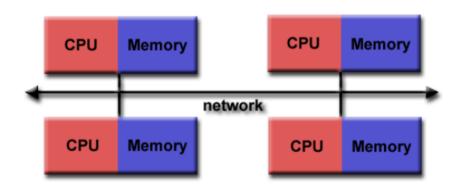


Parallel Computer Memory Architectures

• General Characteristics:

- Distributed memory systems require a communication network to connect inter-processor memory.
- Processors have their own local memory
- Memory addresses in one processor do not map to another processor, so there is no concept of global address space across all processors.
- The network "fabric" used for data transfer varies widely, though it can be as simple as Ethernet.

Parallel Computer Memory Architectures . . .

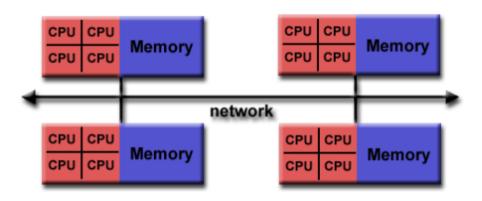


Hybrid Distributed-Shared Memory

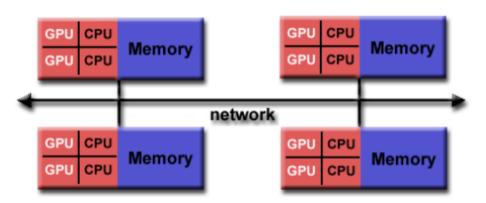
General Characteristics:

- The largest and fastest computers in the world today employ both shared and distributed memory architectures.
- The shared memory component can be a shared memory machine and/or graphics processing units (GPU).
- The distributed memory component is the networking of multiple shared memory/GPU machines
- Network communications are required to move data from one machine to another.
- Current trends seem to indicate that this type of memory architecture will continue to prevail and increase at the high end of computing for the foreseeable future.

Hybrid Distributed-Shared Memory . . .



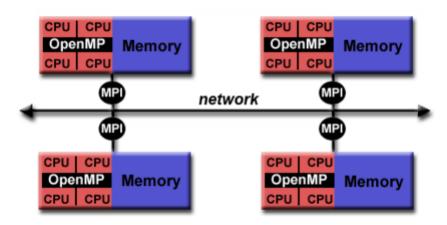
Hybrid Distributed-Shared Memory . . .



Hybrid Parallel Programming Models

- A hybrid model combines more than one of the previously described programming models.
- Currently, a common example of a hybrid model is the combination of the message passing model (MPI) with the threads model (OpenMP).
 - Threads perform computationally intensive kernels using local, on-node data
 - Communications between processes on different nodes occurs over the network using MPI
- This hybrid model lends itself well to the most popular hardware environment of clustered multi/many-core machines.
- Another similar and increasingly popular example of a hybrid model is using MPI with CPU-GPU
 - MPI tasks run on CPUs using local memory and communicating with each other over a network.
 - Computationally intensive kernels are off-loaded to GPUs on-node.
 - Data exchange between node-local memory and GPUs uses CUDA

Hybrid Parallel Programming Models . . .



- An Interface Specification:
- MPI = Message Passing Interface
- MPI primarily addresses the message-passing parallel programming model
- Simply stated, the goal of the Message Passing Interface is to provide a widely used standard for writing message passing programs. The interface attempts to be:
- The MPI standard has gone through a number of revisions, with the most recent version being MPI-4.x
- Programming Model:
 - Originally, MPI was designed for distributed memory architectures
 - Today, MPI runs on virtually any hardware platform:
 - Distributed Memory
 - Shared Memory
 - Hybrid

Implementation

- MPI is a library of function/subroutine calls
- MPI is not a language
- There is no such thing as an MPI compiler
- The C or Fortran compiler you invoke knows nothing about what MPI actually does
 - only knows prototype/interface of the function/subroutine calls

Reasons for Using MPI:

- Standardization MPI is the only message passing library that can be considered a standard.
- Portability There is little or no need to modify your source code when you port your application to a different platform that supports the MPI standard.
- Functionality There are over 430 routines defined in MPI-3, which includes the majority of those in MPI-2 and MPI-1.
- Availability A variety of implementations are available, both vendor and public domain.

MPI Implementations and Compilers

- Although the MPI programming interface has been standardized, actual library implementations will differ.
- MPI library implementations on LC systems vary, as do the compilers they are built for. These are summarized in the table below:

MPI Library	Where?	Compilers
MPICH	Linux clusters	GNU, Intel, PGI, Clang
Open MPI	Linux clusters	GNU, Intel, PGI, Clang
Intel MPI	Linux clusters	Intel, GNU
IBM BG/Q MPI	BG/Q clusters	IBM, GNU
IBM Spectrum MP	Coral Early Access and	IBM, GNU, PGI, Clang
	Sierra clusters	

Message Passing Model

- The message passing model is based on the notion of processes
 - an think of a process as an instance of a running program, together with the program's data
- In the message passing model, parallelism is achieved by having many processes co-operate on the same task
- Each process has access only to its own data
 - i.e all variables are private
- Processes communicate with each other by sending and receiving messages
 - typically library calls from a conventional sequential language

Single-Program-Multiple-Data(SPMD)

- Most message passing programs use the Single-Program-Multiple-Data (SPMD) model
- All processes run (their own copy of) the same program
- Each process has a separate copy of the data
- To make this useful, each process has a unique identifier
- Processes can follow different control paths through the program, depending on their process ID
- Usually run one process per processor / core

```
Writing MPI programs
//hello.c program
#include "mpi.h"
#include <stdio.h>
int main(int argc , char argv){
MPI_Init(&argc, &argv);
 print("Hell world!\n");
 MPI Finalize();
 return 0:
```

Writing MPI programs

- #include "mpi.h" provides basic MPI definitions and type
- MPI Init starts MPI
- MPI_Finalize exits MPI
- Note that all non MPI routines are local; thus the printf run on each process

Compiling, linking & Running MPI programs

- Both openmpi & MPICH implementation provides the commands mpicc for compiling and linking.
 - mpicc -o hello hello.c
- mpirun is common with several MPI implementations to run the progrma.
 - mpirun -np 4 hello

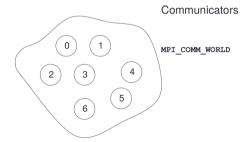
Finding out about the environment

- Two of the first questions asked in a parallel program are:
 - How many processes are there? and
 - Who am I?
- How many is answered with MPI_Comm_size
- Who am I is answered with MPI_Comm_rank.
- The rank is a number between zero and size-1.

```
MPI_Comm_size( MPI_COMM_WORLD, &numprocs );
MPI_Comm_rank( MPI_COMM_WORLD, &myid );
```

Finding out about the environment ...

- MPI_COMM_WORLD is a communicator.
- Communicators are used to separate the communication of different modules of the application.
- Communicators are essential for writing reusable libraries.



Emulating General Message Passing (C)

```
main (int argc, char **argv)
if (controller process)
Controller( /* Arguments */ );
else
Worker
( /* Arguments */ );
```

Messages

- A message transfers a number of data items of a certain type from the memory of one process to the memory of another process
- A message typically contains:
 - the ID of the sending processor
 - the ID of the receiving processor
 - the type of the data items
 - the number of data items
 - the data itself
 - a message type identifier

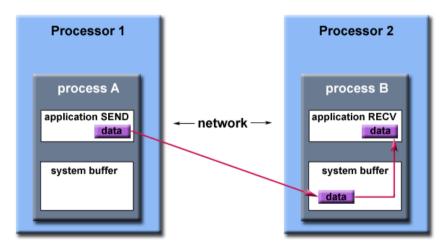
Blocking Point-to-Point Communications

- We have considered two processes:
 - one sender
 - one receiver
- This is called point-to-point communication
 - simplest form of message passing
 - relies on matching send and receive
- Close analogy to sending personal emails

Blocking point-to-point communication

- **Sending** and **receiving** are the two foundational concepts of MPI.
- Almost every single function in MPI can be implemented with basic send and receive calls.
- MPI's send and receive calls operate:
 - First, process A decides a message needs to be sent to process B.
 - Process A then packs up all of its necessary data into a buffer(envelopes) for process B.
 - The location of the message is defined by the process's rank.
 - MPI allows senders and receivers to also specify message IDs with the message (known as tags)

Blocking point-to-point communication ...



Blocking point-to-point communication . . .

• Let's look at the prototypes for the MPI sending functions.

```
MPI_Send(
    void* data,
    int count,
    MPI_Datatype datatype,
    int destination,
    int tag,
    MPI Comm communicator)
```

Blocking point-to-point communication . . .

Let's look at the prototypes for the MPI receiving functions.

```
MPI_Recv(
    void* data,
    int count,
    MPI_Datatype datatype,
    int source,
    int tag,
    MPI_Comm communicator,
    MPI Status* status)
```

Elementary MPI datatypes

MPI Datatype	C datatype
MPI_CHAR	signed char
MPI_SHORT	signed short char
MPI_INT	signed int
MPI_LONG	signed long int
MPI_UNSIGNED_CHAR	unsigned char
MPI_UNSIGNED_SHOR	T unsigned short int
MPI_FLOAT	float
MPI_DOUBLE	double
MPI_LONG_DOUBLE	long double
MPI_BYTE	char

```
MPI send / recv program
// Find out rank, size
int world rank; int world_size; int number;
MPI Comm rank(MPI COMM WORLD, &world rank);
MPI Comm size(MPI COMM WORLD, &world size);
if (world rank == 0) {
    number = -1:
    MPI Send(&number, 1, MPI INT, 1, 0, MPI COMM WORLD);
} else if (world rank == 1) {
    MPI Recv(&number, 1, MPI INT, 0, 0, MPI COMM WORLD,
             MPI STATUS IGNORE);
    printf("Process 1 received number %d from process 0\n", number);
```

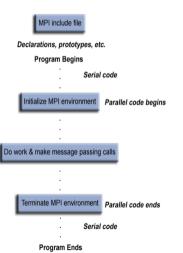
Six Function MPI

- MPI is very simple. These six functions allow you to write many programs:
 - MPI_Init
 - MPI_Finalize
 - MPI_Comm_size
 - MPI_Comm_rank
 - MPI_Send
 - MPI_Recv

Communication modes

- Sending a message can either be synchronous or asynchronous
- A synchronous send is not completed until the message has started to be received
- An asynchronous send completes as soon as the message has gone
- Receives are usually synchronous the receiving process must wait until the message arrives

General MPI Program Structure:



Example in C Language

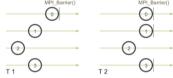
```
// required MPI include file
  #include "mpi.h"
  #include <stdio.h>
   int main(int argc, char *argv[]) {
   int numtasks, rank, len, rc;
   char hostname[MPI MAX PROCESSOR NAME];
  MPI Init(&argc,&argv); // initialize MPI
  MPI Comm size(MPI COMM WORLD, &numtasks); // get number of tasks
  MPI Comm rank(MPI COMM WORLD, &rank); // get my rank
  MPI Get processor name(hostname, &len);
  printf ("Number of tasks= %d My rank= %d Running on %s\n",
           numtasks, rank, hostname);
  MPI Finalize(): // done with MPI
```

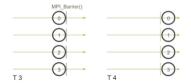
Collective Communications

- MPI also provides primitives for collective communication.
- There are many instances where communication between groups of processes is required
- Collective communication involves many tasks of the application.
- Two simple collective operations:
 - MPI_Bcast spreads data from the root task to all tasks in the communicator comm.
 - MPI_Reduce combines data from all processes in the communicator (using operation), and returns the result to the task root.

MPI Broadcast

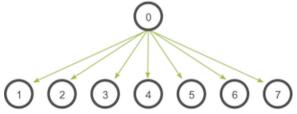
- One of the things to remember about collective communication is that it implies a synchronization point among processes
- MPI has a special function that is dedicated to synchronizing processes:
 - MPI_Barrier(MPI_Comm communicator)





Broadcasting with MPI_Bcast

- A broadcast is one of the standard collective communication techniques.
- During a broadcast, one process sends the same data to all processes in a communicator.
- One of the main uses of broadcasting is to send out user input to a parallel program, or send out configuration parameters to all processes.
- The communication pattern of a broadcast looks like this:



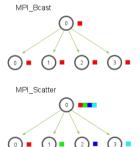
Broadcasting with MPI_Bcast . . .

- process zero is the root process, and it has the initial copy of data. All of the other processes receive the copy of data.
- In MPI, broadcasting can be accomplished by using MPI_Bcast
- The function prototype looks like this:

```
MPI_Bcast(
    void* data,
    int count,
    MPI_Datatype datatype,
    int root,
    MPI Comm communicator)
```

MPI_Scatter

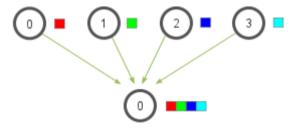
- MPI_Scatter is a collective routine that is very similar to MPI_Bcast
- MPI_Scatter involves a designated root process sending data to all processes in a communicator.
- MPI_Bcast sends the same piece of data to all processes while MPI_Scatter sends chunks of an array to different processes.



MPI Gather

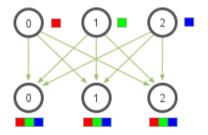
- MPI_Gather is the inverse of MPI_Scatter
- MPI_Gather takes elements from many processes and gathers them to one single process.

MPI_Gather



MPI_Allgather

- MPI_Allgather will gather all of the elements to all the processes.
- In the most basic sense, MPI_Allgather is an MPI_Gather followed by an MPI_Bcast MPI_Allgather



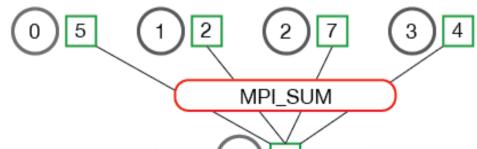
MPI Reduce and Allreduce

- Data reduction involves reducing a set of numbers into a smaller set of numbers via a function.
- Reducing this list of numbers with the sum function would produce sum([1, 2, 3, 4, 5]) = 15.
- The reduction operations defined by MPI include:
 - MPI_MAX Returns the maximum element.
 - MPI_MIN Returns the minimum element.
 - MPI_SUM Sums the elements.
 - MPI_PROD Multiplies all elements.
 - MPI_LAND Performs a logical and across the elements.
 - MPI_LOR Performs a logical or across the elements.
 - MPI_BAND Performs a bitwise and across the bits of the elements.
 - MPI_BOR Performs a bitwise or across the bits of the elements.
 - MPI_MAXLOC Returns the maximum value and the rank of the process that owns it.
 - MPI MINLOC Returns the minimum value and the rank of the process that owns it.

MPI Reduce and Allreduce

- MPI_Reduce is called with a root process of 0 and using MPI_SUM as the reduction operation.
- The four numbers are summed to the result and stored on the root process.

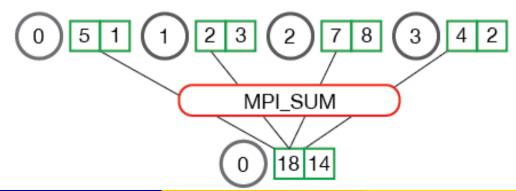
MPI_Reduce



MPI Reduce

• The illustration below shows reduction of multiple numbers per process.

MPI_Reduce



MPI Allreduce

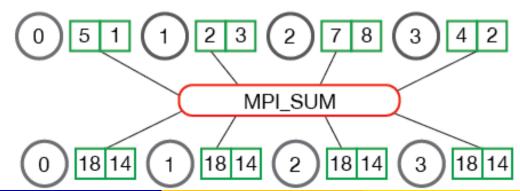
- Many parallel applications will require accessing the reduced results across all processes rather than the root process
- In a similar complementary style of MPI_Allgather to MPI_Gather, MPI_Allreduce will
 reduce the values and distribute the results to all processes.
- The function prototype is the following:

```
MPI_Allreduce(
    void* send_data,
    void* recv_data,
    int count,
    MPI_Datatype datatype,
    MPI_Op op,
    MPI Comm communicator)
```

MPI Allreduce . . .

• The following illustrates the communication pattern of MPI_Allreduce:

MPI_Allreduce



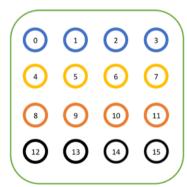
Groups and Communicators

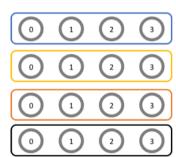
- For simple applications, it's not unusual to do everything using MPI_COMM_WORLD
- For more complex use cases, it might be helpful to have more communicators.
- the first and most common function used to create new communicators:

```
MPI_Comm_split(
    MPI_Comm comm,
    int color,
    int key,
    MPI_Comm* newcomm)
```

Groups and Communicators

Split a Large Communicator Into Smaller Communicators





Groups and Communicators

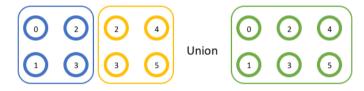
```
• Let's look at the code for this:
/ Get the rank and size in the original communicator
int world rank, world size:
MPI Comm_rank(MPI_COMM_WORLD, &world_rank);
MPI Comm size (MPI COMM WORLD, &world size);
int color = world rank / 4; // Determine color based on row
// Split the communicator based on the color and use the
// original rank for ordering
MPI Comm row comm;
MPI Comm split(MPI_COMM_WORLD, color, world_rank, &row_comm);
```

Groups and Communicators

- Another function is MPI_Comm_create.
- Its signature is almost identical:

```
MPI_Comm_create(
    MPI_Comm comm,
    MPI_Group group,
    MPI_Comm* newcomm)
```

Groups and Communicators





Intersection

Groups and Communicators

• In MPI, it's easy to get the group of processes in a communicator with the API call, MPI_Comm_group.

```
MPI_Comm_group(
     MPI_Comm comm,
     MPI_Group* group)
```

Groups and Communicators

• Once you have a group or two, performing operations on them is straightforward. Getting the union looks like this:

```
MPI_Group_union(
     MPI_Group group1,
     MPI_Group group2,
     MPI_Group* newgroup)
```

Groups and Communicators

• And you can probably guess that the intersection looks like this:

```
MPI_Group_intersection(
     MPI_Group group1,
     MPI_Group group2,
     MPI_Group* newgroup)
```

Groups and Communicators

• this function takes an MPI_Group object and creates a new communicator that has all of the same processes as the group.

```
MPI_Comm_create_group(
          MPI_Comm comm,
          MPI_Group group,
          int tag,
          MPI_Comm* newcomm)
)
```

Download example and slides

clone the source code with git or

git clone https://github.com/mesfind/openmpi.git

② Use the this link to download manually https://github.com/mesfind/openmpi

Run the examples on HPC

- omiple the source with parallel version
- 2 modified the submit.sh file
- submit the job with qsub submit.sh

Exercises: Message-Passing Programming

Installing MPI locally on your ubuntu

- sudo apt-get install gcc
- sudo apt-get install openmpi-bin
- sudo apt-get install libopenmpi-dev

Exercises: Message-Passing Programming

Hello World

- Write an MPI program which prints the message "Hello World".
- Compile and run on several processes in parallel
- Modify your program so that each process prints out both its rank and the total number of processes P that the code is running on, i.e. the size of MPI_COMM_WORLD.
- Modify your program so that only the master process (i.e. rank 0) prints out a message (very useful when you run with hundreds of processes)
- What happens if you omit the final MPI procedure call in your program?

Exercises: Message-Passing Programming

Sum of Array in c

• Write a program that sums all rows in an array using MPI parallelism.